### WIND RIVER

### **VxWorks**<sup>®</sup>

KERNEL API REFERENCE Volume 2: Routines

6.6

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VxWorks Kernel API Reference, Volume 2: Routines, 6.6

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The *VxWorks Kernel API Reference* is a two-volume set that provides reference entries describing the facilities available in the VxWorks kernel. For reference entries that describe facilities for VxWorks process-based application development, see the *VxWorks Application API Reference*. For reference entries that describe VxWorks drivers, see the *VxWorks Drivers API Reference*.

### Volume 1: Libraries

Volume 1 provides reference entries for each of the VxWorks kernel libraries, arranged alphabetically. Each entry lists the routines found in the library, including a one-line synopsis of each and a general description of their use.

Individual reference entries for each of the available functions in these libraries is provided in Volume 2.

### Volume 2: Routines

Volume 2 (this book) provides reference entries for each of the routines found in the VxWorks kernel libraries documented in Volume 1.

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hrfsDevCreate() – create an HRFS device 309 **hrfsDiskFormat()** – format a disk with HRFS 310 **hrfsFormat()** – format the HRFS file system via a path 310 **hrfsFormatFd()** – format the HRFS file system via a file descriptor 311 **hrfsFormatLibInit()** – prepare to use the HRFS formatter 312 hrfsTimeCondense() - condense time in HRFS\_TM to time in msec 313 hrfsTimeGet() – return # of milliseconds since midnight Jan 1, 1970 313 hrfsTimeSplit() – split time in msec into HRFS\_TM format hrfsUpgrade() – upgrade the HRFS file system to the latest version 314 **i()** – print a summary of each task's TCB 315 i0() – return the contents of register i0 (also i1-i7) (SimSolaris) 316 i8042vxbRegister() – register i8042vxb driver 316 ichAtaBlkRW() – read or write sectors to a ATA/IDE disk. 317 ichAtaCmd() – issue a RegisterFile command to ATA/ATAPI device. 317 ichAtaConfig() – configure an ATA drive (hard disk or cdrom drive) 321 **ichAtaConfigInit()** – initialize the hard disk driver 322 ichAtaCtrlReset() - reset the specified ATA/IDE disk controller 322 ichAtaDevCreate() – create a device for a ATA/IDE disk 323 **ichAtaDevIdentify()** – identify device 324 ichAtaDmaRW() – read/write a number of sectors on the current track in DMA mode 324 ichAtaDmaToggle() – turn on or off an individual controllers dma support 325 **ichAtaDrv()** – Initialize the ATA driver 325 ichAtaDumptest() – a quick test of the dump functionality for ATA driver 326 **ichAtaInit()** – initialize ATA device. 327 ichAtaParamRead() – Read drive parameters 327 ichAtaPiInit() – init a ATAPI CD-ROM disk controller 328 **ichAtaRW()** – read/write a data from/to required sector. 328 ichAtaRawio() – do raw I/O access 329 ichAtaShow() – show the ATA/IDE disk parameters 329 ichAtaShowInit() – initialize the ATA/IDE disk driver show routine 330 ichAtaStatusChk() – Check status of drive and compare to requested status. 330 ichAtaXbdDevCreate() – create an XBD device for a ATA/IDE disk 331 ichAtaXbdRawio() – do raw I/O access 331 ichAtapiBytesPerSectorGet() – get the number of Bytes per sector. 332 ichAtapiBytesPerTrackGet() – get the number of Bytes per track. 332 ichAtapiCtrlMediumRemoval() – Issues PREVENT/ALLOW MEDIUM REMOVAL packet command 333 ichAtapiCurrentCylinderCountGet() – get logical number of cylinders in the drive. 333 ichAtapiCurrentHeadCountGet() – get the number of read/write heads in the drive. 334 ichAtapiCurrentMDmaModeGet() – get the enabled Multi word DMA mode. 334 ichAtapiCurrentPioModeGet( ) - get the enabled PIO mode. 335 ichAtapiCurrentRwModeGet() – get the current Data transfer mode. 335 ichAtapiCurrentSDmaModeGet() – get the enabled Single word DMA mode. 336 **ichAtapiCurrentUDmaModeGet()** – get the enabled Ultra DMA mode. 336 ichAtapiCylinderCountGet() – get the number of cylinders in the drive. 337 ichAtapiDriveSerialNumberGet() – get the drive serial number. 338

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ichAtapiDriveTypeGet() – get the drive type. 338 ichAtapiFeatureEnabledGet() – get the enabled features. 339 **ichAtapiFeatureSupportedGet()** – get the features supported by the drive. 340 **ichAtapiFirmwareRevisionGet()** – get the firm ware revision of the drive. 340 **ichAtapiHeadCountGet( )** – get the number heads in the drive. 341 ichAtapiInit() – init ATAPI CD-ROM disk controller 341 **ichAtapiIoctl()** – Control the drive. 342 ichAtapiMaxMDmaModeGet() – get the Maximum Multi word DMA mode the drive supports. 346 ichAtapiMaxPioModeGet() – get the Maximum PIO mode that drive can support. 347 ichAtapiMaxSDmaModeGet() – get the Maximum Single word DMA mode the drive supports 347 ichAtapiMaxUDmaModeGet() – get the Maximum Ultra DMA mode the drive can support. 348 ichAtapiModelNumberGet() – get the model number of the drive. 349 ichAtapiPktCmd() – execute an ATAPI command with error processing 349 **ichAtapiPktCmdSend()** – Issue a Packet command. 350 ichAtapiRead10() – read one or more blocks from an ATAPI Device. 350 ichAtapiReadCapacity() – issue a READ CD-ROM CAPACITY command to a ATAPI device 351 ichAtapiReadTocPmaAtip() – issue a READ TOC command to a ATAPI device 351 ichAtapiRemovMediaStatusNotifyVerGet() – get the Media Stat Notification Version. 352 **ichAtapiScan()** – issue SCAN packet command to ATAPI drive. 352 ichAtapiSeek() – issues a SEEK packet command to drive. 353 ichAtapiSetCDSpeed() – issue SET CD SPEED packet command to ATAPI drive. 353 ichAtapiStartStopUnit() – Issues START STOP UNIT packet command 354 ichAtapiStopPlayScan() – issue STOP PLAY/SCAN packet command to ATAPI drive. 354 ichAtapiTestUnitRdy() – issue a TEST UNIT READY command to a ATAPI drive 355 ichAtapiVersionNumberGet() – get the ATA/ATAPI version number of the drive. 355 **index()** – find the first occurrence of a character in a string 356 **infinity()** – return a very large double 357 **infinityf()** – return a very large float 357 358 **inflate()** – inflate compressed code **intCRGet()** – read the contents of the cause register (MIPS) 358 359 **intCRSet()** – write the contents of the cause register (MIPS) **intConnect()** – connect a C routine to a hardware interrupt 359 **intContext()** – determine if executing in interrupt context 362 **intCount()** – get the current interrupt nesting depth 363 **intCpuLock()** – lock out interrupts on local CPU 363 **intCpuUnlock()** – cancel local CPU interrupt lock 364 intDisable() – disable corresponding interrupt bits (MIPS, PowerPC, ARM) 364 **intDisconnect()** – disconnect a C routine from a hardware interrupt 365 **intEnable()** – enable corresponding interrupt bits (MIPS, PowerPC, ARM) 366 intHandlerCreate() – construct an interrupt handler for a C routine (MC680x0, x86, MIPS, SimSolaris) 366 intHandlerCreateI86() – construct an interrupt handler for a C routine (x86) 367 intLevelSet() – set the interrupt level (MC680X0, x86, ARM, SimSolaris, SimNT and SH) 368 **intLock()** – lock out interrupts 369 intLockLevelGet() – get the current interrupt lock-out level (MC680x0, x86, ARM, SH, SimSolaris, SimNT)

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intLockLevelSet() - set the current interrupt lock-out level (MC680x0, x86, ARM, SH, SimSolaris, SimNT) 372 **intSRGet()** – read the contents of the status register (MIPS) 372 intSRSet() – update the contents of the status register (MIPS) 373 intStackEnable() – enable or disable the interrupt stack usage (x86) 373 intUninitVecSet() – set the uninitialized vector handler (ARM) 374 **intUnlock()** – cancel interrupt locks 374 intVecBaseGet() – get the vector (trap) base address (MC680x0, x86, MIPS, ARM, SimSolaris, SimNT) 375 intVecBaseSet() – set the vector (trap) base address (MC680x0, x86, MIPS, ARM, SimSolaris, SimNT) 375 intVecGet() – get an interrupt vector (MC680x0, x86, MIPS, SH, SimSolaris, SimNT) 376 **intVecGet2()** – get a CPU vector, gate type(int/trap), and gate selector (x86) 377 intVecSet() – set a CPU vector (trap) (MC680x0, x86, MIPS, SH, SimSolaris, SimNT) 378 **intVecSet2()** – set a CPU vector, gate type(int/trap), and selector (x86) 382 intVecTableWriteProtect() - write-protect exception vector table (MC680x0, x86, ARM, SimSolaris, SimNT) 382 ioGlobalStdGet() – get the file descriptor for global input/output/error 383 ioGlobalStdSet() – set file descriptor for global input/output/error 383 ioHelp() – print a synopsis of I/O utility functions 384 ioTaskStdGet() – get the file descriptor for task standard input/output/error 384 **ioTaskStdSet()** – set the file descriptor for task standard input/output/error 385 **ioctl()** – perform an I/O control function 386 **iosDevAdd()** – add a device to the kernel I/O system 387 **iosDevDelDrv()** – invoke device delete driver if reference counter reaches 0. 387 iosDevDelete() - delete a device from the kernel I/O system 388 **iosDevFind()** – find an I/O device in the kernel device list 389 iosDevShow() – display the list of devices in the system 390 iosDrvInstall() – install a kernel I/O driver 390 **iosDryRemove()** – remove a kernel I/O driver 391 iosDrvShow() – display a list of system drivers 392 iosFdEntryGet() – get an unused FD\_ENTRY from the pool 392 iosFdEntryReturn() - return an FD\_ENTRY to the pool 392 iosFdMaxFiles() - return maximum files for current RTP 393 iosFdShow() – display a list of file descriptor names in the system 393 iosInit() – initialize the kernel I/O system 394 **iosRtpFdShow()** – show the per-RTP *fd* table 394 **iosShowInit()** – initialize the I/O system show facility 395 **irint()** – convert a double-precision value to an integer 395 irintf() - convert a single-precision value to an integer 396 **iround()** – round a number to the nearest integer 396 **iroundf()** – round a number to the nearest integer 397 **isatty()** – return whether the underlying driver is a *tty* device 397 isrCreate() - create an ISR object 398 isrDelete() – delete an ISR object 399 **isrIdSelf()** – get the ISR ID of the currently running ISR 400 isrInfoGet( ) - get information about an ISR object 400

isrInvoke() – invoke the handler routine of an ISR object 401 isrShow() – show information about an ISR object 402 kernelCpuEnable() – enable a CPU 403 **kernelInit()** – initialize the kernel 403 **kernellsCpuldle()** – determine whether the specified CPU is idle 404 **kernelIsSystemIdle()** – determine whether all enabled processors are idle 405 **kernelRoundRobinInstall()** – install VxWorks Round Robin implementation 405 **kernelTimeSlice()** – enable round-robin selection 406 **kernelVersion()** – return the WIND kernel revision string 407 kill() – send a signal to a task (POSIX) 407 **1()** – disassemble and display a specified number of instructions 408 **10()** – return the contents of register 10 (also 11-17) (SimSolaris) 409 ld() – load an object module into memory 409 **ledClose()** – discard the line-editor ID 410 ledControl() – change the line-editor ID parameters 411 **ledLibInit()** – initialize the line editing facilities 411 ledOpen() – create a new line-editor ID 412 ledRead() – read a line with line-editing 412 link() – link a file 413 **lio\_listio()** – initiate a list of asynchronous I/O requests (POSIX) 413 lkAddr() – list symbols whose values are near a specified value 414 **lkup()** – list symbols 415 11() – generate a long listing of directory contents 415 **llr()** – do a long listing of directory and all its subdirectories contents 416 **InPciRegister()** – register with the VxBus subsystem 417 loadModule() – load an object module into memory 417 **loadModuleAt()** – load an object module into memory 418 **log10f()** – compute a base-10 logarithm (ANSI) 423 log2() – compute a base-2 logarithm 424 **log2f()** – compute a base-2 logarithm 424 logFdAdd() – add a logging file descriptor 425 **logFdDelete()** – delete a logging file descriptor 425 **logFdSet()** – set the primary logging file descriptor 426 **logInit()** – initialize message logging library 426 **logMsg()** – log a formatted error message 427 logTask() – message-logging support task 428 **logf()** – compute a natural logarithm (ANSI) 428 loginDefaultEncrypt() – default password encryption routine 429 **loginEncryptInstall()** – install an encryption routine 429 **loginInit()** – initialize the login table 430 loginPrompt() – display a login prompt and validate a user entry 431 loginStringSet() – change the login string 431 loginUserAdd() – add a user to the login table 432 **loginUserDelete()** – delete a user entry from the login table 433 loginUserShow() – display the user login table 434

**loginUserVerify()** – verify a user name and password in the login table 434 **logout()** – log out of the VxWorks system 435 **ls()** – generate a brief listing of a directory 435 **lseek()** – set a file read/write pointer 436 **lsr()** – list the contents of a directory and any of its subdirectories 437 lstAdd() – add a node to the end of a list 437 lstConcat() – concatenate two lists 438 **lstCount()** – report the number of nodes in a list 438 **lstDelete()** – delete a specified node from a list 439 **lstExtract()** – extract a sublist from a list 439 **lstFind()** – find a node in a list 440 **lstFirst()** – find first node in list 440 **lstFree()** – free up a list 441 **lstGet()** – delete and return the first node from a list 441 **lstInit()** – initialize a list descriptor 442 **lstInsert()** – insert a node in a list after a specified node 442 **lstLast()** – find the last node in a list 443 **lstNStep()** – find a list node *nStep* steps away from a specified node 443 **lstNext()** – find the next node in a list 444 lstNth() – find the Nth node in a list 444 **lstPrevious()** – find the previous node in a list 445 **m()** – modify memory 445 m6845vxbRegister() – register m6845vxb driver 446 m85xxCCSRRegister() - register m85xxLAWBAR driver 446 **mRegs()** – modify registers 447 mach() – return the contents of system register mach (also macl, pr) (SH) 448 **malloc()** – allocate a block of memory from the system memory partition (ANSI) 448 memAddToPool() - add memory to the system memory partition 449 **memDevCreate()** – create a memory device 449 **memDevCreateDir()** – create a memory device for multiple files 451 memDevDelete() – delete a memory device 451 **memDrv()** – install a memory driver 452 memEdrBlockMark() – mark or unmark selected blocks 452 memEdrBlockShow() – print memory block information 453 **memEdrFreeQueueFlush()** – flush the free queue 454 **memEdrPartShow()** – show partition information in the kernel 454 memEdrRtpBlockMark() - mark or unmark selected allocated blocks in an RTP 455 **memEdrRtpBlockShow()** – print memory block information of an RTP 455 **memEdrRtpPartShow()** – show partition information of an RTP 456 **memFindMax()** – find the largest free block in the system memory partition (kernel heap) 457 **memInfoGet()** – get heap information 457 **memOptionsGet()** – get the options of the system memory partition (kernel heap) 458 **memOptionsSet()** – set the options for the system memory partition (kernel heap) 458 **memPartAddToPool()** – add memory to a memory partition 459 memPartAlignedAlloc() – allocate aligned memory from a partition 460

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memPartAlloc() – allocate a block of memory from a partition 460 **memPartCreate()** – create a memory partition 461 memPartDelete() – delete a partition and free associated memory 462 **memPartFindMax()** – find the size of the largest available free block 462 **memPartFree()** – free a block of memory in a partition 463 memPartInfoGet() – get partition information 464 memPartOptionsGet( ) - get the options of a memory partition 464 memPartOptionsSet() – set the options for a memory partition 465 memPartRealloc() – reallocate a block of memory in a specified partition 466 **memPartShow()** – show blocks and statistics for a given memory partition 466 **memPartSmCreate()** – create a shared memory partition (VxMP Option) 467 memShow() - show blocks and statistics for the current heap partition 468 **memShowInit()** – initialize the memory partition show facility 471 memalign() – allocate aligned memory from system memory partition (kernel heap) 471 miiBusCreate() – create an miiBus attached to a parent bridge 472 miiBusDelete() – delete an miiBus and all its child devices miiBusGet() – get the miiBus that goes with a given VxBus instance 473 miiBusListAdd() – Add a PHY to the MII monitor list 474 **miiBusListDel()** – Remove a PHY to the MII monitor list 474 miiBusMediaAdd() – add an entry to an miiBus's media list 475 miiBusMediaDefaultSet() - set the default media for an miiBus 475 miiBusMediaDel() – delete an entry to an miiBus's media list 476 miiBusMediaListGet() – obtain a pointer to the bus's media list 476 miiBusMediaUpdate() – invoke a PHY's parent's media update callback 477 miiBusModeGet() – get the current media mode and link status 477 miiBusModeSet() – set the current media mode 478 miiBusRead() – read a PHY register 478 miiBusRegister() – register with the vxBus subsystem 479 **miiBusWrite()** – write value to a PHY register 479 **mkdir()** – make a directory 480 mlock() – lock specified pages into memory (POSIX) 480 **mlockall()** – lock all pages used by a process into memory (POSIX) 481 mmapShow() - show information about memory mapped objects in the system 481 mmuPhysToVirt() – translate a physical address to a virtual address (ARM) 483 mmuPro32LibInit() – initialize module 483 mmuPro32Page0UnMap() - unmap the page zero for NULL pointer detection 484 mmuPro36LibInit() – initialize module 484 **mmuPro36Page0UnMap()** – unmap the page zero for **NULL** pointer detection 485 **mmuPro36PageMap()** – map 36bit physical memory page to virtual memory page 485 mmuPro36Translate() – translate a virtual address to a 36bit physical address 486 **mmuShLibInit()** – Initialize the SH MMU library. 486 mmuVirtToPhys() - translate a virtual address to a physical address (ARM) 487 moduleCheck() – verify checksums on all modules loaded in the system 488 **moduleCreate()** – create and initialize a module 489 **moduleCreateHookAdd()** – add a routine to be called when a module is added 489

moduleCreateHookDelete() – delete a previously added module create hook routine 490 moduleDelete() - delete module ID information 490 **moduleFindByGroup()** – find a module by group number 491 **moduleFindByName()** – find a module by name 492 **moduleFindByNameAndPath()** – find a module by filename and path 492 **moduleFlagsGet()** – get the flags associated with a module ID 493 moduleIdListGet() - get a list of loaded modules 493 moduleInfoGet( ) - get information about an object module 494 **moduleNameGet()** – get the name associated with a module ID 494 moduleSegFirst() - find the first segment in a module 495 **moduleSegGet()** – get (delete and return) the first segment from a module 495 moduleSegNext() - find the next segment in a module 496 **moduleShow()** – show information about loaded modules 497 **mountdInit()** – initialize the mount daemon 498 mqPxDescObjIdGet() - returns the OBJ\_ID associated with a mqd\_t descriptor 499 **mqPxLibInit()** – initialize the POSIX message queue library 500 **mqPxShow()** – display message queue internals 500 **mqPxShowInit()** – initialize the POSIX message queue show facility 501 **mq\_close()** – close a message queue (POSIX) 501 **mq\_getattr()** – get message queue attributes (POSIX) 502 503 **mq\_notify()** – notify a task that a message is available on a queue (POSIX) **mq\_open()** – open a message queue (POSIX) 504 **mq\_receive()** – receive a message from a message queue (POSIX) 506 **mq\_send()** – send a message to a message queue (POSIX) 507 **mq\_setattr()** – set message queue attributes (POSIX) 508 **mq\_unlink()** – remove a message queue (POSIX) 509 **msgQClose()** – close a named message queue 510 msgQCreate() – create and initialize a message queue 511 **msgODelete()** – delete a message queue 512 msgQEvStart() – start the event notification process for a message queue 513 msgQEvStop() – stop the event notification process for a message queue 514 **msgQInfoGet()** – get information about a message queue 515 **msgQInitialize()** – initialize a pre-allocated message queue 517 msgQNumMsgs() – get the number of messages queued to a message queue 518 **msgOOpen()** – open a message queue 519 **msgQOpenInit()** – initialize the message queue open facility 521 **msgQReceive()** – receive a message from a message queue 521 **msgQSend()** – send a message to a message queue 523 **msgQShow()** – show information about a message queue 524 **msgOShowInit()** – initialize the message queue show facility 525 msgQSmCreate() – create and initialize a shared memory message queue (VxMP Option) 526 msgQUnlink() – unlink a named message queue 527 **munlock()** – unlock specified pages (POSIX) 528 munlockall() – unlock all pages used by a process (POSIX) 528 **mv()** – mv file into other directory. 529

530 **nanosleep()** – suspend the current task until the time interval elapses (POSIX) **netHelp()** – print a synopsis of network routines 530 **nfsAuthUnixGet()** – get the NFS UNIX authentication parameters 531 nfsAuthUnixPrompt() – modify the NFS UNIX authentication parameters 532 **nfsAuthUnixSet()** – set the NFS UNIX authentication parameters 533 **nfsAuthUnixShow()** – display the NFS UNIX authentication parameters 533 **nfsChkFilePerms()** – check the NFS file permissions with a given permission. 534 534 nfsDevInfoGet() – read configuration information from the requested device **nfsDevListGet()** – create list of all the NFS devices in the system 535 **nfsDevShow()** – display the mounted NFS devices 535 **nfsDrvNumGet()** – Get driver number of NFS device 536 **nfsErrnoSet()** – set NFS status 536 nfsExport() – specify a file system to be NFS exported 537 **nfsExportShow()** – display the exported file systems of a remote host 538 **nfsHelp()** – display the NFS help menu 538 539 **nfsIdSet()** – set the ID number of the NFS UNIX authentication parameters nfsMntDump() – display all NFS file systems mounted on a particular host 540 **nfsMount()** – mount an NFS file system 540 nfsMountAll() - mount all file systems exported by a specified host 541 **nfsStatusGet()** – Get the statistics of the NFS server 542 **nfsUnexport()** – remove a file system from the list of exported file systems 542 **nfsUnmount()** – unmount an NFS device 543 **nfsdHashTableParamsSet()** – sets up the parameters for the NFS hash table 543 **nfsdInit()** – initialize the NFS server 544 nfsdStatusShow() – show the status of the NFS server 545 nicRegister() – register with the VxBus subsystem 546 **npc()** – return the contents of the next program counter (SimSolaris) 546 nseRegister() – register with the VxBus subsystem 547 nvRamSegDefGet() – get segment allocation from BSP 547 **o0()** – return the contents of register o0 (also o1-o7) (SimSolaris) 548 objClassTypeGet() – get an object's class type 549 **objContextGet()** – return the object's context value 549 **objContextSet()** – set the object's context value 550 **objHandleShow()** – show information on the object referenced by an object handle 550 objHandleTblShow() - show information on an RTP's handle table 551 **objNameGet()** – get an object's name 551 objNameLenGet() - get an object's name length 552 **objNameToId()** – find object with matching name string and type 552 **objOwnerGet()** – return the object's owner 554 objOwnerSet() – change the object's owner 554 objShow() – show information on an object 555 objShowAll() - show all information on an object 555 **open()** – open a file 557 **opendir()** – open a directory for searching (POSIX) 559 **operator\_delete()** – default run-time support for memory deallocation (C++) 560

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# CPToUtf16()

NAME	<b>CPToUtf16()</b> – Convert a Unicode codepoint to a UTF-16 encoding	
SYNOPSIS	<pre>int CPToUtf16   (     const unsigned long codePoint,     unsigned short * utf16,     const int length, /* Length is in 16-bit words */     const int littleEndian   )</pre>	
DESCRIPTION	This routine converts an unsigned long representing the value of a Unicode codepoint to the UTF-16 encoding.	
RETURNS	If positive, the return value indicates the number of words used to encode this codepoint. If non-positive, the return value of UC_FORMAT indicates that the value to be converted is not a legitimate Unicode codepoint. A return value of UC_BUFFER indicates that the output string is of insufficient length to hold the converted character.	
ERRNO	Not Available	
SEE ALSO	utfLib	

## CPToUtf8()

NAME	<b>CPToUtf8()</b> – Convert a Unicode codepoint to a UTF-8 encoding	
SYNOPSIS	<pre>int CPToUtf8  (     const unsigned long codePoint,     unsigned char * utf8,     const int length  )</pre>	
DESCRIPTION	This routine converts an unsigned long representing the value of a Unicode codepoint to the UTF-8 encoding.	
RETURNS	If positive, the return value indicates the number of bytes used to encode this codepoint. If non-positive, the return value of <b>UC_FORMAT</b> indicates that the value to be converted is not a legitimate Unicode codepoint. A return value of <b>UC_BUFFER</b> indicates that the output string is of insufficient length to hold the converted character.	

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VxWorks Kernel API Reference, 6.6 CPUSET\_ATOMICCLR()

Not Available ERRNO

utfLib SEE ALSO

### **CPUSET ATOMICCLR()**

NAME CPUSET_ATOMICCLR() – atomically clear a CPU from	a CPU set
---	-----------

SYNOPSIS

CPUSET\_ATOMICCLR (

n

)

```
cpuset
             /* CPU set to operate on */
             /* index of CPU to clear */
```

DESCRIPTION This macro atomically clears CPU index *n* from the *cpuset* variable. The status of other CPU indices in the set, whether set or cleared, is not a affected by this action. This action is the reverse of what CPUSET\_ATOMICSET does. Atomic clearing of a CPU in a set is necessary when the set is likely to be manipulated by more than one task or ISR.

> While this macro does not enforce any restrictions, it is expected that cpuset is always a cpuset\_t type variable and the CPU index is always an unsigned integer between 0 and the number of CPUs either enabled or configured in the system. APIs that expect a cpuset\_t variable as an argument describe the restrictions that apply.

RETURNS	N/A

N/A ERRNO

SEE ALSO cpuset, CPUSET\_CLR, CPUSET\_ATOMICSET

### CPUSET\_ATOMICCOPY()

```
NAME
               CPUSET_ATOMICCOPY() – atomically copy a CPU set value
```

SYNOPSIS	CPUSET_ATOMICCLR	
	(	
	cpusetDst,	/* cpuset to copy to */
	cpusetSrc	/* cpuset to copy from */
	)	

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DESCRIPTION	This macro atomically copies the bit sets from <i>cpusetSrc</i> cpuset and stores the copy ir <i>cpusetDst</i> variable.	
	While this macro does not enforce any restrictions, it is expected that <i>cpusetSrc</i> and <i>cpusetDst</i> are cpuset_t type variables. APIs that expect a cpuset_t variable as an argument describe the restrictions that apply.	
RETURNS	N/A	
ERRNO	N/A	
SEE ALSO	cpuset	

# **CPUSET\_ATOMICSET()**

NAME	<b>CPUSET_ATOMICSET()</b> – atomically set a CPU in a CPU set	
SYNOPSIS	CPUSET_ATOMICSET ( cpuset /* CPU set to operate on */ n /* index of CPU to set */ )	
DESCRIPTION	This macro atomically sets CPU index <i>n</i> in the <i>cpuset</i> variable. It is the atomic version of <b>CPUSET_SET</b> . The status of other CPU indices in the set, whether set or cleared, is not affected by this action. For example, to set CPU0 and CPU1 in a set, this macro needs to be used twice specifying n=0 and then n=1. Atomic setting of a CPU in a set is necessary when the set is likely to be manipulated by more than one task or ISR.	
	While this macro does not enforce any restrictions, it is expected that cpuset is always a cpuset_t type variable and the CPU index is always an unsigned integer between 0 and the number of CPUs either enabled or configured in the system. APIs that expect a cpuset_t variable as an argument describe the restrictions that apply.	
RETURNS	N/A	
ERRNO	N/A	
SEE ALSO	cpuset, CPUSET_SET, CPUSET_ATOMICCLR	

VxWorks Kernel API Reference, 6.6 CPUSET\_CLR()

## **CPUSET\_CLR()**

NAME	<b>CPUSET_CLR()</b> – clear a CPU from a CPU set		
SYNOPSIS	CPUSET_CLR ( cpuset /* CPU set to operate on */ n /* index of CPU to clear */ )		
DESCRIPTION	This macro clears CPU index <i>n</i> in the <i>cpuset</i> variable. The status of other CPU indices in the set, whether set or cleared, is not affected by this action. This action is the reverse of what <b>CPUSET_SET</b> does.		
	While this macro does not enforce any restrictions, it is expected that cpuset is always a cpuset_t type variable and the CPU index is an unsigned integer between 0 and the number of CPUs either enabled or configured in the system. APIs that expect a cpuset_t variable as an argument describe the restrictions that apply.		
RETURNS	N/A		
ERRNO	N/A		
SEE ALSO	cpuset, CPUSET_ZERO, CPUSET_ISZERO, CPUSET_SET, vxCpuConfiguredGet()		

# **CPUSET\_ISSET()**

NAME	<b>CPUSET_ISSET()</b> – determine if a CPU is set in a CPU set	
SYNOPSIS	CPUSET_ISSET ( cpuset /* CPU set to operate on */ n /* index of CPU to query */ )	
DESCRIPTION	This macro resolves to <b>TRUE</b> if the index of CPU <i>n</i> is set in <i>cpuset</i> . Otherwise it returns <b>FALSE</b> .	
	While this macro does not enforce any restrictions, it is expected that cpuset is always a cpuset_t type variable.	
RETURNS	Macro resolves to TRUE or FALSE.	

ERRNO N/A

SEE ALSO cpuset, CPUSET\_SET, CPUSET\_SETALL, CPUSET\_SETALL\_BUT\_SELF

### CPUSET\_ISZERO()

NAME	<b>CPUSET_ISZERO()</b> – determine if all CPUs are cleared from a CPU set	
SYNOPSIS	CPUSET_ISZERO ( cpuset /* CPU set to operate on */ )	
DESCRIPTION	This macro returns <b>TRUE</b> if variable <i>cpuset</i> is empty of CPU indices. Otherwise it returns <b>FALSE</b> . While this macro does not enforce any restrictions, it is expected that cpuset is always a	
	cpuset_t type variable.	
RETURNS	Macro resolves to TRUE or FALSE.	
ERRNO	N/A	
SEE ALSO	cpuset, CPUSET_ZERO	

### CPUSET\_SET()

NAME CPUSET\_SET() – set a CPU in a CPU set

CPUSET_SET	
cpuset,	/* CPU set to operate on */
n )	/* index of CPU to set */
	CPUSET_SET ( cpuset, n )

**DESCRIPTION** This macro sets CPU index *n* in the *cpuset* variable. The status of other CPU indices in the set, whether set or cleared, is not affected by this action. For example, to set CPU0 and CPU1 in a set, this macro needs to be used twice specifying n=0 and then n=1.

While this macro does not enforce any restrictions, it is expected that cpuset is always a cpuset\_t type variable and the CPU index is an unsigned integer between 0 and the number

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of CPUs either enabled or configured in the system. APIs that expect a cpuset\_t variable as an argument describe the restrictions that apply.

RETURNS	N/A
ERRNO	N/A

SEE ALSO cpuset, CPUSET\_SETALL, CPUSET\_SETALL\_BUT\_SELF, CPUSET\_ISSET

## CPUSET\_SETALL()

NAME	<b>CPUSET_SETALL()</b> – set all CPUs in a CPU set	
SYNOPSIS	CPUSET_SETALL ( cpuset /* CPU set to operate on */ )	
DESCRIPTION	This macro sets the <i>cpuset</i> variable with the index of every CPU that is configured in the system.	
	While this macro does not enforce any restrictions, it is expected that cpuset is always a cpuset_t type variable.	
RETURNS	N/A	
ERRNO	N/A	
SEE ALSO	cpuset, CPUSET_SET, CPUSET_SETALL_BUT_SELF, CPUSET_ISSET, vxCpuConfiguredGet( )	

# CPUSET\_SETALL\_BUT\_SELF()

#### NAME CPUSET\_SETALL\_BUT\_SELF() – set all CPUs except self in CPU set

SYNOPSIS CPUSET\_SETALL\_BUT\_SELF
(
cpuset /\* CPU set to operate on \*/
)

2 Routines CPUSET\_ZERO()

DESCRIPTION	This macro sets the <i>cpuset</i> variable with the index of every CPU that is configured in the system excluding the index of the CPU that is calling this macro. Users must be aware that after a cpuset_t variable is set using this macro, the CPU index of the calling CPU, the one that is excluded from the cpuset, may no longer be the correct one when the variable is subsequently used, because a scheduling event may have taken the CPU away from the caller of the macro. The following scenario describes how this can take place:
	<ul> <li>Task 1, running on CPU0, sets myCpuSet using CPUSET_SETALL_BUT_SELF. This causes myCpuSet to be set with the index of every configured CPU except CPU0.</li> </ul>
	- An interrupt occurs and causes a scheduling event that makes Task 2 run on CPU0, effectively making Task 1 wait for a CPU to become available.
	- Another scheduling event occurs freeing CPU1 allowing Task 1 to resume execution. At this point the excluded index in myCpuSet is that of CPU0, not that of CPU1, which is the CPU on which task1 now runs.
	Conceptually this is the same issue as explained in the reference entry for <b>vxCpuIndexGet()</b> . One solution to this problem is for a task to use <b>taskCpuLock()</b> to prevent it from migrating to another CPU while is makes use of a cpuset that has been set with <b>CPUSET_SETALL_BUT_SELF</b> . This issue does not exist for ISRs since these never migrate from one CPU to another while they are executing.
	While this macro does not enforce any restrictions, it is expected that cpuset is always a cpuset_t type variable.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	cpuset, CPUSET_SET, CPUSET_SETALL, CPUSET_ISSET, vxCpuIndexGet( ), vxCpuConfiguredGet( )

## **CPUSET\_ZERO()**

 NAME
 CPUSET\_ZERO() - clear all CPUs from a CPU set

 SYNOPSIS
 CPUSET\_ZERO
 <br/>( cpuset /\* CPU set to operate on \*/
 )

 DESCRIPTION
 This macro clears all CPU indices from the *cpuset* variable. This action

**PTION** This macro clears all CPU indices from the *cpuset* variable. This action is the reverse of what **CPUSET\_SETALL** does.

#### VxWorks Kernel API Reference, 6.6 Sysctl()

	While this macro does not enforce any restrictions, it is expected that cpuset is always a cpuset_t type variable.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	cpuset, CPUSET_CLR, CPUSET_ISZERO, CPUSET_SETALL
	Sysctl()
NAME	<b>Sysctl( )</b> – get or set values for kernel state variables from the C shell
SYNOPSIS	STATUS Sysctl ( char *cmd )
DESCRIPTION	Sysctl is a C shell utility that can be used to get or set the values of various run-time parameters in the system. The underlying calls for this utility are the same as for sysctl and sysctlbyname. Sysctl can be used to get or set simple INTEGER type variables. It does not

have the ability to retrive complex data such as structures and executing procedures. Structures such as "icmp stats" are displayed as hex dumps. To configure these complex variables, one must use either "sysctl" or "sysctlbyname".

- USAGE To get the value of a variable, use Sysctl ("net.inet.ip.forwarding")
  - To set the value of a variable, use Sysctl ("net.inet.ip.forwarding=0")

To get a listing of all the variables registered in the system, Sysctl ("-A")

Options: -A Equivalent to -o -a (for compatibility).

-a List all the currently available non-opaque values. This option is ignored if one or more variable names are specified on the command line.

	-b	Force the value of the variable(s) to be output in raw, binary format. No names are printed and no terminating newlines are output. This is mostly useful with a single variable.
	-е	Separate the name and the value of the variable(s) with =. This is useful for producing output which can be fed back to the sysctl utility. This option is ignored if either -N or -n is specified, or a variable is being set.
	-N	Show only variable names, not their values.
	-n	Show only variable values, not their names.
	-0	Show opaque variables (which are normally suppressed). The format and length are printed, as well as a hex dump of the first sixteen bytes of the value.
	-X	Equivalent to -x -a (for compatibility).
	-x	As -o, but prints a hex dump of the entire value instead of just the first few bytes.
RETURNS	ОК	upon success or ERROR if an error occurred
ERRNO	N/.	A

SEE ALSO sysctl

# VX\_MEM\_BARRIER\_R()

NAME	VX_MEM_BARRIER_R() – Read Memory Barrier
SYNOPSIS	void VX_MEM_BARRIER_R ( void )
DESCRIPTION	This routine is the read memory barrier that guarantees all load memory operations before the read memory barrier has occurred before any subsequent load operations after the read barrier.

A read barrier does not necessarily have the same effect for store operations; hence a read memory barrier should not be used to ensure ordering of store operations.

Generally, a read memory barrier is paired with a write memory barrier in a multicore system to ensure proper interactions between CPUs. Here is an example of read and write barrier used in a system to ensure proper reading and writing of memory.

Example:

```
CPU 1 CPU 2
aa = value1
VX_MEM_BARRIER_W()
bb = value2
cc = aa
VX_MEM_BARRIER_R()
dd = bb
```

In the above example, the write of aa must be enforced before CPU 2 can load the value *value1* of aa. Otherwise, the wrong value of aa might be loaded. The read barrier is necessary to ensure that the read of aa is performed before performing the read of bb.

SEE ALSO	vxAtomicLib
ERRNOS	N/A
RETURNS	N/A

#### VX\_MEM\_BARRIER\_RW()

NAME	VX_MEM_BARRIER_RW() – Read/Write Memory Barrier
SYNOPSIS	void VX_MEM_BARRIER_RW ( void )
DESCRIPTION	This routine provides the read/write memory barrier. A read/write memory barrier is a general barrier that enforces ordering for both reads and writes. Hence, this read/write memory barrier can be used as a substitution for both read and write barriers.
RETURNS	N/A
ERRNOS	N/A
SEE ALSO	vxAtomicLib

#### VX\_MEM\_BARRIER\_W() VX\_MEM\_BARRIER\_W() – Write memory barrier NAME SYNOPSIS void VX\_MEM\_BARRIER\_W ( void ) DESCRIPTION This routine is the write memory barrier that guarantees all store memory operations before the write barrier has occurred before any subsequent store operations after the barrier. A write barrier does not necessarily have the same effect for load operations; hence a write memory barrier should not be used to ensure ordering of load operations. This routine is essential to ensure proper accesses to memory that are shared between CPUs in an SMP system since this enforces the ordering of memory accesses. N/A RETURNS N/A ERRNOS vxAtomicLib SEE ALSO

## a0()

NAME	<b>a0()</b> – return the contents of register <b>a0</b> (also <b>a1</b> - <b>a7</b> ) (MC680x0)
SYNOPSIS	<pre>int a0   (    int taskId /* task ID, 0 means default task */ )</pre>
DESCRIPTION	This command extracts the contents of register <b>a0</b> from the TCB of a specified task. If <i>taskId</i> is omitted or zero, the last task referenced is assumed.
	Similar routines are provided for all address registers (a0 - a7): a0() - a7().
	The stack pointer is accessed via <b>a7()</b> .
RETURNS	The contents of register <b>a0</b> (or the requested register).
ERRNO	Not Available

VxWorks Kernel API Reference, 6.6 access()

**SEE ALSO dbgArchLib**, the VxWorks programmer guides.

## access()

NAME	<b>access( )</b> – determine accessibility of a file
SYNOPSIS	<pre>int access   (    const char *path, /* path of the file */    int amode /* access mode to query */   )</pre>
DESCRIPTION	The <b>access()</b> function checks the file named by the pathname pointed to by the <i>path</i> argument for accessibility according to the bit pattern contained in <i>amode</i> , This allows a process, RTP to verify that it has permission to access this file.
	The value of <i>amode</i> is either the bitwise inclusive OR of the access permissions to be checked ( <b>R_OK</b> , <b>W_OK</b> , <b>X_OK</b> ) or the existence test, <b>F_OK</b> .
	If any access permissions are to be checked, each will be checked individually. If the process has appropriate privileges, it may indicate success even if none of the related permission bits is set.
	These constants are defined in <i>unistd.h</i> as follows:
	R_OK Test for read permission.
	W_OK Test for write permission.
	X_OK Test for execute or search permission.
	F_OK Check existence of file
RETURNS	If the requested access is permitted, <b>access()</b> succeeds and returns <b>OK</b> , 0. Otherwise, <b>ERROR</b> , -1 is returned and errno is set to indicate the error.
ERRNO	<b>ENOENT</b> Either <i>path</i> is an empty string or <b>NULL</b> pointer.
	ELOOP Circular symbolic link of <i>path</i> , or too many links.
	EMFILE Maximum number of files already open.

#### **S\_iosLib\_DEVICE\_NOT\_FOUND** (ENODEV) No valid device name found in *path*.

others

Other errors reported by device driver of *path*.

SEE ALSO fsPxLib

# acosf()

NAME	acosf() – compute an arc cosine (ANSI)
SYNOPSIS	<pre>float acosf   (    float x /* number between -1 and 1 */   )</pre>
DESCRIPTION	This routine computes the arc cosine of $x$ in single precision. If $x$ is the cosine of an angle $T$ , this function returns $T$ .
RETURNS	The single-precision arc cosine of $x$ in the range 0 to pi radians.
ERRNO	Not Available
SEE ALSO	mathALib

# adrSpaceInfoGet()

NAME	<b>adrSpaceInfoGet( )</b> – get status of the address space library
SYNOPSIS	STATUS adrSpaceInfoGet ( ADR_SPACE_INFO * pInfo /* address space info pointer */ )
DESCRIPTION	This routine initializes an <b>ADR_SPACE_INFO</b> structure parameter, <i>pInfo</i> , with the current state of the address space library.
	The following information is returned:

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VxWorks Kernel API Reference, 6.6 adrSpacePageUnmap()

<b>physAllocUnit</b> allocation unit (page) size in physical space.
physTotalPages total system RAM pages.
physFreePages unmapped system RAM pages.
physMaxSize largest unmapped system RAM block.
kernelAllocUnit allocation unit (page) size in kernel region.
<b>kernelTotalPages</b> total pages in the kernel region .
kernelFreePages unmapped pages in the kernel region.
kernelMaxSize largest unmapped block in kernel region.
userAllocUnit allocation unit (page) size in user region.
userTotalPages total pages in the user region.
userFreePages unmapped pages in the user region.
<b>userMaxSize</b> largest unmapped block in user region.
OK, or ERROR in case of failure.
Not Available
adrSpaceLib, adrSpaceShow

# adrSpacePageUnmap()

#### **NAME** adrSpacePageUnmap() – unmap a set of virtual pages

SYNOPSIS STATUS adrSpacePageUnmap

RETURNS

SEE ALSO

ERRNO

	( VM_CONTEXT_ID vmContext, /* VM context ID */ VIRT_ADDR virtAdr, /* base virtual address */ UINT numPages, /* pages to unmap */ UINT options /* unmap options */ )		
DESCRIPTION	This routine unmaps <i>numVirtPages</i> virtual pages, starting at <i>virtAdr</i> . The associated physical pages are returned to the physical page pool while the virtual pages remain allocated.		
RETURNS	OK on success, and ERROR otherwise.		
ERRNOS	The routine may set the following errnos:		
	<b>S_adrSpaceLib_SIZE_IS_INVALID</b> The <i>numPages</i> parameter is 0.		
	<b>S_adrSpaceLib_PARAMETER_NOT_ALIGNED</b> The parameter <i>virtAdr</i> is not aligned on a page boundary.		
	<b>S_pgMgrLib_VIRT_ADDR_OUT_OF_RANGE</b> The specified virtual address range, starting at <i>virtAdr</i> , and of size <i>numVirtPages</i> is not in the range of the RTP's virtual region.		
SEE ALSO	adrSpaceLib, pgMgrPageMap(), pgMgrPageFree().		

# adrSpaceRAMAddToPool()

NAME	adrSpaceRAMAddToPool() – add specified memory block to RAM pool
SYNOPSIS	STATUS adrSpaceRAMAddToPool ( PHYS_ADDR startAddr, /* start adress of RAM memory block */ UINT ramSize /* in bytes */ )
DESCRIPTION	This routine adds the specified RAM to the system global RAM page pool. The address range of the RAM must not partially or fully be part of the physical page pool, and must not be mapped.
	The length and start address are expected to be MMU page aligned. The smallest amount of space that can be added is one page.
RETURNS	OK on success and ERROR otherwise.
ERRNO	Possible errnos generated by this routine include:

 VxWorks Kernel API Reference, 6.6

 adrSpaceRAMReserve()

 S\_adrSpaceLib\_SIZE\_IS\_INVALID

 ramSize is less than a page.

 S\_adrSpaceLib\_PARAMETER\_NOT\_ALIGNED

 start or end address is not page size aligned.

 S\_adrSpaceLib\_PHYSICAL\_OVERLAP

 block overlaps with memory already in the RAM pool

 S\_adrSpaceLib\_ADDRESS\_OUT\_OF\_RANGE

 address out of 4GBytes range.

 SEE ALSO
 adrSpaceLib, adrSpaceShow

## adrSpaceRAMReserve()

NAME	<pre>adrSpaceRAMReserve() - reserve memory from the RAM pool</pre>
SYNOPSIS	PHYS_ADDR adrSpaceRAMReserve ( PHYS_ADDR startAddr, /* start adress of RAM memory block */ UINT ramSize /* in bytes */ )
DESCRIPTION	This routine reserves memory from the system global RAM page pool. If the <i>startAddr</i> parameter is NONE, the pages will be allocated from the free pages available in the RAM pool. If the <i>startAddr</i> parameter is not NONE this routine will attempt to get the pages at that address. If the pages are not available, the routine fails.
	The length and start address are expected to be MMU page aligned. The smallest amount of space that can be reserved is one page.
RETURNS	physical address on success, NONE otherwise.
ERRNO	Possible errnos generated by this routine include:
	<b>S_adrSpaceLib_SIZE_IS_INVALID</b> <i>ramSize</i> is less than a page.
	S_adrSpaceLib_PARAMETER_NOT_ALIGNED start or end address is not page size aligned.
SEE ALSO	adrSpaceLib, adrSpaceShow

# adrSpaceShow()

NAME	adrSpaceShow() – display information about address spaces managed by adrSpaceLib
SYNOPSIS	STATUS adrSpaceShow
	( UINT level /* verbosity level: 0 = summary, 1 = details */ )
DESCRIPTION	This routine displays information about various address space regions managed by the address space library, <b>adrSpaceLib</b> . When the parameter, <i>level</i> , is not 0, more detailed information is displayed.
	The information is displayed under the following headings:
	RAM Physical Address Space Info Displays information about the physical address space of the system RAM.
	User Region Info Displays information about the user virtual address space for the creation of RTP's, shared libraries and shared data regions.
	Kernel Region Info Displays information about the kernel virtual space, including memory-mapped IO space.
EXAMPLE	-> adrSpaceShow 1
	RAM Physical Address Space Info:
	Allocation unit size: 0x1000 Total number of units: 131072 (536870912 bytes) Number of allocated units: 96127 (393736192 bytes) Largest contiguous free block: 143134720 Number of free units: 34945 (143134720 bytes) 1 block(s) of 0x08881000 bytes (0x1777f000-0x1fffffff)
	User Region (RTP/SL/SD) Virtual Space Info:
	Allocation unit size:       0x1000         Total number of units:       720896       (2952790016 bytes)         Number of allocated units:       0       (0 bytes)         Largest contiguous free block:       1610612736         Number of free units:       720896       (2952790016 bytes)         1 block(s) of 0x5000000 bytes       (0x9000000-0xdffffff)         1 block(s) of 0x6000000 bytes       (0x2000000-0x7fffffff)
	Kernel Region Virtual Space Info:
	Allocation unit size: 0x1000 Number reserved of units: 327680 (1342177280 bytes) 1 block(s) of 0x20000000 bytes (0x0000000-0x1fffffff)

VxWorks Kernel API Reference, 6.6 adrSpaceVirtReserve()

```
1 block(s) of 0x10000000 bytes (0x8000000-0x8ffffff) value = 0 = 0x0
```

**RETURNS** OK, always

ERRNO Not Available

SEE ALSO adrSpaceShow, adrSpaceLib

# adrSpaceVirtReserve()

NAME	<pre>adrSpaceVirtReserve() - reserve memory from the virtual space</pre>
SYNOPSIS	VIRT_ADDR adrSpaceVirtReserve ( VIRT_ADDR startAddr, /* start adress of virtual memory block */ UINT numPages /* in bytes */ )
DESCRIPTION	This routine reserves virtual memory pages from the user region. If the <i>startAddr</i> parameter is NONE, the pages will be allocated from the free pages available in the region. If the <i>startAddr</i> parameter is not NONE this routine will attempt to get the pages at that address. If the pages are not available, the routine fails.
	The length and start address are expected to be MMU page aligned. The smallest amount of space that can be reserved is one page.
RETURNS	virtual address on success, NONE otherwise.
ERRNO	Possible errnos generated by this routine include:
	<b>S_adrSpaceLib_SIZE_IS_INVALID</b> <i>ramSize</i> is less than a page.
	S_adrSpaceLib_PARAMETER_NOT_ALIGNED start or end address is not page size aligned.
SEE ALSO	adrSpaceLib, adrSpaceShow

# aimCacheInit()

NAME	<b>aimCacheInit( )</b> – initialize cache aim with supplied parameters
SYNOPSIS	STATUS aimCacheInit ( CACHECONFIG * cacheConfig )
DESCRIPTION	This routine is called by the bsp from an architecture-specific initialization routine. It collects attribute information for all caches and publishes the attributes. It decides which AIM functions are to be called from the VxWorks API. It calculates maximum indices, counts, rounding factors, and so on, and creates local copies specific to the AIM routines in use.
RETURNS	ERROR if an invalid cache operation is requested, otherwise OK.
ERRNO	N/A
SEE ALSO	aimCacheLib

# aimFppLibInit()

NAME	<pre>aimFppLibInit( ) - Initialize the AIM FPU library</pre>
SYNOPSIS	void aimFppLibInit ( void )
DESCRIPTION	Initialize the AIM FPU library.
RETURNS	N/A
ERRNO	
SEE ALSO	aimFppLib

VxWorks Kernel API Reference, 6.6 aimMmuLibInit()

# aimMmuLibInit()

NAME	aimMmuLibInit() – initialize the AIM
SYNOPSIS	STATUS aimMmuLibInit (void)
DESCRIPTION	aimMmuLibInit() performs AIM-specific initialization of the MMU subsystem.
	If _WRS_NONGLOBAL_NULL_PAGE is defined, the null page pte entries are forced to be non-global. This avoids a potential conflict within the MMU between a virtual address that is (non-globally) mapped in one vm context, then accessed in a different vm context that does not have the address mapped.
RETURNS	OK, or ERROR if any problems are encountered during init.
ERRNO	
SEE ALSO	aimMmuLib

# aioShow()

NAME	aioShow() – show AIO requests
SYNOPSIS	STATUS aioShow ( int drvNum /* drv num to show (IGNORED) */ )
DESCRIPTION	This routine displays the outstanding AIO requests.
CAVEAT	The <i>drvNum</i> parameter is not used.
RETURNS	OK, always.
ERRNO	N/A.
SEE ALSO	aioPxShow

NAME

aioSysInit() – initialize the AIO system driver

SYNOPSIS	STATUS aioSysInit	
	int numTasks, int taskPrio, int taskStackSize )	/* number of system tasks */ /* AIO task priority */ /* AIO task stack size */

- **DESCRIPTION** This routine initializes the AIO system driver. It should be called once after the AIO library has been initialized. It spawns *numTasks* system I/O tasks to be executed at *taskPrio* priority level, with a stack size of *taskStackSize*. It also starts the wait task and sets the system driver as the default driver for AIO. If *numTasks*, *taskPrio*, or *taskStackSize* is 0, a default value (AIO\_IO\_TASKS\_DFLT, AIO\_IO\_PRIO\_DFLT, or AIO\_IO\_STACK\_DFLT, respectively) is used.
- **RETURNS** OK if successful, otherwise ERROR.

ERRNO N/A.

SEE ALSO aioSysDrv

### aio\_cancel()

 NAME
 aio\_cancel() - cancel an asynchronous I/O request (POSIX)

 SYNOPSIS
 int aio\_cancel

 (
 int fildes, /\* file descriptor \*/
 struct aiocb \* pAiocb /\* AIO control block \*/
 )

 DESCRIPTION
 This routine attempts to cancel one or more asynchronous I/O request(s) currently
 outstanding against the file descriptor *fildes. pAiocb* points to the asynchronous I/O control

asynchronous I/O requests associated with *fildes* are cancelled.

Normal signal delivery occurs for AIO operations that are successfully cancelled. If there are requests that cannot be cancelled, then the normal asynchronous completion process takes place for those requests when they complete.

block for a particular request to be cancelled. If *pAiocb* is **NULL**, all outstanding cancelable

2

VxWorks Kernel API Reference, 6.6 aio\_error()

	Operations that are cancelled successfully have a return status of -1 and an error status of <b>ECANCELED</b> .
RETURNS	AIO_CANCELED if requested operations were cancelled, AIO_NOTCANCELED if at least one operation could not be cancelled, AIO_ALLDONE if all operations have already completed, or ERROR if an error occurred.
ERRNO	EBADF Invalid, or closed file descriptor.
SEE ALSO	aioPxLib, aio_return(), aio_error()

# aio\_error()

NAME	<b>aio_error()</b> – retrieve error status of asynchronous I/O operation (POSIX)
SYNOPSIS	<pre>int aio_error   (    const struct aiocb * pAiocb /* AIO control block */ )</pre>
DESCRIPTION	This routine returns the error status associated with the I/O operation specified by <i>pAiocb</i> . If the operation is not yet completed, the error status will be <b>EINPROGRESS</b> .
RETURNS	<b>EINPROGRESS</b> if the AIO operation has not yet completed, <b>OK</b> if the AIO operation completed successfully, the error status if the AIO operation failed, otherwise <b>ERROR</b> .
ERRNO	EINVAL
SEE ALSO	aioPxLib

	aio_fsync( )
NAME	<b>aio_fsync( )</b> – asynchronous file synchronization (POSIX)
SYNOPSIS	<pre>int aio_fsync   (    int op, /* operation */    struct aiocb * pAiocb /* AIO control block */   )</pre>
DESCRIPTION	This routine asynchronously forces all I/O operations associated with the file, indicated by <b>aio_fildes</b> , queued at the time <b>aio_fsync()</b> is called to the synchronized I/O completion state. <b>aio_fsync()</b> returns when the synchronization request has be initiated or queued to the file or device.
	The value of <i>op</i> is either <b>O_DSYNC</b> or <b>O_SYNC</b> .
	If the call fails, the outstanding I/O operations are not guaranteed to have completed. If it succeeds, only the I/O that was queued at the time of the call is guaranteed to the relevant completion state.
	The <b>aio_sigevent</b> member of the <i>pAiocb</i> defines an optional signal to be generated on completion of <b>aio_fsync()</b> .
RETURNS	<b>OK</b> if queued successfully, otherwise <b>ERROR</b> .
ERRNO	EINVAL EBADF

SEE ALSO aioPxLib, aio\_error(), aio\_return()

# aio\_read()

NAME	aio_read() – initiate an asynchronous read (POSIX)	
SYNOPSIS	int aio_read ( struct aiocb * pAiocb /* AIO control block */ )	
DESCRIPTION	This routine asynchronously reads data based on the following parameters specified members of the AIO control structure <i>pAiocb</i> . It reads <b>aio_nbytes</b> bytes of data from t <b>aio_fildes</b> into the buffer <b>aio_buf</b> .	

2

VxWorks Kernel API Reference, 6.6 aio\_return()

	The requested operation takes place at the absolute position in the file as specified by <b>aio_offset</b> .
	<b>aio_reqprio</b> can be used to lower the priority of the AIO request; if this parameter is nonzero, the priority of the AIO request is <b>aio_reqprio</b> lower than the calling task priority.
	The call returns when the read request has been initiated or queued to the device. <b>aio_error()</b> can be used to determine the error status and of the AIO operation. On completion, <b>aio_return()</b> can be used to determine the return status.
	<b>aio_sigevent</b> defines the signal to be generated on completion of the read request. If this value is zero, no signal is generated.
RETURNS	OK if the read queued successfully, otherwise ERROR.
ERRNO	EBADF EINVAL
SEE ALSO	aioPxLib, aio_error(), aio_return(), read()

# aio\_return()

NAME	<b>aio_return()</b> – retrieve return status of asynchronous I/O operation (POSIX)	
SYNOPSIS	ssize_t aio_return ( struct aiocb * pAiocb /* AIO control block */ )	
DESCRIPTION	This routine returns the return status associated with the I/O operation specified by <i>pAiocb</i> . The return status for an AIO operation is the value that would be returned by the corresponding <b>read()</b> , <b>write()</b> , or <b>fsync()</b> call. <b>aio_return()</b> may be called only after the AIO operation has completed ( <b>aio_error()</b> returns a valid error codenot <b>EINPROGRESS</b> ). Furthermore, <b>aio_return()</b> may be called only once; subsequent calls will fail.	
RETURNS	The return status of the completed AIO request, or ERROR.	
ERRNO	EINVAL EINPROGRESS	
SEE ALSO	aioPxLib	

# aio\_suspend()

NAME	<b>aio_suspend()</b> – wait for asynchronous I/O request(s) (POSIX)	
SYNOPSIS	<pre>int aio_suspend   (    const struct aiocb *const list[], /* AIO requests */    int</pre>	
DESCRIPTION	This routine suspends the caller until one of the following occurs:	
	- at least one of the previously submitted asynchronous I/O operations referenced by <i>list</i> has completed,	
	- a signal interrupts the function, or	
	- the time interval specified by <i>timeout</i> has passed (if <i>timeout</i> is not NULL).	
RETURNS	<b>OK</b> if an AIO request completes, otherwise <b>ERROR</b> .	
ERRNO	EAGAIN EINTR	
SEE ALSO	aioPxLib	

# aio\_write()

NAME	aio_write() – initiate an asynchronous write (POSIX)	
SYNOPSIS	int aio_write ( struct aiocb * pAiocb /* AIO control block */ )	
DESCRIPTION	This routine asynchronously writes data based on the following parameters specified by members of the AIO control structure <i>pAiocb</i> . It writes <b>aio_nbytes</b> of data to the file <b>aio_fildes</b> from the buffer <b>aio_buf</b> .	
	The requested operation takes place at the absolute position in the file as specified by <b>aio_offset</b> .	
	<b>aio_reqprio</b> can be used to lower the priority of the AIO request; if this parameter is nonzero, the priority of the AIO request is <b>aio_reqprio</b> lower than the calling task priority.	

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The call returns when the write request has been initiated or queued to the device.aio\_error() can be used to determine the error status and of the AIO operation. On<br/>completion, aio\_return() can be used to determine the return status.aio\_sigevent defines the signal to be generated on completion of the write request. If this<br/>value is zero, no signal is generated.RETURNSOK if write queued successfully, otherwise ERROR.ERRNOEBADF<br/>EINVALSEE ALSOaioPxLib, aio\_error(), aio\_return(), write()

## alarm()

NAME	<b>alarm()</b> – set an alarm clock for delivery of a signal	
SYNOPSIS	unsigned int alarm ( unsigned int secs )	
DESCRIPTION	This routine arranges for a <b>SIGALRM</b> signal to be delivered to the calling task after <i>secs</i> seconds.	
	If secs is zero, no new alarm is scheduled. In all cases, any previously set alarm is cancelled.	
RETURNS	Time remaining until a previously scheduled alarm was due to be delivered, zero if there was no previous alarm, or <b>ERROR</b> in case of an error.	
ERRNO	EINVAL	
	ENOSYS	
	EAGAIN	
	S_memLib_NOT_ENOUGH_MEMORY	
SEE ALSO	timerLib	

# anRegister()

NAME	<b>anRegister()</b> – register with the VxBus subsystem	
SYNOPSIS	void anRegister(void)	
DESCRIPTION	This routine registers the AN983 driver with VxBus as a child of the PCI bus type.	
RETURNS	N/A	
ERRNO	N/A	
SEE ALSO	an983VxbEnd	

# asinf()

NAME	asinf() – compute an arc sine (ANSI)
SYNOPSIS	float asinf ( float x /* number between -1 and 1 */ )
DESCRIPTION	This routine computes the arc sine of <i>x</i> in single precision. If <i>x</i> is the sine of an angle <i>T</i> , this function returns <i>T</i> .
RETURNS	The single-precision arc sine of $x$ in the range -pi/2 to pi/2 radians.
ERRNO	Not Available
SEE ALSO	mathALib

# atan2f()

NAME	<b>atan2f( )</b> – compute the arc tangent of $y/x$ (ANSI)
SYNOPSIS	<pre>float atan2f   (    float y, /* numerator */</pre>

# VxWorks Kernel API Reference, 6.6 atanf()

```
float x/* denominator */DESCRIPTIONThis routine returns the principal value of the arc tangent of y/x in single precision.RETURNSThe single-precision arc tangent of y/x in the range -pi to pi.ERRNONot AvailableSEE ALSOmathALib
```

# atanf()

NAME	atanf() – compute an arc tangent (ANSI)
SYNOPSIS	<pre>float atanf   (    float x /* tangent of an angle */ )</pre>
DESCRIPTION	This routine computes the arc tangent of $x$ in single precision. If $x$ is the tangent of an angle $T$ , this function returns $T$ (in radians).
RETURNS	The single-precision arc tangent of $x$ in the range -pi/2 to pi/2.
ERRNO	Not Available
SEE ALSO	mathALib

## atapiParamsPrint()

NAME	<b>atapiParamsPrint( )</b> – Print the drive parameters.
SYNOPSIS	void atapiParamsPrint ( int ctrl, int drive
	)
DESCRIPTION	This user callable routine will read the current parameters from the corresponding drive
-------------	--
	and will print the specified range of parameters on the console.

RETURNS N/A.

ERRNO Not Available

SEE ALSO vxbIntelIchStorage

#### attrib()

attrib() – modify MS-DOS file attributes on a file or directory NAME SYNOPSIS STATUS attrib ( const char \* fileName, /\* file or dir name on which to change flags \*/ const char \* attr /\* flag settings to change \*/ DESCRIPTION This function provides means for the user to modify the attributes of a single file or directory. There are four attribute flags which may be modified: "Archive", "System", "Hidden" and "Read-only". Among these flags, only "Read-only" has a meaning in VxWorks, namely, read-only files can not be modified deleted or renamed. The attr argument string may contain must start with either "+" or "-", meaning the attribute flags which will follow should be either set or cleared. After "+" or "-" any of these four letter will signify their respective attribute flags - "A", "S", "H" and "R". For example, to write-protect a particular file and flag that it is a system file: -> attrib( "bootrom.sys", "+RS") OK, or ERROR if the file can not be opened. RETURNS Not Available ERRNO usrFsLib, dosFsLib, the VxWorks programmer guides. SEE ALSO

## **b()**

```
b() – set or display breakpoints
NAME
SYNOPSIS
                 STATUS b
                      (
                      INSTR * addr,
                                              /* breakpoint addr, or 0 to display */
                      int taskNameOrId, /* task affected; 0 means all tasks
                                                                                               * /
                             count, /* number of passes before hit */
                      int
                                                                                             */
                                               /* TRUE = don't print debugging info,
                      BOOL quiet
                                                /* FALSE = print debugging info */
                      )
DESCRIPTION
                 This routine sets or displays breakpoints. To display the list of currently active breakpoints,
                 call b() without arguments:
                      -> b
                 The list shows the address, task, context type, action, notification status and pass count of
                 each breakpoint. Temporary breakpoints inserted by so() and cret() are also indicated.
                 To set a breakpoint with \mathbf{b}(\mathbf{)}, include the address, which can be specified numerically or
                 symbolically with an optional offset. The other arguments are optional:
                      -> b addr[,task[,count[,quiet]]]
                 If task is zero or omitted, the breakpoint will apply to all breakable tasks. If count is zero or
                 omitted, the breakpoint will occur every time it is hit. If count is specified, the break will not
                 occur until the count +1th time an eligible task hits the breakpoint (i.e., the breakpoint is
                 ignored the first count times it is hit).
                 If quiet is specified, debugging information destined for the console will be suppressed
                 when the breakpoint is hit. This option is included for use by external source code
                 debuggers that handle the breakpoint user interface themselves.
                 Individual tasks can be unbreakable, in which case breakpoints that otherwise would apply
                 to a task are ignored. Tasks can be spawned unbreakable by specifying the task option
                 VX_UNBREAKABLE. Tasks can also be set unbreakable or breakable by resetting
                 VX_UNBREAKABLE with the routine taskOptionsSet().
RETURNS
                 OK, or ERROR if addr is illegal or the breakpoint table is full.
                 N/A
ERRNO
                 dbgLib, bd(), taskOptionsSet(), VxWorks Kernel Programmer's Guide: Kernel Shell, VxWorks
SEE ALSO
                 Command-Line Tools User's Guide 2.2: Host Shell
```

# bcmp()

NAME	<b>bcmp()</b> – compare one buffer to another
SYNOPSIS	<pre>int bcmp ( FAST char *buf1, /* pointer to first buffer */ FAST char *buf2, /* pointer to second buffer */ FAST int nbytes /* number of bytes to compare */ )</pre>
DESCRIPTION	This routine compares the first <i>nbytes</i> characters of <i>buf1</i> to <i>buf2</i> .
RETURNS	<ul> <li>0 if the first <i>nbytes</i> of <i>buf1</i> and <i>buf2</i> are identical,</li> <li>less than 0 if <i>buf1</i> is less than <i>buf2</i>, or</li> <li>greater than 0 if <i>buf1</i> is greater than <i>buf2</i>.</li> </ul>
ERRNO	N/A
SEE ALSO	bLib

# bcopy()

NAME	<b>bcopy()</b> – copy one buffer to another
SYNOPSIS	<pre>void bcopy  (     const char *source, /* pointer to source buffer */     char *destination, /* pointer to destination buffer */     int nbytes /* number of bytes to copy */  )</pre>
DESCRIPTION	This routine copies the first <i>nbytes</i> characters from <i>source</i> to <i>destination</i> . Overlapping buffers are handled correctly. Copying is done in the most efficient way possible, which may include long-word, or even multiple-long-word moves on some architectures. In general, the copy will be significantly faster if both buffers are long-word aligned. (For copying that is restricted to byte, word, or long-word moves, see the manual entries for <b>bcopyBytes()</b> , <b>bcopyWords()</b> , and <b>bcopyLongs()</b> .)
RETURNS	N/A
ERRNO	N/A

VxWorks Kernel API Reference, 6.6 bcopyBytes()

SEE ALSO bLib, bcopyBytes(), bcopyWords(), bcopyLongs()

## bcopyBytes()

NAME	<b>bcopyBytes()</b> – copy one buffer to another one byte at a time
SYNOPSIS	<pre>void bcopyBytes   (    char *source, /* pointer to source buffer */    char *destination, /* pointer to destination buffer */    int nbytes /* number of bytes to copy */   )</pre>
DESCRIPTION	This routine copies the first <i>nbytes</i> characters from <i>source</i> to <i>destination</i> one byte at a time. This may be desirable if a buffer can only be accessed with byte instructions, as in certain byte-wide memory-mapped peripherals.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	bLib, bcopy( )

# bcopyLongs()

NAME	<b>bcopyLongs( )</b> – copy one buffer to another one long word at a time
SYNOPSIS	<pre>void bcopyLongs   (    char *source, /* pointer to source buffer */    char *destination, /* pointer to destination buffer */    int nlongs /* number of longs to copy */   )</pre>
DESCRIPTION	This routine copies the first <i>nlongs</i> characters from <i>source</i> to <i>destination</i> one long word at a time. This may be desirable if a buffer can only be accessed with long instructions, as in certain long-word-wide memory-mapped peripherals. The source and destination must be long-aligned.
RETURNS	N/A

ERRNO N/A

SEE ALSO bLib, bcopy()

#### bcopyWords()

**NAME bcopyWords()** – copy one buffer to another one word at a time

SYNOPSIS void bcopyWords
(
char \*source, /\* pointer to source buffer \*/
char \*destination, /\* pointer to destination buffer \*/
int nwords /\* number of words to copy \*/
)

**DESCRIPTION** This routine copies the first *nwords* words from *source* to *destination* one word at a time. This may be desirable if a buffer can only be accessed with word instructions, as in certain word-wide memory-mapped peripherals. The source and destination must be word-aligned.

RETURNS	N/A
ERRNO	N/A

SEE ALSO bLib, bcopy()

## bd()

NAME	<b>bd( )</b> – delete a breakpoint	
SYNOPSIS	<pre>STATUS bd (     INSTR * addr, /* address of breakpoint to delete     int taskNameOrId /* task affected; 0 means all tasks )</pre>	*/ */
DESCRIPTION	This routine deletes a specified breakpoint, based on its address. To execute, enter:	
	-> bd addr [,task]	

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VxWorks Kernel API Reference, 6.6 bdall()

If *task* is omitted or zero, the breakpoint will be removed for all tasks. If the breakpoint applies to all tasks, removing it for only a single task will be ineffective. It must be removed for all tasks and then set for just those tasks desired. Temporary breakpoints inserted by the routines **so()** or **cret()** can also be deleted.

**RETURNS** OK, or **ERROR** if there is no breakpoint at the specified address.

ERRNO N/A

**SEE ALSO dbgLib**, **b()**, **bdall()**, *VxWorks Kernel Programmer's Guide: Kernel Shell*, *VxWorks Command-Line Tools User's Guide 2.2: Host Shell* 

## bdall()

NAME	bdall() – delete all breakpoints
SYNOPSIS	STATUS bdall ( int taskNameOrId /* task affected; 0 means all tasks */ )
DESCRIPTION	This routine removes all breakpoints.
	To execute, enter:
	-> bdall [task]
	If <i>task</i> is specified, all breakpoints that apply to that task are removed. If <i>task</i> is omitted, all breakpoints for all tasks are removed. Temporary breakpoints inserted by <b>so()</b> or <b>cret()</b> are not deleted; use <b>bd()</b> instead.
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>task</i> cannot be found.
ERRNO	N/A
SEE ALSO	<b>dbgLib</b> , <b>b()</b> , <b>bd()</b> , <i>VxWorks Kernel Programmer's Guide: Kernel Shell</i> , <i>VxWorks Command-Line Tools User's Guide</i> 2.2: Host Shell

#### bfStrSearch()

bfStrSearch() – Search using the Brute Force algorithm NAME SYNOPSIS char \* bfStrSearch ( char \* pattern, /\* pattern to search for \*/ int patternLen, /\* length of the pattern \*/ char \* buffer, /\* text buffer to search in \*/ int bufferLen, /\* length of the text buffer \*/ BOOL caseSensitive /\* case-sensitive search? \*/ ) The Brute Force algorithm is the simplest string search algorithm. It performs comparisons DESCRIPTION between a character in the pattern and a character in the text buffer from left to right. After each attempt it shifts the pattern by one position to the right. The Brute Force algorithm requires no pre-processing and no extra space. It has a O(Pattern Length x Text Buffer Length) worst-case time complexity. RETURNS A pointer to the located pattern, or a NULL pointer if the pattern is not found Not Available ERRNO strSearchLib SEE ALSO

#### bfill()

NAME	<b>bfill()</b> – fill a buffer with a specified character
SYNOPSIS	<pre>void bfill   (    FAST char *buf, /* pointer to buffer */    int nbytes, /* number of bytes to fill */    FAST int ch /* char with which to fill buffer */   )</pre>
DESCRIPTION	This routine fills the first <i>nbytes</i> characters of a buffer with the character <i>ch</i> . Filling is done in the most efficient way possible, which may be long-word, or even multiple-long-word stores, on some architectures. In general, the fill will be significantly faster if the buffer is long-word aligned. (For filling that is restricted to byte stores, see the manual entry for <b>bfillBytes()</b> .)
RETURNS	N/A

VxWorks Kernel API Reference, 6.6 bfillBytes()

ERRNO N/A

SEE ALSO bLib, bfillBytes()

## **bfillBytes()**

bfillBytes() - fill buffer with a specified character one byte at a time NAME SYNOPSIS void bfillBytes ( /\* pointer to buffer \*/ FAST char \*buf, nbytes, /\* number of bytes to fill \*/ int FAST int ch /\* char with which to fill buffer \*/ ) DESCRIPTION This routine fills the first *nbytes* characters of the specified buffer with the character *ch* one byte at a time. This may be desirable if a buffer can only be accessed with byte instructions, as in certain byte-wide memory-mapped peripherals. N/A RETURNS N/A ERRNO bLib, bfill() SEE ALSO

#### bh( )

NAME	<b>bh( )</b> – set a hardware break	point
SYNOPSIS	STATUS bh (	
	INSTR * addr,	/* where to set breakpoint, or */
		<pre>/* 0 = display all breakpoints */</pre>
	int access,	/* access type (arch dependant) */
	int taskNameOrId	, /* task affected; 0 means all tasks */
	int count,	/* number of passes before hit */
	BOOL quiet	<pre>/* TRUE = don't print debugging info, */</pre>
	_	/* FALSE = print debugging info */
	)	

DESCRIPTION	This routine is used to set a hardware breakpoint. If the architecture allows it, this function adds the breakpoint to the list of breakpoints and set the hardware breakpoint register(s). For more information, see the manual entry for <b>b()</b> .
NOTE	The types of hardware breakpoints vary with the architectures. Generally, a hardware breakpoint can be a data breakpoint or an instruction breakpoint.
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>addr</i> is illegal or the hardware breakpoint table is full.
ERRNO	N/A
SEE ALSO	<b>dbgLib</b> , <b>b()</b> , <i>VxWorks Kernel Programmer's Guide: Kernel Shell</i> , <i>VxWorks Command-Line Tools User's Guide 2.2: Host Shell</i>

## binvert()

NAME	<b>binvert()</b> – invert the order of bytes in a buffer
SYNOPSIS	<pre>void binvert   (   FAST char * buf, /* pointer to buffer to invert */   int     nbytes /* number of bytes in buffer */ )</pre>
DESCRIPTION	This routine inverts an entire buffer, byte by byte. For example, the buffer {1, 2, 3, 4, 5} would become {5, 4, 3, 2, 1}.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	bLib

## bmsStrSearch()

NAME bmsStrSearch() – Search using the Boyer-Moore-Sunday (Quick Search) algorithm SYNOPSIS char \* bmsStrSearch ( char \* pattern, /\* pattern to search for \*/ VxWorks Kernel API Reference, 6.6 bmtPhyRegister()

int	patternLen,	/*	length of the pattern */
char *	buffer,	/*	text buffer to search in */
int	bufferLen,	/*	length of the text buffer */
BOOL	caseSensitive	/*	case-sensitive search? */
)			

DESCRIPTIONThe Boyer-Moore-Sunday algorithm is a more efficient simplification of the Boyer-Moore<br/>algorithm. It performs comparisons between a character in the pattern and a character in the<br/>text buffer from left to right. After each mismatch it uses bad character heuristic to shift the<br/>pattern to the right. For more details on the algorithm, refer to "A Very Fast Substring Search<br/>Algorithm", Daniel M. Sunday, Communications of the ACM, Vol. 33 No. 8, August 1990,<br/>pp. 132-142.It has a O(Pattern Length x Text Buffer Length) worst-case time complexity. But empirical<br/>results have shown that this algorithm is one of the fastest in practice.

**RETURNS**A pointer to the located pattern, or a NULL pointer if the pattern is not found**ERRNO**Not Available

SEE ALSO strSearchLib

## bmtPhyRegister()

NAME	<b>bmtPhyRegister()</b> – register with the VxBus subsystem
SYNOPSIS	void bmtPhyRegister(void)
DESCRIPTION	This routine registers the BCM52xx driver with VxBus as a child of the MII bus type.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	bcm52xxPhy

#### bootBpAnchorExtract()

bootBpAnchorExtract() – extract a backplane address from a device field NAME SYNOPSIS STATUS bootBpAnchorExtract ( char \*string, /\* string containing adrs field \*/ char \*\*pAnchorAdrs /\* pointer where to return anchor address \*/ This routine extracts the optional backplane anchor address field from a boot device field. DESCRIPTION The anchor can be specified for the backplane driver by appending to the device name (i.e., "bp") an equal sign (=) and the address in hexadecimal. For example, the "boot device" field of the boot parameters could be specified as: boot device: bp=800000 In this case, the backplane anchor address would be at address 0x800000, instead of the default specified in **config.h**. This routine picks off the optional trailing anchor address by replacing the equal sign (=) in the specified string with an EOS and then scanning the remainder as a hex number. This number, the anchor address, is returned via the *pAnchorAdrs* pointer. RETURNS 1 if the anchor address in *string* is specified correctly, 0 if the anchor address in *string* is not specified, or -1 if an invalid anchor address is specified in *string*. Not Available ERRNO bootParseLib SEE ALSO

#### bootChange()

NAME bootChange() – change the boot line

**SYNOPSIS** void bootChange (void)

**DESCRIPTION** This command changes the boot line used in the boot ROMs. This is useful during a remote login session. After changing the boot parameters, you can reboot the target with the **reboot()** command, and then terminate your login (~.) and remotely log in again. As soon as the system has rebooted, you will be logged in again.

VxWorks Kernel API Reference, 6.6 bootLeaseExtract()

	This command stores the new boot line in non-volatile RAM, if the target has it.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	<b>usrLib</b> , the VxWorks programmer guides.

#### bootLeaseExtract()

**NAME bootLeaseExtract()** – extract the lease information from an Internet address

```
SYNOPSIS int bootLeaseExtract
```

```
(
char *string, /* string containing addr field */
u_long *pLeaseLen, /* pointer to storage for lease duration */
u_long *pLeaseStart /* pointer to storage for lease origin */
)
```

**DESCRIPTION** This routine extracts the optional lease duration and lease origin fields from an Internet address field for use with DHCP. The lease duration can be specified by appending a colon and the lease duration to the netmask field. For example, the "inet on ethernet" field of the boot parameters could be specified as:

```
inet on ethernet: 90.1.0.1:ffff0000:1000
```

If no netmask is specified, the contents of the field could be:

inet on ethernet: 90.1.0.1::fffffff

In the first case, the lease duration for the address is 1000 seconds. The second case indicates an infinite lease, and does not specify a netmask for the address. At the beginning of the boot process, the value of the lease duration field is used to specify the requested lease duration. If the field not included, the value of **DHCP\_DEFAULT\_LEASE** is used instead.

The lease origin is specified with the same format as the lease duration, but is added during the boot process. The presence of the lease origin field distinguishes addresses assigned by a DHCP server from addresses entered manually. Addresses assigned by a DHCP server may be replaced if the bootstrap loader uses DHCP to obtain configuration parameters. The value of the lease origin field at the beginning of the boot process is ignored.

This routine extracts the optional lease duration by replacing the preceding colon in the specified string with an EOS and then scanning the remainder as a number. The lease duration and lease origin values are returned via the *pLeaseLen* and *pLeaseStart* pointers, if those parameters are not **NULL**.

RETURNS	2 if both lease values are specified correctly in <i>string</i> , or
	-2 if one of the two values is specified incorrectly.
	If only the lease duration is found, it returns:
	1 if the lease duration in <i>string</i> is specified correctly,
	0 if the lease duration is not specified in <i>string</i> , or
	-1 if an invalid lease duration is specified in <i>string</i> .
	· · ·

ERRNO Not Available

SEE ALSO bootParseLib

## bootNetmaskExtract()

NAME	<b>bootNetmaskExtract()</b> – extract the net mask field from an Internet address
SYNOPSIS	<pre>STATUS bootNetmaskExtract   (    char *string, /* string containing addr field */    int *pNetmask /* pointer where to return net mask */   )</pre>
DESCRIPTION	This routine extracts the optional subnet mask field from an Internet address field. Subnet masks can be specified for an Internet interface by appending to the Internet address a colon and the net mask in hexadecimal. For example, the "inet on ethernet" field of the boot parameters could be specified as:
	inet on ethernet: 90.1.0.1:ffff0000
	In this case, the network portion of the address (normally just 90) is extended by the subnet mask (to 90.1). This routine extracts the optional trailing subnet mask by replacing the colon in the specified string with an EOS and then scanning the remainder as a hex number. This number, the net mask, is returned via the <i>pNetmask</i> pointer.
	This routine also handles an empty netmask field used as a placeholder for the lease duration field (see <b>bootLeaseExtract()</b> ). In that case, the colon separator is replaced with an EOS and the value of netmask is set to 0.
RETURNS	1 if the subnet mask in <i>string</i> is specified correctly, 0 if the subnet mask in <i>string</i> is not specified, or -1 if an invalid subnet mask is specified in <i>string</i> .
ERRNO	Not Available

VxWorks Kernel API Reference, 6.6 bootParamsPrompt()

SEE ALSO bootParseLib

# bootParamsPrompt( )

NAME	<pre>bootParamsPrompt() – prompt for boot line parameters</pre>
SYNOPSIS	<pre>void bootParamsPrompt   (    char *string /* default boot line */   )</pre>
DESCRIPTION	This routine displays the current value of each boot parameter and prompts the user for a new value. Typing a RETURN leaves the parameter unchanged. Typing a period (.) clears the parameter.
	The parameter <i>string</i> holds the initial values. The new boot line is copied over <i>string</i> . If there are no initial values, <i>string</i> is empty on entry.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	bootLib

## bootParamsShow()

NAME	<b>bootParamsShow()</b> – display boot line parameters
SYNOPSIS	<pre>void bootParamsShow   (    char *paramString /* boot parameter string */ )</pre>
DESCRIPTION	This routine displays the boot parameters in the specified boot string one parameter per line.
RETURNS	N/A
ERRNO	Not Available

SEE ALSO bootLib

# bootStringToStruct()

NAME	<b>bootStringToStruct()</b> – interpret the boot parameters from the boot line
SYNOPSIS	<pre>char * bootStringToStruct   (     char * bootString, /* boot line to be parsed */     FAST BOOT_PARAMS * pBootParams /* where to return parsed boot line */   )</pre>
DESCRIPTION	This routine parses the ASCII string and returns the values into the provided parameters. For a description of the format of the boot line, see the manual entry for <b>bootLib</b>
RETURNS	A pointer to the last character successfully parsed plus one (points to EOS, if <b>OK</b> ). The entire boot line is parsed.
ERRNO	Not Available
SEE ALSO	bootParseLib

# bootStringToStructAdd()

NAME	<b>bootStringToStructAdd()</b> – interpret the boot parameters from the boot line
SYNOPSIS	<pre>char * bootStringToStructAdd   (     char * bootString, /* boot line to be parsed */     FAST BOOT_PARAMS * pBootParams /* where to return parsed boot line */   )</pre>
DESCRIPTION	This routine parses the ASCII string <i>bootString</i> and returns the values into the provided parameters <i>pBootParams</i> . The fields of <i>pBootParams</i> may be previously set to default values. For a description of the format of the boot line, see the manual entry for <b>bootLib</b>
RETURNS	A pointer to the last character successfully parsed plus one (points to EOS, if <b>OK</b> ). The entire boot line is parsed.

VxWorks Kernel API Reference, 6.6 bootStructToString()

ERRNO Not Available

SEE ALSO bootParseLib

# bootStructToString()

NAME	<pre>bootStructToString() - construct a boot line</pre>
SYNOPSIS	<pre>STATUS bootStructToString   (     char *paramString, /* where to return the encoded boot line */     FAST BOOT_PARAMS *pBootParams /* boot line structure to be encoded */     )</pre>
DESCRIPTION	This routine encodes a boot line using the specified boot parameters. For a description of the format of the boot line, see the manual entry for <b>bootLib</b> .
RETURNS	OK.
ERRNO	Not Available
SEE ALSO	bootLib

# bswap()

NAME	<b>bswap( )</b> – swap buffers
SYNOPSIS	<pre>void bswap   (   FAST char * buf1, /* pointer to first buffer */   FAST char * buf2, /* pointer to second buffer */   FAST int nbytes /* number of bytes to swap */   )</pre>
DESCRIPTION	This routine exchanges the first <i>nbytes</i> of the two specified buffers.
RETURNS	N/A
ERRNO	N/A

SEE ALSO bLib

#### bzero()

NAME	<b>bzero()</b> – zero out a buffer
SYNOPSIS	<pre>void bzero   (    char * buffer, /* buffer to be zeroed */    int nbytes /* number of bytes in buffer */   )</pre>
DESCRIPTION	This routine fills the first <i>nbytes</i> characters of the specified buffer with 0.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	bLib

# **c( )**

NAME	<b>c( )</b> – continue from a breakpoin	nt
SYNOPSIS	STATUS c ( int taskNameOrId, INSTR * addr, INSTR * addr1 )	<pre>/* task that should proceed from breakpoint */ /* address to continue at; 0 = next instruction */ /* address for npc; 0 = instruction next to pc */</pre>
DESCRIPTION	This routine continues the exec	ution of a task that has stopped at a breakpoint.
	To execute, enter:	
	-> c [task [,addr[,addr	r1]]]

If *task* is omitted or zero, the last task referenced is assumed. If *addr* is non-zero, the program counter is changed to *addr*; if *addr1* is non-zero, the next program counter is changed to *addr1*, and the task is continued.

VxWorks Kernel API Reference, 6.6 cacheArchClearEntry()

CAVEAT	When a task is continued, <b>c()</b> does not distinguish between a stopped task or a task stopped by the debugger. Therefore, its use should be restricted to only those tasks being debugged.
NOTE	The next program counter, <i>addr1</i> , is currently supported only by SPARC.
RETURNS	OK, or ERROR if the specified task does not exist.
ERRNO	N/A
SEE ALSO	dbgLib, s( ), cret( ), tr( ), VxWorks Kernel Programmer's Guide: Kernel Shell, VxWorks Command-Line Tools User's Guide 2.2: Host Shell

# cacheArchClearEntry()

NAME	<b>cacheArchClearEntry()</b> – clear an entry from a cache (68K, x86)
SYNOPSIS	STATUS cacheArchClearEntry ( CACHE_TYPE cache, /* cache to clear entry for */ void * address /* entry to clear */ )
DESCRIPTION	This routine clears a specified entry from the specified cache.
	For 68040 processors, this routine clears the cache line from the cache in which the cache entry resides.
	For the MC68060 processor, when the instruction cache is cleared (invalidated) the branch cache is also invalidated by the hardware. One line in the branch cache cannot be invalidated so each time the branch cache is entirely invalidated.
For 3 and line proc the s	For 386 family processors do not have a cache, thus it does nothing. The 486, P5(Pentium), and P6(PentiumPro, II, III) family processors do have a cache but does not support a line by line cache control, thus it performs WBINVD instruction. The P7(Pentium4) family processors support the line by line cache control with CLFLUSH instruction, thus flushes the specified cache line.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the cache type is invalid or the cache control is not supported.
ERRNO	Not Available
SEE ALSO	cacheArchLib

#### cacheArchLibInit( )

cacheArchLibInit() – initialize the cache library NAME SYNOPSIS STATUS cacheArchLibInit ( CACHE MODE instMode, /\* instruction cache mode \*/ CACHE\_MODE dataMode /\* data cache mode \*/ ) This routine initializes the cache library for the following processor cache families: DESCRIPTION Motorola 68K, Intel x86, PowerPC ARM, and the Solaris and Windows simulators. It initializes the function pointers and configures the caches to the specified cache modes. 68K PROCESSORS The caching modes vary for members of the 68K processor family: CACHE WRITETHROUGH 68020: (instruction cache only) 68030: CACHE\_WRITETHROUGH CACHE\_BURST\_ENABLE CACHE\_BURST\_DISABLE CACHE\_WRITEALLOCATE (data cache only) CACHE\_NO\_WRITEALLOCATE (data cache only) 68040: CACHE\_WRITETHROUGH CACHE COPYBACK (data cache only) CACHE\_INH\_SERIAL (data cache only) CACHE\_INH\_NONSERIAL (data cache only) CACHE BURST ENABLE (data cache only) CACHE\_NO\_WRITEALLOCATE (data cache only) 68060: CACHE\_WRITETHROUGH CACHE COPYBACK (data cache only) CACHE\_INH\_PRECISE (data cache only) CACHE\_INH\_IMPRECISE (data cache only) CACHE\_BURST\_ENABLE (data cache only) The write-through, copy-back, serial, non-serial, precise and non precise modes change the state of the data transparent translation register (DTTR0) CM bits. Only DTTR0 is modified, since it typically maps DRAM space.

**X86 PROCESSORS** The caching mode **CACHE\_WRITETHROUGH** is available for the 486 family processors. The caching mode **CACHE\_COPYBACK** becomes available for the P5(Pentium) family processors. The caching mode **(CACHE\_COPYBACK | CACHE\_SNOOP\_ENABLE**) becomes available for the P6(PentiumPro, II, III) and P7(Pentium4) family processors.

#### POWER PC PROCESSORS

Modes should be set before caching is enabled. If two contradictory flags are set (for example, enable/disable), no action is taken for any of the input flags.

VxWorks Kernel API Reference, 6.6 cacheAuLibInit()

ARM PROCESSORS	The caching capabilities and modes vary for members of the ARM processor family. All caches are provided on-chip, so cache support is mostly an architecture issue, not a BSP issue. However, the memory map is BSP-specific and some functions need knowledge of the memory map, so they have to be provided in the BSP. ARM7TDMI (In ARM or Thumb state) No cache or MMU at all. Dummy routine provided, so that
	ARM710A Combined instruction and data cache. Actually a write-through cache, but separate write-buffer effectively makes this a copy-back cache if the write-buffer is enabled. Use write-through/copy-back argument to decide whether to enable write buffer. Data and instruction cache modes must be identical.
	ARM810 Combined instruction and data cache. Write-through and copy-back cache modes, but separate write-buffer effectively makes even write-through a copy-back cache as all writes are buffered, when cache is enabled. Data and instruction cache modes must be identical.
	ARMSA110 Separate instruction and data caches. Write-through and copy-back cache mode for data, but separate write-buffer effectively makes even write-through a copy-back cache as all writes are buffered, when cache is enabled.
RETURNS	ОК
ERRNO	Not Available
SEE ALSO	cacheArchLib

# cacheAuLibInit()

NAME	cacheAuLibInit(	) – initialize the Au	cac	he library
SYNOPSIS	STATUS cacheAul ( CACHE_MODE CACHE_MODE UINT32 UINT32 UINT32 UINT32 )	LibInit instMode, dataMode, iCacheSize, iCacheLineSize, dCacheSize, dCacheLineSize	/* /*	instruction cache mode */ data cache mode */

DESCRIPTION	This routine initializes the function pointers for the Au cache library. The board support
	package can select this cache library by assigning the function pointer sysCacheLibInit to
	cacheAuLibInit( ).

ERRNO	Not Available

SEE ALSO cacheAuLib

## cacheClear()

NAME	cacheClear() – clear all or some entries from a cache
SYNOPSIS	STATUS cacheClear ( CACHE_TYPE cache, /* cache to clear */ void * address, /* virtual address */ size_t bytes /* number of bytes to clear */ )
DESCRIPTION	This routine flushes and invalidates all or some entries in the specified cache.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the cache type is invalid or the cache control is not supported.
ERRNO	S_cacheLib_INVALID_CACHE the cache type specified is invalid.
SEE ALSO	cacheLib

## cacheDisable()

NAME	<pre>cacheDisable() – disable the specified cache</pre>	
SYNOPSIS	STATUS cacheDisable ( CACHE_TYPE cache /* cache to disable */ )	
DESCRIPTION	This routine flushes the cache and disables the instruction or data cache.	

VxWorks Kernel API Reference, 6.6 cacheDmaFree()

 

 RETURNS
 OK, or ERROR if the cache type is invalid or the cache control is not supported.

 ERRNO
 S\_cacheLib\_INVALID\_CACHE the cache type specified is invalid.

 SEE ALSO
 cacheLib

)

#### cacheDmaFree()

NAME	<pre>cacheDmaFree( ) - free the buffer acquired with cacheDmaMalloc(</pre>
SYNOPSIS	STATUS cacheDmaFree ( void * pBuf /* pointer to malloc/free buffer */ )
DESCRIPTION	This routine frees the buffer returned by <b>cacheDmaMalloc()</b> .
RETURNS	<b>OK</b> , or <b>ERROR</b> if the cache control is not supported.
ERRNO	N/A
SEE ALSO	cacheLib

## cacheDmaMalloc()

NAME	cacheDmaMalloc() – allocate a cache-safe buffer for DMA devices and drivers
SYNOPSIS	<pre>void * cacheDmaMalloc    (     size_t bytes /* number of bytes to allocate */ )</pre>
DESCRIPTION	This routine returns a pointer to a section of memory that will not experience any cache coherency problems. Function pointers in the CACHE_FUNCS structure provide access to DMA support routines.
RETURNS	A pointer to the cache-safe buffer, or NULL.
ERRNO	N/A

SEE ALSO cacheLib

#### cacheDrvFlush()

NAME cacheDrvFlush() – flush the data cache for drivers

SYNOPSIS STATUS cacheDrvFlush

```
(
CACHE_FUNCS * pFuncs, /* pointer to CACHE_FUNCS */
void * address, /* virtual address */
size_t bytes /* number of bytes to flush */
)
```

**DESCRIPTION** This routine flushes the data cache entries using the function pointer from the specified set.

**RETURNS** OK, or ERROR if the cache control is not supported.

ERRNO N/A

SEE ALSO cacheLib

#### cacheDrvInvalidate()

NAME	cacheDrvInvalidate() – invalidate data cache for drivers	
SYNOPSIS	<pre>STATUS cacheDrvInvalidate   (    CACHE_FUNCS * pFuncs, /* pointer to CACHE_FUNCS */    void * address, /* virtual address */    size_t bytes /* no. of bytes to invalidate */   )</pre>	
DESCRIPTION	This routine invalidates the data cache entries using the function pointer from the specified set.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if the cache control is not supported.	
ERRNO	N/A	
SEE ALSO	cacheLib	

# cacheDrvPhysToVirt( )

NAME	<pre>cacheDrvPhysToVirt() - translate a physical address for drivers</pre>
SYNOPSIS	<pre>void * cacheDrvPhysToVirt   (    CACHE_FUNCS * pFuncs, /* pointer to CACHE_FUNCS */    void * address /* physical address */   )</pre>
DESCRIPTION	This routine performs a physical-to-virtual address translation using the function pointer from the specified set.
RETURNS	The virtual address that maps to the physical address argument.
ERRNO	N/A
SEE ALSO	cacheLib

# cacheDrvVirtToPhys()

NAME	cacheDrvVirtToPhys() – translate a virtual address for drivers	
SYNOPSIS	<pre>void * cacheDrvVirtToPhys   (    CACHE_FUNCS * pFuncs, /* pointer to CACHE_FUNCS */    void * address /* virtual address */   )</pre>	
DESCRIPTION	This routine performs a virtual-to-physical address translation using the function pointer from the specified set.	
RETURNS	The physical address translation of a virtual address argument.	
ERRNO	N/A	
SEE ALSO	cacheLib	

## cacheEnable()

NAME	cacheEnable() – enable the specified cache
SYNOPSIS	STATUS cacheEnable ( CACHE_TYPE cache /* cache to enable */ )
DESCRIPTION	This routine invalidates the cache tags and enables the instruction or data cache.
RETURNS	OK, or ERROR if the cache type is invalid or the cache control is not supported.
ERRNO	S_cacheLib_INVALID_CACHE the cache type specified is invalid.
SEE ALSO	cacheLib

# cacheFlush()

NAME	<b>cacheFlush()</b> – flush all or some of a specified cache	
SYNOPSIS	<pre>STATUS cacheFlush   (    CACHE_TYPE cache, /* cache to flush */    void * address, /* virtual address */    size_t bytes /* number of bytes to flush */   )</pre>	
DESCRIPTION	This routine flushes (writes to memory) all or some of the entries in the specified cache. Depending on the cache design, this operation may also invalidate the cache tags. For write-through caches, no work needs to be done since RAM already matches the cached entries. Note that write buffers on the chip may need to be flushed to complete the flush	
RETURNS	<b>OK</b> , or <b>ERROR</b> if the cache type is invalid or the cache control is not supported.	
ERRNO	S_cacheLib_INVALID_CACHE the cache type specified is invalid.	
SEE ALSO	cacheLib	

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# cacheForeignClear()

NAME	cacheForeignClear() – clear foreign data from selected cache	
SYNOPSIS	<pre>STATUS cacheForeignClear (     CACHE_TYPE cache, /* cache to clear */     VIRT_ADDR virtAddr, /* virtual address */     PHYS_ADDR physAddr, /* physical address */     size_t bytes /* number of bytes to flush */ )</pre>	
DESCRIPTION	This routine performs a clear of the requested area of memory from the cache. Unlike <b>cacheClear()</b> , this routine does not assume that the provided virtual address is valid within the current address space. The called routine may clear more data than is requested.	
RETURNS	OK, or ERROR if the requested cache operation failed.	
ERRNO	S_cacheLib_INVALID_CACHE the cache type specified is invalid.	
SEE ALSO	cacheLib	

# cacheForeignFlush()

NAME	cacheForeignFlush() – flush foreign data from selected cache	
SYNOPSIS	<pre>STATUS cacheForeignFlush   (    CACHE_TYPE cache, /* cache to flush */    VIRT_ADDR virtAddr, /* virtual address */    PHYS_ADDR physAddr, /* physical address */    size_t bytes /* number of bytes to flush */   )</pre>	
DESCRIPTION	This routine performs a flush of the requested area of memory from the cache. Unlike <b>cacheFlush()</b> , this routine does not assume that the provided virtual address is valid withir the current address space. This routine may flush more data than is requested, in order to ensure that the required data has been flushed from the cache.	
RETURNS	OK, or ERROR if the requested cache operation failed.	

ERRNO S\_cacheLib\_INVALID\_CACHE the cache type specified is invalid.

SEE ALSO cacheLib

#### cacheForeignInvalidate()

NAME cacheForeignInvalidate() – invalidate foreign data from selected cache

- **DESCRIPTION** This routine performs an invalidate of the requested area of memory from the cache. Unlike **cacheInvalidate()**, this routine does not assume that the provided virtual address is valid within the current address space. Unlike the flush and clear functions, it is a programming error to invalidate more data from the cache than is requested. For this reason, if the arhitecture does not provide its own foreign invalidation routine, this function emulates the operation using **cacheClear()**.
- **RETURNS** OK, or **ERROR** if the requested cache operation failed.

ERRNO S\_cacheLib\_INVALID\_CACHE the cache type specified is invalid.

SEE ALSO cacheLib

#### cacheInvalidate()

NAME cacheInvalidate() – invalidate all or some of a specified cache

SYNOPSIS	STATUS cacheInvalidate (	
	CACHE_TYPE cache, void * address, size_t bytes )	/* cache to invalidate */ /* virtual address */ /* number of bytes to invalidate */

VxWorks Kernel API Reference, 6.6 cacheLibInit()

DESCRIPTION	This routine invalidates all or some of the entries in the specified cache. Depending on the cache design, the invalidation may be similar to the flush, or one may invalidate the tags directly.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the cache type is invalid or the cache control is not supported.
ERRNO	S_cacheLib_INVALID_CACHE the cache type specified is invalid.
SEE ALSO	cacheLib

# cacheLibInit( )

NAME	<b>cacheLibInit()</b> – initialize the cache library for a processor architecture	
SYNOPSIS	STATUS cacheLibInit ( CACHE_MODE instMode, /* inst cache mode */ CACHE_MODE dataMode /* data cache mode */ )	
DESCRIPTION	This routine initializes the function pointers for the appropriate cache library. For architectures with more than one cache implementation, the board support package must select the appropriate cache library with <b>sysCacheLibInit</b> . Systems without cache coherency problems (i.e., bus snooping) should NULLify the flush and invalidate function pointers in the <b>cacheLib</b> structure to enhance driver and overall system performance. This can be done in <b>sysHwInit()</b> .	
RETURNS	<b>OK</b> , or <b>ERROR</b> if there is no cache library installed.	
ERRNO	N/A	
SEE ALSO	cacheLib	

## cacheLock()

NAME cacheLock() – lock all or part of a specified cache

SYNOPSIS STATUS cacheLock

( CACHE\_TYPE cache, /\* cache to lock \*/ void \* address, /\* virtual address \*/ size t bytes /\* number of bytes to lock \*/ ) DESCRIPTION This routine locks all (global) or some (local) entries in the specified cache. Cache locking is useful in real-time systems. Not all caches can perform locking. RETURNS **OK**, or **ERROR** if the cache type is invalid or the cache control is not supported. ERRNO S\_cacheLib\_INVALID\_CACHE the cache type specified is invalid. cacheLib SEE ALSO

#### cachePipeFlush()

NAME cachePipeFlush() – flush processor write buffers to memory

- SYNOPSIS STATUS cachePipeFlush (void)
- **DESCRIPTION** This routine forces the processor output buffers to write their contents to RAM. A cache flush may have forced its data into the write buffers, then the buffers need to be flushed to RAM to maintain coherency.
- **RETURNS** OK, or ERROR if the cache control is not supported.

ERRNO N/A

SEE ALSO cacheLib

#### cacheR10kLibInit( )

NAME cacheR10kLibInit() - initialize the R10000 cache library
SYNOPSIS STATUS cacheR10kLibInit
(
CACHE\_MODE instMode, /\* instruction cache mode \*/
CACHE\_MODE dataMode, /\* data cache mode \*/

VxWorks Kernel API Reference, 6.6 cacheR4kLibInit()

UINT32	iCacheSize,
UINT32	iCacheLineSize,
UINT32	dCacheSize,
UINT32	dCacheLineSize,
UINT32	sCacheSize,
UINT32	sCacheLineSize
)	

**DESCRIPTION** This routine initializes the function pointers for the R10000 cache library. The board support package can select this cache library by assigning the function pointer *sysCacheLibInit* to **cacheR10kLibInit()**.

RETURNS OK.

ERRNO Not Available

SEE ALSO cacheR10kLib

## cacheR4kLibInit( )

NAME	cacheR4kLibInit() – initialize the R4000 cache library	
SYNOPSIS	<pre>STATUS cacheR4kLibInit   (    CACHE_MODE instMode, /* instruction cache mode */    CACHE_MODE dataMode, /* data cache mode */    UINT32 iCacheSize,    UINT32 iCacheLineSize,    UINT32 dCacheSize,    UINT32 dCacheLineSize,    UINT32 sCacheSize,    UINT32 sCacheSize,    ) </pre>	
DESCRIPTION	This routine initializes the function pointers for the R4000 cache library. The board support package can select this cache library by assigning the function pointer <i>sysCacheLibInit</i> to <b>cacheR4kLibInit()</b> .	
RETURNS	OK.	
ERRNO	Not Available	
SEE ALSO	cacheR4kLib	

## cacheR5kLibInit( )

NAME	cacheR5kLibInit() – initialize the R5000 cache library
SYNOPSIS	<pre>STATUS cacheR5kLibInit   (   CACHE_MODE instMode, /* instruction cache mode */   CACHE_MODE dataMode, /* data cache mode */   UINT32 iCacheSize,   UINT32 iCacheLineSize,   UINT32 dCacheSize,   UINT32 dCacheLineSize,   UINT32 sCacheSize,   UINT32 sCacheLineSize   )</pre>
DESCRIPTION	This routine initializes the function pointers for the R5000 cache library. The board support package can select this cache library by assigning the function pointer <i>sysCacheLibInit</i> to <b>cacheR5kLibInit()</b> .
RETURNS	OK.
ERRNO	Not Available
SEE ALSO	cacheR5kLib

## cacheR7kLibInit( )

NAME	cacheR7kLibInit() -	- initialize the R7000 cache library	
------	---------------------	--------------------------------------	--

SYNOPSIS	STATUS cacheR7kLibInit			
	(			
	CACHE_MODE	instMode,	/* instruction cache mode */	
	CACHE_MODE	dataMode,	/* data cache mode */	
	UINT32	iCacheSize,		
	UINT32	iCacheLineSize,		
	UINT32	dCacheSize,		
	UINT32	dCacheLineSize,		
	UINT32	sCacheSize,		
	UINT32	sCacheLineSize,		
	UINT32	tCacheSize,		
	UINT32	tCacheLineSize		
	)			

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VxWorks Kernel API Reference, 6.6 cacheSh7750LibInit()

**DESCRIPTION** This routine initializes the function pointers for the R7000 cache library. The board support package can select this cache library by assigning the function pointer *sysCacheLibInit* to **cacheR7kLibInit()**.

RETURNS OK.

ERRNO Not Available

SEE ALSO cacheR7kLib

#### cacheSh7750LibInit( )

NAME cacheSh7750LibInit() - initialize the SH7750 cache library
SYNOPSIS STATUS cacheSh7750LibInit
(
CACHE\_MODE instMode, /\* instruction cache mode \*/
CACHE\_MODE dataMode /\* data cache mode \*/
)

**DESCRIPTION** This routine initializes the cache library for the Renesas SH7750 processor. It initializes the function pointers and configures the caches to the specified cache modes. Modes should be set before caching is enabled. If two complementary flags are set (enable/disable), no action is taken for any of the input flags.

The following caching modes are available for the SH7750, SH7750R and SH7770 processors:

SH7750 :	CACHE_WRITETHROUGH	
	CACHE_COPYBACK	(copy-back cache for P0/P3, data cache only)
	CACHE_COPYBACK_P1	(copy-back cache for P1, data cache only)
	CACHE_RAM_MODE	(use half of cache as RAM, data cache only)
	CACHE_A25_INDEX	(use A25 as MSB of cache index)
	CACHE_DMA_BYPASS_P0	(allocate DMA buffer to P2, free it to P0)
	CACHE_DMA_BYPASS_P1	(allocate DMA buffer to P2, free it to P1)
	CACHE_DMA_BYPASS_P3	(allocate DMA buffer to P2, free it to P3)

SH7750R: CACHE\_WRITETHROUGH

# 2 Routines cacheSh7750LibInit()

		CACHE_COPYBACK	(copy-back cache for P0/P3, data cache only)
		CACHE_COPYBACK_P1	(copy-back cache for P1, data cache only)
		CACHE_RAM_MODE	(use half of cache as RAM, data cache only)
		CACHE_2WAY_MODE	(use RAM in 2way associ. mode, data cache only)
		CACHE A25 INDEX	(use A25 as MSB of cache index)
		CACHE_DMA_BYPASS_P0	(allocate DMA buffer to P2, free it to P0)
		CACHE_DMA_BYPASS_P1	(allocate DMA buffer to P2, free it to P1)
		CACHE_DMA_BYPASS_P3	(allocate DMA buffer to P2, free it to P3)
	SH7770 :	CACHE_SH4A_MODE Cache Writethrough	(SH4A cache support)
		CACHE_COPYBACK	(copy-back cache for P0/P3, data cache only)
		CACHE_COPYBACK_P1	(copy-back cache for P1, data cache only)
		CACHE_2WAY_MODE	(use RAM in 2way associ. mode)
		CACHE_DMA_BYPASS_P0	(allocate DMA buffer to P2, free it to P0)
		CACHE_DMA_BYPASS_P1	(allocate DMA buffer to P2, free it to P1)
		CACHE_DMA_BYPASS_P3	(allocate DMA buffer to P2, free it to P3)
	The CACHE_DM. MMU. If none of returns a cache-sa CACHE_DMA_BYI space, and cacheE partition is located selecting CACHE_	A_BYPASS_Px modes allow to al CACHE_DMA_BYPASS_Px mod fe buffer on logical space, which <b>PASS_P0</b> is selected, <b>cacheDmaM</b> <b>DmaFree(</b> ) releases the buffer to l d on P0, cache-safe buffers can be <b>DMA_BYPASS_P0</b> .	locate "cache-safe" buffers without des is specified, <b>cacheDmaMalloc()</b> is created by the MMU. If <b>alloc()</b> returns a cache-safe buffer on P2 P0 space. Namely, if the system memory allocated and freed without MMU, by
RETURNS	<b>OK</b> , or <b>ERROR</b> if s	pecified cache mode is invalid.	
ERRNO			
SEE ALSO	cacheSh7750Lib		

VxWorks Kernel API Reference, 6.6 cacheStoreBufDisable()

## cacheStoreBufDisable()

NAME	<pre>cacheStoreBufDisable() – disable the store buffer (MC68060 only)</pre>
SYNOPSIS	void cacheStoreBufDisable (void)
DESCRIPTION	This routine resets the ESB bit of the Cache Control Register (CACR) to disable the store buffer.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	cacheArchLib

## cacheStoreBufEnable()

NAME	cacheStoreBufEnable() – enable the store buffer (MC68060 only)
SYNOPSIS	void cacheStoreBufEnable (void)
DESCRIPTION	This routine sets the ESB bit of the Cache Control Register (CACR) to enable the store buffer. To maximize performance, the four-entry first-in-first-out (FIFO) store buffer is used to defer pending writes to write through or cache-inhibited imprecise pages.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	cacheArchLib

# cacheTextLocalUpdate()

NAME	<b>cacheTextLocalUpdate()</b> – synchronize the caches on local cpu only		
SYNOPSIS	STATUS cacheTextLocalUpdate		
	void * address, /* virtual address */		

```
      size_t bytes
      /* number of bytes to sync */

      DESCRIPTION
      This routine flushes the data cache, then invalidates the instruction cache. This operation forces the instruction cache to fetch code that may have been created via the data path. The operation is limited to the local CPU (i.e., no CPC is performed).

      RETURNS
      OK, or ERROR if the cache control is not supported.

      ERRNO
      N/A

      SEE ALSO
      cacheLib
```

#### cacheTextUpdate( )

NAME	cacheTextUpdate() – synchronize the instruction and data caches
SYNOPSIS	<pre>STATUS cacheTextUpdate   (    void * address, /* virtual address */    size_t bytes /* number of bytes to sync */   )</pre>
DESCRIPTION	This routine flushes the data cache, then invalidates the instruction cache. This operation forces the instruction cache to fetch code that may have been created via the data path.
RETURNS	OK, or ERROR if the cache control is not supported.
ERRNO	N/A
SEE ALSO	cacheLib

#### cacheTx49LibInit( )

NAME cacheTx49LibInit() - initialize the Tx49 cache library
SYNOPSIS STATUS cacheTx49LibInit
(
CACHE\_MODE instMode, /\* instruction cache mode \*/
CACHE\_MODE dataMode, /\* data cache mode \*/

VxWorks Kernel API Reference, 6.6 cacheUnlock()

UINT32	iCacheSize,	/*	instruction cache size */
UINT32	iCacheLineSize,	/*	instruction cache line size */
UINT32	dCacheSize,	/*	data cache size */
UINT32	dCacheLineSize	/*	data cache line size */
)			

**DESCRIPTION** This routine initializes the function pointers for the Tx49 cache library. The board support package can select this cache library by assigning the function pointer *sysCacheLibInit* to **cacheTx49LibInit()**.

RETURNS OK.

ERRNO Not Available

SEE ALSO cacheTx49Lib

## cacheUnlock()

NAME	cacheUnlock() – unlock all or part of a specified cache
SYNOPSIS	<pre>STATUS cacheUnlock   (    CACHE_TYPE cache, /* cache to unlock */    void * address, /* virtual address */    size_t bytes /* number of bytes to unlock */   )</pre>
DESCRIPTION	This routine unlocks all (global) or some (local) entries in the specified cache. Not all caches can perform unlocking.
RETURNS	OK, or ERROR if the cache type is invalid or the cache control is not supported.
ERRNO	S_cacheLib_INVALID_CACHE the cache type specified is invalid.
SEE ALSO	cacheLib
# calloc()

NAME	<b>calloc( )</b> – allocate space for an array (ANSI)
SYNOPSIS	<pre>void * calloc   (    size_t elemNum, /* number of elements */    size_t elemSize /* size of elements */   )</pre>
DESCRIPTION	This routine allocates a block of memory for an array that contains <i>elemNum</i> elements of size <i>elemSize</i> . This space is initialized to zeros.
RETURNS	A pointer to the block, or NULL if the call fails.
ERRNO	Possible errnos generated by this routine include:
	<b>S_memLib_NOT_ENOUGH_MEMORY</b> There is no free block large enough to satisfy the allocation request.
SEE ALSO	<b>memPartLib</b> , American National Standard for Information Systems -, Programming Language - C, ANSI X3.159-1989: General Utilities ( <b>stdlib.h</b> )

# cbioBlkCopy()

NAME	cbioBlkCopy() – block to block (sector to sector) transfer routine
SYNOPSIS	<pre>STATUS cbioBlkCopy   (    CBIO_DEV_ID dev, /* CBIO handle */    block_t srcBlock, /* source start block */    block_t dstBlock, /* destination start block */    block_t numBlocks /* number of blocks to copy */   )</pre>
DESCRIPTION	This routine verifies the CBIO device is valid and if so calls the devices block to block transfer routine which makes copies of one or more blocks on the lower layer (hardware, subordinate CBIO, or <b>BLK_DEV</b> ). It is optimized for block to block copies on the subordinate layer.

If the CBIO\_DEV\_ID passed to this routine is not a valid CBIO handle, ERROR will be returned with errno set to S\_cbioLib\_INVALID\_CBIO\_DEV\_ID

VxWorks Kernel API Reference, 6.6 cbioBlkRW()

**RETURNS** OK if successful or ERROR if the handle is invalid, or if the CBIO device routine returns ERROR.

ERRNO Not Available

SEE ALSO cbioLib

## cbioBlkRW()

NAME	cbioBlkRW() – transfer blocks to or from memory
SYNOPSIS	<pre>STATUS cbioBlkRW   (    CBIO_DEV_ID dev, /* CBIO handle */    block_t startBlock, /* starting block of transfer */    block_t numBlocks, /* number of blocks to transfer */    addr_t buffer, /* address of the memory buffer */    CBIO_RW rw, /* direction of transfer R/W */    cookie_t * pCookie /* pointer to cookie */    )</pre>
DESCRIPTION	This routine verifies the CBIO device is valid and if so calls the devices block transfer routine. The CBIO device performs block transfers between the device and memory.
	If the <b>CBIO_DEV_ID</b> passed to this routine is not a valid CBIO handle, <b>ERROR</b> will be returned with errno set to <b>S_cbioLib_INVALID_CBIO_DEV_ID</b>
RETURNS	OK if successful or ERROR if the handle is invalid, or if the CBIO device routine returns ERROR.
ERRNO	Not Available
SEE ALSO	cbioLib

## cbioBytesRW()

NAME	<b>cbioBytesRW( )</b> – transfer by	ytes to or from memory
SYNOPSIS	STATUS cbioBytesRW (	
	CBIO_DEV_ID dev,	/* CBIO handle */

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	<pre>block_t startBlock, /* startEng block of the transfer */ off_t offset, /* offset into block in bytes */ addr_t buffer, /* address of data buffer */ size_t nBytes, /* number of bytes to transfer */ CBIO_RW rw, /* direction of transfer R/W */ cookie_t * pCookie /* pointer to cookie */ )</pre>
DESCRIPTION	This routine verifies the CBIO device is valid and if so calls the devices byte transfer routine which transfers between a user buffer and the lower layer (hardware, subordinate CBIO, or <b>BLK_DEV</b> ). It is optimized for byte transfers. If the <b>CBIO_DEV_ID</b> passed to this routine is not a valid CBIO handle, <b>ERROR</b> will be returned with errno set to <b>S_cbioLib_INVALID_CBIO_DEV_ID</b>
RETURNS	OK if successful or ERROR if the handle is invalid, or if the CBIO device routine returns ERROR.
ERRNO	Not Available
SEE ALSO	cbioLib

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# cbioDevCreate()

NAME	cbioDevCreate() – Initialize a CBIO device (Generic)
SYNOPSIS	<pre>CBIO_DEV_ID cbioDevCreate   (    caddr_t ramAddr, /* where it is in memory (0 = KHEAP_ALLOC) */    size_t ramSize /* pool size */   )</pre>
DESCRIPTION	This routine will create an empty <b>CBIO_DEV</b> structure and return a handle to that structure ( <b>CBIO_DEV_ID</b> ). This routine is intended to be used by CBIO modules only. See <b>cbioLibP.h</b>
RETURNS	CBIO_DEV_ID or NULL if ERROR.
ERRNO	Not Available
SEE ALSO	cbioLib

VxWorks Kernel API Reference, 6.6 cbioDevVerify()

# cbioDevVerify()

NAME	<pre>cbioDevVerify() - verify CBIO_DEV_ID</pre>
SYNOPSIS	STATUS cbioDevVerify ( CBIO_DEV_ID device /* CBIO_DEV_ID to be verified */ )
DESCRIPTION	The purpose of this function is to determine if the device complies with the CBIO interface. It can be used to verify a CBIO handle before it is passed to <b>dosFsLib</b> , <b>rawFsLib</b> , <b>usrFdiskPartLib</b> , or other CBIO modules which expect a valid CBIO interface.
	The device handle provided to this function, <i>device</i> is verified to be a CBIO device. If <i>device</i> is not a CBIO device <b>ERROR</b> is returned with errno set to <b>S_cbioLib_INVALID_CBIO_DEV_ID</b>
	The dcacheCbio and dpartCbio CBIO modules (and <b>dosFsLib</b> ) use this function internally, and therefore this function need not be otherwise invoked when using compliant CBIO modules.
RETURNS	<b>OK</b> or <b>ERROR</b> if not a CBIO device, if passed a <b>NULL</b> address, or if the check could cause an unaligned access.
ERRNO	Not Available
SEE ALSO	cbioLib, dosFsLib, dcacheCbio(), dpartCbio()

### cbioIoctl()

NAME	<b>cbioIoctl()</b> – perform ioctl operation on device
SYNOPSIS	<pre>STATUS cbioloctl   (    CBIO_DEV_ID dev, /* CBIO handle */    int command, /* ioctl command to be issued */    addr_t arg /* arg - specific to ioctl */   )</pre>
DESCRIPTION	This routine verifies the CBIO device is valid and if so calls the devices $I/O$ control operation routine.
	CBIO modules expect the following <b>ioctl()</b> codes:

	- CBIO_RESET - reset the CBIO device. When the third argument to the locil call accompanying CBIO_RESET is NULL, the code verifies that the disk is inserted and is ready, after getting it to a known state. When the 3rd argument is a non-zero, it is assumed to be a BLK_DEV pointer and CBIO_RESET will install a new subordinate block device. This work is performed at the BLK_DEV to CBIO layer, and all layers shall account for it. A CBIO_RESET indicates a possible change in device geometry, and the CBIO_PARAMS members will be reinitialized after a CBIO_RESET.
	- CBIO_STATUS_CHK - check device status of CBIO device and lower layer
	- CBIO_DEVICE_LOCK - Prevent disk removal
	- CBIO_DEVICE_UNLOCK - Allow disk removal
	- CBIO_DEVICE_EJECT - Unmount and eject device
	- CBIO_CACHE_FLUSH - Flush any dirty cached data
	- CBIO_CACHE_INVAL - Flush & Invalidate all cached data
	- CBIO_CACHE_NEWBLK - Allocate scratch block
	If the CBIO_DEV_ID passed to this routine is not a valid CBIO handle, ERROR will be returned with errno set to S_cbioLib_INVALID_CBIO_DEV_ID
RETURNS	OK if successful or ERROR if the handle is invalid, or if the CBIO device routine returns ERROR.
ERRNO	Not Available
SEE ALSO	cbioLib

\* 4 \*\*

# cbioLibInit( )

NAME cbioLibInit() – Initialize CBIO Library

SYNOPSIS STATUS cbioLibInit(void)

**DESCRIPTION** This function initializes the CBIO library, and will be called when the first CBIO device is created, hence it does not need to be called during system initialization. It can be called multiple times, but will do nothing after the first call.

RETURNS OK or ERROR

ERRNO Not Available

VxWorks Kernel API Reference, 6.6 cbioLock()

SEE ALSO cbioLib

# cbioLock()

NAME	<b>cbioLock( )</b> – obtain CBIO device semaphore.
SYNOPSIS	<pre>STATUS cbioLock   (    CBIO_DEV_ID dev, /* CBIO handle */    int timeout /* timeout in ticks */ )</pre>
DESCRIPTION	If the CBIO_DEV_ID passed to this routine is not a valid CBIO handle, ERROR will be returned with errno set to S_cbioLib_INVALID_CBIO_DEV_ID
RETURNS	OK or ERROR if the CBIO handle is invalid or semTake fails.
ERRNO	Not Available
SEE ALSO	cbioLib

### cbioModeGet( )

NAME	<b>cbioModeGet()</b> – return the mode setting for CBIO device
SYNOPSIS	<pre>int cbioModeGet   (    CBIO_DEV_ID dev /* CBIO handle */ )</pre>
DESCRIPTION	If the <b>CBIO_DEV_ID</b> passed to this routine is not a valid CBIO handle, <b>ERROR</b> will be returned with errno set to <b>S_cbioLib_INVALID_CBIO_DEV_ID</b> This routine is not protected by a semaphore.
	This routine confirms if the current layer is a CBIO to BLKDEV wrapper or a CBIO to CBIO layer. Depending on the current layer it either returns the mode from <b>BLK_DEV</b> or calls <b>cbioModeGet()</b> recursively.
RETURNS	O_RDONLY, O_WRONLY, or O_RDWR or ERROR
ERRNO	Not Available

SEE ALSO cbioLib

### cbioModeSet()

NAME	<b>cbioModeSet( )</b> – set mode for CBIO device
SYNOPSIS	<pre>STATUS cbioModeSet   (    CBIO_DEV_ID dev, /* CBIO handle */    int mode /* O_RDONLY, O_WRONLY, or O_RDWR */   )</pre>
DESCRIPTION	Valid modes are O_RDONLY, O_WRONLY, or O_RDWR.
	If the <b>CBIO_DEV_ID</b> passed to this routine is not a valid CBIO handle, <b>ERROR</b> will be returned with errno set to <b>S_cbioLib_INVALID_CBIO_DEV_ID</b> This routine is not protected by a semaphore.
	This routine confirms if the current layer is a CBIO to BLKDEV wrapper or a CBIO to CBIO layer. Depending on the current layer it either sets the mode of the <b>BLK_DEV</b> or calls <b>cbioModeSet()</b> recursively.
RETURNS	<b>OK</b> or <b>ERROR</b> if mode is not set.
ERRNO	Not Available
SEE ALSO	cbioLib

# cbioParamsGet()

NAME	cbioParamsGet() – fill in CBIO_PARAMS structure with CBIO device parameters
SYNOPSIS	STATUS cbioParamsGet ( CBIO_DEV_ID dev, /* CBIO handle */ CBIO_PARAMS * pCbioParams /* pointer to CBIO_PARAMS */ )
DESCRIPTION	If the <b>CBIO_DEV_ID</b> passed to this routine is not a valid CBIO handle, <b>ERROR</b> will be returned with errno set to <b>S cbioLib INVALID CBIO DEV ID</b>

VxWorks Kernel API Reference, 6.6 cbioRdyChgdGet()

**RETURNS** OK or ERROR if the CBIO handle is invalid.

ERRNO Not Available

SEE ALSO cbioLib

#### cbioRdyChgdGet()

cbioRdyChgdGet() - determine ready status of CBIO device NAME SYNOPSIS int cbioRdyChgdGet ( CBIO\_DEV\_ID dev /\* CBIO handle \*/ ) For example DESCRIPTION switch (cbioRdyChgdGet (cbioDeviceId)) { case TRUE: printf ("Disk changed.\n"); break; case FALSE: printf ("Disk has not changed.\n"); break; case ERROR: printf ("Not a valid CBIO device.\n"); break; default: break; } If the CBIO\_DEV\_ID passed to this routine is not a valid CBIO handle, ERROR will be returned with errno set to S\_cbioLib\_INVALID\_CBIO\_DEV\_ID This routine is not protected by a semaphore. This routine will check down to the driver layer to see if any lower layer has its ready changed bit set to TRUE. If so, this routine returns TRUE. If no lower layer has its ready changed bit set to TRUE, this layer returns FALSE. TRUE if device ready status has changed, else FALSE if the ready status has not changed, RETURNS else ERROR if the CBIO\_DEV\_ID is invalid.

ERRNO Not Available

SEE ALSO cbioLib

#### cbioRdyChgdSet()

cbioRdyChgdSet() - force a change in ready status of CBIO device NAME SYNOPSIS STATUS cbioRdyChgdSet ( CBIO\_DEV\_ID dev, /\* CBIO handle \*/ BOOL status /\* TRUE/FALSE \*/ ) Pass TRUE in status to force READY status change. DESCRIPTION If the CBIO\_DEV\_ID passed to this routine is not a valid CBIO handle, ERROR will be returned with errno set to S\_cbioLib\_INVALID\_CBIO\_DEV\_ID If status is not passed as TRUE or FALSE, ERROR is returned. This routine is not protected by a semaphore. This routine sets readyChanged bit of passed CBIO\_DEV. OK or ERROR if the device is invalid or status is not TRUE or FALSE. RETURNS ERRNO Not Available cbioLib SEE ALSO

#### cbioShow()

NAME	<b>cbioShow()</b> – print information about a CBIO device
SYNOPSIS	STATUS cbioShow ( CBIO_DEV_ID dev /* CBIO handle */ )
DESCRIPTION	This function will display on standard output all information which is generic for all CBIO devices. See the CBIO modules particular device show routines for displaying implementation-specific information.
	It takes two arguments:
	A <b>CBIO_DEV_ID</b> which is the CBIO handle to display or <b>NULL</b> for the most recent device.
RETURNS	OK or ERROR if no valid CBIO_DEV is found.
ERRNO	Not Available

VxWorks Kernel API Reference, 6.6 cbioUnlock()

SEE ALSO cbioLib, dcacheShow(), dpartShow()

### cbioUnlock()

NAME	<b>cbioUnlock( )</b> – release CBIO device semaphore.
SYNOPSIS	STATUS cbioUnlock ( CBIO_DEV_ID dev /* CBIO handle */ )
DESCRIPTION	If the CBIO_DEV_ID passed to this routine is not a valid CBIO handle, ERROR will be returned with errno set to S_cbioLib_INVALID_CBIO_DEV_ID
RETURNS	<b>OK</b> or <b>ERROR</b> if the CBIO handle is invalid or the semGive fails.
ERRNO	Not Available
SEE ALSO	cbioLib

# cbioWrapBlkDev()

NAME	cbioWrapBlkDev() – create CBIO wrapper atop a BLK_DEV device
SYNOPSIS	CBIO_DEV_ID cbioWrapBlkDev ( BLK_DEV * pDevice /* BLK_DEV * device pointer */ )
DESCRIPTION	The purpose of this function is to make a blkIo ( <b>BLK_DEV</b> ) device comply with the CBIO interface via a wrapper.
	The device handle provided to this function, <i>device</i> is verified to be a blkIo device. A lean CBIO to <b>BLK_DEV</b> wrapper is then created for a valid blkIo device. The returned <b>CBIO_DEV_ID</b> device handle may be used with <b>dosFsDevCreate()</b> , <b>dcacheDevCreate()</b> , and any other routine expecting a valid <b>CBIO_DEV_ID</b> handle.
	To verify a blkIo pointer we see that all mandatory functions are not NULL.
	Note that if a valid <b>CBIO_DEV_ID</b> is passed to this function, it will simply be returned without modification.

	The <b>dosFsLib</b> , dcacheCbio, and dpartCbio CBIO modules use this function internally, and therefore this function need not be otherwise invoked when using those CBIO modules.
RETURNS	a CBIO device pointer, or NULL if not a blkIo device
ERRNO	Not Available
SEE ALSO	cbioLib, dosFsLib, dcacheCbio(), dpartCbio()

# cbrt()

NAME	<b>cbrt( )</b> – compute a cube root
SYNOPSIS	<pre>double cbrt   (    double x /* value to compute the cube root of */   )</pre>
DESCRIPTION	This routine returns the cube root of $x$ in double precision.
RETURNS	The double-precision cube root of <i>x</i> .
ERRNO	Not Available
SEE ALSO	mathALib

# cbrtf()

NAME	<b>cbrtf()</b> – compute a cube root
SYNOPSIS	float cbrtf ( float x /* argument */ )
DESCRIPTION	This routine returns the cube root of $x$ in single precision.
RETURNS	The single-precision cube root of <i>x</i> .
ERRNO	Not Available

SEE ALSO mathALib

# cd()

NAME	<b>cd()</b> – change the default directory
SYNOPSIS	STATUS cd ( const char * name /* new directory name */ )
DESCRIPTION	This command sets the default directory to <i>name</i> . The default directory is a device name, optionally followed by a directory local to that device.
NOTE	This is a target resident function, which manipulates the target I/O system. It must be preceded with the @ letter if executed from the Host Shell (windsh), which has a built-in command of the same name that operates on the Host's I/O system.
	To change to a different directory, specify one of the following:
	- an entire path name with a device name, possibly followed by a directory name. The entire path name will be changed.
	- a directory name starting with a ~ or / or \$. The directory part of the path, immediately after the device name, will be replaced with the new directory name.
	- a directory name to be appended to the current default directory. The directory name will be appended to the current default directory.
	An instance of "" indicates one level up in the directory tree.
	Note that when accessing a remote file system via RSH or FTP, the VxWorks network device must already have been created using <b>netDevCreate()</b> .
WARNING	The <b>cd()</b> command does very little checking that <i>name</i> represents a valid path. If the path is invalid, <b>cd()</b> may return <b>OK</b> , but subsequent calls that depend on the default path will fail.
EXAMPLES	The following example changes the directory to device <b>/fd0</b> /:
	-> cd "/fd0/"
	This example changes the directory to device <b>wrs:</b> with the local directory <b>~leslie/target</b> :
	-> cd "wrs:~leslie/target"
	After the previous command, the following changes the directory to <b>wrs:~leslie/target/config</b> :
	-> cd "config"

	After the previous command, the following changes the directory to <b>wrs:~leslie/target/demo</b> :
	-> cd "/demo"
	After the previous command, the following changes the directory to <b>wrs:/etc</b> .
	-> cd "/etc"
	Note that ~ can be used only on network devices (RSH or FTP).
RETURNS	OK or ERROR.
RETURNS ERRNO	OK or ERROR. Not Available

## cdromFsDevCreate()

NAME	cdromFsDevCreate() – create a CD-ROM filesystem (cdromFs) I/O device.
SYNOPSIS	CDROM_VOL_DESC_ID cdromFsDevCreate ( char * devName, /* device name */ device_t device /* underlying block device handle */ )
DESCRIPTION	This routine creates an instance of a cdromFs device in the I/O system. As input, this function requires an eXtended Block Device (XBD) identifier (device_t) for the CD drive on which to create a cdromFs I/O device. Thus, <b>xxxXbdDevCreate()</b> , for example, should have already been called to create the XBD device. Alternatively, <b>xxxBlkDevCreate()</b> , for example, can be called to create a legacy <b>BLK_DEV</b> driver followed by <b>xbdBlkDevCreate()</b> to create an XBD wrapper around the <b>BLK_DEV</b> device.
RETURNS	CDROM_VOL_DESC_ID, or NULL if error.
ERRNO	S_memLib_NOT_ENOUGH_MEMORY
SEE ALSO	cdromFsLib, cdromFsInit()

#### cdromFsDevDelete()

NAME	cdromFsDevDelete() – delete a CD-ROM filesystem (cdromFs) I/O device
SYNOPSIS	STATUS cdromFsDevDelete ( CDROM_VOL_DESC_ID pVolDesc /* ptr to CDROM_VOL_DESC */ )
DESCRIPTION	This routine deletes the specified volume. This involves removing the "device" from the I/O system, and freeing all resources associated with the volume.
RETURNS	OK if specified volume was successfully deleted, otherwise ERROR
ERRNO	Not Available
SEE ALSO	cdromFsLib, cdromFsInit(), cdromFsDevCreate()

# cdromFsInit()

NAME	cdromFsInit() – initialize the VxWorks CD-ROM file system
SYNOPSIS	STATUS cdromFsInit ( UINT32 commonBufferSize /* common buffer size */ )
DESCRIPTION	This routine initializes the VxWorks CD-ROM file system. It is automatically called when the INCLUDE_CDROMFS component is configured into the system.
RETURNS	<b>OK</b> or <b>ERROR</b> , if driver can not be installed.
ERRNO	S_iosLib_DRIVER_GLUT
SEE ALSO	cdromFsLib, cdromFsDevCreate(), iosLib.h

cdromFsVersionDisplay()
-------------------------

NAME	cdromFsVersionDisplay() – display the cdromFs version number
SYNOPSIS	<pre>void cdromFsVersionDisplay   (    int level /* level of display, not used */ )</pre>
DESCRIPTION	This routine displays the cdromFs version number. This routine has been deprecated.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	cdromFsLib, cdromFsVersionNumGet( ), cdromFsVolConfigShow( )

#### cdromFsVersionNumGet()

NAME	cdromFsVersionNumGet() – return the cdromFs version number
SYNOPSIS	uint32_t cdromFsVersionNumGet ( void )
DESCRIPTION	This routine returns the cdromFs version number. This routine has been deprecated.
RETURNS	the cdromFs version number.
ERRNO	Not Available
SEE ALSO	cdromFsLib, cdromFsVersionDisplay()

# cdromFsVolConfigShow()

NAME cdromFsVolConfigShow() – show the volume configuration information

SYNOPSIS VOID cdromFsVolConfigShow

VxWorks Kernel API Reference, 6.6 ceilf()

```
(
void * arg /* device name or CDROM_VOL_DESC * */
)
```

# **DESCRIPTION** This routine retrieves the volume configuration for the named **cdromFsLib** device and prints it to standard output. The information displayed is retrieved from the **BLK\_DEV** structure for the specified device.

RETURNS N/A

ERRNO Not Available

SEE ALSO cdromFsLib, N/A

# ceilf()

NAME	<b>ceilf()</b> – compute the smallest integer greater than or equal to a specified value (ANSI)
SYNOPSIS	<pre>float ceilf   (    float v /* value to find the ceiling of */   )</pre>
DESCRIPTION	This routine returns the smallest integer greater than or equal to $v$ , in single precision.
RETURNS	The smallest integral value greater than or equal to $v$ , in single precision.
ERRNO	Not Available
SEE ALSO	mathALib

#### cfree()

NAME	<b>cfree()</b> – free a block of memory from the system memory partition (kernel heap)
SYNOPSIS	STATUS cfree ( char * pBlock /* pointer to block of memory to free */ )

DESCRIPTION	This routine returns to the free memory pool a block of memory previously allocated with <b>calloc()</b> .
	It is an error to free a memory block that was not previously allocated.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the block is invalid.
ERRNO	Possible errnos generated by this routine include:
	S_memLib_BLOCK_ERROR The block of memory to free is not valid.
SEE ALSO	memPartLib

## checkStack()

NAME	<b>checkStack()</b> – print a summary of each task's stack usage
SYNOPSIS	<pre>void checkStack   (    int taskNameOrId /* task name or task ID; 0 = summarize all */ )</pre>
DESCRIPTION	This command displays a summary of stack usage for a specified task, or for all tasks if no argument is given. The summary includes the total stack size (SIZE), the current number of stack bytes used (CUR), the maximum number of stack bytes used (HIGH), and the number of bytes never used at the top of the stack (MARGIN = SIZE - HIGH).
	Both the execution and the exception stack information are displayed. The exception stack is used by the task when it gets an exception or by a process task when it enters a system call and executes kernel code.
	For example:
	-> checkStack tShell0 NAME ENTRY TID SIZE CUR HIGH MARGIN
	tShell0       shellTask       0x60351ba8       77824       6272       14144       63680         (Exception Stack)       12096       0       680       11416         value = 1614093224 = 0x60351ba8
	The maximum stack usage is determined by scanning down from the top of the stack for the first byte whose value is not 0xee. In VxWorks, when a task is spawned, all bytes of a task's stack are initialized to 0xee. The task's stack will not be filled with 0xee if the task option <b>VX_NO_STACK_FILL</b> is specified or if the kernel configuration parameter

VX\_GLOBAL\_NO\_STACK\_FILL is set to TRUE.

VxWorks Kernel API Reference, 6.6 chkdsk()

DEFICIENCIES	It is possible for a task to write beyond the end of its stack, but not write into the last part of its stack. This will not be detected by <b>checkStack()</b> .
RETURNS	N/A
ERRNO	N/A
SEE ALSO	usrLib, taskSpawn(), the VxWorks programmer guides.

### chkdsk()

NAME	chkdsk() – perform consistency checking on a MS-DOS file system
SYNOPSIS	STATUS chkdsk ( const char * pDevName, /* device name */ u_int repairLevel, /* how to fix errors */ u_int verbose /* verbosity level */ )
DESCRIPTION	This function invokes the integral consistency checking built into the <b>dosFsLib</b> file system, via FIOCHKDSK ioctl. During the test, the volume will be un-mounted and re-mounted, invalidating file descriptors to prevent any application code from accessing the volume during the test. If the drive was exported, it will need to be re-exported again as its file descriptors were also invalidated. Furthermore, the test will emit messages describing any inconsistencies found on the disk, as well as some statistics, depending upon the value of the <i>verbose</i> argument. Depending upon the value of <i>repairLevel</i> , the inconsistencies will be repaired, and changes written to disk.
	These are the values for <i>repairLevel</i> :
	0 Same as <b>DOS_CHK_ONLY</b> (1)
	DOS_CHK_ONLY (1) Only report errors, do not modify disk.
	DOS_CHK_REPAIR (2) Repair any errors found.
	These are the values for <i>verbose</i> :
	0 similar to <b>DOS_CHK_VERB_1</b>
	<b>DOS_CHK_VERB_SILENT</b> (0xff00) Do not emit any messages, except errors encountered.

	DOS_CHK_VERB_1 (0x0100) Display some volume statistics when done testing, as well as errors encountered during the test.
	DOS_CHK_VERB_2 (0x0200) In addition to the above option, display path of every file, while it is being checked. This option may significantly slow down the test process.
	Note that the consistency check procedure will <i>unmount</i> the file system, meaning the all currently open file descriptors will be deemed unusable.
RETURNS	OK or ERROR if device can not be checked or could not be repaired.
ERRNO	Not Available
SEE ALSO	usrFsLib, dosFsLib, the VxWorks programmer guides.

# chmod()

NAME	<b>chmod()</b> – change the permission mode of a file
SYNOPSIS	<pre>int chmod   (     const char *path, /* path of the file */     mode_t mode /* mode bits to change */   )</pre>
DESCRIPTION	The chmod utility changes or assigns the mode of a file. The mode of a file specifies its permissions and other attributes. Note that this routine receives <b>path of the file whose mode needs to be changed</b> as the first argument compairing to fchmod routine.
	The value of <i>mode</i> is bitwise inclusive OR of the permissions to be assigned
	These permission constants are defined in <i>sys/stat.h</i> as follows:
	S_IRUSR Read permission, owner.
	S_IWUSR Write permission, owner.
	S_IXUSR Execute/search permission, owner.
	S_IRWXU Read/write/execute permission, owner.

	S_IRGRP Read permission, group.
	S_IWGRP Write permission, group.
	S_IXGRP Execute/search permission, group.
	S_IRWXG Read/write/execute permission, group.
	S_IROTH Read permission, other.
	S_IWOTH Write permission, other.
	S_IXOTH Execute/search permission, other.
	S_IRWXO Read/write/execute permission, other.
RETURNS	If it succeeds, returns <b>OK</b> , 0. Otherwise, <b>ERROR</b> , -1 is returned, errno is set to indicate the error and no change is done to the file.
	The following example changes the mode of the file "myFile" to owner Read/write/execute, group Read and other Read:
	<pre>status = chmod ("myFile", S_IRWXU   S_IRGRP   S_IROTH );</pre>
ERRNO	<b>ENOENT</b> Either <i>path</i> is an empty string or <b>NULL</b> pointer.
	ELOOP Circular symbolic link of <i>path</i> , or too many links.
	EMFILE Maximum number of files already open.
	<b>S_iosLib_DEVICE_NOT_FOUND (ENODEV)</b> No valid device name found in <i>path</i> .
	others Other errors reported by device driver of <i>path</i> .
SEE ALSO	fsPxLib

NAME	<b>clock_getres( )</b> – get the clock resolution (POSIX)
SYNOPSIS	<pre>int clock_getres   (    clockid_t clock_id, /* clock ID */    struct timespec * res /* where to store resolution */   )</pre>
DESCRIPTION	This routine gets the clock resolution, in nanoseconds, based on the rate returned by <b>sysClkRateGet()</b> . If <i>res</i> is non-NULL, the resolution is stored in the location pointed to.
RETURNS	0 ( <b>OK</b> ), or -1 ( <b>ERROR</b> ) if <i>clock_id</i> is invalid.
ERRNO	EINVAL
SEE ALSO	clockLib, clock_settime(), sysClkRateGet(), clock_setres()

# clock\_gettime()

NAME	<b>clock_gettime()</b> – get the current time of the clock (POSIX)
SYNOPSIS	<pre>int clock_gettime   (     clockid_t clock_id, /* clock ID */     struct timespec * tp /* where to store current time */   )</pre>
DESCRIPTION	This routine gets the current value <i>tp</i> for the clock.
RETURNS	0 ( <b>OK</b> ), or -1 ( <b>ERROR</b> ) if <i>clock_id</i> is invalid or <i>tp</i> is <b>NULL</b> .
ERRNO	EINVAL EFAULT
SEE ALSO	clockLib

	clock_nanosleep()
NAME	<b>clock_nanosleep()</b> – high resolution sleep with specifiable clock
SYNOPSIS	<pre>int clock_nanosleep   (     clockid_t clock_id,     int flags,     const struct timespec * rqtp,     struct timespec * rmtp   )</pre>
DESCRIPTION	If the flag <b>TIMER_ABSTIME</b> is not set in <i>flags</i> , this function causes the current thread to be delayed until either the time interval specified by <i>rqtp</i> has elapsed, or a signal is delivered to the calling thread and its action is to invoke a signal handler, or the process is terminated. The clock used to measure the time is the clock specified by <i>clock_id</i> .
	If the flag <b>TIMER_ABSTIME</b> is set in <i>flags</i> , this function causes the current thread to be delayed until either the time value of the clock specified by <i>clock_id</i> reaches the absolute time specified by <i>rqtp</i> , or a signal is delivered to the calling thread whose action is to invoke a signal handler, or the process is terminated. If at the time of the call, the time value specified by <i>rqtp</i> is less than or equal to the time value of <i>clock_id</i> , this function returns immediately without delaying the calling process.
	The delay caused by this function may be longer than requested because <i>rqtp</i> is rounded up to an integer multiple of the timer resolution, or because of the scheduling of other tasks by the system. Except for the case of being interrupted by a signal, the suspension time for the relative delay (i.e. if <b>TIMER_ABSTIME</b> is not set) is not less than the time interval <i>rqtp</i> , as measured by the corresponding clock.
	If a signal is caught by the calling task while sleeping for a relative time delay (i.e. flag <b>TIMER_ABSTIME</b> is not set in the <i>flags</i> argument), and the <i>rmtp</i> argument is non- <b>NULL</b> , the timespec structure referenced by <i>rmtp</i> is updated to contain the amount of time remaining in the interval. This is the requested sleep time minus the time actually slept.
	This function only supports CLOCK_REALTIME and CLOCK_MONOTONIC clocks.
RETURNS	0 (OK), or -1 (ERROR) if unsuccessful.
ERRNO	<b>EINVAL</b> <i>tp</i> is outside the supported range, or the <i>tp</i> nanosecond value is less than 0 or equal to or greater than 1,000,000,000.
	EINTR The sleep was interrupted by receiving a signal .
	ENOTSUP The <i>clock_id</i> value is not supported.

clockLib, clock\_getres() SEE ALSO

#### clock\_setres()

NAME	<pre>clock_setres() - set the clock resolution</pre>
SYNOPSIS	<pre>int clock_setres   (     clockid_t clock_id, /* clock ID */     struct timespec * res /* resolution to be set */   )</pre>
DESCRIPTION	This routine is obsolete. It will always return <b>OK</b> .
NOTE	Non-POSIX.
RETURNS	OK always.
ERRNO	EINVAL
SEE ALSO	clockLib, clock_getres( ), sysClkRateSet( )

# clock\_settime()

NAME	<b>clock_settime()</b> – set the clock to a specified time (POSIX)
SYNOPSIS	<pre>int clock_settime   (    clockid_t clock_id, /* clock ID */    const struct timespec * tp /* time to set */   )</pre>
DESCRIPTION	This routine sets the clock to the value $tp$ , which should be a multiple of the clock resolution. If $tp$ is not a multiple of the resolution, it is truncated to the next smallest multiple of the resolution.
RETURNS	0 ( <b>OK</b> ), or -1 ( <b>ERROR</b> ) if <i>clock_id</i> is invalid, <i>tp</i> is outside the supported range, or the <i>tp</i> nanosecond value is less than 0 or equal to or greater than 1,000,000,000.
ERRNO	EINVAL

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SEE ALSO clockLib, clock\_getres()

#### close()

NAME	<b>close( )</b> – close a file
SYNOPSIS	STATUS close ( int fd /* file descriptor to close */ )
DESCRIPTION	This routine closes the specified file and frees the file descriptor. It calls the device driver to do the work.
RETURNS	The status of the driver close routine, or <b>ERROR</b> if the file descriptor is invalid.
ERRNO	EBADF Invalid file descriptor. Others Other errors generated by device drivers.
SEE ALSO	ioLib

## closedir()

NAME	<b>closedir( )</b> – close a directory (POSIX)
SYNOPSIS	STATUS closedir ( DIR *pDir /* pointer to directory descriptor */ )
DESCRIPTION	This routine closes a directory which was previously opened using <b>opendir()</b> . The $pDir$ parameter is the directory descriptor pointer that was returned by <b>opendir()</b> .
RETURNS	OK or ERROR, the result of the close() command.
ERRNO	EBADF Invalid file descriptor.

Others Other errors generated by device drivers.

SEE ALSO dirLib, opendir(), readdir(), rewinddir()

# cnsAppRegister()

**NAME cnsAppRegister()** – Registers an application with the CNS library.

SYNOPSIS STATUS cnsAppRegister

	( char * pName, CNS_APP_READ_FUNCPTR readFunc, CNS_APP_WRITE_FUNCPTR writeFunc, CNS_MEDIATYPE_ADD_FUNCPTR mediaTypeAddFunc, CNS_MEDIATYPE_REMOVE_FUNCPTR mediaTypeRmFunc, CNS_DATA_PARSE_FUNCPTR dataParse )
DESCRIPTION	This routine registers an application that uses CNS services. The information to be passed during registration includes the application name and function pointers to access the application's local objects.
ARGUMENTS	<b>pName</b> identifies the application.
	readFunc writeFunc haveServer - Specifies whether the application implements its own read server.
	<b>mediaTypeAddFunc</b> - Reuired only if <b>haveServer</b> is set to <b>TRUE</b> . It points to a function to notify the application of the addition of a new media type.
	<b>mediaTypeRmFunc</b> - Reuired only if <b>haveServer</b> is set to <b>TRUE</b> . It points to a function to notify the application of the removal of a media type.
RETURNS	<b>OK</b> or <b>ERROR</b> if no more media can be added or <b>pAppInfo</b> is <b>NULL</b> or one or more of the fields of <b>cnsAppInfo_t</b> are <b>NULL</b> .

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VxWorks Kernel API Reference, 6.6 cnsClose()

**ERRNO** ERRNO for CNS-internal errors are TBD.

SEE ALSO cnsLib

#### cnsClose()

NAME	<b>cnsClose()</b> – Close or create and open named communication medium for read/write.
SYNOPSIS	STATUS cnsClose ( cnsMediaId_t * pMediaId )
DESCRIPTION	This routine closes the named communication medium.
ARGUMENTS	<b>pMediaId</b> is a pointer to an instance of the <b>cnsMediaId_t</b> , which describes the media being closed. <b>pMediaId-&gt;connId</b> cannot be 0.
RETURNS	<b>OK</b> or <b>ERROR</b> if the named media cannot be closed.
ERRNO	<b>cnsClose</b> maintains the ERRNO of the underlying media for media close errors. ERRNO for CNS-internal errors are TBD.
SEE ALSO	cnsLib

# cnsCompLibInit( )

NAME	<b>cnsCompLibInit()</b> – Initialize the CNS COMP library.
SYNOPSIS	STATUS cnsCompLibInit ( void )
DESCRIPTION	This routine initializes the CNS COMP library.

ARGUMENTS N/A

**RETURNS** OK or ERROR if media registration fails.

**ERRNO** ERRNO for CNS-internal errors are TBD.

SEE ALSO cnsCompLib

## cnsDefaultMediaTypeSet()

NAME	<b>cnsDefaultMediaTypeSet()</b> – Set the default media type.
SYNOPSIS	STATUS cnsDefaultMediaTypeSet ( char * pName )
DESCRIPTION	This routine sets the default media type to that specified by <b>pName</b> . The default media type is used in cases where channel access routines are called without specifying the media type.
ARGUMENTS	<b>pName</b> specifies the media type to which the default is set.
	Returns: ERROR if the specified media type does not exist or OK.
RETURNS	Not Available
ERRNO	S_cnsLib_MEDIATYPE_INVALID if the specified media type does not exist. S_cnsLib_GEN_ERROR
SEE ALSO	cnsLib

# cnsLibInit()

NAME	<b>cnsLibInit()</b> – Initialize the CNS library.
SYNOPSIS	STATUS cnsLibInit ( ulong_t maxMediaCount )
DESCRIPTION	This routine initializes the CNS library.
ARGUMENTS	<b>maxMediaCount</b> , if non-0, specifies the maximum number of media types to support. If 0, <b>cnsLib</b> uses the default <b>CNS_MAX_MEDIA_TYPES</b> , which is currently defined in ' <b>cnsCfg.h</b> .
RETURNS	<b>OK</b> or <b>ERROR</b> if write error occurs.
ERRNO	ERRNO for CNS-internal errors are TBD.
SEE ALSO	cnsLib

# cnsMediaRegister()

NAME	<b>cnsMediaRegister()</b> – Registers a communication media with the CNS.
SYNOPSIS	STATUS cnsMediaRegister ( cnsMediaInfo_t * pMediaInfo )
DESCRIPTION	This routine registers a communication media with the CNS. The information to be passed during registration includes the medium type name and function pointers that are used to access the communication medium.
ARGUMENTS	<b>pMediaInfo</b> is a pointer to an instance of the <b>cnsMediaInfo_t</b> , which specifies the medium name and the pointers to medium access functions.

RETURNS	OK or ERROR if no more media can be added or <b>pMediaInfo</b> is NULL or one or more of the fields of <b>cnsMediaInfo_t</b> are NULL.
ERRNO	ERRNO for CNS-internal errors are TBD.
SEE ALSO	cnsLib

#### cnsMediaTypeRemove()

**NAME** cnsMediaTypeRemove() – Remove a media type from an application's media list.

**DESCRIPTION cnsMediaRegister()** calls this routine everytime a new media type is removed to the media list.

**ARGUMENTS pApp**, points to an instance of the **cnsApp\_t** structure, which describes the application.

type identifies the media type being added.

RETURNS OK or ERROR.

**ERRNO** ERRNO for CNS-internal errors are TBD.

TODO: Must implement this.

STATUS cnsMediaTypeRemove
 (
 cnsApp\_t \* pApp,
 lont type
 )
{}

SEE ALSO cnsLib

# cnsMediumTypeNext()

NAME	<b>cnsMediumTypeNext( )</b> – Return the name of the media type next in the list.
SYNOPSIS	char * cnsMediumTypeNext ( char * pName )
DESCRIPTION	This routine returns the name pf the media type next to that specified by <b>pName</b> . If <b>pName</b> is <b>NULL</b> or if <b>pName</b> points to an empty string, the name of the first media type in the list is returned.
ARGUMENTS	<b>pName</b> specifies the name of the media type whose successor in the list is being sought.
RETURNS	NULL if the media type specified by <b>pName</b> is the last in the list; the name of the media type next to that specified by <b>pName</b> otherwise.
ERRNO	Not Available
SEE ALSO	cnsLib

## cnsMsgEncode()

NAME	<b>cnsMsgEncode( )</b> – Encode a message as understood by CNS.
SYNOPSIS	STATUS cnsMsgEncode ( long type, /* CNS_MSGTYPE_READREQ */ char * pData, char * pFormat, ulong_t * pFormatLen
DESCRIPTION	) This routine formats a message ad understood by CNS. A CNS message starts with the message type and is followed by optional data field which is specific to the message data.
	There are five message types: read request messafe, write request message, write with ack request message, reply <b>OK</b> message, and reply <b>ERROR</b> message.

	The request messages (read, write, and wite with ack) are typically sent by clients. The reply messages are sent by notification processor applications. A read request message is always followed by a reply message. The reply message starts with a status string (ok or error) and is followed by associated data.
ARGUMENTS	<ul> <li>type specifies the message data type. The value of type can be one of following:</li> <li>o CNS_REQ_READ - specifies a read request.</li> <li>o CNS_REQ_WRITE - specifies a write request.</li> <li>o CNS_REQ_WRITE_WACK - specifies write request with ack back.</li> <li>o CNS_REPLY_OK - specifies a reply to a successful read request or an ack to a successful write request.</li> <li>o CNS_REPLY_ERR - specifies a failed read request or a failed write request.</li> </ul>
	<b>pData</b> points to the data associated with the message type.
	For read requests, the data consists typically of information of what is to be read. For example, for CSM, the data consists of the component's name and the object to be read.
	For write requests, the data consists typically of information of what is to be modified and the modification data.
	<b>pformat</b> points to a buffer to contain the formatted data.
	<b>pFormatLen</b> specifies the capacity of <b>pFormat</b> during input, and specifies the actual length of the format when the <b>cnsMsgEncode( )</b> returns.
	create, if TRUE, specifies that the medium is to be created.
RETURNS	<b>OK</b> or <b>ERROR</b> if the named media cannot be opened/created.
ERRNO	<b>cnsMsgEncode</b> maintains the ERRNO of the underlying media for media open errors. ERRNO for CNS-internal errors are TBD.
SEE ALSO	cnsLib

	cnsOpen()
NAME	<b>cnsOpen()</b> – Open or create and open named communication medium for read/write.
SYNOPSIS	STATUS cnsOpen ( cnsMediaId_t * pMediaId, BOOL create, long * pConnState )
DESCRIPTION	Depending upon the value of <b>create</b> , this routine opens or creates and then opens the named communication medium for read/write.
ARGUMENTS	<b>pMedia</b> is a pointer to an instance of the <b>cnsMediaId_t</b> , which describes the media being created/opened. If open/create of the media succeeds, the <b>mediaId</b> and <b>connId</b> fields are updated to identify the open medium type and connection ID.
	create, if TRUE, specifies that the medium is to be created.
RETURNS	<b>OK</b> or <b>ERROR</b> if the named media cannot be opened/created.
ERRNO	<b>cnsOpen</b> maintains the ERRNO of the underlying media for media open errors. ERRNO for CNS-internal errors are TBD.
SEE ALSO	cnsLib

# cnsRead()

NAME	<b>cnsRead()</b> – Read fro	m a communication medium.
SYNOPSIS	<pre>STATUS cnsRead   (     cnsMediaId_t *     char *     char *     ulong_t     ulong_t *    )</pre>	pMediaId, pReqInfo, pBuf, bufLen, pReadBytes

DESCRIPTION	This routine reads from a communication medium.
ARGUMENTS	<b>pMediaId</b> is a pointer to an instance of the <b>cnsMediaId_t</b> , which describes the media being read from. <b>connId</b> of <b>pMediaId</b> cannot be 0. If <b>pName</b> is provided, then <b>cnsRead</b> verifies that the name matches the corresponding <b>mediaId</b> and <b>connId</b> . If they dont match, <b>ERROR</b> is returned.
	<b>pBuf</b> points to a buffer to which the read data is to be copied.
	<b>bufLen</b> specifies the capacity of <b>pBuf</b> in bytes.
	<b>pReadBytes</b> points to a buffer to which the number of actually read bytes is to be copied.
RETURNS	<b>OK</b> or <b>ERROR</b> if error occured while reading, or if <b>pMediaId-&gt;connId</b> is 0 or if <b>pMediaId-&gt;pName</b> is not <b>NULL</b> and it is not consistent with the corresponding <b>pMediaId-&gt;connId</b> and <b>pMediaId-&gt;mediaId</b> .
ERRNO	<b>cnsRead</b> maintains the ERRNO of the underlying media for media read errors. ERRNO for CNS-internal errors are TBD.
SEE ALSO	cnsLib

# cnsWrite()

**DESCRIPTION** This routine writes to a communication medium.

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ARGUMENTS	<b>pMediaId</b> is a pointer to an instance of the <b>cnsMediaId_t</b> , which describes the media being written to. <b>connId</b> of <b>pMediaId</b> cannot be 0. If <b>pName</b> is provided, then <b>cnsWrite</b> verifies that the name matches the corresponding <b>mediaId</b> and <b>connId</b> . If they dont match, <b>ERROR</b> is returned.
	<b>pBuf</b> points to a buffer from which the data to be written is to be copied.
	<b>bufLen</b> specifies the capacity of <b>pBuf</b> in bytes.
	<b>pReplyLen</b> if greater than 0, it implies that the caller expects an acknowledgement back. The message type is then set to <b>CNS_REQ_WRITE_WACK</b> .
RETURNS	OK or ERROR if error occured while reading, or if <b>pMediaId-&gt;connId</b> is 0 or if <b>pMediaId-&gt;pName</b> is not NULL and it is not consistent with the corresponding <b>pMediaId-&gt;connId</b> and <b>pMediaId-&gt;mediaId</b> .
ERRNO	<b>cnsWrite</b> maintains the ERRNO of the underlying media for media write errors. ERRNO for CNS-internal errors are TBD.
SEE ALSO	cnsLib

## commit()

NAME	<b>commit( )</b> – commit current transaction to disk.
SYNOPSIS	STATUS commit ( const char * pDevName /* name of the device to commit */ )
DESCRIPTION	This command is for transactional based file systems only such as HRFS. It is a shortcut for the ioctl function FIOCOMMITFS which commits the current transaction to disk to make changes permanment.
EXAMPLE	-> commit "/ata0a" /* commit transaction on "/fd0" */

RETURNS	<b>OK</b> , or <b>ERROR</b> if the device is not formatted with a file system that does not support the FIOCOMMITFS ioctl function or <i>pDevName</i> is not valid.
ERRNO	Not Available

**SEE ALSO** usrFsLib, hrFsLib, VxWorks Kernel Programmer's Guide: Kernel Shell

# copy()

NAME	<b>copy()</b> – copy <i>in</i> (or stdin) to <i>out</i> (or stdout)
SYNOPSIS	<pre>STATUS copy   (     const char * in, /* name of file to read (if NULL assume stdin) */     const char * out /* name of file to write (if NULL assume stdout) */   )</pre>
DESCRIPTION	This command copies from the input file to the output file, until an end-of-file is reached.
EXAMPLES	<pre>The following example displays the file dog, found on the default file device:     -&gt; copy <dog (default="" ,="" -="" ^d)="" an="" console="" copies="" ct0="" device="" dog,="" eof="" example="" file="" from="" is="" on="" the="" this="" to="" typed:="" until=""> copy &gt;/ct0/dog This example copies the file dog, found on the default file device, to device /ct0/:     -&gt; copy <dog>/ct0/dog This example makes a conventional copy from the file named file1 to the file named file2:     -&gt; copy "file1", "file2" Remember that standard input and output are global; therefore, spawning the first three constructs will not work as expected.</dog></dog></pre>
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>in</i> or <i>out</i> cannot be opened/created, or if there is an error copying from <i>in</i> to <i>out</i> .
ERRNO	Not Available
SEE ALSO	usrFsLib, copyStreams(), tyEOFSet(), cp(), xcopy(), the VxWorks programmer guides.

# copyStreams()

NAME	<b>copyStreams()</b> – copy from/to specified streams
SYNOPSIS	<pre>STATUS copyStreams   (    int inFd, /* file descriptor of stream to copy from */    int outFd /* file descriptor of stream to copy to */   )</pre>
DESCRIPTION	This command copies from the stream identified by $inFd$ to the stream identified by $outFd$ until an end of file is reached in $inFd$ . This command is used by <b>copy()</b> .
RETURNS	<b>OK</b> , or <b>ERROR</b> if there is an error reading from <i>inFd</i> or writing to <i>outFd</i> .
ERRNO	Not Available
SEE ALSO	usrFsLib, copy(), the VxWorks programmer guides.

# coreDumpClose()

NAME	coreDumpClose() – close a core dump
SYNOPSIS	STATUS coreDumpClose ( CORE_DUMP_ID coreDumpId /* ID returned by coreDumpOpen() */ )
DESCRIPTION	This routine frees resources allocated by <b>coreDumpOpen()</b> .
RETURNS	<b>OK</b> or <b>ERROR</b> if <i>coreDumpId</i> is invalid
ERRNO	N/A
SEE ALSO	coreDumpUtilLib
# coreDumpCopy()

NAME	<b>coreDumpCopy()</b> – copy a core dump to the given path
SYNOPSIS	STATUS coreDumpCopy ( UINT32 coreDumpIndex, /* core dump index, or 0 for all */ char * destPath /* destination path */ )
DESCRIPTION	This routine copies the core dump specified by <i>coreDumpIndex</i> to <i>destPath</i> . If <i>coreDumpIndex</i> is 0, then all the core dumps available on device are copied. If <i>destPath</i> is <b>NULL</b> , then the destination path is current directory.
	<i>coreDumpIndex</i> is the index of the core dump returned by <b>coreDumpNextGet()</b> routine.
RETURNS	OK, or ERROR if a core dump index is invalid, or if the copy failed.
ERRNO	N/A
SEE ALSO	coreDumpUtilLib

# coreDumpCreateHookAdd()

NAME	<b>coreDumpCreateHookAdd()</b> – add a routine to be called at every core dump create
SYNOPSIS	<pre>STATUS coreDumpCreateHookAdd    (     FUNCPTR createHook /* routine to be called when a core dump is created */    )</pre>
DESCRIPTION	This routine adds a specified routine to a list of routines that will be called whenever a core dump is created. Upon creation, all routines specified by <b>coreDumpCreateHookAdd()</b> will be called.
	The routine should be declared as follows:
	STATUS createHook (void)
RETURNS	OK, or ERROR if the table of core dump create routines is full.
ERRNO	S_coreDumpLib_CORE_DUMP_HOOK_TABLE_FULL core dump create hook table is full

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SEE ALSO coreDumpHookLib, coreDumpCreateHookDelete()

### coreDumpCreateHookDelete( )

NAME	<pre>coreDumpCreateHookDelete() - delete a previously added core dump create routine</pre>
SYNOPSIS	STATUS coreDumpCreateHookDelete ( FUNCPTR createHook /* routine to be deleted from list */ )
DESCRIPTION	This routine removes a specified routine from the list of routines to be called at each core dump create.
RETURNS	OK, or ERROR if the routine is not in the table of core dump create routines.
ERRNO	S_coreDumpLib_CORE_DUMP_HOOK_NOT_FOUND core dump create hook can not be found
SEE ALSO	coreDumpHookLib, coreDumpCreateHookAdd()

### coreDumpDevFormat()

NAME	coreDumpDevFormat() – format the core dump device
SYNOPSIS	STATUS coreDumpDevFormat ( UINT32 coreDumpMax /* Maximum number of core dump on device */ )
DESCRIPTION	This routine formats the core dump device. Formatting a core dump device consists of erasing all the core dumps available on device. If an <b>erase()</b> routine is specified for the core dump device, then this routine will be called to erase the whole device (useful for flash devices for example); otherwise the only part of the core dump storage will be erased using the underlying device write command.
	Once the core dump device has been erased, <b>coreDumpDevFormat()</b> reformat it to support the given number of core dump.
NOTE	It is not possible to erase only one core dump from the device, the full device is erased.

**RETURNS** OK, or ERROR if the format operation failed.

ERRNO N/A

SEE ALSO coreDumpLib

#### coreDumpDevShow()

NAME coreDumpDevShow() – display information on core dump device

SYNOPSIS STATUS coreDumpDevShow (void)

**DESCRIPTION** This routine displays basic information on the core dump device. It displays the current number of core dumps stored on the device, the maximum number of core dumps that can be stored on the device. The total size of the core dump device and the free size on this device.

- RETURNS OK, or ERROR.
- ERRNO N/A
- SEE ALSO coreDumpShow

#### coreDumpInfoGet()

NAME	<b>coreDumpInfoGet( )</b> – get information on a core dump
Synopsis	CORE_DUMP_INFO * coreDumpInfoGet ( UINT32 coreDumpIndex /* core dump index */ )
DESCRIPTION	This routine retrieves information on a given core dump. It allocates a <b>CORE_DUMP_INFO</b> structure and fills it with the retrieved information, a pointer to this structure is then returned, and it will be up to the caller to free this structure when the information will have been consumed.
	<i>coreDumpIndex</i> is the index of the core dump returned by <b>coreDumpNextGet()</b> routine.
	typedef struct core_dump_info /* core dump information */

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```
UINT32
                         coreDumpIndex;/* core dump index */
                         valid; /* core dump validity */
   BOOL
                                    /* core dump errno */
   UINT32
                         errnoVal;
   char
                         name[MAX CORE DUMP LEN];
                                      /* name of the core dump */
   size t
                                      /* size of the core dump */
                         size;
   CORE_DUMP_TYPE
                        type;
                                      /* origin of the core dump */
                         taskId;
                                     /* task Id (kernel core dump) */
   int
                         taskName[MAX_CORE_DUMP_TASK_LEN];
   char
                                      /* name of task */
   UINT32
                                      /* process Id (process core dump)
                        rtpId;
*/
   char
                         rtpName[MAX_CORE_DUMP_RTP_LEN];
                                      /* path to RTP */
   int
                         excNum;
                                      /* exception number (Not valid for
* /
                                      /* on-demand & Panic core dumps) */
   UTNT32
                         pc;
                                      /* exception program counter */
                                      /* exception stack pointer */
   UINT32
                        sp;
   UINT32
                        fp;
                                      /* exception frame pointer */
   time t
                        time;
                                      /* generation calendar time */
   UINT32
                        ticks;
                                      /* VxWorks time stamp */
   CORE DUMP CKSUM STATUS cksumStatus; /* core dump checksum status */
           infoString[MAX_CORE_DUMP_INFO_LEN];
   char
                                /* information string */
                               /* Indicate validity of exception */
   BOOL
                    excIsValid;
                                 /* information (excNum, pc, sp, fp) */
```

} CORE\_DUMP\_INFO;

The *valid* bit in **CORE\_DUMP\_INFO** structure indicates if the core dump was successfully written on the storage or if there was an error writing it; if a core dump has been only partially written (because the storage is too small for example) then it will be marked as invalid and the *size* field will represent the size that has been actually written on the storage. If the core dump is marked as not valid, the *errnoVal* field contains the errno that has been set while generating the core dump.

The *type* field indicates the origin of the core dump. Here are the possible values:

typedef enum	/* Core Dump Type */
{	
CORE_DUMP_USER,	/* 0: user coredump (on-demand) */
CORE_DUMP_KERNEL_INIT,	/* 1: fatal error during kernel intialization
*/	
CORE_DUMP_INTERRUPT,	/* 2: Obsolete / Kept for backward compat */
CORE_DUMP_KERNEL_PANIC,	/* 3: Obsolete / Kept for backward compat */
CORE_DUMP_KERNEL_TASK,	/* 4: kernel task error */
CORE_DUMP_RTP,	/* 5: process coredump */
CORE_DUMP_KERNEL	/* 6: VxWorks kernel error */
<pre>} CORE_DUMP_TYPE;</pre>	

The *cksumStatus* field indicates the status of the core dump checksum. Here are the possible values:

typedef enum	/*	Co	re	Dump	Che	ecksum	Sta	tus *	/
{									
CORE_DUMP_NO_CKSUM,	/*	0:	Nc	o cksu	um a	availak	ble	(N/A)	*/

```
CORE_DUMP_CKSUM_OK,
CORE_DUMP_CKSUM_ERROR
} CORE_DUMP_CKSUM_STATUS;
```

/\* 1: Core dump checksum status OK \*/
/\* 2: Core dump checksum status ERROR \*/

If the core dump checksum status is equal to CORE\_DUMP\_NO\_CKSUM, then this means that the checksum facility was not enabled on the target (CORE\_DUMP\_CKSUM\_ENABLE parameter of INCLUDE\_CORE\_DUMP component).

**RETURNS** A pointer to a **CORE\_DUMP\_INFO** structure, or **NULL** if failed to read core dump information from the core dump device.

ERRNO N/A

SEE ALSO coreDumpUtilLib

#### coreDumpIsAvailable( )

- NAME coreDumpIsAvailable() is a core dump available for retrieval
- SYNOPSIS BOOL coreDumpIsAvailable (void)
- **DESCRIPTION** This routine can be called to determine if at least one core dump is available on core dump device for retrieval.
- **RETURNS** TRUE if there is at least one core dump available for retrieval or FALSE if there is no core dump available on device.

ERRNO N/A

SEE ALSO coreDumpUtilLib

#### coreDumpMemDump()

**NAME** coreDumpMemDump() – dump an area of memory in VxWorks core dump

VxWorks Kernel API Reference, 6.6 coreDumpMemFilterAdd()

DESCRIPTION	This routine can be called by a user to dump an additional area of memory in a core dump. The area of memory to dump is specified by <i>buffer</i> and <i>size</i> arguments. The <i>vaddr</i> allows to specify at which address in the core dump, the area of memory must be mapped.
	This routine must be called from a VxWorks core dump creation hook.
	No verification is done to check that an area of memory is already available in a core dump, an area of memory can be dumped several times in the core dump.
	If a core dump memory filter has been installed and filter the provided memory area or part of it, the filtered area will not be written in core dump.
RETURNS	<b>OK</b> or <b>ERROR</b> if core dump generation has failed
ERRNOS	S_coreDumpLib_CORE_DUMP_GENERATE_NOT_RUNNING This routine is not called from a core dump creation hook
SEE ALSO	coreDumpLib

# coreDumpMemFilterAdd()

NAME	<pre>coreDumpMemFilterAdd() – add a memory region filter</pre>		
SYNOPSIS	<pre>STATUS coreDumpMemFilterAdd   (    void * addr, /* address of memory region to filter */    size_t size /* size of memory region to filter */   )</pre>		
DESCRIPTION	This routine adds a filter to exclude a memory region from core dump.		
RETURNS	OK, or ERROR if the maximum number of memory region filter is reached		
ERRNO	S_coreDumpLib_CORE_DUMP_FILTER_TABLE_FULL Core dump memory filter table is full		
SEE ALSO	coreDumpMemFilterLib		

NAME	coreDumpMemFilterDelete() – delete a memory region filter
SYNOPSIS	<pre>STATUS coreDumpMemFilterDelete   (     void * addr, /* address of memory region to filter */     size_t size /* size of memory region to filter */   )</pre>
DESCRIPTION	This routine deletes the core dump memory region filter specified by the <i>addr</i> and <i>size</i> arguments.
RETURNS	OK, or ERROR if the specified filter does not exist.
ERRNO	S_coreDumpLib_CORE_DUMP_FILTER_NOT_FOUND Filter not found in core dump filter table
SEE ALSO	coreDumpMemFilterLib

# coreDumpNextGet( )

NAME	<b>coreDumpNextGet( )</b> – get the next core dump on device
SYNOPSIS	STATUS coreDumpNextGet ( UINT32 currentIdx, /* current core dump index or 0 to */ /* get first core dump */ UINT32 * pNextIdx /* where to store next core dump */ /* index */ )
DESCRIPTION	This routine retrieves the index of the next core dump available on the device; this routine can be used to walk through the core dump list. The first call of this routine must be performed with <i>currentIdx</i> equal to zero to get the index of the first core dump found on the device; the returned index is stored at the memory location pointed to by <i>pNextIdx</i> . The returned index can then be used by other routines like <b>coreDumpOpen()</b> , <b>coreDumpInfoGet()</b> or by <b>coreDumpNextGet()</b> to get the index of the following core dump. If there is no next core dump, the returned index is set to 0.
RETURNS	<b>OK</b> or <b>ERROR</b> if there was an error reading the core dump list.
ERRNO	N/A

VxWorks Kernel API Reference, 6.6 coreDumpOpen()

SEE ALSO coreDumpUtilLib

# coreDumpOpen()

NAME	<b>coreDumpOpen()</b> – open an existing core dump for retrieval
SYNOPSIS	CORE_DUMP_ID coreDumpOpen ( UINT32 coreDumpIndex /* core dump index */ )
DESCRIPTION	This routine opens an existing core dump. The core dump content can then be retrieved using <b>coreDumpRead()</b> . The core dump must then be closed using <b>coreDumpClose()</b> . Some memory is allocated dynamically by <b>coreDumpOpen()</b> routine, it will be freed by <b>coreDumpClose()</b> routine.
	<i>coreDumpIndex</i> is the index of the core dump returned by <b>coreDumpNextGet()</b> routine.
RETURNS	A core dump ID, or NULL if open has failed.
ERRNO	N/A
SEE ALSO	coreDumpUtilLib

# coreDumpRead()

NAME	<b>coreDumpRead()</b> – read from a core file
SYNOPSIS	<pre>int coreDumpRead   (     CORE_DUMP_ID coreDumpId, /* ID returned by coreDumpOpen() */     void * buffer, /* where to store read data */     size_t size /* number of bytes to read */     )</pre>
DESCRIPTION	This routine reads a number of bytes (less than or equal to <i>size</i> ) from the specified core dump ID and places them in <i>buffer</i> .
RETURNS	The number of bytes read (between 1 and <i>size</i> , 0 if end of file), or <b>ERROR</b> if the core dump ID does not exist, of if there was an error reading the core dump.

ERRNO N/A

SEE ALSO coreDumpUtilLib

#### coreDumpShow()

**NAME** coreDumpShow() – display information on core dumps

SYNOPSIS STATUS coreDumpShow (

```
UINT32 coreDumpIndex, /* core dump index */
UINT32 level /* core dump show level */
```

# **DESCRIPTION** This routine displays basic information on the given core dump. If *coreDumpIndex* is 0, then it displays basic information on all core dumps available on the core dump device. If *level* is 1, then detailled information are displayed.

The core dump display contains the following fields:

Field	Meaning
NAME	Core dump name on device
IDX	Core dump index for retrieval
VALID	Indicate if core dump is valid or not
ERRNO	If core dump is not valid, errno set during core dump generation
SIZE	Size of the core dump (compressed size if compression is enabled)
CKSUM	Core dump checksum status
TYPE	Type of the core dump
TASK	Name of the task at the origin of the core dump

The following table describes the different core dump types:

Field	Meaning
USER	User coredump
KERNEL_INIT	Fatal error during kernel intialization
KERNEL_TASK	Kernel task error
RTP	Process coredump
KERNEL	VxWorks kernel error
UNKNOWN	Unknown coredump type

The following table describes the different core dump checksum status:

Field	Meaning
N/A	No cksum available

VxWorks Kernel API Reference, 6.6 coreDumpUsrGenerate()

	Field	Meaning					
	ОК	Core dun	np checksun	n status <b>OK</b>			
	ERROR	Core dun	np checksun	n status ERR	ROR		
EXAMPLE	-> coreDumpShow NAME IDX	0,1 VALID	ERRNO	SIZE	CKSUM	TYPE	TASK
	vxcore1.z 1	Y	N/A (	)x000ef05b	OK	KERNEL_TASK	t1
	Core Dump detai	lled info	ormation:				
	Time: Task: Process: Description: Exception numbe Program counter Stack pointer: Frame pointer: value = 0 = 0x0	r: :	THU JAN 01 "t1" (0x61 "(Kernel)' fatal kerr 0xb 0x6003823e 0x604d3da8 0x604d3fb0	L 00:01:42 11e0a20) ' (0x6017a5 nel task-le	1970 ( 500) evel ex	(ticks = 6174)	
RETURNS	<b>OK</b> , or <b>ERROR</b> if the failed to retrieve is dump index.	ne core du nformatio	mp device is n on a core	s not initializ dump, or if	zed, not the <i>core</i>	formatted, if <b>co</b> <i>DumpIndex</i> is n	oreDumpShow() not a valid core
ERRNO	N/A						
SEE ALSO	coreDumpShow						

# coreDumpUsrGenerate()

NAME	coreDumpUsrGenerate() – generate a user (on-demand) core dump
SYNOPSIS	STATUS coreDumpUsrGenerate (void)
DESCRIPTION	This routine generates a user (on-demand) core dump. When this routine is called, the core dump will be generated, stored on the configured storage and the target will be rebooted.
RETURNS	<b>OK</b> or <b>ERROR</b> if core dump generation failed.
ERRNO	N/A
SEE ALSO	coreDumpLib

# cosf()

NAME	<b>cosf( )</b> – compute a cosine (ANSI)
SYNOPSIS	<pre>float cosf   (    float x /* angle in radians */   )</pre>
DESCRIPTION	This routine returns the cosine of $x$ in single precision. The angle $x$ is expressed in radians.
RETURNS	The single-precision cosine of <i>x</i> .
ERRNO	Not Available
SEE ALSO	mathALib

# coshf()

NAME	<b>coshf( )</b> – compute a hyperbolic cosine (ANSI)
SYNOPSIS	<pre>float coshf   (    float x /* value to compute the hyperbolic cosine of */   )</pre>
DESCRIPTION	This routine returns the hyperbolic cosine of $x$ in single precision.
RETURNS	The single-precision hyperbolic cosine of $x$ if the parameter is greater than 1.0, or NaN if the parameter is less than 1.0.
	Special cases: If <i>x</i> is +INF, -INF, or NaN, <b>coshf( )</b> returns <i>x</i> .
ERRNO	Not Available
SEE ALSO	mathALib

# **cp()**

NAME	<b>cp()</b> – copy file into other file/directory.
SYNOPSIS	<pre>STATUS cp (     const char * src, /* source file or wildcard pattern */     const char * dest /* destination file name or directory */ )</pre>
DESCRIPTION	This command copies from the input file to the output file. If destination name is directory, a source file is copied into this directory, using the last element of the source file name to be the name of the destination file.
	This function is very similar to <b>copy()</b> , except it is somewhat more similar to the UNIX "cp" program in its handling of the destination.
	<i>src</i> may contain a wildcard pattern, in which case all files matching the pattern will be copied to the directory specified in <i>dest</i> . This function does not copy directories, and is not recursive. To copy entire subdirectories recursively, use <b>xcopy()</b> .
EXAMPLES	-> cp( "/sd0/FILE1.DAT","/sd0/dir2/f001.dat") -> cp( "/sd0/dir1/file88","/sd0/dir2") -> cp( "/sd0/*.tmp","/sd0/junkdir")
RETURNS	<b>OK</b> or <b>ERROR</b> if destination is not a directory while <i>src</i> is a wildcard pattern, or if any of the files could not be copied.
ERRNO	Not Available
SEE ALSO	usrFsLib, xcopy(), the VxWorks programmer guides.

# cplusCallNewHandler()

NAME	<b>cplusCallNewHandler()</b> – call the allocation failure handler (C++)
SYNOPSIS	extern void cplusCallNewHandler (void)
DESCRIPTION	This function provides a procedural-interface to the new-handler. It can be used by user-defined new operators to call the current new-handler. This function is specific to VxWorks and may not be available in other C++ environments.
RETURNS	N/A

ERRNO Not Available

SEE ALSO cplusLib

# cplusCtors()

NAME	<b>cplusCtors()</b> – call static constructors (C++)
SYNOPSIS	extern "C" void cplusCtors ( const char * moduleName /* name of loaded module */ )
DESCRIPTION	This function is used to call static constructors under the manual strategy (see <b>cplusXtorSet()</b> ). <i>moduleName</i> is the name of an object module that was "munched" before loading. If <i>moduleName</i> is 0, then all static constructors, in all modules loaded by the VxWorks module loader, are called.
EXAMPLES	The following example shows how to initialize the static objects in modules called "applx.out" and "apply.out".
	-> cplusCtors "applx.out" value = 0 = 0x0 -> cplusCtors "apply.out" value = 0 = 0x0
	The following example shows how to initialize all the static objects that are currently loaded, with a single invocation of <b>cplusCtors()</b> :
	-> cplusCtors value = 0 = 0x0
WARNING	<b>cplusCtors()</b> should only be called once per module otherwise unpredictable behavior may result.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	cplusLib, cplusXtorSet( )

|--|

NAME	<b>cplusCtorsLink()</b> – call all linked static constructors (C++)
SYNOPSIS	extern "C" void cplusCtorsLink (void)
DESCRIPTION	This function calls constructors for all of the static objects linked with a VxWorks bootable image. When creating bootable applications, this function should be called from <b>usrRoot()</b> to initialize all static objects. Correct operation depends on correctly munching the C++ modules that are linked with VxWorks.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	cplusLib

# cplusDemanglerSet()

NAME	cplusDemangler	Set() – change C++ demangling mode (C++)
SYNOPSIS	extern "C" void ( int mode )	cplusDemanglerSet
DESCRIPTION	This command se	ts the C++ demangling mode to <i>mode</i> . The default mode is 2.
	There are three de numeric codes:	emangling modes, <i>complete</i> , <i>terse</i> , and <i>off</i> . These modes are represented by
	Mode	Code
	off	0
	terse	1
	complete	2
	In complete mode and the function's	e, when C++ function names are printed, the class name (if any) is prefixed s parameter type list is appended.

In terse mode, only the function name is printed. The class name and parameter type list are omitted.

In off mode, the function name is not demangled.

1 1 1

EXAMPLES	The following example shows how one function name would be printed under each demangling mode:		
	Mode off terse complete	Printed symbol _member5classFPFl_PvPFPv_v _member foo::_member(void* (*)(long),void (*)(void*))	
RETURNS	N/A		
ERRNO	Not Available		
SEE ALSO	cplusLib		

# cplusDemanglerStyleSet()

NAME	<b>cplusDemanglerStyleSet( )</b> – change C++ demangling style (C++)
SYNOPSIS	extern "C" void cplusDemanglerStyleSet ( DEMANGLER_STYLE style )
DESCRIPTION	This command sets the C++ demangling mode to <i>style</i> . The available demangler styles are enumerated in <b>demangler.h</b> . The default demangling style depends on the toolchain used to build the kernel. For example if the Diab toolchain is used to build the kernel then the default demangler style is <b>DMGL_STYLE_DIAB</b> .
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	cplusLib

# cplusDtors()

**NAME** cplusDtors() – call static destructors (C++)

SYNOPSIS extern "C" void cplusDtors

VxWorks Kernel API Reference, 6.6 cplusDtorsLink()

```
(
const char * moduleName
)
```

- **DESCRIPTION** This function is used to call static destructors under the manual strategy (see **cplusXtorSet()**). *moduleName* is the name of an object module that was "munched" before loading. If *moduleName* is 0, then all static destructors, in all modules loaded by the VxWorks module loader, are called.
- **EXAMPLES** The following example shows how to destroy the static objects in modules called "applx.out" and "apply.out":

```
-> cplusDtors "applx.out"
value = 0 = 0x0
-> cplusDtors "apply.out"
value = 0 = 0x0
```

The following example shows how to destroy all the static objects that are currently loaded, with a single invocation of **cplusDtors()**:

-> cplusDtors value = 0 = 0x0

- **WARNING** cplusDtors() should only be called once per module otherwise unpredictable behavior may result.
- RETURNS N/A
- ERRNO Not Available
- SEE ALSO cplusLib, cplusXtorSet()

#### cplusDtorsLink( )

NAME	<b>cplusDtorsLink( )</b> – call all linked static destructors (C++)
SYNOPSIS	extern "C" void cplusDtorsLink (void)
DESCRIPTION	This function calls destructors for all of the static objects linked with a VxWorks bootable image. When creating bootable applications, this function should be called during system shutdown to decommission all static objects. Correct operation depends on correctly munching the C++ modules that are linked with VxWorks.
RETURNS	N/A

ERRNO Not Available

SEE ALSO cplusLib

# cplusLibInit( )

NAME	<b>cplusLibInit( )</b> – initialize the C++ library (C++)
SYNOPSIS	extern "C" STATUS cplusLibInit (void)
DESCRIPTION	This routine initializes the C++ library and forces all C++ run-time support to be linked with the bootable VxWorks image. If the configuration macro INCLUDE_CPLUS is defined, cplusLibInit() is called automatically from the root task, usrRoot(), in usrConfig.c.
RETURNS	OK or ERROR.
ERRNO	Not Available
SEE ALSO	cplusLib

# cplusXtorGet( )

NAME	<b>cplusXtorGet()</b> – get the c++ Xtors strategy
SYNOPSIS	extern "C" int cplusXtorGet (void)
DESCRIPTION	This function can be used to retrieve the current value of the C++ Xtors strategy.
RETURNS	1 for automatic or 0 for manual
ERRNO	Not Available
SEE ALSO	cplusLib, cplusXtorSet()

2

	cplusXtorSet()
NAME	<b>cplusXtorSet( )</b> – change C++ static constructor calling strategy (C++)
SYNOPSIS	extern "C" void cplusXtorSet ( int strategy )
DESCRIPTION	This command sets the C++ static constructor calling strategy to <i>strategy</i> . The default strategy is 1.
	There are two static constructor calling strategies: <i>automatic</i> and <i>manual</i> . These modes are represented by numeric codes:
	StrategyCodemanual0automatic1
	Under the manual strategy, a module's static constructors and destructors are called by <b>cplusCtors()</b> and <b>cplusDtors()</b> , which are themselves invoked manually.
	Under the automatic strategy, a module's static constructors are called as a side-effect of loading the module using the VxWorks module loader. A module's static destructors are called as a side-effect of unloading the module.
NOTE	The manual strategy is applicable only to modules that are loaded by the VxWorks module loader. Static constructors and destructors contained by modules linked with the VxWorks image are called using <b>cplusCtorsLink()</b> and <b>cplusDtorsLink()</b> .
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	cplusLib, cplusXtorGet()

### cpsr()

NAME cpsr() – return the contents of the current processor status register (ARM)

DESCRIPTION	This command extracts the contents of the status register from the TCB of a specified task. If <i>taskId</i> is omitted or zero, the last task referenced is assumed.
RETURNS	The contents of the current processor status register.
ERRNO	Not Available
SEE ALSO	dbgArchLib, VxWorks Programmer's Guide: Debugging

#### cpuPwrMgrEnable()

cpuPwrMgrEnable() - Set the CPU light power management to ON/OFF NAME SYNOPSIS STATUS cpuPwrMgrEnable ( BOOL enable ) DESCRIPTION This routine enables or disables the light CPU power manager based on the *enable* argument that is passed. When *enable* is **TRUE** the power manager is enabled. When *enable* is **FALSE** the power manager is disable. When it is enabled the power manager puts the CPU in the C1 state (non-executing state) when the VxWorks kernel becomes idle, that is, when there are no ISRs to process or tasks to dispatch. The "C1 state" is a term borrowed from the Advance Configuration and Power Interface (ACPI) specification. When in a non-executing state the CPU reduces its power consumption. The light power manager is automatically enabled when the VxWorks kernel boots. Therefore this routine need not be called unless a kernel application wishes to disable the power manager or to re-enable it after having disabled it. The only reason that could explain this routine returning ERROR is if the light power manager's initialization routine failed, which more than likely indicates that another power manager is configured in the VxWorks kernel. There can only be one power manager in a system. RETURNS OK or ERROR if the power manager could not be enabled. ERRNO N/A SEE ALSO cpuPwrLightLib, cpuPwrMgrIsEnabled()

VxWorks Kernel API Reference, 6.6 cpuPwrMgrIsEnabled()

### cpuPwrMgrIsEnabled()

NAME cpuPwrMgrIsEnabled() – Get the CPU power management status

SYNOPSIS BOOL cpuPwrMgrIsEnabled (void)

**DESCRIPTION** This routine returns the status of the light power manager. **TRUE** is returned when it is enabled. Otherwise **FALSE** is returned.

RETURNS TRUE or FALSE

ERRNO N/A

SEE ALSO cpuPwrLightLib, cpuPwrMgrEnable()

#### creat()

NAME	<b>creat()</b> – create a file
SYNOPSIS	<pre>int creat   (     const char *name, /* name of the file to create */     int flag /* file permissions */   )</pre>
DESCRIPTION	This routine creates a file called <i>name</i> and opens it with a specified <i>flag</i> . This routine determines on which device to create the file; it then calls the create routine of the device driver to do most of the work. Therefore, much of what transpires is device/driver-dependent.
	The parameter <i>flag</i> is set to <b>O_RDONLY</b> (0), <b>O_WRONLY</b> (1), <b>O_RDWR</b> (2) for the duration of time the file is open.
	The parameter <i>flag</i> can be set to <b>O_SYNC</b> , on dosFs volumes, indicating that each write should be immediately written to the backing media. This flag synchronizes the FAT and the directory entries.
	On NFS and POSIX compliant file systems such as HRFS, the parameter <i>flag</i> refers instead to the UNIX style file permission bits.
NOTE	For more information about situations when there are no file descriptors available, see the reference entry for <b>iosInit()</b> .

RETURNS	A file descriptor number, or <b>ERROR</b> if a filename is not specified, the device does not exist, no file descriptors are available, or the driver returns <b>ERROR</b> .
ERRNO	ELOOP Circular symbolic link, too many links.
	EMFILE Maximum number of files already open.
	S_iosLib_DEVICE_NOT_FOUND (ENODEV) No valid device name found in path.
	others Other errors reported by device drivers.
SEE ALSO	ioLib, open()

# cret()

NAME	<b>cret( )</b> – continue until the current subroutine returns
SYNOPSIS	STATUS cret ( int taskNameOrId /* task to continue, 0 = default */ )
DESCRIPTION	This routine places a breakpoint at the return address of the current subroutine of a specified task, then continues execution of that task.
	To execute, enter:
	-> cret [task]
	If <i>task</i> is omitted or zero, the last task referenced is assumed.
	When the breakpoint is hit, information about the task will be printed in the same format as in single-stepping. The breakpoint is automatically removed when hit, or if the task hits another breakpoint first.
RETURNS	<b>OK</b> , or <b>ERROR</b> if there is no such task or the breakpoint table is full.
ERRNO	N/A
SEE ALSO	<b>dbgLib, so( ), c( ), b( )</b> , VxWorks Kernel Programmer's Guide: Kernel Shell, VxWorks Command-Line Tools User's Guide 2.2: Host Shell

# **d()**

NAME	d() – display memory
SYNOPSIS	<pre>void d   (    void * adrs, /* address to display (if 0, display next block */    int nunits, /* number of units to print (if 0, use default) */    int width /* width of displaying unit (1, 2, 4, 8) */   )</pre>
DESCRIPTION	This command displays the contents of memory, starting at <i>adrs</i> . If <i>adrs</i> is omitted or zero, <b>d()</b> displays the next memory block, starting from where the last <b>d()</b> command completed.
	Memory is displayed in units specified by <i>width</i> . If <i>nunits</i> is omitted or zero, the number of units displayed defaults to last use. If <i>nunits</i> is non-zero, that number of units is displayed and that number then becomes the default. If <i>width</i> is omitted or zero, it defaults to the previous value. If <i>width</i> is an invalid number, it is set to 1. The valid values for <i>width</i> are 1, 2, 4, and 8. The number of units <b>d()</b> displays is rounded up to the nearest number of full lines.
RETURNS	N/A
ERRNO	N/A
ERRNO	N/A
SEE ALSO	usrLib, m(), the VxWorks programmer guides.

# d0()

NAME	d0() – return the contents of register d0 (also d1 - d7) (MC680x0)	
SYNOPSIS	<pre>int d0   (    int taskId /* task ID, 0 means default task */   )</pre>	
DESCRIPTION	This command extracts the contents of register <b>d0</b> from the TCB of a specified task. If <i>taskId</i> is omitted or zero, the last task referenced is assumed.	
	Similar routines are provided for all data registers ( $d0 - d7$ ): $d0() - d7()$ .	

**RETURNS** The contents of register **d0** (or the requested register).

ERRNO Not Available

**SEE ALSO dbgArchLib**, \the VxWorks programmer guides

### dbgBpTypeBind()

NAME dbgBpTypeBind() – bind a breakpoint handler to a breakpoint type (MIPS R3000, R4000, R4650) SYNOPSIS STATUS dbgBpTypeBind ( bpType, /\* breakpoint type \*/ int FUNCPTR routine /\* function to bind \*/ ) DESCRIPTION Dynamically bind a breakpoint handler to breakpoints of type 0 - 7. By default only breakpoints of type zero are handled with the vxWorks breakpoint handler (see dbgLib). Other types may be used for Ada stack overflow or other such functions. The installed handler must take the same parameters as excExcHandle() (see excLib). RETURNS OK, or ERROR if *bpType* is out of bounds. ERRNO Not Available SEE ALSO dbgArchLib, dbgLib, excLib

### dbgHelp()

b

NAME	<b>dbgHelp( )</b> – display debugg	ging help menu
SYNOPSIS	void dbgHelp (void)	
DESCRIPTION	This routine displays a summary of <b>dbgLib</b> utilities with a short description of each, similar to the following:	
	dbgHelp dbgInit	Print this list Install debug facilities
	b	Display breakpoints

Set breakpoint

addr[,task[,count]]

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	е	addr[,eventNo[,task[,:	<pre>func[,arg]]]] Set eventpoint (WindView)</pre>
	bd	addr[,task]	Delete breakpoint
	bdall	[task]	Delete all breakpoints
	C	[task[,addr[,addr1]]]	Continue from breakpoint
	cret	[task]	Continue to subroutine return
	S	[task[,addr[,addr1]]]	Single step
	SO	[task]	Single step/step over subroutine
	1	[adr[,nInst]]	List disassembled memory
	tt	[task]	Do stack trace on task
	bh	addr[,access[,task[,co	ount[,quiet]]]] set hardware breakpoint
			(if supported by the architecture)
DETUDNO	NI / A		
NETUNINS	1N/A		
EDDNO	NI / A		
	1 <b>N</b> / A		
SEE ALSO	dbgLib VxWor	ks Kernel Proorammer's Gu	ide: Kernel Shell

# dbgInit()

NAME	<b>dbgInit( )</b> – initialize the shell debugging package
SYNOPSIS	STATUS dbgInit (void)
DESCRIPTION	This routine initializes the shell debugging package and enables the basic breakpoint and single-step functions.
	This routine also enables the shell abort function.
NOTE	The debugging package should be initialized before any debugging routines are used. If the configuration macro INCLUDE_DEBUG is defined, <b>dbgInit(</b> ) is called by the VxWorks root task at initialisation time.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the debug facility cannot be initialized.
ERRNO	N/A
SEE ALSO	dbgLib, VxWorks Kernel Programmer's Guide: Kernel Shell

#### dcacheDevCreate() dcacheDevCreate() – Create a disk cache NAME SYNOPSIS CBIO\_DEV\_ID dcacheDevCreate ( CBIO\_DEV\_ID subDev, /\* block device handle \*/ char \* pRamAddr, /\* where it is in memory (NULL = KHEAP\_ALLOC) \*/ int memSize, /\* amount of memory to use \*/ char \* pDesc /\* dowing 1 /\* device description string \*/ ) DESCRIPTION This routine creates a CBIO layer disk data cache instance. The disk cache unit accesses the disk through the subordinate CBIO device driver, provided with the *subDev* argument. A valid block device BLK\_DEV handle may be provided instead of a CBIO handle, in which case it will be automatically converted into a CBIO device by using the wrapper functionality from cbioLib. Memory which will be used for caching disk data may be provided by the caller with *pRamAddr*, or it will be allocated by **dcacheDevCreate()** from the common system memory pool, if memAddr is passed as NULL. memSize is the amount of memory to use for disk caching, if 0 is passed, then a certain default value will be calculated, based on available memory. *pDesc* is a string describing the device, used later by **dcacheShow()**, and is useful when there are many cached disk devices. A maximum of 16 disk cache devices are supported at this time. RETURNS disk cache device handle, or NULL if there is not enough memory to satisfy the request, or the *blkDev* handle is invalid. Not Available ERRNO dcacheCbio SEE ALSO

#### dcacheDevDisable()

NAME	dcacheDevDisable() – Disable the disk cache for this device
SYNOPSIS	STATUS dcacheDevDisable ( CBIO_DEV_ID dev /* CBIO device handle */ )

VxWorks Kernel API Reference, 6.6 dcacheDevEnable()

DESCRIPTION	This function disables the cache by setting the bypass count to zero and storing the old value, if there is already an old value then we won't repeat the process though. RETURNS <b>OK</b> if cache is successfully disabled or <b>ERROR</b> .
RETURNS	Not Available
ERRNO	Not Available
SEE ALSO	dcacheCbio

# dcacheDevEnable()

NAME	dcacheDevEnable( ) – Reenable the disk cache
SYNOPSIS	STATUS dcacheDevEnable ( CBIO_DEV_ID dev /* CBIO device handle */ )
DESCRIPTION	This function re-enables the cache if we disabled it. If we did not disable it, then we cannot re-enable it. RETURNS <b>OK</b> if cache is sucessfully enabled or <b>ERROR</b> .
RETURNS	Not Available
ERRNO	Not Available
SEE ALSO	dcacheCbio

### dcacheDevMemResize()

NAME	dcacheDevMemResize() – set a new size to a disk cache device
SYNOPSIS	STATUS dcacheDevMemResize ( CBIO_DEV_ID dev, /* device handle */ size_t newSize /* new cache size in bytes */ )

**DESCRIPTION** This routine is used to resize the dcache layer. This routine is also useful after a disk change event, for example a PCMCIA disk swap. The routine **pccardDosDevCreate()** in **pccardLib.c** uses this routine for that function. This should be invoked each time a new disk is inserted on media where the device geometry could possibly change. This function will re-read all device geometry data from the block driver, carve out and initialize all cache descriptors and blocks.

RETURNS **OK** or **ERROR** if the device is invalid or if the device geometry is invalid (**EINVAL**) or if there is not enough memory to perform the operation.

- RETURNS Not Available
- ERRNO Not Available
- SEE ALSO dcacheCbio

#### dcacheDevTune()

NAME	<b>dcacheDevTune( )</b> – modify tunable disk cache parameters	
SYNOPSIS	<pre>STATUS dcacheDevTune   (    CBI0_DEV_ID dev, /* device handle */    int dirtyMax, /* max # of dirty cache blocks allowed */    int bypassCount, /* request size for bypassing cache */    int readAhead, /* how many blocks to read ahead */    int syncInterval /* how many seconds between disk updates */   )</pre>	
DESCRIPTION	This function allows the user to tune some disk cache parameters to obtain better performance for a given application or workload pattern. These parameters are checked f sanity before being used, hence it is recommended to verify the actual parameters being s with <b>dcacheShow()</b> .	
	Following is the description of each tunable parameter:	
	<i>bypassCount</i> In order to achieve maximum performance, Disk Cache is bypassed for very large requests. This parameter sets the threshold number of blocks for bypassing the cache, resulting usually in the data being transferred by the low level driver directly to/from application data buffers (also known as cut-through DMA). Passing the value of 0 in this argument preserves the previous value of the associated parameter.	
	<i>syncInterval</i> The Disk Cache provides a low priority task that will update all modified blocks onto the disk periodically. This parameters controls the time between these updates in	

VxWorks Kernel API Reference, 6.6 dcacheHashTest( )

> seconds. The longer this period, the better throughput is likely to be achieved, while risking to loose more data in the event of a failure. For removable devices this interval is fixed at 1 second. Setting this parameter to 0 results in immediate writes to disk when requested, resulting in minimal data loss risk at the cost of somewhat degraded performance.

	readAhead In order to avoid accessing the disk in small units, the Disk Cache will read many contiguous blocks once a block which is absent from the cache is needed. Increasing this value increases read performance, but a value which is too large may cause blocks which are frequently used to be removed from the cache, resulting in a low Hit Ratio, and increasing the number of Seeks, slowing down performance dramatically. Passing the value of 0 in this argument preserves the pervious value of the associated parameter.
	<i>dirtyMax</i> Routinely the Disk Cache will keep modified blocks in memory until it is specifically instructed to update these blocks to the disk, or until the specified time interval between disk updates has elapsed, or until the number of modified blocks is large enough to justify an update. Because the disk is updated in an ordered manner, and the blocks are written in groups when adjacent blocks have been modified, a larger dirtyMax parameter will minimize the number of Seek operation, but a value which is too large may decrease the Hit Ratio, thus degrading performance. Passing the value of 0 in this argument preserves the pervious value of the associated parameter.
RETURNS	<b>OK</b> or <b>ERROR</b> if device handle is invalid. Parameter value which is out of range will be silently corrected.
ERRNO	Not Available

dcacheCbio, dcacheShow() SEE ALSO

### dcacheHashTest()

NAME	dcacheHashTest() – test hash table integrity
SYNOPSIS	void dcacheHashTest ( CBIO_DEV_ID dev )
DESCRIPTION	none
RETURNS	Not Available

ERRNO Not Available

SEE ALSO dcacheCbio

### dcacheShow()

NAME	dcacheShow() – print information about disk cache
SYNOPSIS	<pre>void dcacheShow   (    CBIO_DEV_ID dev, /* device handle */    int verbose /* 1 - display state of each cache block */   )</pre>
DESCRIPTION	This routine displays various information regarding a disk cache, namely current disk parameters, cache size, tunable parameters and performance statistics. The information is displayed on the standard output.
	The <i>dev</i> argument is the device handle, if it is NULL, all disk caches are displayed.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	dcacheCbio

# devs()

NAME	devs() – list all system-known devices
SYNOPSIS	void devs (void)
DESCRIPTION	This command displays a list of all devices known to the I/O system.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	usrLib, iosDevShow(), the VxWorks programmer guides.

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### dirList()

NAME	dirList() – list contents of a directory (multi-purpose)
SYNOPSIS	<pre>STATUS dirList (     int fd, /* file descriptor to write on */     const char * dirString, /* name of the directory to be listed */     BOOL doLong, /* if TRUE, do long listing */     BOOL doTree /* if TRUE, recurse into subdirs */ )</pre>
DESCRIPTION	This command is similar to UNIX ls. It lists the contents of a directory in one of two formats. If <i>doLong</i> is <b>FALSE</b> , only the names of the files (or subdirectories) in the specified directory are displayed. If <i>doLong</i> is <b>TRUE</b> , then the file name, size, date, and time are displayed. If <i>doTree</i> flag is <b>TRUE</b> , then each subdirectory encountered will be listed as well (i.e. the listing will be recursive).
	The <i>dirName</i> parameter specifies the directory to be listed. If <i>dirName</i> is omitted or <b>NULL</b> , the current working directory will be listed. <i>dirName</i> may contain wildcard characters to list some of the directory's contents.
LIMITATIONS	- With <b>dosFsLib</b> file systems, MS-DOS volume label entries are not reported.
	- Although an output format very similar to UNIX "ls" is employed, some information items have no particular meaning on some file systems.
	- Some file systems which do not support the POSIX compliant <b>dirLib()</b> interface, can not support the <i>doLong</i> and <i>doTree</i> options.
RETURNS	OK or ERROR.
ERRNO	Not Available
SEE ALSO	usrFsLib, dirLib, ls(), ll(), lsr(), llr(), the VxWorks programmer guides.

# diskFormat()

NAME	diskFormat() – format a disk with dosFs
SYNOPSIS	STATUS diskFormat ( const char * pDevName /* name of the device to initialize */ )

DESCRIPTION	This command in now obsolete. Use dosfsDiskFormat or <b>dosFsVolFormat()</b> instead
	This command formats a disk and creates the dosFs file system on it. The device must already have been created by the device driver and dosFs format component must be included.
EXAMPLE	-> diskFormat "/fd0"
RETURNS	OK, or ERROR if the device cannot be opened or formatted.
ERRNO	Not Available
SEE ALSO	usrFsLib, dosFsLib, the VxWorks programmer guides.

### diskInit()

NAME	diskInit() – initialize a file system on a block device
SYNOPSIS	STATUS diskInit ( const char *pDevName /* name of the device to initialize */ )
DESCRIPTION	This function is now obsolete.
RETURNS	Not Available
ERRNO	Not Available
SEE ALSO	usrFsLib

### dosFsCacheCreate()

NAME	dosFsCacheCreate() – create c	cache for a DosFS volume
SYNOPSIS	STATUS dosFsCacheCreate (	(*
	char * Jaha Garlan Jila	/* VOLUME Mame //
	char ^ dataCacheAddr,	/^ memory address (NULL = KHEAP_ALLOC) //
	u_int dataCacheSize,	/* size of cache */
	char * dirCacheAddr,	<pre>/* memory address (NULL = KHEAP_ALLOC) */</pre>

VxWorks Kernel API Reference, 6.6 dosFsCacheDelete()

```
u_int dirCacheSize, /* size of cache */
char * fatCacheAddr, /* memory address (NULL = KHEAP_ALLOC) */
u_int fatCacheSize /* size of cache */
)

DESCRIPTION none
RETURNS Not Available
ERRNO Not Available
SEE ALSO dosFsCacheLib
```

#### dosFsCacheDelete()

NAME	<b>dosFsCacheDelete( )</b> – delete the disk cache for a dosFs volume
SYNOPSIS	STATUS dosFsCacheDelete ( const char * volName /* dosFs volume name */ )
DESCRIPTION	This routine removes the disk cache for <i>volName</i> and frees the allocated memory if it was requested from the system memory pool.
RETURNS	STATUS.
ERRNO	Not Available

SEE ALSO dosFsCacheLib

#### dosFsCacheInfo()

**DESCRIPTION** This routine allows the user to retrieve the cache's **bypass** and **readAhead** parameters. Reading more than **bypass** sectors at once may trigger a cache flush and data will be retrieved directly from disk instead of cache. When reading data, DosFS will try to read **readAhead** sectors at a time. *volName* is a **NULL** terminated character string identifying the name of the DosFS device.

*type* identifies the volume's cache for which to retrieve information. Valid values are **DOS\_DATA\_CACHE**, **DOS\_DIR\_CACHE** and **DOS\_FAT\_CACHE**. *pSettings* points to a structure containing parameters for the **bypass** and **readAhead** values of the cache.

- RETURNS OK on success, ERROR otherwise
- ERRNO Not Available
- SEE ALSO dosFsCacheLib

#### dosFsCacheLibInit()

NAME	<b>dosFsCacheLibInit()</b> – initialize dosFsCache library.
SYNOPSIS	<pre>void dosFsCacheLibInit    (    u_int defaultDataCacheSize,    u_int defaultDirCacheSize,    u_int defaultFatCacheSize )</pre>
DESCRIPTION	none
RETURNS	N/A. /NOMANUAL
ERRNO	Not Available
SEE ALSO	dosFsCacheLib

VxWorks Kernel API Reference, 6.6 dosFsCacheOptionsGet()

# dosFsCacheOptionsGet()

NAME dosFsCacheOptionsGet() – get this dosFs volume's disk cache options

SYNOPSIS	UINT dosFsCacheOptionsGet ( char * volName /* dosFs volume name */ )
DESCRIPTION	This routine gets the cache options for the dosFs volume <i>volName</i> .
RETURNS	value of the volume's options.
ERRNO	Not Available

SEE ALSO dosFsCacheLib

### dosFsCacheOptionsSet()

NAME	dosFsCacheOptionsSet() – set this dosFs volume's disk cache options
SYNOPSIS	UINT dosFsCacheOptionsSet ( char * volName, /* dosFs volume name */ UINT options /* new options */ )
DESCRIPTION	This routine sets the cache options for the dosFs volume <i>volName</i> . Currently the only option available is <b>DOS_CACHE_VOL_NO_DMA</b> , which means that <b>dosFsCacheLib</b> does not need to use DMA-safe buffers on <i>volName</i> when large transfers to and from the media are requested. This option can gain performance on large transfers, and is recommended for USB devices.
RETURNS	new value of the volume's options.
ERRNO	Not Available
SEE ALSO	dosFsCacheLib

### dosFsCacheShow()

NAME	<b>dosFsCacheShow()</b> – show information regarding a dosFs volume's cache
SYNOPSIS	<pre>void dosFsCacheShow   (    const char * volName, /* volume name */    u_int level /* verbosity level */   )</pre>
DESCRIPTION	This routine displays information regarding an specific dosFs volume <i>volName</i> cache. If <i>level</i> is zero, a summary is displayed, otherwise the current state of the internal hash table and LRU list is displayed.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	dosFsShow

# dosFsCacheTune()

NAME	dosFsCacheTune() – tune a cache's settings
SYNOPSIS	STATUS dosFsCacheTune ( char * volName, /* Name of the DosFS volume */ DOS_CACHE_TYPE type, /* Identify which cache to tune */ DOS_CACHE_INFO * pSettings /* New cache settings */ )
DESCRIPTION	This routine allows the user to tune the cache's <b>bypass</b> and <b>readAhead</b> parameters. Reading more than <b>bypass</b> sectors at once may trigger a cache flush and data will be retrieved directly from disk instead of cache. When reading data, DosFS will try to read <b>readAhead</b> sectors at a time.
	<i>volName</i> is a <b>NULL</b> terminated character string identifying the name of the DosFS device. <i>type</i> identifies which cache associated with the volume to tune. Valid values are <b>DOS_DATA_CACHE</b> , <b>DOS_DIR_CACHE</b> and <b>DOS_FAT_CACHE</b> . <i>pSettings</i> is a structure containing parameters for the <b>bypass</b> and <b>readAhead</b> values of the cache.
RETURNS	OK on success, ERROR otherwise

VxWorks Kernel API Reference, 6.6 dosFsChkDsk()

ERRNO Not Available

SEE ALSO dosFsCacheLib

#### dosFsChkDsk()

dosFsChkDsk() - make volume integrity checking. NAME STATUS dosFsChkDsk SYNOPSIS FAST DOS\_FILE\_DESC\_ID pFd, /\* file descriptor of root dir \*/ u int params /\* check level and verbosity \*/ ) DESCRIPTION This library does not make integrity check process itself, but instead uses routine provided by dosChkLib. This routine prepares parameters and invokes checking routine via preinitialized function pointer. If dosChkLib is not configured into vxWorks, this routine returns ERROR. Ownership on device should be taken by an upper level routine. STATUS as returned by volume checking routine or RETURNS

- ERROR, if such routine is not installed.
- ERRNO S\_dosFsLib\_UNSUPPORTED.
- SEE ALSO dosFsLib

#### dosFsClose()

NAME	dosFsClose() – close a dosFs file
SYNOPSIS	STATUS dosFsClose ( DOS_FILE_DESC_ID pFd /* file descriptor pointer */ )
DESCRIPTION	This routine closes the specified dosFs file. If file contains excess clusters beyond EOF they are freed, when last file descriptor is being closed for that file.
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RETURNS	OK, or ERROR if directory couldn't be flushed or entry couldn't be found.
ERRNO	S_dosFsLib_INVALID_PARAMETER S_dosFsLib_DELETED S_dosFsLib_FD_OBSOLETE /NOMANUAL
SEE ALSO	dosFsLib

## dosFsDefaultCacheSizeSet()

NAME	dosFsDefaultCacheSizeSet() – set the default disk cache size
SYNOPSIS	<pre>void dosFsDefaultCacheSizeSet   (    UINT newDataDefaultSize, /* new default size for dosFs data cache */    UINT newDirDefaultSize, /* new default size for dosFs dir cache */    UINT newFatDefaultSize /* new default size for dosFs FAT cache */   )</pre>
DESCRIPTION	This routine sets the default disk cache size to be used for the next dosFs instantiations.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	dosFsCacheLib

## dosFsDefaultDataCacheSizeGet()

- NAME dosFsDefaultDataCacheSizeGet() get the default data cache size
- SYNOPSIS UINT dosFsDefaultDataCacheSizeGet (void)
- **DESCRIPTION** This routine gets the default data cache size.
- **RETURNS** value of the default data cache size.

ERRNO Not Available

SEE ALSO dosFsCacheLib

VxWorks Kernel API Reference, 6.6 dosFsDefaultDirCacheSizeGet()

#### dosFsDefaultDirCacheSizeGet()

NAME dosFsDefaultDirCacheSizeGet() – get the default directory cache size

- SYNOPSIS UINT dosFsDefaultDirCacheSizeGet (void)
- **DESCRIPTION** This routine gets the default directory cache size.
- **RETURNS** value of the default directory cache size.
- ERRNO Not Available

SEE ALSO dosFsCacheLib

### dosFsDefaultFatCacheSizeGet()

- NAME dosFsDefaultFatCacheSizeGet() get the default FAT cache size
- SYNOPSIS UINT dosFsDefaultFatCacheSizeGet (void)
- **DESCRIPTION** This routine gets the default FAT cache size.
- **RETURNS** value of the default FAT cache size.
- ERRNO Not Available
- SEE ALSO dosFsCacheLib

#### dosFsDevCreate()

NAME dosFsDevCreate() – create file system device.

SYNOPSIS STATUS dosFsDevCreate

**DESCRIPTION** This routine associates an XBD device with a logical I/O device name and prepare it to perform file system functions. It takes an XBD device handle, typically created by **xbdBlkDevCreate()** or **xbdPartitionDevCreate()**, and defines it as a dosFs volume. As a result, when high-level I/O operations (e.g., **open()**, **write()**) are performed on the device, the calls will be routed through **dosFsLib**. The *device* parameter is the handle of the underlying partition or block device XBD.

The argument *maxFiles* specifies the number of files that can be opened at once on the device.

The volume structure integrity can be automatically checked during volume mounting. Parameter *dosDevCreateOptions* defines checking level (**DOS\_CHK\_ONLY** or **DOS\_CHK\_REPAIR**), that can be bitwise or-ed with check verbosity level value (**DOS\_CHK\_VERB\_SILENT**, **DOS\_CHK\_VERB\_1** or **DOS\_CHK\_VERB\_2**).

If the value of *dosDevCreateOptions* is 0, the default checking level is used. The default level is (DOS\_CHK\_ONLY | DOS\_CHK\_VERB\_2).

To suppress the automatic check disk, bitwise or (**DOS\_CHK\_NONE**) or set *dosDevCreateOptions* to NONE.

Disk checking is normally suppressed on volumes marked clean. To force a disk-check, bitwise or (DOS\_CHK\_FORCE).

The volume may be configured to request **DOS\_WRITE\_THROUGH** writes for some or all of the disk operations. Additional bits of parameter *dosDevCreateOptions* define the volume's write-through setting. The default (zero) is to use copyback writes (**DOS\_WRITE**) for all write operations. The default is the fastest configuration.

To writethrough all FAT table write operations, or in DOS\_WRITE\_THROUGH\_FAT To writethrough all directory entry write operations, or in DOS\_WRITE\_THROUGH\_DIR To writethrough all user data buffers, or in DOS\_WRITE\_THROUGH\_USR To writethrough both FAT and DIRENT operations, or them together. (DOS\_WRITE\_THROUGH\_DIR | DOS\_WRITE\_THROUGH\_FAT | DOS\_CHK\_NONE)

User data writes will still use copyback **DOS\_WRITE** operations when using (**DOS\_WRITE\_THROUGH\_DIR** | **DOS\_WRITE\_THROUGH\_FAT**)

To write-through all write operations, including all user data, or in **DOS\_WRITE\_THROUGH\_ALL**. This is the slowest operation and all write operation made by the file system will be **DOS\_WRITE\_THROUGH**.

To enable Unicode filenames, or in **DOS\_FILENAMES\_UNICODE**. Case insensitivity (if enabled) currently applies only to ASCII values even when Unicode is turned on. For instance, a German eszet is never considered the same as two uppercase S characters, but two uppercase S characters can match two lowercase S characters because these are both in the first 128 character codes.

**NOTE** Setting parameter *dosDevCreateOptions* to NONE (-1) will both disable the automated chkdsk and force copyback (**DOS\_WRITE**) operation. Unicode filenames will not be enabled.

VxWorks Kernel API Reference, 6.6 dosFsDevDelete( )

Note that during a call to dosFsDevCreate() actual disk accesses are deferred to the time when **open()** or **creat()** are first called. That is also when the automatic disk checking will take place. Therefore this function will succeed in cases where a removable disk is not present in the drive. RETURNS OK, or ERROR if the device name is already in use or insufficient memory. /NOMANUAL Not Available ERRNO dosFsLib

### dosFsDevDelete()

SEE ALSO

NAME	dosFsDevDelete() – delete a dosFs volume	
SYNOPSIS	STATUS dosFsDevDelete ( DOS_VOLUME_DESC_ID pVolDesc /* pointer to volume descriptor */ )	
DESCRIPTION	This routine deletes a dosFs volume.	
RETURNS	OK on success, ERROR otherwise	
ERRNO	Not Available	
SEE ALSO	dosFsLib	

## dosFsDiskProbe()

NAME	<b>dosFsDiskProbe( )</b> – probe if a device contains a valid dosFs	
SYNOPSIS	STATUS dosFsDiskProbe ( device_t xbdDevice /* XBD device to probe */ )	
DESCRIPTION	This routine probes if a device (or a partition) contains a valid DOS	

FS.

RETURNS	OK if successful. ERROR otherwise.
	/NOMANUAL
ERRNO	Not Available
SEE ALSO	dosFsLib

# dosFsFdFree()

NAME	dosFsFdFree() – free a file descriptor
SYNOPSIS	void dosFsFdFree ( DOS_FILE_DESC_ID pFd )
DESCRIPTION	This routine marks a file descriptor as free and decreases reference count of a referenced file handle.
RETURNS	N/A. /NOMANUAL
ERRNO	Not Available
SEE ALSO	dosFsLib

# dosFsFdGet()

NAME	<b>dosFsFdGet( )</b> – get an available file descriptor	
SYNOPSIS	DOS_FILE_DESC_ID dosFsFdGet ( DOS_VOLUME_DESC_ID pVolDesc )	
DESCRIPTION	This routine obtains a free dosFs file descriptor.	
RETURNS	Pointer to file descriptor, or NULL, if none available.	

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VxWorks Kernel API Reference, 6.6 dosFsFmtLiblnit()

ERRNO S\_dosFsLib\_NO\_FREE\_FILE\_DESCRIPTORS /NOMANUAL

SEE ALSO dosFsLib

## dosFsFmtLibInit( )

NAME	dosFsFmtLibInit() – initialize the MS-DOS formatting library	
SYNOPSIS	void dosFsFmtLibInit(void)	
DESCRIPTION	This function is called to optionally enable the formatting functionality from <b>dosFsLib</b> .	
RETURNS	Not Available	
ERRNO	Not Available	
SEE ALSO	dosFsFmtLib, dosFsLib, /NOMANUAL	

## dosFsFmtTest()

NAME	dosFsFmtTest() – UNITEST CODE
SYNOPSIS	void dosFsFmtTest ( int size )
DESCRIPTION	/NOMANUAL
RETURNS	Not Available
ERRNO	Not Available
SEE ALSO	dosFsFmtLib

# dosFsHdlrInstall()

NAME	dosFsHdlrInstall() – install handler.	
SYNOPSIS	<pre>STATUS dosFsHdlrInstall   (    DOS_HDLR_DESC_ID hdlrsList, /* appropriate list */    DOS_HDLR_DESC_ID hdlr /* ptr on handler descriptor */   )</pre>	
DESCRIPTION	This library does not directly access directory structure, nor FAT, rather it uses particular handlers to serve such accesses. This function is intended for use by the <b>dosFsLib</b> sub-modules only.	
	This routine installs a handler into DOS FS handlers list. There are two such lists: FAT Handlers List ( <i>dosFatHdlrsList</i> ) and Directory Handlers List ( <i>dosDirHdlrsList</i> ). Each handler must provide its unique Id (see <b>dosFsLibP.h</b> ) and pointer to appropriate list to install it to. All lists are sorted by Id-s in ascending order. Every handler is tried to be mounted on each new volume in accordance to their order in list, until succeeded. So preferable handlers, that supports the same type of volumes must have less Id values.	
RETURNS	STATUS.	
	/NOMANUAL	
ERRNO	Not Available	
SEE ALSO	dosFsLib	

# dosFsIoctl()

NAME	<b>dosFsIoctl( )</b> – do device specific cont	trol function
SYNOPSIS	STATUS dosFsIoctl ( FAST DOS_FILE_DESC_ID pFd, int functi int arg )	<pre>/* fd of file to control */ .on, /* function code */ /* some argument */</pre>
DESCRIPTION	This routine performs the following is	octl functions.

Any ioctl function codes, that are not supported by this routine are passed to the underlying XBD module for handling.

VxWorks Kernel API Reference, 6.6 dosFsLastAccessDateEnable()

There are some **ioctl()** functions, that suppose to receive as result a 32-bit numeric value (FIONFREE, FIOWHERE and so on), however disks and files with size grater, than 4GB are supported. In order to solve this contradiction new **ioctl()** functions are provided. They have the same name as basic functions, but with suffix **64**: FIONFREE64, FIOWHERE64 and so on. These functions gets pointer to **long long** as an argument. Also FIOWHERE64 returns value via argument, but not as **ioctl()** returned value. If an ioctl fails, the task's status (see **errnoGet()**) indicates the nature of the error.

- **RETURNS** OK or current position in file for FIOWHERE, or **ERROR** if function failed or driver returned error, or if function supposes 32 bit result value, but actual result overloads this restriction.
- ERRNO S\_dosFsLib\_INVALID\_PARAMETER S\_dosFsLib\_VOLUME\_NOT\_AVAILABLE S\_dosFsLib\_FD\_OBSOLETE S\_dosFsLib\_DELETED S\_dosFsLib\_32BIT\_OVERFLOW /NOMANUAL

SEE ALSO dosFsLib

### dosFsLastAccessDateEnable( )

NAME	dosFsLastAccessDateEnable() – enable last access date updating for this volume	
SYNOPSIS	<pre>STATUS dosFsLastAccessDateEnable   (    DOS_VOLUME_DESC_ID dosVolDescId, /* dosfs volume ID to alter */    BOOL enable /* TRUE = enable update, FALSE = disable update */   )</pre>	
DESCRIPTION	This function enables or disables updating of the last access date directory entry field on open-read-close operations for the given dosFs volume. The last access date file indicates the last date that a file has been read or written. When the optional last access date field update is enabled, read operations on a file will cause a write to the media.	
RETURNS	OK or ERROR if the volume is invalid or enable is not TRUE or FALSE.	
ERRNO	Not Available	
SEE ALSO	dosFsLib	

### dosFsLibInit( )

dosFsLibInit() - prepare to use the dosFs library NAME SYNOPSIS STATUS dosFsLibInit ( int maxFiles, int options ) This routine initializes the dosFs library. This initialization is enabled when the DESCRIPTION configuration macro INCLUDE\_DOSFS is defined. This routine installs dosFsLib as a driver in the I/O system driver table, and allocates and sets up the necessary structures. The driver number assigned to **dosFsLib** is placed in the global variable *dosFsDrvNum*. OK or ERROR, if driver can not be installed. RETURNS Not Available ERRNO SEE ALSO dosFsLib

## dosFsMonitorDevCreate( )

NAME	<b>dosFsMonitorDevCreate( )</b> – create a dosFs volume through the fs monitor
SYNOPSIS	<pre>STATUS dosFsMonitorDevCreate   (    device_t xbdId, /* XBD for the device on which to mount. */    char * pDevName /* Name of the DOS FS device (mount point). */ )</pre>
DESCRIPTION	This routine creates an DOS FS device.
RETURNS	OK if successful. ERROR otherwise. /NOMANUAL
ERRNO	Not Available
SEE ALSO	dosFsLib

	dosFsOpen()
NAME	dosFsOpen() – open a file on a dosFs volume
SYNOPSIS	<pre>DOS_FILE_DESC_ID dosFsOpen   (    DOS_VOLUME_DESC_ID pVolDesc, /* pointer to volume descriptor */    char * pPath, /* dosFs full path/filename */    int flags, /* file open flags */    int mode /* file open permissions (mode) */ )</pre>
DESCRIPTION	<ul> <li>This routine opens the file <i>name</i> with the specified mode (O_RDONLY/O_WRONLY/O_RDWR/CREATE/TRUNC). The directory structure is searched, and if the file is found a dosFs file descriptor is initialized for it. Extended flags are provided by DOS FS for more efficiency:</li> <li>DOS_O_CONTIG_CHK - to check file for contiguity.</li> <li>DOS_O_CASENS - force the file name lookup in case insensitive manner, (if directory format provides such opportunity)</li> <li>If this is the very first open for the volume, configuration data will be read from the disk automatically (via dosFsVolMount()).</li> </ul>
RETURNS	A pointer to a dosFs file descriptor, or <b>ERROR</b> if the volume is not available, or there are no available dosFs file descriptors, or there is no such file and <b>O_CREAT</b> was not specified, or file can not be opened with such permissions.
ERRNO	S_dosFsLib_INVALID_PARAMETER S_dosFsLib_READ_ONLY S_dosFsLib_FILE_NOT_FOUND S_dosFsLib_FILE_EXISTS /NOMANUAL
SEE ALSO	dosFsLib

# dosFsShow()

NAME	<b>dosFsShow()</b> – display dosFs volume configuration data.
SYNOPSIS	STATUS dosfsShow (
	void * pDevName, /* name of device */

```
u_int level /* detail level */
)
```

**DESCRIPTION** This routine obtains the dosFs volume configuration for the named device, formats the data, and displays it on the standard output.

If no device name is specified, the current default device is described.

**RETURNS** OK or ERROR, if no valid device specified.

ERRNO Not Available

SEE ALSO dosFsShow

# dosFsVolDescGet()

NAME	<b>dosFsVolDescGet()</b> – convert a device name into a DOS volume descriptor pointer.
SYNOPSIS	DOS_VOLUME_DESC_ID dosFsVolDescGet ( void * pDevNameOrPVolDesc, /* device name or pointer to dos vol desc */ u_char ** ppTail /* return ptr for name, used in iosDevFind */ ) This routine validates <i>pDevNameOrPVolDesc</i> to be a DOS volume descriptor pointer else a path to a DOS device. This routine uses the standard iosLib function iosDevFind() to
	obtain a pointer to the device descriptor. If device is eligible, <i>ppTail</i> is filled with the pointer to the first character following the device name. Note that ppTail is passed to <b>iosDevFind()</b> . <i>ppTail</i> may be passed as <b>NULL</b> , in which case it is ignored.
RETURNS	A DOS_VOLUME_DESC_ID or NULL if not a DOSFS device.
ERRNO	S_dosFsLib_INVALID_PARAMETER
SEE ALSO	dosFsLib

### dosFsVolFormat( )

NAME

dosFsVolFormat() – format an MS-DOS compatible volume

SYNOPSIS

**DESCRIPTION** This utility routine performs the initialization of file system data structures on a disk. It supports FAT12 for small disks, FAT16 for medium size and FAT32 for large volumes. The *device* argument is a device name known to the I/O system.

The *opt* argument is a bit-wise or'ed combination of options controlling the operation of this routine as follows:

#### DOS\_OPT\_DEFAULT

If the current volume boot block is reasonably intact, use existing parameters, else calculate parameters based only on disk size, possibly reusing only the volume label and serial number.

#### DOS\_OPT\_PRESERVE

Attempt to preserve the current volume parameters even if they seem to be somewhat unreliable.

#### DOS\_OPT\_BLANK

Disregard the current volume parameters, and calculate new parameters based only on disk size.

#### DOS\_OPT\_QUIET

Do not produce any diagnostic output during formatting.

#### DOS\_OPT\_FAT16

Format the volume with FAT16. Valid on volumes up to 2 Gbytes big. For larger ones, FAT32 must be used.

#### DOS\_OPT\_FAT32

Format the volume with FAT32, even if the disk is smaller than 2 Gbytes, but is larger then 512 Mbytes.

#### DOS\_OPT\_VXLONGNAMES

Note that this option is deprecated. Calling dosFsVolFormat with this option will result in an error.

The third argument, *pPromptFunc* is an optional pointer to a function that may interactively prompt the user to change any of the modifiable volume parameters before formatting:

void formatPromptFunc( DOS\_VOL\_CONFIG \*pConfig );

The \**pConfig* structure upon entry to **formatPromptFunc( )** will contain the initial volume parameters, some of which can be changed before it returns. *pPromptFunc* should be **NULL** if no interactive prompting is required.

- **COMPATIBILITY** Although this routine tries to format the disk to be compatible with Microsoft implementations of the FAT and FAT32 file systems, there may be differences which are not under WRS control. For this reason, it is highly recommended that any disks which are expected to be interchanged between vxWorks and Windows should be formatted under Windows to provide the best interchangeability. The WRS implementation is more flexible, and should be able to handle the differences when formatting is done on Windows, but Windows implementations may not be able to handle minor differences between their implementation and ours.
- **AVAILABILITY** This function is an optional part of the MS-DOS file system, and may be included in a target system if it is required to be able to format new volumes.
- **RETURNS** OK or ERROR if was unable to format the disk.
- ERRNO Not Available
- SEE ALSO dosFsFmtLib

### dosFsVolFormatFd()

dosFsVolFormatFd() – format an MS-DOS compatible volume via an opened FD NAME SYNOPSIS STATUS dosFsVolFormatFd ( fd, /\* rawFs file descriptor \*/ int int opt, /\* bit-wise or'ed options \*/ FUNCPTR pPromptFunc /\* interactive parameter change callback \*/ ) DESCRIPTION Refer to **dosFsVolFormat()**'s documentation. It should also be noted that the file descriptor parameter will be closed regardless of the routine's outcome. RETURNS **OK** or **ERROR** if was unable to format the disk. ERRNO Not Available

SEE ALSO dosFsFmtLib

## dosFsVolIsFat12()

dosFsVolIsFat12() – determine if a MSDOS volume is FAT12 or FAT16 NAME int dosFsVolIsFat12 SYNOPSIS u\_char \* pBootBuf /\* boot parameter block buffer \*/ ) DESCRIPTION This routine is the container for the logic which determines if a dosFs volume is using FAT12 or FAT16. Two methods are implemented. Both methods use information from the volumes boot parameter block fields found in the boot sector. The first FAT determination method follows the recommendations outlined in the Microsoft document: "Hardware White Paper Designing Hardware for Microsoft Operating Systems FAT: General Overview of On-Disk Format Version 1.02, May 5, 1999 Microsoft Corporation" This method is used in the hopes that greater compatability with MSDOS formatted media will be achieved. The Microsoft recommended method for FAT type determination between FAT12 and FAT16 is done via the count of clusters on the volume. The Microsoft recommended approach is as follows: 1.) Determine the count of sectors occupied by the root directory entries for this volume, rounding up: rootDirSecs = ((rootEntCount \* dirEntSz) + (bytesPerSec-1)) / bytesPerSec; Where dirEntSz is 32 for MSDOS 8.3, and 64 for VXLONGNAMES. 2.) Determine the count of sectors occupied by the volumes data region: dataRgnSecs = totalSecs - (reservedSecs + (nFats \* fatSecs) + rootDirSecs); 3.) determine the count of clusters, rounding down: countOfClusts = dataSecs / secsPerClust; /\* Note: this rounds down. \*/ Note: countOfClusts represents the count of data clusters, starting at two. 4.) determine the FAT types based on the count of clusters on the volume, if (countOfClusts < 4085) /\* Microsoft recommends using "less than" \*/ /\* Volume is FAT12 \*/

```
}
else
{
/* Volume is FAT16 */
}
```

An alternate method is used when mounting a known VxWorks DOSFS-1.0 volume. This method is used for greater backward compatability with VxWorks DOSFS-1.0 volumes. See also: SPR#34704. The VxWorks dosFs1 method deviates from the Microsoft currently recommened method.

This is the VxWorks DOSFS1 method per **dosFsVolDescFill()**, **dosFsLib.c**, dosFs 1.0, revision history: "03l,16mar99,dgp". Using the identical method here will help ensure backward compatablitity when mounting volumes formatted by the VxWorks dosFs1.0 code.

The VxWorks DOSFS 1.0 approach is as follows:

1.) Get starting sector of the root directory:

```
rootSec = reservedSecs + (nFats * secsPerFat);
```

2.) Get the size of the root dir in bytes:

rootBytes = (nRootEnts \* dirEntSz):

Where dirEntSz is 32 for MSDOS 8.3, and 64 for VXLONGNAMES.

3.) Get the starting sector of the data area:

dataSec = rootSec + ((rootBytes + bytesPerSec-1) / bytesPerSec);

4.) Get the number of "FAT entries":

```
countOfClusts =
```

(((totalSecs - dataSecs) / secsPerClust) + DOS\_MIN\_CLUST);

5.) Choose the FAT type based on the count of clusters, note DOSFS1 uses less than or equal here.

```
if (countOfClusts <= 4085) /* VxDosFs1 uses less than or equal to. */
```

VxWorks Kernel API Reference, 6.6	
dosFsVolUnmount( )	

	By mimicking the dosFs 1.0 approach, we should be able to mount all dosFs 1.0 volumes correctly. By using the microsoft recommened approach in all other cases, we should be as compatable as possible with Microsoft OS's.
	The volumes Boot Parameter Block fields MUST be validated for sanity before this routine is called.
	pBootBuf is not verified, DO NOT pass this routine a <b>NULL</b> pointer. This routine is also used by <b>dosFsFmtLib.c</b>
RETURNS	<b>TRUE</b> if the FAT type is FAT12, <b>FALSE</b> if the FAT type is FAT16, or <b>ERROR</b> if the data is invalid.
	/NOMANUAL
ERRNO	Not Available
SEE ALSO	dosFsLib

# dosFsVolUnmount()

NAME	dosFsVolUnmount() – unmount a dosFs volume
SYNOPSIS	<pre>STATUS dosFsVolUnmount   (    void * pDevNameOrPVolDesc /* device name or ptr to */</pre>
DESCRIPTION	This routine is called when I/O operations on a volume are to be discontinued. This is the preferred action prior to changing a removable disk.
	All buffered data for the volume is written to the device (if possible, with no error returned if data cannot be written), any open file descriptors are marked as obsolete, and the volume is marked as not currently mounted.
	When a subsequent <b>open()</b> operation is initiated on the device, new volume will be mounted automatically.
	Once file descriptors have been marked as obsolete, any attempt to use them for file operations will return an error. (An obsolete file descriptor may be freed by using <b>close()</b> . The call to <b>close()</b> will return an error, but the descriptor will in fact be freed).
	This routine may also be invoked by calling ${\bf ioctl}($ ) with the FIOUNMOUNT function code.
	This routine must not be called from interrupt level.

SEE ALSO	dosFsLib
ERRNO	Not Available
	/NOMANUAL
RETURNS	OK, or ERROR if the volume was not mounted.

# dosFsVolumeOptionsGet()

NAME	dosFsVolumeOptionsGet() – get this volume's disk options
SYNOPSIS	UINT dosFsVolumeOptionsGet ( char * volName /* dosFs volume name */ )
DESCRIPTION	This routine gets the volume options for the dosFs volume <i>volName</i> . It replaces the routine <b>dosFsCacheOptionsGet()</b> .
	volName is a NULL terminated character string identifying the DosFS volume.
RETURNS	<b>DOS_VOLUME_VOL_NO_DMA</b> if that option is enabled; 0 if no options are enabled; <b>ERROR</b> on failure
ERRNO	Not Available
SEE ALSO	dosFsLib

# dosFsVolumeOptionsSet()

NAME	<pre>dosFsVolumeOptionsSet() - set this volume's disk options</pre>
SYNOPSIS	UINT dosFsVolumeOptionsSet ( char * volName, /* dosFs volume name */ UINT options /* new options */ )

VxWorks Kernel API Reference, 6.6 dosFsXbdBlkCopy()

DESCRIPTION	This routine sets the volume options for the dosFs volume <i>volName</i> . Currently the only option available is <b>DOS_VOLUME_VOL_NO_DMA</b> , which means that DosFS does not need to use DMA-safe buffers on <i>volName</i> when large transfers to and from the media are requested. This option may gain performance on large transfers, and is recommended for USB devices. It replaces the routine <b>dosFsCacheOptionsSet()</b> .
	<i>volName</i> is a <b>NULL</b> terminated character string that corresponds to the name of the desired DosFS volume. Valid settings for <i>options</i> are currently <b>DOS_VOLUME_VOL_NO_DMA</b> or 0.
RETURNS	new value of the volume's options, or <b>ERROR</b> on failure.
ERRNO	Not Available
SEE ALSO	dosFsLib

# dosFsXbdBlkCopy()

NAME	<b>dosFsXbdBlkCopy()</b> – copy blocks on the underlying XBD block device.
SYNOPSIS	<pre>STATUS dosFsXbdBlkCopy  (     DOS_VOLUME_DESC * pVolDesc, /* volume descriptor */     sector_t srcBlock,     sector_t dstBlock,     sector_t numBlocks  )</pre>
DESCRIPTION	none
RETURNS	STATUS. /NOMANUAL
ERRNO	Not Available
SEE ALSO	dosFsLib

# dosFsXbdBlkRead()

NAME	dosFsXbdBlkRead() – read blocks from the underlying XBD block device.
SYNOPSIS	<pre>STATUS dosFsXbdBlkRead   (    DOS_VOLUME_DESC * pVolDesc, /* volume descriptor */    sector_t startBlock, /* starting block of transfer */    sector_t numBlocks, /* number of blocks to transfer */    addr_t buffer /* address of the memory buffer */ )</pre>
DESCRIPTION	none
RETURNS	STATUS. /NOMANUAL
ERRNO	Not Available
SEE ALSO	dosFsLib

## dosFsXbdBlkWrite()

NAME	<b>dosFsXbdBlkWrite()</b> – write blocks to the underlying XBD block device.
SYNOPSIS	<pre>STATUS dosFsXbdBlkWrite   (    DOS_VOLUME_DESC * pVolDesc,  /* volume descriptor */    sector_t startBlock,  /* starting block of write */    sector_t numBlocks,  /* number of blocks to write */    addr_t buffer,  /* address of the memory buffer */    DOS_RW operation /* DOS_WRITE/DOS_WRITE_THROUGH */   )</pre>
DESCRIPTION	none
RETURNS	STATUS. /NOMANUAL
ERRNO	Not Available
SEE ALSO	dosFsLib

# dosFsXbdBytesRW()

NAME	<b>dosFsXbdBytesRW()</b> – read/write bytes to/from the underlying XBD block device.
SYNOPSIS	<pre>STATUS dosFsXbdBytesRW   (     DOS_VOLUME_DESC * pVolDesc, /* volume descriptor */     sector_t startBlock, /* starting block of the transfer */     off_t offset, /* offset into block in bytes */     addr_t buffer, /* address of data buffer */     size_t nBytes, /* number of bytes to transfer */     u_int operation /* DOS_READ/DOS_WRITE/WRITE_THROUGH */    )</pre>
DESCRIPTION	none
RETURNS	STATUS. /NOMANUAL
ERRNO	Not Available
SEE ALSO	dosFsLib

# dosFsXbdIoctl()

NAME	dosFsXbdIoctl() – Misc control operations
SYNOPSIS	STATUS dosFsXbdIoctl ( DOS_VOLUME_DESC * pVolDesc, UINT32 command, addr_t arg )
DESCRIPTION	This performs the requested old CBIO <b>ioctl()</b> operations. RETURNS <b>OK</b> or <b>ERROR</b> and may otherwise set errno. /NOMANUAL
RETURNS	Not Available
ERRNO	Not Available

SEE ALSO dosFsLib

#### dosPathParse()

**NAME** dosPathParse() – parse a full pathname into an array of names.

size\_t sizeArray

```
SYNOPSIS int dosPathParse
(
u_char * path,
PATH_ARRAY * pnamePtrArray,
```

**DESCRIPTION** This routine is similar to **pathParse( )**, but on the contrary it does not allocate additional buffers nor changes the path string.

Parses a path in directory tree which has directory names separated by / or "s. It fills the supplied array of structures with pointers to directory and file names and correspondence name length. All occurrences of *I*/, and .. are right removed from path. All tail dots and spaces are broken from each name, that is name like "abc. .." is treated as just "abc".

For instance, "/usr/vw/data/../dir/file" gets parsed into



Note that UTF-8 bytes that are not representing ASCII characters ., /, etc., never compare equal to ., /, etc., so that no special work is required here for Unicode.

VxWorks Kernel API Reference, 6.6 dosSetVolCaseSens()

In the future, the "../" erasure trick is likely to vanish. Do not rely on the fact that you can currently access "**nosuchdir/../file.txt**" when "nosuchdir" does not exist.

- **RETURNS** number of levels in path.
- ERRNO S\_dosFsLib\_ILLEGAL\_PATH S\_dosFsLib\_ILLEGAL\_NAME
- SEE ALSO dosFsLib

## dosSetVolCaseSens()

NAME	dosSetVolCaseSens() – set case sensitivity of volume
SYNOPSIS	STATUS dosSetVolCaseSens ( DOS_VOLUME_DESC_ID pVolDesc, BOOL sensitivity )
DESCRIPTION	Pass <b>TRUE</b> to setup a case sensitive volume. Pass <b>FALSE</b> to setup a case insensitive volume. Note this affects rename lookups only.
RETURNS	TRUE if pVolDesc pointed to a DOS volume.
ERRNO	Not Available
SEE ALSO	dosFsLib

# dosfsDiskFormat()

NAME	dosfsDiskFormat() – format a disk with dosFs
SYNOPSIS	STATUS dosfsDiskFormat ( const char * pDevName /* name of the device to initialize */ )

DESCRIPTION	This command formats a disk and creates the dosFs file system on it. The device must already have been created by the device driver and dosFs format component must be included.
EXAMPLE	-> dosfsDiskFormat "/fd0"
RETURNS	OK, or ERROR if the device cannot be opened or formatted.
ERRNO	Not Available

**SEE ALSO** usrFsLib, dosFsLib, the VxWorks programmer guides.

## dosfsDiskToHost16()

NAME	<pre>dosfsDiskToHost16() - convert uint16_t from on-disk to host format</pre>
SYNOPSIS	uint16_t dosfsDiskToHost16 ( uint8_t * pSrc )
DESCRIPTION	This routine converts a uint16_t from on-disk format to host's endian-ness.
RETURNS	uint16_t in hosts's endian-ness /NOMANUAL
ERRNO	Not Available
SEE ALSO	dosFsLib

## dosfsDiskToHost32()

NAME	dosfsDiskToHost32() – convert uint32_t from on-disk to host format
SYNOPSIS	uint32_t dosfsDiskToHost32 ( uint8_t * pSrc )

**DESCRIPTION** This routine converts a uint32\_t from on-disk format to host's endian-ness.

VxWorks Kernel API Reference, 6.6 dosfsHostToDisk16()

SEE ALSO	dosFsLib
ERRNO	Not Available
	/NOMANUAL
RETURNS	uint32_t in host's endian-ness

# dosfsHostToDisk16()

NAME	<pre>dosfsHostToDisk16( ) - convert uint16_t from host to on-disk format</pre>
SYNOPSIS	<pre>void dosfsHostToDisk16   (     uint16_t src,     uint8_t * pDest )</pre>
DESCRIPTION	This routine converts a uint16_t from host's memory to on-disk format.
RETURNS	N/A /NOMANUAL
ERRNO	Not Available
SEE ALSO	dosFsLib

# dosfsHostToDisk32()

NAME	dosfsHostToDisk32() – convert uint32_t from host to on-disk format
SYNOPSIS	<pre>void dosfsHostToDisk32   (     uint32_t src,     uint8_t * pDest   )</pre>
DESCRIPTION	This routine converts a uint32_t from host's memory to on-disk format.
RETURNS	N/A

#### /NOMANUAL

ERRNO Not Available

SEE ALSO dosFsLib

### dpartDevCreate()

NAME dpartDevCreate() – Initialize a partitioned disk

**DESCRIPTION** To handle a partitioned disk, this function should be called, with *subDev* as the handle returned from **dcacheDevCreate()**, It is recommended that for efficient operation a single disk cache be allocated for the entire disk and shared by its partitions.

*nPart* is the maximum number of partitions which are expected for the particular disk drive. Up to 24 (C-Z) partitions per disk are supported.

#### PARTITION DECODE FUNCTION

An external partition table decode function is provided via the *pPartDecodeFunc* argument, which implements a particular style and format of partition tables, and fill in the results into a table defined as Pn array of **PART\_TABLE\_ENTRY** types. See **dpartCbio.h** for definition of **PART\_TABLE\_ENTRY**. The prototype for this function is as follows:

STATUS parDecodeFunc ( (CBIO\_DEV\_ID dev, /\* device from which to read blocks \*/ PART\_TABLE\_ENTRY \*pPartTab, /\* table where to fill results \*/ int nPart /\* # of entries in <pPartTable> \*/ )

**RETURNS** CBIO\_DEV\_ID or NULL if error creating CBIO device.

ERRNO Not Available

SEE ALSO dpartCbio, dosFsDevCreate().

# dpartPartGet()

NAME	<b>dpartPartGet( )</b> – retrieve handle for a partition	
SYNOPSIS	CBIO_DEV_ID dpartPartGet ( CBIO_DEV_ID masterHandle, /* CBIO handle of the master partition */ int partNum /* partition number from 0 to nPart */ )	
DESCRIPTION	This function retrieves a CBIO handle into a particular partition of a partitioned device. This handle is intended to be used with <b>dosFsDevCreate( )</b> .	
RETURNS	<b>CBIO_DEV_ID</b> or <b>NULL</b> if partition is out of range, or <i>masterHandle</i> is invalid.	
ERRNO	Not Available	
SEE ALSO	dpartCbio, dosFsDevCreate()	

# dshmMuxHwAddrToOff()

NAME	dshmMuxHwAddrToOff() – translate a local address to a shared memory offset	
SYNOPSIS	uint32_t dshmMuxHwAddrToOff ( const uint_t hw, const void * const addr )	
DESCRIPTION	This routine takes a pointer in shared memory and translates it into a shared memory offset that can be passed via a message to a remote node. This is needed when nodes do not see the shared memory at the same address locally.* It allows then to pass addresses as offset from a common point of reference.	
RETURNS	the shared memory offset corresponding to the local address	
ERRNO	Not Available	
SEE ALSO	dshmMuxLib	

# dshmMuxHwGet()

NAME	dshmMuxHwGet() – obtain an hardware registration handle based on name
SYNOPSIS	int dshmMuxHwGet ( const char * const name )
DESCRIPTION	This routine returns the handle of a previously registered hardware based on the name used at registration time.
RETURNS	a hardware registration handle, or -1 if it does not exist
ERRNO	Not Available
SEE ALSO	dshmMuxLib

# dshmMuxHwLocalAddrGet()

NAME	dshmMuxHwLocalAddrGet() – obtain address of the local node
SYNOPSIS	uint16_t dshmMuxHwLocalAddrGet ( const uint_t hw )
DESCRIPTION	none
RETURNS	the address of the local node, 0x0 to 0xfffe, 0xffff if error
ERRNO	Not Available
SEE ALSO	dshmMuxLib

## dshmMuxHwNodesNumGet()

NAME	dshmMuxHwNodesNumGet() – obtain the maximum number of nodes on a hardware bus
SYNOPSIS	<pre>int dshmMuxHwNodesNumGet   (     const uint_t hw )</pre>
DESCRIPTION	This routine returns the maximum number of nodes that can exist on a particular hardware bus registered with the MUX.
RETURNS	the maximum number of nodes, or -1 if $hw$ is invalid
ERRNO	Not Available
SEE ALSO	dshmMuxLib

# dshmMuxHwOffToAddr()

NAME	dshmMuxHwOffToAddr() – translate a shared memory offset to a local address
SYNOPSIS	<pre>void * dshmMuxHwOffToAddr   (     const uint_t hw,     const uint32_t offset   )</pre>
DESCRIPTION	This routine takes an offset in shared memory that might have been transmitted via a message and translates it to a local address that can be used as a pointer.
RETURNS	the local address corresponding to the shared memory offset
ERRNO	Not Available
SEE ALSO	dshmMuxLib

	dshmMuxHwRegister()	
NAME	dshmMuxHwRegister() – register a hardware bus with the MUX	
SYNOPSIS	<pre>int dshmMuxHwRegister   (     const DSHM_HW_ID id, /* bus ID */     const char * const name, /* bus name */     const uint_t maxNodes, /* max number of nodes on the bus */     const DSHM_HW_HOOKS * const pHooks /* bus methods (see dshmMuxLib.h) */     )</pre>	
DESCRIPTION	This routine registers a bus supporting DSHM with the MUX, allowing peers to register and thus services to communicate between each other over the bus.	
	An OS-dependant hardware bus <i>id</i> must be provided, as well as a bus <i>name</i> , both of which must be unique on the local target. This routine pre- allocates memory for the potential services in the system. Finally, entry points in the bus controller <i>pHooks</i> are registered at this point as well.	
	These hooks include: shared memory allocation and freeing, message transmission, message broadcasting, atomic set (test-and-set, compare-and-swap, read-modify-write), atomic clear (might be left to NULL if not needed), shared memory offset to local address translation and vice-versa, retrieval of local node address on the bus, allocation of virtual memory region on the bus, and a fast copy routine.	
RETURNS	a hardware registration handle, or -1 if failure	
ERRNOS	S_dshm_MUX_HW_NAME_EXISTS This bus name is already in use.	
	S_dshm_MUX_HW_TABLE_FULL	
	and memory allocation failure errnos.	
SEE ALSO	dshmMuxLib	

## dshmMuxHwTasClearGet()

NAME dshmMuxHwTasClearGet() – obtain the TAS clear routine on this bus

SYNOPSIS DSHM\_TAS\_CLEAR dshmMuxHwTasClearGet

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VxWorks Kernel API Reference, 6.6 dshmMuxHwTasGet()

```
(
const uint_t hw
)
```

SEE ALSO	dshmMuxLib, dshm/adapt/types.h
ERRNO	Not Available
RETURNS	the atomic set routine, NULL if invalid hardware identifier
DESCRIPTION	none

### dshmMuxHwTasGet()

NAME	dshmMuxHwTasGet() – obtain the test-and-set routine on this bus
SYNOPSIS	DSHM_TAS dshmMuxHwTasGet ( const uint_t hw )
DESCRIPTION	none
RETURNS	the atomic set routine, NULL if invalid hardware identifier
ERRNO	Not Available
SEE ALSO	dshmMuxLib, dshm/adapt/types.h

## dshmMuxLibInit()

```
NAME dshmMuxLibInit() - initialize the DSHM MUX

SYNOPSIS void dshmMuxLibInit

(

const uint_t maxHwReg, /* max hw types */

const uint_t maxSvc /* max services */

)
```

DESCRIPTIONPrepare the mux to receive registration for hardware buses and services. The maximum<br/>amount for each is a configuration parameter. Services cannot register before their<br/>hardware bus does a prior registration.<br/>The maximum number of hardware buses maxHwReg and services maxSvc can be\\*\* any<br/>positive integer number, but in practice should always be very small numbers.RETURNSN/AERRNONot AvailableSEE ALSOdshmMuxLib

## dshmMuxMemAlloc()

NAME	dshmMuxMemAlloc() – allocate shared memory from a specific hardware
SYNOPSIS	<pre>void * dshmMuxMemAlloc   (     const uint_t hw,     int * const pSize,     const int min  )</pre>
DESCRIPTION	This routine allocates a specific slab of memory owned by specified hardware $hw$ . The allocated slab will be of size * <i>pSize</i> or closest available if no slab of contiguous memory of greater or equal size is available. Slab is guaranteed to be at least <i>min</i> in size. To allow allocation of any size available, a value of zero can be passed to <i>min</i> .
NOTE	Depending on the hardware, a certain alignment may be forced by the underlying allocation mechanism.
RETURNS	Address of slab if successful or NULL otherwise.
ERRNOS	Any error from the installed allocation routine for the hardware.
SEE ALSO	dshmMuxLib

## dshmMuxMemFree()

NAME dshmMuxMemFree() – free allocated shared memory from a specific hardware

SYNOPSIS	STATUS dshmMuxMemFree
	(
	const uint_t hw,
	void * const pMem
	)

**DESCRIPTION** This routine frees a previously allocated slab of memory owned by specified hardware *hw*.

**RETURNS** OK if slab and hardware are valid, **ERROR** otherwise.

**ERRNOS** Any error from the installed free routine for the hardware.

SEE ALSO dshmMuxLib

## dshmMuxMsgRecv()

NAME	dshmMuxMsgRecv() – receive a message		
SYNOPSIS	<pre>STATUS dshmMuxMsgRecv (     const uint_t hw, /* hardware on which to receive */     DSHM(msg), /* message container */     int unused1, /* possibe future expansion: use 0 */     int unused2 /* possibe future expansion: use 0 */ )</pre>		
DESCRIPTION	This routine processes a message and calls the specified service processing routine previously installed.		
	This routine should only be called by the hardware receive loop, called either via interrupt or in poll mode.		
	When this routine calls the service handler, it owns the reference lock on the service object. The service handler is responsible for unlocking the object via a call to <b>dshmMuxSvcObjRelease()</b> .		
RETURNS	OK if receive is successful or ERROR otherwise.		

 ERRNOS
 S\_dshm\_MUX\_SERVICE\_NOT\_REGISTERED

 any error from the installed receive routine for the service.

SEE ALSO dshmMuxLib

### dshmMuxMsgSend()

dshmMuxMsgSend() - transmit a message NAME SYNOPSIS STATUS dshmMuxMsgSend ( const uint\_t hw, /\* hardware on which to transmit \*/ DSHM(msg), /\* message to transmit \*/ unused1, /\* possibe future expansion: use 0 \*/ int unused2 /\* possibe future expansion: use 0 \*/ int ) This routine will send a message to a specified destination on a specified hardware based DESCRIPTION on information provided in the *msg* argument. A broadcast address can be provided. OK if send is successful or ERROR otherwise. RETURNS S\_dshm\_MUX\_NODE\_NOT\_REGISTERED ERRNOS

#### dshmMuxSvcNodeJoin()

dshmMuxLib

SEE ALSO

NAME	dshmMuxSvcNodeJoin() – signal services that a node has joined the system
SYNOPSIS	<pre>void dshmMuxSvcNodeJoin   (     const uint_t hw, /* on which hw this node sits */     const uint16_t addr /* node's unique address on hw */ )</pre>
DESCRIPTION	This routine is to be called when a node joins the system to allow registered service to perform some action based on the event. This can be something like setting up colleague links, allocating memory for the peer, etc.

This routine should be called by hardware interfaces when they detect that a remote node has appeared on the shared bus.

RETURNS N/A

ERRNO Not Available

SEE ALSO dshmMuxLib

## dshmMuxSvcNodeLeave()

NAME	<b>dshmMuxSvcNodeLeave( )</b> – signal services that a node has left the system				
SYNOPSIS	<pre>void dshmMuxSvcNodeLeave   (     const uint_t hw, /* hardware handle obtained from registration */     const uint16_t addr /* node's unique address on the bus */   )</pre>				
DESCRIPTION	This routine is to be called when a node leaves the system to allow registered services to perform some action based on the event. This allows services to update their view of the system and reclaim resources used for interacting with that particular node.				
	This routine should be called by hardware interfaces when they detect that a remote node has disappeared on the shared bus.				
RETURNS	N/A				
ERRNOS	N/A				
SEE ALSO	dshmMuxLib				

## dshmMuxSvcObjGet()

NAME	dshmMuxSvcObjGet() – retrieve a service object and protect it against deletion
SYNOPSIS	<pre>void * dshmMuxSvcObjGet   (     const uint_t hw, /* hw on which the service exists */     const uint_t svc /* service number, well-known */   )</pre>

DESCRIPTION	This routine retrieves the reference to the object of a previously registered service <i>svc</i> on the hardware bus <i>hw</i> . It will also prevent the object from being deleted.
RETURNS	A pointer to the object or NULL if the service or hardware are not registered.
ERRNOS	S_dshm_MUX_SERVICE_NOT_REGISTERED
SEE ALSO	dshmMuxLib

# dshmMuxSvcObjRelease()

NAME	dshmMuxSvcObjRelease() – allows modifications to be made on a service object
SYNOPSIS	<pre>void dshmMuxSvcObjRelease   (     const uint_t hw,     const uint_t svc   )</pre>
DESCRIPTION	This routine releases a previously obtained reference to the object of a registered service $svc$ on the hardware bus $hw$ .
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	dshmMuxLib

# dshmMuxSvcRegister()

NAME dshmMuxSvcRegister() – register a service with the MUX

SYNOPSIS	STATUS dshmMuxSvcRegister			
	(			
	const uint_t	hw,	/*	hw on which to register service */
	const uint_t	svc,	/*	service number, well-known */
	const void * const	pObj,	/*	pointer to service object */
	const DSHM_SVC_HOOKS *	const pHooks	5	/* service receive routine */
	)			

VxWorks Kernel API Reference, 6.6 dshmMuxSvcWithdraw()

DESCRIPTION	Register a service with well-known number <i>svc</i> with the MUX. A receive hook called when a message arrives on any hardware registered must be provided via the <i>rx</i> argument. An argument to the <i>rx</i> method can register via <i>arg</i> and will be passed to <i>rx</i> when it is called.
RETURNS	OK if service if registered successfully or ERROR otherwise.
ERRNOS	S_dshm_MUX_SERVICE_TABLE_FULL
SEE ALSO	dshmMuxLib

## dshmMuxSvcWithdraw()

NAME	dshmMuxSvcWithdraw() – remove service from MUX
SYNOPSIS	STATUS dshmMuxSvcWithdraw ( const uint_t hw, const uint_t svc )
DESCRIPTION	This routine removes the service with well-known name <i>svc</i> from the MUX, for the hardware bus <i>hw</i> . The <b>stop</b> callback registered with the service will be invoked at this point. After removal, messages intended for it will be discarded by the hardware processing loop.
RETURNS	OK if node withdrawal successful or ERROR otherwise.
ERRNOS	S_dshm_MUX_SERVICE_NOT_REGISTERED
SEE ALSO	dshmMuxLib

# dshmMuxWidtdrawComplete()

NAME	dshmMuxWidtdrawComplete() – signal service has finished withdrawing
SYNOPSIS	STATUS dshmMuxSvcWithdrawComplete ( const uint_t hw, const uint_t svc )
DESCRIPTION	To be called by the service when the service determines it can remove itself from the system completely.
-------------	---
WARNING	This routine must be called with the reference to the object NOT OWNED by the calling thread. This routine cannot be called from within the <b>stop</b> callback registered with the service.
RETURNS	Not Available
ERRNO	Not Available
SEE ALSO	dshmMuxLib

#### dsiDataPoolShow()

NAME	dsiDataPoolShow() – display DSI's data pool statistics
SYNOPSIS	void dsiDataPoolShow (void)
DESCRIPTION	This routine displays the statistics for the allocated/available clusters in the DSI data pool. That pool is used to transfer data packets between DSI sockets.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	dsiSockLib, netPoolShow()

#### dsiSysPoolShow()

NAME dsiSysPoolShow() – display DSI's system pool statistics

SYNOPSIS void dsiSysPoolShow (void)

**DESCRIPTION** This routine displays the statistics for the allocated/available clusters in the DSI system pool. That pool is used by DSI sockets and their protocols' control blocks.

RETURNS N/A

VxWorks Kernel API Reference, 6.6 e( )

ERRNO Not Available

SEE ALSO dsiSockLib, netPoolShow()

## **e( )**

NAME	e() – set or display eventpoints (WindView)
SYNOPSIS	<pre>STATUS e ( INSTR * addr, /* where to set eventpoint, or */</pre>
DESCRIPTION	<ul> <li>This routine sets "eventpoints"that is, breakpoint-like instrumentation markers that can be inserted in code to generate and log an event for use with WindView. Event logging must be enabled with wvEvtLogEnable() for the eventpoint to be logged.</li> <li>eventId selects the evenpoint number that will be logged: it is in the user event ID range (0-25536).</li> <li>If addr is NULL, then all eventpoints and breakpoints are displayed. If taskNameOrId is 0, then this event is logged in all tasks. The evtRtn routine is called when this eventpoint is hit. If evtRtn returns OK, then the eventpoint is logged; otherwise, it is ignored. If evtRtn is a NULL pointer, then the eventpoint (which are set with the b() command) except in how the system responds when the eventpoint is hit. An eventpoint typically records an event and continues immediately (if evtRtn is supplied, this behavior may be different). Eventpoints cannot be used at interrupt level.</li> <li>To delete an eventpoint, use bd().</li> </ul>
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>addr</i> is odd or nonexistent in memory, or if the breakpoint table is full.
ERRNO	N/A
SEE ALSO	dbgLib, wvEvent(), VxWorks Kernel Programmer's Guide: Kernel Shell

# edi()

NAME	edi() – return the contents of register edi (also esi - eax) (x86)
SYNOPSIS	<pre>int edi   (    int taskId /* task ID, 0 means default task */ )</pre>
DESCRIPTION	This command extracts the contents of register <b>edi</b> from the TCB of a specified task. If <i>taskId</i> is omitted or zero, the last task referenced is assumed.
	Similar routines are provided for all general registers (edi - eax): edi() - eax().
	The stack pointer is accessed via <b>eax()</b> .
RETURNS	The contents of register <b>edi</b> (or the requested register).
ERRNO	Not Available
SEE ALSO	dbgArchLib, VxWorks Programmer's Guide: Debugging

# edi()

NAME	edi() – return the contents of register edi (also esi - eax) (x86/SimNT)
SYNOPSIS	<pre>int edi   (    int taskId /* task ID, 0 means default task */ )</pre>
DESCRIPTION	This command extracts the contents of register <b>edi</b> from the TCB of a specified task. If <i>taskId</i> is omitted or zero, the last task referenced is assumed.
	Similar routines are provided for all address registers (edi - eax): edi() - eax().
	The stack pointer is accessed via <b>eax()</b> .
RETURNS	The contents of register <b>edi</b> (or the requested register).
ERRNO	Not Available
SEE ALSO	dbgArchLib, VxWorks Programmer's Guide: Debugging

#### edrBootCountGet()

NAME	<pre>edrBootCountGet() - returns the current boot count</pre>
SYNOPSIS	int edrBootCountGet (void)
DESCRIPTION	This function returns the number of times the system has been rebooted since ED&R was first initialized.
RETURNS	The current boot count, or ERROR if it cannot be determined.
ERRNO	Not Available
SEE ALSO	edrLib

### edrBootShow()

NAME	edrBootShow() – show all stored boot type ED&R records
SYNOPSIS	<pre>STATUS edrBootShow   (    int start, /* starting point */    int count /* number of records to show */   )</pre>
DESCRIPTION	This command displays all records stored in the ED&R log which have a facility of BOOT. The command accepts a <i>start</i> and <i>count</i> as documented in <b>edrShow()</b> .
RETURNS	OK or ERROR
ERRNO	Not Available
SEE ALSO	edrShow, edrShow()

## edrClear()

NAME	edrClear() – a synonym for edrErrorLogClear
SYNOPSIS	STATUS edrClear (void)
DESCRIPTION	This functions provides command line interface to clear the entire ED&R error log.
RETURNS	OK, or ERROR if there was some problem clearing the log
ERRNO	Not Available
SEE ALSO	edrShow

## edrErrLogAttach()

NAME	edrErrLogAttach() – attach to an existing log
SYNOPSIS	EDR_ERR_LOG * edrErrLogAttach ( void * pAddr /* address of an existing log */ )
DESCRIPTION	This routine attaches to a previously created error log, but only if <i>pAddr</i> represents a valid existing log.
RETURNS	pointer to log, or NULL if the pLog doesn't represent a valid log instance.
ERRNO	Not Available
SEE ALSO	edrErrLogLib, edrErrLogCreate()

## edrErrLogClear()

NAME edrErrLogClear() – clear the log's contents

SYNOPSIS BOOL edrErrLogClear

VxWorks Kernel API Reference, 6.6 edrErrLogCreate()

```
(
EDR_ERR_LOG * pLog /* pointer to log */
)
```

**DESCRIPTION** This routine resets all of the the log's node nodes to the unallocated state.

**RETURNS TRUE**, or **FALSE** if the *pLog* doesn't represent a valid log instance.

ERRNO Not Available

SEE ALSO edrErrLogLib, edrErrLogCreate(), edrErrLogAttach()

### edrErrLogCreate()

NAME	edrErrLogCreate() – create a new log
SYNOPSIS	EDR_ERR_LOG * edrErrLogCreate ( void * pAddr, /* address to overlay the log at */ int size, /* length (in bytes) of the log area */ int recordSize /* length (in bytes) of each record */ )
DESCRIPTION	Creates a new log at address <i>pAddr</i> for a length of <i>size</i> bytes. Each record in the log is created as size <i>recordSize</i> . The size of the record may be no less than <b>EDR_ERR_LOG_MIN_PAYLOAD_SIZE</b> . The log has a fixed size overhead of approximately 500 bytes plus an overhead of approximately 150 bytes per record. If the architecture supports an MMU, the log is write protected after being created.
NOTE	The memory provided for the error log must be page aligned and be a multiple of a page size.
RETURNS	A pointer to the log, or <b>NULL</b> if the log could not be created, or <i>pAddr</i> is <b>NULL</b> .
ERRNO	Not Available
SEE ALSO	edrErrLogLib, edrErrLogAttach()

NAME	edrErrLogIterCreate() – create an iterator for traversing the log
SYNOPSIS	<pre>BOOL edrErrLogIterCreate   (    EDR_ERR_LOG * pLog, /* pointer to log */    EDR_ERR_LOG_ITER * pIter, /* pointer to iter for construction */    int start, /* starting position */    int count /* number of nodes to enumerate */   )</pre>
DESCRIPTION	This routine creates an iterator suitable for traversing the set of committed nodes in the log The starting postion (record number) to begin traversing the log is specified by <i>start</i> . The maximum number of records to iterate over is specified by <i>count</i> .
RETURNS	TRUE, or FALSE if <i>pLog</i> points to a corrupt or invalid log.
ERRNO	Not Available
SEE ALSO	edrErrLogLib, edrErrLogIterNext()

# edrErrLogIterNext()

NAME	edrErrLogIterNext() – returns the next committed node
SYNOPSIS	EDR_ERR_LOG_NODE * edrErrLogIterNext ( EDR_ERR_LOG_ITER * pIter /* pointer to iterator */ )
DESCRIPTION	This routine returns the next committed node in the log, or <b>NULL</b> if there are no more nodes that can be enumerated.
RETURNS	A pointer to next node or NULL.
ERRNO	Not Available
SEE ALSO	edrErrLogLib, edrErrLogIterCreate()

## edrErrLogMaxNodeCount()

NAME	edrErrLogMaxNodeCount() – return the maximum number of nodes in the log		
SYNOPSIS	<pre>int edrErrLogMaxNodeCount   (    EDR_ERR_LOG * pLog /* pointer to log */ )</pre>		
DESCRIPTION	This routine returns the maximum number of nodes that can be held within the specified log.		
RETURNS	node count, or <b>ERROR</b> if <i>pLog</i> is not a valid error log.		
ERRNO	Not Available		
SEE ALSO	edrErrLogLib		

## edrErrLogNodeAlloc()

NAME	edrErrLogNodeAlloc() – allocate a node from the error log			
SYNOPSIS	EDR_ERR_LOG_NODE * edrErrLogNodeAlloc ( EDR_ERR_LOG * pLog /* pointer to log */ )			
DESCRIPTION	This routine allocates the next available node from the specified log. Once the node information has been written, the node must be committed to the log using <b>edrErrLogNodeCommit()</b> .			
	If the architecture supports an MMU, once an node is allocated it is unprotected (ie. writable) until it is commit using <b>edrErrLogNodeCommit()</b> .			
RETURNS	A node instance, or NULL if there are no free nodes.			
ERRNO	Not Available			
SEE ALSO	edrErrLogLib, edrErrLogNodeCommit()			

# edrErrLogNodeCommit()

NAME	<pre>edrErrLogNodeCommit( ) - commits a previously allocated node</pre>				
SYNOPSIS	BOOL edrErrLogNodeCommit ( EDR_ERR_LOG * pLog, /* pointer to log */ EDR_ERR_LOG_NODE * pNode /* pointer to node to commit */ )				
DESCRIPTION	This routine commits the previously allocated node <i>pNode</i> to the log specified by <i>pLog</i> . Once a node becomes committed, it may be re-allocated by <b>edrErrLogNodeAlloc()</b> if necessary.				
	If the architecture supports an MMU, the node is write protected after being committed.				
RETURNS	<b>TRUE</b> if <i>pNode</i> is currently allocated, or <b>FALSE</b> if <i>pLog</i> or <i>pNode</i> don't represent valid instances of a log or a node respectively.				
ERRNO	Not Available				
SEE ALSO	edrErrLogLib, edrErrLogNodeAlloc()				

# edrErrLogNodeCount()

NAME	edrErrLogNodeCount() – return the number of committed nodes in the log				
SYNOPSIS	<pre>int edrErrLogNodeCount   (    EDR_ERR_LOG * pLog /* pointer to log */ )</pre>				
DESCRIPTION	This routine returns the number of committed nodes within the specified log.				
RETURNS	THe node count, or <b>ERROR</b> if <i>pLog</i> is not a valid error log.				
ERRNO	Not Available				
SEE ALSO	edrErrLogLib				

	edrErrorInject()				
NAME	edrErrorInject() – injects an error into the ED&R subsystem				
SYNOPSIS	<pre>STATUS edrErrorInject (     int kind, /* severity   facility   option */     const char * fileName, /* name of source file */     int lineNumber, /* line number of source code */     const REG_SET * pRegSet, /* current register values */     const EXC_INFO * pExcInfo, /* CPU-specific exception info */     void * address, /* faulting address */     const char * msg /* additional text string */ )</pre>				
DESCRIPTION	Warning: This function should not normally be called directly, rather, one of the macros in <b>edrLib.h</b> such as <b>EDR_KERNEL_FATAL_INJECT</b> should be used instead.				
	eurlisht such as EDR_KERNEL_FATAL_INJECT should be used instead.         This function takes all the supplied arguments and stores them in an error record, along with numerous other bits of useful information, such as:         - the OS version         - the CPU type and number         - the time at which the error occured         - the current OS context (task, interrupt, exception, RTP)         - a small memory map of the running system         - a code fragment from around the faulting instruction         - a stack trace of the currently active stack         The type of record being injected is represented by the kind parameter. The kind parameter is a bitwise OR of the following three items:         Severity:         EDR_SEVERITY_FATAL       - a fatal event         EDR_SEVERITY_NONFATAL       - a fotal event         EDR_SEVERITY_WARNING       - a warning event				
	Facility:EDR_FACILITY_KERNEL- VxWorks kernel eventsEDR_FACILITY_INTERRUPT- interrupt handler eventsEDR_FACILITY_INIT- system startup eventsEDR_FACILITY_BOOT- system boot eventsEDR_FACILITY_REBOOT- system restart eventsEDR_FACILITY_REBOOT- system restart eventsEDR_FACILITY_RTP- RTP system eventsEDR_FACILITY_USER- user generated events				

Options:

	EDR_EXCLUDE_REGISTERS- don't include registersEDR_EXCLUDE_TRACEBACK- don't include stack traceEDR_EXCLUDE_EXCINFO- don't include exc infoEDR_EXCLUDE_DISASSEMBLY- don't include code dissemblyEDR_EXCLUDE_MEMORYMAP- don't include memory map				
	From an injection point of view, only the options have an effect on how the record is generated. The severity and facility values are merely stored in the record for subsequent use by the show commands.				
	If the ED&R subsystem is not yet initialised, then the error-record cannot be written to the log.				
LIMITATIONS	Since this function may well be called in an exception handling context, it must be able to run in the limited stack environment for exception handlers. The stack may be as small as a single VM page, i.e. 4K. Thus no large stack-based arrays or other structures should be used by the injection hooks.				
SMP CONSIDERATIONS					
	This API is spinlock and intCpuLock restricted.				
RETURNS	OK if the error was stored correctly, or ERROR if some failure occurs during storage				
ERRORS	S_edrLib_NOT_INITIALIZED if the library was not initialized				
	S_edrLib_PROTECTION_FAILURE if the memory could not be protected				
SEE ALSO	edrLib				

# edrErrorInjectHookAdd()

NAME	edrErrorInjectHookAdd() – adds a hook which gets called on error-injection		
SYNOPSIS	STATUS edrErrorInjectHookAdd ( EDR_ERRINJ_HOOK_FUNCPTR injectHook )		
DESCRIPTION	This function adds a hook to <b>edrLib</b> that gets called whenever an error is injected. The hook function (which must be of type <b>EDR_ERRINJ_HOOK_FUNCPTR</b> ) is invoked with a subset of		

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	the parameters passed to <b>edrErrorInject()</b> . The parameters passed to the hook function are as follows:			
	int const char * int void * const char *	kind fileName lineNumber address msg	<pre>/* severity   facility /* name of source file /* line number of source code /* faulting address /* additional text string</pre>	* / * / * /
IMPORTANT NOTE	The hook function is called directly from <b>edrErrorInject()</b> and so may be invoked in an interrupt or exception context. Hook functions should therefore make no blocking calls to the VxWorks kernel API, not should they use an excessive amount of stack space.			
SMP CONSIDERATIONS This API is spinlock and intCpuLock restricted.				
RETURNS	<b>OK</b> or <b>ERROR</b> if the hook table is full			
ERRORS	S_edrLib_INJECT_HOOK_TABLE_FULL if the hook table is full.			

SEE ALSO edrLib, edrErrorInjectHookDelete

## edrErrorInjectHookDelete()

NAME	edrErrorInjectHookDelete() – removes an existing error-inject hook				
SYNOPSIS	STATUS edrErrorInjectHookDelete ( EDR_ERRINJ_HOOK_FUNCPTR injectHook )				
DESCRIPTION	This function removes a hook which was added using <b>edrErrorInjectHookAdd()</b> .				
SMP CONSIDERATIO	NS This API is spinlock and intCpuLock restricted.				
RETURNS	<b>OK</b> or <b>ERROR</b> if the hook was not found in the hook table.				
ERRORS	S_edrLib_INJECT_HOOK_NOT_FOUND if the hook was not found.				
SEE ALSO	edrLib, edrErrorInjectHookAdd				

# edrErrorInjectPrePostHookAdd()

NAME	<pre>edrErrorInjectPrePostHookAdd() - adds a hook which gets called before and after error-injection STATUS edrErrorInjectPrePostHookAdd</pre>				
Synopsis					
DESCRIPTION	This function adds a hook to <b>edrLib</b> that gets called before and after an error is injected. The hook function (which must be of type FUNCPTR) is invoked with a single integer argument indicating whether it is pre-injection (EDR_HOOK_TYPE_PRE), or post-injection (EDR_HOOK_TYPE_POST).				
	These hook points are not generally for use by clients of <b>edrLib</b> , but are provided as a means to, for example, kick a hardware watchdog before entering <b>edrErrorInject()</b> in order to ensure that <b>edrLib</b> has sufficient time to inject an error-record into the error-log before the watchdog reboots the board.				
	It should have the following form:-				
	<pre>void prePostHook   (     int prePost /* EDR_HOOK_TYPE_PRE, EDR_HOOK_TYPE_POST */   );</pre>				
IMPORTANT NOTE	The hook function is called directly from <b>edrErrorInject()</b> and so may be invoked in an interrupt or exception context. Hook functions should therefore make no blocking calls to the VxWorks kernel API, not should they use an excessive amount of stack space.				
SMP CONSIDERATIO	DNS				
	This API is spinlock and intCpuLock restricted.				
RETURNS	<b>OK</b> or <b>ERROR</b> if the hook table is full				
ERRORS	<b>S_edrLib_PP_HOOK_TABLE_FULL</b> if the hook table is full.				
SEE ALSO	edrLib, edrErrorInjectPrePostHookDelete				

### edrErrorInjectPrePostHookDelete()

NAME	edrErrorInjectPrePostHookDelete() – removes the existing pre/post hook				
SYNOPSIS	STATUS edrErrorInjectPrePostHookDelete ( VOIDFUNCPTR hook )				
DESCRIPTION	This function removes a hook which was added using <b>edrErrorInjectPrePostHookAdd()</b> .				
SMP CONSIDERATIO	<b>EXAMPLE</b> API is spinlock and intCpuLock restricted.				
RETURNS	OK or ERROR if the hook was not found in the hook table.				
ERRORS	S_edrLib_PP_HOOK_NOT_FOUND if the hook was not found.				
SEE ALSO	edrLib, edrErrorInjectPrePostHookAdd				

#### edrErrorInjectTextHookAdd()

NAME	edrErrorInjectTextHookAdd() – adds a hook which gets called on record creation				
SYNOPSIS	STATUS edrErrorInjectTextHookAdd ( EDR_ERRINJ_TEXT_HOOK_FUNCPTR textHook )				
DESCRIPTION	This function adds a hook to <b>edrLib</b> that gets called whenever an error is created. The function (which must be of type <b>EDR_ERRINJ_TEXT_HOOK_FUNCPTR</b> ) is invoked w pointer and length to a buffer which can be filled with textual information. The parar passed to the hook function are as follows:				
	char * int const char * int void *	p size kind fileName lineNumber address	<pre>/* pointer to buffer /* size of buffer /* severity   facility /* name of source file /* line number of source cod /* faulting address</pre>	*/ */ */ le */ */	

The hook must ensure the string written is null terminated. The return value of the hook must be the number of bytes stored, including the trailing null.

IMPORTANT NOTE	The hook function is called directly from <b>edrErrorInject( )</b> and so may be invoked in an
	interrupt or exception context. Hook functions should therefore make no blocking calls to
	the VxWorks kernel API, not should they use an excessive amount of stack space.

#### SMP CONSIDERATIONS

RETURNS	<b>OK</b> or <b>ERROR</b> if the hook table is full
ERRORS	S_edrLib_TEXT_HOOK_TABLE_FULL if the hook table is full
SEE ALSO	edrLib, edrErrorInjectTextHookDelete

### edrErrorInjectTextHookDelete()

NAME	<pre>edrErrorInjectTextHookDelete( ) - removes the existing text writing hook</pre>
SYNOPSIS	STATUS edrErrorInjectTextHookDelete ( EDR_ERRINJ_TEXT_HOOK_FUNCPTR textHook )
DESCRIPTION	This function removes a hook which was added using <b>edrErrorInjectTextHookAdd()</b> .
SMP CONSIDERATIO	<b>NS</b> This API is spinlock and intCpuLock restricted.
RETURNS	<b>OK</b> or <b>ERROR</b> if the hook was not found in the hook table.
ERRORS	S_edrLib_TEXT_HOOK_NOT_FOUND if the hook was not found.
SEE ALSO	edrLib, edrErrorInjectTextHookAdd

### edrErrorLogClear()

NAME edrErrorLogClear() – clears the ED&R error log

SYNOPSIS STATUS edrErrorLogClear (void)

VxWorks Kernel API Reference, 6.6 edrErrorRecordCount()

DESCRIPTION	This function clears all error records out of the error log. It is destructive (records cannot be undeleted) and should be used with utmost care.
RETURNS	<b>OK</b> or <b>ERROR</b> if the log was not able to be cleared
ERRORS	S_edrLib_NOT_INITIALIZED if the library was not initialized
	S_edrLib_PROTECTION_FAILURE if the memory could not be protected
SEE ALSO	edrLib

### edrErrorRecordCount()

NAME	edrErrorRecordCount() – returns the number of error-records in the log
SYNOPSIS	int edrErrorRecordCount (void)
DESCRIPTION	This function returns the total number of ED&R records which are present in the error log.
RETURNS	the number of error-records in the log, or ERROR if the library has not been initialized
ERRORS	S_edrLib_NOT_INITIALIZED if the library was not initialized
SEE ALSO	edrLib

### edrErrorRecordDecode()

NAME	edrErrorRecordDecode(	) – decode	one error-record	
SYNOPSIS	STATUS edrErrorRecord ( EDR_ERROR_RECORD* char * int )	Decode pER, pBuf, bufSize	/* pointer to error record /* pointer to output buffer /* size of output buffer	*/ */ */
DESCRIPTION	This routine decodes a sin the record is decoded to a	ngle error-1 stdout.	record into the provided buffer. If	f no buffer is provided,

**RETURNS** OK, or **ERROR** if the record can't be decoded

ERRNO Not Available

SEE ALSO edrShow

#### edrFatalShow()

NAME	edrFatalShow() – show all stored fatal type ED&R records
SYNOPSIS	<pre>STATUS edrFatalShow   (    int start, /* starting point */    int count /* number of records to show */   )</pre>
DESCRIPTION	This command displays all records stored in the ED&R log which have a severity of FATAL. The command accepts a <i>start</i> and <i>count</i> as documented in <b>edrShow()</b> .
RETURNS	OK or ERROR
ERRNO	Not Available
SEE ALSO	edrShow, edrShow()

## edrFlagsGet()

NAME	edrFlagsGet() – return the ED&R flags which are currently set
SYNOPSIS	int edrFlagsGet(void)
DESCRIPTION	This routine returns all the ED&R flags which have been set to "on". An identical API is provided in the RTP space.
RETURNS	an integer with the appropriate bits set
ERRNO	Not Available
SEE ALSO	edrSysDbgLib

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## edrHelp()

SEE ALSO	edrShow
ERRNO	Not Available
RETURNS	ОК
DESCRIPTION	This routine provides the on-line help for the ED&R show commands
SYNOPSIS	STATUS edrHelp (void)
NAME	edrHelp() – prints helpful information on ED&R

## edrHookShow()

edrHookShow() – show the list of installed ED&R hook routines
STATUS edrHookShow (void)
This routine shows all the hook routines installed in the various ED&R hook tables, in the order in which they were installed.
N/A
Not Available
edrShow, edrInjectHookShow( ), edrInjectTextHookShow( ), edrInjectPrePostHookShow( )

## edrInfoShow()

NAME	edrInfoShow() – show all stored info type ED&R records
SYNOPSIS	STATUS edrInfoShow
	int start, /* starting point */

 int count
 /\* number of records to show \*/

 DESCRIPTION
 This command displays all records stored in the ED&R log which have a severity of INFO. The command accepts a *start* and *count* as documented in edrShow().

 RETURNS
 OK or ERROR

 ERRNO
 Not Available

 SEE ALSO
 edrShow, edrShow()

#### edrInitShow()

NAME	edrInitShow() – show all stored init type ED&R records
SYNOPSIS	<pre>STATUS edrInitShow   (    int start, /* starting point */    int count /* number of records to show */   )</pre>
DESCRIPTION	This command displays all records stored in the ED&R log which have a facility of INIT. The command accepts a <i>start</i> and <i>count</i> as documented in <b>edrShow()</b> .
RETURNS	OK or ERROR
ERRNO	Not Available
SEE ALSO	edrShow, edrShow()

#### edrInjectHookShow()

NAME	edrInjectHookShow() – show the list of error injection hook routines
SYNOPSIS	STATUS edrInjectHookShow (void)
DESCRIPTION	This routine shows all the error injection routines installed in the ED&R inject hook table, in the order in which they were installed.

VxWorks Kernel API Reference, 6.6 edrInjectPrePostHookShow()

RETURNS N/A

ERRNO Not Available

SEE ALSO edrShow, edrInjectHookAdd()

## edrInjectPrePostHookShow()

NAME	edrInjectPrePostHookShow() – show the list of pre/post injection hook routines
SYNOPSIS	STATUS edrInjectPrePostHookShow (void)
DESCRIPTION	This routine shows all the pre/post error injection routines installed in the ED&R inject hook table, in the order in which they were installed.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	edrShow, edrInjectPrePostHookAdd( )

# edrInjectTextHookShow()

NAME	edrInjectTextHookShow() – show the list of text injection hook routines
SYNOPSIS	STATUS edrInjectTextHookShow (void)
DESCRIPTION	This routine shows all the text error injection routines installed in the ED&R inject hook table, in the order in which they were installed.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	edrShow, edrInjectTextHookAdd( )

#### edrIntShow()

NAME	edrIntShow() – show all stored interrupt type ED&R records
SYNOPSIS	<pre>STATUS edrIntShow   (    int start, /* starting point */    int count /* number of records to show */   )</pre>
DESCRIPTION	This command displays all records stored in the ED&R log which have a facility of INTERRUPT. The command accepts a <i>start</i> and <i>count</i> as documented in <b>edrShow()</b> .
RETURNS	OK or ERROR
ERRNO	Not Available
SEE ALSO	edrShow, edrShow( )

#### edrIsDebugMode()

NAME	edrIsDebugMode()	) – is the ED&R	debug mod	e flag set?

- SYNOPSIS BOOL edrIsDebugMode(void)
- **DESCRIPTION** This routine returns **TRUE** if the ED&R debug flag is set.
- **RETURNS** TRUE if debug flag is set, FALSE otherwise.

ERRNO Not Available

SEE ALSO edrSysDbgLib

#### edrKernelShow()

NAME edrKernelShow() – show all stored kernel type ED&R records

SYNOPSIS STATUS edrKernelShow

VxWorks Kernel API Reference, 6.6 edrLibInit()

```
(
int start, /* starting point */
int count /* number of records to show */
)
```

**DESCRIPTION** This command displays all records stored in the ED&R log which have a facility of KERNEL. The command accepts a *start* and *count* as documented in **edrShow()**.

RETURNS OK or ERROR

ERRNO Not Available

SEE ALSO edrShow, edrShow()

## edrLibInit()

NAME	edrLibInit() – initializes edrLib	
SYNOPSIS	<pre>STATUS edrLibInit   (   BOOL isNew, /* should the PM log area be re-initialized? */   int recordSize /* size of each ED&amp;R record */   )</pre>	
DESCRIPTION	This function initializes the ED&R susbsystem. The parameter <i>isNew</i> indicates whether the persistent memory (PM) region used to hold the ED&R records is to be re-initialized. If <i>isNew</i> is <b>FALSE</b> , <b>edrLibInit()</b> will attempt to open an existing ED&R log PM region and check that it is valid. If the PM region is invalid, an error is returned. Once the PM region is opened successfully, the PM region is write protected. The parameter <i>recordSize</i> is used whenever the log is re-initilaised. Its value represents the payload size for each of the records in the log. A value of zero specifies that the built in system default size is to be used.	
RETURNS	OK if the log was successfully created or opened, ERROR otherwise.	
ERRORS	<ul> <li>S_edrLib_PMREGION_ERROR if the PM region is invalid</li> <li>S_edrLib_PROTECTION_FAILURE if the PM region cannot be protected</li> <li>S_edrLib_ERRLOG_CORRUPTED if the log appears to be corrupted</li> <li>S_edrLib_ERRLOG_INCOMPATIBLE if the log is a newer version</li> </ul>	

#### edrRebootShow()

NAME	edrRebootShow() – show all stored reboot type ED&R records
SYNOPSIS	<pre>STATUS edrRebootShow   (    int start, /* starting point */    int count /* number of records to show */   )</pre>
DESCRIPTION	This command displays all records stored in the ED&R log which have a facility of REBOOT. The command accepts a <i>start</i> and <i>count</i> as documented in <b>edrShow()</b> .
RETURNS	OK or ERROR
ERRNO	Not Available
SEE ALSO	edrShow, edrShow( )

## edrRtpShow()

NAME	edrRtpShow() – show all stored rtp type ED&R records
SYNOPSIS	<pre>STATUS edrRtpShow   (    int start, /* starting point */    int count /* number of records to show */   )</pre>
DESCRIPTION	This command displays all records stored in the ED&R log which have a facility of RTP. The command accepts a <i>start</i> and <i>count</i> as documented in <b>edrShow()</b> .
RETURNS	OK or ERROR
ERRNO	Not Available
SEE ALSO	edrShow, edrShow( )

### edrShow()

NAME	edrShow() – displays the ED&R error log to stdout		
SYNOPSIS	<pre>STATUS edrShow   (    int start, /* starting point */    int count, /* number of records to show */    int facility, /* limit to specified facility */    int severity /* limit to specified severity */   )</pre>		
DESCRIPTION	This command display all or part of the stored ED&R error log to stdout. The command takes four parameters, the start, count, facility, and severity specifiers.		
	The <i>start</i> parameter specifies the starting record at which the display should begin. If <i>start</i> is a positive number, the display begins at <i>start</i> records from the beginning of the log. If <i>start</i> is a negative number, the display begins at <i>start</i> records from the end of the log.		
	The <i>count</i> parameter denotes the number of records to display. A value of zero will display all records.		
	The <i>facility</i> and <i>severity</i> parameters limit the display to only those records which match the specified facility and severity. A value of zero will match all facilities and severities.		
RETURNS	OK or ERROR if the ED&R library was not initialized		
ERRNO	S_edrLib_NOT_INITIALIZED		
SEE ALSO	edrShow		

# edrSystemDebugModeGet()

NAME	edrSystemDebugModeGet() – indicates if the system is in debug mode
SYNOPSIS	BOOL edrSystemDebugModeGet (void)
DESCRIPTION	This routine returns the current setting of the system mode debug flag if it has been set otherwise it assumes it is off.
RETURNS	TRUE if the <b>debug mode</b> boot flag is set, FALSE if not
ERRNO	Not Available

SEE ALSO edrSysDbgLib

### edrSystemDebugModeInit()

NAME	<pre>edrSystemDebugModeInit() – initialise the system mode debug flag</pre>
SYNOPSIS	STATUS edrSystemDebugModeInit (void)
DESCRIPTION	This routine reads the <b>SYSFLG_SYS_MODE_DEBUG</b> flag from the boot-flags supplied in the system boot line and set the state of the debug flag.
RETURNS	TRUE if the <b>debug mode</b> boot flag is set, FALSE if not
ERRNO	Not Available
SEE ALSO	edrSysDbgLib

# edrSystemDebugModeSet()

NAME	edrSystemDebugModeSet() – modifies the system debug mode flag
SYNOPSIS	void edrSystemDebugModeSet ( BOOL mode )
DESCRIPTION	This routine sets the system debug flag, as maintained by ED&R. It over-rides any setting given in the boot flags.
RETURNS	n/a
ERRNO	Not Available
SEE ALSO	edrSysDbgLib

## edrUserShow()

NAME	edrUserShow() – show all stored user type ED&R records	
SYNOPSIS	<pre>STATUS edrUserShow   (    int start, /* starting point */    int count /* number of records to show */   )</pre>	
DESCRIPTION	This command displays all records stored in the ED&R log which have a facility of USER. The command accepts a <i>start</i> and <i>count</i> as documented in <b>edrShow()</b> .	
RETURNS	OK or ERROR	
ERRNO	Not Available	
SEE ALSO	edrShow, edrShow( )	

# eflags()

NAME	eflags() – return the contents of the status register (x86)	
SYNOPSIS	<pre>int eflags   (    int taskId /* task ID, 0 means default task */ )</pre>	
DESCRIPTION	This command extracts the contents of the status register from the TCB of a specified task. If <i>taskId</i> is omitted or zero, the last task referenced is assumed.	
RETURNS	The contents of the status register.	
ERRNO	Not Available	
SEE ALSO	dbgArchLib, VxWorks Programmer's Guide: Debugging	

## eflags()

NAME	eflags() – return the contents of the status register (x86/SimNT)	
SYNOPSIS	<pre>int eflags   (    int taskId /* task ID, 0 means default task */ )</pre>	
DESCRIPTION	This command extracts the contents of the status register from the TCB of a specified task. If <i>taskId</i> is omitted or zero, the last task referenced is assumed.	
RETURNS	The contents of the status register.	
ERRNO	Not Available	
SEE ALSO	dbgArchLib, VxWorks Programmer's Guide: Debugging	

## elPciRegister()

NAME	elPciRegister() – register with the VxBus subsystem
SYNOPSIS	void elPciRegister(void)
DESCRIPTION	This routine registers the elPci driver with VxBus as a child of the PCI bus type.
RETURNS	N/A
ERRNO	N/A

SEE ALSO tc3c905VxbEnd

## eneRegister()

**NAME** eneRegister() – register with the VxBus subsystem

SYNOPSIS void eneRegister(void)

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VxWorks Kernel API Reference, 6.6 envGet()

**DESCRIPTION** This routine registers the NE2000 driver with VxBus as a child of the PLB bus type.

RETURNS N/A

ERRNO N/A

SEE ALSO ne2000VxbEnd

#### envGet()

NAME	<b>envGet( )</b> – return a pointer to the environment of a task
SYNOPSIS	<pre>char ** envGet   (    int taskId /* task for which the environment is to be returned */ )</pre>
DESCRIPTION	This routine returns a pointer to the environment of a task. If the task has a private environment set, this is what is being referred to then. If <i>taskId</i> is <b>NULL</b> , then the calling task's environment is retrieved.
RETURNS	a pointer to the task's environment, or the pointer to the global environment if the task has none.
ERRNOS	S_objLib_OBJ_ID_ERROR
SEE ALSO	envLib

#### envLibInit()

NAME	<b>envLibInit( )</b> – initialize environment variable facility
SYNOPSIS	STATUS envLibInit ( BOOL installHooks )
DESCRIPTION	If <i>installHooks</i> is <b>TRUE</b> , task create and delete hooks are installed that will optionally create and destroy private environments for the task being created or destroyed, depending on the state of <b>VX_PRIVATE_ENV</b> in the task options word. If <i>installHooks</i> is <b>FALSE</b> and a task

	requires a private environment, it is the application's responsibility to create and destroy the private environment, using <b>envPrivateCreate()</b> and <b>envPrivateDestroy()</b> .
	<i>installHooks</i> is controlled by the configuration parameter <b>ENV_VAR_USE_HOOKS</b> .
RETURNS	OK, or ERROR if an environment cannot be allocated or the hooks cannot be installed.
ERRNOS	N/A
SEE ALSO	envLib

### envPrivateCreate()

NAME	envPrivateCreate() – create a private environment	
SYNOPSIS	STATUS envPrivat ( int taskId, int envSourc )	<pre>seCreate    /* task to have private environment */ se /* -1 = make an empty private environment */    /* 0 = copy global env to new private env */    /* task id = copy the specified task's env */</pre>
DESCRIPTION	<b>RIPTION</b> This routine creates a private set of environment variables for a specified task, environment variable task create hook is not installed.	
	Based on the <i>envSource</i> argument, the environment is created in one of three ways:	
	envSource	Copy Behavior
	 -1 0 a task id	create an empty environment for <i>taskId</i> copy global environment to <i>taskId</i> 's new private environment Given a task ID, copy the task's private environment for <i>taskId</i>
NOTE	This API does not protect against the tasks from deletion while the environment information is being copied from one task to another. The user should take care not to delete the <i>taskId</i> nor the <i>envSource</i> task id while this routine is being used.	
RETURNS	OK, or ERROR if memory is insufficient or source environment is NULL.	
ERRNOS	S_objLib_OBJ_ID_ERROR taskId is invalid.	
SEE ALSO	envLib, envLibIni	it( ), envPrivateDestroy( )

## envPrivateDestroy()

NAME	envPrivateDestroy() – destroy a private environment
SYNOPSIS	<pre>STATUS envPrivateDestroy   (    int taskId /* task with private env to destroy */ )</pre>
DESCRIPTION	This routine destroys a private set of environment variables that were created with <b>envPrivateCreate()</b> . Calling this routine is unnecessary if the environment variable task create hook is installed and the task was spawned with <b>VX_PRIVATE_ENV</b> .
RETURNS	OK, or ERROR if the task does not exist.
ERRNOS	S_objLib_OBJ_ID_ERROR taskId is invalid.
SEE ALSO	envLib, envPrivateCreate()

### envShow()

NAME	envShow() – display the environment for a task	
SYNOPSIS	<pre>void envShow   (    int taskId /* task for which environment is printed */ )</pre>	
DESCRIPTION	This routine prints to standard output all the environment variables for a specified task or the global environment. If <i>taskld</i> is <b>NULL</b> , then the calling task's environment is displayed.	
RETURNS	N/A	
ERRNOS	S_objLib_OBJ_ID_ERROR	
SEE ALSO	envLib	

## errnoGet()

NAME	<b>errnoGet( )</b> – get the error status value of the calling task
SYNOPSIS	int errnoGet (void)
DESCRIPTION	This routine gets the current error status value. It is provided for compatibility with previous versions of VxWorks.
RETURNS	The current error status value.
ERRNO	N/A
SEE ALSO	errnoLib, errnoSet( ), errnoOfTaskGet( )

### errnoOfTaskGet()

NAME	errnoOfTaskGet() – get the error status value of a specified task
SYNOPSIS	<pre>int errnoOfTaskGet   (    int taskId /* task ID, 0 means current task */ )</pre>
DESCRIPTION	This routine gets the error status most recently set in the TCB of a specified task. If <i>taskId</i> is zero, the calling task is assumed.
	This routine is provided primarily for debugging purposes. Normally, tasks access <b>errno</b> directly to set and get their own error status values.
RETURNS	The error status of the specified task, or <b>ERROR</b> if the task does not exist.
ERRNO	N/A
SEE ALSO	errnoLib, errnoSet( ), errnoGet( )

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#### errnoOfTaskSet()

NAME	errnoOfTaskSet() – set the error status value of a specified task
SYNOPSIS	<pre>STATUS errnoOfTaskSet   (    int taskId, /* task ID, 0 means current task */    int errorValue /* error status value */   )</pre>
DESCRIPTION	This routine sets the error status value in the TCB for a specified task. If <i>taskId</i> is zero, the calling task is assumed.
	This routine is provided primarily for debugging purposes. Normally, tasks access <b>errno</b> directly to set and get their own error status values.
RETURNS	OK, or ERROR if the task does not exist.
ERRNO	N/A
SEE ALSO	errnoLib, errnoSet(), errnoOfTaskGet()

# errnoSet()

NAME	<b>errnoSet( )</b> – set the error status value of the calling task
SYNOPSIS	STATUS errnoSet ( int errorValue /* error status value to set */ )
DESCRIPTION	This routine sets the current <b>errno</b> with a specified error status. It is provided for compatibility with previous versions of VxWorks.
RETURNS	OK, or ERROR if the interrupt nest level is too deep.
ERRNO	N/A
SEE ALSO	errnoLib, errnoGet( ), errnoOfTaskSet( )

NAME	etsecRegister() – register with the VxBus subsystem
SYNOPSIS	void etsecRegister(void)
DESCRIPTION	This routine registers the ETSEC driver with VxBus as a child of the PLB bus type.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	vxbEtsecEnd

#### eventClear()

**NAME** eventClear() – Clear the calling task's events register

SYNOPSIS STATUS eventClear (void)

**DESCRIPTION** This routine clears the calling task's events register. Since events can be received at any time, the caller cannot assume its events register is actually cleared by the time this routine returns unless interrupts are locked when this routine is called.

RETURNS OK on success or ERROR

ERRNO S\_intLib\_NOT\_ISR\_CALLABLE Routine was called from an ISR.

SEE ALSO eventLib, eventReceive()

#### eventReceive()

NAME	<pre>eventReceive() – Wait for event(s</pre>	)							
SYNOPSIS	STATUS eventReceive								
	UINT32 events,	/*	events	task	is	waiting	to	occur	*/

```
UINT8 options, /* user options */
int timeout, /* ticks to wait */
UINT32 *pEventsReceived /* events occurred are returned through this */
)
```

DESCRIPTION

Pends calling task until one or all wanted *events* have been received. When the specified events have been received, they are copied from the events register to the variable pointed to by *pEventsReceived*, and the events register is cleared (by default).

The *options* parameter is used to control various aspects of this routine's behaviour. One of which is to specify if the caller wishes to wait for all events to be received or only one. One of the following must be specified:

EVENTS\_WAIT\_ANY (0x1)

Wait for any one of the wanted events.

EVENTS\_WAIT\_ALL (0x0)

Wait for all wanted events.

Another option is to specify if the events written to *pEventsReceived* are only those received and wanted or all events received. Note that an event can be received at any time, including before **eventReceive()** is called. By default this routine returns only wanted events unless the following option is specified:

#### EVENTS\_RETURN\_ALL (0x2)

Causes the routine to return all received events whether they are wanted (as specified in *events*) or not. It also causes all events to be cleared from the task's events register.

The third option available allows the caller to specify if the received unwanted events are to be cleared from the calling task's events register. They are cleared by default unless the following option is specified:

#### EVENTS\_KEEP\_UNWANTED (0x4)

Tells this routine not to clear unwanted events. In cases where the **EVENTS\_RETURN\_ALL** option is used, all events are cleared even if this option is selected. Wanted events are always cleared hence this option has not effect on them.

Lastly, it is possible to retrieve events that have already been received without affecting the events register by selecting the following option:

#### EVENTS\_FETCH (0x80)

If this option is specified, the contents of the calling task's events register are copied to the location pointed to by *pEventsReceived* and the routine returns immediately. The events are not cleared from the register. The *events* and *timeout* arguments are ignored and so are all other options specified.

The *timeout* parameter specifies the number of ticks to wait for wanted events. It can also have the following special values:

#### NO\_WAIT (0)

Return immediately, even if no events have arrived.

The received events are copied to the location pointed to by *pEventsReceived* even when the routine returns **ERROR** unless a **NULL** pointer is passed.

#### **WARNING** This routine may not be used from interrupt level because ISRs do not have events registers.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

RETURNS OK on success or ERROR

#### ERRNO S\_eventLib\_TIMEOUT

Wanted events not received before specified timeout expired.

#### S\_eventLib\_NOT\_ALL\_EVENTS

Wanted events not received at the time of the call. This error can only occur if **NO\_WAIT** is specified in *timeout* 

#### S\_objLib\_OBJ\_DELETED

Task is waiting for events from a resource that has been destroyed. See **semEvLib** and **msgQEvLib** documentation for more information.

#### S\_intLib\_NOT\_ISR\_CALLABLE

Function was called from ISR.

#### S\_eventLib\_ZERO\_EVENTS The *events* parameter has been passed a value of 0.

SEE ALSO eventLib, semEvLib, msgQEvLib, eventSend()

#### eventSend()

NAME	eventSend() – Send event(s)	
SYNOPSIS	<pre>STATUS eventSend   (    int taskId, /* task events will be sent to */    UINT32 events /* events to send */   )</pre>	

**DESCRIPTION** Sends specified *events* to a task. Passing a taskId of **NULL** causes the calling task to send events to itself. This routine can be used by an ISR to send events to a task.

Because an event is actually a bit in the 32 bit word *events*, the sending process consists of bitwise ORing *events* with the present contents of the destination task's events register. Therefore the process is said to be non-destructive since the events that may already be present in the task's events register are not affected.

Sending an event to a task that already has the event in its events register does not alter the contents of the register.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

RETURNS	OK on success or ERROR.
ERRNO	<b>S_objLib_OBJ_ID_ERROR</b> <i>taskId</i> is invalid.
	S_eventLib_NULL_TASKID_AT_INT_LEVEL Routine was called from ISR with a taskId of NULL.

SEE ALSO eventLib, eventReceive()

#### excConnect( )

NAME	<b>excConnect()</b> – connect a C routine to an exception vector (PowerPC)					
SYNOPSIS	STATUS excConnect ( VOIDFUNCPTR * vector, /* exception vector to attach to */ VOIDFUNCPTR routine /* routine to be called */ )					
DESCRIPTION	This routine connects a specified C routine to a specified exception vector. An exception stub is created and in placed at <i>vector</i> in the exception table. The address of <i>routine</i> is stored in the exception stub code. When an exception occurs, the processor jumps to the exception stub code, saves the registers, and calls the C routines.					
	The routine can be any normal C code, except that it must not invoke certain operating system functions that may block or perform I/O operations.					
	The registers are saved to an Exception Stack Frame (ESF) placed on the stack of the task that has produced the exception. The structure of the ESF used to save the registers is defined in h/arch/ppc/esfPpc.h.					
	The only argument passed by the exception stub to the C routine is a pointer to the ESF containing the registers values. The prototype of this C routine is described below:					
----------	--					
	<pre>void excHandler (ESFPPC *);</pre>					
	When the C routine returns, the exception stub restores the registers saved in the ESF and continues execution of the current task.					
RETURNS	OK, always.					
ERRNO	Not Available					
SEE ALSO	excArchLib, excIntConnect(), excVecSet()					

### excCrtConnect()

NAME	<b>excCrtConnect()</b> – connect a C routine to a critical exception vector (PowerPC 403)	
SYNOPSIS	STATUS excCrtConnect ( VOIDFUNCPTR * vector, /* exception vector to attach to */ VOIDFUNCPTR routine /* routine to be called */ )	
DESCRIPTION	This routine connects a specified C routine to a specified critical exception vector. An exception stub is created and in placed at <i>vector</i> in the exception table. The address of <i>routine</i> is stored in the exception stub code. When an exception occurs, the processor jumps to the exception stub code, saves the registers, and call the C routines.	
	The routine can be any normal C code, except that it must not invoke certain operating system functions that may block or perform I/O operations.	
	The registers are saved to an Exception Stack Frame (ESF) which is placed on the stack of the task that has produced the exception. The ESF structure is defined in <b>h/arch/ppc/esfPpc.h</b> .	
	The only argument passed by the exception stub to the C routine is a pointer to the ESF containing the register values. The prototype of this C routine is as follows:	
	<pre>void excHandler (ESFPPC *);</pre>	
	When the C routine returns, the exception stub restores the registers saved in the ESF and continues execution of the current task.	
RETURNS	OK, always.	
ERRNO	Not Available	

VxWorks Kernel API Reference, 6.6 excHookAdd()

SEE ALSO excArchLib, excIntConnect(), excIntCrtConnect, excVecSet()

### excHookAdd()

NAME	<b>excHookAdd()</b> – specify a routine to be called with exceptions	
SYNOPSIS	void excHookAdd ( FUNCPTR excepHook /* routine to call when exceptions occur */ )	
DESCRIPTION	This routine specifies a routine that will be called when hardware exceptions occur. The specified routine is called after normal exception handling, which includes displaying information about the error. Upon return from the specified routine, the task that incurred the error is suspended.	
	The exception handling routine should be declared as:	
	<pre>void myHandler     (     int task, /* ID of offending task */     int vecNum, /* exception vector number */     <esfxx> *pEsf /* pointer to exception stack frame */     ) where task is the ID of the task that was running when the exception occurred. ESFxx is</esfxx></pre>	
	architecture-specific and can be found by examining <b>/target/h/arch/</b> arch/esfarch.h; for example, the PowerPC uses ESFPPC.	
	This facility is normally used by <b>dbgLib()</b> to activate its exception handling mechanism. If an application provides its own exception handler, it will supersede the <b>dbgLib</b> mechanism.	
RETURNS	N/A	
ERRNOS	N/A	
SEE ALSO	excLib	

# excInit()

NAME	<b>excInit()</b> – initialize the exception handling package
SYNOPSIS	STATUS excInit ( UINT maxIsrJobs /* must be a power of two */ )
DESCRIPTION	This routine initializes the interrupt-level job deferral facility. This facility provides the ability to defer function execution to task level from interrupt level.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the tExcTask cannot be spawned.
ERRNOS	N/A
SEE ALSO	excLib

# excIntConnect()

NAME	<b>excIntConnect( )</b> – connect a C routine to an asynchronous exception vector (PowerPC, ARM)	
SYNOPSIS	STATUS excIntConnect ( VOIDFUNCPTR * vector, /* exception vector to attach to */ VOIDFUNCPTR routine /* routine to be called */ )	
DESCRIPTION	This routine connects a specified C routine to a specified asynchronous exception vector.	
	When the C routine is invoked, interrupts are still locked. It is the responsibility of the C routine to re-enable the interrupt.	
	The routine can be any normal C code, except that it must not invoke certain operating system functions that may block or perform I/O operations.	
NOTE	On PowerPC, the vector is typically the external interrupt vector 0x500 and the decrementer vector 0x900. An interrupt stub is created and placed at <i>vector</i> in the exception table. The address of <i>routine</i> is stored in the interrupt stub code. When the asynchronous exception occurs the processor jumps to the interrupt stub code, saves only the requested registers, and calls the C routines.	

	Before saving the requested registers, the interrupt stub switches from the current task stack to the interrupt stack. For nested interrupts, no stack-switching is performed, because the interrupt is already set.
NOTE	On the ARM, the address of <i>routine</i> is stored in a function pointer to be called by the stub installed on the IRQ exception vector following an asynchronous exception. This routine is responsible for determining the interrupt source and despatching the correct handler for that source.
	Before calling the routine, the interrupt stub switches to SVC mode, changes to a separate interrupt stack and saves necessary registers. In the case of a nested interrupt, no SVC stack switch occurs.
RETURNS	OK, always.
ERRNO	Not Available
SEE ALSO	excArchLib, excConnect(), excVecSet()

# excIntCrtConnect()

NAME	<b>excIntCrtConnect()</b> – connect a C routine to a critical interrupt vector (PowerPC 403)
SYNOPSIS	<pre>STATUS excIntCrtConnect   (     VOIDFUNCPTR * vector, /* exception vector to attach to */     VOIDFUNCPTR routine /* routine to be called */   )</pre>
DESCRIPTION	This routine connects a specified C routine to a specified asynchronous critical exception vector such as the critical external interrupt vector $(0x100)$ , or the watchdog timer vector $(0x1020)$ . An interrupt stub is created and placed at <i>vector</i> in the exception table. The address of <i>routine</i> is stored in the interrupt stub code. When the asynchronous exception occurs, the processor jumps to the interrupt stub code, saves only the requested registers, and calls the C routines.
	When the C routine is invoked, interrupts are still locked. It is the C routine's responsibility to re-enable interrupts.
	The routine can be any normal C routine, except that it must not invoke certain operating system functions that may block or perform I/O operations.
	Before the requested registers are saved, the interrupt stub switches from the current task stack to the interrupt stack. In the case of nested interrupts, no stack switching is performed, because the interrupt stack is already set.

RETURNS OK, always.

Not Available ERRNO

excArchLib, excConnect(), excCrtConnect, excVecSet() SEE ALSO

### excJobAdd()

NAME excJobAdd() - request a task-level function call from interrupt level

SYNOPSIS

STATUS	excJobAdd	1	
(			
VO	DFUNCPTR	func,	
int	5	arg1,	

int arg2, int arg3, int arg4, arg5, int int arg6

)

- This routine allows interrupt level code to request a function call to be executed by the DESCRIPTION tExcTask at task-level.
- Care must be taken when pushing jobs to tExcTask. Jobs that may block, hang, or generate WARNING exceptions must be avoided since blocking or suspension of the tExcTask may cause other parts of the system to misbehave.
- RETURNS OK. Otherwise ERROR if job posting fails.

ERRNO N/A

excLib SEE ALSO

### excVecGet()

NAME excVecGet() – get a CPU exception vector (PowerPC, ARM)

SYNOPSIS FUNCPTR excVecGet

```
FUNCPTR * vector /* vector offset */
```

**DESCRIPTION** This routine returns the address of the C routine currently connected to *vector*.

**RETURNS** The address of the C routine.

ERRNO Not Available

SEE ALSO excArchLib, excVecSet()

### excVecInit()

NAME excVecInit() – initialize the exception/interrupt vectors

SYNOPSIS STATUS excVecInit (void)

**DESCRIPTION** This routine sets all exception vectors to point to the appropriate default exception handlers. These handlers will safely trap and report exceptions caused by program errors or unexpected hardware interrupts.

#### MC680x0:

All vectors from vector 2 (address 0x0008) to 255 (address 0x03fc) are initialized. Vectors 0 and 1 contain the reset stack pointer and program counter.

#### MIPS:

All MIPS exception, trap, and interrupt vectors are set to default handlers.

#### **x86**:

All vectors from vector 0 (offset (0x0000) to 255 (offset 0x07f8) are initialized to default handlers. A global variable excDoBell controls the bell that takes 660 microsecs in the default exception show routine. The default value is **TRUE**. To turn the bell off, set it **FALSE**.

#### **PowerPC**:

There are 48 vectors and only vectors that are used are initialized.

#### SH:

There are 256 vectors, initialized with the default exception handler (for exceptions) or the unitialized interrupt handler (for interrupts). On SH-2, vectors 0 and 1 contain the power-on reset program counter and stack pointer. Vectors 2 and 3 contain the manual reset program counter and stack pointer. On SH-3 and SH-4 processors the vector table is located at (vbr + 0x800), and the (exception code / 8) value is used as vector offset.

	The first two vectors are reserved for special use: "trapa #0" (offset 0x0) to implement software breakpoint, and "trapa #1' (offset 0x4) to detect integer zero divide exception.
	<b>ARM</b> : All exception vectors are initialized to default handlers except 0x14 (Address) which is now reserved on the ARM and 0x1C (FIQ), which is not used by VxWorks.
	<b>SimSolaris/SimNT</b> : This routine does nothing on both simulators and always returns <b>OK</b> .
NOTE	This routine is usually called from the system start-up routine, <b>usrInit()</b> , in <b>usrConfig.c</b> . It must be called before interrupts are enabled.
RETURNS	OK, always.
ERRNO	Not Available
SEE ALSO	excArchLib, excLib

# excVecSet()

NAME	excVecSet() – set a CPU exception vector (PowerPC, ARM)
SYNOPSIS	<pre>void excVecSet   (    FUNCPTR * vector, /* vector offset */    FUNCPTR function /* address to place in vector */   )</pre>
DESCRIPTION	This routine specifies the C routine that will be called when the exception corresponding to <i>vector</i> occurs. This routine does not create the exception stub; it simply replaces the C routine to be called in the exception stub.
NOTE ARM	On the ARM, there is no <b>excConnect()</b> routine, unlike the PowerPC. The C routine is attached to a default stub using <b>excVecSet()</b> .
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	excArchLib, excVecGet( ), excConnect( ), excIntConnect( )

# exit()

NAME	exit() – exit a task (ANSI)	
SYNOPSIS	<pre>void exit   (    int code /* code stored in TCB for delete hooks */   )</pre>	
DESCRIPTION	This routine is called by a task to cease to exist as a task. It is called implicitly when the "main" routine of a spawned task is exited. The <i>code</i> parameter will be stored in the <b>WIND_TCB</b> for possible use by the delete hooks, or post-mortem debugging.	
SMP CONSIDERATIONS		
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.	
RETURNS	N/A	
ERRNO	N/A	
SEE ALSO	<b>taskLib</b> , <b>taskDelete( )</b> , American National Standard for Information Systems -, Programming Language - C, ANSI X3.159-1989: Input/Output ( <b>stdlib.h</b> ), The VxWorks Programmer's Guide	

# expf()

NAME	<pre>expf() – compute an exponential value (ANSI)</pre>
SYNOPSIS	<pre>float expf   (    float x /* exponent */   )</pre>
DESCRIPTION	This routine returns the exponential of $x$ in single precision.
RETURNS	The single-precision exponential value of <i>x</i> .
ERRNO	Not Available
SEE ALSO	mathALib

### fabsf()

NAME	fabsf() – compute an absolute value (ANSI)		
SYNOPSIS	float fabsf ( float v /* number to return the absolute value of */ )		
DESCRIPTION	This routine returns the absolute value of $v$ in single precision.		
RETURNS	The single-precision absolute value of <i>v</i> .		
ERRNO	Not Available		
SEE ALSO	mathALib		

### fastStrSearch()

fastStrSearch() - Search by optimally choosing the search algorithm NAME SYNOPSIS char \* fastStrSearch ( char \* pattern, /\* pattern to search for \*/ int patternLen, /\* length of the pattern \*/ char \* buffer, /\* text buffer to search in \*/ int bufferLen, /\* length of the text buffer \*/ BOOL caseSensitive /\* case-sensitive search? \*/ ) DESCRIPTION Depending on the pattern size, this function uses either the Boyer-Moore-Sunday algorithm or the Brute Force algorithm. The Boyer-Moore-Sunday algorithm requires pre-processing, therefore for small patterns it is better to use the Brute Force algorithm. A pointer to the located pattern, or a NULL pointer if the pattern is not found RETURNS ERRNO Not Available strSearchLib SEE ALSO

VxWorks Kernel API Reference, 6.6 fccRegister()

# fccRegister()

SEE ALSO	fccVxbEnd	
ERRNO	N/A	
RETURNS	N/A	
DESCRIPTION	This routine registers the FCC driver with VxBus as a child of the PCI bus type.	
SYNOPSIS	void fccRegister(void)	
NAME	<b>fccRegister()</b> – register with the VxBus subsystem	

# fchmod()

NAME	<b>fchmod()</b> – change the permission mode of a file	
SYNOPSIS	<pre>int fchmod   (    int fd,    mode_t mode   )</pre>	
DESCRIPTION	The fchmod function changes or assigns the mode of a file. The mode of a file specifies its permissions and other attributes. Note that this routine receives <b>open file descriptor</b> as the first argument compairing to chmod routine.	
	The value of <i>mode</i> is bitwise inclusive OR of the permissions to be assigned	
	These permission constants are defined in <i>sys/stat.h</i> as follows:	
	S_IRUSR Read permission, owner.	
	S_IWUSR Write permission, owner.	
	S_IXUSR Execute/search permission, owner.	
	S_IRWXU Read/write/execute permission, owner.	
	S_IRGRP Read permission, group.	

	S_IWGRP Write permission, group.
	S_IXGRP Execute/search permission, group.
	S_IRWXG Read/write/execute permission, group.
	S_IROTH Read permission, other.
	S_IWOTH Write permission, other.
	S_IXOTH Execute/search permission, other.
	S_IRWXO Read/write/execute permission, other.
RETURNS	If it succeeds, returns <b>OK</b> , 0. Otherwise, <b>ERROR</b> , -1 is returned, errno is set to indicate the error and no change is done to the file.
	The following example changes the mode of the file "myFile" to owner Read/write/execute, group Read and other Read:
	fd = open ("myFile", O_RDONLY, 0 ); status = fchmod (fd, S_IRWXU   S_IRGRP   S_IROTH );
ERRNO	<b>EBADF</b> The <i>fd</i> argument is not a valid open file.
	others Other errors reported by device driver.
SEE ALSO	fsPxLib

# fcntl()

NAME	<b>fcntl()</b> – perform control functions over open files	
SYNOPSIS	<pre>int fcntl   (    int fd,    int command,  )</pre>	

VxWorks Kernel API Reference, 6.6 fdatasync()

DESCRIPTION	The <b>fcntl()</b> function provides for control over open files. The <i>fd</i> argument is an open file descriptor. The <b>fcntl()</b> function may take a third argument whose data type, value and use depend upon the value of <i>command</i> which specifies the operation to be performed by <b>fcntl()</b> .
RETURNS	Not Available
ERRNO	EMFILE Ran out of file descriptors
	EBADF Bad file descriptor number.
	ENOSYS Device driver does not support the ioctl command.
	ENXIO Device and its driver are removed. <b>close()</b> should be called to release this file descriptor.
	Other Other errors reported by device driver.
SEE ALSO	ioLib

# fdatasync()

NAME	<b>fdatasync( )</b> – synchronize a file data	
SYNOPSIS	<pre>int fdatasync   (    int fd /* file descriptor of the file to datasync */ )</pre>	
DESCRIPTION	The function forces all currently queued I/O operations associated with the file indicated by $fd$ to the synchronized I/O completion state.	
	The functionality is as described for $fsync()$ with the exception that all I/O operations are completed as defined for synchronised I/O data integrity completion.	
RETURNS	Upon successful completion, <b>OK</b> , 0 is returned. Otherwise, <b>ERROR</b> , -1 returned and errno is set to indicate the error. If the <b>fdatasync()</b> function fails, outstanding I/O operations are not guaranteed to have been completed.	

ERRNO

SEE ALSO fsPxLib, fsync()

# fdprintf()

NAME	<b>fdprintf()</b> – write a formatted string to a file descriptor		
SYNOPSIS	<pre>int fdprintf  (     int fd, /* file descriptor to write to */     const char * fmt, /* format string to write */     /* optional arguments to format */ )</pre>		
DESCRIPTION	This routine writes a formatted string to a specified file descriptor. Its function and syntax are otherwise identical to <b>printf()</b> .		
SMP CONSIDERATIC	This API is spinlock and intCpuLock restricted.		
RETURNS	The number of characters output, or <b>ERROR</b> if there is an error during output.		
ERRNO	Not Available		
SEE ALSO	fioLib, printf()		

# fecRegister()

NAME	<b>fecRegister()</b> – register with the VxBus subsystem	
SYNOPSIS	void fecRegister(void)	
DESCRIPTION	This routine registers the FEC driver with VxBus as a child of the PCI bus type.	
RETURNS	N/A	
ERRNO	N/A	
SEE ALSO	fecVxbEnd	

VxWorks Kernel API Reference, 6.6 feiRegister()

# feiRegister()

NAME	<b>feiRegister()</b> – register with the VxBus subsystem	
SYNOPSIS	void feiRegister(void)	
DESCRIPTION	This routine registers the Intel 8255x driver with VxBus as a child of the PCI bus type.	
RETURNS	N/A	
ERRNO	N/A	
SEE ALSO	fei8255xVxbEnd	

# ffsLsb()

NAME	ffsLsb() – find least significant bit set	
SYNOPSIS	<pre>int ffsLsb   (    UINT32 i /* value in which to find first set bit */ )</pre>	
DESCRIPTION	This routine finds the least significant bit set in the 32 bit argument passed to it and returns the index of that bit. Bits are numbered starting at 1 from the least significant bit. A return value of zero indicates that the value passed is zero.	
RETURNS	index of least significant bit set, or zero	
ERRNO	N/A	
SEE ALSO	ffsLib	

# ffsMsb()

NAME	ffsMsb() - find	l most significant bit set
		~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~

SYNOPSIS int ffsMsb

 ( UINT32 i /\* value in which to find first set bit \*/ )
 DESCRIPTION This routine finds the most significant bit set in the 32 bit argument passed to it and returns the index of that bit. Bits are numbered starting at 1 from the least significant bit. A return value of zero indicates that the value passed is zero.
 RETURNS index of most significant bit set, or zero
 ERRNO N/A
 SEE ALSO ffsLib

### fileUploadPathClose()

NAME	fileUploadPathClose() – close the event-destination file
SYNOPSIS	<pre>void wvFileUploadPathClose   (    UPLOAD_ID pathId /* generic upload-path descriptor */ )</pre>
DESCRIPTION	This routine closes the file associated with <i>pathId</i> that is serving as a destination for event data.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	wvFileUploadPathLib, wvFileUploadPathCreate()

### fioBaseLibInit()

NAME fioBaseLibInit() – initialize the formatted I/O support library

SYNOPSIS void fioBaseLibInit (void)

**DESCRIPTION** This routine initializes the formatted I/O support library. It should be called once in **usrRoot()** when formatted I/O functions such as **printf()** and **scanf()** are used.

VxWorks Kernel API Reference, 6.6 fioFormatV()

RETURNS N/A

ERRNO Not Available

SEE ALSO fioBaseLib

### fioFormatV()

NAME

fioFormatV() – convert a format string

```
SYNOPSIS int fioFormatV

(

FAST const char *fmt, /* format string */

va_list vaList, /* pointer to varargs list */

FUNCPTR outRoutine, /* handler for args as they're formatted */

int outarg /* argument to routine */

)
```

**DESCRIPTION** This routine is used by the **printf()** family of routines to handle the actual conversion of a format string. The first argument is a format string, as described in the entry for **printf()**. The second argument is a variable argument list *vaList* that was previously established.

As the format string is processed, the result will be passed to the output routine whose address is passed as the third parameter, *outRoutine*. This output routine may output the result to a device, or put it in a buffer. In addition to the buffer and length to output, the fourth argument, *outarg*, will be passed through as the third parameter to the output routine. This parameter could be a file descriptor, a buffer address, or any other value that can be passed in an "int".

The output routine should be declared as follows:

STATUS outRoutine
 (
 char \*buffer, /\* buffer passed to routine \*/
 int nchars, /\* length of buffer \*/
 int outarg /\* arbitrary arg passed to fmt routine \*/

The output routine should return **OK** if successful, or **ERROR** if unsuccessful.

**RETURNS** The number of characters output, or **ERROR** if the output routine returned **ERROR**.

ERRNO Not Available

SEE ALSO fioBaseLib

# fioLibInit( )

NAME	<pre>fioLibInit() - initialize the formatted I/O support library</pre>
SYNOPSIS	void fioLibInit (void)
DESCRIPTION	This routine initializes the formatted I/O support library. It should be called once in <b>usrRoot()</b> when formatted I/O functions such as <b>printf()</b> and <b>scanf()</b> are used.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	fioLib

# fioRdString()

NAME	<pre>fioRdString() - read a string from a file</pre>	
SYNOPSIS	<pre>int fioRdString   (     int fd, /* fd of device to read */     FAST char string[], /* buffer to receive input */     int maxbytes /* max no. of chars to read */   )</pre>	
DESCRIPTION	This routine puts a line of input into <i>string</i> . The specified input file descriptor is read until <i>maxbytes</i> , an <b>EOF</b> , an EOS, or a newline character is reached. A newline character or <b>EOF</b> is replaced with EOS, unless <i>maxbytes</i> characters have been read.	
SMP CONSIDERATIO	INS	
	This API is spinlock and intCpuLock restricted.	
RETURNS	The length of the string read, including the terminating EOS; or <b>EOF</b> if a read error occurred or end-of-file occurred without reading any other character.	
ERRNO	Not Available	
SEE ALSO	fioLib	

# fioRead()

NAME	fioRead() – read a buffer	
SYNOPSIS	<pre>int fioRead   (    int fd, /* file descriptor of file to read */    char * buffer, /* buffer to receive input */    int maxbytes /* maximum number of bytes to read */   )</pre>	
DESCRIPTION	This routine repeatedly calls the routine <b>read()</b> until <i>maxbytes</i> have been read into <i>buffer</i> . If <b>EOF</b> is reached, the number of bytes read will be less than <i>maxbytes</i> .	
SMP CONSIDERATIO	This API is spinlock and intCpuLock restricted.	
RETURNS	The number of bytes read, or <b>ERROR</b> if there is an error during the read operation.	
ERRNO	Not Available	
SEE ALSO	fioLib, read()	

# floorf()

NAME	<b>floorf()</b> – compute the largest integer less than or equal to a specified value (ANSI)
SYNOPSIS	<pre>float floorf   (    float v /* value to find the floor of */   )</pre>
DESCRIPTION	This routine returns the largest integer less than or equal to <i>v</i> , in single precision.
RETURNS	The largest integral value less than or equal to $v$ , in single precision.
ERRNO	Not Available
SEE ALSO	mathALib

### fmodf()

NAME	<b>fmodf( )</b> – compute the remainder of $x/y$ (ANSI)
SYNOPSIS	<pre>float fmodf   (    float x, /* numerator */    float y /* denominator */   )</pre>
DESCRIPTION	This routine returns the remainder of $x/y$ with the sign of $x$ , in single precision.
RETURNS	The single-precision modulus of $x/y$ .
ERRNO	Not Available
SEE ALSO	mathALib

# formatTrans()

NAME	formatTrans() – Format a transaction disk.
SYNOPSIS	<pre>STATUS formatTrans   (     int fd,</pre>
DESCRIPTION	The <i>fd</i> argument should be the result of opening a rawFS disk. (Because we just use ordinary read/write/fstat calls, this can also be any ordinary file descriptor for a file on a disk; we will just pretend that that file is a whole disk. But in general it should be a rawFS.)
	The <i>blkshift</i> argument is used to increase the logical block size of the underlying device. This reduces the TRFS overhead, but increases the minimum size of any file on the dosFS atop this TRFS layer. That is, the TRFS overhead goes down but the file system overhead goes up. In general this should just be 0.
	The <i>overhead</i> parameter is 10 times the percent of the disk itself to use as TRFS transactional space, i.e., an argument of 50 gives 5%, while an argument of 100 gives 10%. A value of 0 results in the default of 5%.
	The <i>type</i> parameter should be one of:

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FORMAT\_REGULAR (0): puts transaction master records (TMRs at the beginning and end of the volume.
 FORMAT\_TFFS (1): moves the first TMR to sector 1, leaving sector 0 available for other purposes (such as TFFS internals).
 FORMAT\_DOS (2): not actually supported (yet?).
 RETURNS
 OK if all went well, ERROR otherwise.
 ERRNO
 EINVAL – invalid arguments or inappropriate underlying device E2BIG – underlying device too big other errno set by rawFS
 SEE ALSO
 xbdTrans, usrFormatTrans()

# fpathconf()

NAME	<b>fpathconf()</b> – determine the current value of a configurable limit
SYNOPSIS	<pre>long fpathconf   (     int fd, /* file descriptor of the file */     int name /* Value to query */   )</pre>
DESCRIPTION	The <b>fpathconf()</b> and <b>pathconf()</b> functions provide a method for the application to determine the current value of a configurable limit or option (variable) that is associated with a file or directory.
RETURNS	The current value is returned if valid with the query. Otherwise, <b>ERROR</b> , -1 returned and errno may be set to indicate the error. There are many reasons to return <b>ERROR</b> . If the variable corresponding to name has no limit for the path or file descriptor, both <b>pathconf()</b> and <b>fpathconf()</b> return -1 without changing errno.
ERRNO	
SEE ALSO	fsPxLib, pathconf()

# fppInit()

NAME	<pre>fppInit() - initialize floating-point coprocessor support</pre>
SYNOPSIS	void fppInit (void)
DESCRIPTION	This routine initializes floating-point coprocessor support and must be called before using the floating-point coprocessor. This is done automatically by the root task, <b>usrRoot()</b> , in <b>usrConfig.c</b> when the configuration macro <b>INCLUDE_HW_FP</b> is defined.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	fppLib

### fppProbe( )

NAME **fppProbe()** – probe for the presence of a floating-point coprocessor SYNOPSIS STATUS fppProbe (void) DESCRIPTION This routine determines whether there is a floating-point coprocessor in the system. The implementation of this routine is architecture-dependent: MC680x0, x86, SH-4: This routine sets the illegal coprocessor opcode trap vector and executes a coprocessor instruction. If the instruction causes an exception, **fppProbe()** returns **ERROR**. Note that this routine saves and restores the illegal coprocessor opcode trap vector that was there prior to this call. The probe is only performed the first time this routine is called. The result is stored in a static and returned on subsequent calls without actually probing. MIPS: This routine simply reads the R-Series status register and reports the bit that indicates whether coprocessor 1 is usable. This bit must be correctly initialized in the BSP. ARM: This routine currently returns **ERROR** to indicate no floating-point coprocessor support.

VxWorks Kernel API Reference, 6.6 fppRestore()

#### SimNT, SimSolaris: This routine currently returns OK.

- **RETURNS** OK, or ERROR if there is no floating-point coprocessor.
- ERRNO Not Available
- SEE ALSO fppArchLib

### fppRestore()

NAME	<b>fppRestore()</b> – restore the floating-point coprocessor context	
SYNOPSIS	void fppRestore ( EP_CONTEXT * pEpContext /* where to restore context from */	
	)	

**DESCRIPTION** This routine restores the floating-point coprocessor context. The context restored is:

### &MC680x0:

- registers fpcr, fpsr, and fpiar
- registers f0 f7
- internal state frame (if NULL, the other registers are not saved.)

#### &MIPS:

- register fpcsr
- registers fp0 fp31

#### &SH-4:

- registers fpcsr and fpul
- registers fr0 fr15
- registers xf0 xf15

#### &x86:

108 byte old context with fsave and frstor instruction

- control word, status word, tag word,
- instruction pointer,
- instruction pointer selector,
- last FP instruction op code,
- data pointer,
- data pointer selector,

512 byte new context with fxsave and fxrstor instruction

- control word, status word, tag word,
- last FP instruction op code,
- instruction pointer,
- instruction pointer selector,
- data pointer,
- data pointer selector,
- registers st/mm0 st/mm7 (10 bytes \* 8)
- registers xmm0 xmm7 (16 bytes \* 8)

#### &ARM:

- currently, on this architecture, this routine does nothing.

#### &SimSolaris:

- register fsr
- registers f0 f31

### &SimNT:

- this routine does nothing on Windows simulator.

RETURNS	N/A

ERRNO Not Available
---------------------

SEE ALSO fppArchLib, fppSave()

# fppSave()

NAME	<b>fppSave()</b> – save the floating-point coprocessor context
SYNOPSIS	<pre>void fppSave   (    FP_CONTEXT * pFpContext /* where to save context */   )</pre>
DESCRIPTION	This routine saves the floating-point coprocessor context. The context saved
	&MC680x0: - registers fpcr, fpsr, and fpiar - registers f0 - f7

is:

- internal state frame (if NULL, the other registers are not saved.)

#### &MIPS:

- register **fpcsr**
- registers fp0 fp31

#### &SH-4:

- registers **fpcsr** and **fpul**
- registers fr0 fr15
- registers xf0 xf15

#### &x86:

108 byte old context with fsave and frstor instruction

- control word, status word, tag word,
- instruction pointer,
- instruction pointer selector,
- last FP instruction op code,
- data pointer,
- data pointer selector,
- registers st/mm0 st/mm7 (10 bytes \* 8)
- 512 byte new context with fxsave and fxrstor instruction
- control word, status word, tag word,
- last FP instruction op code,
- instruction pointer,
- instruction pointer selector,
- data pointer,
- data pointer selector,
- registers st/mm0 st/mm7 (10 bytes \* 8)
- registers xmm0 xmm7 (16 bytes \* 8)

#### &ARM:

- currently, on this architecture, this routine does nothing.

#### &SimSolaris:

- register fsr
- registers f0 f31

#### &SimNT:

- this routine does nothing on Windows simulator. Floating point registers are saved by Windows.

RETURNS N/A

ERRNO Not Available

SEE ALSO fppArchLib, fppRestore()

### fppShowInit( )

NAME fppShowInit() - initialize the floating-point show facility SYNOPSIS void fppShowInit (void) DESCRIPTION This routine links the floating-point show facility into the VxWorks system. It is called automatically when the floating-point show facility is configured into VxWorks using either of the following methods: If you use the configuration header files, define INCLUDE\_SHOW\_ROUTINES in config.h. If you use the Tornado project facility, select INCLUDE\_HW\_FP\_SHOW. \_ N/A RETURNS Not Available EBBNO fppShow SEE ALSO

### fppTaskRegsGet()

NAME fppTaskRegsGet() – Gets FPU context for a task
SYNOPSIS STATUS fppTaskRegsGet
(
int task,
FPREG\_SET \*pFpRegSet
)
DESCRIPTION Gets the FPU context for a task.

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VxWorks Kernel API Reference, 6.6 fppTaskRegsGet()

RETURNS	OK on success. ERROR otherwise
ERRNO	S_coprocLib_INVALID_OPERATION S_coprocLib_INVALID_ARGUMENT S_coprocLib_NO_COPROC_SUPPORT
SEE ALSO	aimFppLib

# fppTaskRegsGet( )

NAME	<b>fppTaskRegsGet()</b> – get the floating-point registers from a task TCB
SYNOPSIS	STATUS fppTaskRegsGet ( int task, /* task to get info about */ FPREG_SET * pFpRegSet /* ptr to floating-point register set */ )
DESCRIPTION	This routine copies a task's floating-point registers and/or status registers to the locations whose pointers are passed as parameters. The floating-point registers are copied into an array containing all the registers.
NOTE	This routine only works well if <i>task</i> is not the calling task. If a task tries to discover its own registers, the values will be stale (that is, left over from the last task switch).
RETURNS	<b>OK</b> , or <b>ERROR</b> if there is no floating-point support or there is an invalid state.
ERRNO	Not Available
SEE ALSO	fppArchLib, fppTaskRegsSet( )

# fppTaskRegsSet( )

 NAME
 fppTaskRegsSet() – Sets FPU context for a task

 SYNOPSIS
 STATUS

 (int
 task,

F	PREG_	_SET	*pFpRegSet
)			

DESCRIPTION	Sets the FPU context for a task.
RETURNS	OK on success. ERROR otherwise
ERRNO	S_coprocLib_INVALID_OPERATION S_coprocLib_INVALID_ARGUMENT S_coprocLib_NO_COPROC_SUPPORT
SEE ALSO	aimFppLib

### fppTaskRegsSet( )

fppTaskRegsSet( ) - set the floating-point registers of a task NAME SYNOPSIS STATUS fppTaskRegsSet ( task, /\* task to set registers for \*/ int FPREG\_SET \* pFpRegSet /\* ptr to floating-point register set \*/ ) DESCRIPTION This routine loads the specified values into the TCB of a specified task. The register values are copied from the array at *pFpRegSet*. RETURNS OK, or ERROR if there is no floating-point support or there is an invalid state. ERRNO Not Available fppArchLib, fppTaskRegsGet( ) SEE ALSO

### fppTaskRegsShow()

NAME **fppTaskRegsShow()** – print the contents of a task's floating-point registers

SYNOPSIS void fppTaskRegsShow

```
(

int task /* task to display floating point registers for */

)

DESCRIPTION This routine prints to standard output the contents of a task's floating-point registers.
```

RETURNS N/A

ERRNO Not Available

SEE ALSO fppShow

### free()

NAME	<b>free( )</b> – free a block of memory from the system memory partition (ANSI)
SYNOPSIS	<pre>void free   (    void * ptr /* pointer to block of memory to free */ )</pre>
DESCRIPTION	This routine returns to the free memory pool (kernel heap) a block of memory previously allocated with <b>malloc()</b> , <b>calloc()</b> , <b>memalign()</b> , <b>realloc()</b> or <b>valloc()</b> . If <i>ptr</i> is a null pointer, no action occurs.
RETURNS	N/A
ERRNO	Possible errnos generated by this routine include:
	S_memLib_BLOCK_ERROR The block of memory to free is not valid.
SEE ALSO	<b>memPartLib</b> , <b>malloc()</b> , <b>calloc()</b> , <b>memPartFree()</b> , <i>American National Standard for Information</i> Systems -, Programming Language - C, ANSI X3.159-1989: General Utilities ( <b>stdlib.h</b> )

### fsEventUtilInit()

NAME **fsEventUtilInit()** – Initialize the file system event utility library

SYNOPSIS STATUS fsEventUtilLibInit

( void )

- DESCRIPTIONnoneRETURNSOK on success, ERROR on failureERRNONot Available
- SEE ALSO fsEventUtilLib

### fsMonitorInit()

- NAME fsMonitorInit() Initialize the fsMonitor
- SYNOPSIS STATUS fsMonitorInit(void)
- **DESCRIPTION** This routine initializes the fsMonitor library.
- RETURNS OK or ERROR
- ERRNO Not Available
- SEE ALSO fsMonitor

### fsPathAddedEventRaise()

NAME	fsPathAddedEventRaise() – Raise a "path added" event
SYNOPSIS	void fsPathAddedEventRaise ( char * coreIOPath )
DESCRIPTION	This routine raises an event with the event reporting framework when the specified path has been added to core I/O by a file system. This routine will cause the wait for path handler(s) to run.
RETURNS	Not Available

2

VxWorks Kernel API Reference, 6.6 fsPathAddedEventSetup()

ERRNO Not Available

SEE ALSO fsEventUtilLib

### fsPathAddedEventSetup()

NAME	<b>fsPathAddedEventSetup()</b> – Setup to wait for a path
SYNOPSIS	STATUS fsPathAddedEventSetup ( FS_PATH_WAIT_STRUCT * pWaitData, char * path )
DESCRIPTION	This routine registers with the ERF to wait for the specified path to be added to core I/O by a file system. This is mainly used my file system formatters when the eject the current file system and wait for rawFs to instantiate.
RETURNS	OK on success or ERROR on error.

ERRNO Not Available

SEE ALSO fsEventUtilLib

### fsWaitForPath()

NAME	fsWaitForPath() – wait for a path
SYNOPSIS	STATUS fsWaitForPath ( FS_PATH_WAIT_STRUCT *pWaitData )
DESCRIPTION	This routine waits for a path to be added to core I/O. The function, fsPathAddedEventSetup, must be prior for pWaitData to be setup. This function simply waits on the semaphore provided in pWaitData which will be given by the wait for path handler.
RETURNS	OK on success or ERROR on failure

ERRNO Not Available

SEE ALSO fsEventUtilLib

### fsmGetDriver()

NAME	<b>fsmGetDriver()</b> – Get the XBD name of a mapping based on the path
SYNOPSIS	STATUS fsmGetDriver ( char *volume, /* core I/O pathname */ devname_t driver /* xbd driver name */ )
DESCRIPTION	This routine gets the XBD name which is currently mapped to volume, if such a mapping exists.
	The <i>volume</i> parameter specifies the pathname for which a driver name is to be retrieved.
	The <i>driver</i> parameter is the resultant of the name mapping, if found.
RETURNS	<b>OK</b> if the mapping is found, or <b>ERROR</b> if no such mapping exists.
ERRNO	Not Available
SEE ALSO	fsMonitor

### fsmGetVolume()

NAME	<b>fsmGetVolume( )</b> – get the pathname based on an XBD name mapping		
SYNOPSIS	STATUS fsmGetVolume ( char *driver, /* XBD driver name */ fsmName_t volume /* core I/O pathname */ )		
DESCRIPTION	This routine retrieves the pathname associated with an XBD name if such a mapping exists.		
	The <i>driver</i> parameter specifies a driver name for which a pathname is to be retrieved.		
	The volume parameter specifies the resultant pathname.		

VxWorks Kernel API Reference, 6.6 fsmNameInstall()

**RETURNS** OK if the mapping is retrieved, or **ERROR** if it does not exist.

ERRNO Not Available

SEE ALSO fsMonitor

### fsmNameInstall()

NAME fsmNameInstall() - Add a mapping between an XBD name and a pathname SYNOPSIS STATUS fsmNameInstall ( char \*driver, /\* XBD driver name \*/ char \*volume /\* core I/O pathname \*/ ) This routine creates a mapping between the driver name and the pathname specified by DESCRIPTION volume. This mapping will persist until removed by fsmNameUninstall. The *driver* parameter specifies the XBD name to be mapped. The *volume* parameter specifies the Core I/O path that driver is to be mapped into. RETURNS OK if successful or ERROR if the name cannot be added Not Available ERRNO fsMonitor SEE ALSO

### fsmNameMap()

NAME	<b>fsmNameMap( )</b> – map an XBD name to a Core I/O pat	
SYNOPSIS	STATUS fsmNameMap (	
	devname_t xbdName, /* XBD name */	
	fsmName_t volName  /* core I/O path */ )	

DESCRIPTION	This function maps an XBD name to a path in Core I/O, either by an explicit mapping specified by <b>fsmNameInstall()</b> or by using the XBD name to create a Core I/O pathname. This function always succeeds.
	The <i>xbdName</i> parameter specifies the name of the device to be mapped.
	The <i>volName</i> parameter specifies the resultant mapped name.
RETURNS	ОК
ERRNO	Not Available
SEE ALSO	fsMonitor

### fsmNameUninstall()

NAME	<b>fsmNameUninstall()</b> – remove an XBD name to pathname mapping
SYNOPSIS	STATUS fsmNameUninstall ( char *driver /* driver name */ )
DESCRIPTION	This routine removes a name mapping added by fsmNameInstall. After invocation of this routine, the F/S monitor will create a pathname based on the XBD name instead of using a name mapping.
	The <i>driver</i> parameter specifies the name of the XBD for which a mapping is to be removed. All occurences of driver are removed.
RETURNS	<b>OK</b> if the name mapping is removed or <b>ERROR</b> if the mapping is not found.
ERRNO	Not Available
SEE ALSO	fsMonitor

# fsmProbeInstall()

NAME fsmProbeInstall() – install F/S probe and instantiator functions

SYNOPSIS STATUS fsmProbeInstall

```
      fsmProbeFunc probe, /* probe routine to install */

      fsmInstFunc inst /* instantiator routine to install */

      )

      DESCRIPTION

      This routine installs the file system probe and instantiator functions. When a new file system is discovered, the file system monitor will call these functions to test for, and instantiate a particular file system type. If either function fails, then the file system is not created.

      The probe parameter specifies the probe function to be used.

      The inst parameter specifies the instantiator function to be used when the probe function succeeds.

      RETURNS
      OK on success, or ERROR if an error is detected.
```

ERRNO Not Available

SEE ALSO fsMonitor

### fsmProbeUninstall()

NAME	<pre>fsmProbeUninstall() - remove a file system probe</pre>
SYNOPSIS	STATUS fsmProbeUninstall ( fsmProbeFunc probe /* probe routine to uninstall */ )
DESCRIPTION	This routine removes all probe-instantiator pairs that match the probe parameter. The <i>probe</i> parameter specifies the probe function of the probe-instantiator pair to be removed.
RETURNS	0 if a probe is removed or <b>ERROR</b> if the probe is not found
ERRNO	Not Available
SEE ALSO	fsMonitor

### fsmUnmountHookAdd()

NAME	fsmUnmountHookAdd() – Add an unmount hook function
SYNOPSIS	STATUS fsmUnmountHookAdd ( FUNCPTR fn )
DESCRIPTION	This routine adds a hook routine to run when a vnode based file system unmounts. The <i>fn</i> parameter specifies the hook function.
RETURNS	<b>OK</b> if there is space in the table and <i>fn</i> is not <b>NULL</b> /NOMANUAL
ERRNO	Not Available
SEE ALSO	fsMonitor

# fsmUnmountHookDelete()

NAME	fsmUnmountHookDelete() – Remove an unmount hook function
SYNOPSIS	STATUS fsmUnmountHookDelete ( FUNCPTR fn )
DESCRIPTION	This routine removes a hook routine to run when a vnode based file system unmounts. The <i>fn</i> parameter specifies the hook function.
RETURNS	<b>OK</b> if <i>fn</i> is found in the table. <b>ERROR</b> otherwise /NOMANUAL
ERRNO	Not Available
SEE ALSO	fsMonitor

VxWorks Kernel API Reference, 6.6 fsmUnmountHookRun()

### fsmUnmountHookRun()

fsmUnmountHookRun() - Runs the unmount hook functions NAME SYNOPSIS void fsmUnmountHookRun ( DEV\_HDR \*pDev ) DESCRIPTION This routine is called my the vnode layer when a file system unmounts to run the unmount hook functions. The *pDev* parameter is passed to each hook function /NOMANUAL Not Available RETURNS Not Available ERRNO fsMonitor SEE ALSO

### fstat()

NAME	<b>fstat( )</b> – get file status information (POSIX)
SYNOPSIS	STATUS fstat ( int fd, /* file descriptor for file to check */ struct stat *pStat /* pointer to stat structure */ )
DESCRIPTION	This routine obtains various characteristics of a file (or directory). The file must already have been opened using <b>open()</b> or <b>creat()</b> . The <i>fd</i> parameter is the file descriptor returned by <b>open()</b> or <b>creat()</b> .
	The <i>pStat</i> parameter is a pointer to a <b>stat</b> structure (defined in <b>stat.h</b> ). This structure must be allocated before <b>fstat()</b> is called.
	Upon return, the fields in the <b>stat</b> structure are updated to reflect the characteristics of the file.
RETURNS	OK or ERROR, the result of the ioctl() command to the filesystem driver.
ERRNO	EBADF Bad file descriptor number.
----------	---
	S_ioLib_UNKNOWN_REQUEST (ENOSYS) Device driver does not support the ioctl command.
	Other Other errors reported by device driver.
SEE ALSO	dirLib, stat( ), ls( )

## fstatfs()

NAME	<b>fstatfs()</b> – get file status information (POSIX)
SYNOPSIS	<pre>STATUS fstatfs   (     int fd, /* file descriptor for file to check */     struct statfs *pStat /* pointer to statfs structure */   )</pre>
DESCRIPTION	This routine obtains various characteristics of a file system. A file in the file system must already have been opened using <b>open()</b> or <b>creat()</b> . The <i>fd</i> parameter is the file descriptor returned by <b>open()</b> or <b>creat()</b> .
	The <i>pStat</i> parameter is a pointer to a <b>statfs</b> structure (defined in <b>stat.h</b> ). This structure must be allocated before <b>fstat()</b> is called.
	Upon return, the fields in the <b>statfs</b> structure are updated to reflect the characteristics of the file system. Note that for DosFS, the fields <b>f_files</b> and <b>f_ffree</b> are meaningless and are set to -1.
RETURNS	<b>OK</b> or <b>ERROR</b> , from the <b>ioctl()</b> command.
ERRNO	EBADF Bad file descriptor number.
	S_ioLib_UNKNOWN_REQUEST (ENOSYS) Device driver does not support the ioctl command.
	Other Other errors reported by device driver.
SEE ALSO	dirLib, statfs(), ls()

## fsync()

NAME	<b>fsync()</b> – synchronize a file
SYNOPSIS	<pre>int fsync (     int fd /* file descriptor of the file to sync */ )</pre>
DESCRIPTION	This function moves all modified data and attributes of the file descriptor $fd$ to a storage device. When <b>fsync()</b> returns, all in-memory modified copies of buffers associated with $fd$ have been written to the physical medium. It forces all outstanding data operations to synchronized file integrity completion.
RETURNS	Upon successful completion, <b>OK</b> , 0 is returned. Otherwise, <b>ERROR</b> , -1 returned and errno is set to indicate the error. If the <b>fsync()</b> function fails, outstanding I/O operations are not guaranteed to have been completed.
ERRNO	
SEE ALSO	fsPxLib, fdatasync()

## ftruncate()

NAME	ftruncate() – truncate a file (POSIX)
SYNOPSIS	<pre>int ftruncate   (    int fildes, /* fd of file to truncate */    off_t length /* length to truncate file */   )</pre>
DESCRIPTION	This routine truncates a file to a specified size.
RETURNS	0 (OK) or -1 (ERROR) if unable to truncate file.
ERRNO	EROFS File resides on a read-only file system. EBADF File is open for reading only.
	EINVAL File descriptor refers to a file on which this operation is impossible.

SEE ALSO ftruncate

## g0()

NAME	<b>g0( )</b> – return the contents of register g0 (also g1-g7) (SimSolaris)
SYNOPSIS	<pre>int g0   (    int taskId /* task ID, 0 means default task */   )</pre>
DESCRIPTION	This command extracts the contents of global register g0 from the TCB of a specified task. If <i>taskId</i> is omitted or 0, the current default task is assumed. Similar routines are provided for all global registers (g0 - g7): <b>g0()</b> - <b>g7()</b> .
RETURNS	The contents of register g0 (or the requested register).
ERRNO	Not Available
SEE ALSO	dbgArchLib, VxWorks Programmer's Guide: Debugging

## geiRegister()

NAME	geiRegister() – register with the VxBus subsystem
SYNOPSIS	void geiRegister(void)
DESCRIPTION	This routine registers the gei driver with VxBus as a child of the PCI bus type.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	gei825xxVxbEnd

## getOptServ()

NAME getOptServ() – parse parameter string into argc, argv format

SYNOPSIS	<pre>STATUS getOptServ   (     char * ParamString,     const char * progName, /* program name value for argv[0] */     int * argc,     char * argvloc[],     int argvlen   )</pre>
DESCRIPTION	none
RETURNS	OK if all arguments were successfully stored; otherwise, ERROR.
ERRNO	Not Available
SEE ALSO	getopt

## getenv()

NAME	<b>getenv()</b> – get an environment variable (ANSI)
SYNOPSIS	char *getenv ( FAST const char *name /* env variable to get value for */ )
DESCRIPTION	This routine searches the environment list (see the UNIX BSD 4.3 manual entry for <b>environ(5V)</b> ) for a string of the form "name=value" and returns the value portion of the string, if the string is present; otherwise it returns a <b>NULL</b> pointer.
RETURNS	A pointer to the string value, or a NULL pointer.
ERRNOS	N/A
SEE ALSO	<b>envLib, envLibInit(), putenv()</b> , UNIX BSD 4.3 manual entry , for <b>environ(5V)</b> , American National Standard for Information Systems -, Programming Language - C, ANSI X3.159-1989: General Utilities ( <b>stdlib.h</b> )

## getopt()

NAME	getopt() – parse argc/argv argument vector (POSIX)
SYNOPSIS	<pre>int getopt   (     int nargc,     char * const *nargv,     const char *ostr   )</pre>
DESCRIPTION	Decodes arguments passed in an argc/argv[] vector
	The parameters nargc and nargv are the argument count and argument array as passed to <b>main()</b> . The argument ostr is a string of recognized option characters; if a character is followed by a colon, the option takes an argument.
	The variable optind is the index of the next element of the nargv[] vector to be processed. It shall be initialized to 1 by the system, and <b>getopt()</b> shall update it when it finishes with each element of nargv[]. When an element of nargv[] contains multiple option characters, it is unspecified how <b>getopt()</b> determines which options have already been processed.
	The <b>getopt()</b> function shall return the next option character (if one is found) from nargv that matches a character in ostr, if there is one that matches. If the option takes an argument, <b>getopt()</b> shall set the variable optarg to point to the option-argument as follows:
	If the option was the last character in the string pointed to by an element of nargy, then optarg shall contain the next element of nargy, and optind shall be incremented by 2. If the resulting value of optind is greater than nargc, this indicates a missing option-argument, and <b>getopt()</b> shall return an error indication.
	Otherwise, optarg shall point to the string following the option character in that element of nargy, and optind shall be incremented by 1.
	If, when <b>getopt( )</b> is called:
	nargv[optind] is a null pointer nargv[optind] is not the character - nargv[optind] points to the string "-"
	getopt() shall return -1 without changing optind. If:
	nargv[optind] points to the string ""
	getopt() shall return -1 after incrementing optind.
	If <b>getopt()</b> encounters an option character that is not contained in ostr, it shall return the question-mark (?) character. If it detects a missing option-argument, it shall return the colon character (:) if the first character of ostr was a colon, or a question-mark character (?) otherwise. In either case, <b>getopt()</b> shall set the variable optopt to the option character that caused the error. If the application has not set the variable opterr to 0 and the first

	character of ostr is not a colon, <b>getopt()</b> shall also print a diagnostic message to stderr in the format specified for the getopts utility.
	The <b>getopt()</b> function need not be reentrant. A function that is not required to be reentrant is not required to be thread-safe.
RETURNS	The <b>getopt()</b> function shall return the next option character specified on the command line.
	A colon (:) shall be returned if <b>getopt()</b> detects a missing argument and the first character of ostr was a colon (:).
	A question mark (?) shall be returned if <b>getopt()</b> encounters an option character not in ostr or detects a missing argument and the first character of ostr was not a colon (:).
	Otherwise, <b>getopt()</b> shall return -1 when all command line options are parsed.
ERRNO	Not Available
SEE ALSO	getopt, POSIX

# getoptInit()

NAME	<b>getoptInit( )</b> – initialize the getopt state structure
SYNOPSIS	<pre>void getoptInit   (    GETOPT_ID pArg /* Pointer to getopt structure to be initialized */ )</pre>
DESCRIPTION	This function initializes the structure, <i>pGetOpt</i> that is used to maintain the getopt state. This structure is passed to <b>getopt_r()</b> which is a reentrant threadsafe version of the standard <b>getopt()</b> call. This function must be called before calling <b>getopt_r()</b>
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	getopt

## getopt\_r()

NAME	<b>getopt_r()</b> – parse argc/argv argument vector (POSIX)
SYNOPSIS	<pre>int getopt_r   (     int nargc,     char * const *nargv,     const char *ostr,     GETOPT_ID pGetOpt  )</pre>
DESCRIPTION	This function is a reentrant version of the <b>getopt()</b> function. The non-reentrant version keeps the getopt state in global variables across multiple calls made by the application, while this reentrant version keeps the state in the structure provided by the caller, thus allowing multiple callers to use getopt simultaneously without requiring any synchronization between them.
	The parameters <i>nargc</i> and <i>nargv</i> are the argument count and argument array as passed to <b>main()</b> . The argument <i>ostr</i> is a string of recognized option characters; if a character is followed by a colon, the option takes an argument. The argument <i>pGetOpt</i> points to the structure allocated by the caller to keep track of the getopt state. Prior to calling getopt_r, it is the caller responsibility to initialize this structure by calling <b>getoptInit()</b> .
	The variable pGetOpt->optind is the index of the next element of the nargv[] vector to be processed. <b>getopt_r()</b> shall update it when it finishes with each element of nargv[]. When an element of nargv[] contains multiple option characters, it is unspecified how <b>getopt_r()</b> determines which options have already been processed.
	The <b>getopt_r()</b> function shall return the next option character (if one is found) from nargy that matches a character in ostr, if there is one that matches. If the option takes an argument, <b>getopt_r()</b> shall set the variable pGetOpt->optarg to point to the option-argument as follows:
	If the option was the last character in the string pointed to by an element of nargv, then pGetOpt->optarg shall contain the next element of nargv, and pGetOpt->optind shall be incremented by 2. If the resulting value of pGetOpt->optind is greater than nargc, this indicates a missing option-argument, and <b>getopt_r(</b> ) shall return an error indication.
	Otherwise, pGetOpt->optarg shall point to the string following the option character in that element of nargy, and pGetOpt->optind shall be incremented by 1.
	If, when <b>getopt_r( )</b> is called:
	nargv[pGetOpt->optind] is a null pointer nargv[pGetOpt->optind] is not the character - nargv[pGetOpt->optind] points to the string "-"
	getopt_r() shall return -1 without changing pGetOpt->optind. If:
	nargv[pGetOpt->optind] points to the string ""

	getopt_r() shall return -1 after incrementing pGetOpt->optind.
	If <b>getopt_r()</b> encounters an option character that is not contained in ostr, it shall return the question-mark (?) character. If it detects a missing option-argument, it shall return the colon character (:) if the first character of ostr was a colon, or a question-mark character (?) otherwise. In either case, <b>getopt_r()</b> shall set the variable pGetOpt->optopt to the option character that caused the error. If the application has not set the variable pGetOpt->opterr to 0 and the first character of ostr is not a colon, <b>getopt_r()</b> shall also print a diagnostic message to stderr in the format specified for the getopts utility.
	This function is reentrant and thread-safe.
RETURNS	The <b>getopt_r()</b> function shall return the next option character specified on the command line.
	A colon (:) shall be returned if <b>getopt_r()</b> detects a missing argument and the first character of ostr was a colon (:).
	A question mark (?) shall be returned if <b>getopt_r()</b> encounters an option character not in ostr or detects a missing argument and the first character of ostr was not a colon (:).
	Otherwise, <b>getopt_r()</b> shall return -1 when all command line options are parsed.
ERRNO	Not Available
SEE ALSO	getopt, POSIX

## **h( )**

NAME	<b>h( )</b> – display or set the size of shell history
SYNOPSIS	<pre>void h   (    int size /* 0 = display, &gt;0 = set history to new size */ )</pre>
DESCRIPTION	This command displays or sets the size of VxWorks shell history. If no argument is specified, shell history is displayed. If <i>size</i> is specified, that number of the most recent commands is saved for display. The value of <i>size</i> is initially 20.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	usrLib, shellHistory(), ledLib, the VxWorks programmer guides.

```
hashFuncIterScale() - iterative scaling hashing function for strings
NAME
SYNOPSIS
                int hashFuncIterScale
                    (
                           elements, /* number of elements in hash table */
                    int
                   H_NODE_STRING *pHNode, /* pointer to string keyed hash node */
                    int seed /* seed to be used as scalar */
                    )
DESCRIPTION
                This hashing function interprets the key as a pointer to a null terminated string. A seed of
                13 or 27 appears to work well. It calculates the hash as follows:
                 for (tkey = pHNode->string; *tkey != '\0'; tkey++)
                hash = hash * seed + (unsigned int) *tkey;
                hash &= (elements - 1);
RETURNS
                integer between 0 and (elements - 1)
                N/A
ERRNO
                hashLib
SEE ALSO
```

#### hashFuncModulo()

hashFuncIterScale()

NAME	hashFuncModulo() – hashing function using remainder technique	
SYNOPSIS	<pre>int hashFuncModulo   (     int elements, /* number of elements in hash table */     H_NODE_INT *pHNode, /* pointer to integer keyed hash node */     int divisor /* divisor */ )</pre>	
DESCRIPTION	This hashing function interprets the key as a 32 bit quantity and applies the standard hashing function: $h(k) = K \mod D$ , where D is the passed divisor. The result of the hash function is masked to the appropriate number of bits to ensure the hash is not greater than (elements - 1).	
RETURNS	integer between 0 and (elements - 1)	

VxWorks Kernel API Reference, 6.6 hashFuncMultiply()

ERRNO N/A

SEE ALSO hashLib

## hashFuncMultiply()

NAME	<pre>hashFuncMultiply() – multiplicative hashing function</pre>	
SYNOPSIS	<pre>int hashFuncMultiply   (    int elements, /* number of elements in hash table */    H_NODE_INT *pHNode, /* pointer to integer keyed hash node */    int multiplier /* multiplier */   )</pre>	
DESCRIPTION	This hashing function interprets the key as a unsigned integer quantity and applies the standard hashing function: $h(k) = \text{leading N}$ bits of (B * K), where N is the appropriate number of bits such that the hash is not greater than (elements - 1). The overflow of B * K is discarded. The value of B is passed as an argument. The choice of B is similar to that of the seed to a linear congruential random number generator. Namely, B's value should take on a large number (roughly 9 digits base 10) and end inx21 where x is an even number. (Don't ask it involves statistics mambo jumbo)	
RETURNS	integer between 0 and (elements - 1)	
ERRNO	N/A	
SEE ALSO	hashLib	

## hashKeyCmp()

NAME	hashKeyCmp() – compare keys a	s 32 bit identifiers
SYNOPSIS	BOOL hashKeyCmp ( H_NODE_INT *pMatchHNode, H_NODE_INT *pHNode, int keyCmpArg )	<pre>/* hash node to match */ /* hash node in table to compare to */ /* argument ignored */</pre>

DESCRIPTION	This routine compares hash node keys as 32 bit identifiers. The argument keyCmpArg is unneeded by this comparator.
RETURNS	TRUE if keys match or, FALSE if keys do not match.
ERRNO	N/A
SEE ALSO	hashLib

## hashKeyStrCmp()

NAME	hashKeyStrCmp() – compare keys based on strings they point to	
SYNOPSIS	<pre>BOOL hashKeyStrCmp ( H_NODE_STRING *pMatchHNode, /* hash node to match */ H_NODE_STRING *pHNode, /* hash node in table to compare to */ int keyCmpArg /* argument ignored */ )</pre>	
DESCRIPTION	This routine compares keys based on the strings they point to. The strings must be null terminated. The routine <b>strcmp()</b> is used to compare keys. The argument keyCmpArg is unneeded by this comparator.	
RETURNS	TRUE if keys match or, FALSE if keys do not match.	
ERRNO	N/A	
SEE ALSO	hashLib	

### hashTblCreate()

NAME hashTblCreate() – create a hash table

SYNOPSIS	HASH_ID hashTblCreate	
	(	
	int sizeLog2,	/* number of elements in hash table log 2 */
	FUNCPTR keyCmpRtn,	/* function to test keys for equivalence */
	FUNCPTR keyRtn,	/* hashing function to generate hash from key */
	int keyArg	/* argument to hashing function */
	)	

VxWorks Kernel API Reference, 6.6 hashTblDelete()

DESCRIPTION	This routine creates a hash table 2^sizeLog2 number of elements. The hash table is carved from the caller's heap via malloc (2). To accommodate the list structures associated with the table, the actual amount of memory allocated will roughly eight times the number of elements requested. Additionally, two routines must be specified to dictate the behavior of the hashing table. The first routine, keyCmpRtn, is the key comparator function and the second routine, keyRtn, is the hashing function.
	The hashing function's role is to disperse the hash nodes added to the table as evenly throughout the table as possible. The hashing function receives as its parameters the number of elements in the table, a pointer to the <b>HASH_NODE</b> structure, and finally the keyArg parameter passed to this routine. The keyArg may be used to seed the hashing function. The hash function returns an index between 0 and (elements - 1). Standard hashing functions are available in this library.
	The keyCmpRtn parameter specifies the other function required by the hash table. This routine tests for equivalence of two HASH_NODES. It returns a boolean, TRUE if the keys match, and FALSE if they differ. As an example, a hash node may contain a HASH_NODE followed by a key which is an unsigned integer identifiers, or a pointer to a string, depending on the application. Standard hash node comparators are available in this library.
RETURNS	HASH_ID, or NULL if hash table could not be created.
ERRNO	Possible errnos generated by this routine include:
	<b>S_memLib_NOT_ENOUGH_MEMORY</b> There is not enough memory large enough to satisfy the allocation request.
SEE ALSO	hashLib, hashFuncIterScale( ), hashFuncModulo( ), hashFuncMultiply( ), hashKeyCmp( ), hashKeyStrCmp( )

## hashTblDelete()

NAME	hashTblDelete( ) – delete a hash table
SYNOPSIS	STATUS hashTblDelete ( HASH_ID hashId /* id of hash table to delete */ )
DESCRIPTION	This routine deletes the specified hash table and frees the associated memory. The hash table is marked as invalid.
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>hashId</i> is invalid.

SEE ALSO	hashLib, hashTblDestroy( ), hashTblTerminate( )	
	S_memLib_BLOCK_ERROR The block of memory to free is not valid.	
ERRNO	Possible errnos generated by this routine include:	

#### hashTblDestroy()

NAME	hashTblDestroy() – destroy a hash table	
SYNOPSIS	STATUS hashTblDestroy ( HASH_ID hashId, /* id of hash table to destroy */ BOOL dealloc /* deallocate associated memory */ )	
DESCRIPTION	This routine destroys the specified hash table and optionally frees the associated memory. The hash table is marked as invalid.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>hashId</i> is invalid.	
ERRNO	Possible errnos generated by this routine include: <b>S_memLib_BLOCK_ERROR</b> The block of memory to free is not valid.	
SEE ALSO	hashLib, hashTblDelete(), hashTblTerminate()	

## hashTblEach()

NAME	hashTblEach() – call a routine for each node in a hash table	
SYNOPSIS	HASH_NODE *hashTblEach ( HASH_ID hashId, /* hash table to call routine for */ FUNCPTR routine, /* the routine to call for each hash node */ int routineArg /* arbitrary user-supplied argument */ )	

**DESCRIPTION** This routine calls a user-supplied routine once for each node in the hash table. The routine should be declared as follows:

	BOOL routine (pNode, arg) HASH_NODE * pNode; /* pointer to a hash table node */ int arg; /* arbitrary user-supplied argument */	
	The user-supplied routine should return <b>TRUE</b> if <b>hashTblEach()</b> is to continue calling it with the remaining nodes, or <b>FALSE</b> if it is done and <b>hashTblEach()</b> can exit.	
RETURNS	NULL if traversed whole hash table, or pointer to HASH_NODE that hashTblEach endec with.	
ERRNO	N/A	
SEE ALSO	hashLib	

#### hashTblFind()

NAME	hashTblFind() – find a hash node that matches the specified key	
SYNOPSIS	HASH_NODE *hashTblFind ( FAST HASH_ID hashId, /* id of hash table from which to find node */ HASH_NODE *pMatchNode, /* pointer to hash node to match */ int keyCmpArg /* parameter to be passed to key comparator */ )	
DESCRIPTION	This routine finds the hash node that matches the specified key.	
RETURNS	pointer to HASH_NODE, or NULL if no matching hash node is found.	
ERRNO	N/A	
SEE ALSO	hashLib	

SEE ALSO hashLib

#### hashTblInit()

NAME	hashTblInit( ) – initialize a hash table	
SYNOPSIS	STATUS hashTblInit (	
	HASH_ID hashId, SL LIST *pTblMem,	/* id of hash table to initialize */ /* pointer to memory of sizeLog2 SL LISTs */

DESCRIPTIONThis routine initializes a hash table. Normally, creation and initialization of the hash tab should be done via the routine hashTblCreate(). However, if control over the memory allocation is necessary, this routine is used instead. All parameters are required with the exception of keyArg, which is optional. Refer to hashTblCreate() for a description of parameters.RETURNSOK, or ERROR if number of elements is negative, hashId is NULL, or the routines passed a NULL.ERRNON/ASEE ALSOhashLib, hashTblCreate()		<pre>int sizeLog2, /* number of elements in hash table log 2 */ FUNCPTR keyCmpRtn, /* function to test keys for equivalence */ FUNCPTR keyRtn, /* hashing function to generate hash from key */ int keyArg /* argument to hashing function */ )</pre>		
All parameters are required with the exception of keyArg, which is optional. Refer to hashTblCreate() for a description of parameters.         RETURNS       OK, or ERROR if number of elements is negative, hashId is NULL, or the routines passed a NULL.         ERRNO       N/A         SEE ALSO       hashLib, hashTblCreate()	DESCRIPTION	This routine initializes a hash table. Normally, creation and initialization of the hash table should be done via the routine <b>hashTblCreate()</b> . However, if control over the memory allocation is necessary, this routine is used instead.		
RETURNS       OK, or ERROR if number of elements is negative, hashId is NULL, or the routines passed a NULL.         ERRNO       N/A         SEE ALSO       hashLib, hashTblCreate()		All parameters are required with the exception of <b>keyArg</b> , which is optional. Refer to <b>hashTblCreate()</b> for a description of parameters.		
ERRNON/ASEE ALSOhashLib, hashTblCreate()	RETURNS	OK, or ERROR if number of elements is negative, hashId is NULL, or the routines passed are NULL.		
SEE ALSO hashLib, hashTblCreate()	ERRNO	N/A		
	SEE ALSO	hashLib, hashTblCreate()		

## hashTblPut()

NAME	hashTblPut() – put a hash node into the specified hash table	
SYNOPSIS	STATUS hashTblPut ( HASH_ID hashId, /* id of hash table in which to put node */ HASH_NODE *pHashNode /* pointer to hash node to put in hash table */ )	
DESCRIPTION	This routine puts the specified hash node in the specified hash table. Identical nodes will be kept in FIFO order in the hash table.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>hashId</i> is invalid.	
ERRNO	N/A	
SEE ALSO	hashLib, hashTblRemove()	

#### hashTblRemove()

NAME	hashTblRemove() – remove a hash node from a hash table	
SYNOPSIS	STATUS hashTblRemove ( HASH_ID hashId, /* id of hash table to to remove node from */ HASH_NODE *pHashNode /* pointer to hash node to remove */ )	
DESCRIPTION	This routine removes the hash node that matches the specified key.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>hashId</i> is invalid.	
ERRNO	N/A	
SEE ALSO	hashLib	

## hashTblTerminate()

NAME	hashTblTerminate() – terminate a hash table	
SYNOPSIS	STATUS hashTblTerminate ( HASH_ID hashId /* id of hash table to terminate */ )	
DESCRIPTION	This routine terminates the specified hash table. The memory for the table is not freed. The hash table is marked as invalid.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>hashId</i> is invalid.	
ERRNO	N/A	
SEE ALSO	hashLib, hashTblDestroy( ), hashTblDelete( )	

## help()

**NAME help()** – print a synopsis of selected routines

SYNOPSIS void help (void)

DESCRIPTION Thi

This command prints the following list of the calling sequences for commonly used routines, mostly contained in **usrLib**.

help		Print this list
dbgHelp		Print debug help info
edrHelp		Print ED&R help info
ioHelp		Print I/O utilities help info
nfsHelp		Print nfs help info
netHelp		Print network help info
rtpHelp		Print process help info
spyHelp		Print task histogrammer help info
timexHelp		Print execution timer help info
h	[n]	Print (or set) shell history
i	[task]	Summary of tasks' TCBs
ti	task	Complete info on TCB for task
sp	adr,args	Spawn a task, pri=100, opt=0x19, stk=20000
taskSpawn	name, pri, opt, stk	,adr,args Spawn a task
td -	task	Delete a task
ts	task	Suspend a task
tr	task	Resume a task
tw	task	Print pending task detailed info
W	[task]	Print pending task info
d	[adr[,nunits[,wid	lth]]] Display memory
m	adr[,width]	Modify memory
mRegs	[reg[,task]]	Modify a task's registers interactively
рс	[task]	Return task's program counter
version		Print VxWorks version info, and boot line
iam	"user"[,"passwd"]	] Set user name and passwd
whoami		Print user name
devs		List devices
ld	[syms[,noAbort][	,"name"]] Load std in into memory
		(syms = add symbols to table:
		-1 = none, 0 = globals, 1 = all)
lkup	["substr"]	List symbols in system symbol table
lkAddr	address	List symbol table entries near address
checkStack	[task]	List task stack sizes and usage
printErrno	value	Print the name of a status value
period	<pre>secs,adr,args</pre>	Spawn task to call function periodically
repeat	n,adr,args	Spawn task to call function n times (0=forever)
shConfig	["config"]	Display or set shell configuration variables
strFree	[address]	Free strings allocated within the shell
NOTE: Argur	nents specifying ·	<task> can be either task ID or name.</task>

RETURNS

N/A

VxWorks Kernel API Reference, 6.6 histLoad()

ERRNO N/A

**SEE ALSO** usrLib, ioHelp(), netHelp(), spyHelp(), the VxWorks programmer guides.

#### histLoad()

NAME	histLoad() – load history into the current shell session interpreter(s)	
SYNOPSIS	<pre>void histLoad   (     char * loadFile, /* file path to load the history from */     BOOL allInterp /* whether to save for all interpreters */ )</pre>	
DESCRIPTION	This command loads the shell history for the current shell session. If <i>allInterp</i> is set to <b>TRUE</b> , the load is done for all registered interpreters, otherwise only the history corresponding to the current interpreter is loaded.	
	The full path of the history file (including its name) is pointed by <i>loadFile</i> which must be <b>MAX_FILENAME_LENGTH</b> bytes long at most (including EOS). If <i>loadFile</i> is set to <b>NULL</b> , the system loads the history from a file named shellHistory.dat in the current directory.	
RETURNS	N/A	
ERRNO	Not Available	
SEE ALSO	usrShellHistLib, shellInterpLib, the VxWorks programmer guides.	

### histSave()

NAME SYNOPSIS	<b>histSave( )</b> – save history of the current shell session interpreter(s)	
	<pre>void histSave   (     char * saveFile, /* file path to save the history to */     BOOL allInterp /* whether to save for all interpreters */   )</pre>	
DESCRIPTION	This command saves the shell history for the current shell session. If <i>allInterp</i> is set to <b>TRUE</b> , the save is done for all registered interpreters, otherwise only the current interpreter history is saved.	

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	MAX_FILENAME_LENGTH bytes long at most (including EOS). If <i>saveFile</i> is set to NULL, the history is saved into a file named shellHistory.dat in the current directory.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	usrShellHistLib, shellInterpLib, the VxWorks programmer guides.

The full path of the save file (including its name) is pointed to by *saveFile* which must be

#### hookAddToHead()

hookAddToHead() – add a hook routine at the start of a hook table NAME SYNOPSIS STATUS hookAddToHead ( /\* routine to be added to table \*/ void \* hook, void \* table[], /\* table to which to add \*/ int maxEntries /\* max entries in table \*/ ) DESCRIPTION This routine adds a hook routine into a given hook table. The routine is added at the head (i.e. first entry) of the table. Existing hooks are shifted down to make way for the new hook. The last entry of the table is always NULL. Hooks are executed from the lowest to highest index of the table. Hence this routine should be used if hooks should be executed in LIFO order (i.e. last hook added executes first). Examples of LIFO hook execution are task delete hooks. This routine does not guard against duplicate entries. NOTE RETURNS OK, or ERROR if hook table is full. ERRNO S\_hookLib\_HOOK\_TABLE\_FULL hookLib SEE ALSO

#### hookAddToTail()

NAME	hookAddToTail() – add a hook routine to the end of a hook table	
SYNOPSIS	<pre>STATUS hookAddToTail   (    void * hook,</pre>	
DESCRIPTION	This routine adds a hook routine into a given hook table. The routine is added at the first <b>NULL</b> entry in the table. In other words new hooks are appended to the list of hooks already present.	
NOTE	This routine does not guard against duplicate entries.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if hook table is full.	
ERRNO	S_hookLib_HOOK_TABLE_FULL	
SEE ALSO	hookLib	

## hookDelete()

NAME	hookDelete() – delete a hook from a hook table	
SYNOPSIS	<pre>STATUS hookDelete  (     void * hook,</pre>	
DESCRIPTION	Deletes a previously added hook (if found) from a given hook table. Entries following the deleted hook are moved up to fill the vacant spot created.	
RETURNS	OK, or ERROR if hook could not be found.	
ERRNO	S_hookLib_HOOK_NOT_FOUND	
SEE ALSO	hookLib	

NAME	hookFind() – Search a hook table for a given hook	
SYNOPSIS	<pre>BOOL hookFind   (    void * hook,   /* routine to be deleted from table */    void * table[],   /* table from which to delete */    int maxEntries  /* max entries in table */   )</pre>	
DESCRIPTION	This function searches through a given hook table for a certain hook function. If found <b>TRUE</b> is returned, otherwise <b>FALSE</b> is returned.	
RETURNS	TRUE, or FALSE if the hook was not found.	
ERRNO	N/A.	
SEE ALSO	hookLib	

## hookShow()

NAME	<b>hookShow()</b> – show the hooks in the given hook table
SYNOPSIS	<pre>void hookShow   (    FUNCPTR table[], /* table from which to delete */    int maxEntries /* max entries in table */   )</pre>
DESCRIPTION	Shows the contents of a hook table symbolically.
RETURNS	N/A.
ERRNO	Not Available
SEE ALSO	hookShow

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#### hrfsAdvFormat()

NAME	hrfsAdvFormat() – format the HRFS file system using advanced options
SYNOPSIS	<pre>STATUS hrfsAdvFormat     (         char * path,</pre>
DESCRIPTION	This routine formats a disk or partition referenced by the path to the media.
	The <i>path</i> argument should be a valid path to the disk or partition to be formatted.
	The <i>diskSize</i> argument is used to specify how many bytes of the media the HRFS file system should occupy. It can be used to prevent HRFS for using the end portion of the media. In general this value should be 0 to specify that the entire media is to be used.
	The <i>blkSize</i> parameter is used to specify what block size, in byte, HRFS should use. This block size must be a power of 2, greater than the physical sector size, and be within 512 to 8196 bytes inclusively. In general this value should be specified as 0 so the formatter can determine the most efficient block size to use for the media size.
	The <i>numInodes</i> parameter is used to specify the absolute maximum number of files and directories the file system can ever have. Note this does not include the root directory which the formatter creates automatically. Specifying a value of 0 will tell the HRFS formatter to allow for the maximum number of files/directories the file system can have based on the amount of data blocks. I.e. One inode per data block.
	The <i>majorVersion</i> and <i>minorVersion</i> parameters are used to specify which particular version of the file system layout should be used when formatting the disk. A value of zero for both the major and minor version is used to indicate that the latest version of the file system should be used.
	The options parameter is used to specify additional formatting options. It is currently unused
RETURNS	OK on success or ERROR on failure.
ERRNO	Not Available
SEE ALSO	hrfsFormatLib

#### hrfsAdvFormatFd()

NAME	<b>hrfsAdvFormatFd( )</b> – format the HRFS file system using advanced options via a file descriptor	
SYNOPSIS	<pre>STATUS hrfsAdvFormatFd (     int fd,</pre>	
DESCRIPTION	This routine formats a disk or partition referenced by an open file on the media. The file is closed when formatting is complete.	
	The <i>fd</i> argument should be a valid file descriptor representing the root directory of the current file system. This file descriptor will be marked closed and invalid upon return from this function reguardless of outcome.	
	The <i>diskSize</i> argument is used to specify how many bytes of the media the HRFS file system should occupy. It can be used to prevent HRFS for using the end portion of the media. In general this value should be 0 to specify that the entire media is to be used.	
	The <i>blkSize</i> parameter is used to specify what block size, in byte, HRFS should use. This block size must be a power of 2, greater than the physical sector size, and be within 512 to 8196 bytes inclusively. In general this value should be specified as 0 so the formatter can determine the most efficient block size to use for the media size.	
	The <i>numInodes</i> parameter is used to specify the absolute maximum number of files and directories the file system can ever have. Note this does not include the root directory which the formatter creates automatically. Specifying a value of 0 will tell the HRFS formatter to allow for the maximum number of files/directories the file system can have based on the amount of data blocks. I.e. One inode per data block.	
	The <i>majorVersion</i> and <i>minorVersion</i> parameters are used to specify which particular version of the file system layout should be used when formatting the disk. A value of zero for both the major and minor version is used to indicate that the latest version of the file system should be used.	
	The options parameter is used to specify additional formatting options. It is currently unused	
RETURNS	OK on success or ERROR on failure	
ERRNO	Not Available	

VxWorks Kernel API Reference, 6.6 hrfsAscTime()

SEE ALSO hrfsFormatLib

#### hrfsAscTime()

NAME	<pre>hrfsAscTime() – convert "broken-down" HRFS time to string</pre>	
SYNOPSIS	<pre>int hrfsAscTime   (    HRFS_TM * pHrfsTm, /* Buffer contain time to convert in HRFS format */    char * pBuffer, /* Place to write ASCII time format data */    size_t bufLength /* Size of the supplied ASCII time buffer */   )</pre>	
DESCRIPTION	This routine converts the "broken-down" HRFS time pointed to by <i>pHrfsTm</i> into a string of the form: SUN SEP 16 01:03:52 1973\en\e0 The string is copied into <i>pBuffer</i> . Note that the field pHrfsTm->msec is not displayed.	
RETURNS	the number of bytes copied to <i>pBuffer</i> .	
ERRNO	Not Available	
SEE ALSO	hrFsTimeLib	

#### hrfsChkDsk()

NAME	hrfsChkDsk() – check the HRFS file system		
SYNOPSIS	<pre>STATUS hrfsChkDsk (     char * path, /* path to check */     int verbLevel, /* verbosity level */     int flags /* additional control information */ )</pre>		
DESCRIPTION	This routine is the HRFS consistency checker. It checks to see if the file system referenced by the path is stable and consistent.		
	WARNING. This function can only run on an inactive volume. Any currently opened files will be closed as this routine will eject the current		

	file system. The volume will also be unaccessible while the consistency checker executes.
	The <i>path</i> argument should be a valid path to the HRFS formatted disk or partition to be checked.
	The <i>verbLevel</i> argument is used to specify how much information is outputted to the console. A value of one indicates maximum verbosity. A value of zero indicates minimum verbosity.
	The <i>flags</i> parameter is used to specify additional control information to the consistency checker. If the <b>HRFS_CHKDSK_FLAG_UPGRADE</b> bit is set, the checker will attempt to upgrade the file system to the newest version. All other flags are ignored if this bit is set. If the <b>HRFS_CHKDSK_FLAG_REWIND_INODE_JOURNAL</b> bit is set, the checker will attempt to rewind the inode journal. That is, if the inode journal is not empty and is marked as being out of sync with the other disk structures, inodes contained in the inode journal are copied back into the inode table on disk. This has the effect of restoring the inodes to the previous transaction.
RETURNS	<b>OK</b> if media contains no errors or <b>ERROR</b> if one of more problems are detected.
ERRNO	Not Available
SEE ALSO	hrfsChkDskLib

#### hrfsDevCreate()

NAME	hrfsDevCreate() – create an HRFS	device
SYNOPSIS	<pre>HRFS_DEV_ID hrfsDevCreate   (     char * pDevName,     device_t xbdId,     int numBufs,     int maxFiles,     int defCommitPolicy,     int defCommitPeriod periodic) */   )</pre>	<pre>/* Name of the HRFS device (mount point). */ /* XBD for the device on which to mount. */ /* # of [struct buf] to allocate. */ /* Maximum # of simultaneously open files */ /* Initial commit policy */ /* Initial commit period (if policy is</pre>
DESCRIPTION	This routine creates an HRFS devic	e.
RETURNS	HRFS_DEV_ID if created and instal	led in Core I/O, NULL if not.

VxWorks Kernel API Reference, 6.6 hrfsDiskFormat()

ERRNO Not Available

SEE ALSO hrFsLib

#### hrfsDiskFormat()

NAME	hrfsDiskFormat() – format a disk with HRFS
SYNOPSIS	<pre>STATUS hrfsDiskFormat (     const char * pDevName, /* name of the device to initialize */     int files, /* the maximum number of files to support */     UINT32 majorVer, /* major version of fs to format */     UINT32 minorVer, /* minor version of fs to format */     UINT32 options /* formatter options */ )</pre>
DESCRIPTION	This command formats a disk and creates the HRFS file system on it. The device must already have been created by the device driver and HRFS format component must be included.
EXAMPLE	-> hrfsDiskFormat "/fd0", 0 /* format "/fd0" with HRFS */ /*allowing maximum files */ -> hrfsDiskFormat "/fd0", 100 /* format "/fd0" with HRFS */ /*allowing 100 files */
RETURNS	OK, or ERROR if the device cannot be opened or formatted.
ERRNO	Not Available
SEE ALSO	usrFsLib, hrFsLib, the VxWorks programmer guides.

## hrfsFormat()

NAME	hrfsFormat() – format the HRFS file system via a path
SYNOPSIS	STATUS hrfsFormat ( char * path, /* path to format */ UINT64 diskSize, /* size of disk in bytes */ UINT32 blkSize /* size of block in bytes */
	UINT32 blkSize, /* size of block in bytes */

	UINT32 numInodes /* number of Inodes */ )
DESCRIPTION	This routine formats a disk or partition referenced by the path to the media.
	The <i>path</i> argument should be a valid path to the disk or partition to be formatted.
	The <i>diskSize</i> argument is used to specify how many bytes of the media the HRFS file system should occupy. It can be used to prevent HRFS for using the end portion of the media. In general this value should be 0 to specify that the entire media is to be used.
	The <i>blkSize</i> parameter is used to specify what block size, in bytes, HRFS should use. This block size must be a power of 2, greater than the physical sector size, and be within 512 to 8196 bytes inclusively. In general this value should be specified as 0 so the formatter can determine the most efficient block size to use for the media size.
	The <i>numInodes</i> parameter is used to specify the absolute maximum number of files and directories the file system can ever have. Note this does not include the root directory which the formatter creates automatically. Specifying a value of 0 will tell the HRFS formatter to allow for the maximum number of files/directories the file system can have based on the amount of data blocks. I.e. One inode per data block.
RETURNS	OK on success or ERROR on failure.
ERRNO	Not Available
SEE ALSO	hrfsFormatLib

#### hrfsFormatFd()

this function reguardless of outcome.

 NAME
 hrfsFormatFd() – format the HRFS file system via a file descriptor

 SYNOPSIS
 STATUS hrfsFormatFd

 (
 int
 fd,
 /\* open file descriptor on disk \*/

 UINT64
 diskSize,
 /\* size of disk in bytes \*/

 UINT32
 blkSize,
 /\* size of block in bytes \*/

 UINT32
 numInodes
 /\* number of Inodes \*/

 DESCRIPTION
 This routine formats a disk or partition referenced by an open file on the media. The file is closed when formatting is complete.

 The fd argument should be a valid file descriptor representing the root directory of the current file system. This file descriptor will be marked closed and invalid upon return from

The *diskSize* argument is used to specify how many bytes of the media the HRFS file system should occupy. It can be used to prevent HRFS for using the end portion of the media. In general this value should be 0 to specify that the entire media is to be used.

The *blkSize* parameter is used to specify what block size, in bytes, HRFS should use. This block size must be a power of 2, greater than the physical sector size, and be within 512 to 8196 bytes inclusively. In general this value should be specified as 0 so the formatter can determine the most efficient block size to use for the media size.

The *numInodes* parameter is used to specify the absolute maximum number of files and directories the file system can ever have. Note this does not include the root directory which the formatter creates automatically. Specifying a value of 0 will tell the HRFS formatter to allow for the maximum number of files/directories the file system can have based on the amount of data blocks. I.e. One inode per data block.

RETURNS	OK on success or	<b>ERROR</b> on failure

ERRNO Not Available

SEE ALSO hrfsFormatLib

#### hrfsFormatLibInit()

NAME hrfsFormatLibInit() – prepare to use the HRFS formatter

SYNOPSIS STATUS hrfsFormatLibInit (void)

**DESCRIPTION** This routine initializes the HRFS formatter library. This initialization is enabled when the configuration macro INCLUDE\_HRFS\_FORMAT is defined.

RETURNS OK always

ERRNO Not Available

SEE ALSO hrfsFormatLib

### hrfsTimeCondense()

NAME	hrfsTimeCondense() – condense time in HRFS_TM to time in msec
SYNOPSIS	<pre>INT64 hrfsTimeCondense   (    HRFS_TM * pHrfsTm /* Pointer to where HRFS_TM format time is stored */   )</pre>
DESCRIPTION	This routine condenses the "broken-down" time pointed to by <i>pHrfsTm</i> into the number of milliseconds since midnight Jan 1, 1970.
RETURNS	# of milliseconds since midnight Jan 1, 1970
ERRNO	Not Available
SEE ALSO	hrFsTimeLib

## hrfsTimeGet()

NAME hrfsTimeGet() – return # of milliseconds since midnight Jan 1, 1970

- SYNOPSIS hrfsTime\_t hrfsTimeGet (void)
- **DESCRIPTION** This routine returns the number of milliseconds since midnight, January 1, 1970.
- **RETURNS** # of milliseconds since midnight, January 1, 1970

ERRNO Not Available

SEE ALSO hrFsTimeLib

#### hrfsTimeSplit()

NAME	<pre>hrfsTimeSplit( ) - split time in msec into HRFS_TM format</pre>						
SYNOPSIS	STATUS hrfs	FimeSplit					
	( INT64	milliSeconds,	/* milliseconds	to convert	to HRFS_TM	format */	

 HRFS\_TM \* pHrfsTm
 /\* Buffer to store HRFS\_TM format time \*/

 DESCRIPTION
 This routine splits the time specified in milliseconds into the "broken-down" format of the HRFS\_TM structure. Should the equivalent number of seconds exceed what can be represent by a signed integer, this routine will currently return ERROR, and the split will not have occurred.

 RETURNS
 OK success, or ERROR if the split did not occur

 ERRNO
 Not Available

 SEE ALSO
 hrFsTimeLib

#### hrfsUpgrade()

NAME	hrfsUpgrade() – upgrade the HRFS file system to the latest version
SYNOPSIS	STATUS hrfsUpgrade ( char *path /* path to upgrade */ )
DESCRIPTION	This routine is the HRFS consistency checker. It checks to see if the file system referenced by the path is stable and consistent.
	WARNING. This function can only run on an inactive volume. Any currently opened files will closed as this routine will eject the current file system. The volume will also be unaccessible while the consistency checker executes.
	The <i>path</i> argument should be a valid path to the HRFS formatted disk or partition to be checked.
RETURNS	<b>OK</b> if media was upgraded without errors. <b>ERROR</b> if one of more problems are detected.
ERRNO	Not Available
SEE ALSO	hrfsChkDskLib

#### i()

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tNbioLog

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tExcTask tJobTask

tLogTask

tShell0

186438

187390

logTask

188848

shellTask

SYNOPSIS       void i         (int taskNameOrId /* task name or task ID, 0 = summarize all */         DESCRIPTION       This command displays a synopsis of all the tasks in the system. The ti() routine provide more complete information on a specific task.         Both i() and ti() use taskShow(); see the documentation for taskShow() for a description of the output format.         EXAMPLE       -> i <ul> <li>NAME</li> <li>ENTRY</li> <li>TID</li> <li>PRI STATUS</li> <li>PC</li> <li>SP</li> <li>ERNO</li> <li>DELAY</li> <li>tExcTask</li> <li>excTask</li> <li>602adf00</li> <li>0</li> <li>PEND</li> <li>60164ad0</li> <li>602add08</li> <li>0</li> <li>0</li> </ul> tLogTask       logTask       60351ba8       1       READY       6015fe68       6034fde8       0       0	i() – print a summary of each task's TCB					
DESCRIPTION       This command displays a synopsis of all the tasks in the system. The ti() routine provide more complete information on a specific task.         Both i() and ti() use taskShow(); see the documentation for taskShow() for a description of the output format.         EXAMPLE       -> i         Image: NAME       ENTRY       TID       PRI       STATUS       PC       SP       ERRNO       DELAY         tExcTask       excTask       602adf00       0       PEND       60164ad0       602add08       0       0         tLogTask       logTask       602b56c8       0       PEND       60164ad0       602b54d0       0       0	<pre>void i   (    int taskNameOrId /* task name or task ID, 0 = summarize all */ )</pre>					
Both i() and ti() use taskShow(); see the documentation for taskShow() for a description of the output format.         EXAMPLE       -> i         NAME       ENTRY       TID       PRI       STATUS       PC       SP       ERRNO       DELAY         tExcTask       excTask       602adf00       0       PEND       60164ad0       602add08       0       0         tLogTask       logTask       602b56c8       0       PEND       60164ad0       602b54d0       0       0	This command displays a synopsis of all the tasks in the system. The <b>ti( )</b> routine provides more complete information on a specific task.					
EXAMPLE       -> i         NAME       ENTRY       TID       PRI       STATUS       PC       SP       ERRNO       DELAY         tExcTask       excTask       602adf00       0       PEND       60164ad0       602add08       0       0         tLogTask       logTask       602b56c8       0       PEND       60164ad0       602b54d0       0       0         tShell0       shellTask       60351ba8       1       READY       6015fe68       6034fde8       0       0	'n					
NAMEENTRYTIDPRISTATUSPCSPERRNODELAYtExcTaskexcTask602adf000PEND60164ad0602add0800tLogTasklogTask602b56c80PEND60164ad0602b54d000tShell0shellTask60351ba81READY6015fe686034fde800						
tExcTask         excTask         602adf00         0         PEND         60164ad0         602add08         0         0           tLogTask         logTask         602b56c8         0         PEND         60164ad0         602b54d0         0         0           tShell0         shellTask         60351ba8         1         READY         6015fe68         6034fde8         0         0						
tLexclaskexclask602a01000FEND60164ad0602add0800tLogTasklogTask602b56c80PEND60164ad0602b54d000tShell0shellTask60351ba81READY6015fe686034fde800						
tShellO shellTask 60351ba8 1 READY 6015fe68 6034fde8 0 0						
tWdbTask wdbTask 60338308 3 PEND 601579f4 60337ff0 0 0						
tNetTask netTask 602bf6a8 50 PEND 601579f4 602bf4e0 0 0 value = 0 = 0x0						
<b>CAVEAT</b> This command should be used only as a debugging aid, since the information is obsolete b the time it is displayed.	уy					
SMP CONSIDERATIONS This command displays a "CPU #" column instead of the "DELAY" column. The "CPU #" column provides information on the CPU a task is executing on, or "-" if a task is not running on a CPU.						
SMP EXAMPLE -> i						
NAME ENTRY TID PRI STATUS PC SP ERRNO CP # 	U					

247210

2a5830

2a8990

2ac220

2be530

0 PEND

0 PEND

0 PEND

0 PEND

1 READY

1e0198 249390

2a5770

2a8870

2ac110

2bc780

1e0198

1dd86c

1e0198

1e8ec8

0

0

0

0

0

# VxWorks Kernel API Reference, 6.6 *i0( )*

SEE ALSO	usrLib, ti(), taskShow(), the VxWorks programmer guides.						
ERRNO	N/A						
RETURNS	N/A						
	- value = 0 =	0x0					
	0 tIdleTask1 idleTaskEntr 250630 287 READY 1dfb28 2505b0						0
	- tIdleTask0	idleTaskEntr	24d010 28	7 READY	1dfb1c	24cf90	0
	miiBusMoni>	134794	29d010 25	4 DELAY	1e65a0	29cf80	0

## i0()

NAME	i0() – return the contents of register i0 (also i1-i7) (SimSolaris)		
SYNOPSIS	<pre>int i0   (    int taskId /* task ID, 0 means default task */ )</pre>		
DESCRIPTION	This command extracts the contents of in register i0 from the TCB of a specified task. If <i>taskId</i> is omitted or 0, the current default task is assumed.		
	Similar routines are provided for all in registers (i0 - i7): i0() - i7().		
	The frame pointer is accessed via i6.		
RETURNS	The contents of register i0 (or the requested register).		
ERRNO	Not Available		
SEE ALSO	dbgArchLib, VxWorks Programmer's Guide: Debugging		

## i8042vxbRegister()

#### NAME i8042vxbRegister() – register i8042vxb driver

SYNOPSIS void i8042vxbRegister(void)

DESCRIPTION	This routine registers the i8042vxb driver and device recognition data with the vxBus subsystem.
NOTE	This routine is called early during system initialization, and *MUST NOT* make calls to OS facilities such as memory allocation and I/O.
RETURNS	N/A
ERRNO	
SEE ALSO	vxbI8042Kbd

#### ichAtaBlkRW()

**NAME** ichAtaBlkRW() – read or write sectors to a ATA/IDE disk.

SYNOPSIS STATUS ichAtaBlkRW

(	
ATA_DEV	*pDev,
sector_t	startBlk,
UINT32	nBlks,
char	*pBuf,
int	direction
)	

- **DESCRIPTION** Read or write sectors to a ATA/IDE disk. *startBlk* is the start Block, *nBlks* is the number of blocks, *pBuf* is data buffer pointer and *direction* is the direction either to read or write. It should be **O\_WRONLY** for data write to drive or **O\_RDONLY** for read data from drive.
- **RETURNS** OK, ERROR if the command didn't succeed.
- ERRNO Not Available
- SEE ALSO vxbIntelIchStorage

#### ichAtaCmd()

**NAME** ichAtaCmd() – issue a RegisterFile command to ATA/ATAPI device.

SYNOPSIS STATUS ichAtaCmd

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```
(
int ctrl, /* Controller number. 0 or 1 */
int drive, /* Drive number. 0 or 1 */
int cmd, /* Command Register */
int arg0, /* argument0 */
int arg1, /* argument1 */
int arg2, /* argument2 */
int arg3, /* argument3 */
int arg4, /* argument4 */
int arg5 /* argument5 */
)
```

#### DESCRIPTION

This function executes ATA command to ATA/ATAPI devices specified by arguments *ctrl* and *drive. cmd* is command to be executed and other arguments *arg0* to *arg5* are interpreted for differently in each case depending on the *cmd* command. Some commands (like **ATA\_CMD\_SET\_FEATURE**) have sub commands the case in which *arg0* is interpreted as subcommand and *arg1* is subcommand specific.

In general these arguments *arg0* to *arg5* are interpreted as command registers of the device as mentioned below.

- arg0 Feature Register
- arg1 Sector count
- arg2 Sector number
- arg3 CylLo
- arg4 CylHi
- arg5 sdh Register

As these registers are interpreted for different purpose for each command, arguments are not named after registers.

The following commands are valid in this function and the validity of each argument for different commands. Each command is tabulated in the form

COMM	AND ARG0	ARG1	ARG2	ARG3	ARG4	ARG5
ATA_	CMD_INITP	2			<u>,</u>	
	0	0	0	0	0	0
ATA_	CMD_RECALIB	0	0	0	0	0
ATA_	PI_CMD_SRST 0	0	0	0	0	0
ATA_	CMD_EXECUTE_1 0	DEVICE_DIAGN 0	OSTIC 0	0	0	0
ATA_ or	CMD_SEEK cylinder LBA high	head LBA low	0	0	0	0

ATA_CMD_SET_FEATURE FR SC (SUBCOMMAND) (SubCommand Specific Valu	0 e)	0	0	0
ATA_CMD_SET_MULTI sectors per block 0	0	0	0	0
ATA_CMD_IDLE SC 0 (Timer Period)	0	0	0	0
ATA_CMD_STANDBY SC 0 (Timer Period)	0	0	0	0
ATA_CMD_STANDBY_IMMEDIATE 0 0	0	0	0	0
ATA_CMD_SLEEP 0 0	0	0	0	0
ATA_CMD_CHECK_POWER_MODE 0 0	0	0	0	0
ATA_CMD_IDLE_IMMEDIATE 0 0	0	0	0	0
ATA_CMD_SECURITY_DISABLE_PA ATA_ZERO ATA_ZERO	SSWORD ATA_ZERO	ATA_ZERO	ATA_ZERO	ATA_ZERO
ATA_CMD_SECURITY_ERASE_PREP. 0 0	ARE 0	0	0	0
ATA_CMD_SECURITY_ERASE_UNIT ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO	ATA_ZERO	ATA_ZERO
ATA_CMD_SECURITY_FREEZE_LOC	к 0	0	0	0
ATA_CMD_SECURITY_SET_PASSWO	RD 0	0	0	0
ATA_CMD_SECURITY_UNLOCK	0	0	0	0
ATA_CMD_SMART (not impleme FR SC (SUBCOMMAND) (SubCommand Specific Value)	nted) SN (SubCommand Specific Va	ATA_ZERO alue)	ATA_ZERO	ATA_ZERO
ATA_CMD_GET_MEDIA_STATUS 0 0	0	0	0	0
ATA_CMD_MEDIA_EJECT 0 0	0	0	0	0
ATA_CMD_MEDIA_LOCK	0	0	0	0
ATA_CMD_MEDIA_UNLOCK	0	0	0	0
ATA_CMD_CFA_ERASE_SECTORS	0	0	0	0
ATA_CMD_CFA_WRITE_SECTORS_W ATA_ZERO SC	ITHOUT_ERASE ATA_ZERO	ATA_ZERO	ATA_ZERO	ATA_ZERO
ATA_CMD_CFA_WRITE_SECTORS_W ATA_ZEROSC	ITHOUT_ERASE ATA ZERO	ATA ZERO	ATA ZERO	ATA ZERO

# VxWorks Kernel API Reference, 6.6 ichAtaCmd( )

ATA_CMD_CFA_TRA ATA_ZERO	NSLATE_SECTO ATA_ZERO	DR SN	cylLo	cylHi	DH
ATA_CMD_CFA_REQ ATA_ZERO	UEST_EXTENDI ATA_ZERO	ED_ERROR_CODE ATA_ZERO	E ATA_ZERO	ATA_ZERO	ATA_ZERO
ATA_CMD_SET_MAX FR (SUBCOMMAND	ATA_ZERO	ATA_ZERO	ATA_ZERO	ATA_ZERO	ATA_ZERO

The following are the subcommands valid for ATA\_CMD\_SET\_MAX and are tabulated as below

SUBCOMMAND(in ARG0) ARG1	ARG2	ARG3	ARG4	ARG5
ATA_SUB_SET_MAX_ADDR SC sector (SET_MAX_VOLATILE or SET_MAX_NON_VOLATILE	ESS no cylLo	cylHi	head + modebit	
ATA_SUB_SET_MAX_SET_ ATA_ZERO	PASS ATA_ZERO	ATA_ZERO	ATA_ZERO	ATA_ZERO
ATA_SUB_SET_MAX_LOCK ATA_ZERO	ATA_ZERO	ATA_ZERO	ATA_ZERO	ATA_ZERO
ATA_SUB_SET_MAX_UNLC ATA_ZERO	CK ATA_ZERO	ATA_ZERO	ATA_ZERO	ATA_ZERO
ATA_SUB_SET_MAX_FREE ATA_ZERO	ZE_LOCK ATA_ZERO	ATA_ZERO	ATA_ZERO	ATA_ZERO

In ATA\_CMD\_SET\_FEATURE subcommand only arg0 and arg1 are valid, all other are ATA\_ZERO.

SUBCOMMAND (ARG0)	ARG1
ATA_SUB_ENABLE_8BIT	ATA_ZERO
ATA_SUB_ENABLE_WCACHE	ATA_ZERO
ATA_SUB_SET_RWMODE	mode
ATA_SUB_ENB_ADV_POW_MNGMNT	(see page no 168 table 28 in atapi specs ) 0x90
ATA_SUB_ENB_POW_UP_STDBY	ATA_ZERO
ATA_SUB_POW_UP_STDBY_SPIN	ATA_ZERO
ATA_SUB_BOOTMETHOD	ATA_ZERO
ATA_SUB_ENA_CFA_POW_MOD1	ATA_ZERO
ATA_SUB_DISABLE_NOTIFY	ATA_ZERO
ATA_SUB_DISABLE_RETRY	ATA_ZERO
ATA_SUB_SET_LENGTH	ATA_ZERO
ATA_SUB_SET_CACHE	ATA_ZERO
ATA_SUB_DISABLE_LOOK	ATA_ZERO
ATA_SUB_ENA_INTR_RELEASE	ATA_ZERO
---------------------------	----------
ATA_SUB_ENA_SERV_INTR	ATA_ZERO
ATA_SUB_DISABLE_REVE	ATA_ZERO
ATA_SUB_DISABLE_ECC	ATA_ZERO
ATA_SUB_DISABLE_8BIT	ATA_ZERO
ATA_SUB_DISABLE_WCACHE	ATA_ZERO
ATA_SUB_DIS_ADV_POW_MNGMT	ATA_ZERO
ATA_SUB_DISB_POW_UP_STDBY	ATA_ZERO
ATA_SUB_ENABLE_ECC	ATA_ZERO
ATA_SUB_BOOTMETHOD_REPORT	ATA_ZERO
ATA_SUB_DIS_CFA_POW_MOD1	ATA_ZERO
ATA_SUB_ENABLE_NOTIFY	ATA_ZERO
ATA_SUB_ENABLE_RETRY	ATA_ZERO
ATA_SUB_ENABLE_LOOK	ATA_ZERO
ATA_SUB_SET_PREFETCH	ATA_ZERO
ATA_SUB_SET_4BYTES	ATA_ZERO
ATA_SUB_ENABLE_REVE	ATA_ZERO
ATA_SUB_DIS_INTR_RELEASE	ATA_ZERO
ATA_SUB_DIS_SERV_INTR	ATA_ZERO

- **RETURNS** OK, ERROR if the command didn't succeed.
- ERRNO Not Available

SEE ALSO vxbIntelIchStorage

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#### ichAtaConfig()

**DESCRIPTION** This routine configures an ATA hard disk. Parameters:

VxWorks Kernel API Reference, 6.6 ichAtaConfigInit()

	<i>drive</i> the drive number of the hard disk; 0 is <b>C:</b> and 1 is <b>D:</b> .
	<pre>devName the mount point for all partitions which are expected to be present on the disk, separated with commas, for example "/ata0,/ata1" or "C:,D:". Blanks are not allowed in this string.</pre>
RETURNS	OK or ERROR.
ERRNO	Not Available
SEE ALSO	<b>vxbIntelIchStorage</b> , <b>src/config/usrAta.c</b> , <i>VxWorks Programmer's Guide: I/O System</i> , <i>Local File Systems</i> , <i>Intel i386/i486/Pentium</i>

### ichAtaConfigInit( )

- NAME ichAtaConfigInit() initialize the hard disk driver
- SYNOPSIS void ichAtaConfigInit (void)
- **DESCRIPTION** This routine is called from **usrConfig.c** to initialize the hard drive.
- **RETURNS** Not Available
- ERRNO Not Available
- SEE ALSO vxbIntelIchStorage

## ichAtaCtrlReset()

NAME	ichAtaCtrlReset() – reset the specified ATA/IDE disk controller
SYNOPSIS	STATUS ichAtaCtrlReset ( int ctrl )
DESCRIPTION	This routine resets the ATA controller specified by ctrl. The device control register is written with SRST=1

**RETURNS** OK, ERROR if the command didn't succeed.

ERRNO Not Available

SEE ALSO vxbIntelIchStorage

#### ichAtaDevCreate()

NAME	ichAtaDevCreate() – create a device for a ATA/IDE disk
SYNOPSIS	<pre>BLK_DEV * ichAtaDevCreate   (     int ctrl, /* ATA controller number, 0 is the primary controller */     int drive, /* ATA drive number, 0 is the master drive */     UINT32 nBlocks, /* number of blocks on device, 0 = use entire disc */     UINT32 blkOffset /* offset BLK_DEV nBlocks from the start of the drive */     )</pre>
DESCRIPTION	This routine creates a device for a specified ATA/IDE or ATAPI CDROM disk.
	<i>ctrl</i> is a controller number for the ATA controller; the primary controller is 0. The maximum is specified via <b>ATA_MAX_CTRLS</b> .
	<i>drive</i> is the drive number for the ATA hard drive; the master drive is 0. The maximum is specified via <b>ATA_MAX_DRIVES</b> .
	The <i>nBlocks</i> parameter specifies the size of the device in blocks. If <i>nBlocks</i> is zero, the whole disk is used.
	The <i>blkOffset</i> parameter specifies an offset, in blocks, from the start of the device to be used when writing or reading the hard disk. This offset is added to the block numbers passed by the file system during disk accesses. (VxWorks file systems always use block numbers beginning at zero for the start of a device.)
RETURNS	A pointer to a block device structure ( <b>BLK_DEV</b> ) or <b>NULL</b> if memory cannot be allocated for the device structure.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorage, dosFsMkfs(), dosFsDevInit(), rawFsDevInit()

VxWorks Kernel API Reference, 6.6 ichAtaDevIdentify()

#### ichAtaDevIdentify()

NAME	<pre>ichAtaDevIdentify() - identify device</pre>
SYNOPSIS	STATUS ichAtaDevIdentify ( int ctrl, int dev )
DESCRIPTION	This routine checks whether the device is connected to the controller, if it is, this routine determines drive type. The routine set <b>type</b> field in the corresponding <b>ATA_DRIVE</b> structure. If device identification failed, the routine set <b>state</b> field in the corresponding <b>ATA_DRIVE</b> structure to <b>ATA_DEV_NONE</b> .
RETURNS	TRUE if a device present, FALSE otherwise
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorage

### ichAtaDmaRW()

ichAtaDmaRW() - read/write a number of sectors on the current track in DMA mode NAME SYNOPSIS STATUS ichAtaDmaRW ( int ctrl, int drive, UINT32 cylinder, UINT32 head, UINT32 sector, void \* buffer, UINT32 nSecs, int direction, sector\_t startBlk ) Read/write a number of sectors on the current track in DMA mode DESCRIPTION OK, ERROR if the command didn't succeed. RETURNS Not Available ERRNO

SEE ALSO vxbIntelIchStorage

#### ichAtaDmaToggle()

NAME	ichAtaDmaToggle() – turn on or off an individual controllers dma support
SYNOPSIS	void ichAtaDmaToggle ( int ctrl )
DESCRIPTION	This routine lets you toggle the DMA setting for an individual controller. The controller number is passed in as a parameter, and the current value is toggled.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the parameters are invalid.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorageShow

#### ichAtaDrv()

NAME	ichAtaDrv() – Initialize	the ATA driver
SYNOPSIS	<pre>STATUS ichAtaDrv   (     int ctrl,     int drives,     int vector,     int level,     int configType,     int semTimeout,     int wdgTimeout   )</pre>	<pre>/* controller no. 0,1 */ /* number of drives 1,2 */ /* interrupt vector */ /* interrupt level */ /* configuration type */ /* timeout seconds for sync semaphore */ /* timeout seconds for watch dog */</pre>

**DESCRIPTION** This routine initializes the ATA/ATAPI device driver, initializes IDE host controller and sets up interrupt vectors for requested controller. This function must be called once for each controller, before any access to drive on the controller, usually which is called by **usrRoot()** in **usrConfig.c**.

If it is called more than once for the same controller, it returns **OK** with a message display **Host controller already initialized**, and does nothing as already required initialization is done.

Additionally it identifies devices available on the controller and initializes depending on the type of the device (ATA or ATAPI). Initialization of device includes reading parameters of the device and configuring to the defaults.

**RETURNS** OK, or **ERROR** if initialization fails.

ERRNO Not Available

SEE ALSO vxbIntelIchStorage, ichAtaDevCreate()

#### ichAtaDumptest()

NAME	ichAtaDumptest() – a quick test of the dump functionality for ATA driver
SYNOPSIS	<pre>void ichAtaDumptest   (    device_t d,    sector_t sector,    UINT32 blocks,    char *data  )</pre>
DESCRIPTION	<i>device_t</i> device id of the device to dump to. This can be any XBD device. Could be the XBD of the disk device itself, or could be the xbd of a partition overlayed on the drive.
	<i>sector</i> sector offset to begin dump relative to start of xbd. <i>blocks</i> number of blocks to dump to device * <i>data</i> buffer that contains data to dump
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorageShow

## ichAtaInit()

NAME	ichAtaInit() – initialize ATA device.
SYNOPSIS	STATUS ichAtaInit ( int ctrl, int drive )
DESCRIPTION	This routine issues a soft reset command to ATA device for initialization.
RETURNS	OK, ERROR if the command didn't succeed.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorage

## ichAtaParamRead()

NAME	ichAtaParamRead() – Read drive parameters
SYNOPSIS	<pre>STATUS ichAtaParamRead   (    int ctrl,    int drive,    void *buffer,    int command )</pre>
DESCRIPTION	Read drive parameters.
RETURNS	OK, ERROR if the command didn't succeed.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorage

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VxWorks Kernel API Reference, 6.6 ichAtaPilnit()

#### ichAtaPiInit()

ichAtaPiInit() - init a ATAPI CD-ROM disk controller NAME SYNOPSIS STATUS ichAtaPiInit ( int ctrl, int drive ) This routine resets a ATAPI CD-ROM disk controller. DESCRIPTION RETURNS OK, ERROR if the command didn't succeed. Not Available ERRNO vxbIntelIchStorage SEE ALSO

#### ichAtaRW()

ichAtaRW() – read/write a data from/to required sector. NAME SYNOPSIS STATUS ichAtaRW ( int ctrl, int drive, UINT32 cylinder, UINT32 head, UINT32 sector, \* buffer, void UINT32 nSecs, direction, int sector\_t startBlk ) Read/write a number of sectors on the current track DESCRIPTION OK, ERROR if the command didn't succeed. RETURNS Not Available ERRNO vxbIntelIchStorage SEE ALSO

## ichAtaRawio()

NAME	ichAtaRawio() – do raw I/O access
SYNOPSIS	STATUS ichAtaRawio ( int ctrl, int drive, ATA_RAW *pAtaRaw )
DESCRIPTION	This routine is called to perform raw I/O access. <i>drive</i> is a drive number for the hard drive: it must be 0 or 1. The <i>pAtaRaw</i> is a pointer to the structure <b>ATA_RAW</b> which is defined in <b>ichAtaDrv.h</b> .
RETURNS	OK, or ERROR if the parameters are not valid.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorage

# ichAtaShow()

NAME	ichAtaShow() – show the ATA/IDE disk parameters
SYNOPSIS	STATUS ichAtaShow ( int ctrl, int drive )
DESCRIPTION	This routine shows the ATA/IDE disk parameters. Its first argument is a controller number, 0 or 1; the second argument is a drive number, 0 or 1.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the parameters are invalid.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorageShow

### ichAtaShowInit( )

NAME	ichAtaShowInit() – initialize the ATA/IDE disk driver show routine		
SYNOPSIS	STATUS ichAtaShowInit (void)		
DESCRIPTION	This routine links the ATA/IDE disk driver show routine into the VxWorks system. It is called automatically when this show facility is configured into VxWorks using either of the following methods:		
	- If you use the configuration header files, define INCLUDE_SHOW_ROUTINES in <b>config.h</b> .		
	- If you use the Tornado project facility, select INCLUDE_ATA_SHOW.		
RETURNS	N/A		
ERRNO	Not Available		
SEE ALSO	vxbIntelIchStorageShow		

## ichAtaStatusChk()

NAME	ichAtaStatusChk() – Check status of drive and compare to requested status.
SYNOPSIS	STATUS ichAtaStatusChk ( ATA_CTRL * pCtrl, UINT8 mask, UINT8 status )
DESCRIPTION	Wait until the drive is ready.
RETURNS	<b>OK</b> , <b>ERROR</b> if the drive status check times out.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorage

#### ichAtaXbdDevCreate()

ichAtaXbdDevCreate() - create an XBD device for a ATA/IDE disk NAME

SYNOPSIS device\_t ichAtaXbdDevCreate ( int ctrl, /\* ATA controller number, 0 is the primary controller \*/ drive, /\* ATA drive number, 0 is the master drive \*/ int nBlocks, /\* number of blocks on device, 0 = use entire UINT32 disc \*/ UINT32 blkOffset, /\* offset BLK\_DEV nBlocks from the start of the drive \*/ const char \* name /\* name of xbd device to create \*/ )

- DESCRIPTION Use the existing code to create a standard block dev device, then create an XBD device associated with the BLKDEV.
- a device identifier upon success, or NULLDEV otherwise RETURNS

ERRNO

NAME

SEE ALSO vxbIntelIchStorage

#### ichAtaXbdRawio()

ichAtaXbdRawio() - do raw I/O access SYNOPSIS STATUS ichAtaXbdRawio ( device\_t device, sector\_t sector, UINT32 numSecs, char \*data, int direction )

DESCRIPTION This routine is called to perform raw I/O access.

> *device* is the XBD device identifier for the drive sector starting sector for I/O operation *numSecs* number of sectors to read/write *data* pointer to data buffer *dir* read or write

The *pAtaRaw* is a pointer to the structure **ATA\_RAW** which is defined in **ichAtaDrv.h**.

VxWorks Kernel API Reference, 6.6 ichAtapiBytesPerSectorGet()

**RETURNS OK**, or **ERROR** if the parameters are not valid.

ERRNO Not Available

SEE ALSO vxbIntelIchStorage

## ichAtapiBytesPerSectorGet()

NAME	ichAtapiBytesPerSectorGet() – get the number of Bytes per sector.
SYNOPSIS	UINT16 ichAtapiBytesPerSectorGet ( int ctrl, int drive )
DESCRIPTION	This function will return the number of Bytes per sector. This function will return correct values for drives of ATA/ATAPI-4 or less as this field is retired for the drives compliant to ATA/ATAPI-5 or higher.
RETURNS	Bytes per sector.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorageShow

## ichAtapiBytesPerTrackGet()

NAME	<pre>ichAtapiBytesPerTrackGet( ) - get the number of Bytes per track.</pre>		
SYNOPSIS	UINT16 ichAtapiBytesPerTrackGet ( int ctrl, int drive )		
DESCRIPTION	This function will return the number of Bytes per track. This function will return correct values for drives of ATA/ATAPI-4 or less as this feild is retired for the drives compliant to ATA/ATAPI-5 or higher.		

RETURNS	Bytes per track.
	by teo per truck

ERRNO Not Available

SEE ALSO vxbIntelIchStorageShow

## ichAtapiCtrlMediumRemoval()

NAME	ichAtapiCtrlMediumRemoval() – Issues PREVENT/ALLOW MEDIUM REMOVAL packet command	
SYNOPSIS	STATUS ichAtapiCtrlMediumRemoval ( ATA_DEV * pAtapiDev, int arg0 )	
DESCRIPTION	This function issues a command to drive to PREVENT or ALLOW MEDIA removal. Argument <i>arg0</i> selects to LOCK_EJECT or UNLOCK_EJECT.	
	To lock media eject <i>arg0</i> should be <b>LOCK_EJECT</b> To unload media eject <i>arg0</i> should be <b>UNLOCK_EJECT</b>	
RETURN	OK or ERROR	
RETURNS	Not Available	
ERRNO	Not Available	
SEE ALSO	vxbIntelIchStorage	

# ichAtapiCurrentCylinderCountGet()

NAME	<b>ichAtapiCurrentCylinderCountGet()</b> – get logical number of cylinders in the drive.
SYNOPSIS	UINT16 ichAtapiCurrentCylinderCountGet ( int ctrl, int drive
	)

VxWorks Kernel API Reference, 6.6 ichAtapiCurrentHeadCountGet()

DESCRIPTION	This function will return the number of logical cylinders in the drive. This value represents the no of cylinders that can be addressed.
RETURNS	Cylinder count.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorageShow

## ichAtapiCurrentHeadCountGet()

NAME	<pre>ichAtapiCurrentHeadCountGet() - get the number of read/write heads in the drive.</pre>
SYNOPSIS	UINT8 ichAtapiCurrentHeadCountGet ( int ctrl, int drive )
DESCRIPTION	This function will return the number of heads in the drive from device structure.
RETURNS	Number of heads.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorageShow

## ichAtapiCurrentMDmaModeGet()

NAME	ichAtapiCurrentMDmaModeGet() – get the enabled Multi word DMA mode.		
SYNOPSIS	UINT8 ichAtapiCurrentMDmaModeGet ( int ctrl, int drive )		
DESCRIPTION	This function is used to get drive MDMA mode enable in the ATA/ATAPI drive specified by <i>ctrl</i> and <i>drive</i> from drive structure. The following bit is set for corresponding mode selected.		

SEE ALSO	vxb]	IntelIchStorageShow
ERRNO	Not	Available
RETURNS	Enal	bled Multi word DMA mode.
	-	Bit0 Multi DMA mode 0 is Selected
	-	Bit1 Multi DMA mode 1 is Selected
	-	Bit2 Multi DMA mode 2 is Selected

#### ichAtapiCurrentPioModeGet()

NAME	ichAtapiCurrentPioModeGet() – get the enabled PIO mode.
SYNOPSIS	UINT8 ichAtapiCurrentPioModeGet ( int ctrl, int drive )
DESCRIPTION	This function is used to get drive current PIO mode enabled in the ATA/ATAPI drive specified by <i>ctrl</i> and <i>drive</i> from drive structure.
RETURNS	Enabled PIO mode.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorageShow

#### ichAtapiCurrentRwModeGet()

**NAME** ichAtapiCurrentRwModeGet() – get the current Data transfer mode.

SYNOPSIS	UINT8 ichAtapiCurrentRwModeGet
	(
	int ctrl,
	int drive
	)

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VxWorks Kernel API Reference, 6.6 ichAtapiCurrentSDmaModeGet()

DESCRIPTION	This function will return the current Data transfer mode if it is PIO 0,1,2,3,4 mode, SDMA 0,1,2 mode, MDMA 0,1,2 mode or UDMA 0,1,2,3,4,5 mode.
RETURNS	current PIO mode.

ERRNO Not Available

SEE ALSO vxbIntelIchStorageShow

#### ichAtapiCurrentSDmaModeGet()

NAME	ichAtapiCurrentSDmaModeGet() – get the enabled Single word DMA mode.
SYNOPSIS	UINT8 ichAtapiCurrentSDmaModeGet ( int ctrl, int drive )
DESCRIPTION	This function is used to get drive SDMA mode enable in the ATA/ATAPI drive specified by $ctrl$ and $drive$ from drive structure
RETURNS	Enabled Single word DMA mode.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorageShow

## ichAtapiCurrentUDmaModeGet()

NAME ichAtapiCurrentUDmaModeGet() - get the enabled Ultra DMA mode.
SYNOPSIS UINT8 ichAtapiCurrentUDmaModeGet
(
int ctrl,
int drive
)

This function is used to get drive UDMA mode enable in the ATA/ATAPI drive specified

ing bit is set for corresponding mode	

- by ctrl and drive from drive structure The following bit is set for corresponding mode selected. Bit4 Ultra DMA mode 4 is Selected \_
  - Bit3 Ultra DMA mode 3 is Selected \_
  - Bit2 Ultra DMA mode 2 is Selected \_
  - Bit1 Ultra DMA mode 1 is Selected \_
  - Bit0 Ultra DMA mode 0 is Selected \_
- Enabled Ultra DMA mode. RETURNS
- Not Available ERRNO

DESCRIPTION

vxbIntelIchStorageShow SEE ALSO

#### ichAtapiCylinderCountGet()

NAME	<pre>ichAtapiCylinderCountGet() - get the number of cylinders in the drive.</pre>
SYNOPSIS	UINT16 ichAtapiCylinderCountGet ( int ctrl, int drive )
DESCRIPTION	This function is used to get cyclinder count of the ATA/ATAPI drive specified by <i>ctrl</i> and <i>drive</i> from drive structure.
RETURNS	Cylinder count.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorageShow

# ichAtapiDriveSerialNumberGet()

NAME	ichAtapiDriveSerialNumberGet() – get the drive serial number.
SYNOPSIS	<pre>char * ichAtapiDriveSerialNumberGet   (    int ctrl,    int drive   )</pre>
DESCRIPTION	This function is used to get drive serial number of the ATA/ATAPI drive specified by <i>ctrl</i> and <i>drive</i> from drive structure. It returns a pointer to character array of 20 bytes length which contains serial number in ascii.
RETURNS	Drive serial number.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorageShow

## ichAtapiDriveTypeGet()

NAME	ichAtapiDriveTypeGet() – get the drive type.	
SYNOPSIS	UINT8 ichAtapiDriveTypeGet ( int ctrl, int drive )	
DESCRIPTION	This function rou following table in 0x00h 0x01h 0x02h 0x02h 0x03h 0x04h 0x05h 0x06h 0x07h 0x08b	tine will return the type of the drive if it is CD-ROM or Printer etc. The adicates the type depending on the return value. Direct-access device Sequential-access device Printer Device Processor device Write-once device CD-ROM device Scanner device Optical memory device
	0x09h	Communications device

SEE ALSO	vxbIntelIchStora	geShow
ERRNO	Not Available	
RETURNS	drive type.	
	0x0Fh 0x1Fh	Optical Card Reader/Writer Device Unknown or no device type
	0x0Eh	Reduced Block Command Devices
	0x0Dh	Encloser Services Device
	0x0Ch	Array Controller Device

ichAtapiFeatureEnabledGet()

NAME	ichAtapiFeatureEnabledGet() – get the enabled features.	
SYNOPSIS	UINT32 ichAtapi ( int ctrl, int drive )	FeatureEnabledGet
DESCRIPTION	This function is used to get drive Features Enabled by the ATA/ATAPI drive specified by <i>ctrl</i> and <i>drive</i> from drive structure. It returns a 32 bit value whose bits represents the feature Enabled. The following table gives the cross reference for the bits.	
	Bit 21	Power-up in Standby Feature
	Bit 20	Removable Media Status
		Notification Feature
	Bit 19	Adavanced Power Management
		Feature
	Bit 18	CFA Feature
	Bit 10	Host protected Area Feature
	Bit 4	Packet Command Feature
	Bit 3	Power Management Feature
	Bit 2	Removable Media Feature
	Bit 1	Security Mode Feature
	Bit 0	SMART Feature
RETURNS	enabled features.	
ERRNO	Not Available	

VxWorks Kernel API Reference, 6.6 ichAtapiFeatureSupportedGet()

SEE ALSO vxbIntelIchStorageShow

	ichAtap	iFeatureSupportedGet( )	
NAME	ichAtapiFeatu	<b>reSupportedGet( )</b> – get the features supported by the drive.	
SYNOPSIS	UINT32 ichAtapiFeatureSupportedGet ( int ctrl, int drive )		
DESCRIPTION	This function is <i>ctrl</i> and <i>drive</i> from supported. The	This function is used to get drive Feature supported by the ATA/ATAPI drive specified by <i>ctrl</i> and <i>drive</i> from drive structure. It returns a 32 bit value whose bits represents the features supported. The following table gives the cross reference for the bits.	
	Bit 21 Bit 20 Bit 19 Bit 18 Bit 10 Bit 4 Bit 3 Bit 2 Bit 1 Bit 0	Power-up in Standby Feature Removable Media Status Notification Feature Adavanced Power Management Feature CFA Feature Host protected Area Feature Packet Command Feature Power Management Feature Removable Media Feature Security Mode Feature SMART Feature	
RETURNS	Supported feat	rures.	
ERRNO	Not Available		
SEE ALSO	vxbIntelIchSto	orageShow	

# ichAtapiFirmwareRevisionGet()

NAME	<b>ichAtapiFirmwareRevisionGet()</b> – get the firm ware revision of the drive.
SYNOPSIS	char * ichAtapiFirmwareRevisionGet (
	int ctrl,

	int drive )
DESCRIPTION	This function is used to get drive Firmware revision of the ATA/ATAPI drive specified by <i>ctrl</i> and <i>drive</i> from drive structure. It returns a pointer to character array of 8 bytes length which contains serial number in ascii.
RETURNS	firmware revision.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorageShow

## ichAtapiHeadCountGet( )

NAME	ichAtapiHeadCountGet() – get the number heads in the drive.
SYNOPSIS	UINT8 ichAtapiHeadCountGet ( int ctrl, int drive )
DESCRIPTION	This function is used to get head count of the ATA/ATAPI drive specified by <i>ctrl</i> and <i>drive</i> from drive structure.
RETURNS	Number of heads in the drive.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorageShow

## ichAtapiInit()

 NAME
 ichAtapiInit() – init ATAPI CD-ROM disk controller

 SYNOPSIS
 STATUS ichAtapiInit

( int ctrl, VxWorks Kernel API Reference, 6.6 ichAtapiloctl()

int drive

**DESCRIPTION** This routine resets the ATAPI CD-ROM disk controller.

**RETURNS** OK, ERROR if the command didn't succeed.

ERRNO Not Available

SEE ALSO vxbIntelIchStorage

#### ichAtapiIoctl()

NAME ichAtapiIoctl() – Control the drive.

SYNOPSIS STATUS ichAtapiIoctl

(					
int			function,	/*	The IO operation to do */
int			ctrl,	/*	Controller number of the drive */
int			drive,	/*	Drive number */
int		password	[16],	/*	Password to set. NULL if not
applicable	e */				
int			arg0,	/*	1st arg to pass. NULL if not
applicable	> */				
UINT32	2 *		arg1,	/*	Ptr to 2nd arg. NULL if not
applicable	e */				
UINT8	* *		ppBuf	/*	The data buffer */
)					

**DESCRIPTION** This routine is used to control the drive like setting the password, putting in power save mode, locking/unlocking the drive, ejecting the medium etc. The argument *function* defines the ioctl command, *password*, and integer array is the password required or set password value for some commands. Arguments *arg0*, pointer *arg1*, pointer to pointer *buffer ppBuf* are commad specific.

The following commands are supported for various functionality.

#### IOCTL\_DIS\_MASTER\_PWD

Disable the master password. where 4th parameter is the master password.

#### IOCTL\_DIS\_USER\_PWD

Disable the user password.

#### IOCTL\_ERASE\_PREPARE

Prepare the drive for erase incase the user password lost, and it is in max security mode.

**IOCTL\_ENH\_ERASE\_UNIT\_USR** Erase in enhanced mode supplying the user password.

IOCTL\_ENH\_ERASE\_UNIT\_MSTR

Erase in enhanced mode supplying the master password.

IOCTL\_NORMAL\_ERASE\_UNIT\_MSTR Erase the drive in normal mode supplying the master password.

IOCTL\_NORMAL\_ERASE\_UNIT\_USR Erase the drive in normal mode supplying the user password.

**IOCTL\_FREEZE\_LOCK** Freeze lock the drive.

**IOCTL\_SET\_PASS\_MSTR** Set the master password.

**IOCTL\_SET\_PASS\_USR\_MAX** Set the user password in Maximum security mode.

**IOCTL\_SET\_PASS\_USR\_HIGH** Set the user password in High security mode.

**IOCTL\_UNLOCK\_MSTR** Unlock the master password.

IOCTL\_UNLOCK\_USR Unlock the user password.

**IOCTL\_CHECK\_POWER\_MODE** Find the drive power saving mode.

IOCTL\_IDLE\_IMMEDIATE

Idle the drive immediatly. this will get the drive from the standby or active mode to idle mode immediatly.

IOCTL\_SLEEP

Set the drive in sleep mode. this is the highest power saving mode. to return to the normal active or IDLE mode, drive need an hardware reset or power on reset or device reset command.

IOCTL\_STANDBY\_IMMEDIATE

Standby the drive immediatly.

IOCTL\_EJECT\_DISK

Eject the media of an ATA drive. Use IOsystem ioctl function for ATAPI drive.

IOCTL\_GET\_MEDIA\_STATUS

Find the media status.

#### IOCTL\_ENA\_REMOVE\_NOTIFY

Enable the drive's removable media notification feature set.

#### VxWorks Kernel API Reference, 6.6 ichAtapiloctl( )

The following table describes these arguments validity. These are tabulated in the following form

FUNCTION password [16]	arg0	*arg1	**ppBuf
	_		
10CTL_DIS_MASTER_PWI password	d ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_DIS_USER_PWD password	ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_ERASE_PREPARE ATA_ZERO	ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_ENH_ERASE_UNI password	I_USR ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_ENH_ERASE_UNI password	I_MSTR ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_NORMAL_ERASE_U password	UNIT_MSTR ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_NORMAL_ERASE_U password	UNIT_USR ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_FREEZE_LOCK ATA_ZERO	ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_SET_PASS_MSTR password	ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_SET_PASS_USR_I password	MAX ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_SET_PASS_USR_1 password	HIGH ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_UNLOCK_MSTR password	ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_UNLOCK_USR password	ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_READ_NATIVE_M ATA_ZERO (A A	AX_ADDRESS - it TA_SDH_IBM or L TA_SDH_LBA ) ( L L L L L	returns address in BA/CHS add BA 27:24 / Head BA 23:16 / cylHi BA 15:8 / cylLow BA 7:0 / sector no	<arg1> ATA_ZERO )</arg1>
IOCTL_SET_MAX_ADDRES ATA_ZERO SET_MAX SET_MAX	SS – <ar _VOLATILE or L _NON_VOLATILE</ar 	gl> is pointer to LB. BA address	A address ATA_ZERO
IOCTL_SET_MAX_SET_PA password	ASS ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_SET_MAX_LOCK ATA_ZERO	ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_SET_MAX_UNLOCI ATA_ZERO	K ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_SET_MAX_FREEZ	E_LOCK ATA ZERO	ATA ZERO	ATA ZERO

IOCTL_CHECK_POWER_MODE - re ATA_ZERO ATA_ZERO	eturns power mode in returns power	<arg1> ATA_ZERO</arg1>
power modes :-1) 0x00 E 2) 0x80 E 3) 0xff Device in Ad	Device in standby mod Device in Idle mode ctive or Idle mode	le
IOCTL_IDLE_IMMEDIATE ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_SLEEP ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_STANDBY_IMMEDIATE ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_ENB_POW_UP_STDBY ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_ENB_SET_ADV_POW_MNGMNT ATA_ZERO arg0	ATA_ZERO	ATA_ZERO
NOTE:- arg0 value - 1). for minimum 2). for minimum 3). for maximum	power consumption wi power consumption wi performance 0xFEh	th standby 0x01h thout standby 0x01h
IOCTL_DISABLE_ADV_POW_MNGMNT ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_EJECT_DISK ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_LOAD_DISK ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_MEDIA_LOCK ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_MEDIA_UNLOCK ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_GET_MEDIA_STATUS - returns ATA_ZERO ATA_ZERO	s status in <arg1> status</arg1>	ATA_ZERO
NOTE: value in <arg1> is 0x04 -Command aborted 0x02 -No media in drive 0x08 -Media change is re 0x20 -Media changed 0x40 -Write Protected</arg1>	equested	
IOCTL_ENA_REMOVE_NOTIFY ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_DISABLE_REMOVE_NOTIFY ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_SMART_DISABLE_OPER ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_SMART_ENABLE_ATTRIB_AUTO ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_SMART_DISABLE_ATTRIB_AUTO ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_SMART_ENABLE_OPER ATA_ZERO ATA_ZERO	ATA_ZERO	ATA_ZERO
IOCTL_SMART_OFFLINE_IMMED ATA_ZERO SubCommand	ATA_ZERO	ATA_ZERO

(refer to ref1 page no 190) IOCTL\_SMART\_READ\_DATA - returns pointer to pointer <ppBuf> of read data ATA\_ZERO ATA\_ZERO ATA\_ZERO read data IOCTL\_SMART\_READ\_LOG\_SECTOR - returns pointer to pointer <ppBuf>of read data ATA\_ZERO no of sector to log Address read data be read IOCTL\_SMART\_RETURN\_STATUS ATA\_ZERO ATA\_ZERO ATA ZERO ATA\_ZERO IOCTL\_SMART\_SAVE\_ATTRIB ATA\_ZERO ATA\_ZERO ATA\_ZERO ATA\_ZERO IOCTL\_SMART\_WRITE\_LOG\_SECTOR ATA ZERO no of to be written Log Sector address write data NOTE: - <ppBuf> contains pointer to pointer data buffer to be written IOCTL\_CFA\_ERASE\_SECTORS ATA\_ZERO PackedCHS/LBA ATA\_ZERO sector count IOCTL\_CFA\_REQUEST\_EXTENDED\_ERROR\_CODE ATA\_ZERO ATA\_ZERO ATA\_ZERO ATA ZERO IOCTL\_CFA\_TRANSLATE\_SECTOR - <ppbuf> returns pointer to data pointer. ATA ZERO ATA ZERO PackedLBA/CHS read data IOCTL\_CFA\_WRITE\_MULTIPLE\_WITHOUT\_ERASE ATA ZERO sector count PackedCHS/LBA write data NOTE: -<pbuf> contains pointer to data pointer. IOCTL\_CFA\_WRITE\_SECTORS\_WITHOUT\_ERASE ATA\_ZERO sector count PackedCHS/LBA write data

#### RETURNS OK or ERROR

ERRNO Not Available

SEE ALSO vxbIntelIchStorage

#### ichAtapiMaxMDmaModeGet( )

DESCRIPTION	This fund drive spe correspo	ction is used to get drive maximum MDMA mode supported by the ATA/ATAPI ecified by <i>ctrl</i> and <i>drive</i> from drive structure The following bits are set for nding modes supported.
	Bit2	Multi DMA mode 2 and below are supported
	Bit0	Multi DMA mode 1 and below are supported Multi DMA mode 0 is supported
RETURNS	Maximu	m Multi word DMA mode.
ERRNO	Not Avai	ilable
SEE ALSO	vxbIntel	IchStorageShow

#### ichAtapiMaxPioModeGet()

NAME	ichAtapiMaxPioModeGet() – get the Maximum PIO mode that drive can support.
SYNOPSIS	UINT8 ichAtapiMaxPioModeGet ( int ctrl, int drive )
DESCRIPTION	This function is used to get drive maximum PIO mode supported by the ATA/ATAPI drive specified by <i>ctrl</i> and <i>drive</i> from drive structure
RETURNS	maximum PIO mode.
ERRNO	Not Available

SEE ALSO vxbIntelIchStorageShow

## ichAtapiMaxSDmaModeGet()

NAME ichAtapiMaxSDmaModeGet() – get the Maximum Single word DMA mode the drive supports

SYNOPSIS UINT8 ichAtapiMaxSDmaModeGet

VxWorks Kernel API Reference, 6.6 ichAtapiMaxUDmaModeGet()

( int ctrl, int drive )

- **DESCRIPTION** This function is used to get drive maximum SDMA mode supported by the ATA/ATAPI drive specified by *ctrl* and *drive* from drive structure
- **RETURNS** Maximum Single word DMA mode.

ERRNO Not Available

SEE ALSO vxbIntelIchStorageShow

## ichAtapiMaxUDmaModeGet()

NAME	<b>ichAtapiMaxUDmaModeGet( )</b> – get the Maximum Ultra DMA mode the drive can support.
SYNOPSIS	UINT8 ichAtapiMaxUDmaModeGet ( int ctrl, int drive )
DESCRIPTION	This function is used to get drive maximum UDMA mode supported by the ATA/ATAPI drive specified by <i>ctrl</i> and <i>drive</i> from drive structure. The following bits are set for corresponding modes supported.
	Bit4 Ultra DMA mode 4 and below are supported
	Bit3 Ultra DMA mode 3 and below are supported
	Bit2 Ultra DMA mode 2 and below are supported
	Bit1 Ultra DMA mode 1 and below are supported
	Bit0 Ultra DMA mode 0 is supported
RETURNS	Maximum Ultra DMA mode.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorageShow

#### ichAtapiModelNumberGet( ) ichAtapiModelNumberGet() – get the model number of the drive. NAME SYNOPSIS char \* ichAtapiModelNumberGet ( int ctrl, int drive ) This function is used to get drive Model Number of the ATA/ATAPI drive specified by ctrl DESCRIPTION and *drive* from drive structure. It returns a pointer to character array of 40 bytes length which contains serial number in ascii. pointer to the model number. RETURNS ERRNO Not Available

SEE ALSO vxbIntelIchStorageShow

#### ichAtapiPktCmd()

NAME	ichAtapiPktCmd() – execute an ATAPI command with error processing
SYNOPSIS	UINT8 ichAtapiPktCmd ( ATA_DEV * pAtapiDev, ATAPI_CMD * pComPack )
DESCRIPTION	This routine executes a single ATAPI command, checks the command completion status and tries to recover if an error encountered during command execution at any stage.
RETURN	SENSE_NO_SENSE if success, or ERROR if not successful for any reason.
RETURNS	Not Available
ERRNO	S_ioLib_DEVICE_ERROR
SEE ALSO	vxbIntelIchStorage

VxWorks Kernel API Reference, 6.6 ichAtapiPktCmdSend()

# ichAtapiPktCmdSend()

NAME	ichAtapiPktCmdSend() – Issue a Packet command.
SYNOPSIS	UINT8 ichAtapiPktCmdSend ( ATA_DEV * pAtapiDev, ATAPI_CMD * pComPack )
DESCRIPTION	This function issues a packet command to specified drive. See library file description for more details.
RETURN	SENSE_NO_SENSE if success, or ERROR if not successful for any reason
RETURNS	Not Available
ERRNO	S_ioLib_DEVICE_ERROR
SEE ALSO	vxbIntelIchStorage

# ichAtapiRead10()

NAME	ichAtapiRead10() – read one or more blocks from an ATAPI Device.
SYNOPSIS	<pre>STATUS ichAtapiRead10   (    ATA_DEV * pAtapiDev,    UINT32 startBlk,    UINT32 nBlks,    UINT32 transferLength,    char * pBuf   )</pre>
DESCRIPTION	This routine reads one or more blocks from the specified device, starting with the specified block number.
	The name of this routine relates to the SFF-8090i (Mt. Fuji), used for DVD-ROM, and indicates that the entire packet command uses 10 bytes, rather than the normal 12.
RETURNS	OK, ERROR if the read command didn't succeed.
ERRNO	Not Available

SEE ALSO vxbIntelIchStorage

## ichAtapiReadCapacity()

NAME	<b>ichAtapiReadCapacity()</b> – issue a READ CD-ROM CAPACITY command to a ATAPI device
SYNOPSIS	STATUS ichAtapiReadCapacity ( ATA_DEV * pAtapiDev )
DESCRIPTION	This routine issues a READ CD-ROM CAPACITY command to a specified ATAPI device.
RETURN	<b>OK</b> , or <b>ERROR</b> if the command fails.
RETURNS	Not Available
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorage

# ichAtapiReadTocPmaAtip()

ichAtapiReadTocPmaAtip() – issue a READ TOC command to a ATAPI device
<pre>STATUS ichAtapiReadTocPmaAtip   (    ATA_DEV * pAtapiDev,    UINT32 transferLength,    char * resultBuf )</pre>
This routine issues a READ TOC command to a specified ATAPI device.
<b>OK</b> , or <b>ERROR</b> if the command fails.
Not Available
Not Available

VxWorks Kernel API Reference, 6.6 ichAtapiRemovMediaStatusNotifyVerGet()

SEE ALSO vxbIntelIchStorage

## ichAtapiRemovMediaStatusNotifyVerGet()

NAME	$ichAtapiRemovMediaStatusNotifyVerGet () - get \ the \ Media \ Stat \ Notification \ Version.$
SYNOPSIS	UINT16 ichAtapiRemovMediaStatusNotifyVerGet ( int ctrl, int drive )
DESCRIPTION	This function will return the removable media status notification version of the drive.
RETURNS	Version Number.
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorageShow

# ichAtapiScan()

NAME	ichAtapiScan() – issue SCAN packet command to ATAPI drive.
SYNOPSIS	STATUS ichAtapiScan ( ATA_DEV * pAtapiDev, UINT32 startAddressField, int function )
DESCRIPTION	This function issues SCAN packet command to ATAPI drive. The <i>function</i> argument should be 0x00 for fast forward and 0x10 for fast reversed operation.
RETURN	OK or ERROR
RETURNS	Not Available
ERRNO	Not Available

SEE ALSO vxbIntelIchStorage

## ichAtapiSeek()

ichAtapiSeek() - issues a SEEK packet command to drive. NAME SYNOPSIS STATUS ichAtapiSeek ( ATA\_DEV \* pAtapiDev, UINT32 addressLBA ) DESCRIPTION This function issues a SEEK packet command (not ATA SEEK command) to the specified drive. OK or ERROR RETURN Not Available RETURNS Not Available ERRNO

SEE ALSO vxbIntelIchStorage

#### ichAtapiSetCDSpeed()

NAME	ichAtapiSetCDSpeed() – issue SET CD SPEED packet command to ATAPI drive.	
SYNOPSIS	STATUS ichAtapiSetCDSpeed ( ATA_DEV * pAtapiDev, int readDriveSpeed, int writeDriveSpeed )	
DESCRIPTION	This function issues SET CD SPEED packet command to ATAPI drive while reading and writing of ATAPI drive(CD-ROM) data. The arguments <i>readDriveSpeed</i> and <i>writeDriveSpeat</i> are in Kbytes/Second.	
RETURN	OK or ERROR	

VxWorks Kernel API Reference, 6.6 ichAtapiStartStopUnit()

**RETURNS** Not Available

ERRNO Not Available

SEE ALSO vxbIntelIchStorage

# ichAtapiStartStopUnit()

NAME	<pre>ichAtapiStartStopUnit( ) - Issues START STOP UNIT packet command</pre>
SYNOPSIS	STATUS ichAtapiStartStopUnit ( ATA_DEV * pAtapiDev, int arg0 )
DESCRIPTION	<ul><li>This function issues a command to drive to MEDIA EJECT and MEDIA LOAD. Argument <i>arg0</i> selects to EJECT or LOAD.</li><li>To eject media <i>arg0</i> should be <b>EJECT_DISK</b> To load media <i>arg0</i> should be <b>LOAD_DISK</b></li></ul>
RETURN	OK or ERROR
RETURNS	Not Available
ERRNO	Not Available
SEE ALSO	vxbIntelIchStorage

# ichAtapiStopPlayScan()

NAME	ichAtapiStopPlayScan() – issue STOP PLAY/SCAN packet command to ATAPI drive.
SYNOPSIS	STATUS ichAtapiStopPlayScan ( ATA_DEV * pAtapiDev )
RETURN	OK or ERROR
RETURNS	Not Available

ERRNO Not Available

SEE ALSO vxbIntelIchStorage

#### ichAtapiTestUnitRdy()

NAME ichAtapiTestUnitRdy() – issue a TEST UNIT READY command to a ATAPI drive

- SYNOPSIS STATUS ichAtapiTestUnitRdy ( ATA\_DEV \* pAtapiDev )
- **DESCRIPTION** This routine issues a TEST UNIT READY command to a specified ATAPI drive.
- **RETURNS** OK, or ERROR if the command fails.
- ERRNO Not Available
- SEE ALSO vxbIntelIchStorage

#### ichAtapiVersionNumberGet()

**NAME** ichAtapiVersionNumberGet() – get the ATA/ATAPI version number of the drive.

SYNOPSIS UINT32 ichAtapiVersionNumberGet ( int ctrl, int drive

**DESCRIPTION** This function will return the ATA/ATAPI version number of the drive. Most significant 16 bits represent the Major Version Number and the Lease significant 16 bits represents the minor Version Number.

Major Version Number

Bit 22	ATA/ATAPI-6
Bit 21	ATA/ATAPI-5
Bit 20	ATA/ATAPI-4
Bit 19	ATA-3

VxWorks Kernel API Reference, 6.6 index()

Bit 18 ATA-2

Minor version Number (bit 15 through bit 0)

0001h	Obsolete
0002h	Obsolete
0003h	Obsolete
0004h	ATA-2 published, ANSI X3.279-1996
0005h	ATA-2 X3T10 948D prior to revision 2k
0006h	ATA-3 X3T10 2008D revision 1
0007h	ATA-2 X3T10 948D revision 2k
0008h	ATA-3 X3T10 2008D revision 0
0009h	ATA-2 X3T10 948D revision 3
000Ah	ATA-3 published, ANSI X3.298-199x
000Bh	ATA-3 X3T10 2008D revision 6
000Ch	ATA-3 X3T13 2008D revision 7 and 7a
000Dh	ATA/ATAPI-4 X3T13 1153D revision 6
000Eh	ATA/ATAPI-4 T13 1153D revision 13
000Fh	ATA/ATAPI-4 X3T13 1153D revision 7
0010h	ATA/ATAPI-4 T13 1153D revision 18
0011h	ATA/ATAPI-4 T13 1153D revision 15
0012h	ATA/ATAPI-4 published, ANSI NCITS 317-1998
0013h	Reserved
0014h	ATA/ATAPI-4 T13 1153D revision 14
0015h	ATA/ATAPI-5 T13 1321D revision 1
0016h	Reserved
0017h	ATA/ATAPI-4 T13 1153D revision 17
0018h-FFFFh	Reserved

**RETURNS** ATA/ATAPI version number.

ERRNO Not Available

SEE ALSO vxbIntelIchStorageShow

# index()

NAME	<b>index()</b> – find the first occurrence of a character in a string			
SYNOPSIS	char *index ( FAST const char * s, /* string in which to find character FAST int c /* character to find in string )	*/ */		
SEE ALSO	bLib, strchr().			
-------------	--			
ERRNO	N/A			
RETURNS	A pointer to the located character, or <b>NULL</b> if <i>c</i> is not found.			
DESCRIPTION	This routine finds the first occurrence of character <i>c</i> in string <i>s</i> .			

# infinity()

SEE ALSO	mathALib
ERRNO	Not Available
RETURNS	The double-precision representation of positive infinity.
DESCRIPTION	This routine returns a very large double.
SYNOPSIS	double infinity (void)
NAME	infinity() – return a very large double

# infinityf()

NAME	<pre>infinityf() - return a very large float</pre>
SYNOPSIS	float infinityf (void)
DESCRIPTION	This routine returns a very large float.
RETURNS	The single-precision representation of positive infinity.
ERRNO	Not Available
SEE ALSO	mathALib

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# inflate()

NAME	inflate() – inflate compressed code
SYNOPSIS	<pre>int inflate   (    Byte * src,    Byte * dest,    int nBytes   )</pre>
DESCRIPTION	This routine inflates <i>nBytes</i> of data starting at address <i>src</i> . The inflated code is copied starting at address <i>dest</i> . Two sanity checks are performed on the data being decompressed. First, we look for a magic number at the start of the data to verify that it is really a compressed stream. Second, the entire data is optionally checksummed to verify its integrity. By default, the checksum is not verified in order to speed up the booting process. To turn on checksum verification, set the global variable <b>inflateCksum</b> to <b>TRUE</b> in the BSP.
RETURNS	OK or ERROR.
ERRNO	Not Available
SEE ALSO	inflateLib

# intCRGet()

SEE ALSO	intArchLib, intCRSet()
ERRNO	Not Available
RETURNS	The contents of the cause register.
DESCRIPTION	This routine reads and returns the contents of the MIPS cause register.
SYNOPSIS	int intCRGet (void)
NAME	<b>intCRGet()</b> – read the contents of the cause register (MIPS)

# intCRSet()

NAME	<b>intCRSet()</b> – write the contents of the cause register (MIPS)
SYNOPSIS	<pre>void intCRSet   (    int value /* value to write to cause register */ )</pre>
DESCRIPTION	This routine writes the contents of the MIPS cause register.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	intArchLib, intCRGet()

# intConnect()

NAME	<b>intConnect()</b> – connect a C routine to a hardware interrupt
SYNOPSIS	STATUS intConnect ( VOIDFUNCPTR * vector, /* interrupt vector to attach to */ VOIDFUNCPTR routine, /* routine to be called */ int parameter /* parameter to be passed to routine */ )
DESCRIPTION	This routine connects a specified C routine to a specified interrupt vector. The address of <i>routine</i> is generally stored at <i>vector</i> so that <i>routine</i> is called with <i>parameter</i> when the interrupt occurs. The routine is invoked in supervisor mode at interrupt level. A proper C environment is established, the necessary registers saved, and the stack set up.
	The routine can be any normal C code, except that it must not invoke certain operating system functions that may block or perform I/O operations.
	This routine generally simply calls <b>intHandlerCreate()</b> and <b>intVecSet()</b> . The address of the handler returned by <b>intHandlerCreate()</b> is what actually goes in the interrupt vector.
	This routine takes an interrupt vector as a parameter, which is the byte offset into the vector table. Macros are provided to convert between interrupt vectors and interrupt numbers, see <b>intArchLib</b> .

VxWorks Kernel API Reference, 6.6 intConnect()

```
NOTE ARM ARM processors generally do not have on-chip interrupt controllers. Control of interrupts is a BSP-specific matter. This routine calls a BSP-specific routine to install the handler such that, when the interrupt occurs, routine is called with parameter.
```

**NOTE X86** Refer to the special x86 routine **intHandlerCreateI86()**.

NOTE SH The on-chip interrupt controller (INTC) design of SH architecture depends on the processor type, but there are some similarities. The number of external interrupt inputs are limited, so it may necessary to multiplex some interrupt requests. However most of them are auto-vectored, thus have only one vector to an external interrupt input. As a framework to handle this type of multiplexed interrupt, you can use your original intConnect code by hooking it to \_func\_intConnectHook pointer. If \_func\_intConnectHook is set, the SH version of **intConnect()** simply calls the hooked routine with same arguments, then returns the status of hooked routine. A **sysLib** sample is shown below:

```
#include <intLib.h>
#include <iv.h>
                            /* INUM_INTR_HIGH for SH7750/SH7700 */
#define SYS_INT_TBL_SIZE (255 - INUM_INTR_HIGH)
typedef struct
    {
    VOIDFUNCPTR routine; /* routine to be called */
    int parameter; /* parameter to be passed */
    } SYS_INT_TBL;
LOCAL SYS_INT_TBL sysIntTbl [SYS_INT_TBL_SIZE]; /* local vector table */
LOCAL int sysInumVirtBase = INUM_INTR_HIGH + 1;
STATUS sysIntConnect
   VOIDFUNCPTR *vec, /* interrupt vector to attach to
VOIDFUNCPTR routine, /* routine to be called */
int param
   (
                                                                           * /
    int param
                            /* parameter to be passed to routine */
    )
    {
    FUNCPTR intDrvRtn;
    if (vec >= INUM_TO_IVEC (0) && vec < INUM_TO_IVEC (sysInumVirtBase))
        {
        /* do regular intConnect() process */
        intDrvRtn = intHandlerCreate ((FUNCPTR)routine, param);
        if (intDrvRtn == NULL)
            return ERROR;
        /* make vector point to synthesized code */
        intVecSet ((FUNCPTR *)vec, (FUNCPTR)intDrvRtn);
        }
    else
```

```
{
        int index = IVEC_TO_INUM (vec) - sysInumVirtBase;
        if (index < 0 || index >= SYS_INT_TBL_SIZE)
            return ERROR;
        sysIntTbl [index].routine = routine;
        sysIntTbl [index].parameter = param;
        }
    return OK;
    }
void sysHwInit (void)
    {
    . . .
    _func_intConnectHook = (FUNCPTR) sysIntConnect;
    }
LOCAL void sysVmeIntr (void)
    {
    volatile UINT32 vec = *VME VEC REGISTER; /* get VME interrupt vector */
    int i = vec - sysInumVirtBase;
    if (i >= 0 && i < SYS_INT_TBL_SIZE && sysIntTbl[i].routine != NULL)
        (*sysIntTbl[i].routine)(sysIntTbl[i].parameter);
    else
        logMsg ("uninitialized VME interrupt: vec = %d\n", vec,0,0,0,0,0);
    }
void sysHwInit2 (void)
    {
    int i;
    /* initialize VME interrupts dispatch table */
    for (i = 0; i < SYS_INT_TBL_SIZE; i++)</pre>
        {
        sysIntTbl[i].routine = (VOIDFUNCPTR)NULL;
        sysIntTbl[i].parameter = NULL;
        }
    /* connect generic VME interrupts handler */
    intConnect (INT_VEC_VME, sysVmeIntr, NULL);
    . . .
    }
```

The used vector numbers of SH processors are limited to certain ranges, depending on the processor type. The **sysInumVirtBase** should be initialized to a value higher than the last used vector number, defined as **INUM\_INTR\_HIGH**. It is typically safe to set **sysInumVirtBase** to (**INUM\_INTR\_HIGH** + 1).

The **sysIntConnect()** routine simply acts as the regular **intConnect()** if *vector* is smaller than **INUM\_TO\_IVEC** (sysInumVirtBase), so **sysHwInit2()** connects a common VME

interrupt dispatcher **sysVmeIntr** to the multiplexed interrupt vector. If *vector* is equal to or greater than **INUM\_TO\_IVEC** (sysInumVirtBase), the **sysIntConnect()** fills a local vector entry in sysIntTbl[] with an individual VME interrupt handler, in a coordinated manner with **sysVmeIntr**.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted.

- **RETURNS** OK, or ERROR if the interrupt handler cannot be built.
- ERRNO Not Available

SEE ALSO intArchLib, intHandlerCreate(), intVecSet(), intDisconnect()

#### intContext()

NAME	<pre>intContext() - determine if executing in interrupt context</pre>	
SYNOPSIS	BOOL intContext (void)	
DESCRIPTION	This routine returns <b>TRUE</b> only if the caller is executing in an interrupt context. If executing in a task context <b>FALSE</b> is returned.	
SMP CONSIDERATIONS		
	In a VxWorks SMP system it is possible for the CPUs to be in different contexts (task or interrupt) at the same time. Therefore this routine returns the status for the CPU the caller is executing on.	
RETURNS	TRUE or FALSE.	
ERRNO	N/A	

SEE ALSO intLib

	intCount()
NAME	<b>intCount()</b> – get the current interrupt nesting depth
SYNOPSIS	int intCount (void)
DESCRIPTION	This routine returns the number of interrupts that are currently nested.
SMP CONSIDERAT	INNS In a VxWorks SMP system it is possible for the CPUs to be in different contexts (task or interrupt) at the same time. Therefore this routine returns the nested interrupt count for the CPU the caller is executing on.
RETURNS	The number of nested interrupts.
ERRNO	N/A
SEE ALSO	intLib

## intCpuLock()

NAME intCpuLock() – lock out interrupts on local CPU

SYNOPSIS int intCpuLock (void)

**DESCRIPTION** This routine disables interrupts on the CPU the calling task or ISR is running on. The returned value is a lock-out key to be used in a subsequent call to **intCpuUnlock()** to release the lock. Execution of interrupts on other CPUs in the SMP system is not affected by this routine. Because of this behaviour this routine is not a suitable mutual exclusion mechanism unless all tasks and/or ISRs participating in the mutual exclusion scenario have a single CPU affinity to the very same CPU.

Calling this routine on the uniprocessor version of VxWorks is equivalent to calling **intLock()**.

Invoking a VxWorks system routine after having locked interrupts using **intCpuLock()** on VxWorks SMP is not permitted and will cause the call to abort and an error to be reported. Not all VxWorks APIs enforce this restriction. Only those that are **intCpuLock restricted**. The reference entries in the VxWorks Kernel API Reference manual specifies when this restriction applies. Since the **intCpuLock()** behaviour in the uniprocessor version of VxWorks is identical to the **intLock()** API behaviour, the concept of **intCpuLock()** restricted APIs only applies to VxWorks SMP.

VxWorks Kernel API Reference, 6.6 intCpuUnlock()

RETURNS	An architecture-dependent lock-out key for the interrupt level prior to the call.
ERRNO	Not Available

intArchLib, intCpuUnlock(), taskCpuLock(), intLockLevelSet()

## intCpuUnlock()

SEE ALSO

NAME intCpuUnlock() – cancel local CPU interrupt lock SYNOPSIS void intCpuUnlock ( int lockKey /\* lock-out key returned by preceding intCpuLock() \*/ ) This routine removes the lock established using intCpuLock (). It re-enables interrupts on DESCRIPTION the CPU the calling task or ISR is running on. Calling this routine on the uniprocessor version of VxWorks is equivalent to calling **intUnlock()**. The parameter *lockKey* is an architecture-dependent lock-out key returned by a preceding intCpuLock() call. RETURNS N/A Not Available ERRNO intArchLib, intCpuLock(), taskCpuUnlock() SEE ALSO

#### intDisable()

NAME	intDisable() – disable corresponding interrupt bits (MIPS, PowerPC, ARM)
SYNOPSIS	<pre>int intDisable   (    int level /* new interrupt bits (0x0 - 0xff00) */ )</pre>
DESCRIPTION	On MIPS and PowerPC architectures, this routine disables the corresponding interrupt bits from the present status register.

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NOTE ARM	ARM processors generally do not have on-chip interrupt controllers. Control of interrupts is a BSP-specific matter. This routine calls a BSP-specific routine to disable a particular interrupt level, regardless of the current interrupt mask level.
NOTE MIPS	For MIPS, the macros <b>SR_IBIT1</b> - <b>SR_IBIT8</b> define bits that may be set.
RETURNS	OK or ERROR. (MIPS: The previous contents of the status register).
ERRNO	Not Available
SEE ALSO	intArchLib, intEnable()

# intDisconnect( )

NAME	intDisconnect() – disconnect a C routine from a hardware interrupt
SYNOPSIS	STATUS intDisconnect ( VOIDFUNCPTR * vector, /* interrupt vector to dettach from */ VOIDFUNCPTR routine, /* routine to disconnect */ int parameter /* parameter to be matched */ )
DESCRIPTION	This routine disconnects a specified C routine that has a specified <i>parameter</i> from a specified interrupt vector.
	The caller of this routine must first disable the source of interrupts before calling this routine.
SMP CONSIDERATIO	NS
	This API is spinlock and intCpuLock restricted.
RETURNS	OK, ERROR if the interrupt handler cannot be disconnected.
ERRNO	The following are the possible errnos returned when ISR object is supported.
	S_intLib_NOT_ISR_CALLABLE this routine must not be called from an ISR
SEE ALSO	intLib, intConnect()

## intEnable()

NAME	intEnable() – enable corresponding interrupt bits (MIPS, PowerPC, ARM)
SYNOPSIS	<pre>int intEnable   (    int level /* new interrupt bits (0x00 - 0xff00) */   )</pre>
DESCRIPTION	This routine enables the input interrupt bits on the present status register of the MIPS and PowerPC processors.
NOTE ARM	ARM processors generally do not have on-chip interrupt controllers. Control of interrupts is a BSP-specific matter. This routine calls a BSP-specific routine to enable the interrupt. For each interrupt level to be used, there must be a call to this routine before it will be allowed to interrupt.
NOTE MIPS	For MIPS, it is strongly advised that the level be a combination of <b>SR_IBIT1 - SR_IBIT8</b> .
RETURNS	<b>OK</b> or <b>ERROR</b> . (MIPS: The previous contents of the status register).
ERRNO	Not Available
SEE ALSO	intArchLib, intDisable()

## intHandlerCreate()

**NAME** intHandlerCreate() – construct an interrupt handler for a C routine (MC680x0, x86, MIPS, SimSolaris)

SYNOPSIS	S FUNCPTR intHandlerCreate						
	( FUNCPTR int )	routine, parameter	/* /*	routine to be called parameter to be passed	to	routine	*/ */

**DESCRIPTION** This routine builds an interrupt handler around the specified C routine. This interrupt handler is then suitable for connecting to a specific vector address with **intVecSet()**. The interrupt handler is invoked in supervisor mode at interrupt level. A proper C environment is established, the necessary registers saved, and the stack set up.

The routine can be any normal C code, except that it must not invoke certain operating system functions that may block or perform I/O operations.

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SMP CONSIDERATIONS		
	This API is spinlock and intCpuLock restricted.	
RETURNS	A pointer to the new interrupt handler, or NULL if memory is insufficient.	
ERRNO	Not Available	
SEE ALSO	intArchLib, intConnect()	

## intHandlerCreateI86()

NAME	<pre>intHandlerCreateI86() - const</pre>	ruct an interrupt handler for	a C routine (x86)
SYNOPSIS	<pre>FUNCPTR intHandlerCreateI80   (    FUNCPTR routine,    int parameter,    FUNCPTR routineBoi,    int parameterBoi,    FUNCPTR routineEoi,    int parameterEoi   )</pre>	/* routine to be called /* parameter to be passe /* BOI routine to be cal /* parameter to be passe /* EOI routine to be cal /* parameter to be passe	*/ d to routine */ led */ d to routineBoi */ led */ d to routineEoi */
DESCRIPTION	This routine builds an interrup handler is then suitable for con interrupt handler is invoked in is established, the necessary re	t handler around a specified necting to a specific vector ac supervisor mode at interrupt gisters saved, and the stack se	C routine. This interrupt ldress with <b>intVecSet()</b> . The level. A proper C environment et up.
	The routine can be any normal system functions that may bloc	C code, except that it must n k or perform I/O operations	ot invoke certain operating
IMPLEMENTATION	This routine builds an interrup	t handler of the following for	rm in allocated memory:
	00 e8 kk kk kk kk 05 50 06 52 07 51	call _intEnt pushl %eax pushl %edx pushl %ecx	* tell kernel * save regs
	08       68       pp       pp       pp       pp         13       e8       rr       rr       rr       rr         18       68       pp       pp       pp       pp         23       e8       rr       rr       rr       rr         28       68       pp       pp       pp       pp         23       e8       rr       rr       rr       rr         28       68       pp       pp       pp       pp         33       e8       rr       rr       rr       rr         38       33       c4       0c	<pre>pushl \$_parameterBoi call _routineBoi pushl \$_parameter call _routine pushl \$_parameterEoi call _routineEoi addl \$12, %e popl %ecx popl %edx</pre>	<pre>* push BOI param * call BOI routine * push param * call C routine * push EOI param * call EOI routine sp * pop param * restore regs</pre>
	43 58	popl %eax	

44 e9 kk kk kk jmp \_intExit \* exit via kernel

Third and fourth parameter of **intHandlerCreateI86()** are the BOI routine address and its parameter that are inserted into the code as "routineBoi" and "parameterBoi". Fifth and sixth parameter of **intHandlerCreateI86()** are the EOI routine address and its parameter that are inserted into the code as "routineEoi" and "parameterEoi". The BOI routine detects if this interrupt is stray/spurious/phantom by interrogating the interrupt controller, and returns from the interrupt if it is. The EOI routine issues End Of Interrupt signal to the interrupt controller, if it is required by the controller. Each interrupt controller has its own BOI and EOI routine. They are located in the BSP, and their address and parameter are taken by the intEoiGet function (set to **sysIntEoiGet()** in the BSP). The Tornado 2, and later, BSPs should use the BOI and EOI mechanism with intEoiGet function pointer.

To keep the Tornado 101 BSP backward compatible, the function pointer intEOI is not removed. If intEoiGet is NULL, it should be set to the **sysIntEoiGet()** routine in the BSP, **intHandlerCreate()** and the intEOI function pointer (set to **sysIntEOI()** in the Tornado 101 BSP) is used.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted.

- **RETURNS** A pointer to the new interrupt handler, or NULL if memory is insufficient.
- ERRNO Not Available

SEE ALSO intArchLib

#### intLevelSet()

NAME	intLevelSet() – se	et the interrupt level (MC680X0, x86, ARM, SimSolaris, SimNT and SH)	
SYNOPSIS	<pre>int intLevelSet    (    int level    )</pre>	/* new interrupt level mask */	
DESCRIPTION	This routine changes the interrupt mask in the status register to take on the value by <i>level</i> . Interrupts are locked out at or below that level. The value of <i>level</i> must l following range:		
	MC680x0: SH: ARM: SimSolaris:	0 - 7 0 - 15 BSP-specific 0 - 1	

	x86:	interrupt controller specific
	On x86 systems, t controller manage always.	here are no interrupt level in the processor and the external interrupt as the interrupt level. Therefore this routine does nothing and returns <b>OK</b>
NOTE SIMNT	This routine does	nothing.
WARNING	Do not call VxWo re-enable interrup	rks system routines with interrupts locked. Violating this rule may ts unpredictably.
RETURNS	The previous inte	rrupt level.
ERRNO	Not Available	
SEE ALSO	intArchLib	

## intLock()

**NAME** intLock() – lock out interrupts

- SYNOPSIS int intLock (void)
- **DESCRIPTION** This routine disables interrupts. It can be called from either interrupt or task level. The **intLock()** routine returns an architecture-dependent lock-out key representing the interrupt level prior to the call; this key can be passed to **intUnlock()** to re-enable interrupts.

For MC680x0, x86, and SH architectures, interrupts are disabled at the level set by **intLockLevelSet()**. The default lock-out level is the highest interrupt level (MC680x0 = 7, x86 = 1, SH = 15).

For SimSolaris architecture, interrupts are masked. Lock-out level returned is 1 if interrupts were already locked, 0 otherwise.

For SimNT, a windows semaphore is used to lock the interrupts. Lock-out level returned is 1 if interrupts were already locked, 0 otherwise.

For MIPS processors, interrupts are disabled at the master lock-out level; this means no interrupt can occur even if unmasked in the IntMask bits (15-8) of the status register.

For ARM processors, interrupts (IRQs) are disabled by setting the I bit in the CPSR. This means no IRQs can occur.

# VxWorks Kernel API Reference, 6.6 intLock()

	For PowerPC pro offset 0x500) is d interrupted by an	ocessors, there is only one interrupt vector. The external interrupt (vector isabled when <b>intLock()</b> is called; this means that the processor cannot be ny external event.		
IMPLEMENTATION	The lock-out key	is implemented differently for different architectures:		
	MC680x0: MIPS: x86: PowerPC: ARM SH: SimSolaris: SimNT:	interrupt field mask status register interrupt enable flag (IF) bit from EFLAGS register MSR register value I bit from the CPSR status register 1 or 0 1 or 0		
WARNINGS	Invoking a VxWo re-enabled for an higher priority ta while another tas	orks system routine with interrupts locked may result in interrupts being a unspecified period of time. If the called routine blocks, or results in a ask becoming eligible for execution (READY), interrupts will be re-enabled sk executes, or while the kernel is idle.		
	To prevent interrupts from being re-enabled for the case where a called routine results in a higher priority task becoming eligible for execution, the <b>taskLock()</b> primitive can be used to disable rescheduling. Note that if a task blocks or suspends, the scheduler will always select the highest priority ready task to execute (or become idle) regardless of whether the task has locked preemption via <b>taskLock()</b> , and thus interrupts will be re-enabled.			
	The interrupt loc task disables inte calling task to blo be restored when code fragment, ir	k level is an attribute of a task, i.e. it's part of the task context. Thus, if a rrupts and subsequently invokes a VxWorks system routine that causes the ock or cause a higher priority task to be ready, the interrupt lock level will a the task is later rescheduled for execution. For example, in the following interrupts will be disabled after returning from the <b>taskDelay()</b> invocation:		

```
lockKey = intLock ();
... (work with interrupts locked out)
taskDelay (sysClkRateGet() * 10); /* delay for 10 seconds */
... (work with interrupts locked out)
```

Finally, the above descriptions only applies for tasks since ISRs are not permitted to block or suspend, and task rescheduling only occurs when an ISR completes execution.

EXAMPLES	<pre>lockKey = intLock ();</pre>
	(work with interrupts locked out)

intUnlock (lockKey);

To lock out interrupts and task scheduling as well (see WARNINGS above):

```
if (taskLock() == OK)
    {
        lockKey = intLock ();
        ... (critical section)
        intUnlock (lockKey);
        taskUnlock();
     }
else
     {
        ... (error message or recovery attempt)
     }
```

#### SMP CONSIDERATIONS

This routine is not available in VxWorks SMP. Refer to the VxWorks SMP Migration Guide for suitable alternatives.

RETURNS An architecture-de	ependent lock-out ke	ey for the interrupt le	evel prior to the call.
----------------------------	----------------------	-------------------------	-------------------------

ERRNO	Not Available
-------	---------------

SEE ALSO intArchLib, intUnlock(), taskLock(), intLockLevelSet()

#### intLockLevelGet()

NAME	<pre>intLockLevelGet() - get the current interrupt lock-out level (MC680x0, x86, ARM, SH, SimSolaris, SimNT)</pre>
SYNOPSIS	int intLockLevelGet (void)
DESCRIPTION	This routine returns the current interrupt lock-out level, which is set by <b>intLockLevelSet()</b> and stored in the globally accessible variable <b>intLockMask</b> . This is the interrupt level currently masked when interrupts are locked out by <b>intLock()</b> . The default lock-out level (MC680x0 = 7, x86 = 1, SH = 15) is initially set by <b>kernelInit()</b> when VxWorks is initialized.
NOTE SIMNT	This routine does nothing.
RETURNS	The interrupt level currently stored in the interrupt lock-out mask. (ARM = <b>ERROR</b> always)
ERRNO	Not Available
SEE ALSO	intArchLib, intLockLevelSet()

#### intLockLevelSet()

NAME intLockLevelSet() – set the current interrupt lock-out level (MC680x0, x86, ARM, SH, SimSolaris, SimNT)

SYNOPSIS void intLockLevelSet
(
int newLevel /\* new interrupt level \*/
)

**DESCRIPTION** This routine sets the current interrupt lock-out level and stores it in the globally accessible variable **intLockMask**. The specified interrupt level is masked when interrupts are locked by **intLock()**. The default lock-out level (MC680x0 = 7, x86 = 1, SH = 15) is initially set by **kernelInit()** when VxWorks is initialized.

#### NOTE SIMSOLARIS, SIMNT

This routine does nothing.

- NOTE ARM On the ARM, this call establishes the interrupt level to be set when **intLock()** is called.
- RETURNS N/A

ERRNO Not Available

SEE ALSO intArchLib, intLockLevelGet(), intLock(), taskLock()

### intSRGet()

SEE ALSO	intArchLib, intSRSet()
ERRNO	Not Available
RETURNS	The previous contents of the status register.
DESCRIPTION	This routine reads and returns the contents of the MIPS status register.
SYNOPSIS	int intSRGet (void)
NAME	<b>intSRGet()</b> – read the contents of the status register (MIPS)

# intSRSet()

NAME	intSRSet() – update the contents of the status register (MIPS)
SYNOPSIS	<pre>int intSRSet   (    int value /* value to write to status register */ )</pre>
DESCRIPTION	This routine updates and returns the previous contents of the MIPS status register.
RETURNS	The previous contents of the status register.
ERRNO	Not Available
SEE ALSO	intArchLib, intSRGet()

## intStackEnable()

NAME	<b>intStackEnable()</b> – enable or disable the interrupt stack usage (x86)
SYNOPSIS	STATUS intStackEnable ( BOOL enable /* TRUE to enable, FALSE to disable */ )
DESCRIPTION	This routine enables or disables the interrupt stack usage and is only callable from the task level. An Error is returned for any other calling context. The interrupt stack usage is disabled in the default configuration for the backward compatibility. Routines that manipulate the interrupt stack, are located in the file <b>i86/windALib.s</b> . These routines include <b>intStackEnable()</b> , <b>intEnt()</b> and <b>intExit()</b> .
RETURNS	<b>OK</b> , or <b>ERROR</b> if it is not in the task level.
ERRNO	Not Available
SEE ALSO	intArchLib

VxWorks Kernel API Reference, 6.6 intUninitVecSet()

## intUninitVecSet()

NAME	intUninitVecSet() – set the uninitialized vector handler (ARM)
SYNOPSIS	<pre>void intUninitVecSet    (    VOIDFUNCPTR routine /* ptr to user routine */ )</pre>
DESCRIPTION	This routine installs a handler for the uninitialized vectors to be called when any uninitialised vector is entered.
RETURNS	N/A.
ERRNO	Not Available
SEE ALSO	intArchLib

# intUnlock()

NAME	<pre>intUnlock() – cancel interrupt locks</pre>	
SYNOPSIS	<pre>void intUnlock   (     int lockKey /* lock-out key returned by preceding intLock() */   )</pre>	
DESCRIPTION	This routine re-enables interrupts that have been disabled by <b>intLock()</b> . The parameter <i>lockKey</i> is an architecture-dependent lock-out key returned by a preceding <b>intLock()</b> call.	
SMP CONSIDERATIONS		
	This routine is not available in VxWorks SMP. Refer to the VxWorks SMP Migration Guide for suitable alternatives.	
RETURNS	N/A	
ERRNO	Not Available	
SEE ALSO	intArchLib, intLock()	

NAME	intVecBaseGet() – get the vector (trap) base address (MC680x0, x86, MIPS, ARM, SimSolaris, SimNT)
SYNOPSIS	FUNCPTR *intVecBaseGet (void)
DESCRIPTION	This routine returns the current vector base address, which is set with <b>intVecBaseSet()</b> .
RETURNS	The current vector base address (MIPS = 0 always, ARM = 0 always, SimSolaris = 0 always and SimNT = 0 always).
ERRNO	Not Available
SEE ALSO	intArchLib, intVecBaseSet()

# intVecBaseSet( )

NAME	intVecBaseSet() – set the vector (trap) base address (MC680x0, x86, MIPS, ARM, SimSolaris, SimNT)
SYNOPSIS	<pre>void intVecBaseSet   (   FUNCPTR * baseAddr /* new vector (trap) base address */ )</pre>
DESCRIPTION	This routine sets the vector (trap) base address. The CPU's vector base register is set to the specified value, and subsequent calls to <b>intVecGet()</b> or <b>intVecSet()</b> will use this base address. The vector base address is initially 0, until modified by calls to this routine.
NOTE 68000	The 68000 has no vector base register; thus, this routine is a no-op for 68000 systems.
NOTE MIPS	The MIPS processors have no vector base register; thus this routine is a no-op for this architecture.
NOTE SH77XX	This routine sets <i>baseAddr</i> to vbr, then loads an interrupt dispatch code to (vbr + $0x600$ ). When SH77XX processor accepts an interrupt request, it sets an exception code to INTEVT register and jumps to (vbr + $0x600$ ). Thus this dispatch code is commonly used for all interrupts' handling.

2

All SH77XX processors have INTEVT register at address 0x0fffffd8. The SH7707 processor has yet another INTEVT2 register at address 0x04000000, to identify its enhanced interrupt sources. The dispatch code obtains the address of INTEVT register from a global constant intEvtAdrs. The constant is defined in sysLib, thus the selection of INTEVT/INTEVT2 is configurable at BSP level. The intEvtAdrs is loaded to (vbr + 4) by intVecBaseSet().After fetching the exception code, the interrupt dispatch code applies a new interrupt mask to the status register, and jumps to an individual interrupt handler. The new interrupt mask is taken from intPrioTable[], which is defined in sysALib. The intPrioTable[] is loaded to (vbr + 0xc00) by intVecBaseSet().NOTE ARMThe ARM processors have no vector base register; thus this routine is a no-op for this architecture.NOTE SIMSOLARIS,SMNT This routine does nothing.RETURNSN/AERRNONot AvailableSEE ALSOintArchLib, intVecBaseGet(), intVecGet(), intVecSet()		The exception codes are 12bits width, and interleaved by $0x20$ . VxWorks for SH77XX locates a vector table at (vbr + $0x800$ ), and defines the vector offsets as (exception codes / 8). This vector table is commonly used by all interrupts, exceptions, and software traps.	
After fetching the exception code, the interrupt dispatch code applies a new interrupt mask to the status register, and jumps to an individual interrupt handler. The new interrupt mask is taken from intPrioTable[], which is defined in sysALib. The intPrioTable[] is loaded to (vbr + 0xc00) by intVecBaseSet().NOTE ARMThe ARM processors have no vector base register; thus this routine is a no-op for this architecture.NOTE SIMSOLARISSIMNT This routine does nothing.RETURNSN/AERRNONot AvailableSEE ALSOintArchLib, intVecBaseGet(), intVecGet(), intVecSet()		All SH77XX processors have INTEVT register at address 0xffffffd8. The SH7707 processor has yet another INTEVT2 register at address 0x04000000, to identify its enhanced interrupt sources. The dispatch code obtains the address of INTEVT register from a global constant <b>intEvtAdrs</b> . The constant is defined in <b>sysLib</b> , thus the selection of INTEVT/INTEVT2 is configurable at BSP level. The <b>intEvtAdrs</b> is loaded to (vbr + 4) by <b>intVecBaseSet()</b> .	
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RETURNS       N/A         ERRNO       Not Available         SEE ALSO       intArchLib, intVecBaseGet(), intVecGet(), intVecSet()	NOTE SIMSOLARIS, SIMNT		
RETURNS     N/A       ERRNO     Not Available       SEE ALSO     intArchLib, intVecBaseGet(), intVecGet(), intVecSet()		This routine does nothing.	
ERRNONot AvailableSEE ALSOintArchLib, intVecBaseGet( ), intVecGet( ), intVecSet( )	RETURNS	N/A	
SEE ALSO intArchLib, intVecBaseGet( ), intVecGet( ), intVecSet( )	ERRNO	Not Available	
	SEE ALSO	intArchLib, intVecBaseGet( ), intVecGet( ), intVecSet( )	

# intVecGet()

NAME	<pre>intVecGet() – get an interrupt vector (</pre>	MC680x0, x86, MIPS, SH, SimSolaris, SimNT)
SYNOPSIS	FUNCPTR intVecGet ( FUNCPTR * vector /* vector o )	ffset */
DESCRIPTION	This routine returns a pointer to the exception/interrupt handler attached to a specific vector. The vector is specified as an offset into the CPU's vector table. This vector ta starts, by default, at:	
	MC680x0: MIPS: x86: SH702x/SH703x/SH704x/SH76xx:	0 excBsrTbl in excArchLib 0 excBsrTbl in excArchLib

	SH77xx: SimSolaris:	vbr + 0x800 0
	However, the vector table may be set to for which it is available).	start at any address with <b>intVecBaseSet( )</b> (on CPUs
	This routine takes an interrupt vector as a parameter, which is the byte offset into the ve table. Macros are provided to convert between interrupt vectors and interrupt numbers <b>intArchLib</b> .	
NOTE SIMNT	This routine does nothing and always	returns 0.
RETURNS	A pointer to the exception/interrupt ha	andler attached to the specified vector.
ERRNO	Not Available	
SEE ALSO	intArchLib, intVecSet(), intVecBaseSe	et()

## intVecGet2()

intVecGet2() – get a CPU vector, gate type(int/trap), and gate selector (x86) NAME SYNOPSIS void intVecGet2 ( FUNCPTR \* vector, /\* vector offset \*/
FUNCPTR \* pFunction, /\* address to place in vector \*/
int \* pIdtGate, /\* IDT\_TRAP\_GATE or IDT\_INT\_GATE \*/
int \* pIdtSelector /\* sysCsExc or sysCsInt \*/ int \* pIdtSelector /\* sysCsExc or sysCsInt \*/ ) DESCRIPTION This routine gets a pointer to the exception/interrupt handler attached to a specified vector, the type of the gate, the selector of the gate. The vector is specified as an offset into the CPU's vector table. This vector table starts, by default, at address 0. However, the vector table may be set to start at any address with intVecBaseSet(). N/A RETURNS ERRNO Not Available SEE ALSO intArchLib, intVecBaseSet(), intVecGet(), intVecSet(), intVecSet2()

### intVecSet()

intVecSet() – set a CPU vector (trap) (MC680x0, x86, MIPS, SH, SimSolaris, SimNT) NAME SYNOPSIS void intVecSet ( FUNCPTR \* vector, /\* vector offset \*/ FUNCPTR function /\* address to place in vector \*/ ) This routine attaches an exception/interrupt/trap handler to a specified vector. The vector DESCRIPTION is specified as an offset into the CPU's vector table. This vector table starts, by default, at: MC680x0: 0 MIPS: excBsrTbl in excArchLib x86: 0 SH702x/SH703x/SH704x/SH76xx: excBsrTbl in excArchLib SH77xx: vbr + 0x800SimSolaris: 0 However, the vector table may be set to start at any address with intVecBaseSet() (on CPUs for which it is available). The vector table is set up in usrInit(). This routine takes an interrupt vector as a parameter, which is the byte offset into the vector table. Macros are provided to convert between interrupt vectors and interrupt numbers, see intArchLib. The **intVecSet()** routine puts this generated code into the trap table entry corresponding to vector. Window overflow and window underflow are sacred to the kernel and may not be pre-empted. They are written here only to track changing trap base registers (TBRs). With the "branch anywhere" scheme (as opposed to the branch PC-relative +/-8 megabytes) the first instruction in the vector table must not be a change of flow control nor affect any critical registers. The JMPL that replaces the BA will always execute the next vector's first instruction. NOTE MIPS On MIPS CPUs the vector table is set up statically in software. The specified interrupt handler *function* has to coordinate with an interrupt stack frame NOTE SH77XX which is specially designed for SH77XX version of VxWorks: [ task's stack ] [ interrupt stack ] xxx | high address +---+ yyy \_<----|task'sp| 0 777 INTEVT | -4 | low address | ssr | -8

This interrupt stack frame is formed by a common interrupt dispatch code which is loaded at (vbr + 0x600). You usually do not have to pay any attention to this stack frame, since **intConnect()** automatically appends an appropriate stack manipulation code to your interrupt service routine. The **intConnect()** assumes that your interrupt service routine (ISR) is written in C, thus it also wraps your ISR in minimal register save/restore codes. However if you need a very fast response time to a particular interrupt request, you might want to skip this register save/restore sequence by directly attaching your ISR to the corresponding vector table entry using **intVecSet()**. Note that this technique is only applicable to an interrupt service with NO VxWorks system call. For example it is not allowed to use **semGive()** or **logMsg()** in the interrupt service routine which is directly attached to vector table by intVecSet(). To facilitate the direct usage of intVecSet() by user, a special entry point to exit an interrupt context is provided within the SH77XX version of VxWorks kernel. This entry point is located at address (vbr + intRte1W), here the intRte1W is a global symbol for the vbr offset of the entry point in 16 bit length. This entry point intRte1 assumes that the current register bank is 0 (SR.RB == 0), and r1 and r0 are still saved on the interrupt stack, and it also requires 0x70000000 in r0. Then intRte1 properly cleans up the interrupt stack and executes *rte* instruction to return to the previous interrupt or task context. The following code is an example of intRte1 usage. Here the corresponding intPrioTable[] entry is assumed to be 0x400000X0, namely MD=1, RB=0, BL=0 at the beginning of usrIsr1.

```
.text
    .align 2
    .global _usrIsr1
    .type _usrIsr1,@function
    .extern _usrRtn
    .extern _intRte1W
                          /* intPrioTable[] sets SR to 0x400000X0 */
usrIsr1:
    mov.l r0,@-sp
                         /* must save r0 first (BANK0) */
    mov.l r1,@-sp
                          /* must save r1 second (BANK0) */
   mov.l r2,@-sp
                         /* save rest of volatile registers (BANKO) */
   mov.l r3,@-sp
   mov.l r4,@-sp
   mov.l r5,@-sp
   mov.l r6,@-sp
   mov.l r7,@-sp
   sts.l pr,@-sp
   sts.l mach,@-sp
   sts.l macl,@-sp
   mov.l UsrRtn,r0
```

```
jsr
          @r0
                        /* call user's C routine */
                         /* (delay slot) */
   nop
   lds.l
          @sp+,macl
                         /* restore volatile registers (BANKO) */
   lds.l @sp+,mach
   lds.l @sp+,pr
   mov.l @sp+,r7
   mov.l @sp+,r6
   mov.l @sp+,r5
   mov.l
          @sp+,r4
   mov.l
          @sp+,r3
   mov.l @sp+,r2
                         /* intRtel restores r1 and r0 */
   mov.l
          IntRte1W,r1
   mov.w @r1,r0
   stc
          vbr,r1
   add
          r0,r1
                       /* r0: 0x70000000 */
   mov.l IntRteSR,r0
   jmp
          @r1
                        /* let intRte1 clean up stack, then rte */
   nop
                         /* (delay slot) */
          .align 2
UsrRtn:
          .long _usrRtn
                               /* user's C routine */
IntRteSR: .long 0x7000000
                                /* MD=1, RB=1, BL=1 */
IntRte1W:
          .long intRte1W
```

The **intRte1** sets r0 to status register (SR: 0x70000000), to safely restore SPC/SSR and to clean up the interrupt stack. Note that TLB mishit exception immediately reboots CPU while SR.BL=1. To avoid this fatal condition, VxWorks loads the **intRte1** code and the interrupt stack to a physical address space (P1) where no TLB mishit happens.

Furthermore, there is another special entry point called **intRte2** at an address (vbr + intRte2W). The **intRte2** assumes that SR is already set to 0x70000000 (MD: 1, RB: 1, BL: 1), then it does not restore r1 and r0. While SR value is 0x70000000, you may use r0,r1,r2,r3 in BANK1 as volatile registers. The rest of BANK1 registers (r4,r5,r6,r7) are non-volatile, so if you need to use them then you have to preserve their original values by saving/restoring them on the interrupt stack. So, if you need the ultimate interrupt response time, you may set the corresponding intPrioTable[] entry to NULL and manage your interrupt service only with r0,r1,r2,r3 in BANK1 as shown in the next sample code:

```
mov.l UsrIntCnt,r1
   mov.l X1FFFFFFF, r2
   mov.l
           UsrP1Base,r3
   mov.l
           @r3,r3
                          /* r3: SH P1 TEXT BASE */
   and
           r2,r1
           r3,r1
                          /* r1: _usrIntCnt address in P1 */
   or
   mov.l @r1,r0
           #1,r0
   add
                         /* increment counter */
   mov.l r0,@r1
   mov.l
           IntRte2W,r1
           r2,r1
   and
           r3,r1
                          /* r1: intRte2W address in P1 */
   or
           @r1,r0
   mov.w
   stc
           vbr,r1
   add
           r1,r0
                          /* let intRte2 clean up stack, then rte */
   jmp
           @r0
                          /* (delay slot) */
   nop
           .align 2
UsrIntAck: .long 0xa0001234
                                  /* interrupt acknowledge register */
UsrIntCnt: .long _usrIntCnt
IntRte2W: .long _intRte2W
X1FFFFFFF: .long 0x1ffffff
UsrP1Base: .long _vxShP1TextBase
```

Note that the entire interrupt service is executed under SR.BL=1 in this sample code. It means that any access to virtual address space may reboot CPU, since TLB mishit exception is blocked. Therefore **usrIsr2** has to access **usrIntCnt** and **intRte2W** from P1 region. Also **usrIsr2** itself has to be executed on P1 region, and it can be done by relocating the address of **usrIsr2** to P1 as shown below:

IMPORT void usrIsr2 (void);

intVecSet (vector, (FUNCPTR) usrIsr2);

In conclusion, you have to guarantee that the entire ISR does not access to any virtual address space if you set the corresponding intPrioTable[] entry to **NULL**.

**NOTE SIMNT** This routine does nothing.

RETURNS N/A

ERRNO Not Available

SEE ALSO intArchLib, intVecBaseSet(), intVecGet()

intVacSat2()

	Int vecSet2()	
NAME	intVecSet2() – set a CPU vector, gate type(int/trap), and selector (x86)	
SYNOPSIS	<pre>void intVecSet2   (   FUNCPTR * vector, /* vector offset */ FUNCPTR function, /* address to place in vector */   int idtGate, /* IDT_TRAP_GATE or IDT_INT_GATE */   int idtSelector /* sysCsExc or sysCsInt */ )</pre>	
DESCRIPTION	This routine attaches an exception handler to a specified vector, with the type of the gate and the selector of the gate. The vector is specified as an offset into the CPU's vector table. This vector table starts, by default, at address 0. However, the vector table may be set to start at any address with <b>intVecBaseSet()</b> . The vector table is set up in <b>usrInit()</b> .	
RETURNS	N/A	
ERRNO	Not Available	
SEE ALSO	intArchLib, intVecBaseSet( ), intVecGet( ), intVecSet( ), intVecGet2( )	

### intVecTableWriteProtect()

NAME intVecTableWriteProtect() – write-protect exception vector table (MC680x0, x86, ARM, SimSolaris, SimNT) SYNOPSIS STATUS intVecTableWriteProtect (void) DESCRIPTION If the unbundled Memory Management Unit (MMU) support package (VxVMI) is present, this routine write-protects the exception vector table to protect it from being accidentally corrupted. Note that other data structures contained in the page will also be write-protected. In the default VxWorks configuration, the exception vector table is located at location 0 in memory. Write-protecting this affects the backplane anchor, boot configuration information, and potentially the text segment (assuming the default text location of 0x1000.) All code that manipulates these structures has been modified to write-enable memory for the duration of the operation. If you select a different address for the exception vector table, be sure it resides in a page separate from other writable data structures.

#### NOTE SIMSOLARIS, SIMNT

This routine always returns **ERROR** on simulators.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted.

**RETURNS** OK, or ERROR if memory cannot be write-protected.

ERRNO S\_intLib\_VEC\_TABLE\_WP\_UNAVAILABLE

SEE ALSO intArchLib

### ioGlobalStdGet( )

NAME	${f ioGlobalStdGet}($ ) – get the file descriptor for global input/output/error
SYNOPSIS	<pre>int ioGlobalStdGet   (    int stdFd /* std input (0), output (1), or error (2) */ )</pre>
DESCRIPTION	This routine returns the current underlying file descriptor for global standard input, output, and error.
RETURNS	The underlying global file descriptor, or <b>ERROR</b> if <i>stdFd</i> is not 0, 1, or 2.
ERRNO	N/A.
SEE ALSO	ioLib, ioGlobalStdSet( ), ioTaskStdGet( )

### ioGlobalStdSet()

NAME	<b>ioGlobalStdSet()</b> – set file descriptor for global input/output/error
SYNOPSIS	<pre>STATUS ioGlobalStdSet   (     int stdFd, /* std input (0), output (1), or error (2) */     int newFd /* new underlying file descriptor */   )</pre>

VxWorks Kernel API Reference, 6.6 ioHelp()

DESCRIPTION	This routine changes the assignment of a specified global standard file descriptor <i>stdFd</i> (0, 1, or, 2) to the specified underlying file descriptor <i>newFd</i> . <i>newFd</i> should be a file descriptor open to the desired device or file. All tasks will use this new assignment when doing I/O to <i>stdFd</i> , unless they have specified a task-specific standard file descriptor (see <b>ioTaskStdSet( )</b> ). If <i>stdFd</i> is not 0, 1, or 2, this routine has no effect.
RETURNS	<b>OK</b> , or <b>ERROR</b> if input data is not valid.
ERRNO	EBADF The newFd does not represent a valid open file. EINVAL The stdFd value is not valid.
SEE ALSO	ioLib, ioGlobalStdGet( ), ioTaskStdSet( )

# ioHelp()

NAME	<b>ioHelp()</b> – print a synopsis of I/O utility functions
SYNOPSIS	void ioHelp (void)
DESCRIPTION	This function prints out synopsis for the I/O and File System utility functions.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	usrFsLib, the VxWorks programmer guides.

# ioTaskStdGet( )

NAME	ioTaskStdGet( ) – get the file descriptor for task standard input/output/error
SYNOPSIS	<pre>int ioTaskStdGet   (    int taskId, /* ID of desired task (0 = self) */    int stdFd /* std input (0), output (1), or error (2) */   )</pre>

DESCRIPTION	This routine returns the current underlying file descriptor for task-specific standard input, output, and error.
RETURNS	The underlying file descriptor, or <b>ERROR</b> if <i>stdFd</i> is not 0, 1, or 2, or the routine is called at interrupt level.
ERRNO	N/A.
SEE ALSO	ioLib, ioGlobalStdGet( ), ioTaskStdSet( )

## ioTaskStdSet( )

NAME	<b>ioTaskStdSet()</b> – set the file descriptor for task standard input/output/error
SYNOPSIS	<pre>STATUS ioTaskStdSet   (    int taskId, /* task whose std fd is to be set (0 = self) */    int stdFd, /* std input (0), output (1), or error (2) */    int newFd /* new underlying file descriptor */   )</pre>
DESCRIPTION	This routine changes the assignment of a specified task-specific standard file descriptor $stdFd$ (0, 1, or, 2) to the specified underlying file descriptor $newFd$ . $newFd$ should be a file descriptor open to the desired device or file. The calling task will use this new assignment when doing I/O to $stdFd$ , instead of the system-wide global assignment which is used by default. If $stdFd$ is not 0, 1, or 2, this routine has no effect.
NOTE	This routine has no effect if it is called at interrupt level.
RETURNS	<b>OK</b> , or <b>ERROR</b> if input data is not valid.
ERRNO	EBADF The newFd does not represent a valid open file. EINVAL The stdFd or taskId values are not valid.
SEE ALSO	ioLib, ioGlobalStdGet( ), ioTaskStdGet( )

# ioctl()

NAME	ioctl() – perform an I/O control function
SYNOPSIS	<pre>int ioctl   (    int fd, /* file descriptor */    int function, /* function code */  )</pre>
DESCRIPTION	This routine performs an I/O control function on a device. The control functions used by VxWorks device drivers are defined in the header file <b>ioLib.h</b> . Most requests are passed on to the driver for handling. Since the availability of <b>ioctl()</b> functions is driver-specific, these functions are discussed separately in <b>tyLib</b> , <b>pipeDrv</b> , <b>nfsDrv</b> , <b>dosFsLib</b> , and <b>rawFsLib</b> .
	The following example renames the file or directory to the string "newname":
	<pre>ioctl (fd, FIORENAME, "newname");</pre>
	Note that the function FIOGETNAME is handled by the I/O interface level and is not passed on to the device driver itself. Thus this function code value should not be used by customer-written drivers. This ioctl code is convenient for checking the validity of any <i>fd</i> number. The following example shows how to quickly validate a given <i>fd</i> number as being valid. if (ioctl (fd, FIOGETNAME, NULL) == ERROR) { /* fd is not valid */
RETURNS	The return value of the driver, or <b>ERROR</b> if the file descriptor does not exist.
ERRNO	EBADF Bad file descriptor number.
	S_ioLib_UNKNOWN_REQUEST (ENOSYS) Device driver does not support the ioctl command.
	ENXIO Device and its driver are removed. <b>close( )</b> should be called to release this file descriptor.
	Other Other errors reported by device driver.
SEE ALSO	ioLib, tyLib, pipeDrv, nfsDrv, dosFsLib, rawFsLib, the VxWorks programmer guides.

# iosDevAdd()

NAME	iosDevAdd() – add a device to the kernel I/O system
SYNOPSIS	<pre>STATUS iosDevAdd   (     DEV_HDR *pDevHdr, /* pointer to device's structure */     const char *name, /* name of device */     int drvnum /* no. of servicing driver, */</pre>
DESCRIPTION	This routine adds a device to the I/O system device list, making the device available for subsequent <b>open()</b> and <b>creat()</b> calls.
	The parameter <i>pDevHdr</i> is a pointer to a device header, <b>DEV_HDR</b> (defined in <b>ioLib.h</b> ), which is used as the node in the device list. Usually this is the first item in a larger device structure for the specific device type. The parameters <i>name</i> and <i>drvnum</i> are entered in <i>pDevHdr</i> .
RETURNS	<b>OK</b> , or <b>ERROR</b> if there is already a device with the specified name.
ERRNO	S_iosLib_DUPLICATE_DEVICE_NAME (EINVAL) Device name already in use.
	EINVAL invalid arguments
SEE ALSO	iosLib

# iosDevDelDrv()

NAME	<b>iosDevDelDrv()</b> – invoke device delete driver if reference counter reaches 0.
SYNOPSIS	int iosDevDelDrv ( DEV_HDR *pDevHdr /* pointer to device's structure */ )
DESCRIPTION	This routine invokes device delete driver if reference counter reaches 0.
	If the device was never added to the device list, unpredictable results may occur.

VxWorks Kernel API Reference, 6.6 iosDevDelete()

RETURNS	DELETE_DONE Device deleted successfully. Driver called if being set.
	<b>REFCNT_NOT_ZERO</b> Device driver reference counter is not zero, device delete driver invocation is delayed until the reference counter reaches zero. The device entry is deleted and all file descriptors open on the device are invalidated in this case.
	DELETE_ERROR Error encountered in device delete.
ERRNO	N/A.
SEE ALSO	iosLib

## iosDevDelete()

NAME	iosDevDelete() – delete a device from the kernel I/O system
SYNOPSIS	<pre>int iosDevDelete   (    DEV_HDR *pDevHdr /* pointer to device's structure */ )</pre>
DESCRIPTION	This routine deletes a device from the I/O system device list, making it unavailable to subsequent IO accesses. The driver of the device is intact. <b>iosDrvRemove()</b> will do the same as this function and remove the driver in addition.
	All file descriptors open on the device are invalidated which will fail all subsequent use other than <b>close()</b> on them. They are held, even invalidated, continuously until the holding application closes them and thus releases the resource.
	If a device delete callback function is installed by <b>iosDevDelCallback()</b> , it will be called if the device driver reference counter is zero. Otherwise the callback invocation is delayed.
	If the device was never added to the device list, unpredictable results may occur.
RETURNS	<b>OK (DELETE_DONE)</b> Device deleted successfully.
	<b>REFCNT_NOT_ZERO</b> Device deleted successfully. This code is returned only when the device delete callback function is installed by <b>iosDevDelCallback()</b> . When the device driver reference counter is not zero, the callback invocation is delayed until the reference counter reaches zero when the last device driver returns. However the device entry is already deleted and all file descriptors open on the device are invalidated in this case.

	ERROR (DELETE_ERROR) Error encountered in device delete.
ERRNO	EINVAL Invalid argument. Device already deleted, or not installed, etc.
SEE ALSO	iosLib

# iosDevFind()

NAME	<b>iosDevFind( )</b> – find an I/O device in the kernel device list	
SYNOPSIS	DEV_HDR *iosDevFind ( const char *name, /* name of the device */ const char *(*pNameTail) /* where to put ptr to tail of name */ )	
DESCRIPTION	This routine searches the device list for a device whose name matches the first portion of <i>name</i> . If a device is found, <b>iosDevFind()</b> sets the character pointer pointed to by <i>pNameTail</i> to point to the first character in <i>name</i> , following the portion which matched the device name. It then returns a pointer to the device. If the routine fails, it returns a pointer to the default device (that is, the device where the current working directory is mounted) and sets <i>pNameTail</i> to point to the beginning of <i>name</i> . If there is no default device, <b>iosDevFind()</b> returns <b>NULL</b> .	
RETURNS	A pointer to the device header, or NULL if the device is not found.	
ERRNO	EINVAL Invalid arguments. S_iosLib_DEVICE_NOT_FOUND (ENODEV)	
SEE ALSO	No device found.	

## iosDevShow()

NAME	iosDevShow() – display the list of devices in the system
SYNOPSIS	void iosDevShow (void)
DESCRIPTION	This routine displays a list of all devices in the device list.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	<b>iosShow</b> , <b>devs()</b> , <b>windsh</b> , the VxWorks programmer guides, and the IDE and host tools guides.

## iosDrvInstall()

NAME	iosDrvInstall() – install a kernel I/O driver
SYNOPSIS	<pre>int iosDrvInstall    (    FUNCPTR pCreate, /* pointer to driver create function */    FUNCPTR pRemove, /* pointer to driver remove function */    FUNCPTR pOpen, /* pointer to driver open function */    FUNCPTR pClose, /* pointer to driver close function */    FUNCPTR pRead, /* pointer to driver read function */    FUNCPTR pWrite, /* pointer to driver write function */    FUNCPTR pIoctl /* pointer to driver ioctl function */    )</pre>
DESCRIPTION	This routine should be called once by each I/O driver. It hooks up the various I/O service calls to the driver service routines, assigns the driver a number, and adds the driver to the driver table.
RETURNS	The driver number of the new driver, or <b>ERROR</b> if there is no room for the driver.
ERRNO	S_iosLib_DRIVER_GLUT (ENOMEM) No memory available for data structures.
SEE ALSO	iosLib

# iosDrvRemove()

NAME	iosDrvRemove() – remove a kernel I/O driver
SYNOPSIS	STATUS iosDrvRemove ( int drvnum, /* no. of driver to remove, */ /* returned by iosDrvInstall() */ BOOL forceClose /* if TRUE, force closure of open files */ )
DESCRIPTION	This routine removes an I/O driver (added by <b>iosDrvInstall()</b> ) from the driver table and all device header entries which access the driver.
	The parameter <i>dronum</i> is an indicator of driver to be removed that is the number returned by <b>iosDrvInstall()</b> . If <i>forceClose</i> is true, all open file descriptors on this device will be closed even they are not closed by applications. This is not recommended. If <i>forceClose</i> is not true, file descriptors will be invalidated which will fail all IO other than <b>close()</b> on them. This is the graceful behavior in the driver removal case.
	If a device delete callback function is installed to a device by <b>iosDevDelCallback()</b> , it will be called if the device driver reference counter is zero when the device entry is deleted. Otherwise the callback invocation is delayed.
RETURNS	OK Driver and Device entries removed successfully.
	REFCNT_NOT_ZERO Driver and Device entries are removed successfully as returning OK case. This code is returned only when the device delete callback function is installed by <b>iosDevDelCallback()</b> . When the device driver reference counter is not zero, the callback invocation is delayed until the reference counter reaches zero when the last device driver returns.
	ERROR Error encountered in driver & device delete.
ERRNO	EINVAL invalid arguments
SEE ALSO	iosLib, iosDrvInstall(), iosDevDelete()

## iosDrvShow()

NAME	iosDrvShow() – display a list of system drivers
SYNOPSIS	void iosDrvShow (void)
DESCRIPTION	This routine displays a list of all drivers in the driver list. It now includes the null driver, which was previously omitted from the list.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	iosShow, windsh, the VxWorks programmer guides, and the IDE and host tools guides.

## iosFdEntryGet()

NAME	<pre>iosFdEntryGet() – get an unused FD_ENTRY from the pool</pre>
SYNOPSIS	FD_ENTRY* iosFdEntryGet (void)
DESCRIPTION	Returns an unused FD_ENTRY pointer, or NULL.
RETURNS	Returns a pointer to the item, or NULL if none is available.
ERRNO	ENFILE Too many open files, no internal structures available.
SEE ALSO	iosLib

iosFdEntryReturn()

```
NAME iosFdEntryReturn() - return an FD_ENTRY to the pool
SYNOPSIS STATUS iosFdEntryReturn
(
FD_ENTRY * pFdEntry /* entry to be returned to pool */
)
```
**DESCRIPTION** The FD\_ENTRY argument is returned to the FD\_ENTRY pool.

**RETURNS OK**, always.

ERRNO N/A.

SEE ALSO iosLib

### iosFdMaxFiles()

NAME	iosFdMaxFiles() – return maximum files for current RTP
SYNOPSIS	size_t iosFdMaxFiles (void)
DESCRIPTION	Returns the maximum number of open files for the current RTP.
RETURNS	The maximum number of files for the current RTP. The highest valid file descriptor number will be one less than this value.
ERRNO	N/A.
SEE ALSO	iosLib

### iosFdShow()

NAME	iosFdShow() – display a list of file descriptor names in the system
SYNOPSIS	void iosFdShow (void)
DESCRIPTION	This routine displays a list of all file descriptors in the system.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	<b>iosShow</b> , <b>ioctl()</b> , <b>windsh</b> , the VxWorks programmer guides, and the IDE and host tools guides.

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## iosInit()

NAME	iosInit() – initialize the kernel I/O system
SYNOPSIS	<pre>STATUS iosInit   (    int max_drivers, /* maximum number of drivers allowed */    int max_files, /* max number of files allowed open at once */    const char* nullDevName /* name of the null device (bit bucket) */   )</pre>
DESCRIPTION	This routine initializes the kernel I/O system. It must be called before any other I/O system routine.
RETURNS	OK, or ERROR if memory is insufficient.
ERRNO	N/A.
SEE ALSO	iosLib

# iosRtpFdShow()

NAME	<b>iosRtpFdShow()</b> – show the per-RTP <i>fd</i> table
SYNOPSIS	STATUS iosRtpFdShow ( RTP_ID rtpId )
DESCRIPTION	Primarily a debugging aid, this routine displays the <b>FD_ENTRY</b> pointers for all open file descriptors in a specified RTP.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the argument is invalid.
ERRNO	N/A
SEE ALSO	iosShow

## iosShowInit()

NAME	<b>iosShowInit()</b> – initialize the I/O system show facility
SYNOPSIS	void iosShowInit (void)
DESCRIPTION	This routine links the I/O system show facility into the VxWorks system. It is called automatically when INCLUDE_SHOW_ROUTINES is defined in <b>configAll.h</b> .
RETURNS	N/A
ERRNO	N/A
SEE ALSO	iosShow

## irint()

NAME	irint() – convert a double-precision value to an integer
SYNOPSIS	<pre>int irint   (    double x /* argument */ )</pre>
DESCRIPTION	This routine converts a double-precision value <i>x</i> to an integer using the selected IEEE rounding direction.
CAVEAT	The rounding direction is not pre-selectable and is fixed for round-to-the-nearest.
RETURNS	The integer representation of <i>x</i> .
ERRNO	Not Available
SEE ALSO	mathALib

## irintf()

NAME	irintf() – convert a single-precision value to an integer
SYNOPSIS	<pre>int irintf   (    float x /* argument */   )</pre>
DESCRIPTION	This routine converts a single-precision value $x$ to an integer using the selected IEEE rounding direction.
CAVEAT	The rounding direction is not pre-selectable and is fixed as round-to-the-nearest.
RETURNS	The integer representation of <i>x</i> .
ERRNO	Not Available
SEE ALSO	mathALib

## iround()

NAME	<b>iround()</b> – round a number to the nearest integer
SYNOPSIS	<pre>int iround   (    double x /* argument */   )</pre>
DESCRIPTION	This routine rounds a double-precision value $x$ to the nearest integer value.
NOTE	If $x$ is spaced evenly between two integers, it returns the even integer.
RETURNS	The integer nearest to <i>x</i> .
ERRNO	Not Available
SEE ALSO	mathALib

## iroundf()

NAME	iroundf() – round a number to the nearest integer
SYNOPSIS	<pre>int iroundf   (    float x /* argument */   )</pre>
DESCRIPTION	This routine rounds a single-precision value $x$ to the nearest integer value.
NOTE	If $x$ is spaced evenly between two integers, the even integer is returned.
RETURNS	The integer nearest to <i>x</i> .
ERRNO	Not Available
SEE ALSO	mathALib

## isatty()

NAME	<b>isatty()</b> – return whether the underlying driver is a <i>tty</i> device
SYNOPSIS	BOOL isatty ( int fd /* file descriptor to check */ )
DESCRIPTION	This routine simply invokes the ${\bf ioctl}($ ) function FIOISATTY on the specified file descriptor.
RETURNS	TRUE, or FALSE if the driver does not indicate a <i>tty</i> device.
ERRNO	See ioctl().
SEE ALSO	ioLib

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### isrCreate()

isrCreate() – create an ISR object NAME SYNOPSIS ISR\_ID isrCreate ( char \* name, /\* name of ISR object UINT isrTag, /\* interrupt identifier \*/ \*/ FUNCPTR handlerRtn, /\* handler routine \*/ int parameter, /\* parameter to handler routine \*/ \*/ UINT options /\* not used, must be set to 0 ) DESCRIPTION This routine creates an ISR object and initializes it with the values passed as arguments. It is meant to be used by code which connects routines to interrupt vectors by means other than calling intConnect() since that routine implicitly creates an ISR object. The name argument is the actual name of the ISR object. Any null terminated ASCII string is acceptable. The name is merely used for debugging purposes. Should *name* be NULL, a name is created using a "isrN" pattern where "N" is a decimal number representing the Nth ISR object for which the system had to choose a name at creation. Numbering starts at N =1. ISR objects implicitly created (via **intConnect()**) are named in that manner. The *isrTag* is the means by which an ISR object is associated to an interrupt source. Typically this would be used to store the vector associated with the interrupt but this is not a requirement. The caller can use other identification schemes. This library does not make use of the *isrTag* other than for information providing/displaying purposes. ISR objects implicitly created due to a call to intConnect() have their isrTag set to the value of the vector passed to intConnect(). The *handlerRtn* is the interrupt service routine that gets invoked when the associated interrupt occurs. The invocation includes passing *arg* to the *handlerRtn*. The prototype for the *handlerRtn* should be: STATUS handlerRtn (int arg) The *handlerRtn* and *arg* arguments for ISR objects that are implicitly created map directly to the *routine* and *parameter* arguments of the **intConnect()** call. The *options* argument is not use and must be set to zero by the caller. The **CODING EXAMPLE** section found in the library description illustrates how to use this routine. The ISR\_ID of the ISR object created or NULL if creation failed RETURNS S\_isrLib\_ISR\_NOT\_INIT ERRNO ISR library must first be initialized

### S\_intLib\_NOT\_ISR\_CALLABLE this routine must not be called from an ISR

S\_isrLib\_INVALID\_PARAM options is not valid

SEE ALSO isrLib, isrDelete()

### isrDelete()

NAME	isrDelete() – delete an ISR object
SYNOPSIS	STATUS isrDelete ( ISR_ID isrId /* ID of ISR object to delete */ )
DESCRIPTION	This routine destroys the ISR object specified by <i>isrId</i> and de-allocates the memory used by the object. This routine complements <b>isrCreate()</b> and is meant to be used by code which disconnects routines from interrupt vectors by means other than calling <b>intDisconnect()</b> since that routine implicitly deletes an ISR object.
	<b>WARNINGS</b> Before deleting an ISR object, one must ensure it no longer plays a role in interrupt processing. That is, the ISR object must be disconnected from its interrupt source before being deleted.
	An implicitly created ISR object, one created via <b>intConnect()</b> that is, must never be explicitly deleted using <b>isrDelete()</b> .
RETURNS	OK, ERROR if the ISR object could not be deleted.
ERRNO	S_objLib_OBJ_ID_ERROR isrId is not a valid ISR object
	S_intLib_NOT_ISR_CALLABLE this routine must not be called from an ISR
SEE ALSO	isrLib, isrCreate()

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<pre>isrIdSelf()</pre>
------------------------

isrIdSelf() – get the ISR ID of the currently running ISR NAME

SYNOPSIS ISR\_ID isrIdSelf (void)

This routine returns the ISR ID of the calling ISR. The ISR ID is NULL if the routine is called DESCRIPTION at task level. Calling this routine from a watchdog routine either returns the ISR ID of the system clock ISR or NULL. The latter is returned in cases where the processing of watchdog routines is deferred. This deferral only takes place when the system clock interrupts the kernel while it is in a critical section.

RETURNS	The ISR ID of the calling ISR.	Can be NULL.
---------	--------------------------------	--------------

ERRNO	N/A

isrLib SEE ALSO

### isrInfoGet( )

NAME	isrInfoGet() – get i	information about an	ISR object	
SYNOPSIS	STATUS isrInfoGet ( ISR_ID i: ISR_DESC * p: )	t srId, /* ISR ob <u></u> IsrDesc /* pointer	ject ID */ r to ISR description struct */	
DESCRIPTION	This routine retriev ISR_DESC structure	es information regard pointed to by <i>pIsrDe</i>	ling <i>idrId</i> . The information is return <i>sc</i> .	ed in the
	ISR_DESC is define	d as:		
	typedef struct is { ISR_ID char * UINT UINT UINT64 int FUNCPTR int	<pre>sr_desc isrId; name; isrTag; count; serviceCount; cpuTime; options; handlerRtn; arg;</pre>	<pre>/* ISR_ID /* name /* interrupt tag /* # of times this ISR has beer /* # of times this ISR has retu /* cpu time spent in ISR /* ISR object options /* pointer to handler routine /* parameter to be passed to retuine</pre>	*/ */ 1 invoked */ 1 rned OK */ */ */ putine */

} ISR\_DESC;

Note that because ISR time stamping functionality is not implemented, the *cpuTime* member of the structure is always 0.

**RETURNS** OK, ERROR if *isrld* if not valid or *plsrDesc* is NULL.

ERRNO S\_objLib\_OBJ\_ID\_ERROR isrId is not a valid ISR object

> S\_isrLib\_INVALID\_PARAM pIsrDesc is not valid

SEE ALSO isrLib

### isrInvoke()

NAME isrInvoke() – invoke the handler routine of an ISR object SYNOPSIS STATUS isrInvoke ( ISR\_ID isrId ) DESCRIPTION This routine invokes the handler routine of *isrld* as specified by the *handlerRtn* argument provided when the ISR object was created. For implicitly created ISR objects, which are created as a result of a call to intConnect(), isrInvoke() is automatically called when the interrupt associated with the vector specified in intConnect() occurs. Then, isrInvoke() takes care of calling the user supplied handler routine for that interrupt. That is, the *routine* argument provided when **intConnect()** was called. For explicitly created ISR objects, this routine is meant to be installed as an interrupt handling routine by the creator of the ISR object such that when the associated interrupt occurs, it is automatically dispatched. The **CODING EXAMPLE** section of the library description illustrates how to use this routine in that manner. Routine isrInvoke() is therefore meant to be an intermediate routine between the hardware interrupt and the user provided handler for the interrupt. In essence, it is an instrumentation routine and this allows it to keep statistics on *isrId* such as the number of times its handler ran. A WINDVIEW events is also generated as a result of **isrInvoke()** running. OK, ERROR if *isrld* is not valid RETURNS

VxWorks Kernel API Reference, 6.6 isrShow()

ERRNO S\_objLib\_OBJ\_ID\_ERROR isrId is not a valid ISR object

SEE ALSO isrLib, isrCreate()

### isrShow()

isrShow() - show information about an ISR object NAME SYNOPSIS STATUS isrShow ( ISR\_ID isrId /\* ID of ISR object or NULL for all \*/ ) DESCRIPTION This routine displays information related to an ISR object or all ISR objects. EXAMPLE A summary of a single ISR object is displayed as follows: -> isrShow myIsrId : 0x20f70568 ID Name : isrTestObject2 Interrupt Tag : 1 Count : 0 Service Count:Options:Under Routine:Argument:0x0 A summary of all ISR objects is displayed as follows: -> isrShow ISR ID Name Tag Counts HandlerRtn \_\_\_\_\_ 0x20ffb178 isr1 4 0x20ffb100 isr2 0 0x20f89610 isr3 3 winIntRcv sysNvRamSe > +0x190 0/0 31/30 0/0 3 wdbPipePkt > +0x1ba4 0x20f705e8 isrTest0 > 1 0/0 isrHandlerRtn 0x20f70568 isrTest0 > 1 0/0 isrHandlerRtn The format for the **Counts** column is count/serviceCount. OK or ERROR RETURNS ERRNO S\_objLib\_OBJ\_ID\_ERROR isrId is not a valid ISR object isrShow SEE ALSO

### kernelCpuEnable() kernelCpuEnable() – enable a CPU NAME SYNOPSIS STATUS kernelCpuEnable unsigned int cpuToEnable /\* CPU to enable \*/ ) DESCRIPTION This routine enables the CPU whose index matches the *cpuToEnable* argument. The value must be between 1 and N - 1, where N is the number of CPUs configured in the system. This figure can be obtained by calling vxCpuConfiguredGet(). If successful, the call returns OK. Once a CPU is enabled, it starts dispatching tasks as per the scheduling algorithm. Furthermore, a subsequent call to vxCpuEnabledGet() will include the specified CPU in the set of enabled CPUs. By default VxWorks SMP enables all CPUs configured in the system at boot time. Therefore calling this routine is not required unless the default behaviour is overridden. This routine returns ERROR if the specified CPU is already enabled or is outside the range of configured CPUs. CPU 0 can never be enabled using this routine as it is the bootstrap CPU. If this routine is unable to enable the specified CPU, before the timeout (VX\_ENABLE\_CPU\_TIMEOUT) expires, the routine will return ERROR. This routine always returns ERROR for VxWorks UP. SMP CONSIDERATIONS This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions. OK or ERROR if cpu index is invalid or VX\_ENABLE\_CPU\_TIMEOUT has been reached RETURNS N/A FRRNO SEE ALSO kernelLib, vxCpuConfiguredGet( ), vxCpuEnabledGet( )

### kernelInit()

NAME kernelInit() – initialize the kernel

SYNOPSIS \_\_WRS\_FUNC\_NORETURN void kernelInit

VxWorks Kernel API Reference, 6.6 kernellsCpuldle()

	( UINT32 sanity, /* must match _KERNEL_INIT_PARAMS_VN_AND_SIZE */ _KERNEL_INIT_PARAMS *pParams /* parameters */ )
DESCRIPTION	This routine initializes and starts the kernel. It should be called only once. The parameter <i>rootRtn</i> specifies the entry point of the user's start-up code that subsequently initializes system facilities (i.e., the I/O system, network). Typically, <i>rootRtn</i> is set to <b>usrRoot()</b> .
	Interrupts are enabled for the first time after <b>kernelInit()</b> exits. VxWorks will not exceed the specified interrupt lock-out level during any of its brief uses of interrupt locking as a means of mutual exclusion.
	The system memory partition is initialized by <b>kernelInit()</b> with the size set by <i>pMemPoolStart</i> and <i>pMemPoolEnd</i> . Architectures that support a separate interrupt stack allocate a portion of memory for this purpose, of <i>intStackSize</i> bytes starting at <i>pMemPoolStart</i> .
NOTE SH77XX	The interrupt stack is emulated by software, and it has to be located in a fixed physical address space (P1 or P2) if the on-chip MMU is enabled. If <i>pMemPoolStart</i> is in a logical address space (P0 or P3), the interrupt stack area is reserved on the same logical address space. The actual interrupt stack is relocated to a fixed physical space pointed by VBR.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	kernelLib, intLockLevelSet()

# kernelIsCpuIdle()

NAME	<b>kernelIsCpuIdle( )</b> – determine whether the specified CPU is idle
SYNOPSIS	BOOL kernelIsCpuIdle ( unsigned int cpu /* CPU to query status of */ )
DESCRIPTION	This routine returns TRUE if the specified CPU is idle.
	For the uniprocessor VxWorks environment, this routine returns <b>TRUE</b> if the kernel is spinning in the idle loop, i.e. the kernel is not executing any tasks.
	For SMP, this routine returns <b>TRUE</b> if the specified CPU is executing the idle task. If the specified CPU is not enabled, the CPU is considered to be idle and <b>TRUE</b> is returned. If the

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specified CPU is the one executing the calling task this routine returns FALSE since a CPU cannot be idle and executing a task other than the idle task. When called from ISR this routine returns TRUE if the interrupted task is the idle task. Otherwise FALSE is returned.

This routine is meant to be a debugging and system-monitoring tool.

**RETURNS** TRUE if the specified CPU is idle, FALSE otherwise

ERRNO N/A

SEE ALSO kernelLib, kernelSystemIsIdle()

### kernelIsSystemIdle()

NAME kernelIsSystemIdle() – determine whether all enabled processors are idle

- SYNOPSIS BOOL kernelIsSystemIdle (void)
- **DESCRIPTION** For the uniprocessor VxWorks environment, this routine returns **TRUE** if the kernel is spinning in the idle loop, i.e. the kernel is not executing any tasks.

For SMP, this routine returns **TRUE** if all enabled processors are idle, or more specifically executing their respective idle task. CPUs that are not enabled are considered to be idle. Routine **vxCpuEnabledGet()** can be used to determine the enabled CPUs in the system.

This routine is meant to be a debugging and system-monitoring tool.

- **RETURNS** TRUE if all CPUs are idle, FALSE otherwise
- ERRNO N/A
- SEE ALSO kernelLib, kernelIsCpuIdle(), vxCpuEnabledGet(), kernelCpuEnable()

### kernelRoundRobinInstall()

NAME kernelRoundRobinInstall() – install VxWorks Round Robin implementation

SYNOPSIS STATUS kernelRoundRobinInstall(void)

VxWorks Kernel API Reference, 6.6 kernelTimeSlice()

**DESCRIPTION** This routine allows user custom schedulers to take advantage of the vxWorks implementation of the round robin scheduling policy. This routine should only be used if the component INCLUDE\_CUSTOM\_SCHEDULER is configured and the user wants to take advantage of the VxWorks round robin policy. Below is an example of its usage:

```
usrCustomSchedulerInit (void)
                  {
                 . . .
                 tickAnnounceHook (usrTickFunc); /* register custom hook func */
                 . . .
                  }
               usrTickFunc (int tid)
                  {
                  . . .
                  if (_func_kernelRoundRobinHook)
                     _func_kernelRoundRobinHook (tid);
                  . . .
                  }
             OK, or ERROR if _func_kernelRoundRobinHook has been initialized
RETURNS
```

ERRNO N/A

SEE ALSO kernelLib, usrCustomScheduler.c

### kernelTimeSlice()

NAME	kernelTimeSlice() – enable round-robin selection
SYNOPSIS	STATUS kernelTimeSlice ( int ticks /* time-slice in ticks or 0 to disable round-robin */ )
DESCRIPTION	This routine enables round-robin selection among tasks of same priority and sets the system time-slice to <i>ticks</i> . Round-robin scheduling is disabled by default. A time-slice of zero ticks disables round-robin scheduling.
	A hook routine, <b>kernelRoundRobinHook()</b> , is installed by this routine. <b>kernelRoundRobinHook()</b> is the routine that performs the bulk of the work to schedule tasks in a round-robin fashion. This hook is called at each tick interrupt when <b>kernelTimeSlice()</b> is called for the first time if the system is configured with <b>INCLUDE_VX_NATIVE_SCHEDULER</b> .
	For more information about round-robin scheduling, see the manual entry for kernelLib.

#### SMP CONSIDERATIONS

	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	OK, or ERROR if kernelRoundRobinHook() can not be installed
ERRNO	N/A
SEE ALSO	kernelLib

### kernelVersion()

NAME	<b>kernelVersion()</b> – return the WIND kernel revision string
SYNOPSIS	char *kernelVersion (void)
DESCRIPTION	This routine returns a string which contains the current revision of the WIND kernel. The string is of the form "WIND version x.y", where "x" corresponds to the kernel major revision, and "y" corresponds to the kernel minor revision.
RETURNS	A pointer to a string of format "WIND version x.y"
ERRNO	N/A
SEE ALSO	kernelLib

## kill()

NAME	<b>kill()</b> – send a signal to a task (POSIX)	
SYNOPSIS	<pre>int kill   (     int tid, /* task to send signal to */     int signo /* signal to send to task */ )</pre>	
DESCRIPTION	This routine sends a signal <i>signo</i> to the task specified by <i>tid</i> .	

VxWorks Kernel API Reference, 6.6 I()

**RETURNS** OK (0), or ERROR (-1) if the task ID or signal number is invalid.

ERRNO EINVAL

SEE ALSO sigLib, taskKill()

## 1()

NAME

**I()** – disassemble and display a specified number of instructions

### SYNOPSIS

- **DESCRIPTION** This routine disassembles a specified number of instructions and displays them on standard output. If the address of an instruction is entered in the system symbol table, the symbol will be displayed as a label for that instruction. Also, addresses will be displayed symbolically.

To execute, enter:

-> 1 [address [,count]]

If *address* is omitted or zero, disassembly continues from the previous address. If *count* is omitted or zero, the last specified count is used (initially 10). As with all values entered via the shell, the address may be typed symbolically.

RETURNS	N/A
ERRNO	N/A
SEE ALSO	<b>dbgLib</b> , <b>d()</b> , <i>VxWorks Kernel Programmer's Guide: Kernel Shell</i> , <i>VxWorks Command-Line Tools User's Guide 2.2: Host Shell</i>

# 10( )

NAME	10() – return the contents of register 10 (also 11-17) (SimSolaris)
SYNOPSIS	<pre>int l0   (    int taskId /* task ID, 0 means default task */   )</pre>
DESCRIPTION	This command extracts the contents of local register 10 from the TCB of a specified task. If <i>taskId</i> is omitted or 0, the current default task is assumed. Similar routines are provided for all local registers (10 - 17): <b>10( ) - 17( )</b> .
RETURNS	The contents of register 10 (or the requested register).
ERRNO	Not Available
SEE ALSO	dbgArchLib, VxWorks Programmer's Guide: Debugging

# ld()

NAME	ld() – load an object module into memory	
SYNOPSIS	<pre>MODULE_ID ld   (     int syms, /* -1, 0, or 1 */     BOOL noAbort, /* TRUE = don't abort script on error */     char * name /* name of object module, NULL = standard input */   )</pre>	
DESCRIPTION	N This command loads an object module from a file or from standard input. The object mo must be in architecture object module format (OMF). For most of the architectures, thi ELF format. External references in the module are resolved during loading. The <i>syms</i> parameter determines how symbols are loaded; possible values are:	
	<ul> <li>Add global symbols to the system symbol table.</li> <li>Add global and local symbols to the system symbol table.</li> <li>Add no symbols to the system symbol table.</li> </ul>	

	If there is an error during loading (e.g., externals undefined, too many symbols, etc.), then <b>shellScriptAbort()</b> is called to stop any script that this routine was called from. If <i>noAbort</i> is <b>TRUE</b> , errors are noted but ignored. The normal way of using <b>ld()</b> is to load all symbols ( <i>syms</i> = 1) during debugging and to load only global symbols later.
NOTE	The routine <b>ld()</b> is a <b>shell routine</b> . That is, it is designed to be used only from the shell, and not in code running on the target. In future releases, calling <b>ld()</b> directly from code may not be supported.
COMMON SYMBOLS	On the kernel shell, for the <b>ld()</b> routine only, common symbol behavior is determined by the value of the global variable <b>ldCommonMatchAll</b> . The reasoning for <b>ldCommonMatchAll</b> matches the purpose of the windsh environment variable, <b>LD_COMMON_MATCH_ALL</b> as explained below.
	If <b>ldCommonMatchAll</b> is set to <b>TRUE</b> (equivalent to windsh "LD_COMMON_MATCH_ALL=on"), the loader tries to match a common symbol with an existing one. If a symbol with the same name is already defined, the loader takes its address. Otherwise, the loader creates a new entry. If set to <b>FALSE</b> (equivalent to windsh "LD_COMMON_MATCH_ALL=off"), the loader does not try to find an existing symbol. It creates an entry for each common symbol.
EXAMPLE	The following example loads the ELF file "module" from the default file device into memory, and adds any global symbols to the symbol table: -> ld < module This example loads " <b>test.o</b> " with all symbols: -> ld 1,0, "test.o"
RETURNS	a <b>MODULE_ID</b> , or <b>NULL</b> if there are too many symbols, the object file format is invalid, or there is an error reading the file.
ERRNO	open() errnos, loadModule() errnos.
SEE ALSO	usrLib, loadLib, unld(), reld(), the VxWorks programmer guides.

## ledClose()

NAME ledClose() – discard the line-editor ID

SYNOPSIS STATUS ledClose

```
(
FAST LED_ID ledId /* ID returned by ledOpen */
)
DESCRIPTION This routine frees resources allocated by ledOpen(). The low-level input/output file
descriptors are not closed.
RETURNS OK, or ERROR if ledId is invalid.
ERRNO N/A
SEE ALSO ledLib, ledOpen()
```

### ledControl()

NAME	<b>ledControl()</b> – change the line-editor ID parameters
SYNOPSIS	<pre>void ledControl    (    FAST LED_ID ledId, /* ID returned by ledOpen */    int inFd, /* new input fd (NONE = no change) */    int outFd, /* new output fd (NONE = no change) */    int histSize /* new hist list size (NONE=no change),    (0=display)*/    )</pre>
DESCRIPTION	This routine changes the input/output file descriptor and the size of the history list.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	ledLib

### ledLibInit( )

NAME	<b>ledLibInit()</b> – initialize the line editing facilities
SYNOPSIS	STATUS ledLibInit (void)
DESCRIPTION	This routine initializes the line editing facilities. It is called once from <b>shellLibInit()</b> .

VxWorks Kernel API Reference, 6.6 ledOpen( )

OK or ERROR if there was a problem RETURNS

Not Available ERRNO

ledLib SEE ALSO

SEE ALSO

## ledOpen()

NAME	ledOpen() – create a new line-editor ID
SYNOPSIS	<pre>LED_ID ledOpen   (     int inFd,    /* low-level device input fd */     int outFd,    /* low-level device output fd */     int histSize    /* size of history list */     )</pre>
DESCRIPTION	This routine creates the ID that is used by <b>ledRead()</b> , <b>ledClose()</b> , and <b>ledControl()</b> . Storage is allocated for up to <i>histSize</i> previously read lines.
RETURNS	The line-editor ID, or <b>ERROR</b> if the routine runs out of memory.
ERRNO	N/A
SEE ALSO	ledLib, ledRead(), ledClose(), ledControl().

## ledRead()

NAME	<b>ledRead()</b> – read a line with line-editing
SYNOPSIS	<pre>int ledRead   (    LED_ID ledId, /* ID returned by ledOpen */    char * string, /* where to return line */    UINT maxBytes /* maximum number of chars to read */   )</pre>
DESCRIPTION	This routine handles line-editing and history substitutions. If the low-level input file descriptor is not in <b>OPT_LINE</b> mode, only an ordinary <b>read()</b> routine will be performed.
RETURNS	the number of characters read, or EOF.

ERRNO N/A

SEE ALSO ledLib, ledOpen()

## link()

NAME	link() – link a file
SYNOPSIS	<pre>int link   (     const char *name, /* name of file to be linked */     const char *newname /* name with which to link */   )</pre>
DESCRIPTION	This routine links the name of a file from <i>newname</i> to <i>name</i> .
RETURNS	OK, or ERROR if the file could not be opened or linked.
ERRNO	ENOENT Either name or newname is an empty string.
	EMFILE Maximum number of files already open.
	S_iosLib_DEVICE_NOT_FOUND (ENODEV) No valid device name found in path.
	others Other errors reported by device driver.

SEE ALSO fsPxLib

## lio\_listio()

NAME	<b>lio_listio( )</b> – initiate a list of asynchronous I/O requests (POSIX)
SYNOPSIS	<pre>int lio_listio    (     int mode, /* LIO_WAIT or LIO_NOWAIT */    struct alocb *const list[], /* list of operations */    int</pre>

	struct sigevent * pSig /* signal on completion */ )
DESCRIPTION	This routine submits a number of I/O operations (up to <b>AIO_LISTIO_MAX</b> ) to be performed asynchronously. <i>list</i> is a pointer to an array of <b>aiocb</b> structures that specify the AIO operations to be performed. The array is of size <i>nEnt</i> .
	The <b>aio_lio_opcode</b> field of the <b>aiocb</b> structure specifies the AIO operation to be performed. Valid entries include <b>LIO_READ</b> , <b>LIO_WRITE</b> , and <b>LIO_NOP</b> . <b>LIO_READ</b> corresponds to a call to <b>aio_read()</b> , <b>LIO_WRITE</b> corresponds to a call to <b>aio_write()</b> , and <b>LIO_NOP</b> is ignored.
	The <i>mode</i> argument can be either LIO_WAIT or LIO_NOWAIT. If <i>mode</i> is LIO_WAIT, lio_listio() does not return until all the AIO operations complete and the <i>pSig</i> argument is ignored. If <i>mode</i> is LIO_NOWAIT, the lio_listio() returns as soon as the operations are queued. In this case, if <i>pSig</i> is not NULL and the signal number indicated by <b>pSig-sigev_signo</b> is not zero, the signal <b>pSig-sigev_signo</b> is delivered when all requests have completed.
RETURNS	OK if requests queued successfully, otherwise ERROR.
ERRNO	EINVAL EAGAIN EIO
SEE ALSO	aioPxLib, aio_read(), aio_write(), aio_error(), aio_return().

### lkAddr()

 

 NAME
 IkAddr() – list symbols whose values are near a specified value

 SYNOPSIS
 void 1kAddr ( unsigned int addr /\* address around which to look \*/ )

 DESCRIPTION
 This command lists the symbols in the system symbol table that are near a specified value. The symbols that are displayed include: - symbols whose values are immediately less than the specified value - symbols with the specified value - succeeding symbols, until at least 12 symbols have been displayed This command also displays symbols that are local, i.e., symbols found in the system symbol table only because their module was loaded by Id().

 RETURNS N/A

ERRNO N/A

**SEE ALSO** usrLib, symLib, symEach(), the VxWorks programmer guides.

## lkup()

lkup() – list symbols
<pre>void lkup   (    char *substr /* substring to match */ )</pre>
This command lists all symbols in the system symbol table whose names contain the string <i>substr</i> . If <i>substr</i> is omitted or is 0, a short summary of symbol table statistics is printed. If <i>substr</i> is the empty string (""), all symbols in the table are listed.
This command also displays symbols that are local, i.e., symbols found in the system symbol table only because their module was loaded by <b>ld()</b> .
By default, <b>lkup()</b> displays 22 symbols at a time. This can be changed by modifying the global variable <b>symLkupPgSz</b> . If this variable is set to 0, <b>lkup()</b> displays all the symbols without interruption.
N/A
N/A
usrLib, symLib, symEach(), the VxWorks programmer guides.

## 11()

NAME Il() - generate a long listing of directory contents
SYNOPSIS
STATUS 11
(
const char \* dirName /\* name of directory to list \*/
)

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VxWorks Kernel API Reference, 6.6 IIr()

DESCRIPTION	This command causes a long listing of a directory's contents to be displayed. It is equivalent to:
	-> dirList 1, dirName, 1, 0
	<i>dirName</i> is a name of a directory or file, and may contain wildcards.
NOTE 1	This is a target resident function, which manipulates the target I/O system. It must be preceded with the @ letter if executed from the Host Shell (windsh), which has a built-in command of the same name that operates on the Host's I/O system.
NOTE 2	When used with <b>netDrv</b> devices (FTP or RSH), <b>ll()</b> does not give directory information. It is equivalent to an <b>ls()</b> call with no long-listing option.
RETURNS	OK or ERROR.
ERRNO	Not Available
SEE ALSO	usrFsLib, dirList(), the VxWorks programmer guides.

## llr()

NAME	<pre>llr() – do a long listing of directory and all its subdirectories contents</pre>
SYNOPSIS	STATUS llr ( const char * dirName /* name of directory to list */ )
DESCRIPTION	This command causes a long listing of a directory's contents to be displayed. It is equivalent to: -> dirList 1, dirName, 1, 0 <i>dirName</i> is a name of a directory or file, and may contain wildcards.
NOTE	When used with <b>netDrv</b> devices (FTP or RSH), <b>ll()</b> does not give directory information. It is equivalent to an <b>ls()</b> call with no long-listing option.
RETURNS	OK or ERROR.
ERRNO	Not Available
SEE ALSO	usrFsLib, dirList(), the VxWorks programmer guides.

## InPciRegister()

DETUDNO	N/Δ
DESCRIPTION	This routine registers the PCnet/PCI driver with VxBus as a child of the PCI bus type.
SYNOPSIS	void lnPciRegister(void)
	in encesister() register with the value subsystem

## loadModule()

NAME	loadModule() – load an object module into memory
SYNOPSIS	<pre>MODULE_ID loadModule   (    int fd,</pre>
DESCRIPTION	This routine loads an object module from the specified file, and places the TEXT, DATA, and BSS into memory allocated from the system memory pool (i.e. the heap).
	Calling this function is equivalent to calling <b>loadModuleAt()</b> with <b>NULL</b> for the addresses of TEXT, DATA, and BSS segments. For more details as well as the supported option flags, see the reference entry for <b>loadModuleAt()</b> .
RETURNS	The <b>MODULE_ID</b> , or <b>NULL</b> if there was a problem. <b>NULL</b> is also returned if there are undefined symbols, but the module is not unloaded. (See <b>loadModuleAt( )</b> for more details).
ERRNO	Not Available
SEE ALSO	loadLib, loadModuleAt( )

### loadModuleAt()

loadModuleAt() – load an object module into memory NAME SYNOPSIS MODULE\_ID loadModuleAt ( fd, /\* file descriptor from which to read module \*/ int options, /\* symbols to add to table, other behavior options \*/ int char \*\* ppText, /\* load TEXT segment at addr. pointed to by this \*/ /\* pointer, return load addr. via this ptr \*/ /\* load DATA segment at addr. pointed to by this \*/ char \*\* ppData, /\* pointer, return load addr. via this ptr \*/ /\* load BSS segment at addr. pointed to by this \*/ char \*\* ppBss /\* pointer, return load addr. via this ptr \*/ )

**DESCRIPTION** This routine reads an object module from a file descriptor (the *fd* parameter) and loads the TEXT, DATA and BSS segments into the system memory space. The code is properly relocated according to the relocation commands found in the ELF file.

Unresolved references to external symbols will be linked to symbols found in the system. Symbols in the object module being loaded can optionally be added to a symbol table.

It is also possible to give a specific load directive for each loadable segment, in which case the loader can use memory that the user set aside using **malloc()** or **memalign()**.

The exact loader behavior can be controlled using load flags, as described below. Load flags may be combined (binary OR) to the extent that they are not mutually exclusive.

### LINKING UNRESOLVED EXTERNALS

As the module is loaded, any unresolved external references are resolved by looking up the symbols in the system symbol table. If found, the references to those functions or data in the new module are linked to the symbols found in the system symbol table. If there is more than one possible match in the system symbol table, the symbol encountered first (which is the one added most recently) will be used.

If an unresolved external reference cannot be found in the system symbol table, then an error message ("undefined symbol: ...") is printed for the symbol, but the loading and linking continues. The partially resolved module is not removed, to enable the user to examine the module for debugging purposes. Care should be taken when executing code from the module. Executing code which contains references to unresolved symbols may have unexpected results and may corrupt the system memory.

Even though a module with unresolved symbols remains loaded after this routine returns, **NULL** will be returned to enable the caller to detect the failure programatically. To unload the module, the caller may either call the unload routine with the module name, or look up the module using the module name and then unload the module using the returned **MODULE\_ID**. See the library entries for **moduleLib** and **unldLib** for more details.

#### FULLY LINKED MODULES

The VxWorks kernel loader supports loading of fully linked modules via the **LOAD\_FULLY\_LINKED** flag. Fully linked modules do not contain any relocations and are statically linked on the host to run at a fixed address.

As the loader has no facility to allocate memory at a particular address, it is up to the user to configure his/her board so that the addresses are available for the loader to copy the module to. The information contained in the ELF headers can be used for that purpose. It can be printed by using the "-l" option flag of the GNU "readelf" tool (segment addresses, sizes and alignement requirements). Please also refer to the alignement and memory protection sections below for important information on how to allocate memory for use by the loader.

Since the host linker will store the section/segment addresses and sizes at link time in the object file, there is no need to provide them to the loader. It will pick them up automatically from the ELF headers.

Symbol tables are supported (see the symbol table dedicated section below) but not mandatory.

#### ADDING SYMBOLS TO THE SYMBOL TABLE

The symbols defined in the module to be loaded may be added to the system symbol table; this behaviour is controlled by the value of the *options* parameter:

### LOAD\_NO\_SYMBOLS

Add no symbols to the system symbol table.

#### LOAD\_LOCAL\_SYMBOLS

Add only local symbols to the system symbol table.

### LOAD\_GLOBAL\_SYMBOLS

Add only external symbols to the system symbol table.

### LOAD\_ALL\_SYMBOLS

Add both local and external symbols to the system symbol table.

When the *options* parameter is left unspecified (NULL), the loader defaults to LOAD\_GLOBAL\_SYMBOLS. If the module symbols are added to the system symbol table, modules loaded later may link against these symbols. There is no way to make a module symbols available for debugging and, at the same time, prevent other modules from linking against those symbols.

#### CODE MODULE VISIBILITY

By default any object module loaded in the system will appear as a code module and be visible with commands such as **moduleShow()**. It is however possible to hide a code module with the flag:

### HIDDEN\_MODULE

Do not display the module via **moduleShow()**.

VxWorks Kernel API Reference, 6.6 IoadModuleAt()

**RELOCATION** The relocation commands in the object module are used to relocate the TEXT, DATA, and BSS segments of the module. The ELF file sections are sorted into these three types of segments according to the ELF section flags. The location of each segment can either be specified explicitly or left unspecified, in which case memory will be allocated for the segment from the kernel heap, according to the section category.

To specify where one or more of the segments should be installed, use the parameters *ppText*, *ppData*, and *ppBss*. Each of these can have either of the following values:

#### NULL

No load address is specified, none will be returned.

#### A pointer to LD\_NO\_ADDRESS

No load address is specified; after the load is performed, LD\_NO\_ADDRESS will be replaced by the actual segment load address.

#### A pointer to an address

The load address is specified.

The *ppText*, *ppData*, and *ppBss* parameters specify where to load the TEXT, DATA, and BSS segments, respectively. Each of these parameters is a pointer to a pointer; for example, *\*\*ppText* gives the address where the TEXT segment is to begin.

Note that it is up to the user to reserve a sufficient amount of memory for each segment which address is specified. In particular, alignement requirements need to be kept in mind when reserving memory for a segment (see the alignement-dedicated section below for more information). Finally, remember that the loader will only free memory it allocates. This means the user will have to free memory he/she reserved when, for instance, a module is unloaded.

For any of the three parameters, there are two ways to request that new memory be allocated, rather than specifying the segment starting address: you can either specify the parameter itself as NULL, or you can write the constant LD\_NO\_ADDRESS in place of an address. In the second case, the **loadModuleAt()** routine replaces the LD\_NO\_ADDRESS value with the address actually used for each section (that is, it records the address at *\*ppText*, *\*ppData*, or *\*ppBss*).

The double indirection not only permits reporting the addresses actually used, but also allows you to specify loading a segment at the beginning of memory, since the following cases can be distinguished:

- Allocate memory for a segment (TEXT in this example): *ppText* == **NULL**
- Begin a section at address zero (the TEXT section, below): \*ppText == 0

Note that calling **loadModule()** is equivalent to calling **loadModuleAt()** with all three segment-address parameters set to **NULL**.

COMMON Some host compiler/linker combinations use another storage class internally called COMMON. In the C language, uninitialized global variables are eventually put in the BSS segment. However, in partially linked object modules they are flagged internally as COMMON and the static linker (host) resolves these and places them in BSS as a final step in creating a fully linked object module. However, the kernel loader is used to load partially linked object modules into the kernel, not executable modules. When the VxWorks loader encounters a variable labeled as COMMON, memory for the variable can be allocated (see below), and the variable is entered in the symbol table (if specified) at that address.

Note that most UNIX loaders have an option that forces resolution of the COMMON storage while leaving the module relocatable. For example, with typical BSD UNIX loaders, "-d" serves that purpose (in conjunction with "-r" to generate relocatable output). The GNU linker, ld, belongs to this category. With DIAB, option "-a" has to be passed to the dld linker (again, in conjunction with "-r" to generate relocatable output).

When the kernel loader encounters a variable labeled "COMMON", its behavior depends on the following flags:

### LOAD\_COMMON\_MATCH\_NONE

Allocate memory for the variable with **malloc()** and enter the variable in the target symbol table (if specified) with type **SYM\_COMM** at that address. This is the default. Note that the remark about symbol linking visibility also applies here : unless **LOAD\_NO\_SYMBOLS** is set, **LOAD\_COMMON\_MATCH\_NONE** won't prevent other modules from linking against the module symbols.

#### LOAD\_COMMON\_MATCH\_USER

The loader seeks a matching "user" symbol, i.e. symbols that have been added by dynamically loaded object modules (*vs* symbols statically present at boot time). If no matching symbol exists, it acts like **LOAD\_COMMON\_MATCH\_NONE**. If several matching symbols exist, the symbol most recently added to the target symbol table is used.

### LOAD\_COMMON\_MATCH\_ALL

The loader seeks for any matching symbol. All symbols are considered. If no matching symbol exists, then it acts like LOAD\_COMMON\_MATCH\_NONE. If several matches are found, the preference order is the same as for LOAD\_COMMON\_MATCH\_USER.

#### C++ CONSTRUCTORS SUPPORT

The loader applies the C++ strategy as defined by the C++ runtime library at the time of the load operation. If this strategy is set to "automatic", then the C++ constructors are executed. If this strategy is set to "manual", then the loader does not execute the C++ constructors (in that case, they can be called manually via the **cplusCallCtors()** API). It is possible to prevent this default behavior with the following flags:

#### LOAD\_CPLUS\_XTOR\_AUTO

When this flag is set, the loader always executes the C++ constructors associated to the module.

### LOAD\_CPLUS\_XTOR\_MANUAL

When this flag is set, the loader never executes the C++ constructors associated to the module.

VxWorks Kernel API Reference, 6.6 IoadModuleAt()

Note that if there are undefined symbols in the module, the loader will not run the C++ constructors regardless of the status of the above flags (this is to prevent erratic behavior/application crashes).

Please refer to **unldLib** for the corresponding unloader flags.

### WEAK SYMBOL HANDLING

Most ELF symbols use the standard global/local symbol binding. Some languages, however, make use of another binding called WEAK. When a WEAK symbol is encountered, the loader behavior depends on the following flags:

### LOAD\_WEAK\_MATCH\_ALL

The loader looks for an already existing global definition with the same name in the symbol table. If one can be found, the loader honors the existing definition and ignores the WEAK one. If no match can be found, the loader behaves like **LOAD\_WEAK\_MATCH\_NONE**. This is the default.

#### LOAD\_WEAK\_MATCH\_NONE

The WEAK symbol is registered in the symbol table (if specified) as a global symbol regardless of any existing definition. This behavior matches what was done by the VxWorks 5.x loader.

#### ALIGNMENT CONSIDERATIONS

Please note that memory required to load a segment is more than the bare sum of the size of sections it contains. The memory needed to load a segment is the sum of the sizes of its sections plus the padding required by each section alignment requirements (if a section has no alignment requirements, the architecture default alignment is taken as the section alignment requirement). Sections alignment requirements can be visualized using the GNU "objdump" tool (part of the binutils) with the "-h" option (alignment requirements are represented as powers of 2).

#### MEMORY PROTECTION

When memory protection is available (see **vmBaseLib** for more information), the loader will write-protect the module TEXT segment. For this to properly work, the TEXT segment size will be rounded up to the next page of memory and the TEXT segment address will be requested aligned on a page boundary. The unloader will also accordingly unprotect the TEXT segment when unloading the module.

Please note that the loader will only protect the TEXT segment (and the unloader unprotect it) if it itself allocates memory for it. If **loadModuleAt(**) is used with user-reserved memory instead (or when loading a fully linked module), it is up to the user to properly write-protect/unprotect the module TEXT segment. Here is how to do it :

Allocate aligned TEXT segment memory rounded up to the next page :

protectedTextSize = ROUND\_UP (textSize, VM\_PAGE\_SIZE\_GET());
pText = (char \*) memalign (VM\_PAGE\_SIZE\_GET(), protectedTextSize);

Write-protect the TEXT segment after it has been loaded by loadModuleAt():

### 2 Routines log10f()

	<pre>VM_STATE_SET (NULL, pText, protectedTextSize, VM_STATE_MASK_WRITABLE, VM_STATE_WRITABLE_NOT);</pre>
	Later, you may want to unprotect it :
	<pre>VM_STATE_SET (NULL, pText, protectedTextSize, VM_STATE_MASK_WRITABLE, VM_STATE_WRITABLE);</pre>
EXAMPLES	These examples are of ways to invoke the loader from C code. To use the loader from the shell, use the loader shell commands, defined in <b>usrLib</b> .
	Load a module into memory allocated by the loader:
	<pre>module_id = loadModuleAt (fd, LOAD_GLOBAL_SYMBOLS, NULL, NULL, NULL);</pre>
	Load a module into memory allocated by the loader and retrieve segment addresses:
	<pre>pText = pData = pBss = LD_NO_ADDRESS; module_id = loadModuleAt (fd, LOAD_GLOBAL_SYMBOLS, &amp;pText, &amp;pData, &amp;pBss);</pre>
	Load a module to off-board memory at a specified address:
	pText = 0x800000; /* address of TEXT segment
	<pre>*/     pData = pBss = LD_NO_ADDRESS;</pre>
	<pre>/ module_id = loadModuleAt (fd, LOAD_GLOBAL_SYMBOLS, &amp;pText, &amp;pData, &amp;pBss);</pre>
RETURNS	A <b>MODULE_ID</b> , or <b>NULL</b> if there was a problem.
ERRNO	Possible errnos set by this routine include:
	+ S_loadLib_INVALID_ARGUMENT
	For a complete description of the errnos, see the reference documentation for <b>loadLib</b> .
SEE ALSO	loadLib, unldLib

# log10f()

NAME	log10f() – compute a base-10 logarithm (ANSI)
SYNOPSIS	<pre>float log10f   (    float x /* value to compute the base-10 logarithm of */   )</pre>
DESCRIPTION	This routine returns the base-10 logarithm of $x$ in single precision.

VxWorks Kernel API Reference, 6.6 log2()

**RETURNS** The single-precision base-10 logarithm of *x*.

ERRNO Not Available

SEE ALSO mathALib

## log2()

NAME	<b>log2( )</b> – compute a base-2 logarithm
SYNOPSIS	<pre>double log2   (     double x /* value to compute the base-two logarithm of */   )</pre>
DESCRIPTION	This routine returns the base-2 logarithm of $x$ in double precision.
RETURNS	The double-precision base-2 logarithm of $x$ .
ERRNO	Not Available
SEE ALSO	mathALib

## log2f()

NAME	<b>log2f( )</b> – compute a base-2 logarithm
SYNOPSIS	<pre>float log2f   (    float x /* value to compute the base-2 logarithm of */ )</pre>
DESCRIPTION	This routine returns the base-2 logarithm of $x$ in single precision.
RETURNS	The single-precision base-2 logarithm of <i>x</i> .
ERRNO	Not Available
SEE ALSO	mathALib

### logFdAdd() logFdAdd() – add a logging file descriptor NAME SYNOPSIS STATUS logFdAdd ( int fd /\* file descriptor for additional logging device \*/ ) DESCRIPTION This routine adds to the log file descriptor list another file descriptor *fd* to which messages will be logged. The file descriptor must be a valid open file descriptor. RETURNS OK, or ERROR if the allowable number of additional logging file descriptors (5) is exceeded. Not Available ERRNO SEE ALSO logLib, logFdDelete()

## logFdDelete()

NAME	<b>logFdDelete( )</b> – delete a logging file descriptor
SYNOPSIS	STATUS logFdDelete ( int fd /* file descriptor to stop using as logging device */ )
DESCRIPTION	This routine removes from the log file descriptor list a logging file descriptor added by <b>logFdAdd()</b> . The file descriptor is not closed; but is no longer used by the logging facilities.
RETURNS	OK, or ERROR if the file descriptor was not added with logFdAdd().
ERRNO	Not Available
SEE ALSO	logLib, logFdAdd()

2

## logFdSet()

NAME	<b>logFdSet()</b> – set the primary logging file descriptor
SYNOPSIS	<pre>void logFdSet   (    int fd /* file descriptor to use as logging device */ )</pre>
DESCRIPTION	This routine changes the file descriptor where messages from <b>logMsg()</b> are written, allowing the log device to be changed from the default specified by <b>logInit()</b> . It first removes the old file descriptor (if one had been previously set) from the log file descriptor list, then adds the new <i>fd</i> .
	The old logging file descriptor is not closed or affected by this call; it is simply no longer used by the logging facilities.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	logLib, logFdAdd(), logFdDelete()

## logInit()

NAME	<b>logInit( )</b> – initialize message logging library
SYNOPSIS	STATUS logInit ( int fd, /* file descriptor to use as logging device */ int maxMsgs /* max. number of messages allowed in log queue */ )
DESCRIPTION	This routine specifies the file descriptor to be used as the logging device and the number of messages that can be in the logging queue. If more than <i>maxMsgs</i> are in the queue, they will be discarded. A message is printed to indicate lost messages.
	This routine spawns logTask(), the task-level portion of error logging.
	This routine must be called before any other routine in <b>logLib</b> . This is done by the root task, <b>usrRoot()</b> , in <b>usrConfig.c</b> .
RETURNS	OK, or ERROR if a message queue could not be created or logTask() could not be spawned.

ERRNO Not Available

SEE ALSO logLib

### logMsg()

)

NAME logMsg() – log a formatted error message

```
SYNOPSIS
```

int logMsg
 (
 char \*fmt, /\* format string for print \*/
 int arg1, /\* first of six required args for fmt \*/
 int arg2,
 int arg3,
 int arg4,
 int arg5,
 int arg6

**DESCRIPTION** This routine logs a specified message via the logging task. This routine's syntax is similar to **printf()** -- a format string is followed by arguments to format. However, the **logMsg()** routine takes a char \* rather than a const char \* and requires a fixed number of arguments (6).

The task ID of the caller is prepended to the specified message.

### SPECIAL CONSIDERATIONS

Because **logMsg()** does not actually perform the output directly to the logging streams, but instead queues the message to the logging task, **logMsg()** can be called from interrupt service routines.

However, since the arguments are interpreted by the **logTask()** at the time of actual logging, instead of at the moment when **logMsg()** is called, arguments to **logMsg()** should not be pointers to volatile entities (e.g., dynamic strings on the caller stack).

**logMsg()** checks to see whether or not it is running in interupt context. If it is, it will not block. However, if invoked from a task, it can cause the task to block.

For more detailed information about the use of logMsg(), see the manual entry for logLib.

**EXAMPLE** If the following code were executed by task 20:

```
{
name = "GRONK";
num = 123;
logMsg ("ERROR - name = %s, num = %d.\\n", name, num, 0, 0, 0, 0);
```

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SEE ALSO logLib, printf(), logTask()

## logTask()

NAME	logTask() – message-logging support task
SYNOPSIS	void logTask (void)
DESCRIPTION	This routine prints the messages logged with <b>logMsg()</b> . It waits on a message queue and prints the messages as they arrive on the file descriptor specified by <b>logInit()</b> (or a subsequent call to <b>logFdSet()</b> or <b>logFdAdd()</b> ).
	This task is spawned by logInit().
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	logLib, logMsg()

## logf()

NAME	<b>logf( )</b> – compute a natural logarithm (ANSI)
SYNOPSIS	<pre>float logf    (    float x /* value to compute the natural logarithm of */    )</pre>
DESCRIPTION	This routine returns the logarithm of $x$ in single precision.
**RETURNS** The single-precision natural logarithm of *x*.

ERRNO Not Available

SEE ALSO mathALib

#### loginDefaultEncrypt()

NAME	loginDefaultEncrypt() – default password encryption routine
SYNOPSIS	<pre>STATUS loginDefaultEncrypt   (     char * in, /* input string */     char * out /* encrypted string */   )</pre>
DESCRIPTION	This routine provides default encryption for login passwords. It employs a simple encryption algorithm. It takes as arguments a string <i>in</i> and a pointer to a buffer <i>out</i> . The encrypted string is then stored in the buffer.
	The input strings must be at least 8 characters and no more than 40 characters.
	If a more sophisticated encryption algorithm is needed, this routine can be replaced, as long as the new encryption routine retains the same declarations as the default routine. The routine <b>vxencrypt</b> in <b>host</b> / <i>hostOs</i> /bin should also be replaced by a host version of <i>encryptionRoutine</i> . For more information, see the manual entry for <b>loginEncryptInstall()</b> .
RETURNS	<b>OK</b> , or <b>ERROR</b> if the password is invalid.
ERRNO	Possible errnos set by this routine include:
	<b>S_loginLib_INVALID_PASSWORD</b> <i>in</i> string is not a valid password string.
SEE ALSO	loginLib, loginEncryptInstall(), vxencrypt

### loginEncryptInstall()

NAME loginEncryptInstall() – install an encryption routine

**SYNOPSIS** void loginEncryptInstall

VxWorks Kernel API Reference, 6.6 loginInit()

```
FUNCPTR rtn, /* function pointer to encryption routine */
                    int
                            var
                                   /* argument to the encryption routine (unused) */
                    )
DESCRIPTION
                This routine allows the user to install a custom encryption routine. The custom routine rtn
                must be of the following form:
                STATUS encryptRoutine
                       char *password,
                                                      /* string to encrypt
                                                                                 */
                       char *encryptedPassword /* resulting encryption */
                The encryption result string must remain less or equal to MAX_PASSWORD_LEN.
                When a custom encryption routine is installed, a host version of this routine must be written
                to replace the tool vxencrypt in host/hostOs/bin.
                The custom example above could be installed as follows:
EXAMPLE
                #ifdef INCLUDE SECURITY
                                                                 /* initialize login table
                    loginInit ();
                                                                                               */
                    shellLoginInstall (loginPrompt, NULL);
                                                                /* install shell security
                                                                                               */
                    loginEncryptInstall (encryptRoutine, NULL);
                                                                /* install encrypt. routine */
                #endif
                N/A
RETURNS
                N/A
ERRNO
SEE ALSO
                loginLib, loginDefaultEncrypt(), vxencrypt
```

#### loginInit()

NAME	<b>loginInit( )</b> – initialize the login table
SYNOPSIS	void loginInit (void)
DESCRIPTION	This routine must be called to initialize the login data structure used by routines throughout this module. If the configuration macro INCLUDE_SECURITY is defined, it is called by usrRoot() in usrConfig.c, before any other routines in this module.
RETURNS	N/A
ERRNO	N/A

SEE ALSO loginLib

# loginPrompt()

NAME	loginPrompt() – display a login prompt and validate a user entry
SYNOPSIS	STATUS loginPrompt ( char * userName /* user name, ask if NULL or not provided */ )
DESCRIPTION	This routine displays a login prompt and validates a user entry. If both user name and password match with an entry in the login table, the user is then given access to the VxWorks system. Otherwise, it prompts the user again.
	All control characters are disabled during authentication except CTRL-D, which will terminate the remote login session.
RETURNS	OK if the name and password are valid, or ERROR if there is an EOF or the routine times out.
ERRNO	N/A
SEE ALSO	loginLib

# loginStringSet()

NAME	<b>loginStringSet( )</b> – change the login string
SYNOPSIS	<pre>void loginStringSet   (     char * newString /* string to become new login prompt */   )</pre>
DESCRIPTION	This routine changes the login prompt string to <i>newString</i> . The maximum string length is <b>MAX_LOGIN_NAME_LEN</b> characters.
RETURNS	N/A
ERRNO	N/A

VxWorks Kernel API Reference, 6.6 IoginUserAdd()

SEE ALSO loginLib

# loginUserAdd()

NAME	<b>loginUserAdd()</b> – add a user to the login table
SYNOPSIS	<pre>STATUS loginUserAdd   (     char name[MAX_LOGIN_NAME_LEN + 1], /* user name */     char passwd[MAX_PASSWORD_LEN + 1] /* user password */     )</pre>
DESCRIPTION	This routine adds a user name and password entry to the login table. Note that what is saved in the login table is the user name and the address of <i>passwd</i> , not the actual password.
	The length of user names should not exceed MAX_LOGIN_NAME_LEN, while the length of passwords depends on the encryption routine used. For the default encryption routine, passwords should be at least 8 characters long and no more than MAX_PASSWORD_LEN characters.
	The procedure for adding a new user to login table is as follows:
	(1) Generate the encrypted password by invoking <b>vxencrypt</b> in <b>host</b> / <i>hostOs</i> /bin.
	(2) Add a user by invoking <b>loginUserAdd()</b> in the VxWorks shell with the user name and the encrypted password.
	The password of a user can be changed by first deleting the user entry, then adding the user entry again with the new encrypted password.
EXAMPLE	<pre>-&gt; loginUserAdd "peter", "RRdRd9Qbyz" value = 0 = 0x0 -&gt; loginUserAdd "robin", "bSzyydqbSb" value = 0 = 0x0 -&gt; loginUserShow User Name ========= peter robin value = 0 = 0x0</pre>
	->
RETURNS	OK, or ERROR if the user name has already been entered, or one of the arguments is NULL.
ERRNO	Possible errnos set by this routine include:

SEE ALSO	loginLib, loginUserVerify(), loginUserDelete(), vxencrypt
	S_loginLib_USER_ALREADY_EXISTS The user name <i>name</i> is already registered.
	EINVAL An invalid argument is passed to the routine.

# loginUserDelete()

NAME	<b>loginUserDelete( )</b> – delete a user entry from the login table
SYNOPSIS	<pre>STATUS loginUserDelete   (     char * name, /* user name */     char * passwd /* user password */   )</pre>
DESCRIPTION	This routine deletes an entry in the login table. Both the user name and password must be specified to remove an entry from the login table.
RETURNS	OK, or ERROR if the specified user or password is incorrect.
ERRNO	Possible errnos set by this routine include:
	EINVAL An invalid argument is passed to the routine.
	S_loginLib_UNKNOWN_USER Unknown user name <i>name</i> .
	S_loginLib_INVALID_PASSWORD Invalid password <i>passwd</i> for <i>name</i>
	Encryption routine's errnos (see loginEncryptInstall())
SEE ALSO	loginLib, loginUserAdd()

VxWorks Kernel API Reference, 6.6 IoginUserShow()

# loginUserShow()

NAME	<b>loginUserShow()</b> – display the user login table
SYNOPSIS	void loginUserShow (void)
DESCRIPTION	This routine displays valid user names.
EXAMPLE	-> loginUserShow ()
	User Name ======= peter robin value = 0 = 0x0
RETURNS	N/A
ERRNO	N/A
SEE ALSO	loginLib

# loginUserVerify()

NAME	<b>loginUserVerify()</b> – verify a user name and password in the login table
SYNOPSIS	<pre>STATUS loginUserVerify   (     char * name, /* name of user */     char * passwd /* password of user */   )</pre>
DESCRIPTION	This routine verifies a user entry in the login table.
RETURNS	OK, or ERROR if the user name or password is not found.
ERRNO	Possible errnos set by this routine include:
	EINVAL An invalid argument is passed to the routine.
	S_loginLib_UNKNOWN_USER Unknown user name <i>name</i> .

	S_loginLib_INVALID_PASSWORD Invalid password <i>passwd</i> for <i>name</i>
	Encryption routine's errnos (see loginEncryptInstall())
SEE ALSO	loginLib, loginUserAdd( )

# logout()

NAME	<b>logout( )</b> – log out of the VxWorks system
SYNOPSIS	void logout (void)
DESCRIPTION	This command logs out of the VxWorks shell. If a remote login is active (via <b>rlogin</b> or <b>telnet</b> ), it is stopped, and standard I/O is restored to the console.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	usrLib, rlogin(), telnet(), shellLogout(), the VxWorks programmer guides.

# ls( )

NAME	<b>ls()</b> – generate a brief listing of a directory
SYNOPSIS	STATUS ls ( const char * dirName, /* name of dir to list */ BOOL doLong /* switch on details */ )
DESCRIPTION	This function is simply a front-end for <b>dirList()</b> , intended for brevity and backward compatibility. It produces a list of files and directories, without details such as file size and date, and without recursion into subdirectories.
	<i>dirName</i> is a name of a directory or file, and may contain wildcards. <i>doLong</i> is provided for backward compatibility.

VxWorks Kernel API Reference, 6.6 Iseek()

**NOTE** This is a target resident function, which manipulates the target I/O system. It must be preceded with the @ letter if executed from the Host Shell (windsh), which has a built-in command of the same name that operates on the Host's I/O system.

**RETURNS** OK or ERROR.

ERRNO Not Available

**SEE ALSO usrFsLib**, **dirList()**, the VxWorks programmer guides, the, *VxWorks Command-Line Tools User's Guide*.

#### lseek()

NAME	lseek() – set a file read/write pointer
SYNOPSIS	<pre>off_t lseek   (     int fd, /* file descriptor */     off_t offset, /* new byte offset to seek to */     int whence /* relative file position */   )</pre>
DESCRIPTION	This routine sets the file read/write pointer of file <i>fd</i> to <i>offset</i> . The argument <i>whence</i> , which affects the file position pointer, has three values:
	SEEK_SET (0)set to offsetSEEK_CUR (1)set to current position plus offsetSEEK_END (2)set to the size of the file plus offset
	This routine calls <b>ioctl()</b> with functions FIOWHERE, FIONREAD, and FIOSEEK.
RETURNS	The new offset from the beginning of the file, or ERROR.
ERRNO	See ioctl().
SEE ALSO	ioLib

# lsr()

NAME	<b>lsr()</b> – list the contents of a directory and any of its subdirectories
SYNOPSIS	STATUS lsr ( const char * dirName /* name of dir to list */ )
DESCRIPTION	This function is simply a front-end for <b>dirList()</b> , intended for brevity and backward compatibility. It produces a list of files and directories, without details such as file size and date, with recursion into subdirectories. <i>dirName</i> is a name of a directory or file, and may contain wildcards.
RETURNS	OK or ERROR.
ERRNO	Not Available
SEE ALSO	usrFsLib, dirList(), the VxWorks programmer guides.

# lstAdd()

NAME	lstAdd() – add a node to the end of a list
SYNOPSIS	<pre>void lstAdd   (    LIST *pList, /* pointer to list descriptor */    NODE *pNode /* pointer to node to be added */   )</pre>
DESCRIPTION	This routine adds a specified node to the end of a specified list.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	lstLib

#### lstConcat()

NAME	lstConcat() – concatenate two lists
SYNOPSIS	<pre>void lstConcat   (   FAST LIST *pDstList, /* destination list */   FAST LIST *pAddList /* list to be added to dstList */ )</pre>
DESCRIPTION	This routine concatenates the second list to the end of the first list. The second list is left empty. Either list (or both) can be empty at the beginning of the operation.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	lstLib

# lstCount()

NAME	<b>IstCount( )</b> – report the number of nodes in a list
SYNOPSIS	int lstCount ( LIST *pList /* pointer to list descriptor */ )
DESCRIPTION	This routine returns the number of nodes in a specified list.
RETURNS	The number of nodes in the list.
ERRNO	Not Available
SEE ALSO	lstLib

#### lstDelete()

NAME	<b>lstDelete( )</b> – delete a specified node from a list
SYNOPSIS	<pre>void lstDelete   (   FAST LIST *pList, /* pointer to list descriptor */   FAST NODE *pNode /* pointer to node to be deleted */  )</pre>
DESCRIPTION	This routine deletes a specified node from a specified list.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	lstLib

### lstExtract()

NAME	lstExtract() – extract a sublist from a list
SYNOPSIS	<pre>void lstExtract   (   FAST LIST *pSrcList, /* pointer to source list */ FAST NODE *pStartNode, /* first node in sublist to be extracted */ FAST NODE *pEndNode, /* last node in sublist to be extracted */ FAST LIST *pDstList /* ptr to list where to put extracted list */ )</pre>
DESCRIPTION	This routine extracts the sublist that starts with <i>pStartNode</i> and ends with <i>pEndNode</i> from a source list. It places the extracted list in <i>pDstList</i> .
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	lstLib

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#### lstFind()

NAME	lstFind() – find a node in a list
SYNOPSIS	<pre>int lstFind   (    LIST *pList, /* list in which to search */   FAST NODE *pNode /* pointer to node to search for */   )</pre>
DESCRIPTION	This routine returns the node number of a specified node (the first node is 1).
RETURNS	The node number, or <b>ERROR</b> if the node is not found.
ERRNO	Not Available
SEE ALSO	lstLib

# lstFirst()

NAME	lstFirst() – find first node in list
SYNOPSIS	NODE *lstFirst ( LIST *pList /* pointer to list descriptor */ )
DESCRIPTION	This routine finds the first node in a linked list.
RETURNS	A pointer to the first node in a list, or <b>NULL</b> if the list is empty.
ERRNO	Not Available
SEE ALSO	lstLib

#### lstFree()

NAME	<b>lstFree( )</b> – free up a list
SYNOPSIS	<pre>void lstFree ( LIST *pList /* list for which to free all nodes */ )</pre>
DESCRIPTION	This routine turns any list into an empty list. It also frees up memory used for nodes.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	lstLib, free()

# lstGet()

NAME	<b>lstGet()</b> – delete and return the first node from a list
SYNOPSIS	NODE *lstGet ( FAST LIST *pList /* ptr to list from which to get node */ )
DESCRIPTION	This routine gets the first node from a specified list, deletes the node from the list, and returns a pointer to the node gotten.
RETURNS	A pointer to the node gotten, or NULL if the list is empty.
ERRNO	Not Available
SEE ALSO	lstLib

#### lstInit()

NAME	<b>lstInit( )</b> – initialize a list descriptor
SYNOPSIS	<pre>void lstInit   (   FAST LIST *pList /* ptr to list descriptor to be initialized */ )</pre>
DESCRIPTION	This routine initializes a specified list to an empty list.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	lstLib

#### lstInsert()

NAME	lstInsert() – insert a node in a list after a specified node
SYNOPSIS	<pre>void lstInsert   (   FAST LIST *pList, /* pointer to list descriptor */   FAST NODE *pPrev, /* pointer to node after which to insert */   FAST NODE *pNode /* pointer to node to be inserted */   )</pre>
DESCRIPTION	This routine inserts a specified node in a specified list. The new node is placed following the list node <i>pPrev</i> . If <i>pPrev</i> is <b>NULL</b> , the node is inserted at the head of the list.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	lstLib

# lstLast()

NAME	lstLast() – find the last node in a list
SYNOPSIS	NODE *lstLast ( LIST *pList /* pointer to list descriptor */ )
DESCRIPTION	This routine finds the last node in a list.
RETURNS	A pointer to the last node in the list, or <b>NULL</b> if the list is empty.
ERRNO	Not Available
SEE ALSO	lstLib

# lstNStep()

NAME	<b>lstNStep()</b> – find a list node <i>nStep</i> steps away from a specified node
SYNOPSIS	NODE *lstNStep ( FAST NODE *pNode, /* the known node */ int nStep /* number of steps away to find */ )
DESCRIPTION	This routine locates the node <i>nStep</i> steps away in either direction from a specified node. If <i>nStep</i> is positive, it steps toward the tail. If <i>nStep</i> is negative, it steps toward the head. If the number of steps is out of range, <b>NULL</b> is returned.
RETURNS	A pointer to the node <i>nStep</i> steps away, or <b>NULL</b> if the node is out of range.
ERRNO	Not Available
SEE ALSO	lstLib

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#### lstNext()

NAME	<b>lstNext()</b> – find the next node in a list
SYNOPSIS	NODE *1stNext ( NODE *pNode /* ptr to node whose successor is to be found */ )
DESCRIPTION	This routine locates the node immediately following a specified node.
RETURNS	A pointer to the next node in the list, or NULL if there is no next node.
ERRNO	Not Available
SEE ALSO	lstLib

### lstNth()

NAME	lstNth() – find the Nth node in a list
SYNOPSIS	NODE *lstNth ( FAST LIST *pList, /* pointer to list descriptor */ FAST int nodenum /* number of node to be found */ )
DESCRIPTION	This routine returns a pointer to the node specified by a number <i>nodenum</i> where the first node in the list is numbered 1. Note that the search is optimized by searching forward from the beginning if the node is closer to the head, and searching back from the end if it is closer to the tail.
RETURNS	A pointer to the Nth node, or NULL if there is no Nth node.
ERRNO	Not Available
SEE ALSO	lstLib

#### lstPrevious()

NAME	<b>IstPrevious( )</b> – find the previous node in a list
SYNOPSIS	NODE *1stPrevious ( NODE *pNode /* ptr to node whose predecessor is to be found */ )
DESCRIPTION	This routine locates the node immediately preceding the node pointed to by <i>pNode</i> .
RETURNS	A pointer to the previous node in the list, or <b>NULL</b> if there is no previous node.
ERRNO	Not Available
SEE ALSO	lstLib

#### **m( )**

NAME	<b>m()</b> – modify memory
SYNOPSIS	<pre>void m   (    void * adrs, /* address to change */    int width /* width of unit to be modified (1, 2, 4, 8) */   )</pre>
DESCRIPTION	This command prompts the user for modifications to memory in byte, short word, or long word specified by <i>width</i> , starting at the specified address. It prints each address and the current contents of that address, in turn. If <i>adrs</i> or <i>width</i> is zero or absent, it defaults to the previous value. The user can respond in one of several ways: [RETURN] Do not change this address, but continue, prompting at the next address. <i>number</i> Set the content of this address to <i>number</i> .
	. (dot) Do not change this address, and quit.
	[EOF] Do not change this address, and quit.

VxWorks Kernel API Reference, 6.6 m6845vxbRegister()

All numbers entered and displayed are in hexadecimal.
---

RETURNS	N/A
ERRNO	N/A

**SEE ALSO** usrLib, mRegs(), the VxWorks programmer guides.

#### m6845vxbRegister()

NAME	<b>m6845vxbRegister()</b> – register m6845vxb driver
SYNOPSIS	void m6845vxbRegister(void)
DESCRIPTION	This routine registers the m6845vxb driver and device recognition data with the vxBus subsystem.
NOTE	This routine is called early during system initialization, and *MUST NOT* make calls to OS facilities such as memory allocation and I/O.
RETURNS	N/A
ERRNO	
SEE ALSO	vxbM6845Vga

#### m85xxCCSRRegister()

NAME	m85xxCCSRRegister() – register m85xxLAWBAR driver
SYNOPSIS	void m85xxCCSRRegister (void)
DESCRIPTION	This routine registers the m85xxLAWBAR driver and device recognition data with the vxBus subsystem.
NOTE	This routine is called early during system initialization, and *MUST NOT* make calls to OS facilities such as memory allocation and I/O.
RETURNS	N/A

ERRNO

SEE ALSO

m85xxCCSR

#### mRegs()

NAME	mRegs() – modify registers	
SYNOPSIS	<pre>STATUS mRegs   (     char * regName, /* register name, NULL for all */     int taskNameOrId /* task name or task ID, 0 = default task */   )</pre>	
DESCRIPTION	This command modifies the specified register for the specified task. If <i>taskNameOrId</i> is omitted or zero, the last task referenced is assumed. If the specified register is not found, it prints out the valid register list and returns <b>ERROR</b> . If no register is specified, it sequentially prompts the user for new values for a task's registers. It displays each register and the current contents of that register, in turn. The user can respond in one of several ways:	
	[RETURN] Do not change this register, but continue, prompting at the next register.	
	<i>number</i> Set this register to <i>number</i> .	
	. (dot) Do not change this register, and quit.	
	[EOF] Do not change this register, and quit.	
	All numbers are entered and displayed in hexadecimal, except floating-point values, which may be entered in double precision.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if the task or register does not exist.	
ERRNO	N/A	
SEE ALSO	usrLib, m(), the VxWorks programmer guides.	

#### mach()

NAME	mach() – return the contents of system register mach (also macl, pr) (SH)
SYNOPSIS	<pre>int mach   (    int taskId /* task ID, 0 means default task */   )</pre>
DESCRIPTION	This command extracts the contents of register mach from the TCB of a specified task. If <i>taskId</i> is omitted or zero, the last task referenced is assumed.
	Similar routines are provided for other system registers (macl, pr): macl(), pr(). Note that pc() is provided by usrLib.c.
RETURNS	The contents of register mach (or the requested system register).
ERRNO	Not Available
SEE ALSO	dbgArchLib, VxWorks Programmer's Guide: Debugging

### malloc()

NAME	<b>malloc( )</b> – allocate a block of memory from the system memory partition (ANSI)
SYNOPSIS	void * malloc ( size_t nBytes /* number of bytes to allocate */ )
DESCRIPTION	This routine allocates a block of memory from the free lists of the system memory partition (kernel heap). The size of the block will be equal to or greater than <i>nBytes</i> .
RETURNS	A pointer to the allocated block of memory, or a null pointer if there is an error.
ERRNO	Possible errnos generated by this routine include: <b>S_memLib_NOT_ENOUGH_MEMORY</b> There is no free block large enough to satisfy the allocation request.
SEE ALSO	<b>memPartLib</b> , <b>free( )</b> , <b>calloc( )</b> , <b>valloc( )</b> , <b>memPartAlloc( )</b> , American National Standard for Information Systems -, Programming Language - C, ANSI X3.159-1989: General Utilities ( <b>stdlib.h</b> )

#### memAddToPool()

```
memAddToPool() – add memory to the system memory partition
NAME
SYNOPSIS
                STATUS memAddToPool
                    (
                    FAST char * pPool, /* pointer to memory block */
                    FAST unsigned poolSize /* block size in bytes */
DESCRIPTION
                This routine adds memory to the system memory partition (kernel heap), in addition to the
                amount of memory specified during its creation.
RETURNS
                OK or ERROR.
ERRNO
                Possible errnos generated by this routine include:
                S_memLib_INVALID_ADDRESS
                    pPool is equal to NULL.
                S_memLib_INVALID_NBYTES
                    poolSize value is too small.
                memPartLib, memPartAddToPool()
SEE ALSO
```

#### memDevCreate()

**DESCRIPTION** This routine creates a memory device containing a single file. Memory for the device is simply an absolute memory location beginning at *base*. The *length* parameter indicates the size of memory.

For example, to create the device "/mem/cpu0/", a device for accessing the entire memory of the local processor, the proper call would be:

memDevCreate ("/mem/cpu0/", 0, sysMemTop())

VxWorks Kernel API Reference, 6.6 memDevCreate()

The device is created with the specified name, start location, and size.

To open a file descriptor to the memory, use **open()**. Specify a pseudo-file name of the byte offset desired, or open the "raw" file at the beginning and specify a position to seek to. For example, the following call to **open()** allows memory to be read starting at decimal offset 1000.

-> fd = open ("/mem/cpu0/1000", O\_RDONLY, 0)

Pseudo-file name offsets are scanned with "%d".

**CAVEAT** The FIOSEEK operation overrides the offset given via the pseudo-file name at open time.

**EXAMPLE** Consider a system configured with two CPUs in the backplane and a separate dual-ported memory board, each with 1 megabyte of memory. The first CPU is mapped at VMEbus address 0x00400000 (4 Meg.), the second at bus address 0x00800000 (8 Meg.), the dual-ported memory board at 0x00c00000 (12 Meg.). Three devices can be created on each CPU as follows. On processor 0:

```
-> memDevCreate ("/mem/local/", 0, sysMemTop())
...
-> memDevCreate ("/mem/cpu1/", 0x00800000, 0x00100000)
...
-> memDevCreate ("/mem/share/", 0x00c00000, 0x00100000)
```

#### On processor 1:

```
-> memDevCreate ("/mem/local/", 0, sysMemTop())
...
-> memDevCreate ("/mem/cpu0/", 0x00400000, 0x00100000)
...
-> memDevCreate ("/mem/share/", 0x00c00000, 0x00100000)
```

Processor 0 has a local disk. Data or an object module needs to be passed from processor 0 to processor 1. To accomplish this, processor 0 first calls:

-> copy </disk1/module.o >/mem/share/0

Processor 1 can then be given the load command:

-> ld </mem/share/0

**RETURNS** OK, or ERROR if memory is insufficient or the I/O system cannot add the device.

- ERRNO S\_ioLib\_NO\_DRIVER
- SEE ALSO memDrv

#### memDevCreateDir()

memDevCreateDir() - create a memory device for multiple files NAME SYNOPSIS STATUS memDevCreateDir ( 
 char \*
 name,
 /\* device name
 \*/

 MEM\_DRV\_DIRENTRY \* files,
 /\* array of dir. entries - not copied \*/

 int
 numFiles
 /\* number of entries
 \*/
 ) This routine creates a memory device for a collection of files organised into directories. The DESCRIPTION given array of directory entry records describes a number of files, some of which may be directories, represented by their own directory entry arrays. The structure may be arbitrarily deep. This effectively allows a filesystem to be created and installed in VxWorks, for essentially read-only use. The filesystem structure can be created on the host using the memdrvbuild utility. Note that the array supplied is not copied; a reference to it is kept. This array should not be modified after being passed to memDevCreateDir. RETURNS **OK**, or **ERROR** if memory is insufficient or the I/O system cannot add the device. ERRNO S\_ioLib\_NO\_DRIVER memDrv SEE ALSO

#### memDevDelete()

NAME	memDevDelete() – delete a memory device
Synopsis	STATUS memDevDelete ( char * name /* device name */ )
DESCRIPTION	This routine deletes a memory device containing a single file or a collection of files. The device is deleted with it own name.
	For example, to delete the device created by memDevCreate ("/mem/cpu0/", 0, <b>sysMemTop( )</b> ), the proper call would be:
	<pre>memDevDelete ("/mem/cpu0/");</pre>

VxWorks Kernel API Reference, 6.6 memDrv()

**RETURNS** OK, or ERROR if the device doesn't exist.

ERRNO N/A.

SEE ALSO memDrv

#### memDrv()

NAME	<b>memDrv( )</b> – install a memory driver
SYNOPSIS	STATUS memDrv (void)
DESCRIPTION	This routine initializes the memory driver. It is called automatically when VxWorks is configured with the INCLUDE_MEMDRV component.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the I/O system cannot install the driver.
ERRNO	N/A.
SEE ALSO	memDrv

#### memEdrBlockMark()

NAME	memEdrBlockMark() – mark or unmark selected blocks
SYNOPSIS	<pre>int memEdrBlockMark  (     int partId, /* partition ID selector */     int taskId, /* task ID selector */     BOOL unmark /* TRUE to unmark */    )</pre>
DESCRIPTION	This routine marks blocks selected by partition ID and/or taskId. Passing <b>NULL</b> for either <i>partId</i> or <i>taskId</i> means no filtering is done for that field.
RETURNS	number of newly marked or unmarked blocks
ERRNO	Not Available
SEE ALSO	memEdrLib, memEdrBlockShow( )

# memEdrBlockShow()

NAME	<b>memEdrBlockShow()</b> – print memory block information
SYNOPSIS	STATUS memEdrBlockShow ( int partId, /* partition ID selector */ void * addr, /* address selector */ int taskId, /* task ID selector */ UINT type, /* block type selector */ UINT level, /* detail level */ BOOL continuous /* print in continuous mode */ )
DESCRIPTION	This routine displays memory block information based on various selection criteria. NULL or 0 can be used for partId, addr, taskId and type to exclude the respective field from filtering. The level parameter can be used to enable printing of extended block information (trace) when collection of extended information is enabled.         The following type parameters are accepted:         type       description         0       any block         1       global variable reported by RTC         2       allocated block         3       queued free block         4       marked allocated block         5       unmarked allocated block         1       globack at time. With continuous mode information is also collected 20 at a time, but printing is continuous, with no user intervention enabled. Note that either way, after each batch of 20 blocks the mutex lock is released, allowing other tasks to change the instrumentation database.
RETURNS	OK, or ERROR if getting the info failed.
ERRNO	N/A
SEE ALSO	memEdrShow, memEdrLib, memEdrPartShow( ), memEdrBlockMark( )

#### memEdrFreeQueueFlush()

NAME	memEdrFreeQueueFlush() – flush the free queue
SYNOPSIS	void memEdrFreeQueueFlush (void)
DESCRIPTION	This routine can be used to remove all blocks queued on the free queue, and finalize the free operation. This way memory blocks previously queued will be freed into their respective memory partitions.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	memEdrLib

#### memEdrPartShow()

NAME	<b>memEdrPartShow()</b> – show partition information in the kernel
SYNOPSIS	STATUS memEdrPartShow ( PART_ID partId /* partition ID */ )
DESCRIPTION	This routine displays information about memory partitions in the kernel. If the <i>partId</i> parameter is <b>NULL</b> , it lists all partitions recorded in the kernel's database.
RETURNS	OK, or ERROR if getting the info failed.
ERRNO	N/A
SEE ALSO	memEdrShow, memEdrLib, memEdrBlockShow( )

#### memEdrRtpBlockMark() - mark or unmark selected allocated blocks in an RTP NAME SYNOPSIS int memEdrRtpBlockMark RTP\_ID rtpId, /\* RTP id \*/ int partId, /\* partition ID selector \*/ int taskId, /\* task ID selector \*/ BOOL unmark /\* TRUE to unmark \*/ ) DESCRIPTION This routine marks blocks selected by partition ID and/or task ID. Passing NULL for either *partId* or *taskId* means no filtering is done using that field. This routine only works with RTPs with the memory manager instrumentation (memEdrLib) enabled and the MEDR\_SHOW\_ENABLE environment variable set to TRUE. number of newly marked or unmarked blocks RETURNS N/A ERRNO memEdrRtpShow, memEdrLib, memEdrRtpBlockShow() SEE ALSO

#### memEdrRtpBlockShow()

memEdrRtpBlockMark()

 NAME
 memEdrRtpBlockShow() – print memory block information of an RTP

 SYNOPSIS
 STATUS memEdrRtpBlockShow

 (
 RTP\_ID rtpId, /\* RTP id \*/
 int partId, /\* partition ID selector \*/
 void \* addr, /\* address selector \*/
 int taskId, /\* task ID selector \*/
 UINT type, /\* block type selector \*/
 UINT level, /\* detail level \*/
 BOOL continuous /\* print in continuous mode \*/
 )

 DESCRIPTION
 This routine displays memory block information based on various selector

**DESCRIPTION**This routine displays memory block information based on various selection criteria. NULL<br/>or 0 can be used for *partId*, *addr*, *taskId* and *type* to exclude the respective field from filtering.<br/>The *level* parameter can be used to enable printing of extended block information (trace)<br/>when collection of extended information is enabled.

The following *type* parameters are accepted:

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	type	description
	0	any block
	1	global variable reported by RTC
	2	allocated block (marked or unmarked)
	3	queued free block
	4	marked allocated block
	5	unmarked allocated block
	If more that printed 20 printing is 20 blocks to database.	an 20 blocks match the selection criteria and <i>continuous</i> is <b>FALSE</b> , blocks are at a time. With continuous mode information is also collected 20 at a time, but continuous, with no user intervention enabled. Note that between each batch of the mutex lock is released allowing other tasks to change the instrumentation
	This routin ( <b>memEdrl</b> For symbo option.	ne only works with RTPs with the memory manager instrumentation Lib) enabled and the MEDR_SHOW_ENABLE environment variable set to TRUE. olic information, the RTP has to be spawned with the RTP_WITH_SYMBOLS
RETURNS	OK, or ERROR if getting the info failed.	
ERRNO	N/A	
SEE ALSO	memEdrR	tpShow, memEdrLib, memEdrRtpPartShow( ), memEdrRtpBlockMark( )

# memEdrRtpPartShow()

NAME	<b>memEdrRtpPartShow()</b> – show partition information of an RTP
SYNOPSIS	STATUS memEdrRtpPartShow ( RTP_ID rtpId, /* RTP id */ int partId /* partition ID selector */ )
DESCRIPTION	This routine displays information about memory partitions in an RTP. If the <i>partId</i> parameter is <b>NULL</b> , it lists all partitions recorded in the RTP's database.
	This routine only works with RTPs with the memory manager instrumentation (memEdrLib) enabled and the MEDR_SHOW_ENABLE environment variable set to TRUE.
RETURNS	<b>OK</b> , or <b>ERROR</b> if getting the info failed.
ERRNO	N/A

#### memFindMax()

**NAME** memFindMax() – find the largest free block in the system memory partition (kernel heap)

SYNOPSIS int memFindMax (void)

- **DESCRIPTION**This routine searches for the largest block in the system memory partition (kernel heap )<br/>free list and returns its size. It returns 0 if there is no free block in the system memory<br/>partition. The size returned corresponds to the largest block that can be allocated using the<br/>default alignment value, which is used via calls to malloc(), realloc(), or calloc().Allocation of such a size with an alignment greater than the default alignment value is documented<br/>in the manual entry for memPartLib as the architecture specific boundary.
- **RETURNS** The size, in bytes, of the largest available block.
- ERRNO Not Available
- SEE ALSO memInfo, memPartFindMax()

#### memInfoGet( )

NAME	<b>memInfoGet( )</b> – get heap information
SYNOPSIS	STATUS memInfoGet ( MEM_PART_STATS * pPartStats /* partition stats structure */ )
DESCRIPTION	This routine takes a pointer to a <b>MEM_PART_STATS</b> structure. All fields of the structure are filled in with data from the RTP heap memory partition. For the description of the information provided, see the <b>memPartInfoGet()</b> documentation.
RETURNS	OK if the structure has valid data, otherwise ERROR.
ERRNO	Not Available

VxWorks Kernel API Reference, 6.6 memOptionsGet()

SEE ALSO memInfo, memPartInfoGet( )

#### memOptionsGet()

NAME	<b>memOptionsGet( )</b> – get the options of the system memory partition (kernel heap)
SYNOPSIS	STATUS memOptionsGet ( UINT * pOptions /* pointer to options for kernel heap */ )
DESCRIPTION	This routine sets the parameter <i>pOptions</i> with the options of the system memory partition (kernel heap).
	Heap/memory partition options are discussed in details in the reference entry for the library <b>memLib</b> .
RETURNS	OK or ERROR.
ERRNO	Not Available
SEE ALSO	memLib, memOptionsSet( ), memPartOptionsGet( ), memPartOptionsSet( )

#### memOptionsSet()

NAME	<b>memOptionsSet( )</b> – set the options for the system memory partition (kernel heap)		
SYNOPSIS	STATUS memOptionsSet ( unsigned options /* options for system memory partition (kernel heap) */ )		
DESCRIPTION	This routine sets the debug and error handling options for the system memory partition (kernel heap). For detailed description of these options see the <b>memLib</b> and <b>memPartOptionsSet( )</b> .		
RETURNS	OK or ERROR.		
ERRNO	Not Available		

#### memPartAddToPool()

NAME	<pre>memPartAddToPool( ) - add memory to a memory partition</pre>		
SYNOPSIS	STATUS memPartAddToPool ( FAST PART_ID partId, /* partition to add memory to */ FAST char * pPool, /* pointer to memory block */ FAST unsigned poolSize /* block size in bytes */ )		
DESCRIPTION	This routine adds memory to a specified memory partition already created with <b>memPartCreate()</b> . The memory added need not be contiguous with memory previously assigned to the partition.		
	The size of the memory pool being added has to be large enough to accommodate the section overhead consisting of a section header and some reserved blocks that mark the beginning and the end of the section. This overhead, approximately 64 bytes, is not available for allocation.		
	This routine does not verify that the memory block passed corresponds to valid memory or not. It is the user's responsability to ensure that the block is valid and it does not overlap with other blocks added to the partition.		
RETURNS	OK or ERROR.		
ERRNO	Possible errnos generated by this routine include:		
	S_smObjLib_NOT_INITIALIZED partId is a shared partition but the Shared Memory Allocator component was not initialized.		
	<b>S_memLib_INVALID_ADDRESS</b> <i>pPool</i> is equal to <b>NULL</b> .		
	S_memLib_INVALID_NBYTES poolSize value is too small.		
SEE ALSO	memPartLib, smMemLib, memPartCreate( ), memAddToPool( )		

#### memPartAlignedAlloc()

memPartAlignedAlloc() – allocate aligned memory from a partition NAME void \* memPartAlignedAlloc SYNOPSIS ( FAST PART\_ID partId, /\* memory partition to allocate from \*/ unsigned nBytes, /\* number of bytes to allocate \*/ unsigned alignment /\* boundary to align to \*/ ) This routine allocates a buffer of size *nBytes* from a specified partition. Additionally, it DESCRIPTION ensures that the allocated buffer begins on a memory address evenly divisible by *alignment*. The *alignment* parameter must be a power of 2. A pointer to the newly allocated block, or NULL if the buffer could not be allocated. RETURNS ERRNO Possible errnos generated by this routine include: S memLib INVALID ALIGNMENT *alignment* is not a power of two. S\_memLib\_NOT\_ENOUGH\_MEMORY There is no free block large enough to satisfy the allocation request.

SEE ALSO memPartLib, memalign()

#### memPartAlloc( )

NAME	<b>memPartAlloc()</b> – allocate a block of memory from a partition		
SYNOPSIS	<pre>void * memPartAlloc   (   FAST PART_ID partId, /* memory partition to allocate from */   unsigned nBytes /* number of bytes to allocate */ )</pre>		
DESCRIPTION	This routine allocates a block of memory from a specified partition. The size of the block will be equal to or greater than <i>nBytes</i> . The partition must already be created with <b>memPartCreate()</b> .		
RETURNS	A pointer to a block, or NULL if the call fails.		
ERRNO	Possible errnos generated by this routine include:		

# S\_smObjLib\_NOT\_INITIALIZED partId is a shared partition but the Shared Memory Allocator component was not initialized. S\_memLib\_NOT\_ENOUGH\_MEMORY There is no free block large enough to satisfy the allocation request. SEE ALSO memPartLib, smMemLib, memPartCreate(), malloc()

#### memPartCreate( )

NAME	<b>memPartCreate()</b> – create a memory partition		
SYNOPSIS	<pre>PART_ID memPartCreate   (     char * pPool, /* pointer to memory area */     unsigned poolSize /* size in bytes */   )</pre>		
DESCRIPTION	This routine creates a new memory partition containing a specified memory pool defined by its start address, <i>pPool</i> , and its size in bytes, <i>poolSize</i> . It returns a partition ID, which can be passed to other routines to manage the partition (i.e., to allocate and free memory blocks in the partition). Partitions can be created to manage any number of separate memory pools.		
Empty memory partitions can be created by setting <i>pPool</i> to <b>NULL</b> and <i>poolSize</i> to <b>partitions</b> , it is necessary to add memory blocks to the partition via <b>memPartA</b> before performing any allocation request.			
	Unless creating an empty partition, the memory pool size has to be large enough to accomodate some overhead consisting of a section header and some reserved blocks that mark the beginning and the end of the section. In addition, certain internal data structures used to store free block information are also carved from the pool. This overhead, in total approximately 248 bytes, is not available for allocations.		
	The create routine does not verify that the memory block passed corresponds to valid memory or not. It is the user's responsability to make sure the block is valid.		
NOTE	The descriptor for the new partition object is allocated out of the system memory partition (i.e., with <b>malloc( )</b> ).		
RETURNS	The partition ID, or <b>NULL</b> if there is insufficient memory in the system memory partition (kernel heap) for a new partition descriptor, or <i>poolSize</i> value is too small.		
ERRNO	Possible errnos generated by this routine include:		

VxWorks Kernel API Reference, 6.6 memPartDelete()

**S\_memLib\_INVALID\_NBYTES** *poolSize* value is too small.

SEE ALSO memPartLib, smMemLib

#### memPartDelete()

NAME	<b>memPartDelete( )</b> – delete a partition and free associated memory		
Synopsis	STATUS memPartDelete ( PART_ID partId /* partition to delete */ )		
DESCRIPTION	This routine deletes the memory partition object. It is supported for local memory partition but not for shared memory partition.		
RETURNS	OK or ERROR.		
ERRNO	Possible errnos generated by this routine include:		
	S_memLib_NO_PARTITION_DESTROY feature not supported for shared memory partition.		
SEE ALSO	memPartLib		

#### memPartFindMax()

NAME	<b>memPartFindMax()</b> – find the size of the largest available free block		
SYNOPSIS	int memPartFindMax ( FAST PART_ID partId /* partition ID */ )		
DESCRIPTION	This routine searches for the largest block in the memory partition free list and returns its size. It returns 0 if there is no free block in the memory partition. The size returned corresponds to the largest block that can be allocated using the default alignment value, which is used via calls to <b>memPartAlloc()</b> , or <b>memPartRealloc()</b> . Allocation of such a size with an alignment greater than the default alignment will fail: this may occur when using		

	<b>memPartAlignedAlloc( )</b> . The default alignment value is documented in the manual entry for <b>memPartLib</b> as the architecture specific <b>boundary</b> .	
RETURNS	The size, in bytes, of the largest available block.	
ERRNO	Possible errnos generated by this routine include:	
	<b>S_smObjLib_NOT_INITIALIZED</b> <i>partId</i> is a shared partition but the Shared Memory Allocator component was not initialized.	
SEE ALSO	memInfo, smMemLib, memFindMax( )	

#### memPartFree()

NAME	<b>memPartFree( )</b> – free a block of memory in a partition		
SYNOPSIS	<pre>STATUS memPartFree  (  PART_ID partId, /* memory partition to free a block from */  char * pBlock /* pointer to block of memory to free */ )</pre>		
DESCRIPTION	This routine returns to a partition's free memory lists a block of memory previously allocated with <b>memPartAlloc( )</b> , <b>memPartAlignedAlloc( )</b> or <b>memPartRealloc( )</b> . If <i>pBlock</i> is a null pointer, no action occurs and the function returns <b>OK</b> .		
RETURNS	<b>OK</b> , or <b>ERROR</b> if the block or the partition is invalid.		
ERRNO	Possible errnos generated by this routine include:		
	S_smObjLib_NOT_INITIALIZED partId is a shared partition but the Shared Memory Allocator component was not initialized.		
	<b>S_memLib_BLOCK_ERROR</b> The block of memory to free is not valid.		
	<b>S_memLib_WRONG_PART_ID</b> The block does not belong to the partition.		
SEE ALSO	memPartLib, smMemLib, memPartAlloc( ), memPartAlignedAlloc( ), free( )		

#### memPartInfoGet()

NAME	<b>memPartInfoGet( )</b> – get partition information			
SYNOPSIS	<pre>STATUS memPartInfoGet   (    PART_ID partId, /* partition ID */    MEM_PART_STATS * pPartStats /* partition stats structure */   )</pre>			
DESCRIPTION	This routine takes a partition ID and a pointer to a <b>MEM_PART_STATS</b> structure. All parameters of the structure are filled in with the current partition information whice include:			
	numBytesFree number of free bytes in the partition			
	numBlocksFree number of free blocks in the partition maxBlockSizeFree maximum block size in bytes that is free numBytesAlloc number of allocated bytes in the partition			
	numBlocksAlloc number of allocated blocks in the partition			
	maxBytesAlloc maximum number of allocated bytes at any time (peak usage)			
RETURNS	<b>OK</b> if the structure has valid data, otherwise <b>ERROR</b> .			
ERRNO	Not Available			
SEE ALSO	memInfo, memShow(), memPartShow()			

#### memPartOptionsGet()

NAME	<b>memPartOptionsGet()</b> – get the options of a memory partition		
SYNOPSIS	IOPSIS STATUS memPartOptionsGet		
	( PART_ID partId,	/* partition to set option for */	
```
      UINT * pOptions /* pointer to partition options */

      DESCRIPTION

      This routine sets the parameter pOptions with the options of a specified memory partition.

      RETURNS
      OK, or ERROR if partition is shared or pOptions is a NULL pointer.

      ERRNO
      Possible errnos generated by this routine include:

      S_memLib_FUNC_NOT_AVAILABLE
      partId is a shared partition for which memPartOptionsGet() is not supported.

      SEE ALSO
      memLib, smMemLib, memPartOptionsSet(), memOptionsGet()
```

## memPartOptionsSet()

NAME	<b>memPartOptionsSet()</b> – set the options for a memory partition
SYNOPSIS	<pre>STATUS memPartOptionsSet   (    PART_ID partId, /* partition to set option for */    unsigned options /* memory management options */   )</pre>
DESCRIPTION	This routine sets the debug options for a specified memory partition. Two kinds of errors are detected: attempts to allocate more memory than is available, and bad blocks found when memory is freed. In both cases, the error status is returned. For the supported options see the <b>memLib</b> library reference guide.
RETURNS	OK or ERROR.
ERRNO	Possible errnos generated by this routine include:
	<b>S_smObjLib_NOT_INITIALIZED</b> <i>partId</i> is a shared partition but the Shared Memory Allocator component was not initialized.
SEE ALSO	<pre>memLib, smMemLib, memPartOptionsGet( ), memOptionsSet( )</pre>

#### memPartRealloc( )

NAME	<b>memPartRealloc( )</b> – reallocate a block of memory in a specified partition
SYNOPSIS	<pre>void * memPartRealloc   (    PART_ID partId, /* partition ID */    char * pBlock, /* block to be reallocated */    unsigned nBytes /* new block size in bytes */   )</pre>
DESCRIPTION	This routine changes the size of a specified block of memory and returns a pointer to the new block. The contents that fit inside the new size (or old size if smaller) remain unchanged. The memory alignment of the new block is not guaranteed to be the same as the original block.
	If <i>pBlock</i> is <b>NULL</b> , this call is equivalent to <b>memPartAlloc( )</b> .
	If <i>nBytes</i> is set to zero and <i>pBlock</i> points to a valid allocated block, this call is equivalent to <b>memPartFree()</b> and returns <b>NULL</b> .
RETURNS	A pointer to the new block of memory, <b>NULL</b> if the call fails or <i>nBytes</i> is equal to zero.
ERRNO	Possible errnos generated by this routine include:
	S_memLib_BLOCK_ERROR The block of memory to free is not valid.
	<b>S_smObjLib_NOT_INITIALIZED</b> <i>partId</i> is a shared partition but the Shared Memory Allocator component was not initialized.
	<b>S_memLib_NOT_ENOUGH_MEMORY</b> There is no free block large enough to satisfy the allocation request.
	S_memLib_WRONG_PART_ID The block does not belong to the partition.
SEE ALSO	memLib, smMemLib, realloc()

## memPartShow()

NAME memPartShow() – show blocks and statistics for a given memory partition

SYNOPSIS STATUS memPartShow

SEE ALSO	memShow, memShow(), memPartAddToPool()
	S_smObjLib_NOT_INITIALIZED partId is a shared partition but the Shared Memory Allocator component was not initialized.
ERRNO	Possible errnos generated by this routine include:
RETURNS	OK or ERROR.
DESCRIPTION	) This routine displays statistics about the available and allocated memory in a specified memory partition. For details about usage and information shown by this routine refer to the <b>memShow()</b> documentation.
	( PART_ID partId, /* memory partition ID */ int type /* 0 = statistics, 1 = statistics & list */ /* 2 = statistics & list & extra info */

## memPartSmCreate()

NAME	<b>memPartSmCreate( )</b> – create a shared memory partition (VxMP Option)		
SYNOPSIS	<pre>PART_ID memPartSmCreate   (     char * pPool, /* global address of shared memory area */     unsigned poolSize /* size in bytes */   )</pre>		
DESCRIPTION	This routine creates a shared memory partition that can be used by tasks on all CPUs in the system. It returns a partition ID which can then be passed to generic <b>memPartLib</b> routines to manage the partition (i.e., to allocate and free memory blocks in the partition).		
	pPool is the global address of shared memory dedicated to the partition. The memory area pointed to by $pPool$ must be in the same address space as the shared memory anchor and shared memory pool.		
	poolSize is the size in bytes of shared memory dedicated to the partition.		
	Before this routine can be called, the shared memory objects facility must be initialized (see <b>smMemLib</b> ).		
NOTE	The descriptor for the new partition is allocated out of an internal dedicated shared memory partition. The maximum number of partitions that can be created is <b>SM_OBJ_MAX_MEM_PART</b> .		

VxWorks Kernel API Reference, 6.6 memShow()

Memory pool size is rounded down to a 16-byte boundary.

- **AVAILABILITY** This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.
- **RETURNS** The partition ID, or **NULL** if there is insufficient memory in the dedicated partition for a new partition descriptor.
- ERRNO S\_memLib\_NOT\_ENOUGH\_MEMORY S\_smObjLib\_LOCK\_TIMEOUT

SEE ALSO smMemLib, memLib

#### memShow()

NAME	<b>memShow()</b> – show blocks and statistics for the current heap partition		
SYNOPSIS	<pre>STATUS memShow (     int type /* 0 = summary, 1 = list all blocks in the free list, */</pre>		
DESCRIPTION	This routine displays statistics about the available and allocated memory in the current heap partition. It shows the number of bytes, the number of blocks, and the average block size in both free and allocated memory, as well as the maximum block size of free memory. It also shows the number of blocks currently and cumulatively allocated, the average allocated block size and the maximum number of bytes allocated at any given time (peak usage). The cumulative information wraps around after reaching <b>UINT_MAX</b> (4GB). Part of the heap space is used internally by the partition for its bookkeeping. This amount of memory is displayed as part of the current usage.		
	The <b>memShow()</b> routine called with <i>type</i> 1 or 2 requires that certain internal data structures (binary tree, linked list) are traversed. In case of corrupted internal heap structures an exception may occur, causing the task executing <b>memShow()</b> to get suspended while holding the partition's mutex semaphore. This, in effect, results in the partition being locked indefinitely.		
	The <b>memShow()</b> routine temporarily saves free block information on the stack in order to avoid performing IO while the partition semaphore is taken. To avoid excessive stack requirement, the number of lines listed in the free block section is restricted; however, at least 200 lines are printed. Sections listed are also limited to 50. The user should make sure that <b>memShow()</b> is called in the context of a task that has sufficient stack space (approximately 4k needed for <b>memShow()</b> ).		

-> memShow status	bytes	blocks	avg block	max block
current				
free	2330696	11	211881	1962688
alloc	1858696	4961	374	-
internal	400	2	200	-
cumulative				
alloc	2240928	5419	413	-
peak				
alloc	2202992	-	-	-

**memShow()** can be used to detect memory leaks within the current heap. This can be achieved by comparing values of the current number of allocated bytes before and after the function call(s) you want to verify. The current amount of free memory cannot be used to detect memory leaks since it is updated everytime that memory blocks are allocated internally by the system for the heap partition bookkeeping.

In addition, if *type* is 1, the routine displays a list of all different size of free blocks present in the heap partition. The size corresponds to the amount of usable data plus the overhead required for the block header. The heap partition options are also displayed.

EXAMPLE -> memShow 1 LIST OF FREE BLOCKS: number size \_\_\_\_\_  $\begin{array}{cccc} 1 & 24 \\ 1 & 72 \\ 1 & 104 \\ 1 & 144 \\ 1 & 176 \\ 1 & 232 \\ 1 & 264 \\ 1 & 1440 \\ 1 & 127080 \\ 1 & 238472 \\ 1 & 1962688 \end{array}$ 1 24 OPTIONS: ALLOC\_ERROR\_LOG BLOCK\_ERROR\_LOG BLOCK\_ERROR\_EDR SUMMARY: status bytes blocks avg block max block current 

 free
 2330696
 11
 211881
 1962688

 alloc
 1858696
 4961
 374

 internal
 400
 2
 200

 cumulative
 alloc
 2240928
 5419
 413

2

peak alloc 2202992 - - -

If *type* is 2, the routine also displays the address of each free block. The address of the free blocks is the start address of the free block header, not the start address of the usable data. In addition, a list of all the memory sections that were added to the heap partition with **memPartAddToPool()** or **memAddToPool()** is displayed.

FREE BLOC number  1 1 1	KS: size	 24 0x	addr	_	
number 1 _1 _1	size	 24 0x	addr 	_	
1 1 1		24 0x		-	
1			002075f	0	
1		72 0x	002b6dd	.0	
1		104 Ox	00279a5	8	
T		144 Ox	002b596	8	
1		176 Ox	0027a04	8	
1		232 Ox	0020611	0	
1		264 0x	0027a1a	.8	
1	1	440 0x	:00204c5	8	
1	127	080 0x	002b6±2	8	
1	238	4/2 Ux	002/a2i	0	
Ŧ	1702	000 01	.0042003	0	
MEMORY SE start add	CTIONS A r si	DDED TO ze	THE PA	RTITION:	
0x00206f6	0 41	65792			
0x002003e 0x0020520	8 8	20000 4096			
:					
ALLOC_ERR	OR_LOG				
BLOCK_ERR	OR_LOG OR_EDR				
s by	tes	block	s avg	block	max bloc
2	330696		11	211881	1962688
1	858696	49	61	374	-
1	400		2	200	-
lve	240000	- 4	1.0	410	
2	240928	54	ТА	413	-
2	202992		_	_	_
2			-	_	_
	MEMORY SE start add ox002006f6 0x002003e 0x0020520 ALLOC_ERR BLOCK_ERR BLOCK_ERR BLOCK_ERR 1 uve 2 1 1 1 2 2 2	1 1 1 127 1 238 1 1962 MEMORY SECTIONS A start addr si  0x00206f60 41 0x002003e8 0x00205208 ALLOC_ERROR_LOG BLOCK_ERROR_LOG BLOCK_ERROR_EDR 	1 1440 0x 1 127080 0x 1 238472 0x 1 1962688 0x MEMORY SECTIONS ADDED TO start addr size 	1       1440       0x00204c5         1       127080       0x002b6f2         1       238472       0x0027a2f         1       1962688       0x00420d3         MEMORY SECTIONS ADDED TO THE PA         start addr       size	1       1440       0X00204c58         1       127080       0X002b6f28         1       238472       0x0027a2f0         1       1962688       0x00420d30         MEMORY SECTIONS ADDED TO THE PARTITION:       start addr       size

RETURNS

ERRNO

Not Available

OK or ERROR.

SEE ALSO memShow, memPartShow(), memAddToPool()

### memShowInit()

NAME	<b>memShowInit()</b> – initialize the memory partition show facility
SYNOPSIS	void memShowInit (void)
DESCRIPTION	This routine links the memory partition show facility into the VxWorks system. These routines are included automatically when this show facility is configured into VxWorks using the INCLUDE_MEM_SHOW component.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	memShow

## memalign()

NAME	<b>memalign()</b> – allocate aligned memory from system memory partition (kernel heap)		
SYNOPSIS	<pre>void * memalign   (     unsigned alignment, /* boundary to align to (power of 2) */     unsigned size</pre>		
DESCRIPTION	This routine allocates a buffer of size <i>size</i> from the system memory partition (kernel heap). Additionally, it insures that the allocated buffer begins on a memory address evenly divisible by the specified alignment parameter. The alignment parameter must be a power of 2.		
RETURNS	A pointer to the newly allocated block, or NULL if the buffer could not be allocated.		
ERRNO	Possible errnos generated by this routine include:		
	S_memLib_INVALID_ALIGNMENT alignment is not a power of two.		

#### S\_memLib\_NOT\_ENOUGH\_MEMORY

There is no free block large enough to satisfy the allocation request.

SEE ALSO memLib, memPartAlignedAlloc()

### miiBusCreate()

NAME	<b>miiBusCreate( )</b> – create an miiBus attached to a parent bridge
SYNOPSIS	STATUS miiBusCreate ( VXB_DEVICE_ID pDev, VXB_DEVICE_ID *pBus )
DESCRIPTION	This function allocates a new VxBus instance and configures it to be an miiBus device. The new miiBus device inherits almost all of its properties from its parent bridge device, specified by pDev. An miiBus must be created by all ethernet device instances that have an MII-based transceiver attached, regardless of whether or not the ethernet controller itself has access to the PHY's management registers. For example, with MPC8260 boards, there are two FCC 10/100 ethernet ports which use MII PHY's, but the MDIO pins for the PHY's are typically connected to parallel I/O port D, which is logically distinct from either FCC controller. Nevertheless, FCC1 and FCC2 must both have a child MII bus, even though initially these buses will appear empty. A separate driver is required to provide a 3rd MII bus instance attached to parellel I/O port D in order to actually make the PHY's available to the system. These PHY's should be remapped to the MII buses attached to the FCC ports via remapping entries in <b>hwconf.c</b> .
	The <b>miiBusCreate()</b> function will allocate a private pDrvCtrl structure for the bus device, along with an empty <b>END_MEDIALIST</b> list. (This list will be filled in as PHYs are attached.) Assuming the new instance is successfully allocated and configured, it will then be announced to VxBus. (Note that this just tells VxBus of the new device instance. The call to <b>vxbBusAnnounce()</b> , which tells VxBus to create a new bus instance, has to be done later.)
	If the caller provided a non-NULL pBus pointer, it will be used to return a pointer to the newly created device to the caller.
RETURNS	OK if a bus is successfully created, or ERROR otherwise
ERRNO	N/A
SEE ALSO	miiBus

## miiBusDelete()

NAME	miiBusDelete() – delete an miiBus and all its child devices
SYNOPSIS	STATUS miiBusDelete ( VXB_DEVICE_ID pDev )
DESCRIPTION	This routine is used to shut down an miiBus. All child PHY instances attached to the bus are deleted, and then the bus instance itself is destroyed. This routine should only be called by the parent driver that also called <b>miiBusCreate()</b> .
RETURNS	OK if a bus is successfully destroyed, or ERROR otherwise
ERRNO	N/A
SEE ALSO	miiBus

## miiBusGet()

NAME	<b>miiBusGet( )</b> – get the miiBus that goes with a given VxBus instance		
SYNOPSIS	STATUS miiBusGet ( VXB_DEVICE_ID pDev, VXB_DEVICE_ID *pBus )		
DESCRIPTION	This routine is meant for use by ethernet drivers to locate their child miiBus instances. Normally, each ethernet driver will create an miiBus with <b>miiBusCreate()</b> , and it should save a pointer to the instance that <b>miiBusCreate()</b> returns (in which case it doesn't need to use this function). However, hEnd drivers currently must call <b>miiBusCreate()</b> before their private pDrvCtrl structures are allocated, and hence have nowhere to store the miiBus pointer that <b>miiBusCreate()</b> returns. These drivers can therefore use <b>miiBusGet()</b> to recover and save this pointer later, usually during their <b>DLInit()</b> routines.		
RETURNS	OK if a bus is found, or ERROR otherwise		
ERRNO	N/A		
SEE ALSO	miiBus		

VxWorks Kernel API Reference, 6.6 miiBusListAdd( )

## miiBusListAdd()

NAME	miiBusListAdd( ) – Add a PHY to the MII monitor list
SYNOPSIS	void miiBusListAdd ( VXB_DEVICE_ID pDev )
DESCRIPTION	This function adds a PHY instance to the monitor list so that it can be checked periodically by the monitor task. This function is usually called by a PHY in its VxBus instConnect routine.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	miiBus

## miiBusListDel()

NAME	<pre>miiBusListDel() - Remove a PHY to the MII monitor list</pre>
SYNOPSIS	void miiBusListDel ( VXB_DEVICE_ID pDev )
DESCRIPTION	This function removes a PHY instance from the monitor list.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	miiBus

## miiBusMediaAdd()

NAME	miiBusMediaAdd() – add an entry to an miiBus's media list	
SYNOPSIS	STATUS miiBusMediaAdd ( VXB_DEVICE_ID pDev, UINT32 media )	
DESCRIPTION	This routine is used by a PHY instance to announce support for a given media type to its parent miiBus.	
	If the media type already exists in the list, this routine leaves the list alone and returns <b>ERROR</b> .	
RETURNS	OK if media isn't a duplicate, otherwise ERROR	
ERRNO	N/A	
SEE ALSO	miiBus	

## miiBusMediaDefaultSet()

NAME	miiBusMediaDefaultSet() – set the default media for an miiBus	
SYNOPSIS	STATUS miiBusMediaDefaultSet ( VXB_DEVICE_ID pDev, UINT32 media )	
DESCRIPTION	This routine is used to specify which media type to specify in the endMediaListDefault member of a bus's media list. PHY driver detaches. If the media type doesn't exist in the list, this routine returns <b>ERROR</b> .	
RETURNS	OK if media entry is found in the list, otherwise ERROR	
ERRNO	N/A	
SEE ALSO	miiBus	

## miiBusMediaDel()

NAME	<b>miiBusMediaDel( )</b> – delete an entry to an miiBus's media list	
SYNOPSIS	STATUS miiBusMediaDel ( VXB_DEVICE_ID pDev, UINT32 media )	
DESCRIPTION	This routine is used by a PHY instance to remove an entry for a given media type from its parent miiBus. This is used when a PHY driver detaches.	
	If the media type doesn't exist in the list, this routine leaves the list alone and returns <b>ERROR</b> .	
RETURNS	<b>OK</b> if media entry is found in the list, otherwise <b>ERROR</b>	
ERRNO	N/A	
SEE ALSO	miiBus	

## miiBusMediaListGet()

NAME	miiBusMediaListGet() – obtain a pointer to the bus's media list
SYNOPSIS	STATUS miiBusMediaListGet ( VXB_DEVICE_ID pDev, END_MEDIALIST ** mediaList )
DESCRIPTION	This routine returns a pointer to the bus's <b>END_MEDIALIST</b> structure to the caller. This is used by ethernet drivers that support ifmedia to service the EIOCGIFMEDIALIST ioctl.
RETURNS	OK if media list is found, or ERROR otherwise
ERRNO	N/A
SEE ALSO	miiBus

	miiBusMediaUpdate()	
NAME	miiBusMediaUpdate( ) – invoke a PHY's parent's media update callback	
SYNOPSIS	STATUS miiBusMediaUpdate ( VXB_DEVICE_ID pDev )	
DESCRIPTION	This function is used to notify the parent ethernet driver associated with a PHY device that a link change event has occured. This routine works by calling the miiMediaUpdate method exported by the parent driver. If the parent device has no miiMediaUpdate method, this routine fails.	
RETURNS	OK media update notification succeeds, or ERROR otherwise	
ERRNO	N/A	
SEE ALSO	miiBus	

## miiBusModeGet()

NAME	miiBusModeGet() – get the current media mode and link status	
SYNOPSIS	STATUS miiBusModeGet ( VXB_DEVICE_ID pDev, UINT32 * mode, UINT32 * sts )	
DESCRIPTION	This function queries the PHY currently active on the MII bus, specified by pDev, to determine the current link mode and state. The mode and link state are specified using ifmedia definitions (specified in <b>if_media.h</b> ). These values can be passed directly to calle of the EIOCGIFMEDIA ioctl in drivers that implement ifmedia support.	
	If no active PHY is currently selected, this routine will fail. If a PHY has been selected (by a call to <b>miiBusModeSet()</b> ), its miiModeGet method will be called to query its current setting. Note that this will result in a read of several of the PHY's registers being accessed, including the status register.	
RETURNS	OK if reading the current mode/state succeeds, or ERROR otherwise	

VxWorks Kernel API Reference, 6.6 miiBusModeSet()

ERRNO N/A

SEE ALSO miiBus

## miiBusModeSet()

NAME	miiBusModeSet() – set the current media mode	
SYNOPSIS	STATUS miiBusModeSet ( VXB_DEVICE_ID pDev, UINT32 mode )	
DESCRIPTION	This function sets the desired link mode of the MII bus. A bus could potentially have more than one PHY attached, though only one PHY should support any given mode (i.e. you can have one PHY supporting 100baseTX and another supporting 100baseFX, but you can't have two that both support 100baseTX there wouldn't be any point).	
	This routine will search the bus media list for the desired media and make the PHY that supports it the active PHY. The PHY's miiModeSet method will then be invoked to program the PHY for the desired mode. If the desired mode doesn't exist in the list, or if the PHY doesn't export an miiModeSet method, this routine fails.	
	This routine is typically used by ethernet drivers with ifmedia support to service the EIOCSIFMEDIA ioctl.	
RETURNS	OK if setting the mode succeeds, or ERROR otherwise	
ERRNO	N/A	
SEE ALSO	miiBus	

## miiBusRead()

NAME	miiBusRead() – read a PHY register
SYNOPSIS	STATUS miiBusRead (
	VXB_DEVICE_ID pDev, int phyAddr,

	int p UINT16 * )	bhyReg, 'regVal
DESCRIPTION	This function reads a register from a PHY, specified by pDev, at a given address. This routine works by invoking the miiRead method exported by the parent bridge device to which the bus is attached (e.g. the mottsec driver). If the parent bridge does not export a miiRead method, this routine will fail.	
RETURNS	OK read succeeds, or ERROR otherwise	
ERRNO	N/A	
SEE ALSO	miiBus	

## miiBusRegister()

NAME	<b>miiBusRegister()</b> – register with the vxBus subsystem
SYNOPSIS	void miiBusRegister(void)
DESCRIPTION	This routine registers the miiBus driver with vxBus as a child of all applicable parent bus types, and registers the <b>VXB_BUSID_MII</b> bus type.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	miiBus

## miiBusWrite()

NAME miiBusWrite() – write value to a PHY register
SYNOPSIS STATUS miiBusWrite
(
VXB\_DEVICE\_ID pDev,
int phyAddr,
int phyReg,

VxWorks Kernel API Reference, 6.6 mkdir( )

	UINT16 regVal )
DESCRIPTION	This function writes a value to a register from a PHY, specified by pDev, at a given address. This routine works by invoking the miiWrite method exported by the parent bridge device to which the bus is attached (e.g. the mottsec driver). If the parent bridge does not export an miiWrite method, this routine will fail.
RETURNS	OK write succeeds, or ERROR otherwise
ERRNO	N/A
SEE ALSO	miiBus

## mkdir()

NAME	mkdir() – make a directory	
SYNOPSIS	STATUS mkdir ( const char * dirName /* directory name */ )	
DESCRIPTION	This command creates a new directory in a hierarchical file system. The <i>dirName</i> string specifies the name to be used for the new directory, and can be either a full or relative pathname.	
	This call is supported by the VxWorks NFS and dosFs file systems.	
RETURNS	OK, or ERROR if the directory cannot be created.	
ERRNO	Not Available	
SEE ALSO	usrFsLib, rmdir(), VxWorks Kernel Programmer's Guide: Kernel Shell	

## mlock()

NAME mlock() – lock specified pages into memory (POSIX)

SYNOPSIS int mlock

**DESCRIPTION** This routine guarantees that the specified pages are memory resident. In VxWorks, the *addr* and *len* arguments are ignored, since all pages are memory resident.

**RETURNS** 0 (OK) always.

ERRNO N/A

SEE ALSO mmanPxLib

## mlockall()

NAME	<b>mlockall()</b> – lock all pages used by a process into memory (POSIX)
SYNOPSIS	int mlockall ( int flags )
DESCRIPTION	This routine guarantees that all pages used by a process are memory resident. In VxWorks, the <i>flags</i> argument is ignored, since all pages are memory resident.
RETURNS	0 ( <b>OK</b> ) always.
ERRNO	N/A
SEE ALSO	mmanPxLib

#### mmapShow()

NAME mmapShow() – show information about memory mapped objects in the system

SYNOPSIS STATUS mmapShow (

( char \* name ) 2

VxWorks Kernel API Reference, 6.6 mmapShow()

**DESCRIPTION** This routine displays information about objects mapped in the memory space of processes running in the system. Two types of objects are supported: regular files of supported file systems, and shared memory objects. These objects are mapped in the address space of a process with **mmap()**.

This routine can be used in two ways. With the summary mode, when *name* is **NULL**, this routine lists objects that are mapped in at least one process at the time of the call. Mappings of objects that have been unlinked are shown using this mode.

With the detailed mode, when a file or shared memory object name is specified, this routine lists all memory mappings of the object. For shared memory objects this also includes shared mappings that have been unmapped from all processes by the time this routine is called. Object names that have been unlinked are not accepted.

This routine should be used for debugging purposes only.

**EXAMPLE** The following example shows the output of summary **mmapShow()** using the shell's C-interpreter:

-> mi	mapShow
-------	---------

OBJECT	UNLINKED	RTP ID	RTP	NAME
/pxTestFs/mmapFile1	no	0x61746228	<	
		0x606d6010	<	
in/tmMmanFdLib.vxe				
/pxTestFs/mmapFile2	yes	0x61746228	<	
in/tmMmanFdLib.vxe		0	,	
in /tm/manEdLib		0X00000010	~	
III/ CIUMINAIIFULID.VXe				
value = 0 = 0x0				

The following example shows the output of summary **mmapShow()** using the shell's C-interpreter:

```
-> mmapShow "/pxTestFs/mmapFile1"

ADDRESS LENGTH PROT FLAGS OFFSET RTP ID

0x6306a000 0x00002000 R-- SHARED 0x000000000000 0x61746228

0x630d4000 0x00002000 RW- PRIVATE 0x00000000000 0x606d6010

value = 0 = 0x0
```

For the command-interpreter shell, use the **mmap list** command.

RETURNS N/A

ERRNO Not Available

SEE ALSO mmanShow

## mmuPhysToVirt()

NAME	mmuPhysToVirt() – translate a physical address to a virtual address (ARM)
SYNOPSIS	VIRT_ADDR mmuPhysToVirt ( PHYS_ADDR physAddr /* physical address to be translated */ )
DESCRIPTION	This function converts a physical address to a virtual address using the information contained within the sysPhysMemDesc structure of the BSP. This routine may be used both by the BSP MMU initialization and by the vm(Base)Lib code.
	If the BSP has a default mapping where physical and virtual addresses are not identical, then it must provide routines to the cache and MMU architecture code to convert between physical and virtual addresses. If the mapping described within the sysPhysMemDesc structure is accurate, then the BSP may use this routine. If it is not accurate, then routines must be provided within the BSP that are accurate.
NOTE	This routine simply performs a linear search through the sysPhysMemDesc structure looking for the first entry with an address range that includes the given address. Typically, the performance of this should not be a problem, as this routine will generally be called to translate RAM addresses, and by convention, the RAM entries come first in the structure. If this becomes an issue, the routine could be changed so that a separate structure to sysPhysMemDesc is used, containing the information in a more quickly accessible form. In any case, if this is not satisfactory, the BSP can provide its own routines.
RETURNS	the virtual address
ERRNO	Not Available
SEE ALSO	mmuMapLib, mmuVirtToPhys

### mmuPro32LibInit( )

NAME mmuPro32LibInit() – initialize module

SYNOPSIS STATUS mmuPro32LibInit

VxWorks Kernel API Reference, 6.6 mmuPro32Page0UnMap()

int pageSize /\* system pageSize (must be 4096 for i86) \*/ )

- Build a dummy translation table that will hold the page table entries for the global DESCRIPTION translation table. The mmu remains disabled upon completion. Supervisor Mode Only Not callable from ISR
- RETURNS OK if no error, ERROR otherwise

(

S\_mmuLib\_INVALID\_PAGE\_SIZE ERRNO

mmuPro32Lib SEE ALSO

### mmuPro32Page0UnMap()

NAME	mmuPro32Page0UnMap() – unmap the page zero for NULL pointer detection
SYNOPSIS	STATUS mmuPro32Page0UnMap (void)
DESCRIPTION	This routine unmap the page zero for NULL pointer access detection. Not Callable from user level. Not Callable from ISR.
RETURNS	OK or ERROR if mmuPageUnMap fails
ERRNO	Not Available
SEE ALSO	mmuPro32Lib

#### mmuPro36LibInit( )

```
mmuPro36LibInit() - initialize module
NAME
SYNOPSIS
                STATUS mmuPro36LibInit
                    (
                    int pageSize /* system pageSize (must be 4KB or 2MB) */
                    )
```

 DESCRIPTION
 Build a dummy translation table that will hold the page table entries for the global translation table. The mmu remains disabled upon completion.

 RETURNS
 OK if no error, ERROR otherwise

 ERRNO
 S\_mmuLib\_INVALID\_PAGE\_SIZE

 SEE ALSO
 mmuPro36Lib

#### mmuPro36Page0UnMap()

NAME mmuPro36Page0UnMap() – unmap the page zero for NULL pointer detection

- SYNOPSIS STATUS mmuPro36Page0UnMap (void)
- **DESCRIPTION** This routine unmap the page zero for NULL pointer access detection. Not Callable from user level. Not Callable from ISR.
- **RETURNS** OK or ERROR if mmuPageUnMap fails
- ERRNO Not Available
- SEE ALSO mmuPro36Lib

#### mmuPro36PageMap()

NAME	mmuPro36PageMap() – map 36bit physic	cal memory page to virtual memory page
SYNOPSIS	STATUS mmuPro36PageMap ( MMU_TRANS_TBL * transTbl, VIRT_ADDR virtualAddress, LL_INT physPage )	/* translation table */ /* 32bit virtual address */ /* 36bit physical address */
DESCRIPTION	The 36bit physical page address is entered into the PTE corresponding to the given virtua page. The state of a newly mapped page is undefined.	
RETURNS	<b>OK</b> or <b>ERROR</b> if translation table creation	failed.

VxWorks Kernel API Reference, 6.6 mmuPro36Translate()

ERRNO Not Available

SEE ALSO mmuPro36Lib

#### mmuPro36Translate()

mmuPro36Translate() - translate a virtual address to a 36bit physical address NAME SYNOPSIS STATUS mmuPro36Translate ( VIRT\_ADDR virtAddress, /\* 32bit virtual address \*/ LL\_INT \* physAddress /\* place to return 36bit result \*/ ) DESCRIPTION Traverse the translation table and extract the 36bit physical address for the given virtual address from the PTE corresponding to the virtual address. OK or ERROR if no PTE for given virtual address. RETURNS Not Available ERRNO mmuPro36Lib SEE ALSO

#### mmuShLibInit( )

NAME	<b>mmuShLibInit( )</b> – Initialize the SH MMU library.
SYNOPSIS	STATUS mmuShLibInit ( int pageSize /* minimum vm page size */ )
DESCRIPTION	This routine performs the necessary initialization for the SH MMU library. Initialization consists mainly of initializing processing variables, setting up the processing variables for the AIM and <b>vmLib</b> and calling the AIM init function.
RETURNS	OK or ERROR if unsuccessful.

ERRNO

SEE ALSO mmuShLib

## mmuVirtToPhys()

NAME	<b>mmuVirtToPhys( )</b> – translate a virtual address to a physical address (ARM)
SYNOPSIS	PHYS_ADDR mmuVirtToPhys ( VIRT_ADDR virtAddr /* virtual address to be translated */ )
DESCRIPTION	This function converts a virtual address to a physical address using the information contained within the sysPhysMemDesc structure of the BSP. This routine may be used both by the BSP MMU initialization and by the vm(Base)Lib code.
	If the BSP has a default mapping where physical and virtual addresses are not identical, then it must provide routines to the cache and MMU architecture code to convert between physical and virtual addresses. If the mapping described within the sysPhysMemDesc structure is accurate, then the BSP may use this routine. If it is not accurate, then routines must be provided within the BSP that are accurate.
NOTE	This routine simply performs a linear search through the sysPhysMemDesc structure looking for the first entry with an address range that includes the given address. Typically, the performance of this should not be a problem, as this routine will generally be called to translate RAM addresses, and by convention, the RAM entries come first in the structure. If this becomes an issue, the routine could be changed so that a separate structure to sysPhysMemDesc is used, containing the information in a more quickly accessible form. In any case, if this is not satisfactory, the BSP can provide its own routines.
RETURNS	the physical address
ERRNO	Not Available
SEE ALSO	mmuMapLib, mmuPhysToVirt

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## moduleCheck()

NAME	<b>moduleCheck()</b> – verify checksums on all modules loaded in the system	
SYNOPSIS	STATUS moduleCheck ( int options /* validation options */ )	
DESCRIPTION	This routine verifies the checksums on the sections of all loaded modules. The checksums are compared to original checksums computed when the modules were initialy loaded. If any of the checksums are incorrect, a message is printed to the console, and the routine returns <b>ERROR</b> .	
	By default, only the text section checksums are validated.	
	Bits in the <i>options</i> parameter may be set to control specific checks:	
	MODCHECK_TEXT Validate the checksum for the TEXT sections (default).	
	MODCHECK_DATA Validate the checksum for the DATA sections.	
	MODCHECK_BSS Validate the checksum for the BSS sections.	
	MODCHECK_RODATA Validate the checksum for the RODATA sections.	
	MODCHECK_ALL Validate the checksum for the all sections.	
	MODCHECK_NOPRINT Do not print a message (moduleCheck() still returns ERROR on failure.)	
	See the definitions in <b>moduleLib.h</b>	
CAVEAT	This routine is a not able to check the integrity of a module at the time of its load. It can only detect corruption subsequent to the load.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if a checksum is invalid for any module.	
ERRNO	Not Available	
SEE ALSO	moduleLib	

#### moduleCreate()

NAME	<b>moduleCreate()</b> – create and initialize a module		
SYNOPSIS	MODULE_ID moduleCreate ( char * name, /* module name */ int format, /* object module format */ int flags /* <options> passed to loader (see loadModuleAt()) */ )</options>		
DESCRIPTION	This routine creates a code module descriptor.		
	The arguments specify the name of the object module file, the object module format and an argument specifying which symbols to add to the symbol table. See the <b>loadModuleAt( )</b> description of <i>options</i> for possibles <i>flags</i> values.		
	Space for the new code module descriptor is allocated dynamically.		
	This function is not intended to be used by code outside of the VxWorks kernel libraries. Documentation is provided for reference only.		
RETURNS	The <b>MODULE_ID</b> of the newly created module or <b>NULL</b> if there is an error.		
ERRNO	Not Available		
SEE ALSO	loadModuleAt( ), moduleLib		

### moduleCreateHookAdd()

NAME	<b>moduleCreateHookAdd( )</b> – add a routine to be called when a module is added
SYNOPSIS	STATUS moduleCreateHookAdd ( FUNCPTR moduleCreateHookRtn /* routine called when module is added */ )
DESCRIPTION	This routine adds the specified routine to a list of routines to be called each time <b>moduleCreate()</b> is called. The specified routine should be declared as follows:
	<pre>void moduleCreateHookFunc   (     MODULE_ID moduleId /* the module ID to act upon */ )</pre>

This routine is called after all fields of the module ID have been filled in.

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VxWorks Kernel API Reference, 6.6 moduleCreateHookDelete()

NOTE	Modules do not have information about their object sections or segments when they are
	created. This information is not available until after the entire load process has finished.
	Therefore functions used as module create hooks should not use the section or segment
	information associated with a module.

**RETURNS OK** or **ERROR** if there was a problem.

ERRNO Not Available

SEE ALSO moduleLib, moduleCreateHookDelete()

### moduleCreateHookDelete()

NAME	<pre>moduleCreateHookDelete() – delete a previously added module create hook routine</pre>
SYNOPSIS	STATUS moduleCreateHookDelete ( FUNCPTR moduleCreateHookRtn /* routine called when module is added */ )
DESCRIPTION	This routine removes a specified routine from the list of routines to be called at each <b>moduleCreate()</b> call.
RETURNS	OK, or ERROR if the routine is not in the table of module creation hook routines.
ERRNO	<ul> <li>Possible errnos set by this routine include:</li> <li>+ S_moduleLib_HOOK_NOT_FOUND</li> <li>For a complete description of the errnos, see the reference documentation for moduleLib.</li> </ul>
SEE ALSO	moduleLib, moduleCreateHookAdd()

#### moduleDelete()

NAME	<b>moduleDelete( )</b> – delete module ID information
SYNOPSIS	STATUS moduleDelete ( MODULE_ID moduleId /* module to delete */ )

DESCRIPTION	This routine deletes a module descriptor, freeing any space that was allocated for the use of the module ID.
	This routine does not free space allocated for the object module itself this is done by the unload routine ( <b>unld()</b> or <b>unldByModuleId()</b> ).
	This function is not intended to be used by code outside of the VxWorks kernel libraries. Documentation is provided for reference only.
RETURNS	<b>OK</b> or <b>ERROR</b> if there was a problem.
ERRNO	<ul> <li>Possible errnos set by this routine include:</li> <li><b>S_moduleLib_INVALID_MODULE_ID</b></li> <li>For a complete description of the errnos, see the reference documentation for <b>moduleLib</b>.</li> </ul>
SEE ALSO	moduleLib, unldByModuleId( )

## moduleFindByGroup()

NAME	moduleFindByGroup() – find a module by group number
SYNOPSIS	MODULE_ID moduleFindByGroup ( int groupNumber /* group number to find */ )
DESCRIPTION	This routine searches for a module with a group number matching groupNumber.
RETURNS	A <b>MODULE_ID</b> corresponding to the first module whose group number matches, or <b>NULL</b> if no match is found.
ERRNO	Not Available
SEE ALSO	moduleLib, moduleIdFigure()

VxWorks Kernel API Reference, 6.6 moduleFindByName()

#### moduleFindByName()

moduleFindByName() - find a module by name NAME SYNOPSIS MODULE\_ID moduleFindByName ( char \* moduleName /\* name of module to find \*/ ) DESCRIPTION This routine searches for a module with a name matching *moduleName*. The name is the one that was used when the module was loaded. RETURNS A MODULE\_ID corresponding to the module name, or NULL if no match is found. Not Available ERRNO SEE ALSO moduleLib, moduleFindByNameAndPath()

#### moduleFindByNameAndPath()

NAME	<b>moduleFindByNameAndPath()</b> – find a module by filename and path
SYNOPSIS	MODULE_ID moduleFindByNameAndPath ( char * moduleName, /* file name to find */ char * pathName /* path name to find */ )
DESCRIPTION	This routine searches for a module with a name matching <i>moduleName</i> and path matching <i>pathName</i> . The name and path correspond to the parameters that were passed to the load routine <i>when the module was loaded</i> .
EXAMPLES	If the module was loaded using the following name and path:
	fd = open ("path/to/the/module/to/load/moduleName", O_RDONLY); moduleLoad (fd, LOAD_GLOBAL_SYMBOLS);
	then the call to <b>moduleFindByNameAndPath()</b> would be done as:
	<pre>moduleFindByNameAndPath("moduleName", "path/to/the/module/to/load");</pre>
	The path field should be left empty if the module was loaded without any path specified:
	fd = open ("moduleName", O_RDONLY); moduleLoad (fd, LOAD_GLOBAL_SYMBOLS); moduleFindByNameAndPath("moduleName", "");

**RETURNS** A MODULE\_ID, or NULL if no match is found.

ERRNO Not Available

SEE ALSO moduleLib

## moduleFlagsGet( )

NAME	moduleFlagsGet() – get the flags associated with a module ID
SYNOPSIS	<pre>int moduleFlagsGet   (    MODULE_ID moduleId   )</pre>
DESCRIPTION	This routine returns the flags associated with a module ID. A module's flags correspond to the options passed to the loader when loading the module. See <b>loadModuleAt()</b> reference entry for more information concerning loader flags.
RETURNS	The flags associated with the module ID, or zero if the module ID is invalid.
ERRNO	Possible errnos set by this routine include: + S_moduleLib_INVALID_MODULE_ID For a complete description of the errnos, see the reference documentation for moduleLib.
SEE ALSO	loadModuleAt( ), moduleLib

## moduleIdListGet()

NAME	<pre>moduleIdListGet( ) - get a list of loaded modules</pre>
SYNOPSIS	<pre>int moduleIdListGet   (    MODULE_ID * idList, /* Array of module IDs to be filled in */    int maxModules /* Max modules <idlist> can accommodate */   )</idlist></pre>
DESCRIPTION	This routine provides the calling task with a list of all loaded object modules. An unsorted list of module IDs for no more than <i>maxModules</i> modules is put into <i>idList</i> .

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VxWorks Kernel API Reference, 6.6 moduleInfoGet( )

**RETURNS** The number of modules put into *idList*.

ERRNO Not Available

SEE ALSO moduleLib

## moduleInfoGet( )

NAME	<b>moduleInfoGet( )</b> – get information about an object module
SYNOPSIS	<pre>STATUS moduleInfoGet   (    MODULE_ID moduleId, /* module to return information about */    MODULE_INFO * pModuleInfo /* pointer to module info struct */   )</pre>
DESCRIPTION	This routine fills in a <b>MODULE_INFO</b> structure with information about the specified module. Note that the name field of the <b>MODULE_INFO</b> structure is a fixed length ( <b>NAME_MAX</b> , see <i>vxParams.h</i> for actual value); the name of the module will be truncated to fit in the field, if necessary.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the module ID is invalid.
ERRNO	Possible errnos set by this routine include:
	+ S_moduleLib_INVALID_MODULE_ID
	For a complete description of the errnos, see the reference documentation for <b>moduleLib</b> .
SEE ALSO	moduleLib

### moduleNameGet()

NAME	<b>moduleNameGet()</b> – get the name associated with a module ID
SYNOPSIS	char * moduleNameGet (
	MODULE_ID moduleId )

DESCRIPTION	This routine returns a pointer to the name associated with a module ID. Note: this is a pointer to the module library's copy of the name string, so the memory it points to should not be modified.
RETURNS	A pointer to the module name corresponding to the module ID, or <b>NULL</b> if the module ID is invalid.
ERRNO	<ul> <li>Possible errnos set by this routine include:</li> <li><b>S_moduleLib_INVALID_MODULE_ID</b></li> <li>For a complete description of the errnos, see the reference documentation for <b>moduleLib</b>.</li> </ul>
SEE ALSO	moduleLib

## moduleSegFirst()

NAME	<b>moduleSegFirst( )</b> – find the first segment in a module
SYNOPSIS	<pre>SEGMENT_ID moduleSegFirst   (    MODULE_ID moduleId /* module to get first segment of */ )</pre>
DESCRIPTION	This routine returns the ID of the first segment of a module descriptor.
RETURNS	A pointer to the segment ID, or NULL if the segment list is empty or the module ID is invalid.
ERRNO	<ul> <li>Possible errnos set by this routine include:</li> <li><b>S_moduleLib_INVALID_MODULE_ID</b></li> <li>For a complete description of the errnos, see the reference documentation for <b>moduleLib</b>.</li> </ul>
SEE ALSO	moduleLib, moduleSegGet()

## moduleSegGet()

**NAME** moduleSegGet() – get (delete and return) the first segment from a module

SYNOPSIS SEGMENT\_ID moduleSegGet

	( MODULE_ID moduleId /* module to get segment from */ )
DESCRIPTION	This routine returns the ID of the first segment of a module descriptor, and then removes the segment descriptor from the module's segment list.
	The memory associated with the segment descriptor is not freed.
	This function is not intended to be used by code outside of the VxWorks kernel libraries. Documentation is provided for reference only. Use the routines <b>moduleSegFirst()</b> and <b>moduleSegNext()</b> to retrieve information about a module's segments.
RETURNS	A pointer to the segment ID, or NULL if the segment list is empty.
ERRNO	Possible errnos set by this routine include:
	+ S_moduleLib_INVALID_MODULE_ID
	For a complete description of the errnos, see the reference documentation for <b>moduleLib</b> .
SEE ALSO	moduleLib, moduleSegFirst( ), moduleSegNext( )

## moduleSegNext()

NAME	<b>moduleSegNext( )</b> – find the next segment in a module
Synopsis	SEGMENT_ID moduleSegNext ( SEGMENT_ID segmentId /* segment whose successor is to be found */ )
DESCRIPTION	This routine returns the ID of the segment in the list immediately following <i>segmentId</i> .
RETURNS	A SEGMENT_ID, or NULL if there is no next segment.
ERRNO	Not Available
SEE ALSO	moduleLib, moduleSegFirst ()

## moduleShow()

NAME	<b>moduleShow()</b> – show information about loaded modules
SYNOPSIS	STATUS moduleShow
	( char * moduleNameOrId, /* name or ID of the module to show */ int options /* display options */ )
DESCRIPTION	This routine displays information about currently loaded modules and where they are placed in memory. Different information can be obtained depending on the value of the <i>moduleNameOrId</i> parameter :
	NULL A summary list of all loaded modules will be displayed. For each module are displaye its base name, ID, group number and the start addresses of the text, data and BSS segments.
	A module ID or the name of a loaded module More information about this specific module will be displayed (namely the sizes of th text, data and BSS segments and the total size of the module).
	If the <i>options</i> parameter is set to <b>MODDISPLAY_LONG</b> , module names longer than 15 characters are displayed on their own line (they would otherwise be truncated). In this case if <b>moduleShow()</b> is called with a module ID as an argument, the full module path is also printed.
EXAMPLES	Show the list of all modules loaded (C shell):
	-> moduleShow ()
	MODULE NAME MODULE ID GROUP # TEXT START DATA START BSS START
	versionDotOWith 0x616fc520         2 0x60532000 0x60534000 NO SEGMENT           versionDot015.0 0x616fbe98         3 0x60536000 0x60538000 NO SEGMENT           ctdt.0 0x60534020         4 NO SEGMENT 0x61700000 NO SEGMENT           value = 0 = 0x0         0x60534020
	Display information about a particular module (C shell):
	-> moduleShow (0x616fc520, 0)
	MODULE NAME MODULE ID GROUP # TEXT START DATA START BSS START
	versionDotOWith 0x616fc520 2 0x60532000 0x60534000 NO SEGMENT
	Size of text segment:58Size of data segment:16Size of bss segment:0Total size:74value = 0 = 0x0

Display full names for all modules loaded (C shell):

-> moduleShow (0, 1) MODULE NAME MODULE ID GROUP # TEXT START DATA START BSS START versionDotOWithALongName.o 0x616fc520 2 0x60532000 0x60534000 NO SEGMENT versionDotO15.o 0x616fbe98 3 0x60536000 0x60538000 NO SEGMENT ctdt.o 0x60534020 4 NO SEGMENT 0x61700000 NO SEGMENT value = 0 = 0x0

Display full name and path for a particular module (C shell):

-> moduleShow (0x616fc520, 1) MODULE NAME MODULE ID GROUP # TEXT START DATA START BSS START \_\_\_\_\_ \_\_\_\_\_ versionDotOWithALongName.o 0x616fc520 2 0x60532000 0x60534000 NO SEGMENT Size of text segment: 58 Size of data segment: 16 16 Size of data segment: Size of bss segment: 0 Total size 74 : Module path: huelgoat:/folk/joe/target/proj/linux\_gnu/default value =  $0 = 0 \times 0$ 

It is also possible to pass a module name instead of a module ID to the **moduleShow()** command. Thus :

-> moduleShow ("versionDotOWithALongName.o", 1)

would give the same output as above.

**RETURNS** OK, or ERROR if the module is not found.

ERRNO Not Available

**SEE ALSO** moduleLib, VxWorks Kernel Programmer's Guide: `Target Shell`, Workbench User's Guide: `Wind Shell`

#### mountdInit()

NAME mountdInit() – initialize the mount daemon

SYNOPSIS STATUS mountdInit

```
(
int priority, /* priority of the mount daemon */
int stackSize, /* stack size of the mount daemon */
FUNCPTR authHook, /* hook to run to authorize each request */
int nExports, /* maximum number of exported file systems */
int options /* Currently unused */
```

# DESCRIPTION This routine spawns a mount daemon if one does not already exist. Defaults for the *priority* and *stackSize* arguments are in the global variables **mountdPriorityDefault** and **mountdStackSizeDefault**, and are initially set to MOUNTD\_PRIORITY\_DEFAULT and MOUNTD\_STACKSIZE\_DEFAULT respectively.

Normally, no authorization checking is performed by either mountd or nfsd. To add authorization checking, set *authHook* to point to a routine declared as follows:

nfsstat routine

(		
int	progNum,	/* RPC program number */
int	versNum,	/* RPC program version number */
int	procNum,	/* RPC procedure number */
struct sockaddr_i	n clientAddr,	/* address of the client */
void *	mountdArg	<pre>/* argument of the call */</pre>

The mountdArg will be of type MOUNT3D\_ARGUMENT when versNum is 3 and it will be of type **MOUNTD\_ARGUMENT** when versNum is 1. The user routine must typecast the mountdArg accoringly and use it. The definitions of MOUNT3D\_ARGUMENT & **MOUNTD\_ARGUMENT** are available in **mountd.h** file.

The *authHook* callback must return **OK** if the request is authorized, and any defined NFS error code (usually NFSERR\_ACCESS) if not.

**RETURNS** OK, or **ERROR** if the mount daemon could not be correctly initialized.

ERRNO Not Available

SEE ALSO mountd

#### mqPxDescObjIdGet()

**NAME** mqPxDescObjIdGet() – returns the OBJ\_ID associated with a mqd\_t descriptor

VxWorks Kernel API Reference, 6.6 mqPxLibInit( )

DESCRIPTION	The message queue object identifier ( <b>OBJ_ID</b> ) is returned given a POSIX message queue descriptor (mqd_t).
RETURNS	<b>OBJ_ID</b> , or <b>NULL</b> if the message queue descriptor is invalid.
ERRNO	None
SEE ALSO	mqPxLib

## mqPxLibInit()

NAME	<b>mqPxLibInit()</b> – initialize the POSIX message queue library
SYNOPSIS	<pre>int mqPxLibInit   (    int hashSize /* not used */ )</pre>
DESCRIPTION	This routine initializes the POSIX message queue facility.
RETURNS	OK or ERROR.
ERRNO	N/A
SEE ALSO	mqPxLib

## mqPxShow()

NAME	<b>mqPxShow( )</b> – display message queue internals
SYNOPSIS	STATUS mqPxShow ( mqd_t mqDesc, int level )
DESCRIPTION	This routine displays information on a POSIX message queue
RETURNS	<b>OK</b> or <b>ERROR</b> if the descriptor is invalid.
ERRNO S\_objLib\_OBJ\_ID\_ERROR message queue is invalid.

SEE ALSO mqPxShow

### mqPxShowInit()

**NAME mqPxShowInit()** – initialize the POSIX message queue show facility

SYNOPSIS STATUS mqPxShowInit (void)

**DESCRIPTION** This routine links the POSIX message queue show routine into the VxWorks system. It is called automatically when this show facility is configured into VxWorks using either of the following methods:

If you use the configuration header files, define INCLUDE\_SHOW\_ROUTINES.

If you use the project facility, select INCLUDE\_POSIX\_MQ\_SHOW.

RETURNS OK always

ERRNO N/A

SEE ALSO mqPxShow

### mq\_close()

NAME	<b>mq_close( )</b> – close a message queue (POSIX)
SYNOPSIS	<pre>int mq_close    (     mqd_t mqdes /* message queue descriptor */ )</pre>
DESCRIPTION	This routine is used to indicate that the calling task is finished with the specified message queue <i>mqdes</i> . The <b>mq_close()</b> call deallocates any system resources allocated by the system for use by this task for its message queue. The behavior of a task that is blocked on either a <b>mq_send()</b> or <b>mq_receive()</b> is undefined when <b>mq_close()</b> is called. The <i>mqdes</i> parameter will no longer be a valid message queue ID.

#### SMP CONSIDERATIONS

	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	0 (OK) if the message queue is closed successfully, otherwise -1 (ERROR).
ERRNO	<b>EBADF</b> The <i>mqdes</i> argument is not a valid message queue descriptor.
SEE ALSO	mqPxLib, mq_open( )

# mq\_getattr()

NAME	<b>mq_getattr( )</b> – get message queue attributes (POSIX)
SYNOPSIS	<pre>int mq_getattr   (     mqd_t mqdes, /* message queue descriptor */     struct mq_attr * pMqStat /* buffer in which to return attributes */   )</pre>
DESCRIPTION	This routine gets status information and attributes associated with a specified message queue <i>mqdes</i> . Upon return, the following members of the <b>mq_attr</b> structure referenced by $pMqStat$ will contain the values set when the message queue was opened but with modifications made by subsequent calls to <b>mq_setattr()</b> :
	<b>q_flags</b> May be modified by <b>mq_setattr( )</b> .
	The following members were set at message queue creation:
	mq_maxmsg Maximum number of messages.
	mq_msgsize Maximum message size.
	The following member contains the current state of the message queue.
	mq_curmsgs The number of messages currently in the queue.
RETURNS	0 (OK) if message attributes can be determined, otherwise -1 (ERROR).

ERRNO	EBADF
	The <i>mqes</i> argument is not a valid message queue descriptor.

```
SEE ALSO
                mqPxLib, mq_open(), mq_send(), mq_setattr()
```

### mg notify()

NAME **mq\_notify()** – notify a task that a message is available on a queue (POSIX)

SYNOPSIS int mq\_notify (

)

mqd\_t mqdes, /\* message queue descriptor \*/ const struct sigevent \* pNotification /\* real-time signal \*/

If *pNotification* is not **NULL**, this routine attaches the specified *pNotification* request by the DESCRIPTION calling task to the specified message queue *mqdes* associated with the calling task. The real-time signal specified by *pNotification* will be sent to the task when the message queue changes from empty to non-empty. If a task has already attached a notification request to the message queue, all subsequent attempts to attach a notification to the message queue will fail. A task can get notifications from multiple messages queues.

> If *pNotification* is **NULL** and the task has previously attached a notification request to the message queue, the attached notification request is detached and the queue is available for another task to attach a notification request.

> If a notification request is attached to a message queue and any task is blocked in **mq\_receive()** waiting to receive a message when a message arrives at the queue, then the appropriate **mq\_receive()** will be completed and the notification request remains pending.

#### SMP CONSIDERATIONS

ERRNO

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

RETURNS 0 (OK) if successful, otherwise -1 (ERROR).

#### EBADF

The *mges* argument is not a valid message queue descriptor.

#### EBUSY

A task is already registered for notification by the message queue.

#### EINVAL

This task is trying to remove the registration of another task.

VxWorks Kernel API Reference, 6.6 mq\_open()

SEE ALSO mqPxLib, mq\_open(), mq\_send()

## mq\_open()

NAME	<b>mq_open()</b> – open a message queue (POSIX)
SYNOPSIS	<pre>mqd_t mq_open   (     const char * mqName, /* name of queue to open */     int oflags, /* open flags */     /* extra optional parameters */   )</pre>
DESCRIPTION	This routine establishes a connection between a named message queue and the calling task. After a call to <b>mq_open()</b> , the task can reference the message queue using the address returned by the call. The message queue remains usable until the queue is closed by a successful call to <b>mq_close()</b> .
	The message queue must have a name. <b>NULL</b> and empty strings result in <b>EINVAL</b> . If the <i>name</i> begins with the slash character, then it is treated as a public message queue. All RTPs can open their own references to the public message queue by using its name. If the <i>name</i> does not begin with the slash character, then it is treated as a private message queue and RTPs cannot get access to it.
	The following flag bits can be set in <i>oflags</i> :
	O_RDONLY Open the message queue for receiving messages. The task can use the returned message queue descriptor with <b>mq_receive( )</b> , but not <b>mq_send( )</b> .
	O_WRONLY Open the message queue for sending messages. The task can use the returned message queue descriptor with <b>mq_send()</b> , but not <b>mq_receive()</b> .

#### O\_RDWR

Open the queue for both receiving and sending messages. The task can use any of the functions allowed for O\_RDONLY and O\_WRONLY.

Any combination of the following flags can be specified in *oflags*. These control whether the message queue is created or merely accessed by the mq\_open() call.

#### O\_CREAT

This flag is used to create a message queue if it does not already exist. If O\_CREAT is set and the message queue already exists, then O\_CREAT has no effect except as noted below under O\_EXCL. Otherwise, mq\_open() creates a message queue. The O\_CREAT flag requires a third and fourth argument: *mode*, which is of type **mode\_t**, and *pAttr*, which is of type pointer to an **mq\_attr** structure. The value of *mode* has no effect in this implementation. If *pAttr* is NULL, the message queue is created with a MQ\_NUM\_MSG\_DEFAULT messages of size MQ\_MSG\_SIZE\_DEFAULT. If *pAttr* is non-NULL, the message queue attributes **mq\_maxmsg** and **mq\_msgsize** are set to the values of the corresponding members in the **mq\_attr** structure referred to by *pAttr*; if either attribute is less than or equal to zero, an error is returned and errno is set to EINVAL.

#### O\_EXCL

This flag is used to test whether a message queue already exists. If **O\_EXCL** and **O\_CREAT** are set, **mq\_open()** fails if the message queue name exists.

#### O\_NONBLOCK

The setting of this flag is associated with the open message queue descriptor. If this flag is set, then the **mq\_send()** and **mq\_receive()** do not wait for resources or messages that are not currently available. Instead, they fail with errno set to **EAGAIN**.

The mq\_open() call does not add or remove messages from the queue.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

NOTE	Some POSIX functionality is not yet supported:
	- A message queue cannot be closed with calls to <b>_exit( )</b> or <b>exec( )</b> .
	- A message queue cannot be implemented as a file.
	- Message queue names will not appear in the file system.
RETURNS	A message queue descriptor, otherwise -1 (ERROR).
ERRNO	EEXIST O_CREAT and O_EXCL are set and the message queue already exists.
	ENOENT O_CREAT is not set and the message queue does not exist.
	<b>ENOSPC</b> There is insufficient space for the creation of the new message queue.
	EINVAL
	- The specified <i>name</i> is invalid.
	- An invalid combination of <i>oflags</i> is specified.
	- <b>O_CREAT</b> is specified in <i>oflags</i> , the value of <i>pAttr</i> is not <b>NULL</b> and either <i>mq_maxmsg</i> or <i>mq_msgsize</i> is less than or equal to zero.

SEE ALSO mqPxLib, mq\_send(), mq\_receive(), mq\_close(), mq\_setattr(), mq\_getattr(), mq\_unlink()

# mq\_receive()

NAME	<b>mq_receive( )</b> – receive a message from a message queue (POSIX)
SYNOPSIS	<pre>ssize_t mq_receive   (    mqd_t mqdes, /* message queue descriptor */    void * pMsg, /* buffer to receive message */    size_t msgLen, /* size of buffer, in bytes */    int * pMsgPrio /* if not NULL, priority of message */   )</pre>
DESCRIPTION	This routine receives the oldest of the highest priority message from the message queue specified by <i>mqdes</i> . If the size of the buffer in bytes, specified by the <i>msgLen</i> argument, is less than the <b>mq_msgsize</b> attribute of the message queue, <b>mq_receive()</b> will fail and return an error. Otherwise, the selected message is removed from the queue and copied to <i>pMsg</i> .
	If <i>pMsgPrio</i> is not <b>NULL</b> , the priority of the selected message will be stored in <i>pMsgPrio</i> .
	If the message queue is empty and <b>O_NONBLOCK</b> is not set in the message queue's description associated with <i>mqdes</i> , <b>mq_receive()</b> will block until a message is added to the message queue, or until it is interrupted by a signal. If more than one task is waiting to receive a message when a message arrives at an empty queue, the task of highest priority will be selected to receive the message. If the specified message queue is empty and <b>O_NONBLOCK</b> is set in the message queue's description associated with <i>mqdes</i> , no message is removed from the queue, and <b>mq_receive()</b> returns an error.
	The non-negative size value of the msgLen is not limited, and if a negative value is specified for msgLen, the negativitiy of that value will be ignored.
SMP CONSIDERAT	IONS
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	The length of the selected message in bytes, otherwise -1 (ERROR).
ERRNO	EAGAIN O_NONBLOCK was set in the message queue description associated with <i>mqdes</i> , and the specified message queue is empty.

#### EBADF

The *mqdes* argument is not a valid message queue descriptor open for for reading.

#### EMSGSIZE

The specified message buffer size, *msgLen*, is less than the message size attribute of the message queue.

#### EINVAL

The *pMsg* pointer is invalid.

#### EINTR

A signal was received while blocking on the message queue. This error only occurs for an RTP task.

SEE ALSO mqPxLib, mq\_send()

### mq\_send()

**NAME mq\_send()** – send a message to a message queue (POSIX)

SYNOPSIS

```
int mg_send
  (
    mqd_t mqdes, /* message queue descriptor */
    const void * pMsg, /* message to send */
    size_t msgLen, /* size of message, in bytes */
    int msgPrio /* priority of message */
    )
```

**DESCRIPTION** This routine adds the message pMsg to the message queue mqdes. The msgLen parameter specifies the length of the message in bytes pointed to by pMsg. The value of pMsg must be less than or equal to the **mq\_msgsize** attribute of the message queue, or **mq\_send()** will fail.

If the message queue is not full, **mq\_send()** will behave as if the message is inserted into the message queue at the position indicated by the *msgPrio* argument. A message with a higher numeric value for *msgPrio* is inserted before messages with a lower value. The value of *msgPrio* must be less than **MQ\_PRIO\_MAX**.

If the specified message queue is full and O\_NONBLOCK is not set in the message queue, **mq\_send()** will block until space becomes available to queue the message, or until it is interrupted by a signal. If the message queue is full and O\_NONBLOCK is set in the message queue's descriptions associated with *mqdes*, the message is not queued, and **mq\_send()** returns an error.

#### USE BY INTERRUPT SERVICE ROUTINES

This routine can be called by interrupt service routines as well as by tasks. This is one of the primary means of communication between an interrupt service routine and a task. If

**mq\_send()** is called from an interrupt service routine, it will behave as if the **O\_NONBLOCK** flag were set.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

**RETURNS** 0 (OK), otherwise -1 (ERROR).

#### ERRNO EAGAIN

**O\_NONBLOCK** was set in the message queue description associated with *mqdes*, and the specified message queue is full.

#### EBADF

The *mqdes* argument is not a valid message queue descriptor open for for writing.

#### EMSGSIZE

The specified message length, *msgLen*, exceeds the message size attribute of the message queue.

#### EINVAL

- The value of *msgPrio* is greater than or equal to MQ\_PRIO\_MAX.
- The *pMsg* pointer is invalid.

#### EINTR

The request has been interrupted by a signal.

#### SEE ALSO mqPxLib, mq\_receive()

### mq\_setattr( )

NAME	<b>mq_setattr()</b> – set message queue attributes (POSIX)
SYNOPSIS	<pre>int mq_setattr   (     mqd_t mqdes, /* message queue descriptor */     const struct mq_attr * pMqStat, /* new attributes */     struct mq_attr * pOldMqStat /* old attributes */     )</pre>
DESCRIPTION	This routine sets attributes associated with the specified message queue <i>mqdes</i> .

The message queue attributes corresponding to the following members defined in the **mq\_attr** structure are set to the specified values upon successful completion of the call:

#### mq\_flags

The value of the **O\_NONBLOCK** flag.

If *pOldMqStat* is non-NULL, **mq\_setattr()** will store, in the location referenced by *pOldMqStat*, the previous message queue attributes and the current queue status. These values are the same as would be returned by a call to **mq\_getattr()** at that point.

#### SMP CONSIDERATIONS

	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	0 ( <b>OK</b> ) if attributes are set successfully, otherwise -1 ( <b>ERROR</b> ).
ERRNO	<b>EBADF</b> The <i>mqes</i> argument is not a valid message queue descriptor.
SEE ALSO	mqPxLib, mq_open( ), mq_send( ), mq_getattr( )

# mq\_unlink()

NAME	<b>mq_unlink()</b> – remove a message queue (POSIX)
SYNOPSIS	<pre>int mq_unlink   (     const char * mqName /* name of message queue */ )</pre>
DESCRIPTION	This routine removes the message queue named by the pathname <i>mqName</i> . After a successful call to <b>mq_unlink()</b> , a call to <b>mq_open()</b> on the same message queue will fail if the flag <b>O_CREAT</b> is not set. If one or more tasks have the message queue open when <b>mq_unlink()</b> is called, removal of the message queue is postponed until all references to the message queue have been closed by <b>mq_close()</b> .
SMP CONSIDERATIO	DNS
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	0 (OK) if the message queue is unlinked successfully, otherwise -1 (ERROR).
ERRNO	<b>ENOENT</b> A message queue with the specified name, <i>mqName</i> , does not exist.

VxWorks Kernel API Reference, 6.6 msgQClose()

SEE ALSO mqPxLib, mq\_close(), mq\_open()

# msgQClose()

NAME	<b>msgQClose( )</b> – close a named message queue	
SYNOPSIS	STATUS msgQClose ( MSG_Q_ID msgQId /* message queue ID to close */ )	
DESCRIPTION	This routine closes a named message queue and decrements its reference counter. In the case where the counter becomes zero, the message queue is deleted if:	
	- It has been already removed from the name space by a call to <b>msgQUnlink()</b> .	
	- It was created with the OM_DESTROY_ON_LAST_CALL option.	
	This routine is not ISR callable.	
SMP CONSIDERATIONS		
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.	
RETURNS	OK, or ERROR if unsuccessful.	
ERRNO	S_objLib_OBJ_ID_ERROR The message queue ID is invalid.	
	S_objLib_OBJ_INVALID_ARGUMENT The message queue ID is NULL.	
	S_objLib_OBJ_OPERATION_UNSUPPORTED The message queue is not named.	
	S_objLib_OBJ_DESTROY_ERROR An error was detected while deleting the message queue.	
	S_intLib_NOT_ISR_CALLABLE This routine must not be called from an ISR.	
SEE ALSO	msgQOpen, msgQOpen( ), msgQUnlink( )	

<b>msgQCreate()</b> – create and initialize a message queue	
<pre>MSG_0_ID msgQCreate    (     int maxMsgs,</pre>	
This routine creates a message queue capable of holding up to <i>maxMsgs</i> messages, each up to <i>maxMsgLength</i> bytes long. The routine returns a message queue ID used to identify the created message queue in all subsequent calls to routines in this library. The queue can be created with the following options:	
MSG_Q_FIFO (0x00) Queue pended tasks in FIFO order.	
<b>MSG_Q_PRIORITY</b> (0x01) Queue pended tasks in priority order.	
MSG_Q_EVENTSEND_ERR_NOTIFY (0x02) When a message is sent, if a task is registered for events and the actual sending of events fails, a value of ERROR is returned and errno is set. This option is off by default.	
MSG_Q_INTERRUPTIBLE (0x04) Signal sent to a RTP task pended on a message queue created with this option, would make the task ready and return ERROR with errno set to EINTR. This option has no affect for a kernel task pended on the same message queue created with this option. This option is off by default.	
SMP CONSIDERATIONS	
This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.	
MSG_Q_ID, or NULL if error	
S_memLib_NOT_ENOUGH_MEMORY There is not enough memory to create the queue as specified.	
S_intLib_NOT_ISR_CALLABLE This routine cannot be called from interrupt level.	
S_msgQLib_INVALID_MSG_LENGTH Negative maxMsgLength specified.	

VxWorks Kernel API Reference, 6.6 msgQDelete()

S\_msgQLib\_INVALID\_MSG\_COUNT Negative maxMsgs specified.

S\_msgQLib\_INVALID\_QUEUE\_TYPE Invalid pending queue type specified.

S\_msgQLib\_ILLEGAL\_OPTIONS Illegal option bits were specified.

SEE ALSO msgQLib, msgQSmLib

# msgQDelete()

NAME	msgQDelete() – delete a message queue
SYNOPSIS	<pre>STATUS msgQDelete   (     MSG_Q_ID msgQId /* message queue to delete */   )</pre>
DESCRIPTION	This routine deletes a message queue. All tasks pending on either <b>msgQSend()</b> or <b>msgQReceive()</b> , or pending for the reception of events meant to be sent from the message queue, unblock and return <b>ERROR</b> . When this function returns, <i>msgQld</i> is no longer a valid message queue ID.
SMP CONSIDERATIONS	
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	OK on success or ERROR otherwise
ERRNO	S_objLib_OBJ_ID_ERROR The message queue ID is invalid.
	S_intLib_NOT_ISR_CALLABLE The routine cannot be called from interrupt level.
	S_smObjLib_NO_OBJECT_DESTROY Deleting a shared message queue is not permitted.
	S_objLib_OBJ_OPERATION_UNSUPPORTED Deleting a named message queue is not permitted.
SEE ALSO	msgQLib, msgQSmLib

```
msgQEvStart() - start the event notification process for a message queue
NAME
SYNOPSIS
                 STATUS msgQEvStart
                      (
                      MSG_Q_ID msgQId, /* msg Q for which to register events */
                      UINT32 events, /* 32 possible events
                                                                                     */
                      UINT8 options /* event-related msg Q options
                                                                                     */
                      )
DESCRIPTION
                 This routine turns on the event notification process for a given message queue, registering
                 the calling task on that queue. When a message arrives on the queue and no receivers are
                 pending, the events specified are sent to the registered task. A task can always overwrite its
                 own registration.
                 The events are user-defined. For more information, see the reference entry for eventLib.
                 The options parameter is used for three user options:
                      Specify whether the events are to be sent only once or every time a message arrives until
                      msgQEvStop() is called.
                      Specify if another task can subsequently register itself while the calling task is still
                      registered. If so specified, the existing task registration will be overwritten without any
                      warning.
                      Specify if events are to be sent immediately in the case a message is waiting to be picked
                      up.
                 Here are the possible values to be used in the option field:
                 EVENTS_SEND_ONCE (0x1)
                      The message queue will send the events only once.
                 EVENTS_ALLOW_OVERWRITE (0x2)
                      Subsequent registrations from other tasks can overwrite the current one.
                 EVENTS SEND IF FREE (0x4)
                      The registration process will send events if a message is present on the message queue
                      when msgQEvStart() is called.
                 EVENTS_OPTIONS_NONE (0x0)
                      Must be passed to the options parameter if none of the other three options are used.
WARNING
                 This routine cannot be called from interrupt level.
WARNING
                 Task preemption can allow a msgQDelete() to be performed between the calls to
                 msgQEvStart() and eventReceive(). This would prevent the task from ever receiving the
                 events wanted from the message queue.
```

#### SMP CONSIDERATIONS

	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	OK on success, or ERROR
ERRNO	<b>S_objLib_OBJ_ID_ERROR</b> The message queue ID is invalid.
	S_eventLib_ALREADY_REGISTERED A task is already registered on the message queue.
	S_intLib_NOT_ISR_CALLABLE This routine cannot be called from interrupt level.
	S_eventLib_EVENTSEND_FAILED The user chose to send events immediately and that operation failed.
	S_eventLib_ZERO_EVENTS The user passed in a value of zero to the <i>events</i> parameter.

SEE ALSO msgQEvLib, eventLib, msgQLib, msgQEvStop()

# msgQEvStop()

NAME	<b>msgQEvStop()</b> – stop the event notification process for a message queue	
SYNOPSIS	STATUS msgQEvStop ( MSG_Q_ID msgQId )	
DESCRIPTION	This routine turns off the event notification process for a given message queue. This allows another task to register itself for event notification on that particular message queue. The routine must be called by the task that is already registered on that particular message queue.	
SMP CONSIDERATIONS		
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.	
RETURNS	OK on success, or ERROR.	

ERRNO	S_objLib_OBJ_ID_ERROR The message queue ID is invalid.
	S_intLib_NOT_ISR_CALLABLE The routine was called from interrupt level.
	S_eventLib_TASK_NOT_REGISTERED The routine was not called by the registered task.
SEE ALSO	msgQEvLib, eventLib, msgQLib, msgQEvStart( )

### msgQInfoGet()

NAME msgQInfoGet() – get information about a message queue SYNOPSIS STATUS msgOInfoGet ( MSG\_Q\_ID msgQId, /\* message queue to query \*/ MSG\_Q\_INFO \* pInfo /\* where to return msg info \*/ ) This routine gets information about the state and contents of a message queue. The DESCRIPTION parameter *plnfo* is a pointer to a structure of type **MSG\_Q\_INFO** defined in **msgQLib.h** as follows: typedef struct /\* MSG O INFO \*/ { int numMsgs; /\* OUT: number of messages queued int numTasks; /\* OUT: number of tasks waiting on msg q \* / \* / int sendTimeouts; /\* OUT: count of send timeouts \*/ recvTimeouts; /\* OUT: count of receive timeouts \*/ int options; /\* OUT: options with which msg q was created \*/ maxMsgs; /\* OUT: max messages that can be queued \*/ int int maxMsgLength; /\* OUT: max byte length of each message int \*/ \*/ int taskIdListMax; /\* IN: max tasks to fill in taskIdList int \* taskIdList; /\* PTR: array of task IDs waiting on msg q \*/ msgListMax; /\* IN: max msgs to fill in msg lists \*/ int. char \*\* msgPtrList; /\* PTR: array of msg ptrs queued to msg q \*/ int \* msgLenList; /\* PTR: array of lengths of msgs \*/ } MSG O INFO;

If a message queue is empty, there may be tasks blocked on receiving. If a message queue is full, there may be tasks blocked on sending. This can be determined as follows:

- If **numMsgs** is 0, then **numTasks** indicates the number of tasks blocked on receiving.
- If **numMsgs** is equal to **maxMsgs**, then **numTasks** is the number of tasks blocked on sending.

- If numMsgs is greater than 0 but less than maxMsgs, then numTasks
---

A list of pointers to the messages queued and their lengths can be obtained by setting **msgPtrList** and **msgLenList** to the addresses of arrays to receive the respective lists, and setting **msgListMax** to the maximum number of elements in those arrays. If either list pointer is **NULL**, no data is returned for that array.

No more than **msgListMax** message pointers and lengths are returned, although **numMsgs** is always returned with the actual number of messages queued.

For example, if the caller supplies a **msgPtrList** and **msgLenList** with room for 10 messages and sets **msgListMax** to 10, but there are 20 messages queued, then the pointers and lengths of the first 10 messages in the queue are returned in **msgPtrList** and **msgLenList**, but **numMsgs** is returned with the value 20.

A list of the task IDs of tasks blocked on the message queue can be obtained by setting **taskIdList** to the address of an array to receive the list, and setting **taskIdListMax** to the maximum number of elements in that array. If **taskIdList** is **NULL**, then no task IDs are returned. No more than **taskIdListMax** task IDs are returned, although **numTasks** is always returned with the actual number of tasks blocked.

For example, if the caller supplies a **taskIdList** with room for 10 task IDs and sets **taskIdListMax** to 10, but there are 20 tasks blocked on the message queue, then the IDs of the first 10 tasks in the blocked queue are returned in **taskIdList**, but **numTasks** is returned with the value 20.

The tasks returned in **taskIdList** may be blocked for either send or receive. As noted above this can be determined by examining **numMsgs**.

The variables **sendTimeouts** and **recvTimeouts** are the counts of the number of times **msgQSend()** and **msgQReceive()** respectively returned with a timeout.

The variables **options**, **maxMsgs**, and **maxMsgLength** are the parameters with which the message queue was created.

- WARNING The information returned by this routine is not static and may be obsolete by the time it is examined. In particular, the lists of task IDs or message pointers may no longer be valid. However, the information is obtained atomically; it is an accurate snapshot of the state of the message queue at the time of the call. This information is generally used for debugging purposes only.
- **WARNING** The current implementation of this routine locks out interrupts while obtaining the information. This can compromise the overall interrupt latency of the system. Generally this routine should be used for debugging purposes only.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted.

**RETURNS** OK or ERROR.

ERRNO	S_smObjLib_NOT_INITIALIZED The shared memory message queue library (VxMP Option) was not initialized.
	S_objLib_OBJ_ID_ERROR The message queue ID is invalid.
SEE ALSO	msgQInfo

### msgQInitialize()

msgQInitialize() - initialize a pre-allocated message queue NAME SYNOPSIS MSG\_Q\_ID msgQInitialize ( int maxMsgLength, /\* max bytes in a message \*/ /\* message queue options \*/ int options ) DESCRIPTION This routine initializes a pre-allocated message queue structure and message pool memory. The message pool memory must be capable of holding up to maxMsgs messages, each of up to maxMsgLength bytes long. Parameter pMsgPool points to the buffer to be used for holding queued messages. *pMsgPool* must point to a 4 byte aligned buffer whose size is (*maxMsgs* \* **MSG\_NODE\_SIZE** (*maxMsgLength*)). The queue can be created with the following options: MSG O FIFO (0x00) Queue pended tasks in FIFO order. MSG\_Q\_PRIORITY (0x01) Queue pended tasks in priority order. MSG\_Q\_EVENTSEND\_ERR\_NOTIFY (0x02) When a message is sent, if a task is registered for events and the actual sending of events fails, a value of ERROR is returned and errno is set. This option is off by default.  $MSG_Q_INTERRUPTIBLE$  (0x04) A signal sent to an RTP task pended on a message queue created with this option would make the task ready and return ERROR with errno set to EINTR. This option has no affect for kernel tasks pended on the same message queue created with this option. This option is disabled by default.

VxWorks Kernel API Reference, 6.6 msgQNumMsgs()

#### SMP CONSIDERATIONS

RETURNS

ERRNO

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

The following example illustrates use of the VX\_MSG\_Q macro and this function together to instantiate a message queue statically (without using any dynamic memory allocation):

```
#include <vxWorks.h>
#include <msgQLib.h>

VX_MSG_Q(myMsgQ,100,16); /* declare the msgQ */
MSG_Q_ID myMsgQId; /* MsgQ ID to send/receive messages */
STATUS initializeFunction (void)
{
    if ((myMsgQId = msgQInitialize (myMsgQ, 100, 16, options)) == NULL)
        return (ERROR); /* initialization failed */
    else
        return (OK);
}
The MSG_Q_ID, or NULL on error.
N/A
```

SEE ALSO msgQLib, msgQCreate()

### msgQNumMsgs()

NAME	msgQNumMsgs() – get the number of messages queued to a message queue	
SYNOPSIS	<pre>int msgQNumMsgs   (   FAST MSG_Q_ID msgQId /* message queue to examine */ )</pre>	
DESCRIPTION	This routine returns the number of messages currently queued to a specified message queue.	
SMP CONSIDERATIONS		
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.	
RETURNS	The number of messages queued, or <b>ERROR</b>	

ERRNO	<b>S_smObjLib_NOT_INITIALIZED</b> The shared memory message queue library (VxMP Option) was not initialized.
	S_objLib_OBJ_ID_ERROR The message queue ID is invalid.
SEE ALSO	msgQLib, msgQSmLib

### msgQOpen()

NAME msgQOpen() – open a message queue

**DESCRIPTION** This routine opens a message queue, which means that it searches the name space and returns the **MSG\_Q\_ID** of an existing message queue with *name*. If none is found, it creates a new one with *name* according on the flags set in the mode parameter.

The argument *name* is mandatory. NULL or empty strings are not allowed.

There are two name spaces available in which **msgQOpen()** can perform the search. The name space searched is dependent upon the first character in the *name* parameter. When this character is a forward slash /, the **public** name space is searched; otherwise the **private** name space is searched. Similarly, if a message queue is created, the first character in *name* specifies the name space that contains the message queue.

Message queues created by this routine can not be deleted with **msgQDelete()**. Instead, a **msgQClose()** must be issued for every **msgQOpen()**. Then the message queue is deleted when it is removed from the name space by a call to **msgQUnlink()**. Alternatively, the message queue can be previously removed from the name space, and deleted during the last **msgQClose()**.

A description of the *mode* and *context* arguments follows. See the reference entry for **msgQCreate()** for a description of the remaining arguments.

mode

The mode parameter passed to this routine consists of the opening flags, which can be set using a bitwise-OR:

#### OM\_CREATE

Create a message queue if none is found.

#### OM\_EXCL

When set jointly with the **OM\_CREATE** flag, create a new message queue without trying to open an existing one. The call fails if *name* causes a name clash. This flag has no effect if the flag **OM\_CREATE** is not set.

#### OM\_DELETE\_ON\_LAST\_CLOSE

Only used when a message queue is created. If set, the message queue is deleted during the last **msgQClose()** call, independently of whether **msgQUnlink()** was previously called or not.

#### context

Context value assigned to the created message queue. This value is not actually used by VxWorks. Instead, the context value can be used by OS extensions to implement object permissions, for example.

Unlike private objects, a public message queue is not automatically reclaimed when an application terminates. Note that nevertheless, a **msgQClose()** is issued on every application's outstanding **msgQOpen()**. Therefore, a public message queue can effectively be deleted, if during this process it is closed for the last time, and it is already unlinked or it was created with the **OM\_DELETE\_ON\_LAST\_CLOSE** flag.

This routine is not ISR callable.

#### SMP CONSIDERATIONS

ERRNO

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the
implementation so it is the responsibility of the caller to ensure they are complied with.
Future implementations may enforce these restrictions.

**RETURNS** The MSG\_Q\_ID of the opened message queue, or NULL if unsuccessful.

#### S\_objLib\_OBJ\_INVALID\_ARGUMENT

An invalid option was specified in the mode argument or name is invalid.

#### S\_msgQLib\_INVALID\_MSG\_LENGTH Negative maxMsgLength specified.

#### S\_msgQLib\_INVALID\_MSG\_COUNT Negative maxMsgs specified.

#### S\_objLib\_OBJ\_NOT\_FOUND

The **OM\_CREATE** flag was not set in the *mode* argument and a message queue matching *name* was not found.

#### S\_objLib\_OBJ\_NAME\_CLASH

The **OM\_CREATE** and **OM\_EXCL** flags were set and a name clash was detected when creating the message queue.

S\_intLib\_NOT\_ISR\_CALLABLE This routine must not be called from an ISR.

SEE ALSO msgQOpen, msgQUnlink(), msgQClose()

### msgQOpenInit( )

**NAME** msgQOpenInit() – initialize the message queue open facility

SYNOPSIS void msgQOpenInit (void)

DESCRIPTION This routine links the message queue creation routine with the open facility into the VxWorks system. It is called automatically when the message queue facility is configured into VxWorks by either defining the INCLUDE\_OBJ\_OPEN and INCLUDING\_MSG\_Q components in config.h or selecting INCLUDE\_OBJ\_OPEN and INCLUDING\_MSG\_Q in the project facility.

RETURNS N/A

ERRNO N/A

SEE ALSO msgQOpen

### msgQReceive()

NAME	msgQReceive() – receive a message from a message queue
SYNOPSIS	<pre>int msgQReceive   (    MSG_Q_ID msgQId, /* message queue from which to receive */    char * buffer, /* buffer to receive message */    UINT maxNBytes, /* length of buffer */    int timeout /* ticks to wait */   )</pre>
DESCRIPTION	This routine receives a message from the message queue <i>msgQld</i> . The received message

**DESCRIPTION** This routine receives a message from the message queue *msgQld*. The received message is copied into the specified *buffer*, which is *maxNBytes* in length. If the message is longer than *maxNBytes*, the remainder of the message is discarded (no error indication is returned).

	The <i>timeout</i> parameter specifies the number of ticks to wait for a message to be sent to the queue, if no message is available when <b>msgQReceive()</b> is called. The <i>timeout</i> parameter can also have the following special values:
	NO_WAIT (0) Return immediately, whether a message has been received or not.
	WAIT_FOREVER (-1) Never time out.
WARNING	This routine must not be called by interrupt service routines.
SMP CONSIDERATIO	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	The number of bytes copied to <i>buffer</i> , or <b>ERROR</b>
ERRNO	S_smObjLib_NOT_INITIALIZED The shared memory message queue library (VxMP Option) was not initialized.
	S_objLib_OBJ_ID_ERROR The message queue ID is invalid.
	S_objLib_OBJ_DELETED The message queue was deleted while the calling task was pended.
	S_objLib_OBJ_UNAVAILABLE No message was available and the NO_WAIT timeout was specified.
	<b>S_objLib_OBJ_TIMEOUT</b> A timeout occurred while waiting for a message to become available.
	S_msgQLib_INVALID_MSG_LENGTH The message length exceeds the limit.
	S_intLib_NOT_ISR_CALLABLE This routine cannot be called from interrupt level.
SEE ALSO	msgQLib, msgQSmLib, msgQSend( )

### msgQSend()

NAME

msgQSend() – send a message to a message queue

#### SYNOPSIS STATUS msgQSend ( MSG\_Q\_ID msgQId, /\* message queue on which to send \*/ char \* buffer, /\* message to send \*/ UINT nBytes, /\* length of message \*/ int timeout, /\* ticks to wait \*/ int priority /\* MSG\_PRI\_NORMAL or MSG\_PRI\_URGENT \*/ )

**DESCRIPTION** This routine sends the message in *buffer* of length *nBytes* to the message queue *msgQld*. If any tasks are already waiting to receive messages on the queue, the message is immediately delivered to the first waiting task. If no task is waiting to receive messages, the message is saved in the message queue and, if a task has previously registered to receive events from the message queue, these events are sent in the context of this call. This may result in the unpending of the task waiting for the events. If the message queue fails to send events, and if it was created using the MSG\_Q\_EVENTSEND\_ERR\_NOTIFY option, ERROR is returned even though the message was successfully sent to the queue.

The *timeout* parameter specifies the number of ticks to wait for adding its message to the queue if the message queue is full. The *timeout* parameter can also have the following special values:

#### NO\_WAIT (0)

Return immediately, even if the message has not been sent.

#### WAIT\_FOREVER (-1)

Never time out.

The *priority* parameter specifies the priority of the message being sent. The possible values are:

#### MSG\_PRI\_NORMAL (0)

Normal priority; add the message to the tail of the list of queued messages.

#### MSG\_PRI\_URGENT (1)

Urgent priority; add the message to the head of the list of queued messages.

#### USE BY INTERRUPT SERVICE ROUTINES

This routine can be called by interrupt service routines as well as by tasks. This is one of the primary means of communication between an ISR and a task. When called from an ISR, *timeout* must be **NO\_WAIT**.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

RETURNS OK on success or ERROR otherwise

ERRNO S\_smObjLib\_NOT\_INITIALIZED

The shared memory message queue library (VxMP Option) was not initialized.

**S\_objLib\_OBJ\_ID\_ERROR** The message queue ID is invalid.

S\_objLib\_OBJ\_DELETED The message queue was deleted while the calling task was pended.

S\_objLib\_OBJ\_UNAVAILABLE

No free buffer space was available and the NO\_WAIT timeout was specified.

#### S\_objLib\_OBJ\_TIMEOUT

A timeout occurred while waiting for buffer space to become available.

- S\_msgQLib\_INVALID\_MSG\_LENGTH The message length exceeds the limit.
- S\_msgQLib\_NON\_ZERO\_TIMEOUT\_AT\_INT\_LEVEL The routine was called from an ISR with a non-zero timeout.

#### S\_eventLib\_EVENTSEND\_FAILED

The message queue failed to send events to the registered task. This **errno** value can occur only if the message queue was created with the **MSG\_Q\_EVENTSEND\_ERR\_NOTIFY** option.

SEE ALSO msgQLib, msgQSmLib, msgQEvStart()

### msgQShow()

NAME	msgQShow() – show information about a message queue
SYNOPSIS	<pre>STATUS msgQShow   (    MSG_Q_ID msgQId, /* message queue to display */    int level /* 0 = summary, 1 = details */   )</pre>
DESCRIPTION	This routine displays the state, and optionally the contents, of a message queue.

A summary of the state of the message queue is displayed as follows:

Message Queue Id Task Queuing Message Byte Len Messages Max Messages Queued Receivers Blocked Send timeouts Receive timeouts Options	: 0x3f8c20 : FIFO : 150 : 50 : 0 : 1 : 0 : 0 : 0x1	MSG_Q_FIFO
VxWorks Events		
Registered Task Event(s) to Send Options	: 0x3f5c70 : 0x1 : 0x7	(t1) EVENTS_SEND_ONCE EVENTS_ALLOW_OVERWRITE EVENTS_SEND_IF_FREE

If *level* is 1, then more detailed information is displayed. If messages are queued, they are displayed as follows:

Messages queued: # address length value 1 0x123eb204 4 0x0000001 0x12345678

If tasks are blocked on the queue, they are displayed as follows:

Receivers blocked:

NAME	TID	PRI	DELAY
tExcTask	3fd678	0	21

- RETURNS OK or ERROR.
- ERRNO S\_smObjLib\_NOT\_INITIALIZED The shared memory message queue library (VxMP Option) was not initialized.
- SEE ALSO msgQShow, windsh

### msgQShowInit( )

- NAME msgQShowInit() initialize the message queue show facility
- SYNOPSIS void msgQShowInit (void)
- **DESCRIPTION** This routine links the message queue show facility into the VxWorks system. It is called automatically when the message queue show facility is configured into VxWorks using either of the following methods:

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- Using the configuration header files, define INCLUDE\_SHOW\_ROUTINES in config.h.
- Using the project facility, select INCLUDE\_MSG\_Q\_SHOW.

RETURNS N/A

ERRNO Not Available

SEE ALSO msgQShow

## msgQSmCreate()

NAME	msgQSmCreate() – create and initialize a shared memory message queue (VxMP Option)	
SYNOPSIS	<pre>MSG_Q_ID msgQSmCreate   (    int maxMsgs,</pre>	
DESCRIPTION	This routine creates a shared memory message queue capable of holding up to <i>maxMsg</i> messages, each up to <i>maxMsgLength</i> bytes long. It returns a message queue ID used to identify the created message queue. The queue can only be created with the option <b>MSG_Q_FIFO</b> (0), thus queuing pended tasks in FIFO order.	
	The global message queue identifier returned can be used directly by generic message queue handling routines in msgQLib msgQSend(), msgQReceive(), and msgQNumMsgs() and by the show routines show() and msgQShow().	
	If there is insufficient memory to store the message queue structure in the shared memory message queue partition or if the shared memory system pool cannot handle the requested message queue size, shared memory message queue creation will fail with <b>errno</b> set to <b>S_memLib_NOT_ENOUGH_MEMORY</b> . This problem can be solved by incrementing the value of <b>SM_OBJ_MAX_MSG_Q</b> and/or the shared memory objects dedicated memory size <b>SM_OBJ_MEM_SIZE</b> .	
	Before this routine can be called, the shared memory objects facility must be initialized (see msgQSmLib).	
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.	
RETURNS	MSG_Q_ID, or NULL if error.	

ERRNO	S_intLib_NOT_ISR_CALLABLE Routine has been called from ISR.
	S_objLib_OBJ_ID_ERROR The shared memory message queue partition has not been initialized properly.
	S_memLib_NOT_ENOUGH_MEMORY Can't allocate shared memory message queue object.
	S_msgQLib_INVALID_QUEUE_TYPE Incorrect message queue pend queue type specified.
	S_msgQLib_INVALID_MSG_COUNT Incorrect number (negative) of messages specified.
	S_msgQLib_INVALID_MSG_LENGTH Incorrect length (negative) of messages specified.
	<b>S_smObjLib_LOCK_TIMEOUT</b> Can't get the lock on the shared memory message queue partition in time.
SEE ALSO	msgQSmLib, smObjLib, msgQLib, msgQShow

### msgQUnlink()

NAME msgQUnlink() – unlink a named message queue SYNOPSIS STATUS msgOUnlink ( const char \* name /\* name of message queue to unlink \*/ ) This routine removes a message queue from the name space, and marks it as ready for DESCRIPTION deletion on the last msgQClose(). In the case where there is no outstanding msgQOpen() call, the message queue is deleted immediately. After a message queue is unlinked, subsequent calls to **msgQOpen()** using *name* will not be able to find the message queue, even if it has not been deleted yet. Instead, a new message queue could be created if msgQOpen() is called with the OM\_CREATE flag. This routine is not ISR callable. SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

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RETURNS	OK, or ERROR if unsuccessful.
ERRNO	S_objLib_OBJ_INVALID_ARGUMENT name is NULL or empty.
	S_objLib_OBJ_NOT_FOUND No message queue with <i>name</i> was found.
	S_objLib_OBJ_OPERATION_UNSUPPORTED The message queue is not named.
	S_objLib_OBJ_DESTROY_ERROR An error was detected while deleting the message queue.
	S_intLib_NOT_ISR_CALLABLE This routine must not be called from an ISR.
SEE ALSO	msgQOpen, msgQOpen(), msgQClose()

# munlock()

NAME	<pre>munlock() – unlock specified pages (POSIX)</pre>
SYNOPSIS	<pre>int munlock   (    const void * addr,    size_t len   )</pre>
DESCRIPTION	This routine unlocks specified pages from being memory resident.
RETURNS	0 ( <b>OK</b> ) always.
ERRNO	N/A
SEE ALSO	mmanPxLib

# munlockall()

NAME	<b>munlockall()</b> – unlock all pages used by a process (POSIX)

SYNOPSIS int munlockall (void)

**DESCRIPTION** This routine unlocks all pages used by a process from being memory resident.

**RETURNS** 0 (OK) always.

ERRNO N/A

SEE ALSO mmanPxLib

## **mv()**

NAME	<b>mv()</b> – mv file into other directory.
SYNOPSIS	<pre>STATUS mv (     const char * src, /* source file name or wildcard */     const char * dest /* destination name or directory */ )</pre>
DESCRIPTION	This function is similar to <b>rename()</b> but behaves somewhat more like the UNIX program "mv", it will overwrite files.
	This command moves the <i>src</i> file or directory into a file which name is passed in the <i>dest</i> argument, if <i>dest</i> is a regular file or does not exist. If <i>dest</i> name is a directory, the source object is moved into this directory as with the same name, if <i>dest</i> is <b>NULL</b> , the current directory is assumed as the destination directory. <i>src</i> may be a single file name or a path containing a wildcard pattern, in which case all files or directories matching the pattern will be moved to <i>dest</i> which must be a directory in this case.
EXAMPLES	-> mv( "/sd0/dir1","/sd0/dir2") -> mv( "/sd0/*.tmp","/sd0/junkdir") -> mv( "/sd0/FILE1.DAT","/sd0/dir2/f001.dat")
RETURNS	<b>OK</b> or error if any of the files or directories could not be moved, or if <i>src</i> is a pattern but the destination is not a directory.
ERRNO	Not Available
SEE ALSO	usrFsLib, the VxWorks programmer guides.

	nanosleep()	
NAME	<b>nanosleep()</b> – suspend the current task until the time interval elapses (POSIX)	
SYNOPSIS	<pre>int nanosleep   (     const struct timespec * rqtp, /* time to delay */     struct timespec * rmtp /* premature wakeup (NULL=no result) */   )</pre>	
DESCRIPTION	This routine suspends the current task for a specified time <i>rqtp</i> or until a signal or event notification is made.	
	The suspension may be longer than requested due to the rounding up of the request to the timer's resolution or to other scheduling activities (e.g., a higher priority task intervenes).	
	The <b>timespec</b> structure is defined as follows:	
	<pre>struct timespec {</pre>	
	remaining. If <i>rmtp</i> is <b>NULL</b> , the remaining time is not returned. The <i>rqtp</i> parameter is greater than 0 or less than or equal to 1,000,000,000.	
RETURNS	0 ( <b>OK</b> ), or -1 ( <b>ERROR</b> ) if the routine is interrupted by a signal or an asynchronous event notification, or <i>rqtp</i> is invalid.	
ERRNO	EINVAL EINTR	
SEE ALSO	timerLib, sleep(), taskDelay()	

# netHelp()

**NAME netHelp()** – print a synopsis of network routines

SYNOPSIS void netHelp (void)

**DESCRIPTION** This command prints a brief synopsis of network facilities that are typically called from the shell.

hostAdd	"hostname","inetaddr"	- add a host to remote host table; "inetaddr" must be in standard Internet address format e g "90 0 0 4"
hostShow		- print current remote host table
netDevCreat	e "devname", "hostname",	protocol
		<ul> <li>create an I/O device to access files on the specified host (protocol 0=rsh, 1=ftp)</li> </ul>
routeAdd	"destaddr","gateaddr"	- add route to route table
routeDelete	destaddr", "gateaddr"	- delete route from route table
routeShow		- print current route table
iam	"usr"[,"passwd"]	- specify the user name by which
		you will be known to remote
		hosts (and optional password)
whoami		- print the current remote ID
rlogin	"host"	<ul> <li>log in to a remote host;</li> </ul>
		"host" can be inet address or
		host name in remote host table
ifShow inetstatSho tcpstatShow udpstatShow ipstatShow icmpstatSho arptabShow mbufShow	["ifname"] w v	<ul> <li>show info about network interfaces</li> <li>show all Internet protocol sockets</li> <li>show statistics for TCP</li> <li>show statistics for UDP</li> <li>show statistics for IP</li> <li>show statistics for ICMP</li> <li>show a list of known ARP entries</li> <li>show mbuf statistics</li> </ul>
EXAMPLE: -	> hostAdd "wrs", "90.0.	0.2"
-	<pre>&gt; netDevCreate "wrs:", &gt; iam "fred"</pre>	"wrs", 0
-	<pre>&gt; copy <wrs: etc="" passwd<="" pre=""></wrs:></pre>	/* copy file from host "wrs" */
-	> rlogin "wrs"	/* rlogin to host "wrs" */
N/A		
N/A		

**SEE ALSO** usrLib, the VxWorks programmer guides.

RETURNS

ERRNO

# nfsAuthUnixGet()

NAME	<b>nfsAuthUnixGet()</b> – get the NFS UNIX authentication parameters	
SYNOPSIS	void nfsAuthUnixGet	
	char *machname, /* where to store host machine	*/
	int *pUid, /* where to store user ID	*/
	int *pGid, /* where to store group ID	*/

VxWorks Kernel API Reference, 6.6 nfsAuthUnixPrompt()

 

 int \*pNgids, /\* where to store number of group IDs \*/ int \*gids /\* where to store array of group IDs \*/ )

 DESCRIPTION
 This routine gets the previously set UNIX authentication values.

 RETURNS
 N/A

 ERRNO
 Not Available

 SEE ALSO
 nfsCommon, nfsAuthUnixPrompt(), nfsAuthUnixShow(), nfsAuthUnixSet(), nfsIdSet()

### nfsAuthUnixPrompt()

NAME nfsAuthUnixPrompt() – modify the NFS UNIX authentication parameters

SYNOPSIS void nfsAuthUnixPrompt (void)

**DESCRIPTION** This routine allows UNIX authentication parameters to be changed from the shell. The user is prompted for each parameter, which can be changed by entering the new value next to the current one.

EXAMPLE	-> nfsAuthUnixPrompt
	machine name: yuba
	user ID: 2001 128
	group ID: 100
	num of groups: 1 3
	group #1: 100 100
	group #2: 0 120
	group #3: 0 200
	value = $3 = 0x3$
RETURNS	Not Available
ERRNO	Not Available
SEE ALSO	nfsCommon, nfsAuthUnixShow(), nfsAuthUnixSet(), nfsAuthUnixGet(), nfsIdSet()

# nfsAuthUnixSet()

NAME	nfsAuthUnixSet() – set the NFS UNIX authentication parameters	
SYNOPSIS	<pre>void nfsAuthUnixSet   (    char *machname, /* host machine */    int uid, /* user ID */    int gid, /* group ID */    int ngids, /* number of group IDs */    int *aup_gids /* array of group IDs */   )</pre>	
DESCRIPTION	This routine sets UNIX authentication parameters. It is initially called by <b>usrNetInit()</b> . <b>machname</b> should be set with the name of the mounted system (i.e. the target name itself) to distinguish hosts from hosts on a NFS network.	
RETURNS	N/A	
ERRNO	Not Available	
SEE ALSO	nfsCommon, nfsAuthUnixPrompt( ), nfsAuthUnixShow( ), nfsAuthUnixGet( ), nfsIdSet( )	

# nfsAuthUnixShow()

NAME	nfsAuthUnixShow() – display the NFS UNIX authentication parameters	
SYNOPSIS	void nfsAuthUnixShow (void)	
DESCRIPTION	This routine displays the parameters set by <b>nfsAuthUnixSet( )</b> or <b>nfsAuthUnixPrompt( )</b> .	
EXAMPLE	-> nfsAuthUnixShow machine name = yuba user ID = 2001 group ID = 100 group [0] = 100 value = 1 = 0x1	
RETURNS	N/A	
ERRNO	Not Available	

SEE ALSO nfsCommon, nfsAuthUnixPrompt(), nfsAuthUnixSet(), nfsAuthUnixGet(), nfsIdSet()

### nfsChkFilePerms()

NAME	<b>nfsChkFilePerms()</b> – check the NFS file permissions with a given permission.	
SYNOPSIS	<pre>STATUS nfsChkFilePerms   (     int nfsPerms, /* permissions of the opened file */     int ruid, /* Remote uid */     int rgid, /* Remote gid */     int perm /* permission to be checked for 4:read 2:write 1:execute */    )</pre>	
DESCRIPTION	<ul> <li>This routine compares the NFS file permissions with a given permission.</li> <li>This routine is basically designed for nfsOpen() to verify the target file's permission prior to deleting it due to O_TRUNC.</li> <li>The parameter <i>perm</i> will take 4(read), 2(write), 1(execute), or combinations of them.</li> <li>OK means the file has valid permission whichever group is.</li> <li>Called only by the I/O system.</li> </ul>	
RETURNS	OK, ERROR	
ERRNO	Not Available	
SEE ALSO	nfsCommon	

# nfsDevInfoGet()

NAME nfsDevInfoGet() – read configuration information from the requested device

DESCRIPTION	This routine accesses the NFS device specified in the parameter <i>nfsDevHandle</i> and f in the structure pointed to by <i>pnfsinfo</i> . The calling function should allocate memory for <i>pnfsinfo</i> and for the two character buffers "remFileSys" and "locFileSys", that are part of <i>pnfsInfo</i> . These buffers should have a size of nfsMaxPath	
RETURNS	<b>OK</b> , if <i>pnfsInfo</i> information is valid, otherwise <b>ERROR</b>	
ERRNO	<b>S_objLib_OBJ_UNAVAILABLE</b> This operation is not supported by the available NFS versions.	
SEE ALSO	nfsCommon, nfsDevListGet( )	

# nfsDevListGet()

NAME	nfsDevListGet() – create list of all the NFS devices in the system	
SYNOPSIS	<pre>int nfsDevListGet    (     unsigned long nfsDevList[], /* NFS dev list of handles */     int listSize /* number of elements available in the list */    )</pre>	
DESCRIPTION	This routine fills the array <i>nfsDevlist</i> up to <i>listSize</i> , with handles to NFS devices currently in the system.	
RETURNS	The number of entries filled in the <i>nfsDevList</i> array.	
ERRNO	N/A	
SEE ALSO	nfsCommon, nfsDevInfoGet( )	

# nfsDevShow()

NAME	nfsDevShow() – display the mounted NFS devices	
SYNOPSIS	void nfsDevShow (void)	
DESCRIPTION	This routine displays the device names and their associated NFS file systems.	

VxWorks Kernel API Reference, 6.6 nfsDrvNumGet()

EXAMPLE	-> nfsDevShow device name	file system
	/yubal/ /wrs1/	yuba:/yuba1 wrs:/wrs1
RETURNS	N/A	
ERRNO	Not Available	

SEE ALSO nfsCommon

## nfsDrvNumGet()

NAME	nfsDrvNumGet() – Get driver number of NFS device	
SYNOPSIS	STATUS nfsDrvNumGet ( int version /* Version number of NFS */ )	
DESCRIPTION	This routine returns the NFS driver number for the version requested. If the user specifies NFS version 2, this routine will return the value stored in variable nfs2DrvNum. If the user specifies NFS version 3, this routine will return the value stored in the variable nfs3DrvNum.	
	If the NFS driver of the user-specified version is yet to be initialized, or if initialization failed, nfsDrvNumGet will return ERROR.	
RETURNS	Returns the NFS driver number or ERROR.	
ERRNO	S_objLib_OBJ_UNAVAILABLE This version does not support this operation.	
SEE ALSO	nfsCommon	

# nfsErrnoSet()

NAME nfsErrnoSet() – set NFS status

SYNOPSIS void nfsErrnoSet
( enum nfsstat status )

- **DESCRIPTION** nfsErrnoSet calls errnoSet with the given "nfs stat" or'd with the NFS status prefix.
- **RETURNS** Not Available
- ERRNO Not Available
- SEE ALSO nfsCommon

# nfsExport()

NAME	nfsExport() – specify a file system to be NFS exported
SYNOPSIS	<pre>STATUS nfsExport   (     char * directory, /* Directory to export - FS must support NFS */     int id, /* ID number for file system */     BOOL readOnly, /* TRUE if file system is exported read-only */     int options /* Reserved for future use - set to 0 */    )</pre>
DESCRIPTION	This routine makes a file system available for mounting by a client. The client should be in the local host table (see <b>hostAdd()</b> ), although this is not required.
	The <i>id</i> parameter can either be set to a specific value, or to 0. If it is set to 0, an ID number is assigned sequentially. Every time a file system is exported, it must have the same ID number, or clients currently mounting the file system will not be able to access files.
NOTE	exporting a file system requires at least 512kb of free space available on the file system for creation of configuration files.
	To display a list of exported file systems, use:
	-> nfsExportShow "localhost"
RETURNS	OK, or ERROR if the file system could not be exported.
ERRNO	Not Available
SEE ALSO	mountd, nfsLib, nfsExportShow(), nfsUnexport()

	nfsExportShow()
NAME	<b>nfsExportShow()</b> – display the exported file systems of a remote host
SYNOPSIS	<pre>STATUS nfsExportShow   (    char *hostName /* host machine to show exports for */ )</pre>
DESCRIPTION	This routine displays the file systems of a specified host and the groups that are allowed to mount them.
EXAMPLE	<pre>-&gt; nfsExportShow "wrs" /d0 staff /d1 staff eng /d2 eng /d3 value = 0 = 0x0</pre>
RETURNS	OK or ERROR.
ERRNO	<b>S_hostLib_INVALID_PARAMETER</b> <i>hostName</i> is invalid.
	S_objLib_OBJ_UNAVAILABLE This operation is not supported by the available NFS versions.
	S_nfsLib_NFSERR_NOTSUPP Remote system does not have a compatible mount version.
SEE ALSO	nfsCommon

# nfsHelp()

NAME	<b>nfsHelp( )</b> – display the NFS help men	ıu
SYNOPSIS	void nfsHelp (void)	
DESCRIPTION	This routine displays a summary of N	FS facilities typically called from the shell:
	nfsHelp netHelp nfsMount "host","filesystem"[ nfsUnmount "devname"	Print this list Print general network help list ,"devname"] Create device with file system/directory from host Remove an NFS device

```
Print current UNIX authentication
                    nfsAuthUnixShow
                                                   Prompt for UNIX authentication
                    nfsAuthUnixPrompt
                    nfsIdSet id
                                                   Set user ID for UNIX authentication
                    nfsDevShow
                                                   Print list of NFS devices
                    nfsExportShow "host"
                                                  Print a list of NFS file systems which
                                                   are exported on the specified host
                    mkdir "dirname"
                                                   Create directory
                    rm "file"
                                                   Remove file
                    EXAMPLE: -> hostAdd "wrs", "90.0.0.2"
                              -> nfsMount "wrs", "/disk0/path/mydir", "/mydir/"
                              -> cd "/mydir/"
                              -> nfsAuthUnixPrompt
                                                       /* fill in user ID, etc.
                                                                                       */
                              -> ls
                                                        /* list /disk0/path/mydir
                                                                                       */
                              -> copy < foo
                                                        /* copy foo to standard out */
                              -> ld < foo.o /* load object module foo.o */
-> nfsUnmount "/mydir/" /* remove NFS device /mydir/ */
               N/A
RETURNS
               Not Available
ERRNO
                nfsCommon
SEE ALSO
```

### nfsIdSet()

NAME	nfsIdSet() – set the ID number of the NFS UNIX authentication parameters
SYNOPSIS	void nfsIdSet ( int uid /* user ID on host machine */ )
DESCRIPTION	This routine sets only the UNIX authentication user ID number. For most NFS permission needs, only the user ID needs to be changed. Set <i>uid</i> to the user ID on the NFS server.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	nfsCommon, nfsAuthUnixPrompt( ), nfsAuthUnixShow( ), nfsAuthUnixSet( ), nfsAuthUnixGet( )

2

# nfsMntDump()

NAME	nfsMntDump() – display all NFS file systems mounted on a particular host
SYNOPSIS	STATUS nfsMntDump ( const char *hostName /* host machine */ )
DESCRIPTION	This routine displays all the NFS file systems mounted on a specified host machine.
RETURNS	OK or ERROR.
ERRNO	<pre>S_nfsLib_NFSERR_INVAL hostName is invalid. S_nfsLib_NFSERR_NOTSUPP Remote system does not have a compatible mount version. S_objLib_OBJ_UNAVAILABLE</pre>
	This routine is not supported by the included NFS versions.

### nfsMount()

nfsCommon

SEE ALSO

NAME	nfsMount() – mount an NFS file system	
SYNOPSIS	<pre>STATUS nfsMount   (    const char *host, /* name of remote host */    const char *fileSystem, /* name of remote directory to mount */    const char *localName /* local device name for remote dir */   )</pre>	
DESCRIPTION	This routine mounts a remote file system. It creates a local device <i>localName</i> for a remote file system on a specified host. The host must have already been added to the local host table with <b>hostAdd()</b> .	
RETURNS	<b>OK</b> , or <b>ERROR</b> if the driver is not installed, <i>host</i> is invalid, or memory is insufficient.	
ERRNO	S_nfsLib_NFSERR_INVAL Provided arguments are invalid.	

	S_objLib_OBJ_UNAVAILABLE This operation is not supported by the included NFS versions.
	S_nfsLib_NFSERR_NOTSUPP Remote system does not have a compatible mount version.
SEE ALSO	nfsCommon, nfsUnmount( ), hostAdd( )

# nfsMountAll()

NAME	<b>nfsMountAll()</b> – mount all file systems exported by a specified host	
SYNOPSIS	<pre>STATUS nfsMountAll   (     const char *pHostName, /* name of remote host */     const char *pClientName, /* name of a client specified in access list,     if any */     BOOL quietFlag /* FALSE = print name of each mounted file     system */     )</pre>	
DESCRIPTION	This routine mounts the file systems exported by the host <i>pHostName</i> which are accessible by <i>pClientName</i> . A <i>pClientName</i> entry of <b>NULL</b> will only mount file systems that are accessible by any client. The <b>nfsMount()</b> routine is called to mount each file system. It creates a local device for each mount that has the same name as the remote file system. If the <i>quietFlag</i> setting is <b>FALSE</b> , each file system is printed on standard output after it is mounted successfully.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if any mount fails.	
ERRNO	<ul> <li>S_nfsLib_NFSERR_INVAL Invalid arguments.</li> <li>S_objLib_OBJ_UNAVAILABLE This operation is not supported by the included NFS versions.</li> <li>S_nfcLib_NESERR_NOTSUPR</li> </ul>	
SEE ALSO	nfsCommon, nfsMount()	

### nfsStatusGet()

NAME	nfsStatusGet() – Get the statistics of the NFS server
SYNOPSIS	STATUS nfsStatusGet ( void * serverStats, /* pointer to status struct */ int version /* NFS v2 or NFS v3 */ )
DESCRIPTION	This routine returns the statistics of the NFS procedure calls made by the remote NFS clients <i>serverStats</i> pointer to a memory location where the statistics information will be copied. <i>version</i> Statistics of which NFS version is desired. This parameter takes two values only. 0x01 for NFS Version 2 and 0x02 for NFS version 3.
RETURNS	<b>OK</b> or <b>ERROR</b> if version is invalid.
ERRNO	Not Available
SEE ALSO	nfsdCommon

# nfsUnexport()

NAME	nfsUnexport() – remove a file system from the list of exported file systems
SYNOPSIS	STATUS nfsUnexport ( char * dirName /* Name of the directory to unexport */ )
DESCRIPTION	This routine removes a file system from the list of file systems exported from the target. Any client attempting to mount a file system that is not exported will receive an error (NFSERR_ACCESS).
RETURNS	OK, or ERROR if the file system could not be removed from the exports list.
ERRNO	ENOENT
SEE ALSO	mountd, nfsLib, nfsExportShow(), nfsExport()

### nfsUnmount()

NAME	nfsUnmount() – unmount an NFS device
SYNOPSIS	STATUS nfsUnmount ( const char *localName /* local of nfs device */ )
DESCRIPTION	This routine unmounts file systems that were previously mounted via NFS.
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>localName</i> is not an NFS device or cannot be mounted.
ERRNO	S_nfsLib_NFSERR_INVAL localName is invalid.
	S_nfsDrv_NOT_AN_NFS_DEVICE localName is not an NFS device.
	S_objLib_OBJ_UNAVAILABLE This operation is not supported by the included NFS versions.
SEE ALSO	nfsCommon, nfsMount()

### nfsdHashTableParamsSet()

NAME	nfsdHashTableParamsSet() – sets up the parameters for the NFS hash table
SYNOPSIS	void nfsdHashTableParamsSet ( int bucketSize, int tableLen, char * basePath )
DESCRIPTION	This function must be called prior to the call to nfsExport, and can be called repeatedly to configure different exports differently.
	<i>bucketSize</i> represents the number of bytes available in each hash bucket. The selectable sizes are 512, 1k, 2k, 4k, 8k, 16k. It is probably a good idea to make this value the same size as a disk allocation unit in general. Never smaller, but perhaps larger if you plan on accessing lots of files with very long filenames. Note, that the code always keeps one bucket in

memory, and reads and writes whole buckets at a time. So larger buckets can impact performance. A value of zero preserves the previous value. The default is 1k.

*tableLen* represents the number of buckets. As we all know, prime numbers work best for hashing, so selectable sizes for *tableLen* are all prime numbers that are close to powers of 2. Specifically: 7, 17, 31, 61, 127, 251, 509, 1021, and 2039. When the hash-function is applied to a filename or inode, it is the result of the modulus operation that then determines the bucket the value is placed into. A value of zero preserves the previous value. The default value is 509.

*basePath* allows you to move the hash file to a seperate location. Perhaps a disk volume dedicated to hash files, or a RAM drive if you do not need the inode-to-filename mappings preserved. The default is "" which indicates that the hash files will be saved to the exported volume. nfsExport("/foo",0,0) would result in /foo/nfsHashTbl.cfg and /foo/nfsHashTbl.cfg being created. With basePath set to "/**ramDrv**", the files would be placed in the /**ramDrv** volume and pre-pended with the export path. Specifically, the files created would be /ramDrv/foo\_nfsHashTbl.cfg and /ramDrv/foo\_nfsHashTbl.dat. Passing in NULL preserves the existing value. The default value is "";

*bucketSize* and *tableLen* are only used when creating a hash table for the first time. After that, their values are preserved in nfsHashTbl.cfg. However, basePath (if different from the default) must be called so the NFS server can locate the files.

RETURNS	Not Available	

ERRNO Not Available

SEE ALSO nfsHash

### nfsdInit()

NAME	<b>nfsdInit( )</b> – initialize the NFS s	erver
SYNOPSIS	<pre>STATUS nfsdInit   (     int nServers,     int nExportedFs,     int priority,     FUNCPTR authHook,     FUNCPTR mountAuthHook,     int options     )</pre>	<pre>/* the number of NFS servers to create */ /* maximum number of exported file systems */ /* the priority for the NFS servers */ /* Authentication hook */ /* authentication hook for mount daemon */ /* 3 bits used only */</pre>
DESCRIPTION	This routine initializes the NFS spawned to handle NFS request <i>authHook</i> is a pointer to an auth routine, passed to <b>mountdInit</b> (	server. <i>nServers</i> specifies the number of Tasks to be ts. <i>priority</i> is the priority that those tasks will run at. orization routine. <i>mountAuthHook</i> is a pointer to a similar ). <i>options</i> (only 3 LSBs are used for specifying the NFS

version). Currently options can take the following values. 0x01 to start the NFS V2 service

only. 0x02 to start the NFS V3 service only. 0x00 to start the NFS V2 and V3 services (default)

Normally, no authorization is performed by either mountd or nfsd. If you want to add authorization, set *authHook* to a function pointer to a routine declared as follows:

nfsstat routine
 (
 int progNum, /\* RPC program number \*/
 int versNum, /\* RPC program version number \*/
 int procNum, /\* RPC procedure number \*/
 struct sockaddr\_in clientAddr, /\* address of the client \*/
 NFSD\_ARGUMENT \* nfsdArg /\* argument of the call \*/
 )

The nfsdArg will be of type "NFSD\_ARGUMENT" if versNum is 2. The nfsdArg will be of type "NFS3D\_ARGUMENT" if versNum is 3. The user authentication hook must use the nfsdArg accordingly.

The *authHook* routine should return NFS\_OK if the request is authorized, and NFS3ERR\_ACCES if not. (NFSERR\_ACCESS is not required; any legitimate NFS error code can be returned.)

See **mountdInit()** for documentation on *mountAuthHook*. Note that *mountAuthHook* and *authHook* can point to the same routine. Simply use the *progNum*, *versNum*, and *procNum* fields to decide whether the request is an NFS request or a mountd request.

**RETURNS** OK, or ERROR if the NFS server cannot be started.

ERRNO Not Available

SEE ALSO nfsd, nfsExport(), mountdInit()

### nfsdStatusShow()

 NAME
 nfsdStatusShow() – show the status of the NFS server

 SYNOPSIS
 STATUS nfsdStatusShow

 (int options /\* unused \*/
)

 DESCRIPTION
 This routine shows statistics of procedure calls to the NFS server. This routine takes one
 parameter to specify the NFS server version whose statistics are to be shown. options takes
 one of the following 3 values. 0x01 Display statistics of NFS version 2 server only. 0x02
 Display statistics of NFS version 3 server only. 0x00 Display statistics of NFS version 2 & 3.

VxWorks Kernel API Reference, 6.6 nicRegister()

**RETURNS** OK, or **ERROR** if the information cannot be obtained.

ERRNO Not Available

SEE ALSO nfsdCommon

# nicRegister()

SEE ALSO	ns83902VxbEnd
ERRNO	N/A
RETURNS	N/A
DESCRIPTION	This routine registers the ST-NIC driver with VxBus as a child of the PLB bus type.
SYNOPSIS	void nicRegister(void)
NAME	<b>nicRegister()</b> – register with the VxBus subsystem

### npc()

NAME	<b>npc( )</b> – return the contents of the next program counter (SimSolaris)
SYNOPSIS	<pre>int npc   (    int taskId /* task ID, 0 means default task */ )</pre>
DESCRIPTION	This command extracts the contents of the next program counter from the TCB of a specified task. If <i>taskId</i> is omitted or 0, the current default task is assumed.
RETURNS	The contents of the next program counter.
ERRNO	Not Available
SEE ALSO	dbgArchLib, ti()

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NAME	<b>nseRegister()</b> – register with the VxBus subsystem
SYNOPSIS	void nseRegister(void)
DESCRIPTION	This routine registers the NatSemi driver with VxBus as a child of the PCI bus type.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	ns8381xVxbEnd

# nvRamSegDefGet()

----

\_

NAME	<b>nvKamSegDefGet()</b> – get segment allocation from BSP
SYNOPSIS	<pre>STATUS nvRamSegDefGet   (    VXB_DEVICE_ID   pInst,    HCF_DEVICE *   pHcf,    NVRAM_SEGMENT **   ppSegList,    int *</pre>
DESCRIPTION	This routine reads entries from the hcf record for the NVRam device specified by pInst, and fills them in to a table. Allocation of the table is performed within this routine, using <b>hwMemLib</b> allocation. The table consists of a nSeg field indicating the number of segments, followed by a set of structures containing the start offset, size, allocation name, and allocation unit number. Within the <b>hwconf.c</b> file, resource names are specified with the name "segAddr". Segment size names are "segSz". Driver (or OS module) names are "drvName". Driver unit number names are "drvUnit". Note that each of these four resource names must be specified for each segment, regardless of whether the segment is allocated or not. For segments not allocated, the segment should

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either be omitted, or allocated to a module named "unallocated" unit -1.

The hwconf entries for each segment must be grouped together, and each segment must have all segments fully described: name, unit, addr, size. Failure to meet these restrictions will result in segment information being corrupted or discarded.

RETURNS	<b>OK</b> , or <b>ERROR</b> if the table cannot be allocate	d

ERRNO Not Available

SEE ALSO vxbNonVolLib

### 00()

NAME	<b>o0( )</b> – return the contents of register o0 (also o1-o7) (SimSolaris)
SYNOPSIS	<pre>int o0   (   int taskId /* task ID, 0 means default task */ )</pre>
DESCRIPTION	This command extracts the contents of out register o0 from the TCB of a specified task. If <i>taskId</i> is omitted or 0, the current default task is assumed.
	Similar routines are provided for all out registers (o0 - o7): <b>o0( ) - o7( )</b> .
	The stack pointer is accessed via o6.
RETURNS	The contents of register o0 (or the requested register).
ERRNO	Not Available
SEE ALSO	dbgArchLib, VxWorks Programmer's Guide: Debugging

# objClassTypeGet()

NAME	<b>objClassTypeGet( )</b> – get an object's class type
SYNOPSIS	enum windObjClassType objClassTypeGet ( OBJ_ID objId )
DESCRIPTION	The class type of the specified object is returned.
RETURNS	class type enum, or <b>ERROR</b> if caller has insufficient access rights.
ERRNO	Possible errnos generated by this routine include: <b>S_objLib_OBJ_ID_ERROR</b> Invalid object identifier.
SEE ALSO	objLib

# objContextGet()

NAME	<pre>objContextGet( ) - return the object's context value</pre>
SYNOPSIS	<pre>STATUS objContextGet   (    OBJ_ID objId, /* object to get context from */    void ** pContext /* where to store context value */   )</pre>
DESCRIPTION	The value stored in the object's context field is returned. The context field is typically set by calling <b>objContextSet()</b> or an <b>xxxOpen()</b> routine.
RETURNS	<b>OK</b> , or <b>ERROR</b> if objId is invalid.
ERRNO	Possible errnos generated by this routine include: <b>S_objLib_OBJ_ID_ERROR</b> Invalid object identifier.
SEE ALSO	objLib

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# objContextSet()

NAME	<b>objContextSet( )</b> – set the object's context value
SYNOPSIS	<pre>STATUS objContextSet   (    OBJ_ID objId, /* object to set context on */    void * context /* context value */   )</pre>
DESCRIPTION	This routine sets the object's context field. This value is not actually used by VxWorks. Instead, the context value can be used by OS extensions to implement object permissions, for example.
RETURNS	<b>OK</b> , or <b>ERROR</b> if objId is invalid.
ERRNO	Possible errnos generated by this routine include: <b>S_objLib_OBJ_ID_ERROR</b> Invalid object identifier.
SEE ALSO	objLib

# objHandleShow()

NAME	<b>objHandleShow()</b> – show information on the object referenced by an object handle
SYNOPSIS	STATUS objHandleShow ( OBJ_HANDLE objHandle, /* object handle to get information from */ RTP_ID rtpId /* ID of RTP to which objHandle belongs */ )
DESCRIPTION	This routine displays information regarding the WIND object an object handle references. This routine is intended to be used for debugging purposes.
RETURNS	OK, or ERROR if the information could not be displayed.
ERRNO	N/A
SEE ALSO	objShow, objShowAll( )

# objHandleTblShow() - show information on an RTP's handle table

SYNOPSIS	void objHandleTblShow	
	(	
	RTP_ID rtpId,	
	int disp	
	)	

objHandleTblShow()

DESCRIPTION This routine displays the contents of the supplied *rtpId* handle table. The argument *count* indicates the number of slots in the table to display. In case *count* is zero, all the in-use slots in the table are displayed. This routine is intended to be used only for debugging purposes.

RETURNS	N/A
ERRNO	N/A

NAME

SEE ALSO objShow, rtpShow(), rtpDetailShow()

### objNameGet()

NAME	<b>objNameGet( )</b> – get an object's name	
SYNOPSIS	<pre>STATUS objNameGet   (    OBJ_ID objId, /* pointer to object to get name */    char * nameBuf, /* pointer to name string buffer */    int bufSize /* size, in bytes, of name buffer */   )</pre>	
DESCRIPTION	The specified object's name string is copied into <i>nameBuf</i> .	
RETURNS	OK, or ERROR if object name cannot be retrieved.	
ERRNO	Possible errnos generated by this routine include:	
	<b>S_objLib_OBJ_NAME_TRUNCATED</b> Supplied name buffer is too small. Truncated name has been returned.	
	<b>S_objLib_OBJ_NOT_NAMED</b> Object has not been labeled with a name.	

VxWorks Kernel API Reference, 6.6 objNameLenGet()

S\_objLib\_OBJ\_ID\_ERROR Invalid object identifier.

objLib

SEE ALSO

### objNameLenGet()

NAME	<b>objNameLenGet( )</b> – get an object's name length
SYNOPSIS	int objNameLenGet ( OBJ_ID objId )
DESCRIPTION	The specified object's name length (without the terminating $\0$ character) is returned.
RETURNS	name length or -1 if the object name cannot be retrieved
ERRNO	Possible errnos generated by this routine include: <b>S_objLib_OBJ_ID_ERROR</b> Invalid object identifier. <b>S_objLib_OBJ_OPERATION_UNSUPPORTED</b> Object class does not support the name get operation. <b>S_objLib_OBJ_NOT_NAMED</b> Object has not been labeled with a name.
SEE ALSO	objLib

# objNameToId()

NAME	<pre>objNameToId() - find object</pre>	with matching name string and type
SYNOPSIS	OBJ_ID objNameToId ( enum windObjClassType const char *	classType, name
	)	

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**DESCRIPTION** The object name space is searched for an object with a matching *name* and *classType*. There may exist more than one object of the same type with identical names. In such cases, this routine will return the id of the first object found.

This routine is provided if the INCLUDE\_OBJ\_INFO component is present in the configuration.

Values for the windObjClassType enumerated type are tabulated below:

	Value	Object Class
	0	Invalid
	1	Wind Semaphore
	2	POSIX Semaphore
	3	Wind Message Queue
	4	POSIX Message Queue
	5	Real Time Process
	6	Task
	7	Watchdog Timer
	8	File Descriptor
	9	Page Pool
	10	Page Manager
	11	Group
	12	Virtual Memory Context
	13	Event Trigger
	14	Memory Partition
	15	I2O
	16	device management system
	17	Set
	18	ISR object
	19	POSIX Timer
	20	Shared data region
RETURNS	NULL i	f no match occurs, otherwise <b>OBJ_ID</b> of matching object.
ERRNO	Possibl	le errnos generated by this routine include:
	S_objLi	ib_OBJ_ILLEGAL_CLASS_TYPE
	11	ie specifica object class type is nivana.

SEE ALSO objLib

# objOwnerGet()

NAME	<b>objOwnerGet( )</b> – return the object's owner
SYNOPSIS	OBJ_ID obj0wnerGet ( OBJ_ID objId /* object to get owner from */ )
DESCRIPTION	The ID of the object that owns the specified object is returned.
RETURNS	owner object ID, or <b>NULL</b> if invalid object id, the task does not have access rights, or object ownership is excluded from the system
ERRNO	Possible errnos generated by this routine include:
	S_objLib_OBJ_ID_ERROR Invalid object identifier.
SEE ALSO	objLib

# objOwnerSet()

NAME	<b>objOwnerSet( )</b> – change the object's owner
SYNOPSIS	STATUS objOwnerSet ( OBJ_ID objId, /* object to set owner */ OBJ_ID ownerId /* owner object ID */ )
DESCRIPTION	Set the owner of an object. This routine is used to change the default object ownership hierarchy. The calling task must have access rights to both the object <i>objId</i> whose owner is being changed, and the owner object <i>ownerId</i> which must be a real time process.
	The owner object <i>ownerId</i> must be a real time process.
	If INCLUDE_OBJ_OWNERSHIP is excluded this routine simply returns OK.
RETURNS	ERROR if object ID or owner ID is invalid.
ERRNO	Possible errnos generated by this routine include:

S\_objLib\_OBJ\_ID\_ERROR Invalid object or owner object identifier.

**S\_objLib\_OBJ\_INVALID\_OWNER** Invalid object ownership relationship.

SEE ALSO objLib

### objShow()

NAME	<b>objShow()</b> – show information on an object	
SYNOPSIS	<pre>STATUS objShow   (    OBJ_ID objId, /* object to show information on */    int showType /* show type */   )</pre>	
DESCRIPTION	Call class attached show routine for an object.	
RETURNS	OK, or ERROR if information could not be shown.	
ERRNO	Possible errnos generated by this routine include:	
	S_objLib_OBJ_ID_ERROR Invalid object identifier.	
	S_objLib_OBJ_NO_METHOD Show routine for this class of object not installed.	
SEE ALSO	objLib	

# objShowAll()

NAME	<b>objShowAll()</b> – show all information on an object
SYNOPSIS	<pre>STATUS objShowAll   (    OBJ_ID objId, /* object to show information on */    int showType /* show type */   )</pre>

VxWorks Kernel API Reference, 6.6 objShowAll()

DESCRIPTION This routine displays all information about an object. The generic object information is handled directly by this routine, while the class specific information is handled by the show routine registered for the class. The routine **objShow()** only displays the class-specific information. The *showType* parameter is passed transparently to the class-specific show routine. Typically, a *showType* of 1 is used to enable a detailed information display. If *objld* is not given or is null, objShowAll() displays all information on the system. The following example shows information on a task with TID = 0x188d78: EXAMPLE [vxKernel] -> objShowAll 0x188d78 Generic Object Information \_\_\_\_\_ Type : Task Name : /pubTestTask Attr : 0xc1 (WIND\_OBJ\_NAME\_DYNAMIC WIND\_OBJ\_PUBLIC WIND\_OBJ\_NAMED) refCnt: 2 Ctx : 0x0 (type = 1) Owner Information \_\_\_\_\_ ID : 0x000badc4 Type : Real Time Process Name : (null) Object Handles opened on this object: Object Handle RTP -----0x1e001d 0xdc722c helloworld.vxe Owned Objects \_\_\_\_\_ Object Id Object Type Object Name \_\_\_\_\_ \_ \_\_\_\_ 0x001ff2d8 Binary Semaphore (null) 0x001ff288 Binary Semaphore (userTblSem) 0x001ff238 Mutex Semaphore wdMutexSem 0x001ff1b8 |-Watchdog (null) 0x001ff168 Message Queue (null) Task Specific Information \_\_\_\_\_ ENTRY TID PRI STATUS PC SP ERRNO DELAY NAME /pubTestTask 0x45345 c7cce4 100 PEND 419563 dc9fc8 0 0 value =  $0 = 0 \times 0$ 

	The ownership hierarchy is shown by an indentation of the object type. In the above example, the watchdog object is directly owned by a semaphore object (wdMutexSem), which in turn is owned by the task object (testTask). Hence the watchdog object is indirectly owned by the task object (testTask).
	For an object that is not named, if a symbol table entry whose value matches the object id exists, the symbol name, enclosed in brackets, will be displayed under the object name column. When an object is not named and an exact symbol table match does not exist, "(null)" will be displayed under the object name column.
WARNING	Deleting <i>objId</i> while <b>objShowAll()</b> is gathering information, can lead to unexpected results.
RETURNS	OK, or ERROR if the information could not be displayed.
ERRNO	N/A
SEE ALSO	objShow

# open()

NAME	open() – open a file
SYNOPSIS	<pre>int open   (    const char *name, /* name of the file to open */    int flags, /* access control flag */    int mode /* mode of file to create (UNIX chmod style) */   )</pre>
DESCRIPTION	This routine opens a file for reading, writing, or updating, and returns a file descriptor for that file. The arguments to <b>open()</b> are the filename <i>name</i> and the type of access set in <i>flags</i> and a UNIX chmod-style file mode <i>mode</i> .
	The parameter <i>flags</i> is set to one or a combination of the following access settings by bitwise OR operation for the duration of time the file is open. The following list is just a generic description of supported settings. Their availability and effect with or without combination among them change from device to device. Check the specific device manual for further details.
	O_RDONLY Open for reading only.
	O_WRONLY Open for writing only.

#### O\_RDWR

Open for reading and writing.

#### O\_CREAT

Create a file if not existing.

#### O\_EXCL

Error on open if file exists and O\_CREAT is also set.

#### O\_SYNC

Write on the file descriptor complete as defined by synchronized I/O file integrity completion.

#### O\_DSYNC

Write on the file descriptor complete as defined by synchronized I/O data integrity completion.

#### O\_RSYNC

Read on the file descriptor complete at the same sync level as **O\_DSYNC** and **O\_SYNC** flags.

#### O\_APPEND

If set, the file offset is set to the end of the file prior to each write. So writes are guaranteed at the end. It has no effect on devices other than the regular file system.

#### O\_NONBLOCK

Non-blocking I/O if being set.

#### O\_NOCTTY

Do not assign a ctty on this open, which does not cause the terminal device to become the controlling terminal for the process. Effective only on a terminal device.

#### O\_TRUNC

Open with truncation. If the file exists and is a regular file, and the file is successfully opened, its length is truncated to 0. It has no effect on devices other than the regular file system.

In general, **open()** can only open pre-existing devices and files. However, files can also be created with **open()** by setting **O\_CREAT** and perhaps some other like **O\_RDWR** which depends on the file system implementation. In this case, the file is created with a UNIX chmod-style file mode, as indicated with the parameter *mode*. For example:

fd = open ("/usr/myFile", O\_CREAT | O\_RDWR, 0644);

Files, on dosFs volumes, can be opened with the **O\_SYNC** flag indicating that each write should be immediately written to the backing media. This synchronizes the FAT and the directory entries.

# **NOTE** For more information about situations when there are no file descriptors available, see the reference entry for **iosInit()**.

	Also note that not all device drivers honor the flags or mode values when opening a file. Most simple devices simply ignore them and return an open file descriptor for both reading and writing. Read the device driver manual for information on this.
RETURNS	A file descriptor number, or <b>ERROR</b> if a file name is not specified, the device does not exist, no file descriptors are available, or the driver returns <b>ERROR</b> .
ERRNO	ELOOP Circular symbolic link, too many links.
	EMFILE Maximum number of files already open.
	S_iosLib_DEVICE_NOT_FOUND (ENODEV) No valid device name found in path.
	others Other errors reported by device drivers.
SEE ALSO	ioLib, creat()

# opendir()

NAME	<b>opendir()</b> – open a directory for searching (POSIX)
SYNOPSIS	DIR *opendir ( const char* dirName /* name of directory to open */ )
DESCRIPTION	This routine opens the directory named by <i>dirName</i> and allocates a directory descriptor (DIR) for it. A pointer to the DIR structure is returned. The return of a NULL pointer indicates an error.
	After the directory is opened, <b>readdir()</b> is used to extract individual directory entries. Finally, <b>closedir()</b> is used to close the directory.
WARNING	For remote file systems mounted over <b>netDrv</b> , <b>opendir()</b> fails, because the <b>netDrv</b> implementation strategy does not provide a way to distinguish directories from plain files. To permit use of <b>opendir()</b> on remote files, use NFS rather than <b>netDrv</b> .
RETURNS	A pointer to a directory descriptor, or NULL if there is an error.
ERRNO	N/A.

VxWorks Kernel API Reference, 6.6 operator\_delete()

SEE ALSO dirLib, closedir(), readdir(), rewinddir(), ls()

# operator\_delete()

NAME	<b>operator_delete( )</b> – default run-time support for memory deallocation (C++)
SYNOPSIS	extern void operator delete ( void *pMem /* pointer to dynamically-allocated object */ )
DESCRIPTION	This function provides the default implementation of operator delete. It returns the memory, previously allocated by operator new, to the VxWorks system memory partition.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	cplusLib

### operator\_new()

NAME	<b>operator_new()</b> – default run-time support for operator new (C++)
SYNOPSIS	<pre>extern void * operator new   (    size_t n /* size of object to allocate */    ) throw (std::bad_alloc)</pre>
DESCRIPTION	This function provides the default implementation of operator new. It allocates memory from the system memory partition for the requested object. The value, when evaluated, is a pointer of the type <b>pointer-to</b> - <i>T</i> where <i>T</i> is the type of the new object.
	If allocation fails a new-handler, if one is defined, is called. If the new-handler returns, presumably after attempting to recover from the memory allocation failure, allocation is retried. If there is no new-handler an exception of type "bad_alloc" is thrown.
THROWS	std::bad_alloc if allocation failed.
RETURNS	Pointer to new object.

ERRNO Not Available

SEE ALSO cplusLib

### operator\_new()

**operator\_new()** – default run-time support for operator new (nothrow) (C++) NAME SYNOPSIS extern void \* operator new ( size\_t n, /\* size of object to allocate \*/
const nothrow\_t & /\* supply argument of "nothrow" here \*/ ) throw () DESCRIPTION This function provides the default implementation of operator new (nothrow). It allocates memory from the system memory partition for the requested object. The value, when evaluated, is a pointer of the type **pointer-to-***T* where *T* is the type of the new object. If allocation fails, a new-handler, if one is defined, is called. If the new-handler returns, presumably after attempting to recover from the memory allocation failure, allocation is retried. If the new\_handler throws a bad\_alloc exception, the exception is caught and 0 is returned. If allocation fails and there is no new\_handler 0 is returned. RETURNS Pointer to new object or 0 if allocation fails. ERRNO Not Available SEE ALSO cplusLib

### operator\_new()

NAME	<b>operator_new()</b> – run-time support for operator new with placement (C++)
SYNOPSIS	extern void * operator new ( size_t n, /* size of object to allocate (unused) */ void * pMem /* pointer to allocated memory */ )
DESCRIPTION	This function provides the default implementation of the global new operator, with support for the placement syntax. New-with-placement is used to initialize objects for which

VxWorks Kernel API Reference, 6.6 oprintf()

memory has already been allocated.  $p{\it Mem}$  points to the previously allocated memory. memory.

**RETURNS** *pMem* 

ERRNO Not Available

SEE ALSO cplusLib

### oprintf()

NAME	<b>oprintf()</b> – write a formatted string to an output function
SYNOPSIS	<pre>int oprintf ( FUNCPTR prtFunc, /* pointer to output function */ int prtArg, /* argument for output function */ const char * fmt, /* format string to write */ /* optional arguments to format string */ )</pre>
DESCRIPTION	This routine prints a formatted string via the function specified by <i>prtFunc</i> . The function will receive as parameters a pointer to a buffer, an integer indicating the length of the buffer, and the argument <i>prtArg</i> . If <b>NULL</b> is specified as the output function, the output will be sent to stdout.
	The function and syntax of oprintf are otherwise identical to <b>printf()</b> .
SMP CONSIDERATIO	<b>NS</b> This API is spinlock and intCpuLock restricted.
RETURNS	The number of characters output, not including the NULL terminator.
ERRNO	Not Available
SEE ALSO	fioBaseLib, printf()

### partLibCreate()

partLibCreate() – partition a device NAME SYNOPSIS STATUS partLibCreate ( int fd, int nPart, int size1, int size2, int size3 ) This routine partitions a device. DESCRIPTION RETURNS OK on success, ERROR otherwise Not Available ERRNO SEE ALSO partLib

### passFsDevInit()

NAME	<pre>passFsDevInit() - associate a device with passFs file system functions</pre>
SYNOPSIS	<pre>void * passFsDevInit   (    char * devName /* device name */ )</pre>
DESCRIPTION	This routine associates the name $devName$ with the file system and installs it in the I/O System's device table.
RETURNS	A pointer to the volume descriptor on success, else NULL.
ERRNO	S_iosLib_DUPLICATE_DEVICE_NAME (EINVAL) Device name already in use.
SEE ALSO	passFsLib

# passFsInit()

NAME	<b>passFsInit()</b> – prepare to use the passFs library
SYNOPSIS	STATUS passFsInit ( int passfs, /* number of pass-through file systems */ BOOL cacheEnable /* enable passfs cache ? */ )
DESCRIPTION	This routine initializes the passFs library. It must be called only once, before any other routines in the library. First argument specifies the number of passFs devices that may be open at once, second argument is a boolean that specifies if cache must be enabled or not. This routine installs <b>passFsLib</b> as a driver in the I/O system driver table, allocates and sets up the necessary memory structures, and initializes semaphores. Usually this routine is called from the root task, <b>usrRoot()</b> , in <b>prjConfig()</b> . This initialization is enabled when the configuration component <b>INCLUDE_PASSFS</b> is defined.
NOTE	Maximum number of pass-through file systems is 1.
RETURNS	OK on success, else ERROR.
ERRNO	S_iosLib_DRIVER_GLUT (ENOMEM) No memory available for data structures.
SEE ALSO	passFsLib

# pathconf()

NAME	<b>pathconf()</b> – determine the current value of a configurable limit
SYNOPSIS	<pre>long pathconf  (     const char *path, /* path of the file */     int name /* Value to query */ )</pre>
DESCRIPTION	The <b>fpathconf()</b> and <b>pathconf()</b> functions provide a method for the application

CRIPTION The **fpathconf()** and **pathconf()** functions provide a method for the application to determine the current value of a configurable limit or option (variable) that is associated with a file or directory.

RETURNS	The current value is returned if valid with the query. Otherwise, <b>ERROR</b> , -1 returned and errno may be set to indicate the error. There are many reasons to return <b>ERROR</b> . If the variable corresponding to name has no limit for the path or file descriptor, both <b>pathconf()</b> and <b>fpathconf()</b> return -1 without changing errno.
ERRNO	
SEE ALSO	fsPxLib, fpathconf()

# pause()

NAME	<b>pause()</b> – suspend the task until delivery of a signal (POSIX)
SYNOPSIS	int pause (void)
DESCRIPTION	This routine suspends the task until delivery of a signal.
NOTE	Since the <b>pause()</b> function suspends thread execution indefinitely, there is no successful completion return value.
RETURNS	-1, always.
ERRNO	EINTR
SEE ALSO	sigLib

# pc()

NAME	<b>pc( )</b> – return the contents of the program counter
SYNOPSIS	int pc ( int task /* task ID */ )
DESCRIPTION	This command extracts the contents of the program counter for a specified task from the task's TCB. If <i>task</i> is omitted or 0, the current task is used.
RETURNS	the contents of the program counter.

VxWorks Kernel API Reference, 6.6 pcConDevBind()

 ERRNO
 N/A

 SEE ALSO
 usrLib, ti(), the VxWorks programmer guides.

### pcConDevBind()

pcConDevBind() - bind keyboard or VGA device with console NAME SYNOPSIS TY\_DEV \* pcConDevBind ( int arg, FUNCPTR pFunc, void \* pArg ) This routine is called by the keyboard and VGA drivers to associate themselves with a PC DESCRIPTION console instance. The keyboard driver should normally pass its unit number for arg, and NULL for the remaining two parameters. The VGA driver should pass a pointer to its buffer processing routine and the argument to this routine for these two parameters. RETURNS TY\_DEV pointer associated with console, or NULL. Not Available ERRNO vxbPcConsole SEE ALSO

### pcConDevCreate( )

NAME	a device for the on-board ports		
SYNOPSIS	STATUS pcConDevCreate (	(* none to use for this device	<b>+</b> /
	char ^ name,	/ name to use for this device	^/
	FAST int channel,	/* virtual console number	*/
	int rdBufSize,	/* read buffer size, in bytes	*/
	int wrtBufSize )	/* write buffer size in bytes	*/
DECODUDITION	This routing greates a device	an ana of the naConcole north. Each	nont to b

**DESCRIPTION** This routine creates a device on one of the pcConsole ports. Each port to be used should have only one device associated with it, by calling this routine.

ERRNO Not Available

SEE ALSO vxbPcConsole

### pcConDrv()

NAME	<b>pcConDrv( )</b> – initialize the console driver
SYNOPSIS	STATUS pcConDrv (void)
DESCRIPTION	This routine initializes the console driver, sets up interrupt vectors, and performs hardware initialization of the keybord and display.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the driver cannot be installed.
ERRNO	Not Available
SEE ALSO	vxbPcConsole

### pentiumBtc()

NAME	<b>pentiumBtc()</b> – execute atomic compare-and-exchange instruction to clear a bit		
SYNOPSIS	STATUS pentiumBtc (pFlag) char * pFlag; /* flag address */		
DESCRIPTION	This routine compares a byte specified by the first parameter with <b>TRUE</b> . If it is <b>TRUE</b> , it changes it to 0 and returns <b>OK</b> . If it is not <b>TRUE</b> , it returns <b>ERROR</b> . LOCK and CMPXCHGB are used to get the atomic memory access.		
RETURNS	OK or ERROR if the specified flag is not TRUE		
ERRNO	Not Available		
SEE ALSO	pentiumALib		

### pentiumBts()

**NAME pentiumBts()** – execute atomic compare-and-exchange instruction to set a bit

SYNOPSIS	STATUS pentiumBts	(pFlag)				
	char * pFlag;		/*	flag	address	*/

- **DESCRIPTION** This routine compares a byte specified by the first parameter with 0. If it is 0, it changes it to **TRUE** and returns **OK**. If it is not 0, it returns **ERROR**. LOCK and CMPXCHGB are used to get the atomic memory access.
- **RETURNS** OK or ERROR if the specified flag is not zero.
- ERRNO Not Available

SEE ALSO pentiumALib

### pentiumCr4Get()

NAME	pentiumCr4Get() – get contents of CR4 register
SYNOPSIS	int pentiumCr4Get (void)
DESCRIPTION	This routine gets the contents of the CR4 register. This routine is kept for the backward compatibility, and <b>vxCr4Get()</b> should be used instead. The CR4 is introduced in the Pentium processor, thus this routine just returns in the pre Pentium generation processors.
RETURNS	Contents of CR4 register.
ERRNO	Not Available
SEE ALSO	pentiumALib

### pentiumCr4Set( )

NAME pentiumCr4Set() – sets specified value to the CR4 register

SYNOPSIS void pentiumCr4Set (cr4)

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	<pre>int cr4; /* value to write CR4 register */</pre>
DESCRIPTION	This routine sets a specified value to the CR4 register. This routine is kept for the backward compatibility, and <b>vxCr4Set()</b> should be used instead. The CR4 is introduced in the Pentium processor, thus this routine just returns in the pre Pentium generation processors.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	pentiumALib

# pentiumMcaEnable()

NAME	pentiumMcaEnable() – enable/disable the MCA (Machine Check Architecture)
SYNOPSIS	<pre>void pentiumMcaEnable   (   BOOL enable /* TRUE to enable, FALSE to disable the MCA */ )</pre>
DESCRIPTION	This routine enables/disables 1) the Machine Check Architecture and its Error Reporting register banks 2) the Machine Check Exception by toggling the MCE bit in the CR4. This routine works on either P5, P6 or P7 family.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	pentiumLib

# pentiumMcaShow()

NAME	<b>pentiumMcaShow()</b> – show MCA (Machine Check Architecture) registers	
SYNOPSIS	void pentiumMcaShow (void)	
DESCRIPTION	This routine shows Machine-Check global control registers and Error-Reporting register banks. Number of the Error-Reporting register banks is kept in a variable mcaBanks.	

VxWorks Kernel API Reference, 6.6 pentiumMsrGet()

MCi\_ADDR and MCi\_MISC registers in the Error-Reporting register bank are showed if MCi\_STATUS indicates that these registers are valid.

RETURNS N/A

ERRNO Not Available

SEE ALSO pentiumShow

### pentiumMsrGet()

NAME	pentiumMsrGet() – get the contents of the specified MSR (Model Specific Register)		
SYNOPSIS	<pre>void pentiumMsrGet (addr, pData)     int addr;</pre>		
DESCRIPTION	This routine gets the contents of the specified MSR. The first parameter is an address of the MSR. The second parameter is a pointer of 64Bit variable.		
RETURNS	N/A		
ERRNO	Not Available		
SEE ALSO	pentiumALib		

### pentiumMsrInit( )

NAME	<b>pentiumMsrInit( )</b> – initialize all the MSRs (Model Specific Register)
SYNOPSIS	STATUS pentiumMsrInit (void)
DESCRIPTION	This routine initializes all the MSRs in the processor. This routine works on either P5, P6 or P7 family processors.
RETURNS	OK, or ERROR if RDMSR/WRMSR instructions are not supported.
ERRNO	Not Available

SEE ALSO pentiumLib

### pentiumMsrSet()

pentiumMsrSet() - set a value to the specified MSR (Model Specific Registers) NAME void pentiumMsrSet (addr, pData) SYNOPSIS /\* MSR address \*/ int addr; long long int \* pData; /\* MSR data \*/ This routine sets a value to a specified MSR. The first parameter is an address of the MSR. DESCRIPTION The second parameter is a pointer of 64Bit variable. RETURNS N/A Not Available ERRNO SEE ALSO pentiumALib

### pentiumMsrShow()

- NAME pentiumMsrShow() show all the MSR (Model Specific Register)
- SYNOPSIS void pentiumMsrShow (void)
- **DESCRIPTION** This routine shows all the MSRs in the Pentium and Pentium[234].
- RETURNS N/A
- ERRNO Not Available
- SEE ALSO pentiumShow

VxWorks Kernel API Reference, 6.6 pentiumMtrrDisable()

### pentiumMtrrDisable()

NAMEpentiumMtrrDisable() – disable MTRR (Memory Type Range Register)SYNOPSISvoid pentiumMtrrDisable (void)DESCRIPTIONThis routine disables the MTRR that provide a mechanism for associating the memory types<br/>with physical address ranges in system memory.RETURNSN/AERRNONot AvailableSEE ALSOpentiumLib

### pentiumMtrrEnable()

NAME	pentiumMtrrEnable() – enable MTRR (Memory Type Range Register)
SYNOPSIS	void pentiumMtrrEnable (void)
DESCRIPTION	This routine enables the MTRR that provide a mechanism for associating the memory types with physical address ranges in system memory.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	pentiumLib

### pentiumMtrrGet()

NAME pentiumMtrrGet() - get MTRRs to a specified MTRR table
SYNOPSIS STATUS pentiumMtrrGet
(
MTRR \* pMtrr /\* MTRR table \*/

)
**DESCRIPTION** This routine gets MTRRs to a specified MTRR table with RDMSR instruction. The read MTRRs are CAP register, DEFTYPE register, fixed range MTRRs, and variable range MTRRs.

- **RETURNS** OK, or ERROR if MTRR is being accessed.
- ERRNO Not Available
- SEE ALSO pentiumLib

#### pentiumMtrrSet( )

**NAME pentiumMtrrSet()** – set MTRRs from specified MTRR table with WRMSR instruction.

SYNOPSIS STATUS pentiumMtrrSet ( MTRR \* pMtrr /\* MTRR table \*/

**DESCRIPTION** This routine sets MTRRs from specified MTRR table with WRMSR instruction. The written MTRRs are DEFTYPE register, fixed range MTRRs, and variable range MTRRs.

**RETURNS** OK, or ERROR if MTRR is enabled or being accessed.

ERRNO Not Available

SEE ALSO pentiumLib

#### pentiumP5PmcGet( )

0, and the second parameter is for the Counter 1.

NAME	<b>pentiumP5PmcGet( )</b> – get the contents of	f P5 PMC0 and PMC1
SYNOPSIS	<pre>void pentiumP5PmcGet (pPmc0, pPmc1)     long long int * pPmc0;     long long int * pPmc1;</pre>	/* Performance Monitoring Counter 0 */ /* Performance Monitoring Counter 1 */
DESCRIPTION	This routine gets the contents of both PM PMC1. The first parameter is a pointer of	20 (Performance Monitoring Counter 0) and 64Bit variable to store the content of the Counter

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VxWorks Kernel API Reference, 6.6 pentiumP5PmcGet0()

RETURNS N/A

ERRNO Not Available

SEE ALSO pentiumALib

# pentiumP5PmcGet0()

NAME	pentiumP5PmcGet0() – get the contents of P5 PMC0	
SYNOPSIS	<pre>void pentiumP5PmcGet0 (pPmc0)     long long int * pPmc0;</pre>	/* Performance Monitoring Counter 0 */
DESCRIPTION	This routine gets the contents of PMC0 (Performance Monitoring Counter 0). The parameter is a pointer of 64Bit variable to store the content of the Counter.	
RETURNS	N/A	
ERRNO	Not Available	
SEE ALSO	pentiumALib	

# pentiumP5PmcGet1()

NAME	<b>pentiumP5PmcGet1()</b> – get the contents of P5 PMC1	
SYNOPSIS	<pre>void pentiumP5PmcGet1 (pPmc1)     long long int * pPmc1;</pre>	/* Performance Monitoring Counter 1 */
DESCRIPTION	This routine gets a content of PMC1 (Performation pointer of 64Bit variable to store the content	mance Monitoring Counter 1). Parameter is a t of the Counter.
RETURNS	N/A	
ERRNO	Not Available	
SEE ALSO	pentiumALib	

NAME pentiumP5PmcReset() – reset both PMC0 and PMC1

SYNOPSIS void pentiumP5PmcReset (void)

**DESCRIPTION** This routine resets both PMC0 (Performance Monitoring Counter 0) and PMC1.

RETURNS N/A

ERRNO Not Available

SEE ALSO pentiumALib

## pentiumP5PmcReset0()

NAME	pentiumP5PmcReset0() – reset PMC0
SYNOPSIS	void pentiumP5PmcReset0 (void)
DESCRIPTION	This routine resets PMC0 (Performance Monitoring Counter 0).
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	pentiumALib

## pentiumP5PmcReset1()

NAME	<pre>pentiumP5PmcReset1() - reset PMC1</pre>	
SYNOPSIS	void pentiumP5PmcReset1 (void)	
DESCRIPTION	This routine resets PMC1 (Performance Monitoring Counter 1).	
RETURNS	N/A	

VxWorks Kernel API Reference, 6.6 pentiumP5PmcStart0()

ERRNO Not Available

SEE ALSO pentiumALib

## pentiumP5PmcStart0()

NAME pentiumP5PmcStart0() – start PMC0

- SYNOPSIS STATUS pentiumP5PmcStart0 (pmc0Cesr) int pmc0Cesr; /\* PMC0 control and event select \*/
- **DESCRIPTION** This routine starts PMC0 (Performance Monitoring Counter 0) by writing specified PMC0 events to Performance Event Select Registers. The only parameter is the content of Performance Event Select Register.
- **RETURNS OK** or **ERROR** if PMC0 is already started.

ERRNO Not Available

SEE ALSO pentiumALib

#### pentiumP5PmcStart1()

NAME	pentiumP5PmcStart1() – start PMC1
SYNOPSIS	STATUS pentiumP5PmcStart1 (pmc1Cesr) int pmc1Cesr; /* PMC1 control and event select */
DESCRIPTION	This routine starts PMC1 (Performance Monitoring Counter 0) by writing specified PMC1 events to Performance Event Select Registers. The only parameter is the content of Performance Event Select Register.
RETURNS	<b>OK</b> or <b>ERROR</b> if PMC1 is already started.
ERRNO	Not Available
SEE ALSO	pentiumALib

	pentiumP5PmcStop()
NAME	pentiumP5PmcStop() – stop both P5 PMC0 and PMC1
SYNOPSIS	void pentiumP5PmcStop (void)
DESCRIPTION	This routine stops both PMC0 (Performance Monitoring Counter 0) and PMC1 by clearing two Performance Event Select Registers.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	pentiumALib

### pentiumP5PmcStop0()

NAME	pentiumP5PmcStop0() – stop P5 PMC0
SYNOPSIS	void pentiumP5PmcStop0 (void)
DESCRIPTION	This routine stops only PMC0 (Performance Monitoring Counter 0) by clearing the PMC0 bits of Control and Event Select Register.
RETURNS	N/A
ERRNO	Not Available

SEE ALSO pentiumALib

### pentiumP5PmcStop1()

NAME pentiumP5PmcStop1() – stop P5 PMC1

SYNOPSIS void pentiumP5PmcStop1 (void)

**DESCRIPTION** This routine stops only PMC1 (Performance Monitoring Counter 1) by clearing the PMC1 bits of Control and Event Select Register.

VxWorks Kernel API Reference, 6.6 pentiumP6PmcGet( )

N/A RETURNS

Not Available ERRNO

pentiumALib SEE ALSO

# pentiumP6PmcGet()

NAME	<b>pentiumP6PmcGet( )</b> – get the contents of	FPMC0 and PMC1
SYNOPSIS	<pre>void pentiumP6PmcGet (pPmc0, pPmc1)     long long int * pPmc0;     long long int * pPmc1;</pre>	/* Performance Monitoring Counter 0 */ /* Performance Monitoring Counter 1 */
DESCRIPTION	This routine gets the contents of both PM0 PMC1. The first parameter is a pointer of 0, and the second parameter is for the Cou	C0 (Performance Monitoring Counter 0) and 64Bit variable to store the content of the Counter inter 1.
RETURNS	N/A	
ERRNO	Not Available	
SEE ALSO	pentiumALib	

# pentiumP6PmcGet0()

NAME	<b>pentiumP6PmcGet0()</b> – get the contents of PMC0	
SYNOPSIS	<pre>void pentiumP6PmcGet0 (pPmc0)     long long int * pPmc0; /* Performance Monitoring Counter 0 */</pre>	
DESCRIPTION	This routine gets the contents of PMC0 (Performance Monitoring Counter 0). The parameter is a pointer of 64Bit variable to store the content of the Counter.	
RETURNS	N/A	
ERRNO	Not Available	
SEE ALSO	pentiumALib	

	pentiumP6PmcGet1()	
NAME	<pre>pentiumP6PmcGet1() - get the contents of PMC1</pre>	
SYNOPSIS	<pre>void pentiumP6PmcGet1 (pPmc1)     long long int * pPmc1; /* Performance Monitoring Counter 1 */</pre>	
DESCRIPTION	This routine gets a content of PMC1 (Performance Monitoring Counter 1). Parameter is a pointer of 64Bit variable to store the content of the Counter.	
RETURNS	N/A	
ERRNO	Not Available	
SEE ALSO	pentiumALib	

## pentiumP6PmcReset()

NAME	pentiumP6PmcReset() – reset both PMC0 and PMC1
SYNOPSIS	void pentiumP6PmcReset (void)
DESCRIPTION	This routine resets both PMC0 (Performance Monitoring Counter 0) and PMC1.
RETURNS	N/A
ERRNO	Not Available

SEE ALSO pentiumALib

# pentiumP6PmcReset0()

- NAME pentiumP6PmcReset0() reset PMC0
- SYNOPSIS void pentiumP6PmcReset0 (void)
- **DESCRIPTION** This routine resets PMC0 (Performance Monitoring Counter 0).

VxWorks Kernel API Reference, 6.6 pentiumP6PmcReset1()

RETURNS N/A

ERRNO Not Available

SEE ALSO pentiumALib

# pentiumP6PmcReset1()

NAME	pentiumP6PmcReset1() – reset PMC1
SYNOPSIS	void pentiumP6PmcReset1 (void)
DESCRIPTION	This routine resets PMC1 (Performance Monitoring Counter 1).
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	pentiumALib

# pentiumP6PmcStart()

NAME	pentiumP6PmcStart() – start both PMC0 and PMC1
SYNOPSIS	STATUS pentiumP6PmcStart (pmcEvtSel0, pmcEvtSel1) int pmcEvtSel0; /* Performance Event Select Register 0 */ int pmcEvtSel1; /* Performance Event Select Register 1 */
DESCRIPTION	This routine starts both PMC0 (Performance Monitoring Counter 0) and PMC1 by writing specified events to Performance Event Select Registers. The first parameter is a content of Performance Event Select Register 0, and the second parameter is for the Performance Event Select Register 1.
RETURNS	<b>OK</b> or <b>ERROR</b> if PMC is already started.
ERRNO	Not Available
SEE ALSO	pentiumALib

	pentiumP6PmcStop()
NAME	pentiumP6PmcStop() – stop both PMC0 and PMC1
SYNOPSIS	void pentiumP6PmcStop (void)
DESCRIPTION	This routine stops both PMC0 (Performance Monitoring Counter 0) and PMC1 by clearing two Performance Event Select Registers.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	pentiumALib

### pentiumP6PmcStop1()

NAME	<pre>pentiumP6PmcStop1() - stop PMC1</pre>
SYNOPSIS	void pentiumP6PmcStop1 (void)
DESCRIPTION	This routine stops only PMC1 (Performance Monitoring Counter 1) by clearing the Performance Event Select Register 1. Note, clearing the Performance Event Select Register 0 stops both counters, PMC0 and PMC1.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	pentiumALib

# pentiumPmcGet()

 
 NAME
 pentiumPmcGet() - get the contents of PMC0 and PMC1

 SYNOPSIS
 void pentiumPmcGet (pPmc0, pPmc1) long long int \* pPmc0; /\* Performance Monitoring Counter 0 \*/ long long int \* pPmc1; /\* Performance Monitoring Counter 1 \*/
 VxWorks Kernel API Reference, 6.6 pentiumPmcGet0()

DESCRIPTION	none
-------------	------

RETURNS N/A

ERRNO Not Available

SEE ALSO pentiumLib

# pentiumPmcGet0()

NAME	<b>pentiumPmcGet0()</b> – get the contents of PMC0	
SYNOPSIS	<pre>void pentiumPmcGet0 (pPmc0)     long long int * pPmc0;</pre>	/* Performance Monitoring Counter 0 */
DESCRIPTION	none	
RETURNS	N/A	
ERRNO	Not Available	
SEE ALSO	pentiumLib	

# pentiumPmcGet1()

NAME	<pre>pentiumPmcGet1() - get the contents of PMC1</pre>	
SYNOPSIS	<pre>void pentiumPmcGet1 (pPmc1)     long long int * pPmc1;</pre>	/* Performance Monitoring Counter 1 */
DESCRIPTION	none	
RETURNS	N/A	
ERRNO	Not Available	
SEE ALSO	pentiumLib	

### pentiumPmcReset()

NAME pentiumPmcReset() – reset both PMC0 and PMC1

- SYNOPSIS void pentiumPmcReset (void)
- DESCRIPTION none
- RETURNS N/A
- ERRNO Not Available
- SEE ALSO pentiumLib

### pentiumPmcReset0()

- NAME pentiumPmcReset0() reset PMC0
- SYNOPSIS void pentiumPmcReset0 (void)
- DESCRIPTION none
- RETURNS N/A
- ERRNO Not Available
- SEE ALSO pentiumLib

### pentiumPmcReset1()

- NAME pentiumPmcReset1() reset PMC1
- SYNOPSIS void pentiumPmcReset1 (void)
- DESCRIPTION none
- RETURNS N/A

VxWorks Kernel API Reference, 6.6 pentiumPmcShow()

ERRNO Not Available

SEE ALSO pentiumLib

# pentiumPmcShow()

NAME	<pre>pentiumPmcShow() - show PMCs (Performance Monitoring Counters)</pre>
SYNOPSIS	<pre>void pentiumPmcShow   (    BOOL zap /* 1: reset PMC0 and PMC1 */   )</pre>
DESCRIPTION	This routine shows Performance Monitoring Counter 0 and 1. Monitored events are selected by Performance Event Select Registers in in pentiumPmcStart (). These counters are cleared to 0 if the parameter "zap" is <b>TRUE</b> .
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	pentiumShow

# pentiumPmcStart()

NAME	pentiumPmcStart() – start both PMC0 and PMC1
SYNOPSIS	STATUS pentiumPmcStart (pmcEvtSel0, pmcEvtSel1) int pmcEvtSel0; /* Performance Event Select Register 0 */ int pmcEvtSel1; /* Performance Event Select Register 1 */
DESCRIPTION	none
RETURNS	<b>OK</b> or <b>ERROR</b> if PMC is already started.
ERRNO	Not Available
SEE ALSO	pentiumLib

NAME pentiumPmcStart0() – start PMC0

 
 SYNOPSIS
 STATUS pentiumPmcStart0 (pmcEvtSel0) int pmcEvtSel0;
 /\* PMC0 control and event select \*/

**DESCRIPTION** none

**RETURNS OK** or **ERROR** if PMC is already started.

ERRNO Not Available

SEE ALSO pentiumLib

#### pentiumPmcStart1()

NAME	<pre>pentiumPmcStart1() - start PMC1</pre>
SYNOPSIS	STATUS pentiumPmcStart1 (pmcEvtSel1) int pmcEvtSel1; /* PMC1 control and event select */
DESCRIPTION	none
RETURNS	<b>OK</b> or <b>ERROR</b> if PMC1 is already started.
ERRNO	Not Available
SEE ALSO	pentiumLib

### pentiumPmcStop()

**NAME pentiumPmcStop()** – stop both PMC0 and PMC1

SYNOPSIS void pentiumPmcStop (void)

DESCRIPTION none

VxWorks Kernel API Reference, 6.6 pentiumPmcStop0()

RETURNS N/A

ERRNO Not Available

SEE ALSO pentiumLib

## pentiumPmcStop0()

NAME	pentiumPmcStop0() – stop PMC(
SYNOPSIS	void pentiumPmcStop0 (void)
DESCRIPTION	none
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	pentiumLib

# pentiumPmcStop1()

NAMEpentiumPmcStop1() - stop PMC1SYNOPSISvoid pentiumPmcStop1 (void)DESCRIPTIONnoneRETURNSN/AERRNONot AvailableSEE ALSOpentiumLib

NAME	<b>pentiumSerialize()</b> – execute a serializing instruction CPUID
SYNOPSIS	void pentiumSerialize (void)
DESCRIPTION	This routine executes a serializing instruction CPUID. Serialization means that all modifications to flags, registers, and memory by previous instructions are completed before the next instruction is fetched and executed and all buffered writes have drained to memory.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	pentiumALib

### pentiumTlbFlush()

NAME pentiumTlbFlush() – flush TLBs (Translation Lookaside Buffers)

- SYNOPSIS void pentiumTlbFlush (void)
- **DESCRIPTION** This routine flushes TLBs by loading the CR3 register. All of the TLBs are automatically invalidated any time the CR3 register is loaded. The page global enable (PGE) flag in register CR4 and the global flag in a page-directory or page-table entry can be used to frequently used pages from being automatically invalidated in the TLBs on a load of CR3 register. The only way to deterministically invalidate global page entries is to clear the PGE flag and then invalidate the TLBs.
- RETURNS N/A
- ERRNO Not Available
- SEE ALSO pentiumALib

VxWorks Kernel API Reference, 6.6 pentiumTscGet32()

#### pentiumTscGet32()

NAMEpentiumTscGet32() – get the lower half of the 64Bit TSC (Timestamp Counter)SYNOPSISUINT32 pentiumTscGet32 (void)DESCRIPTIONThis routine gets a lower half of the 64Bit TSC by RDTSC instruction. RDTSC instruction<br/>saves the lower 32Bit in EAX register, so this routine simply returns after executing RDTSC<br/>instruction.RETURNSLower half of the 64Bit TSC (Timestamp Counter)ERRNONot AvailableSEE ALSOpentiumALib

## pentiumTscGet64()

NAME	pentiumTscGet64() – get 64Bit TSC (Timestam	np Counter)
SYNOPSIS	<pre>void pentiumTscGet64 (pTsc)     long long int * pTsc; /*</pre>	Timestamp Counter */
DESCRIPTION	This routine gets 64Bit TSC by RDTSC instruct to store the content of the Counter.	ion. Parameter is a pointer of 64Bit variable
RETURNS	N/A	
ERRNO	Not Available	
SEE ALSO	pentiumALib	

#### pentiumTscReset( )

NAME	pentiumTscReset() – reset the TSC (Timestamp Counter
------	--

SYNOPSIS void pentiumTscReset (void)

**DESCRIPTION** This routine resets the TSC by writing zero to the TSC with WRMSR instruction.

RETURNS N/A

ERRNO Not Available

SEE ALSO pentiumALib

#### period()

```
NAME
                 period() – spawn a task to call a function periodically
SYNOPSIS
                 int period
                     (
                              secs, /* period in seconds
                     int
                                                                        */
                     FUNCPTR func, /* function to call repeatedly */
                     int arg1, /* first of eight args to pass to func */
                     int
                              arg2,
                     int
                              arg3,
                     int
                              arg4,
                     int
                              arg5,
                     int
                              arg6,
                     int
                              arg7,
                     int
                              arg8
                     )
DESCRIPTION
                 This command spawns a task that repeatedly calls a specified function, with up to eight of
                 its arguments, delaying the specified number of seconds between calls.
                 For example, to have i() display task information every 5 seconds, just type:
                     -> period 5, i
                 The task is spawned using the sp() routine. See the description of sp() for details about
NOTE
                 priority, options, stack size, and task ID.
                 A task ID, or ERROR if the task cannot be spawned.
RETURNS
ERRNO
                 sp() errnos.
                 usrLib, periodRun(), sp(), the VxWorks programmer guides.
SEE ALSO
```

VxWorks Kernel API Reference, 6.6 periodRun()

### periodRun()

periodRun() - call a function periodically NAME SYNOPSIS void periodRun ( secs, /\* no. of seconds to delay between calls \*/ int FUNCPTR func, /\* function to call repeatedly \*/ int arg1, /\* first of eight args to pass to func \*/ arg2, int int arg3, arg4, int int arg5, int arg6, int arg7, int arg8 ) DESCRIPTION This command repeatedly calls a specified function, with up to eight of its arguments, delaying the specified number of seconds between calls. Normally, this routine is called only by **period()**, which spawns it as a task. N/A RETURNS N/A ERRNO usrLib, period(), the VxWorks programmer guides. SEE ALSO

#### philDemo()

NAME phi	ilDemo() – entry point for VxWorks/SMP Dijkstra's dining philosophers demo
SYNOPSIS int	philDemo (int arg)
DESCRIPTION This Spe outj	is routine is the entry point for the VxWorks/SMP Dijkstra's dining philosophers demo. ecifying a non-0 value for <i>arg</i> disables the usage of ANSI escape sequences in the console put.
RETURNS OK	always.
ERRNO N/A	A
SEE ALSO phi	il de la constant de

#### pipeDevCreate() pipeDevCreate() - create a pipe device NAME SYNOPSIS STATUS pipeDevCreate ( const char\* name, /\* name of pipe to be created \*/ int nMessages, /\* max. number of messages in pipe \*/ int nBytes /\* size of each message \*/ ) This routine creates a pipe device. It cannot be called from an interrupt service routine. It DESCRIPTION allocates memory for the necessary structures and initializes the device. The pipe device will have a maximum of *nMessages* messages of up to *nBytes* each in the pipe at once. When the pipe is full, a task attempting to write to the pipe will be suspended until a message has been read. Messages are lost if written to a full pipe at interrupt level. OK, or ERROR if the call fails. RETURNS ENXIO ERRNO driver not initialized S\_intLib\_NOT\_ISR\_CALLABLE cannot be called from an ISR **EINVAL** invalid arguments

SEE ALSO pipeDrv

# pipeDevDelete()

NAME	<pre>pipeDevDelete( ) - delete a pipe device</pre>
SYNOPSIS	STATUS pipeDevDelete ( const char * name, /* name of pipe to be deleted */ BOOL force /* if TRUE, force pipe deletion */ )
DESCRIPTION	This routine deletes a pipe device of a given name. The name must match that passed to <b>pipeDevCreate()</b> else <b>ERROR</b> will be returned. This routine frees memory for the necessary structures and deletes the device. It cannot be called from an interrupt service routine.

	A pipe device cannot be deleted until its number of open requests has been reduced to zero by an equal number of close requests and there are no tasks pending in its select list. If the optional force flag is asserted, the above restrictions are ignored, resulting in forced deletion of any select list and freeing of pipe resources.
CAVEAT	Forced pipe deletion can have catastrophic results. Use only as a last resort.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the call fails.
ERRNO	S_intLib_NOT_ISR_CALLABLE cannot be called from an ISR
	ENXIO driver not initialized
	EMFILE pipe still has open files
	EBUSY pipe is selected by at least one pending task
	EINVAL invalid arguments
	ENODEV no device found
SEE ALSO	pipeDrv

# pipeDrv()

NAME	<b>pipeDrv( )</b> – initialize the pipe driver
SYNOPSIS	STATUS pipeDrv (void)
DESCRIPTION	This routine initializes and installs the driver. It must be called before any pipes are created. It is called automatically during initialization when VxWorks is configured with the INCLUDE_PIPES component.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the driver installation fails.
ERRNO	S_iosLib_DRIVER_GLUT (ENOMEM) No memory available for data structures.
SEE ALSO	pipeDrv

# pmFreeSpace()

NAME	<b>pmFreeSpace( )</b> – returns the amount of free space left in the PM arena
SYNOPSIS	<pre>int pmFreeSpace   (    PM_ARENA_DEF arena /* arena definition function */ )</pre>
DESCRIPTION	This function returns the amount of useable free space remaining in the PM arena. Clients of <b>pmLib</b> may request any amount up to this value.
RETURNS	the amount of free space in the arena (in bytes), or ERROR if pmLib has not been initialized
ERRNO	Not Available
SEE ALSO	pmLib

# pmInvalidate()

NAME	<b>pmInvalidate( )</b> – invalidates the entire PM arena
SYNOPSIS	STATUS pmInvalidate ( PM_ARENA_DEF arena /* arena definition function */ )
DESCRIPTION	Warning: THIS ROUTINE WILL RENDER THE ENTIRE PM ARENA INVALID!
	This function should be used with utmost care. It will invalidate the entire PM arena, effectively wiping out all regions and making their contents inaccessible.
	It should only be used when there is a need to wipe the entire PM arena, typically during development.
RETURNS	OK or ERROR
ERRNO	Not Available
SEE ALSO	pmLib

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# pmRegionAddr()

NAME	<b>pmRegionAddr()</b> – returns the address of a persistent heap region
SYNOPSIS	<pre>void *pmRegionAddr   (    PM_ARENA_DEF arena, /* arena definition function */    int region /* region number number */ )</pre>
DESCRIPTION	This function returns a pointer to the virtual address of the start of the data area of a persistent heap region.
RETURNS	a pointer to the region's data area or NULL if the region or arena is invalid
ERRNO	Not Available
SEE ALSO	pmLib

# pmRegionClose()

NAME	<b>pmRegionClose( )</b> – closes a region making it inaccessible to clients
SYNOPSIS	<pre>STATUS pmRegionClose   (    PM_ARENA_DEF arena, /* arena definition function */    int region /* region identifier */   )</pre>
DESCRIPTION	This function makes the given region inaccessible to all clients of <b>pmLib</b> . It does not guarantee to return the memory to the arena's free space, since it may be non-contiguous with the remaining free space, but it will attempt to coalesce it if at all possible.
RETURNS	OK if the region was closed safely, or ERROR if not
ERRNO	Not Available
SEE ALSO	pmLib

#### pmRegionCreate()

pmRegionCreate() - creates a persistent heap region NAME SYNOPSIS int pmRegionCreate PM\_ARENA\_DEF arena, /\* arena definition function \*/ const char \* key, /\* short name for region \*/ unsigned int size, /\* the requested size \*/ int mode /\* initial protection \*/ ) DESCRIPTION This function creates a new region in the PM arena, of the given size, with the given key. The key must be unique -- if a region already exists with the same name, it is considered an error, and this function will fail by returning ERROR. If the requested size is not a multiple of the page size, it will be rounded up to the next multiple of the page size. a positive integer identifying the region, or ERROR if it could not be created, or an existing RETURNS region has the same name Not Available ERRNO SEE ALSO pmLib

### pmRegionOpen()

NAME	<b>pmRegionOpen()</b> – opens an existing persistent heap region
SYNOPSIS	<pre>int pmRegionOpen   (    PM_ARENA_DEF arena, /* arena definition function */    const char * key /* short name for region */   )</pre>
DESCRIPTION	This function opens an existing region in the PM arena. It looks for a region with a name matching the supplied key. If one is found, it will return its region identifier. If no such region is found, it returns <b>ERROR</b> .
RETURNS	a positive integer identifying the region, or ERROR if it could not be located
ERRNO	Not Available

VxWorks Kernel API Reference, 6.6 pmRegionProtect()

SEE ALSO pmLib

# pmRegionProtect( )

NAME	<pre>pmRegionProtect() – makes a PM region read-only</pre>
SYNOPSIS	<pre>STATUS pmRegionProtect   (    PM_ARENA_DEF arena, /* arena definition function */    int region, /* the region identifier */    int mode /* PM_PROT_XXX value */   )</pre>
DESCRIPTION	This function only alters the protection state of the region if the containing arena is set up for RDONLY mode, i.e. it prefers to be immutable most of the time. If the arena was initialised in RDWR mode, then it will remain writeable always, and any attempt to set it (or one of its regions) into RDONLY mode is an error.
RETURNS	OK or ERROR if the region or arena is invalid
ERRNO	Not Available
SEE ALSO	pmLib

# pmRegionSize()

NAME	<b>pmRegionSize( )</b> – return the size of a persistent heap region
SYNOPSIS	<pre>int pmRegionSize   (     PM_ARENA_DEF arena, /* arena definition function */     int region /* persistent heap region number */   )</pre>
DESCRIPTION	This function returns the size of a region within the persistent heap.
RETURNS	the size of the region's data area or ERROR if the region or arena is invalid
ERRNO	Not Available

SEE ALSO pmLib

# pmShow()

NAME	<b>pmShow()</b> – shows the created persistent heap segments
SYNOPSIS	int pmShow ( PM_ARENA_DEF arena /* arena definition function */ )
DESCRIPTION	This function displays the allocated persistent heaps and their headers
RETURNS	OK normally, or ERROR if the PM library is not initialised
ERRNO	Not Available
SEE ALSO	pmLib

# pmValidate()

NAME	<b>pmValidate( )</b> – validates a PM arena
SYNOPSIS	STATUS pmValidate ( PM_ARENA_DEF arena /* arena definition function */ )
DESCRIPTION	This function tests the validity or otherwise of a PM arena.
RETURNS	OK if the arena is valid, or ERROR if it is corrupt or does not appear to be a PM arena
ERRNO	Not Available
SEE ALSO	pmLib

poolBlockAdd()	
1	

NAME	<b>poolBlockAdd()</b> – add an item block to the pool	
SYNOPSIS	ULONG poolBlockAdd ( POOL_ID poolId, /* ID of pool to delete */ void * pBlock, /* base address of block to add */ ULONG size /* size of block to add */ )	
DESCRIPTION	This routine adds an item block to the pool using memory provided by the user. The memory provided must be sufficient for at least one properly aligned item.	
RETURNS	number of items added, or 0 in case of error	
ERRNO	<ul> <li>S_poolLib_INVALID_POOL_ID not a valid pool ID.</li> <li>S_poolLib_INVALID_BLK_ADDR pBlock parameter is NULL.</li> <li>S_poolLib_BLOCK_TOO_SMALL size insufficient for at least one item</li> </ul>	
	size insumeent for at least one nem.	

SEE ALSO poolLib, poolCreate()

# poolCreate()

NAME	<b>poolCreate()</b> – create a pool	
SYNOPSIS	POOL_ID poolCreate	
	Const char * pName, ULONG itmSize, ULONG alignment, ULONG initCnt, ULONG incrCnt, PART_ID partId, ULONG options )	<pre>/* optional name to assign to pool */ /* size in bytes of a pool item (must be &gt; 0) */ /* alignment of a pool item */ /* (must be power of 2, or 0) */ /* initial number of items to put in pool */ /* min no of items to add to pool dynamically */ /* (if 0, no pool expansion is done) */ /* memory partition ID */ /* initial options for pool */</pre>

**DESCRIPTION** This routine creates a pool by allocating an initial block of memory which is guarenteed to contain at least *initCnt* items. The pool will hold items of the specified size and alignment only. The alignment defaults to the architecture specific allocation alignment size, and it must be a power of two value. As items are allocated from the pool, the initial block may be emptied. When a block is emptied and more items are requested, another block of memory is dynamically allocated which is guarenteed to contain *incrCnt* items. If *incrCnt* is zero, no automatic pool expansion is done.

The partition ID parameter can be used to request all item blocks being allocated from a specific memory partition. If this parameter is **NULL**, the item blocks are allocated from the system memory partition.

**POOL OPTIONS** The options parameter can be used to set the following properties of the pool. Options cannot be changed after the pool has been created. The following options are supported:

Option	Description
POOL_THREAD_SAFE	Pool operations are protected with mutex semaphore
POOL_CHECK_ITEM	Items returned to the pool are verified to be valid

**RETURNS** ID of pool or NULL if any zero count or size or insufficient memory.

#### ERRNO S\_poolLib\_ARG\_NOT\_VALID one or more invalid input arguments.

SEE ALSO poolLib, poolDelete()

#### poolDelete( )

NAME **poolDelete()** – delete a pool SYNOPSIS STATUS poolDelete ( POOL\_ID poolId, /\* ID of pool to delete \*/ BOOL force /\* force deletion if there are items in use \*/ ) DESCRIPTION This routine deletes a specified pool and all item blocks allocated for it. Memory provided by the user using **poolBlockAdd()** are not freed. If the pool is still in use (i.e. not all items have been returned to the pool) deletion can be forced with the *force* parameter set to TRUE. RETURNS OK or ERROR if bad pool ID or pool in use.

VxWorks Kernel API Reference, 6.6 poolFreeCount()

ERRNO	<b>S_poolLib_INVALID_POOL_ID</b> not a valid pool ID.
	<b>S_poolLib_POOL_IN_USE</b> can't delete a pool still in use.
SEE ALSO	<pre>poolLib, poolCreate( )</pre>

# poolFreeCount()

NAME	<pre>poolFreeCount() - return number of free items in pool</pre>
SYNOPSIS	ULONG poolFreeCount ( POOL_ID poolId /* ID of pool */ )
DESCRIPTION	This routine returns the number of free items in the specified pool.
RETURNS	number of items, or zero if invalid pool ID.
ERRNO	<b>S_poolLib_INVALID_POOL_ID</b> not a valid pool ID.
SEE ALSO	poolLib, poolTotalCount()

# poolIncrementGet( )

NAME	<pre>poolIncrementGet() - get the increment value used to grow the pool</pre>
SYNOPSIS	ULONG poolIncrementGet ( POOL_ID poolId /* ID of pool */ )
DESCRIPTION	This routine can be used to get the increment value used to grow the pool. The increment specifies how many new items are added to the pool when there are no free items left in the pool.
RETURNS	increment value, or zero if invalid pool ID.

ERRNO	S_poolLib_INVALID_POOL_ID
	not a valid pool ID.

SEE ALSO poolLib, poolIncrementSet()

#### poolIncrementSet( )

poolIncrementSet() - set the increment value used to grow the pool NAME SYNOPSIS STATUS poolIncrementSet ( POOL\_ID poolId, /\* ID of pool \*/ ULONG incrCnt /\* new increment value \*/ )

This routine can be used to set the increment value used to grow the pool. The increment DESCRIPTION specifies how many new items are added to the pool when there are no free items left in the pool.

Setting the increment to zero disables automatic growth of the pool.

- RETURNS OK, or ERROR if poolId is invalid
- S\_poolLib\_INVALID\_POOL\_ID ERRNO not a valid pool ID.
- poolLib, poolIncrementGet() SEE ALSO

### poolItemGet( )

NAME	<b>poolItemGet()</b> – get next free item from pool and return a pointer to it
SYNOPSIS	<pre>void * poolItemGet   (    POOL_ID poolId /* ID of pool from which to get item */ )</pre>
DESCRIPTION	This routine gets the next free item from the specified pool and returns a pointer to it. If the current block of items is empty, the pool increment count is non-zero, and the routine is called from task context then a new block is allocated of the given incremental size and an item from the new block is returned.

VxWorks Kernel API Reference, 6.6 poolItemReturn()

	In the kernel, this routine can be called from interrupt context if the pool was created without the <b>POOL_THREAD_SAFE</b> option. When called from ISR, the pool will not automatically grow and the routine fails if there are no free items in the pool.
RETURNS	pointer to item, or NULL in case of error.
ERRNO	<b>S_poolLib_INVALID_POOL_ID</b> not a valid pool ID.
	<b>S_poolLib_STATIC_POOL_EMPTY</b> no more items available in static pool.
	<b>S_poolLib_INT_CTX_POOL_EMPTY</b> no more items in pool while called from ISR.
SEE ALSO	poolLib, poolItemReturn()

# poolItemReturn()

NAME	<b>poolItemReturn()</b> – return an item to the pool		
SYNOPSIS	<pre>STATUS poolItemReturn (     POOL_ID poolId, /* ID of pool to which to return item */     void * pItem /* pointer to item to return */ )</pre>		
DESCRIPTION	This routine returns the specified item to the specified pool. To enable address verification on the item, the pool should be created with the <b>POOL_CHECK_ITEM</b> option. The verification can be an expensive operation, therefore the <b>POOL_CHECK_ITEM</b> option should be used when error detection is more important than deterministic behaviour of this routine.		
	In the kernel, this routine can be called from an ISR if the pool was created without the <b>POOL_THREAD_SAFE</b> option.		
RETURNS	OK, or ERROR in case of failure.		
ERRNO	S_poolLib_INVALID_POOL_ID not a valid pool ID.		
	S_poolLib_NOT_POOL_ITEM NULL pointer or item does not belong to pool.		

#### S\_poolLib\_UNUSED\_ITEM

item is already in pool free list.

SEE ALSO poolLib, poolItemGet()

#### poolShow()

NAME poolShow() – display pool information

SYNOPSIS void poolShow

```
(
POOL_ID poolId, /* ID of pool from which to get item */
ULONG level /* display info level */
)
```

**DESCRIPTION** This show routine displays information about a pool. If level is 1, it also displays statistics about memory usage efficiency by the pool. Some count values and statistics typically change dynamically, so the displayed values represent a snapshot of the pool status at the time of querying.

If the pool ID passed to this routine is **NULL**, a summary of all pools managed by **poolLib** is displayed (up to 128 pools). The following is an example for a summary info:

#### EXAMPLE

NAME	POOL ID	SIZE	TOTAL	FREE
fdEntries	0x02439ef0	80	450	44
sets	0x02439d00	84	72	7
set_nodes	0x02439a60	12	288	31
mmuPgTables	0x02438f60	4096	1647	3
memEdrPool	0x02338d20	32	294913	26973

The following is an example for a detailed info for a specific pool, with info level 1:

#### EXAMPLE

-> poolShow 0x02438f60, 1 Pool : mmuPgTables Item Size : 4096 Alignment : 0x1000 Increment : 8 Total items : 1647 Free items : 3 Options : THREAD\_SAFE Blocks : 2 VxWorks Kernel API Reference, 6.6 poolTotalCount()

Overhead	:	204	bytes	(0%)		
BLOCK	ADDI	ર	ITEMS		FREE	
0x024e	ea000	)	8	3		3
0x0243	3a000	)	175	5		0

If the pool ID passed to this routine is **NULL**, a summary of all pools managed by **poolLib** is displayed (up to 128 pools).

RETURNS N/A

ERRNO none

SEE ALSO poolShow, poolLib

## poolTotalCount()

NAME	<pre>poolTotalCount() - return total number of items in pool</pre>
SYNOPSIS	ULONG poolTotalCount ( POOL_ID poolId /* ID of pool */ )
DESCRIPTION	This routine returns the total number of items in the specified pool.
RETURNS	number of items, or zero if invalid pool ID.
ERRNO	<b>S_poolLib_INVALID_POOL_ID</b> not a valid pool ID.
SEE ALSO	poolLib, poolFreeCount()

## poolUnusedBlocksFree( )

NAME	<pre>poolUnusedBlocksFree( ) - free blocks that have all items unused</pre>

SYNOPSIS STATUS poolUnusedBlocksFree

	( POOL_ID poolId /* ID of pool to free blocks */ )
DESCRIPTION	This routine allows reducing the memory used by a pool by freeing item blocks that have all items returned to the pool. Execution time of this routine is not deterministic as it depends on the number of free items and the number of blocks in the pool. In case of multi-thread safe pools ( <b>POOL_THREAD_SAFE</b> ), this routine also locks the pool for that time.
	Blocks that were added using <b>poolBlockAdd()</b> are not freed by this routine, even if all items have been returned; only blocks that were automatically allocated during creation or auto-growth from the pool's memory partition are freed.
RETURNS	OK, or ERROR in case of failure
ERRNO	S_poolLib_INVALID_POOL_ID not a valid pool ID.
SEE ALSO	<pre>poolLib, poolBlockAdd( ), poolCreate( )</pre>

# powf()

NAME	<b>powf()</b> – compute the value of a number raised to a specified power (ANSI)
SYNOPSIS	<pre>float powf   (    float x, /* operand */    float y /* exponent */   )</pre>
DESCRIPTION	This routine returns the value of $x$ to the power of $y$ in single precision.
RETURNS	The single-precision value of <i>x</i> to the power of <i>y</i> .
ERRNO	Not Available
SEE ALSO	mathALib

# primesCompute()

NAME	<b>primesCompute( )</b> – entry point for the VxWorks SMP prime number computation demo
SYNOPSIS	<pre>STATUS primesCompute   (    unsigned int maxPrimeNum,    unsigned int numTasks</pre>
DESCRIPTION	This routine is the entry point for the VxWorks SMP prime number computation demo.
	This function will create <i>numTasks</i> computational tasks to compute prime numbers from 2 to <i>maxPrimeNum</i> . Specifying a <i>numTasks</i> of 0 selects "graph" mode. Graph mode will repeatedly compute prime numbers from 2 to <i>maxPrimeNum</i> using 1 to <i>numTasks</i> computational tasks. The compute times are plotted on an ASCII graph on standard output ( <b>STD_OUT</b> ). The x-axis represents the number of tasks used to compute prime numbers, and the y-axis represents the elapsed computation time.
	See the module description for more information.
RETURNS	<b>ERROR</b> if failed to allocate memory for the prime number candidate array or failed to spawn computational tasks. Otherwise <b>OK</b> is returned.
ERRNO	S_memLib_NOT_ENOUGH_MEMORY Out of memory for creation of computational tasks or prime number candidate array
SEE ALSO	primesDemo

# printErr()

NAME	<b>printErr()</b> – write a formatted string to the standard error stream
SYNOPSIS	<pre>int printErr (     const char * fmt, /* format string to write */     /* optional arguments to format */ )</pre>
DESCRIPTION	This routine writes a formatted string to standard error. Its function and syntax are

otherwise identical to **printf()**.

SMP CONSIDERATIONS		
	This API is spinlock and intCpuLock restricted.	
RETURNS	The number of characters output, or <b>ERROR</b> if there is an error during output.	
ERRNO	Not Available	
SEE ALSO	fioBaseLib, printf()	

# printErrno()

NAME	<b>printErrno()</b> – print the definition of a specified error status value
SYNOPSIS	<pre>void printErrno   (    int errNo /* status code whose name is to be printed */ )</pre>
DESCRIPTION	This command displays the error-status string, corresponding to a specified error-status value. It is only useful if the error-status symbol table has been built and included in the system. If <i>errNo</i> is zero, then the current task status is used by calling <b>errnoGet()</b> . This facility is described in <b>errnoLib</b> .
RETURNS	N/A
ERRNO	N/A
SEE ALSO	usrLib, errnoLib, errnoGet(), the VxWorks programmer guides.

# printLogo()

NAME	printLogo() – print the VxWorks logo
SYNOPSIS	void printLogo (void)
DESCRIPTION	This command displays the VxWorks banner seen at boot time. It also displays the VxWorks version number and kernel version number.
RETURNS	N/A

VxWorks Kernel API Reference, 6.6 printf( )

ERRNO N/A
SEE ALSO usrLib, the VxWorks programmer guides.

### printf()

**NAME** printf() – write a formatted string to the standard output stream (ANSI)

SYNOPSIS

```
int printf
 (
    const char * fmt, /* format string to write */
    ... /* optional arguments to format string */
)
```

**DESCRIPTION** This routine writes output to standard output under control of the string *fmt*. The string *fmt* contains ordinary characters, which are written unchanged, plus conversion specifications, which cause the arguments that follow *fmt* to be converted and printed as part of the formatted string.

The number of arguments for the format is arbitrary, but they must correspond to the conversion specifications in *fmt*. If there are insufficient arguments, the behavior is undefined. If the format is exhausted while arguments remain, the excess arguments are evaluated but otherwise ignored. The routine returns when the end of the format string is encountered.

The format is a multibyte character sequence, beginning and ending in its initial shift state. The format is composed of zero or more directives: ordinary multibyte characters (not %) that are copied unchanged to the output stream; and conversion specification, each of which results in fetching zero or more subsequent arguments. Each conversion specification is introduced by the % character. After the %, the following appear in sequence:

- Zero or more flags (in any order) that modify the meaning of the conversion specification.
- An optional minimum field width. If the converted value has fewer characters than the field width, it will be padded with spaces (by default) on the left (or right, if the left adjustment flag, described later, has been given) to the field width. The field width takes the form of an asterisk (\*) (described later) or a decimal integer.
- An optional precision that gives the minimum number of digits to appear for the d, i,
   o, u, x, and X conversions, the number of digits to appear after the decimal-point character for e, E, and f conversions, the maximum number of significant digits for the g and G conversions, or the maximum number of characters to be written from a string in the s conversion. The precision takes the form of a period (.) followed either by an asterisk (\*) (described later) or by an optional decimal integer; if only the period is
specified, the precision is taken as zero. If a precision appears with any other conversion specifier, the behavior is undefined.

- An optional **h** specifying that a following **d**, **i**, **o**, **u**, **x**, and **X** conversion specifier applies to a **short int** or **unsigned short int** argument (the argument will have been promoted according to the integral promotions, and its value converted to **short int** or **unsigned short int** before printing); an optional **h** specifying that a following **n** conversion specifier applies to a pointer to a **short int** argument. An optional **l** (ell) specifying that a following **d**, **i**, **o**, **u**, **x**, and **X** conversion specifier applies to a **long int** or **unsigned long int** argument; or an optional **l** specifying that a following **n** conversion specifier applies to a pointer to a **long int** argument. An optional **ll** (ell-ell) specifying that a following **d**, **i**, **o**, **u**, **x**, and **X** conversion specifier applies to a **long long int** or `unsigned **long long int**' argument; or an optional **l** specifying that a following **n** conversion specifier applies to a pointer to a **long int** argument. An optional **ll** (ell-ell) specifying that a following **d**, **i**, **o**, **u**, **x**, and **X** conversion specifier applies to a **long long int** or `unsigned long long int' argument; or an optional **ll** specifying that a following **n** conversion specifier applies to a pointer to a **long long int** argument. If a **h**, **l** or **ll** appears with any other conversion specifier, the behavior is undefined.
- WARNING: ANSI C also specifies an optional L in some of the same contexts as l above, corresponding to a long double argument. However, the current release of the VxWorks libraries does not support long double data; using the optional L gives unpredictable results.
- A character that specifies the type of conversion to be applied.

As noted above, a field width, or precision, or both, can be indicated by an asterisk (\*). In this case, an **int** argument supplies the field width or precision. The arguments specifying field width, or precision, or both, should appear (in that order) before the argument (if any) to be converted. A negative field width argument is taken as a - flag followed by a positive field width. A negative precision argument is taken as if the precision were omitted.

The flag characters and their meanings are:

.

The result of the conversion will be left-justified within the field. (it will be right-justified if this flag is not specified.)

+

The result of a signed conversion will always begin with a plus or minus sign. (It will begin with a sign only when a negative value is converted if this flag is not specified.)

#### space

If the first character of a signed conversion is not a sign, or if a signed conversion results in no characters, a space will be prefixed to the result. If the **space** and **+** flags both appear, the **space** flag will be ignored.

#

The result is to be converted to an "alternate form." For **o** conversion it increases the precision to force the first digit of the result to be a zero. For **x** (or **X**) conversion, a non-zero result will have "0x" (or "0X") prefixed to it. For **e**, **E**, **f**, **g**, and **g** conversions, the result will always contain a decimal-point character, even if no digits follow it. (Normally, a decimal-point character appears in the result of these conversions only if

no digit follows it). For **g** and **G** conversions, trailing zeros will not be removed from the result. For other conversions, the behavior is undefined.

0

For **d**, **i**, **o**, **u**, **x**, **X**, **e**, **E**, **f**, **g**, and **G** conversions, leading zeros (following any indication of sign or base) are used to pad to the field width; no space padding is performed. If the **0** and **-** flags both appear, the **0** flag will be ignored. For **d**, **i**, **o**, **u**, **x**, and **X** conversions, if a precision is specified, the **0** flag will be ignored. For other conversions, the behavior is undefined.

The conversion specifiers and their meanings are:

d, i

The **int** argument is converted to signed decimal in the style [-]dddd. The precision specifies the minimum number of digits to appear; if the value being converted can be represented in fewer digits, it will be expanded with leading zeros. The default precision is 1. The result of converting a zero value with a precision of zero is no characters.

o, u, x, X

The **unsigned int** argument is converted to unsigned octal (**o**), unsigned decimal (**u**), or unsigned hexadecimal notation (**x** or **X**) in the style **dddd**; the letters abcdef are used for **x** conversion and the letters ABCDEF for **X** conversion. The precision specifies the minimum number of digits to appear; if the value being converted can be represented in fewer digits, it will be expanded with leading zeros. The default precision is 1. The result of converting a zero value with a precision of zero is no characters.

f

The **double** argument is converted to decimal notation in the style [-]**dd.ddd**, where the number of digits after the decimal point character is equal to the precision specification. If the precision is missing, it is taken as 6; if the precision is zero and the **#** flag is not specified, no decimal-point character appears. If a decimal-point character appears, at least one digit appears before it. The value is rounded to the appropriate number of digits.

e, E

The **double** argument is converted in the style **[-]d.ddde+/-dd**, where there is one digit before the decimal-point character (which is non-zero if the argument is non-zero) and the number of digits after it is equal to the precision; if the precision is missing, it is taken as 6; if the precision is zero and the **#** flag is not specified, no decimal-point character appears. The value is rounded to the appropriate number of digits. The **E** conversion specifier will produce a number with **E** instead of **e** introducing the exponent. The exponent always contains at least two digits. If the value is zero, the exponent is zero.

g, G

The **double** argument is converted in style **f** or **e** (or in style **E** in the case of a **G** conversion specifier), with the precision specifying the number of significant digits. If the precision is zero, it is taken as 1. The style used depends on the value converted;

style  $\mathbf{e}$  (or  $\mathbf{E}$ ) will be used only if the exponent resulting from such a conversion is less than -4 or greater than or equal to the precision. Trailing zeros are removed from the fractional portion of the result; a decimal-point character appears only if it is followed by a digit.

С

The **int** argument is converted to an **unsigned char**, and the resulting character is written.

 $\mathbf{s}$ 

The argument should be a pointer to an array of character type. Characters from the array are written up to (but not including) a terminating null character; if the precision is specified, no more than that many characters are written. If the precision is not specified or is greater than the size of the array, the array will contain a null character.

р

The argument should be a pointer to **void**. The value of the pointer is converted to a sequence of printable characters, in hexadecimal representation (prefixed with "0x").

n

The argument should be a pointer to an integer into which the number of characters written to the output stream so far by this call to **fprintf()** is written. No argument is converted.

%

A % is written. No argument is converted. The complete conversion specification is %%.

If a conversion specification is invalid, the behavior is undefined.

If any argument is, or points to, a union or an aggregate (except for an array of character type using **s** conversion, or a pointer using **p** conversion), the behavior is undefined.

In no case does a non-existent or small field width cause truncation of a field if the result of a conversion is wider than the field width, the field is expanded to contain the conversion result.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted.

- **RETURNS** The number of characters written, or a negative value if an output error occurs.
- ERRNO Not Available
- SEE ALSO fioBaseLib, fprintf(), American National Standard for Information Systems -, Programming Language - C, ANSI X3.159-1989: Input/Output (stdio.h)

# proofUtf8()

NAME	proofUtf8() – Determine if a string represents a valid UTF-8 character
SYNOPSIS	<pre>int proofUtf8  (     const unsigned char * utf8,     const int length  )</pre>
DESCRIPTION	This routine checks a string to determine if it contains a valid UTF-8 encoded character, including $0$ .
RETURNS	If positive, the number of encoded bytes used to represent the Unicode character. If non-positive, <b>UC_FORMAT</b> indicates that the string is in an invalid format, and <b>UC_NOSRC</b> indicates that the string contains insufficient characters to represent a valid encoding, given the value of the first character.
ERRNO	Not Available
SEE ALSO	utfLib

## proofUtf8String()

NAME	<pre>proofUtf8String() – determine if a string is valid UTF-8</pre>	
SYNOPSIS	<pre>int proofUtf8String   (    const unsigned char * utf8   )</pre>	
DESCRIPTION	This routine determines if a NULL terminated string is valid UTF-8.	
RETURNS	If positive, the number of Unicode characters represented by a UTF-8 encoding. If non-positive, <b>UC_FORMAT</b> indicates that the string is of invalid format.	
ERRNO	Not Available	
SEE ALSO	utfLib	

### psr()

NAME	<b>psr()</b> – return the contents of the processor status register (SimSolaris)	
SYNOPSIS	<pre>int psr (     int taskId /* task ID, 0 means default task */ )</pre>	
DESCRIPTION	This command extracts the contents of the processor status register from the TCB of a specified task. If <i>taskId</i> is omitted or 0, the default task is assumed.	
RETURNS	The contents of the processor status register.	
ERRNO	Not Available	
SEE ALSO	dbgArchLib, VxWorks Programmer's Guide: Debugging	

# psrShow()

NAMEpsrShow() – display the meaning of a specified PSR value, symbolically (ARM)SYNOPSISSTATUS psrShow<br/>(<br/>UINT32 psrval /\* psr value to show \*/<br/>)DESCRIPTIONThis routine displays the meaning of all fields in a specified PSR value, symbolically.RETURNSOK, always.ERRNONot AvailableSEE ALSOdbgArchLib

VxWorks Kernel API Reference, 6.6 pthread\_attr\_destroy()

# pthread\_attr\_destroy()

NAME	<pre>pthread_attr_destroy() - destroy a thread attributes object (POSIX)</pre>
SYNOPSIS	<pre>int pthread_attr_destroy   (    pthread_attr_t *pAttr /* thread attributes */ )</pre>
DESCRIPTION	Destroy the thread attributes object <i>pAttr</i> . It should not be re-used until it has been reinitialized.
RETURNS	On success zero; on failure the EINVAL error code.
ERRNO	N/A
SEE ALSO	pthreadLib, pthread_attr_init()

# pthread\_attr\_getdetachstate()

NAME	<b>pthread_attr_getdetachstate()</b> – get value of detachstate attribute from thread attributes object (POSIX)		
SYNOPSIS	<pre>int pthread_attr_getdetachstate   (     const pthread_attr_t *pAttr, /* thread attributes */     int *pDetachstate /* current detach state (out) */   )</pre>		
DESCRIPTION	This routine returns the current detach state specified in the thread attributes object <i>pAttr</i> . The value is stored in the location pointed to by <i>pDetachstate</i> . Possible values for the detach state are: <b>PTHREAD_CREATE_DETACHED</b> and <b>PTHREAD_CREATE_JOINABLE</b> .		
RETURNS	zero on success, <b>EINVAL</b> if an invalid thread attribute is passed or if <i>pDetachState</i> is <b>NULL</b> .		
ERRNO	None.		
SEE ALSO	pthreadLib, pthread_attr_init( ), pthread_attr_setdetachstate( )		

2

	pthread_attr_getinheritsched()			
NAME	<b>pthread_attr_getinheritsched()</b> – get current value if inheritsched attribute in thread attributes object (POSIX)			
SYNOPSIS	<pre>int pthread_attr_getinheritsched   (     const pthread_attr_t *pAttr, /* thread attributes object */     int *pInheritsched /* inheritance mode (out) */   )</pre>			
DESCRIPTION	This routine gets the scheduling inheritance value from the thread attributes object <i>pAttr</i> .			
	Possible values are:			
	PTHREAD_INHERIT_SCHED Inherit scheduling parameters from parent thread.			
	<b>PTHREAD_EXPLICIT_SCHED</b> Use explicitly provided scheduling parameters (i.e. those specified in the thread attributes object).			
RETURNS	On success zero; on failure the EINVAL error code.			
ERRNO	N/A			
SEE ALSO	pthreadLib, pthread_attr_init( ), pthread_attr_getschedparam( ), pthread_attr_getschedpolicy( ) pthread_attr_setinheritsched( )			

# pthread\_attr\_getname( )

NAME	<pre>pthread_attr_getname() - get name of thread attribute object</pre>		
SYNOPSIS	<pre>int pthread_attr_getname   (    pthread_attr_t *pAttr,    char **name   )</pre>		
DESCRIPTION	This routine gets the name in the specified thread attributes object, <i>pAttr</i> .		
RETURNS	zero on success, <b>EINVAL</b> if an invalid thread attribute is passed or if <i>name</i> is <b>NULL</b> .		
ERRNO	None.		

VxWorks Kernel API Reference, 6.6 pthread\_attr\_getopt()

SEE ALSO pthreadLib, pthread\_attr\_setname()

## pthread\_attr\_getopt( )

NAME	<pre>pthread_attr_getopt() - get options from thread attribute object</pre>
SYNOPSIS	<pre>int pthread_attr_getopt   (    pthread_attr_t * pAttr,    int *</pre>
DESCRIPTION	This non-POSIX routine gets options from the specified thread attributes object, <i>pAttr</i> . To see the options actually applied to the VxWorks task under thread, use <b>taskOptionsGet()</b> . This routine expects the <i>pOptions</i> parameter to be a valid storage space. See <i>taskLib.h</i> for definitions of task options.
RETURNS	zero on success, <b>EINVAL</b> if an invalid thread attribute is passed or if <i>pOptions</i> is <b>NULL</b> .
ERRNO	None.
SEE ALSO	pthreadLib, pthread_attr_setopt( ), taskOptionsGet( )

# pthread\_attr\_getschedparam()

NAME	<pre>pthread_attr_getschedparam() - get value of schedparam attribute from thre object (POSIX)</pre>	ad attributes
SYNOPSIS	<pre>int pthread_attr_getschedparam   (     const pthread_attr_t *pAttr, /* thread attributes     struct sched_param *pParam /* current parameters (out)   )</pre>	*/ */
DESCRIPTION	Return, via the pointer <i>pParam</i> , the current scheduling parameters from the three object <i>pAttr</i> .	ead attributes
RETURNS	On success zero; on failure the EINVAL error code.	

ERRNO N/A SEE ALSO pthreadLib, pthread\_attr\_init(), pthread\_attr\_setschedparam(), pthread\_getschedparam(), pthread\_setschedparam(), sched\_getparam(), sched\_setparam()

## pthread\_attr\_getschedpolicy()

NAME	<b>pthread_attr_getschedpolicy()</b> – get schedpolicy attribute from thread attributes object (POSIX)	
SYNOPSIS	<pre>int pthread_attr_getschedpolicy   (     const pthread_attr_t *pAttr, /* thread attributes */     int *pPolicy /* current policy (out) */   )</pre>	
DESCRIPTION	This routine returns, via the pointer <i>pPolicy</i> , the current scheduling policy in the thread attributes object specified by <i>pAttr</i> . Possible values for VxWorks systems are <b>SCHED_RR</b> , <b>SCHED_FIFO</b> and <b>SCHED_OTHER</b> .	
RETURNS	On success zero; on failure the EINVAL error code.	
ERRNO	N/A	
SEE ALSO	<pre>pthreadLib, pthread_attr_init(), pthread_attr_setschedpolicy(), pthread_getschedparam(), pthread_setschedparam(), sched_setscheduler(), sched_getscheduler()</pre>	

### pthread\_attr\_getscope( )

NAME	<pre>pthread_attr_getscope( ) - g</pre>	et contention scope fro	om thread attributes (POSIX)
SYNOPSIS	<pre>int pthread_attr_getscop (</pre>	e *pAttr,	/* thread attributes object
	int */ )	*pContentionScope	/* contention scope (out)

VxWorks Kernel API Reference, 6.6 pthread\_attr\_getstackaddr( )

DESCRIPTION	Reads the current contention scope setting from a thread attributes object. For VxWorks the is always <b>PTHREAD_SCOPE_SYSTEM</b> . If the thread attributes object is uninitialized then <b>EINVAL</b> will be returned. The contention scope is returned in the location pointed to by <i>pContentionScope</i> .	
RETURNS	On success zero; on failure the EINVAL error code.	
ERRNO	N/A	
SEE ALSO	pthreadLib, pthread_attr_init( ), pthread_attr_setscope( )	

### pthread\_attr\_getstackaddr( )

NAME	<b>pthread_attr_getstackaddr( )</b> – get value of stackaddr attribute from thread attributes object (POSIX)
SYNOPSIS	<pre>int pthread_attr_getstackaddr   (     const pthread_attr_t *pAttr, /* thread attributes */     void **ppStackaddr /* current stack address (out) */   )</pre>
DESCRIPTION	This routine returns the stack address from the thread attributes object <i>pAttr</i> in the location pointed to by <i>ppStackaddr</i> .
RETURNS	zero on success, <b>EINVAL</b> if an invalid thread attribute is passed or if <i>ppStackaddr</i> is <b>NULL</b> .
ERRNO	None.
SEE ALSO	pthreadLib, pthread_attr_init( ), pthread_attr_getstacksize( ), pthread_attr_setstackaddr( )

## pthread\_attr\_getstacksize( )

NAME pthread\_attr\_getstacksize() – get stack value of stacksize attribute from thread attributes object (POSIX)

SYNOPSIS	int pthread_attr_getstacksize		
	(		
	<pre>const pthread_attr_t *pAttr,</pre>	/* thread attributes	*/

 size\_t
 \*pStacksize /\* current stack size (out) \*/

 DESCRIPTION
 This routine gets the current stack size from the thread attributes object *pAttr* and places it in the location pointed to by *pStacksize*.

 RETURNS
 zero on success, EINVAL if an invalid thread attribute is passed or if *pStackSize* is NULL.

 ERRNO
 None.

 SEE ALSO
 pthreadLib, pthread\_attr\_init(), pthread\_attr\_setstacksize(), pthread\_attr\_getstackaddr()

#### pthread\_attr\_init( )

NAME	<pre>pthread_attr_init() – initialize thread attributes object (POSIX)</pre>
SYNOPSIS	<pre>int pthread_attr_init   (    pthread_attr_t *pAttr /* thread attributes */ )</pre>
DESCRIPTION	This routine initializes a thread attributes object. If <i>pAttr</i> is <b>NULL</b> then this function will return <b>EINVAL</b> .
	The attributes that are set by default are as follows:
	Stack Address NULL - allow the system to allocate the stack.
	<b>Stack Size</b> 0 - use the VxWorks <b>taskLib</b> default stack size.
	Detach State PTHREAD_CREATE_JOINABLE
	Contention Scope PTHREAD_SCOPE_SYSTEM
	Scheduling Inheritance PTHREAD_INHERIT_SCHED
	Scheduling Policy SCHED_OTHER (i.e. active VxWorks native scheduling policy).
	Scheduling Priority Use pthreadLib default priority

VxWorks Kernel API Reference, 6.6 pthread\_attr\_setdetachstate()

	Note that the scheduling policy and priority values are only used if the scheduling inheritance mode is changed to <b>PTHREAD_EXPLICIT_SCHED</b> - see <b>pthread_attr_setinheritsched()</b> for information.
	Additionally, VxWorks-specific attributes are being set as follows:
	Task Name NULL - the task name is automatically generated.
	Task Options VX_FP_TASK
RETURNS	On success zero; on failure the EINVAL error code.
ERRNO	N/A
SEE ALSO	<pre>pthreadLib, pthread_attr_destroy(), pthread_attr_getdetachstate(), pthread_attr_getinheritsched(), pthread_attr_getschedparam(), pthread_attr_getschedpolicy(), pthread_attr_getscope(), pthread_attr_getstackaddr(), pthread_attr_getstacksize(), pthread_attr_setdetachstate(), pthread_attr_setinheritsched(), pthread_attr_setschedparam(), pthread_attr_setschedpolicy(), pthread_attr_setscope(), pthread_attr_setstackaddr(), pthread_attr_setstacksize(), pthread_attr_setscope(), pthread_attr_setstackaddr(),</pre>

#### pthread\_attr\_setopt( ) (VxWorks extension)

# pthread\_attr\_setdetachstate()

NAME	<pre>pthread_attr_setdetachstate() - set detachstate attribute in thread attributes object (POSIX)</pre>
SYNOPSIS	<pre>int pthread_attr_setdetachstate   (    pthread_attr_t *pAttr, /* thread attributes */    int detachstate /* new detach state */ )</pre>
DESCRIPTION	This routine sets the detach state in the thread attributes object <i>pAttr</i> . The new detach state specified by <i>detachstate</i> must be one of <b>PTHREAD_CREATE_DETACHED</b> or <b>PTHREAD_CREATE_JOINABLE</b> . Any other values will cause an error to be returned ( <b>EINVAL</b> ).
RETURNS	On success zero; on failure the EINVAL error code.
ERRNO	N/A

### pthread\_attr\_setinheritsched()

NAME	<b>pthread_attr_setinheritsched()</b> – set inheritsched attribute in thread attribute object (POSIX)
SYNOPSIS	<pre>int pthread_attr_setinheritsched   (    pthread_attr_t *pAttr, /* thread attributes object */    int inheritsched /* inheritance mode */   )</pre>
DESCRIPTION	This routine sets the scheduling inheritance to be used when creating a thread with the thread attributes object specified by $pAttr$ .
	Possible values are:
	PTHREAD_INHERIT_SCHED Inherit scheduling parameters from parent thread.
	<b>PTHREAD_EXPLICIT_SCHED</b> Use explicitly provided scheduling parameters (i.e. those specified in the thread attributes object).
RETURNS	On success zero; on failure the <b>EINVAL</b> error code.
ERRNO	N/A
SEE ALSO	pthreadLib, pthread_attr_getinheritsched( ), pthread_attr_init( ), pthread_attr_setschedparam( ), pthread_attr_setschedpolicy( )

# pthread\_attr\_setname( )

NAME	<pre>pthread_attr_setname() - set name in thread attribute object</pre>
SYNOPSIS	<pre>int pthread_attr_setname   (     pthread_attr_t *pAttr,</pre>
	char *name

VxWorks Kernel API Reference, 6.6 pthread\_attr\_setopt()

SEE ALSO	pthreadLib, pthread_attr_getname( )
ERRNO	None.
RETURNS	zero on success, <b>EINVAL</b> if an invalid thread attribute is passed.
DESCRIPTION	This routine sets the name in the specified thread attributes object, <i>pAttr</i> .

#### pthread\_attr\_setopt( )

NAME	<pre>pthread_attr_setopt() - set options in thread attribute object</pre>
SYNOPSIS	<pre>int pthread_attr_setopt   (    pthread_attr_t * pAttr,    int options  )</pre>
DESCRIPTION	This non-POSIX routine sets options in the specified thread attributes object, <i>pAttr</i> . This allows for specifying a non-default set of options for the VxWorks task acting as a thread. Additional options may be applied to the task once the thread has been created via the <b>taskOptionsSet()</b> API.
	Note that the task options provided through this routine will supersede the default options otherwise applied at thread creation.
	See <i>taskLib.h</i> for definitions of valid task options.
RETURNS	zero on success, <b>EINVAL</b> if an invalid thread attribute is passed.
ERRNO	None.
SEE ALSO	pthreadLib, pthread_attr_getopt( ), taskOptionsSet( )

# pthread\_attr\_setschedparam()

NAME pthread\_attr\_setschedparam() – set schedparam attribute in thread attributes object (POSIX)

SYNOPSIS int pthread\_attr\_setschedparam

SEE ALSO	<pre>pthreadLib, pthread_attr_getschedparam(), pthread_attr_init(), pthread_getschedparam(), pthread_setschedparam(), pthread_attr_setinheritsched(), sched_getparam(), sched_setparam()</pre>
ERRNO	N/A
RETURNS	On success zero; on failure the EINVAL error code.
DESCRIPTION	Set the scheduling parameters in the thread attributes object <i>pAttr</i> . The scheduling parameters are essentially the thread's priority. Note that the <b>PTHREAD_EXPLICIT_SCHED</b> mode must be set (see <b>pthread_attr_setinheritsched()</b> for information) for the priority to take effect.
	( pthread_attr_t *pAttr, /* thread attributes */ const struct sched_param *pParam /* new parameters */ )

# pthread\_attr\_setschedpolicy( )

NAME	<b>pthread_attr_setschedpolicy()</b> – set schedpolicy attribute in thread attributes object (POSIX)
SYNOPSIS	<pre>int pthread_attr_setschedpolicy   (     pthread_attr_t *pAttr, /* thread attributes */     int        policy /* new policy</pre>
DESCRIPTION	Select the thread scheduling policy. The default scheduling policy is to inherit the current system setting. Unlike the POSIX model, scheduling policies under VxWorks are global. If a scheduling policy is being set explicitly, the PTHREAD_EXPLICIT_SCHED mode must be set (see pthread_attr_setinheritsched() for information), and the selected scheduling policy must match the global scheduling policy in place at the time; failure to do so will result in pthread_create() failing with the error EPERM.
	POSIX defines the following policies:
	SCHED_RR Realtime, round-robin scheduling.
	SCHED_FIFO Realtime, first-in first-out scheduling.
	SCHED_OTHER Other, active VxWorks native scheduling policy.

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Although the SCHED\_RR and SCHED\_FIFO policies can be set when the precaution described above is respected, using the SCHED\_OTHER policy instead is always ensured to be successful.

**RETURNS** On success zero; on failure the **EINVAL** error code.

ERRNO N/A

 SEE ALSO
 pthreadLib, pthread\_attr\_getschedpolicy(), pthread\_attr\_init(),

 pthread\_attr\_setinheritsched(), pthread\_getschedparam(), pthread\_setschedparam(),

 sched\_setscheduler(), sched\_getscheduler()

### pthread\_attr\_setscope( )

NAME	<pre>pthread_attr_setscope( ) - set contention scope for thread attributes (POSIX)</pre>
SYNOPSIS	<pre>int pthread_attr_setscope   (    pthread_attr_t *pAttr, /* thread attributes object */    int contentionScope /* new contention scope */   )</pre>
DESCRIPTION	For VxWorks <b>PTHREAD_SCOPE_SYSTEM</b> is the only supported contention scope. If the <b>PTHREAD_SCOPE_PROCESS</b> value is passed to this function this will result in <b>ENOTSUP</b> being returned.
RETURNS	On success zero; on failure the EINVAL or ENOTSUP error code.
ERRNO	N/A
SEE ALSO	pthreadLib, pthread_attr_getscope( ), pthread_attr_init( )

#### pthread\_attr\_setstackaddr( )

NAME	<pre>pthread_attr_setstackaddr() - set sta</pre>	ackaddr attribute in thread attributes	object (POSIX)
SYNOPSIS	int pthread_attr_setstackaddr		
	, pthread_attr_t *pAttr,	/* thread attributes	*/

	void *pStackaddr /* new stack address */ )
DESCRIPTION	This routine sets the stack address in the thread attributes object $pAttr$ to be $pStackaddr$ . On VxWorks this address must be the lowest address of the stack regardless of what the thread considers as the stack base or the stack end.
	No alignment constraints are imposed by the pthread library so the thread's stack can be obtained via a simple call to <b>malloc()</b> or <b>memPartAlloc()</b> .
	The memory area used a stack is not automatically freed when the thread exits. This operation cannot be done via the exiting thread's cleanup stack since the cleanup handler routines use the same stack as the thread. Therefore freeing the stack space must be done by the code which allocated the thread's stack once the thread's task no longer exists in the system.
	The stack size is set using the routine <b>pthread_attr_setstacksize()</b> . Note that failure to set the stack size when a stack address is provided will result in an <b>EINVAL</b> error status returned by <b>pthread_create()</b> .
RETURNS	zero on success, <b>EINVAL</b> if an invalid thread attribute is passed.
ERRNO	None.
SEE ALSO	pthreadLib, pthread_attr_getstacksize( ), pthread_attr_setstacksize( ), pthread_attr_init( )

# pthread\_attr\_setstacksize()

NAME	pthread_attr_setstacksize() – set stacksize attribute in thread attributes object (POSIX)
SYNOPSIS	<pre>int pthread_attr_setstacksize   (    pthread_attr_t *pAttr, /* thread attributes */    size_t stacksize /* new stack size */ )</pre>
DESCRIPTION	This routine sets the thread stack size (in bytes) in the specified thread attributes object, $pAttr$ .
	The stack address is set using the routine <b>pthread_attr_setstackaddr()</b> . Note that failure to set the stack size when a stack address is provided will result in an <b>EINVAL</b> error status returned by <b>pthread_create()</b> .
RETURNS	EINVAL if the stack size is lower than PTHREAD_STACK_MIN or if an invalid thread attribute is passed. Zero otherwise.

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 ERRNO
 None.

 SEE ALSO
 pthreadLib, pthread\_attr\_getstacksize(), pthread\_attr\_setstackaddr(), pthread\_attr\_init(), pthread\_create()

### pthread\_cancel()

NAME	<b>pthread_cancel()</b> – cancel execution of a thread (POSIX)		
SYNOPSIS	<pre>int pthread_cancel   (    pthread_t thread /* thread to cancel */   )</pre>		
DESCRIPTION	This routine sends a cancellation request to the thread specified by <i>thread</i> . Depending on the settings of that thread, it may ignore the request, terminate immediately or defer termination until it reaches a cancellation point.		
	When the thread terminates it performs as if <b>pthread_exit( )</b> had been called with the exit status <b>PTHREAD_CANCELED</b> .		
IMPLEMENTATION NO	IMPLEMENTATION NOTE		
	In VxWorks, asynchronous thread cancellation is accomplished using a signal. The signal <b>SIGCNCL</b> has been reserved for this purpose. Applications should take care not to block or handle this signal.		
RETURNS	On success zero; on failure the ESRCH error code.		
ERRNO	N/A		
SEE ALSO	pthreadLib, pthread_exit(), pthread_setcancelstate(), pthread_setcanceltype(), pthread_testcancel()		
SEE ALSO	pthreadLib, pthread_exit( ), pthread_setcancelstate( ), pthread_setcanceltype( ), pthread_testcancel( )		

# pthread\_cleanup\_pop()

**NAME pthread\_cleanup\_pop()** – pop a cleanup routine off the top of the stack (POSIX)

SYNOPSIS void pthread\_cleanup\_pop

```
(
int run /* execute handler? */
)
```

**DESCRIPTION** This routine removes the cleanup handler routine at the top of the cancellation cleanup stack of the calling thread and executes it if *run* is non-zero. The routine should have been added using the **pthread\_cleanup\_push()** function.

Once the routine is removed from the stack it will no longer be called when the thread exits.

RETURNS N/A

ERRNO N/A

SEE ALSO pthreadLib, pthread\_cleanup\_push(), pthread\_exit()

### pthread\_cleanup\_push()

NAME	<pre>pthread_cleanup_push() - pushes a routine onto the cleanup stack (POSIX)</pre>
SYNOPSIS	<pre>void pthread_cleanup_push   (    void (*routine)(void *), /* cleanup routine */    void *arg /* argument */   )</pre>
DESCRIPTION	This routine pushes the specified cancellation cleanup handler routine, <i>routine</i> , onto the cancellation cleanup stack of the calling thread. When a thread exits and its cancellation cleanup stack is not empty, the cleanup handlers are invoked with the argument <i>arg</i> in LIFO order from the cancellation cleanup stack.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	pthreadLib, pthread_cleanup_pop(), pthread_exit()

## pthread\_cond\_broadcast( )

NAME	<b>pthread_cond_broadcast( )</b> – unblock all threads waiting on a condition (POSIX)
SYNOPSIS	<pre>int pthread_cond_broadcast   (    pthread_cond_t *pCond   )</pre>
DESCRIPTION	This function unblocks all threads blocked on the condition variable $pCond$ . Nothing happens if no threads are waiting on the specified condition variable.
RETURNS	On success zero; on failure the EINVAL error code.
ERRNO	N/A
SEE ALSO	pthreadLib, pthread_condattr_init( ), pthread_condattr_destroy( ), pthread_cond_destroy( ), pthread_cond_init( ), pthread_cond_signal( ), pthread_cond_timedwait( ), pthread_cond_wait( )

# pthread\_cond\_destroy()

NAME	<pre>pthread_cond_destroy() - destroy a condition variable (POSIX)</pre>
SYNOPSIS	<pre>int pthread_cond_destroy   (     pthread_cond_t *pCond /* condition variable */ )</pre>
DESCRIPTION	This routine destroys the condition variable pointed to by <i>pCond</i> . No threads can be waiting on the condition variable when this function is called. If there are threads waiting on the condition variable, then <b>pthread_cond_destroy()</b> returns <b>EBUSY</b> .
RETURNS	On success zero; on failure a non-zero error code. EINVAL EBUSY
ERRNO	
SEE ALSO	pthreadLib, pthread_condattr_init( ), pthread_condattr_destroy( ), pthread_cond_broadcast( ), pthread_cond_init( ), pthread_cond_signal( ), pthread_cond_timedwait( ), pthread_cond_wait( )

# pthread\_cond\_init( )

NAME	<pre>pthread_cond_init() - initialize condition variable (POSIX)</pre>
SYNOPSIS	<pre>int pthread_cond_init   (     pthread_cond_t *pCond, /* condition variable */     pthread_condattr_t *pAttr /* condition variable attributes */   )</pre>
DESCRIPTION	This function initializes a condition variable. A condition variable is a synchronization device that allows threads to block until some predicate on shared data is satisfied. The basic operations on conditions are to signal the condition (when the predicate becomes true), and wait for the condition, blocking the thread until another thread signals the condition.
	A condition variable must always be associated with a mutex to avoid a race condition between the wait and signal operations.
	If <i>pAttr</i> is <b>NULL</b> then the default attributes are used as specified by POSIX; if <i>pAttr</i> is non- <b>NULL</b> then it is assumed to point to a condition attributes object initialized by <b>pthread_condattr_init()</b> , and those are the attributes used to create the condition variable.
RETURNS	On success zero; on failure a non-zero error code:
	EINVAL
ERRNO	
SEE ALSO	pthreadLib, pthread_condattr_init( ), pthread_condattr_destroy( ), pthread_cond_broadcast( ), pthread_cond_destroy( ), pthread_cond_signal( ), pthread_cond_timedwait( ), pthread_cond_wait( )

# pthread\_cond\_signal()

NAME	<pre>pthread_cond_signal() - unblock a thread waiting on a condition (POSIX)</pre>
SYNOPSIS	int pthread_cond_signal ( pthread_cond_t *pCond )
DESCRIPTION	This routine unblocks one thread waiting on the specified condition variable <i>pCond</i> . If no threads are waiting on the condition variable then this routine does nothing; if more than one thread is waiting, then one will be released, but it is not specified which one.

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SEE ALSO	<pre>pthreadLib, pthread_condattr_init(), pthread_condattr_destroy(),     pthread_cond_broadcast(), pthread_cond_destroy(), pthread_cond_init(),     pthread_cond_timedwait(), pthread_cond_wait()</pre>
ERRNO	N/A
RETURNS	On success zero; on failure the <b>EINVAL</b> error code.

#### pthread\_cond\_timedwait( )

NAME pthread\_cond\_timedwait() – wait for a condition variable with a timeout (POSIX) SYNOPSIS int pthread cond timedwait /\* condition variable pthread\_cond\_t \*pCond, \*/ pthread\_mutex\_t \*pMutex, /\* POSIX mutex \*/ const struct timespec \*pAbstime /\* timeout time \*/ ) DESCRIPTION This function atomically releases the mutex *pMutex* and waits for another thread to signal the condition variable *pCond*. As with **pthread\_cond\_wait()**, the mutex must be locked by the calling thread when **pthread\_cond\_timedwait()** is called. If the condition variable is signalled before the system time reaches the time specified by *pAbsTime*, then the mutex is re-acquired and the calling thread unblocked. If the system time reaches or exceeds the time specified by *pAbsTime* before the condition is signalled, then the mutex is re-acquired, the thread unblocked and ETIMEDOUT returned. If the calling thread gets cancelled while pending on the condition variable pthread\_cond\_timedwait() will also re-acquire the mutex prior to executing the cancellation cleanup handlers (if any). The mutex will however be released prior to the thread exiting so that this mutex can be used by other threads. NOTE The timeout is specified as an absolute value of the system clock in a *timespec* structure (see **clock\_gettime()** for more information). This is different from most VxWorks timeouts which are specified in ticks relative to the current time. On success zero; on failure a non-zero error code: RETURNS

#### EINVAL ETIMEDOUT

#### ERRNO

SEE ALSO pthrbeadLib, pthread\_condattr\_init(), pthread\_condattr\_destroy(), pthread\_cond\_broadcast(), pthread\_cond\_destroy(), pthread\_cond\_init(), pthread\_cond\_signal(), pthread\_cond\_wait()

# pthread\_cond\_wait()

NAME	<b>pthread_cond_wait()</b> – wait for a condition variable (POSIX)
SYNOPSIS	<pre>int pthread_cond_wait   (    pthread_cond_t *pCond, /* condition variable */    pthread_mutex_t *pMutex /* POSIX mutex */   )</pre>
DESCRIPTION	This function atomically releases the mutex $pMutex$ and waits for the condition variable $pCond$ to be signalled by another thread. The mutex must be locked by the calling thread when <b>pthread_cond_wait()</b> is called; if it is not then this function returns an error <b>(EINVAL)</b> .
	Before returning to the calling thread, <b>pthread_cond_wait()</b> re-acquires the mutex.
	If the calling thread gets cancelled while pending on the condition variable <b>pthread_cond_wait()</b> will also re-acquire the mutex prior to executing the cancellation cleanup handlers (if any). The mutex will however be released prior to the thread exiting so that this mutex can be used by other threads.
RETURNS	On success zero; on failure the <b>EINVAL</b> error code.
ERRNO	N/A
SEE ALSO	<pre>pthreadLib, pthread_condattr_init(), pthread_condattr_destroy(), pthread_cond_broadcast(), pthread_cond_destroy(), pthread_cond_init(), pthread_cond_signal(), pthread_cond_timedwait()</pre>

VxWorks Kernel API Reference, 6.6 pthread\_condattr\_destroy()

# pthread\_condattr\_destroy()

NAME	<pre>pthread_condattr_destroy() - destroy a condition attributes object (POSIX)</pre>
SYNOPSIS	<pre>int pthread_condattr_destroy   (    pthread_condattr_t *pAttr /* condition variable attributes */   )</pre>
DESCRIPTION	This routine destroys the condition attribute object <i>pAttr</i> . It must not be reused until it is reinitialized.
RETURNS	Always returns zero.
ERRNO	None.
SEE ALSO	pthreadLib, pthread_cond_init( ), pthread_condattr_init( )

# pthread\_condattr\_init( )

NAME	<pre>pthread_condattr_init() – initialize a condition attribute object (POSIX)</pre>
SYNOPSIS	<pre>int pthread_condattr_init   (    pthread_condattr_t *pAttr /* condition variable attributes */ )</pre>
DESCRIPTION	This routine initializes the condition attribute object <i>pAttr</i> and fills it with default values for the attributes.
RETURNS	On success zero; on failure a non-zero error code: EINVAL
ERRNO	
SEE ALSO	pthreadLib, pthread_cond_init( ), pthread_condattr_destroy( )

	pthread_create()		
NAME	<pre>pthread_create( ) - create a thread (POSIX)</pre>		
SYNOPSIS	<pre>int pthread_create   (   pthread_t * pThread, /* Thread ID (out) */   const pthread_attr_t * pAttr, /* Thread attributes object */   void * (*startRoutine)(void *), /* Entry function */   void * arg /* Entry function argument */   )</pre>		
DESCRIPTION	This routine creates a new thread and if successful writes its ID into the location pointed to by <i>pThread</i> . If <i>pAttr</i> is <b>NULL</b> then default attributes are used. The new thread executes <i>startRoutine</i> with <i>arg</i> as its argument.		
	The new thread's cancelability state and cancelability type are respectively set to <b>PTHREAD_CANCEL_ENABLE</b> and <b>PTHREAD_CANCEL_DEFERRED</b> .		
RETURNS	On success zero; on failure one of the following non-zero error codes:		
	<b>EINVAL</b> can be returned when the value specified by <i>pAttr</i> is invalid, when a user-supplied stack address is provided but the stack size is invalid, and when the <i>pThread</i> parameter is null.		
	EAGAIN can be returned when not enough memory is available to either create the thread or create a resource required for the thread.		
	<b>EPERM</b> the explicit scheduling policy does not match the VxWorks scheduling policy currently in effect.		
ERRNO	N/A		
SEE ALSO	pthreadLib, pthread_exit( ), pthread_join( ), pthread_detach( )		
	nthread_detach()		

#### **NAME pthread\_detach()** – dynamically detach a thread (POSIX)

**SYNOPSIS** int pthread\_detach

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(
pthread\_t thread /\* thread to detach \*/
)

**DESCRIPTION** This routine puts the thread *thread* into the detached state. This prevents other threads from synchronizing on the termination of the thread using **pthread\_join()**.

 RETURNS
 On success zero; on failure a non-zero error code:

 EINVAL
 ESRCH

 ERRNO
 N/A

SEE ALSO pthreadLib, pthread\_join()

### pthread\_equal()

NAME	<pre>pthread_equal() - compare thread IDs (POSIX)</pre>		
SYNOPSIS	<pre>int pthread_equal   (    pthread_t t1, /* thread one */    pthread_t t2 /* thread two */   )</pre>		
DESCRIPTION	Tests the equality of the two threads $t1$ and $t2$ .		
RETURNS	Non-zero if <i>t</i> 1 and <i>t</i> 2 refer to the same thread, otherwise zero.		
ERRNO	Not Available		
SEE ALSO	pthreadLib		

### pthread\_exit()

NAME pthread\_exit() – terminate a thread (POSIX)

SYNOPSIS void pthread\_exit

```
(
void *status /* exit status */
)
```

**DESCRIPTION** This function terminates the calling thread. All cleanup handlers that have been set for the calling thread with **pthread\_cleanup\_push()** are executed in reverse order (the most recently added handler is executed first). Termination functions for thread-specific data are then called for all keys that have non-NULL values associated with them in the calling thread (see **pthread\_key\_create()** for more details). Finally, execution of the calling thread is stopped.

The *status* argument is the return value of the thread and can be consulted from another thread using **pthread\_join()** unless this thread was detached (i.e. a call to **pthread\_detach()** had been made for it, or it was created in the detached state).

All threads that remain *joinable* at the time they exit should ensure that **pthread\_join()** is called on their behalf by another thread to reclaim the resources that they hold.

**RETURNS** Does not return.

ERRNO N/A

SEE ALSO pthreadLib, pthread\_cleanup\_push(), pthread\_detach(), pthread\_join(), pthread\_key\_create()

#### pthread\_getschedparam()

NAME	pthread_getschedparam()	– get value	of schedparam attribute from a thread	d (POSIX)
SYNOPSIS	int pthread_getschedpar (	am		
	pthread_t	thread,	/* thread	*/
	int	*pPolicy,	/* current policy (out)	*/
	struct sched_param )	*pParam	/* current parameters (out)	*/
DESCRIPTION	This routine reads the current scheduling parameters and policy of the thread specified by <i>thread</i> . The information is returned via <i>pPolicy</i> and <i>pParam</i> .			
	Note that this routine actually always maps the current VxWorks scheduling policy on one of the two following POSIX scheduling policies: <b>SCHED_FIFO</b> or <b>SCHED_RR</b> . The <b>SCHED_OTHER</b> policy can therefore never be returned even if it has been set via <b>pthread_setschedparam()</b> .			
RETURNS	On success zero; on failure	the ESRCH	error code.	

VxWorks Kernel API Reference, 6.6 pthread\_getspecific()

 

 ERRNO
 N/A

 SEE ALSO
 pthreadLib, pthread\_attr\_getschedparam() pthread\_attr\_getschedpolicy(), pthread\_attr\_setschedparam() pthread\_attr\_setschedpolicy(), pthread\_setschedparam(), sched\_getparam(), sched\_setparam()

#### pthread\_getspecific( )

NAME	<pre>pthread_getspecific() - get thread specific data (POSIX)</pre>
SYNOPSIS	<pre>void *pthread_getspecific   (    pthread_key_t key /* thread specific data key */ )</pre>
DESCRIPTION	This routine returns the value associated with the thread specific data key <i>key</i> for the calling thread.
RETURNS	The value associated with <i>key</i> , or <b>NULL</b> .
ERRNO	N/A
SEE ALSO	pthreadLib, pthread_key_create(), pthread_key_delete(), pthread_setspecific()

#### pthread\_join()

NAME	<pre>pthread_join() - wait for a thread to terminate (POSIX)</pre>		
SYNOPSIS	<pre>int pthread_join   (     pthread_t thread, /* thread to wait for */     void **ppStatus /* exit status of thread (out) */   )</pre>		
DESCRIPTION	This routine will block the calling thread until the thread specified by <i>thread</i> terminates, or is canceled. The thread must be in the joinable state, i.e. it cannot have been detached by a call to <b>pthread_detach()</b> , or created in the detached state.		

If *ppStatus* is not **NULL** and **pthread\_join()** returns successfully, when *thread* terminates its exit status will be stored in the specified location. The exit status will be either the value

	passed to <b>pthread_exit( )</b> , or <b>PTHREAD_CANCELED</b> if the thread was canceled or the thread was deleted by a VxWorks task.
	Only one thread can wait for the termination of a given thread. If another thread is already waiting when this function is called an error will be returned (EINVAL).
	If the calling thread passes its own ID in <i>thread</i> , the call will fail with the error <b>EDEADLK</b> .
NOTE	All threads that remain <i>joinable</i> at the time they exit should ensure that <b>pthread_join()</b> is called on their behalf by another thread to reclaim the resources that they hold.
RETURNS	On success zero; on failure a non-zero error code:
	EINVAL
	ESRCH
	EDEADLK
ERRNO	N/A
SEE ALSO	pthreadLib, pthread_detach(), pthread_exit()

# pthread\_key\_create()

NAME	<pre>pthread_key_create() - create a thread specific data key (POSIX)</pre>		
SYNOPSIS	<pre>int pthread_key_create   (     pthread_key_t *pKey, /* thread specific data key */     void (*destructor)(void *) /* destructor function */     )</pre>		
DESCRIPTION	This routine allocates a new thread specific data key. The key is stored in the location pointed to by <i>key</i> . The value initially associated with the returned key is <b>NULL</b> in all currently executing threads. If the maximum number of keys are already allocated, the function returns an error ( <b>EAGAIN</b> ).		
	The <i>destructor</i> parameter specifies a destructor function associated with the key. When a thread terminates via <b>pthread_exit( )</b> , or by cancellation, <i>destructor</i> is called with the value associated with the key in that thread as an argument. The destructor function is <b>not</b> called if that value is <b>NULL</b> . The order in which destructor functions are called at thread termination time is unspecified.		
	It is the user's responsibility to call <b>pthread_key_delete()</b> when the memory associated with the key is no longer required, and to ensure that no threads access the key after it has been deleted. Failure to do this can return unexpected results, and can cause memory leaks.		

VxWorks Kernel API Reference, 6.6 pthread\_key\_delete()

**RETURNS** On success zero; on failure the **EAGAIN** error code.

ERRNO N/A

SEE ALSO pthreadLib, pthread\_getspecific(), pthread\_key\_delete(), pthread\_setspecific()

# pthread\_key\_delete()

NAME	<pre>pthread_key_delete() - delete a thread specific data key (POSIX)</pre>
SYNOPSIS	<pre>int pthread_key_delete   (    pthread_key_t key /* thread specific data key to delete */ )</pre>
DESCRIPTION	This routine deletes the thread specific data associated with <i>key</i> , and deallocates the key itself. It does not call any destructor associated with the key. Any attempt to use key following the call to <b>pthread_key_delete()</b> results in undefined behavior.
RETURNS	On success zero; on failure the EINVAL error code.
ERRNO	N/A
SEE ALSO	pthreadLib, pthread_key_create()

## pthread\_kill()

NAME	<b>pthread_kill( )</b> – send a signal to a thread (POSIX)
SYNOPSIS	<pre>int pthread_kill   (    pthread_t thread, /* thread to signal */    int sig /* signal to send */   )</pre>
DESCRIPTION	This routine sends signal number <i>sig</i> to the thread specified by <i>thread</i> . The signal is delivered and handled as described for the <b>kill()</b> function.
RETURNS	On success zero; on failure one of the following non-zero error codes: ESRCH, EINVAL

ERRNO N/A

SEE ALSO pthreadLib, kill(), pthread\_sigmask(), sigprocmask(), sigaction(), sigsuspend(), sigwait()

## pthread\_mutex\_destroy()

NAME	<pre>pthread_mutex_destroy() - destroy a mutex (POSIX)</pre>
SYNOPSIS	<pre>int pthread_mutex_destroy   (     pthread_mutex_t *pMutex /* POSIX mutex */   )</pre>
DESCRIPTION	This routine destroys a mutex object, freeing the resources it might hold. The mutex can be safely destroyed when unlocked. On VxWorks a thread may destroy a mutex that it owns (i.e. that the thread has locked). If the mutex is locked by an other thread this routine will return an error ( <b>EBUSY</b> ).
RETURNS	On success zero; on failure a non-zero error code: EINVAL EBUSY
ERRNO	N/A
SEE ALSO	<pre>pthreadLib, semLib, semMLib, pthread_mutex_init(), pthread_mutex_lock(), pthread_mutex_trylock(), pthread_mutex_unlock(), pthread_mutexattr_init(), semDelete()</pre>

# pthread\_mutex\_getprioceiling()

NAME	<b>pthread_mutex_getprioceiling()</b> – get the value of the prioceiling attribute of a mutex (POSIX)
SYNOPSIS	int pthread_mutex_getprioceiling (
	pthread_mutex_t *pMutex, /* POSIX mutex */
	int *pPrioceiling /* current priority ceiling (out) */
	)

VxWorks Kernel API Reference, 6.6 pthread\_mutex\_init()

DESCRIPTION	This function gets the current value of the prioceiling attribute of a mutex. Unless the mutex was created with a protocol attribute value of <b>PTHREAD_PRIO_PROTECT</b> , this value is meaningless.
RETURNS	On success zero; on failure the EINVAL error code.
ERRNO	N/A
SEE ALSO	pthreadLib, pthread_mutex_setprioceiling(), pthread_mutexattr_getprioceiling(), pthread_mutexattr_setprioceiling()

# pthread\_mutex\_init()

NAME	<pre>pthread_mutex_init() - initialize mutex from attributes object (POSIX)</pre>			
SYNOPSIS	<pre>int pthread_mutex_init   (    pthread_mutex_t *pMutex, /* POSIX mutex */    const pthread_mutexattr_t *pAttr /* mutex attributes */   )</pre>			
DESCRIPTION	This routine initializes the mutex object pointed to by <i>pMutex</i> according to the mutex attributes specified in <i>pAttr</i> . If <i>pAttr</i> is <b>NULL</b> , default attributes are used as defined in the POSIX specification. If <i>pAttr</i> is <b>non-NULL</b> then it is assumed to point to a mutex attributes object initialized by <b>pthread_mutexattr_init()</b> , and those are the attributes used to create the mutex.			
RETURNS	On success zero; on failure a non-zero error code: EINVAL			
ERRNO	N/A			
SEE ALSO	<pre>pthreadLib, semLib, semMLib, pthread_mutex_destroy(), pthread_mutex_lock(), pthread_mutex_trylock(), pthread_mutex_unlock(), pthread_mutexattr_init(), semMCreate()</pre>			

	pthread_mutex_lock()		
NAME	<pre>pthread_mutex_lock() - lock a mutex (POSIX)</pre>		
SYNOPSIS	<pre>int pthread_mutex_lock   (    pthread_mutex_t *pMutex /* POSIX mutex */   )</pre>		
DESCRIPTION	This routine locks the mutex specified by <i>pMutex</i> . If the mutex is currently unlocked, it becomes locked, and is said to be owned by the calling thread. In this case <b>pthread_mutex_lock()</b> returns immediately.		
	If the mutex is already locked by another thread, <b>pthread_mutex_lock()</b> blocks the calling thread until the mutex is unlocked by its current owner.		
	If it is already locked by the calling thread, pthread_mutex_lock will deadlock on itself and the thread will block indefinitely.		
RETURNS	On success zero; on failure the EINVAL error code.		
ERRNO	N/A		
SEE ALSO	<pre>pthreadLib, semLib, semMLib, pthread_mutex_init(), pthread_mutex_lock(), pthread_mutex_trylock(), pthread_mutex_unlock(), pthread_mutexattr_init(), semTake()</pre>		

# pthread\_mutex\_setprioceiling()

NAME	pthread_mutex_setpric (POSIX)	oceiling() – dynamic	ally set th	ne prioceiling attrib	ute of a m	nutex
SYNOPSIS	<pre>int pthread_mutex_se   (     pthread_mutex_t     int     int     )</pre>	etprioceiling *pMutex, prioceiling, *pOldPrioceiling	/* POSIX /* new g /* old g	K mutex priority ceiling priority ceiling	(out)	* / * / * /
DESCRIPTION	This function dynamica mutex was created with nothing.	ally sets the value of t h a protocol value of t	the prioce PTHREAE	eiling attribute of a D_PRIO_PROTECT, t	mutex. U his functi	nless the ion does

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VxWorks Kernel API Reference, 6.6 pthread\_mutex\_trylock()

RETURNS	On success zero; on failure a non-zero error code:
	EINVAL
	EPERM
	S_objLib_OBJ_ID_ERROR
	S_semLib_NOT_ISR_CALLABLE
ERRNO	N/A
SEE ALSO	<pre>pthreadLib, pthread_mutex_getprioceiling( ), pthread_mutexattr_getprioceiling( ), pthread_mutexattr_setprioceiling( )</pre>

# pthread\_mutex\_trylock( )

NAME	<b>pthread_mutex_trylock( )</b> – lock mutex if it is available (POSIX)		
SYNOPSIS	<pre>int pthread_mutex_trylock   (    pthread_mutex_t *pMutex /* POSIX mutex */   )</pre>		
DESCRIPTION	This routine locks the mutex specified by <i>pMutex</i> . If the mutex is currently unlocked, it becomes locked and owned by the calling thread. In this case <b>pthread_mutex_trylock()</b> returns immediately.		
	If the mutex is already locked by another thread, <b>pthread_mutex_trylock()</b> returns immediately with the error code <b>EBUSY</b> .		
RETURNS	On success zero; on failure a non-zero error code: EINVAL EBUSY		
ERRNO	N/A		
SEE ALSO	<pre>pthreadLib, semLib, semMLib, pthread_mutex_init(), pthread_mutex_lock(), pthread_mutex_trylock(), pthread_mutex_unlock(), pthread_mutexattr_init(), semTake()</pre>		

#### pthread\_mutex\_unlock( ) pthread\_mutex\_unlock() - unlock a mutex (POSIX) NAME SYNOPSIS int pthread\_mutex\_unlock pthread\_mutex\_t \*pMutex ) DESCRIPTION This routine unlocks the mutex specified by *pMutex*. If the calling thread is not the current owner of the mutex, pthread\_mutex\_unlock() returns with the error code EPERM. RETURNS On success zero; on failure a non-zero error code: EINVAL EPERM S\_objLib\_OBJ\_ID\_ERROR S\_semLib\_NOT\_ISR\_CALLABLE N/A ERRNO pthreadLib, semLib, semMLib, pthread\_mutex\_init(), pthread\_mutex\_lock(), SEE ALSO pthread\_mutex\_trylock(), pthread\_mutex\_unlock(), pthread\_mutexattr\_init(), semGive()

#### pthread\_mutexattr\_destroy( )

NAME	<pre>pthread_mutexattr_destroy() - destroy mutex attributes object (POSIX)</pre>
SYNOPSIS	<pre>int pthread_mutexattr_destroy   (    pthread_mutexattr_t *pAttr /* mutex attributes */   )</pre>
DESCRIPTION	This routine destroys a mutex attribute object. The mutex attribute object must not be reused until it is reinitialized.
RETURNS	On success zero; on failure the EINVAL error code.
ERRNO	N/A

VxWorks Kernel API Reference, 6.6 pthread\_mutexattr\_getprioceiling()

SEE ALSO pthreadLib, pthread\_mutexattr\_getprioceiling(), pthread\_mutexattr\_getprotocol(), pthread\_mutexattr\_init(), pthread\_mutexattr\_setprioceiling(), pthread\_mutexattr\_setprotocol(), pthread\_mutex\_init()

#### pthread\_mutexattr\_getprioceiling()

NAME	<b>pthread_mutexattr_getprioceiling()</b> – get the current value of the prioceiling attribute in a mutex attributes object (POSIX)		
SYNOPSIS	<pre>int pthread_mutexattr_getprioceiling   (     pthread_mutexattr_t *pAttr, /* mutex attributes */     int *pPrioceiling /* current priority ceiling (out) */   )</pre>		
DESCRIPTION	This function gets the current value of the prioceiling attribute in a mutex attributes obje Unless the value of the protocol attribute is <b>PTHREAD_PRIO_PROTECT</b> , this value is ignored.		
RETURNS	On success zero; on failure the EINVAL error code.		
ERRNO	N/A		
SEE ALSO	pthreadLib, pthread_mutexattr_destroy( ), pthread_mutexattr_getprotocol( ), pthread_mutexattr_init( ), pthread_mutexattr_setprioceiling( ), pthread_mutexattr_setprotocol( ), pthread_mutex_init( )		

#### pthread\_mutexattr\_getprotocol( )

NAME	<b>pthread_mutexattr_getprotocol()</b> – get value of protocol in mutex attributes object (POSL		
SYNOPSIS	<pre>int pthread_mutexattr_getprotocol    (</pre>		
	pthread_mutexattr_t *pAttr, /* mutex attributes	*/	
	)	^/	
DESCRIPTION	This function gets the current value of the protocol attribute in a mutex attributes object		
RETURNS	On success zero; on failure the <b>EINVAL</b> error code.		
ERRNO
 N/A

 SEE ALSO
 pthreadLib, pthread\_mutexattr\_destroy(), pthread\_mutexattr\_getprioceiling(), pthread\_mutexattr\_init(), pthread\_mutexattr\_setprioceiling(), pthread\_mutexattr\_setprotocol(), pthread\_mutex\_init()

#### pthread\_mutexattr\_init( )

NAME	<pre>pthread_mutexattr_init( ) - initialize mutex attributes object (POSIX)</pre>
SYNOPSIS	<pre>int pthread_mutexattr_init   (    pthread_mutexattr_t *pAttr /* mutex attributes */ )</pre>
DESCRIPTION	This routine initializes the mutex attribute object $pAttr$ and fills it with default values for the attributes:
	Mutex Protocol PTHREAD_PRIO_INHERIT - the priority of the owner thread is temporarily raised if a higher priority thread is blocked on the mutex.
	Mutex Priority Ceiling 0 - lowest priority.
RETURNS	On success zero; on failure the EINVAL error code.
ERRNO	N/A
SEE ALSO	pthreadLib, pthread_mutexattr_destroy( ), pthread_mutexattr_getprioceiling( ), pthread_mutexattr_getprotocol( ), pthread_mutexattr_setprioceiling( ), pthread_mutexattr_setprotocol( ), pthread_mutex_init( )

#### pthread\_mutexattr\_setprioceiling()

NAME	<pre>pthread_mutexattr_setprioceiling() - set prioceiling attribute in mutex attributes object (POSIX)</pre>
SYNOPSIS	<pre>int pthread_mutexattr_setprioceiling   (</pre>
	pthread_mutexattr_t *pAttr, /* mutex_attributes */

VxWorks Kernel API Reference, 6.6 pthread\_mutexattr\_setprotocol()

	<pre>int prioceiling /* new priority ceiling */ )</pre>
DESCRIPTION	This function sets the value of the prioceiling attribute in a mutex attributes object. Unless the protocol attribute is set to <b>PTHREAD_PRIO_PROTECT</b> , this attribute is ignored.
RETURNS	On success zero; on failure the EINVAL error code.
ERRNO	N/A
SEE ALSO	<pre>pthreadLib, pthread_mutexattr_destroy(), pthread_mutexattr_getprioceiling(), pthread_mutexattr_getprotocol(), pthread_mutexattr_init(), pthread_mutexattr_setprotocol(), pthread_mutex_init()</pre>

## pthread\_mutexattr\_setprotocol( )

NAME	<pre>pthread_mutexattr_setprotocol() - set protocol attribute in mutex attribute object (POSIX)</pre>
SYNOPSIS	<pre>int pthread_mutexattr_setprotocol   (    pthread_mutexattr_t *pAttr, /* mutex attributes */    int protocol /* new protocol */  )</pre>
DESCRIPTION	This function selects the locking protocol to be used when a mutex is created using this attributes object. The protocol to be selected is either <b>PTHREAD_PRIO_INHERIT</b> or <b>PTHREAD_PRIO_PROTECT</b> .
RETURNS	On success zero; on failure a non-zero error code: EINVAL ENOTSUP
ERRNO	N/A
SEE ALSO	<pre>pthreadLib, pthread_mutexattr_destroy(), pthread_mutexattr_getprioceiling(), pthread_mutexattr_getprotocol(), pthread_mutexattr_init(), pthread_mutexattr_setprioceiling(), pthread_mutex_init()</pre>

NAME	<pre>pthread_once() – dynamic package initialization (POSIX)</pre>
SYNOPSIS	<pre>int pthread_once   (    pthread_once_t * pOnceControl, /* once control location */    void (*initFunc)(void) /* function to call */   )</pre>
DESCRIPTION	This routine provides a mechanism to ensure that one, and only one call to a user specified initialization function will occur. This allows all threads in a system to attempt initialization of some feature they need to use, without any need for the application to explicitly prevent multiple calls.
	When a thread makes a call to <b>pthread_once()</b> , the first thread to call it with the specified control variable, <i>pOnceControl</i> , will result in a call to <i>initFunc</i> , but subsequent calls will not. The <i>pOnceControl</i> parameter determines whether the associated initialization routine has been called. The <i>initFunc</i> function is complete when <b>pthread_once()</b> returns.
	The function <b>pthread_once()</b> is not a cancellation point; however, if the function <i>initFunc</i> is a cancellation point, and the thread is canceled while executing it, the effect on <i>pOnceControl</i> is the same as if <b>pthread_once()</b> had never been called.
CAVEAT	If the initialization function does not return then all threads calling <b>pthread_once()</b> with the same control variable will stay blocked as well. It is therefore imperative that the initialization function always returns.
WARNING	If <i>pOnceControl</i> has automatic storage duration or is not initialized to the value <b>PTHREAD_ONCE_INIT</b> , the behavior of <b>pthread_once()</b> is undefined.
	The constant <b>PTHREAD_ONCE_INIT</b> is defined in the <b>pthread.h</b> header file.
RETURNS	On success zero; on failure the EINVAL error code.
ERRNO	None
SEE ALSO	pthreadLib

VxWorks Kernel API Reference, 6.6 pthread\_self()

## pthread\_self()

NAME	<b>pthread_self( )</b> – get the calling thread's ID (POSIX)
SYNOPSIS	pthread_t pthread_self (void)
DESCRIPTION	This function returns the calling thread's ID. If the caller is a native VxWorks task it will be given a POSIX thread persona.
RETURNS	Calling thread's ID.
ERRNO	Not Available
SEE ALSO	pthreadLib

## pthread\_setcancelstate()

NAME	<b>pthread_setcancelstate()</b> – set cancellation state for calling thread (POSIX)	
SYNOPSIS	<pre>int pthread_setcancelstate   (    int state, /* new state */    int *oldstate /* old state (out) */   )</pre>	
DESCRIPTION	This routine sets the cancellation state for the calling thread to <i>state</i> , and, if <i>oldstate</i> is not <b>NULL</b> , returns the old state in the location pointed to by <i>oldstate</i> .	
	The state can be one of the following:	
	PTHREAD_CANCEL_ENABLE Enable thread cancellation.	
	<b>PTHREAD_CANCEL_DISABLE</b> Disable thread cancellation (i.e. thread cancellation requests are ignored).	
RETURNS	On success zero; on failure the EINVAL error code.	
ERRNO	N/A	
SEE ALSO	pthreadLib, pthread_cancel(), pthread_setcanceltype(), pthread_testcancel()	

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NAME	<b>pthread_setcanceltype()</b> – set cancellation type for calling thread (POSIX)
SYNOPSIS	<pre>int pthread_setcanceltype   (    int type, /* new type */    int *oldtype /* old type (out) */   )</pre>
DESCRIPTION	This routine sets the cancellation type for the calling thread to <i>type</i> . If <i>oldtype</i> is not <b>NULL</b> , then the old cancellation type is stored in the location pointed to by <i>oldtype</i> .
	Possible values for <i>type</i> are:
	PTHREAD_CANCEL_ASYNCHRONOUS Any cancellation request received by this thread will be acted upon as soon as it is received.
	<b>PTHREAD_CANCEL_DEFERRED</b> Cancellation requests received by this thread will be deferred until the next cancellation point is reached.
RETURNS	On success zero; on failure the EINVAL error code.
ERRNO	N/A
SEE ALSO	pthreadLib, pthread_cancel(), pthread_setcancelstate(), pthread_testcancel()

# pthread\_setschedparam()

NAME	<b>pthread_setschedparam()</b> – dynamically set schedparam attribute for a thread (POSIX)		
SYNOPSIS	int pthread_setschedparam (		
	pthread_t thread, /* thread */		
	int policy, /* new policy */		
	const struct sched_param *pParam /* new parameters */ )		
DESCRIPTION	This routine will set the scheduling parameters ( <i>pParam</i> ) and policy ( <i>policy</i> ) for the thread specified by <i>thread</i> .		
	In VxWorks the scheduling policy is global and not set on a per-thread basis; if the selected policy is one of <b>SCHED FIFO</b> or <b>SCHED RR</b> and this does not match the current VxWorks		

VxWorks Kernel API Reference, 6.6 pthread\_setspecific()

	scheduling policy then this function will return an error ( <b>EPERM</b> ). If the <i>policy</i> parameter is set to <b>SCHED_OTHER</b> , which always matches the active scheduling policy, only the thread's priority will be changed.
RETURNS	On success zero; on failure one of the following non-zero error codes: <b>EPERM</b> , <b>ESRCH</b> (invalid task ID), <b>EINVAL</b> (scheduling priority is outside valid range)
ERRNO	N/A
SEE ALSO	pthreadLib, pthread_attr_getschedparam( ), pthread_attr_getschedpolicy( ), pthread_attr_setschedparam( ), pthread_attr_setschedpolicy( ), pthread_getschedparam( ), sched_getparam( ), sched_setparam( )

#### pthread\_setspecific( )

NAME	<b>pthread_setspecific()</b> – set thread specific data (POSIX)
SYNOPSIS	<pre>int pthread_setspecific   (     pthread_key_t key, /* thread specific data key */     const void *value /* new value */   )</pre>
DESCRIPTION	Sets the value of the thread specific data associated with key to value for the calling thread.
RETURNS	On success zero; on failure a non-zero error code: EINVAL ENOMEM
ERRNO	N/A
SEE ALSO	pthreadLib, pthread_getspecific( ), pthread_key_create( ), pthread_key_delete( )

#### pthread\_sigmask()

NAME pthread\_sigmask() – change and/or examine calling thread's signal mask (POSIX)

SYNOPSIS int pthread\_sigmask

	( int how, /* method for changing set */ const sigset_t * set, /* new set of signals */ sigset_t * oset /* old set of signals */ )
DESCRIPTION	This routine changes the signal mask for the calling thread as described by the <i>how</i> and <i>set</i> arguments. If <i>oset</i> is not <b>NULL</b> , the previous signal mask is stored in the location pointed to by it.
	The value of <i>how</i> indicates the manner in which the set is changed and consists of one of the following defined in <b>signal.h</b> :
	<b>SIG_BLOCK</b> The resulting set is the union of the current set and the signal set pointed to by <i>set</i> .
	<b>SIG_UNBLOCK</b> The resulting set is the intersection of the current set and the complement of the signal set pointed to by <i>set</i> .
	<b>SIG_SETMASK</b> The resulting set is the signal set pointed to by <i>oset</i> .
RETURNS	On success zero; on failure a <b>EINVAL</b> error code is returned.
ERRNO	N/A
SEE ALSO	pthreadLib, kill(), pthread_kill(), sigprocmask(), sigaction(), sigsuspend(), sigwait()

#### pthread\_testcancel()

**NAME** pthread\_testcancel() – create a cancellation point in the calling thread (POSIX)

SYNOPSIS void pthread\_testcancel (void)

**DESCRIPTION** This routine creates a cancellation point in the calling thread. It has no effect if cancellation is disabled (i.e. the cancellation state has been set to **PTHREAD\_CANCEL\_DISABLE** using the **pthread\_setcancelstate()** function).

If cancellation is enabled, the cancellation type is **PTHREAD\_CANCEL\_DEFERRED** and a cancellation request has been received, then this routine will call **pthread\_exit()** with the exit status set to **PTHREAD\_CANCELED**. If any of these conditions is not met, then the routine does nothing.

RETURNS N/A

VxWorks Kernel API Reference, 6.6 ptyDevCreate()

ERRNO

N/A

SEE ALSO pthreadLib, pthread\_cancel(), pthread\_setcancelstate(), pthread\_setcanceltype()

#### ptyDevCreate()

NAME	<pre>ptyDevCreate( ) - create a pseudo terminal</pre>
SYNOPSIS	<pre>STATUS ptyDevCreate   (     char *name, /* name of pseudo terminal */     int rdBufSize, /* size of terminal read buffer */     int wrtBufSize /* size of write buffer */   )</pre>
DESCRIPTION	This routine creates a master and slave device which can then be opened by the master and slave processes. The master process simulates the "hardware" side of the driver, while the slave process is the application program that normally talks to a <i>tty</i> driver. Data written to the master device can then be read on the slave device, and vice versa.
RETURNS	<b>OK</b> , or <b>ERROR</b> if memory is insufficient.
ERRNO	S_ioLib_NO_DRIVER (ENXIO) The <b>ptyDrv</b> driver is not installed.
	S_iosLib_DUPICATE_DEVICE_NAME (EINVAL) The device name is already in use.
SEE ALSO	ptyDrv

#### ptyDevRemove()

NAME	<pre>ptyDevRemove() – destroy a pseudo terminal</pre>
SYNOPSIS	STATUS ptyDevRemove ( char * pName /* name of pseudo terminal to remove */ )
DESCRIPTION	This routine removes an existing master and slave device and releases all allocated memory. It will close any open files using either device.

RETURNS	OK, or ERROR if terminal not found
ERRNO	S_ioLib_NO_DRIVER (ENXIO) The <b>ptyDrv</b> is not installed.
SEE ALSO	ptyDrv

## ptyDrv()

SEE ALSO	ptyDrv
ERRNO	N/A
RETURNS	OK, or ERROR if the master or slave devices cannot be installed.
DESCRIPTION	This routine initializes the pseudo-terminal driver. It must be called before any other routine in this module.
SYNOPSIS	STATUS ptyDrv (void)
NAME	ptyDrv() – initialize the pseudo-terminal driver

## putenv()

NAME	<b>putenv( )</b> – set an environment variable
SYNOPSIS	STATUS putenv ( char *pEnvString /* string to add to env */ )
DESCRIPTION	This routine sets an environment variable to a value by altering an existing variable or creating a new one. The parameter points to a string of the form "variableName=value". Unlike the UNIX implementation, the string passed as a parameter is copied to a private buffer.
RETURNS	<b>OK</b> , or <b>ERROR</b> if space cannot be malloc'd.

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VxWorks Kernel API Reference, 6.6 pwd()

ERRNOS S\_memLib\_NOT\_ENOUGH\_MEMORY There is no free block large enough to satisfy the allocation request.

SEE ALSO envLib, envLibInit(), getenv()

#### pwd()

NAME	<b>pwd()</b> – print the current default directory
SYNOPSIS	void pwd (void)
DESCRIPTION	This command displays the current working device/directory.
NOTE	This is a target resident function, which manipulates the target I/O system. It must be preceded with the @ letter if executed from the Host Shell (windsh), which has a built-in command of the same name that operates on the Host's I/O system.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	<b>usrFsLib</b> , <b>cd()</b> , the VxWorks programmer guides, the, <i>VxWorks Command-Line Tools User's Guide</i> .

## quiccEngineDrvCtrlShow()

NAME	quiccEngineDrvCtrlShow() – place holder just prints out control structure ptr
SYNOPSIS	<pre>int quiccEngineDrvCtrlShow   (     VXB_DEVICE_ID pInst    )</pre>
DESCRIPTION	none
RETURNS	N/A
ERRNO	
SEE ALSO	quiccEngineUtils

	quiccEngineRegister()
NAME	quiccEngineRegister() – register quiccEngine driver
SYNOPSIS	void quiccEngineRegister(void)
DESCRIPTION	This routine registers the quiccEngine driver and device recognition data with the vxBus subsystem.
NOTE	This routine is called early during system initialization, and *MUST NOT* make calls to OS facilities such as memory allocation and I/O.
RETURNS	N/A
ERRNO	
SEE ALSO	quiccEngineUtils

## r0()

NAME	<b>r0( )</b> – return the contents of register <b>r0</b> (also <b>r1 - r14</b> ) (ARM)
SYNOPSIS	<pre>int r0   (    int taskId /* task ID, 0 means default task */ )</pre>
DESCRIPTION	This command extracts the contents of register <b>r0</b> from the TCB of a specified task. If <i>taskId</i> is omitted or zero, the last task referenced is assumed. Similar routines are provided for registers ( <b>r1 - r14</b> ): <b>r1() - r14()</b> .
RETURNS	The contents of register <b>r0</b> (or the requested register).
ERRNO	Not Available
SEE ALSO	dbgArchLib, VxWorks Programmer's Guide: Debugging

## r0()

NAME	<b>r0( )</b> – return the contents of general register <b>r0</b> (also <b>r1</b> - $r15'$ ) (SH)
SYNOPSIS	<pre>int r0   (    int taskId /* task ID, 0 means default task */ )</pre>
DESCRIPTION	This command extracts the contents of register <b>r0</b> from the TCB of a specified task. If <i>taskId</i> is omitted or zero, the last task referenced is assumed. Similar routines are provided for all general registers ( <b>r1 - r15</b> ): <b>r1() - r15()</b> .
RETURNS	The contents of register r0 (or the requested register).
ERRNO	Not Available
SEE ALSO	dbgArchLib, VxWorks Programmer's Guide: Debugging

## raise()

NAME	<b>raise( )</b> – send a signal to the caller's task
SYNOPSIS	int raise ( int signo /* signal to send to caller's task */ )
DESCRIPTION	This routine sends the signal <i>signo</i> to the task invoking the call.
RETURNS	<b>OK</b> (0), or <b>ERROR</b> (-1) if the signal number or task ID is invalid.
ERRNO	EINVAL
SEE ALSO	sigLib, taskRaise()

#### ramDevCreate()

NAME	<pre>ramDevCreate() – create a RAM disk device</pre>
SYNOPSIS	<pre>BLK_DEV* ramDevCreate   (     char *ramAddr, /* where it is in memory (0 = malloc) */     int bytesPerBlk, /* number of bytes per block */     int blksPerTrack, /* number of blocks per track */     int nBlocks, /* number of blocks on this device */     int blkOffset /* no. of blks to skip at start of device */     )</pre>
DESCRIPTION	This routine creates a RAM disk device.
	Memory for the RAM disk can be pre-allocated separately; if so, the <i>ramAddr</i> parameter should be the address of the pre-allocated device memory. Or, memory can be automatically allocated with <b>malloc()</b> by setting <i>ramAddr</i> to zero.
	The <i>bytesPerBlk</i> parameter specifies the size of each logical block on the RAM disk. If <i>bytesPerBlk</i> is zero, 512 is used.
	The <i>blksPerTrack</i> parameter specifies the number of blocks on each logical track of the RAM disk. If <i>blksPerTrack</i> is zero, the count of blocks per track is set to <i>nBlocks</i> (i.e., the disk is defined as having only one track).
	The <i>nBlocks</i> parameter specifies the size of the disk, in blocks. If <i>nBlocks</i> is zero, a default size is used. The default is calculated using a total disk size of either 51,200 bytes or one-half of the size of the largest memory area available, whichever is less. This default disk size is then divided by <i>bytesPerBlk</i> to determine the number of blocks.
	The <i>blkOffset</i> parameter specifies an offset, in blocks, from the start of the device to be used when writing or reading the RAM disk. This offset is added to the block numbers passed by the file system during disk accesses. (VxWorks file systems always use block numbers beginning at zero for the start of a device.) This offset value is typically useful only if a specific address is given for <i>ramAddr</i> . Normally, <i>blkOffset</i> is 0.
FILE SYSTEMS	Once the device has been created, it must be associated with a name and a file system (dosFs, hrfs, or rawFs). This is accomplished in a two step process. The <b>ramDevCreate()</b> call returns a pointer to a block device structure ( <b>BLK_DEV</b> ). This structure contains fields that describe the physical properties of a disk device and specify the addresses of routines within the <b>ramDrv</b> driver. The <b>BLK_DEV</b> structure address should be passed to an XBD wrapper via <b>xbdBlkDevCreate()</b> along with the name of the device. XBDs are the new and preferred method for interfacing with file systems.
	After the XBD wrapper is created, the file system framework will attempt to identify the type of file system instantiated on the device. If it can not be identified, then it is instantiated with rawFs.

VxWorks Kernel API Reference, 6.6 ramDiskDevCreate()

The desired file system (dosFs or hrfs) can be instantiated on the ram drive using either **dosFsVolFormat()**, **dosfsDiskFormat()**, **hrfsFormat()**, or **hrfsDiskFormat()**. The ram drive to be formatted is identified by the name of the device given in the XBD wrapper.

**EXAMPLE** In the following example, a 208-Kbyte RAM disk is created with automatically allocated memory, 512-byte blocks, 32 blocks per track, and no block offset. The device is then initialized for use with dosFs and assigned the name "/ramDrv":

BLK\_DEV \*pBlkDev;

pBlkDev = ramDevCreate (NULL, 512, 32, 416, 0); xbdBlkDevCreate (pBlkDev, "/ramDrv"); dosFsVolFormat ("/ramDrv:0", DOS\_OPT\_BLANK, NULL);

The names used in **xbdBlkDevCreate()** and **dosFsVolFormat()** are slightly different on purpose. The ":0" is appended to "/**ramDrv**" by **xbdBlkDevCreate()** and represents the whole (unpartitioned) disk.

If the RAM disk memory already contains a disk image created elsewhere, the first argument to **ramDevCreate()** should be the address in memory, and the formatting parameters -- *bytesPerBlk*, *blksPerTrack*, *nBlocks*, and *blkOffset* -- must be identical to those used when the image was created. For example:

pBlkDev = ramDevCreate (0xc0000, 512, 32, 416, 0);

In this case, the file system does not have to be explicitly created as the file system framework will probe the ram drive to determine the type of file system previously instantiated on it. The detected file system will be automatically re-instantiated on the device. This procedure is useful if a RAM disk is to be created at the same address used in a previous boot of VxWorks. The contents of the RAM disk will then be preserved.

If no known file system was detected, the ram drive will default to rawFs.

- **RETURNS** A pointer to a block device structure (**BLK\_DEV**) or **NULL** if memory cannot be allocated for the device structure or for the RAM disk.
- ERRNO N/A.
- SEE ALSO ramDrv, xbdBlkDevCreate(), dosFsVolFormat(), hrfsFormat()

#### ramDiskDevCreate()

#### NAME ramDiskDevCreate() – Initialize a RAM Disk device

SYNOPSIS CBIO\_DEV\_ID ramDiskDevCreate

	( char *pRamAddr, /* where it is in memory (0 = malloc) int bytesPerBlk, /* number of bytes per block int blksPerTrack, /* number of blocks per track int nBlocks, /* number of blocks on this device int blkOffset /* no. of blks to skip at start of devi )	*/ */ */ ce */
DESCRIPTION	This function creates a compact RAM-Disk device that can be directly utiliz without the intermediate disk cache. It can be used for non-volatile RAM a RAM disks.	zed by <b>dosFsLib</b> , s well as volatile
	The RAM size is specified in terms of total number of blocks in the device a in bytes. The minimal block size is 32 bytes. If <i>pRamAddr</i> is <b>NULL</b> , space w from the default memory pool.	nd the block size ill be allocated
CAVEAT	When used with NV-RAM, this module can not eliminate mid-block write which may cause file system corruption not existent in common disk driv	e interruption, es.
RETURNS	a CBIO handle that can be directly used by <b>dosFsDevCreate( )</b> or <b>NULL</b> if amount of RAM is not available.	the requested
ERRNO	Not Available	
SEE ALSO	ramDiskCbio, dosFsDevCreate().	

#### ramDrv()

**NAME** ramDrv() – prepare a RAM disk driver for use (optional)

SYNOPSIS STATUS ramDrv (void)

**DESCRIPTION** This routine performs no real function, except to provide compatibility with earlier versions of **ramDrv** and to parallel the initialization function found in true disk device drivers. It also is used in **usrConfig.c** to link in the RAM disk driver when building VxWorks. It is automatically called when VxWorks is configured with the **INCLUDE\_RAMDRV** component.

**RETURNS** OK, always.

ERRNO N/A.

SEE ALSO ramDrv

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#### rawFsDevInit()

NAME	<pre>rawFsDevInit() - associate a block device with raw volume functions</pre>
SYNOPSIS	<pre>RAW_VOL_DESC *rawFsDevInit   (     char * pVolName, /* volume name to be used with iosDevAdd */     device_t xbd /* XBD device */   )</pre>
DESCRIPTION	This routine takes a block device created by a device driver and defines it as a raw file system volume. As a result, when high-level I/O operations, such as <b>open()</b> and <b>write()</b> are performed, on the device, the calls will be routed through <b>rawFsLib</b> .
	This routine associates <i>pVolName</i> with a device and installs it in the VxWorks I/O System's device table. The driver number used when the device is added to the table is that which was assigned to the raw library during <b>rawFsInit()</b> . (The driver number is kept in the global variable <b>rawFsDrvNum</b> .)
	The <i>xbd</i> is a device_t referring to an XBD device which represents the backing media for this rawFs. The XBD device will not be accessed until an I/O operation occurs on the file system.
RETURNS	A pointer to the volume descriptor (RAW_VOL_DESC), or NULL if there is an error.
ERRNO	Not Available
SEE ALSO	rawFsLib

#### rawFsInit()

NAME	<b>rawFsInit( )</b> – prepare to use the raw volume library
SYNOPSIS	STATUS rawFsInit ( int maxFiles /* max no. of simultaneously open files */ )
DESCRIPTION	This routine initializes the raw volume library. It must be called exactly once, before any other routine in the library. The argument specifies the number of file descriptors that may be open at once. This routine allocates and sets up the necessary memory structures and initializes semaphores.
	This routine also installs raw volume library routines in the VxWorks I/O system driver table. The driver number assigned to <b>rawFsLib</b> is placed in the global variable

**rawFsDrvNum**. This number will later be associated with system file descriptors opened to rawFs devices.

RETURNS OK or ERROR.

ERRNO Not Available

SEE ALSO rawFsLib

#### rawPerfDemo()

**NAME** rawPerfDemo() – entry point for the VxWorks/SMP raw performance demo

SYNOPSIS STATUS rawPerfDemo (void)

**DESCRIPTION** This routine is the entry point for the VxWorks/SMP raw performance demo. It is typically called from the target shell.

This function will create N worker tasks; N = number of CPUs currently enabled in the system. Aggregate raw performance figures are obtained as described in the module description. The aggregate raw performance data is plotted in real-time on an ASCII character graph.

See the module description for more information.

**RETURNS ERROR** if failed to create worker tasks, otherwise OK is returned.

ERRNO S\_memLib\_NOT\_ENOUGH\_MEMORY Out of memory for creation of worker tasks

SEE ALSO rawPerfDemo

#### read()

	size_t maxbytes /* max no. of bytes to read into buffer */ )
DESCRIPTION	This routine reads a number of bytes (less than or equal to <i>maxbytes</i> ) from a specified file descriptor and places them in <i>buffer</i> . It calls the device driver to do the work.
RETURNS	The number of bytes read (between 1 and <i>maxbytes</i> , 0 if end of file), or <b>ERROR</b> if the file descriptor does not exist, the driver does not have a read routines, or the driver returns <b>ERROR</b> . If the driver does not have a read routine, errno is set to <b>ENOTSUP</b> .
ERRNO	EBADF Bad file descriptor number.
	ENOTSUP Device driver does not support the read command.
	ENXIO Device and its driver are removed. <b>close( )</b> should be called to release this file descriptor.
	Other Other errors reported by device driver.
SEE ALSO	ioLib

#### readdir()

NAME	readdir() – read one entry from a directory (POSIX)
SYNOPSIS	struct dirent *readdir ( DIR *pDir /* pointer to directory descriptor */ )
DESCRIPTION	This routine obtains directory entry data for the next file from an open directory. The <i>pDir</i> parameter is the pointer to a directory descriptor (DIR) which was returned by a previous <b>opendir()</b> .
	This routine returns a pointer to a <b>dirent</b> structure which contains the name of the next file. Empty directory entries and MS-DOS volume label entries are not reported. The name of the file (or subdirectory) described by the directory entry is returned in the <b>d_name</b> field of the <b>dirent</b> structure. The name is a single null-terminated string.
	The returned <b>dirent</b> pointer will be <b>NULL</b> , if it is at the end of the directory or if an error occurred. Because there are two conditions which might cause <b>NULL</b> to be returned, the task's error number ( <b>errno</b> ) must be used to determine if there was an actual error. Before

	calling <b>readdir()</b> , set <b>errno</b> to <b>OK</b> . If a <b>NULL</b> pointer is returned, check the new value of <b>errno</b> . If <b>errno</b> is still <b>OK</b> , the end of the directory was reached; if not, <b>errno</b> contains the error code for an actual error which occurred.
RETURNS	A pointer to a <b>dirent</b> structure, or <b>NULL</b> if there is an end-of-directory marker or error from the IO system.
ERRNO	EBADF Bad file descriptor number.
	S_ioLib_UNKNOWN_REQUEST (ENOSYS) Device driver does not support the ioctl command.
	Other Other errors reported by device driver.
SEE ALSO	dirLib, opendir(), readdir_r(), closedir(), rewinddir(), ls()

## readdir\_r()

NAME	<b>readdir_r()</b> – read one entry from a directory (POSIX)
SYNOPSIS	<pre>int readdir_r (</pre>
DESCRIPTION	This routine obtains directory entry data for the next file from an open directory. The <i>pDir</i> parameter is the pointer to a directory descriptor (DIR) which was returned by a previous <b>opendir()</b> .
	The caller must allocate storage pointed to by <i>entry</i> to be large enough for a dirent structure with an array of char d_name member containing at least <b>NAME_MAX</b> .
	On successful return, the pointer returned at <i>result</i> will be the same value as the argument <i>entry</i> . Upon reaching the end of the directory stream, this pointer will have the value <b>NULL</b> .
RETURNS	zero if successful or an error number to indicate failure.
ERRNO	EBADF Bad file descriptor number.
	S_ioLib_UNKNOWN_REQUEST (ENOSYS) Device driver does not support the ioctl command.

Other Other errors reported by device driver.

SEE ALSO dirLib, opendir(), readdir(), closedir(), rewinddir(), ls()

#### realloc()

NAME realloc() – reallocate a block of memory (ANSI) SYNOPSIS void \* realloc ( void \* pBlock, /\* block to reallocate \*/ size\_t newSize /\* new block size \*/ ) This routine changes the size of a specified block of memory and returns a pointer to the DESCRIPTION new block of memory. The contents that fit inside the new size (or old size if smaller) remain unchanged. The memory alignment of the new block is not guaranteed to be the same as the original block. When the INCLUDE\_MEM\_MGR\_FULL component is included this funcition changes to an optimized implementation that attempts to resize the existing block. If *pBlock* is **NULL**, this call is equivalent to **malloc()**. If *newSize* is set to zero and *pBlock* points to a valid allocated block, this call is equivalent to free(). RETURNS A pointer to the new block of memory, **NULL** if the call fails or if *newSize* is equal to zero. ERRNO Possible errnos generated by this routine include: S\_memLib\_NOT\_ENOUGH\_MEMORY There is no free block large enough to satisfy the allocation request. memPartLib, memPartRealloc(), American National Standard for Information Systems -, SEE ALSO Programming Language - C, ANSI X3.159-1989: General Utilities (stdlib.h)

NAME	<b>reboot()</b> – reset network devices and transfer control to boot ROMs
SYNOPSIS	<pre>void reboot   (    int startType /* how the boot ROMS will reboot */   )</pre>
DESCRIPTION	This routine returns control to the boot ROMs after calling a series of preliminary shutdown routines that have been added via <b>rebootHookAdd()</b> , including routines to reset all network devices. After calling the shutdown routines, interrupts are locked, all caches are cleared, and control is transferred to the boot ROMs.
	The bit values for <i>startType</i> are defined in <b>sysLib.h</b> :
	<b>BOOT_NORMAL</b> (0x00) causes the system to go through the countdown sequence and try to reboot VxWorks automatically. Memory is not cleared.
	<b>BOOT_NO_AUTOBOOT</b> (0x01) causes the system to display the VxWorks boot prompt and wait for user input to the boot ROM monitor. Memory is not cleared.
	<b>BOOT_CLEAR</b> (0x02) the same as <b>BOOT_NORMAL</b> , except that memory is cleared.
	<b>BOOT_QUICK_AUTOBOOT</b> (0x04) the same as <b>BOOT_NORMAL</b> , except the countdown is shorter.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	<b>rebootLib</b> , <b>sysToMonitor()</b> , <b>rebootHookAdd()</b> , <b>windsh</b> , the VxWorks programmer guides, and the IDE and host tools guides.

## rebootHookAdd()

NAME	<b>rebootHookAdd()</b> – add a routine to be called at reboot
SYNOPSIS	STATUS rebootHookAdd ( FUNCPTR rebootHook /* routine to be called at reboot */ )

DESCRIPTION	This routine adds the specified routine to a list of routines to be called when VxWorks is rebooted. The specified routine should be declared as follows:
	<pre>void rebootHook   (     int startType /* startType is passed to all hooks */ )</pre>
	Reboot hooks will be called in the order they were added, when the <b>reboot( )</b> function is executed.
	Reboot hooks must follow similar restrictions as with interrupt service routines (ISRs). Reboot hooks must not invoke kernel service routines that may block. For example, calling <b>free()</b> or <b>semTake()</b> while the semaphore is not available, would cause the caller to block. Blocking calls within the reboot hooks may causes the reboot process to <b>reschedule()</b> or hang, potentially leaving the system in an undefined state.
RETURNS	OK, or ERROR if memory is insufficient.
ERRNO	Not Available
SEE ALSO	rebootLib, reboot()

#### reld()

NAME	<b>reld()</b> – reload an object module (shell command)
SYNOPSIS	<pre>MODULE_ID reld   (    void * nameOrId, /* name or ID of the object module file */    int options /* options used for unloading */   )</pre>
DESCRIPTION	This routine unloads a specified object module from the system, and then calls <b>loadModule()</b> to load a new copy of the same name.
	If the file was originally loaded using a complete pathname, then <b>reld()</b> will use the complete name to locate the file. If the file was originally loaded using a partial pathname, then the current working directory must be changed to the working directory in use at the time of the original load.
	Valid values for the options parameter are the same as those allowed for the function <b>unld()</b> .
	This routine is a <b>shell command</b> . That is, it is designed to be used only in the shell, and not in code running on the target. In future releases, calling <b>reld()</b> directly from code may not be supported.

RETURNS	A module ID (type <b>MODULE_ID</b> ), or <b>NULL</b> .
---------	--

ERRNO Not Available

**SEE ALSO** usrLib, loadLib, unld(), ld(), the VxWorks programmer guides.

#### rename()

NAME	<b>rename( )</b> – change the name of a file	
SYNOPSIS	<pre>int rename   (    const char *oldname, /* name of file to rename */    const char *newname /* name with which to rename file */   )</pre>	
DESCRIPTION	This routine changes the name of a file from <i>oldfile</i> to <i>newfile</i> .	
NOTE	Only certain devices support <b>rename()</b> . To confirm that your device supports it, consult the respective <b>xxDrv</b> or xxFs listings to verify that ioctl FIORENAME exists. For example, dosFs, HRFS and NFS support <b>rename()</b> , but <b>netDrv</b> does not.	
RETURNS	OK, or ERROR if the file could not be opened or renamed.	
ERRNO	ENOENT Either oldname or newname is an empty string. ELOOP Circular symbolic link, too many links. EMFILE Maximum number of files already open. S_iosLib_DEVICE_NOT_FOUND (ENODEV) No uplid douise name found in path	
	No valid device name found in path. ENOSYS Device driver does not support the symlink ioctl command. others Other errors reported by device driver.	
SEE ALSO	fsPxLib	

#### repeat()

NAME	<b>repeat( )</b> – spawn a task to call a function repeatedly
SYNOPSIS	<pre>int repeat     (     FAST int n, /* no. of times to call func (0=forever) */     FAST FUNCPTR func, /* function to call repeatedly */     int arg1, /* first of eight args to pass to func */     int arg2,     int arg3,     int arg4,     int arg5,     int arg6,     int arg7,     int arg8     )</pre>
DESCRIPTION	This command spawns a task that calls a specified function $n$ times, with up to eight of its arguments. If $n$ is 0, the routine is called endlessly, or until the spawned task is deleted.
NOTE	The task is spawned using <b>sp()</b> . See the description of <b>sp()</b> for details about priority, options, stack size, and task ID.
RETURNS	A task ID, or <b>ERROR</b> if the task cannot be spawned.
ERRNO	sp() errnos.
SEE ALSO	usrLib, repeatRun(), sp(), the VxWorks programmer guides.

## repeatRun()

NAME	<pre>repeatRun() – call</pre>	a functio	on repeatedly
SYNOPSIS	void repeatRun ( FAST int FAST FUNCPTR int int int int int int int int	n, func, arg1, arg2, arg3, arg4, arg5, arg6, arg7,	<pre>/* no. of times to call func (0=forever) */ /* function to call repeatedly */ /* first of eight args to pass to func */</pre>

	int arg8 )
DESCRIPTION	This command calls a specified function $n$ times, with up to eight of its arguments. If $n$ is 0, the routine is called endlessly.
	Normally, this routine is called only by <b>repeat()</b> , which spawns it as a task.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	usrLib, repeat(), the VxWorks programmer guides.

## rewinddir()

NAME	<b>rewinddir()</b> – reset position to the start of a directory (POSIX)
SYNOPSIS	<pre>void rewinddir   (     DIR *pDir /* pointer to directory descriptor */ )</pre>
DESCRIPTION	This routine resets the position pointer in a directory descriptor (DIR). The <i>pDir</i> parameter is the directory descriptor pointer that was returned by <b>opendir()</b> .
	As a result, the next <b>readdir()</b> will cause the current directory data to be read in again, as if an <b>opendir()</b> had just been performed. Any changes in the directory that have occurred since the initial <b>opendir()</b> will now be visible. The first entry in the directory will be returned by the next <b>readdir()</b> .
RETURNS	N/A
ERRNO	N/A.
SEE ALSO	dirLib, opendir( ), readdir( ), closedir( )

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#### rindex()

NAME	<b>rindex( )</b> – find the last occurrence of a character in a string
SYNOPSIS	<pre>char *rindex   (   FAST const char * s, /* string in which to find character */   int c /* character to find in string */   )</pre>
DESCRIPTION	This routine finds the last occurrence of character $c$ in string $s$ .
RETURNS	A pointer to <i>c</i> , or <b>NULL</b> if <i>c</i> is not found.
ERRNO	N/A
SEE ALSO	bLib

#### rm( )

NAME	rm() – remove a file	
SYNOPSIS	STATUS rm ( const char * fileName /* name of file to remove */ )	
DESCRIPTION	This command is provided for UNIX similarity. It simply calls <b>remove()</b> .	
RETURNS	OK, or ERROR if the file cannot be removed.	
ERRNO	Not Available	
SEE ALSO	usrFsLib, remove(), the VxWorks programmer guides.	

## rmdir()

NAME	rmdir() – remove a directory
SYNOPSIS	STATUS rmdir ( const char * dirName /* name of directory to remove */ )
DESCRIPTION	This command removes an existing directory from a hierarchical file system. The <i>dirName</i> string specifies the name of the directory to be removed, and may be either a full or relative pathname.
	This call is supported by the VxWorks NFS and dosFs file systems.
RETURNS	OK, or ERROR if the directory cannot be removed.
ERRNO	Not Available
SEE ALSO	usrFsLib, mkdir(), the VxWorks programmer guides.

## rngBufGet()

NAME	<b>rngBufGet( )</b> – get characters from a ring buffer
SYNOPSIS	<pre>int rngBufGet   (   FAST RING_ID rngId, /* ring buffer to get data from */   char *buffer, /* pointer to buffer to receive data */   int maxbytes /* maximum number of bytes to get */  )</pre>
DESCRIPTION	This routine copies bytes from the ring buffer <i>rngId</i> into <i>buffer</i> . It copies as many bytes as are available in the ring, up to <i>maxbytes</i> . The bytes copied will be removed from the ring.
RETURNS	The number of bytes actually received from the ring buffer; it may be zero if the ring buffer is empty at the time of the call.
ERRNO	N/A.
SEE ALSO	rngLib

## rngBufPut()

NAME	<b>rngBufPut()</b> – put bytes into a ring buffer
SYNOPSIS	<pre>int rngBufPut   (   FAST RING_ID rngId, /* ring buffer to put data into */   char *buffer, /* buffer to get data from */   int nbytes /* number of bytes to try to put */   )</pre>
DESCRIPTION	This routine puts bytes from <i>buffer</i> into ring buffer <i>ringId</i> . The specified number of bytes will be put into the ring, up to the number of bytes available in the ring.
RETURNS	The number of bytes actually put into the ring buffer; it may be less than number requested, even zero, if there is insufficient room in the ring buffer at the time of the call.
ERRNO	N/A.
SEE ALSO	rngLib

## rngCreate()

NAME	<b>rngCreate()</b> – create an empty ring buffer
SYNOPSIS	RING_ID rngCreate ( int nbytes /* number of bytes in ring buffer */ )
DESCRIPTION	This routine creates a ring buffer of size <i>nbytes</i> , and initializes it. Memory for the buffer is allocated from the system memory partition.
RETURNS	The ID of the ring buffer, or NULL if memory cannot be allocated.
ERRNO	N/A.
SEE ALSO	rngLib

## rngDelete()

NAME	<b>rngDelete( )</b> – delete a ring buffer
SYNOPSIS	<pre>void rngDelete   (   FAST RING_ID ringId /* ring buffer to delete */ )</pre>
DESCRIPTION	This routine deletes a specified ring buffer. Any data currently in the buffer will be lost.
RETURNS	N/A
ERRNO	N/A.
SEE ALSO	rngLib

## rngFlush()

NAME	<b>rngFlush( )</b> – make a ring buffer empty
SYNOPSIS	<pre>void rngFlush   (    FAST RING_ID ringId /* ring buffer to initialize */ )</pre>
DESCRIPTION	This routine initializes a specified ring buffer to be empty. Any data currently in the buffer will be lost.
RETURNS	N/A
ERRNO	N/A.
SEE ALSO	rngLib

#### rngFreeBytes()

NAME	<b>rngFreeBytes()</b> – determine the number of free bytes in a ring buffer
SYNOPSIS	int rngFreeBytes ( FAST RING_ID ringId /* ring buffer to examine */ )
DESCRIPTION	This routine determines the number of bytes currently unused in a specified ring buffer.
RETURNS	The number of unused bytes in the ring buffer.
ERRNO	N/A.
SEE ALSO	rngLib

## rngIsEmpty()

NAME	<b>rngIsEmpty( )</b> – test if a ring buffer is empty
SYNOPSIS	BOOL rngIsEmpty ( RING_ID ringId /* ring buffer to test */ )
DESCRIPTION	This routine determines if a specified ring buffer is empty.
RETURNS	TRUE if empty, FALSE if not.
ERRNO	N/A.
SEE ALSO	rngLib

#### rngIsFull()

NAME	<pre>rngIsFull() - test if a ring buffer is full (no more room)</pre>

SYNOPSIS BOOL rngIsFull

```
(
FAST RING_ID ringId /* ring buffer to test */
)
```

This routine determines if a specified ring buffer is completely full. DESCRIPTION

RETURNS TRUE if full, FALSE if not.

ERRNO N/A.

rngLib SEE ALSO

#### rngMoveAhead()

NAME	<b>rngMoveAhead()</b> – advance a ring pointer by <i>n</i> bytes
SYNOPSIS	<pre>void rngMoveAhead  (  FAST RING_ID ringId, /* ring buffer to be advanced */ FAST int n /* number of bytes ahead to move input pointer */ )</pre>
DESCRIPTION	This routine advances the ring buffer input pointer by <i>n</i> bytes. This makes <i>n</i> bytes available in the ring buffer, after having been written ahead in the ring buffer with <b>rngPutAhead()</b> .
RETURNS	N/A
ERRNO	N/A.
SEE ALSO	rngLib

#### rngNBytes()

NAME	<b>rngNBytes( )</b> – determine the number of bytes in a ring buffer
SYNOPSIS	<pre>int rngNBytes   (    FAST RING_ID ringId /* ring buffer to be enumerated */ )</pre>
DESCRIPTION	This routine determines the number of bytes currently in a specified ring buffer.

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VxWorks Kernel API Reference, 6.6 rngPutAhead()

**RETURNS** The number of bytes filled in the ring buffer.

ERRNO N/A.

SEE ALSO rngLib

#### rngPutAhead()

NAME

rngPutAhead() – put a byte ahead in a ring buffer without moving ring pointers

SYNOPSIS

OPSIS	void rngPutAhead	
	(	
	FAST RING_ID ringl	d, /* ring buffer to put byte in */
	char byte,	/* byte to be put in ring */
	int offse	t /* offset beyond next input byte where to put byte
	*/	
	)	

**DESCRIPTION** This routine writes a byte into the ring, but does not move the ring buffer pointers. Thus the byte will not yet be available to **rngBufGet()** calls. The byte is written *offset* bytes ahead of the next input location in the ring. Thus, an offset of 0 puts the byte in the same position as would **RNG\_ELEM\_PUT** would put a byte, except that the input pointer is not updated.

Bytes written ahead in the ring buffer with this routine can be made available all at once by subsequently moving the ring buffer pointers with the routine **rngMoveAhead()**.

Before calling **rngPutAhead()**, the caller must verify that at least *offset* + 1 bytes are available in the ring buffer.

RETURNS N/A

ERRNO N/A.

SEE ALSO rngLib

#### romStart()

NAME	romStart() – generic ROM initialization
------	---

SYNOPSIS void romStart

```
      (FAST int startType /* start type */)

      DESCRIPTION
      This is the first C code executed after reset.

      This routine is called by the assembly start-up code in romInit(). It clears memory, copies ROM to RAM, and possibly invokes the uncompressor. It then jumps to the entry point of the uncompressed object code.

      RETURNS
      N/A

      ERRNO
      SEE ALSO

      bootInit
      bootInit
```

#### round()

NAME	round() – round a number to the nearest integer
SYNOPSIS	<pre>double round   (     double x /* value to round */    )</pre>
DESCRIPTION	This routine rounds a double-precision value $x$ to the nearest integral value.
RETURNS	The double-precision representation of $x$ rounded to the nearest integral value.
ERRNO	Not Available
SEE ALSO	mathALib

#### roundf()

NAME roundf() - round a number to the nearest integer SYNOPSIS float roundf ( float x /\* argument \*/ ) VxWorks Kernel API Reference, 6.6 rtgRegister()

SEE ALSO	mathALib
ERRNO	Not Available
RETURNS	The single-precision representation of <i>x</i> rounded to the nearest integral value.
DESCRIPTION	This routine rounds a single-precision value $x$ to the nearest integral value.

## rtgRegister()

NAME	<b>rtgRegister()</b> – register with the VxBus subsystem
SYNOPSIS	void rtgRegister(void)
DESCRIPTION	This routine registers the RealTek driver with VxBus as a child of the PCI bus type.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	rtl8169VxbEnd

## rtlRegister()

NAME	<b>rtlRegister()</b> – register with the VxBus subsystem
SYNOPSIS	void rtlRegister(void)
DESCRIPTION	This routine registers the RealTek driver with VxBus as a child of the PCI bus type.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	rtl8139VxbEnd

# rtpDelete()

NAME	<b>rtpDelete( )</b> – terminates a real time process (RTP)
SYNOPSIS	<pre>STATUS rtpDelete   (   RTP_ID rtpId, /* ID of the RTP to be deleted */   int options, /* options for deletion */   int delStatus /* exit status */   )</pre>
DESCRIPTION	This routine terminates an RTP from the system. The termination of the RTP removes all objects (including tasks) owned by the RTP from the system via the resource reclamation facility. Any mapped memory in the RTP will be unmapped and memory will be freed back to the system. Memory allocated to the RTP for the executable file will also be freed back to the system. Note that public objects still in use by other users in the system will be inherited by the kernel, and will not be reclaimed at this point.
	Shared data regions created and mapped in the RTP will be unmapped. If the RTP is the last client reference to the shared data, the termination of the RTP will trigger a deletion of the shared data region. For more information on shared data regions, please refer to <b>sdLib</b> .
	Shared libraries created and mapped by an RTP will also be unmapped. As with shared data regions, if the terminating RTP is the last reference to the shared library, the termination of the RTP will trigger the deletion of the shared library from the system.
	There is currently no user-available values for the <i>options</i> parameter. This parameter should always be set to zero.
	<i>delStatus</i> is the exit status of the RTP that can be extracted by the parent of this RTP through the <b>wait()</b> system call.
	Users may install RTP delete hooks to be called before the termination of the victim RTP's resources. These hooks must reside within the kernel. For hooks that access the RTP's user space, the hooks are responsible for switching into the context of the RTP to perform the access. The delete hooks are called after tasks within the RTP are terminated. To add a delete hook routine, use the following:
	<b>rtpDeleteHookAdd()</b> Add a hook routine to be called during the termination of an RTP.
WARNING	<b>rtpDelete( )</b> may not be called from an Interrupt Service Routine (ISR).
SMP CONSIDERATIO	DNS
	This API is spinlock and intCpuLock restricted.
RETURNS	OK, or ERROR if RTP can not be deleted

VxWorks Kernel API Reference, 6.6 rtpDeleteHookAdd()

ERRNOS	Possible errnos generated are:
	S_rtpLib_INVALID_RTP_ID The specified rtpId parameter is not a valid ID or is in the deleted state already.
	S_objLib_OBJ_ID_ERROR The specified rtpId is invalid for the <b>rtpDelete( )</b> operation.
	<b>S_objLib_OBJ_DELETED</b> The RTP object has already been terminated and the RTP no longer exists.
SEE ALSO	rtpLib, rtpSpawn( ), sdDelete( )

## rtpDeleteHookAdd()

NAME	rtpDeleteHookAdd() – add a routine to be called when RTPs are deleted
SYNOPSIS	STATUS rtpDeleteHookAdd ( RTP_DELETE_HOOK hook, /* routine to be called when RTPs are deleted */ BOOL addToHead /* add routine to head of list? */ )
DESCRIPTION	This routine adds a specified routine to a list of routines that will be called whenever an RTP is deleted. The hook routine should have the following prototype: void deleteHook ( const RTP_ID rtpId /* RTP ID of the RTP about to be deleted */ const int exitCode /* exit code or delete status for the RTP */ ) The second parameter <i>addToHead</i> specifies the order in which the hook is added to the table. If FALSE, the hook is appended to the list of hooks already installed. If addToHead is TRUE, the new hook is added to the head of the list (in other words, it will be the first hook to execute). It is typical for delete hooks to be added at the head of the table. Doing so enables the most recently added hooks to execute first. This is useful when hook routines have dependencies among themselves, and the order in which the hooks execute is important. RTP delete hooks are called from rtpDelete() before any deletion is done. Delete hooks are not expected to return anything (return values if any are not checked).
RETURNS	<b>OK</b> , or <b>ERROR</b> if the table of RTP delete routines is full.
ERRNO	N/A.
# rtpDeleteHookDelete()

NAME	rtpDeleteHookDelete() – delete a previously added RTP delete hook routine				
SYNOPSIS	STATUS rtpDeleteHookDelete ( RTP_DELETE_HOOK hook /* routine to be deleted from list */ )				
DESCRIPTION	This routine removes a specified routine from the list of routines to be called at each RTP delete.				
RETURNS	<b>OK</b> , or <b>ERROR</b> if the routine is not in the table of RTP delete routines.				
ERRNO	N/A.				
SEE ALSO	rtpHookLib, rtpDeleteHookAdd()				

# rtpHelp()

SEE ALSO	usrRtpLib				
ERRNO	N/A				
RETURNS	N/A				
DESCRIPTION	This routine prints the synopsis of the shell commands applying to Real-Time Processes.				
SYNOPSIS	void rtpHelp (void)				
NAME	<b>rtpHelp()</b> – print a synopsis of RTP-related shell commands				

	rtpHookShow()					
NAME	<pre>rtpHookShow() – display all installed RTP hooks</pre>					
SYNOPSIS	void rtpHookShow (void)					
DESCRIPTION	This routine displays the contents of all three RTP hook tables, the pre-create, post-create, and delete hook tables.					
EXAMPLE	The following example shows a hypothetical set of RTP hook table contents					
	-> rtpHookShow					
	RTP Pre-Create Hook Table:					
	hookfunc1					
	hookfunc2 hookfunc3					
	RTP Post-Create Hook Table:					
	hookfunc4					
	hookfunc5 hookfunc6					
	RTP Init-Complete Hook Table:					
	hookfunc7					
	hookfunc8 hookfunc9					
	RTP Delete Hook Table:					
	hookfunc12					
	hookfuncll hookfuncl0					
	value = 1 = 0x1					
	->					
RETURNS	N/A					
ERRNOS	N/A					
SEE ALSO	rtpShow, rtpShow(), hookShow(), syscallHookShow()					

### rtpInfoGet()

NAME rtpInfoGet() –

SYNOPSIS

rtpInfoGet( ) – Get specific information on an RTP

```
STATUS rtpInfoGet
    (
        RTP_ID rtpId, /* RTP ID to get info */
        RTP_DESC * rtpStruct /* Location to store RTP info */
    )
```

**DESCRIPTION** This routine obtains information about an RTP and stores the information in the specified RTP descriptor *rtpStruct*. The information stored in the descriptor, for the most part, is a snapshot copy of the information about the RTP object. The descriptor must have been allocated before calling this function, and the memory for it must come from the the calling task's memory space. To allocate the memory for the descriptor from the calling task's memory space, either use **malloc()** within the calling task or declare the structure as an automatic variable in the calling task, placing it on the calling task's stack.

The rtpStruct structure looks like the following:

typedef struct

```
{
         pathName[VX_RTP_NAME_LENGTH+1]; // pointer to executable path
char
int
         status; // the state of the RTP
                        // option bits, e.g. debug, symtable
UINT32
         options;
         options;
entrAddr;
                        // entry point of ELF file
void *
         initTaskId;
int
                        // the initial task ID
         taskCnt;
                        // number of tasks in the RTP
TNT32
RTP_ID parentId;
                        // RTP ID of the parent
} RTP_DESC;
```

The length of the pathName field is limited to VX\_RTP\_NAME\_LENGTH (255). The errno S\_rtpLib\_RTP\_NAME\_MAX will be set if the RTP's executable pathName exceeds this limit.

The initTaskId will be 0 if the initial task of the RTP was deleted at the time this routine is called. The initTaskId will also be 0 if the caller is a task in a different RTP, as initial task is private to an RTP.

The IDs of the initTaskId and parentId are the WIND kernel IDs, when the routine is used in the kernel. **objShow()** may be called directly to display information on these objects. If the routine is invoked within an RTP, these IDs are opaque IDs local to the RTP. To display information on these IDs, use the **objHandleShow()** routine on the opaque IDs.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted.

RETURNS OK or ERROR

VxWorks Kernel API Reference, 6.6 rtplnitCompleteHookAdd()

ERRNOS	S_objLib_OBJ_ID_ERROR Invalid RTP ID or null <i>rtpStruct</i> parameter.		
	S_rtpLib_RTP_NAME_LENGTH_EXCEEDED The actual path name of the RTP exceeds the VX_RTP_NAME_LENGTH (255).		
SEE ALSO	rtpUtilLib, rtpLib, rtpShow()		

### rtpInitCompleteHookAdd()

rtpInitCompleteHookAdd() – Add routine to be called after RTP init-complete. NAME SYNOPSIS STATUS rtpInitCompleteHookAdd RTP\_INIT\_COMPLETE\_HOOK hook, /\* routine to be called after rtp init \*/ addToHead /\* add routine to head of list? \*/ BOOL ) DESCRIPTION This routine adds a specified routine to a list of routines that will be called after an RTP is created and fully initialized (i.e. just before its initial task starts running). Init-Complete hook routines should have the following prototype: void rtpInitCompleteHook ( const RTP\_ID rtpId /\* ID of the created RTP \*/ ) The second parameter *addToHead* specifies the order in which the hook is added to the table. If FALSE, the hook is appended to the list of hooks already installed. If addToHead is TRUE, the new hook is added to the head of the list (i.e. it will be the first hook to execute). RTP Init-Complete hooks are called as a result of calling **rtpSpawn()**, and are called after the RTP's creation and initialization are fully complete. Init-Complete hooks are used as a notification point for debuggers and other tools that may want to access the newly created RTP's memory space. Init-Complete hooks are a notification point, and can't return an ERROR. Therefore, they should typically not create any objects or perform any actions that are critical to the existence of the RTP. Consider using a post-create hook if this functionality is required. **OK**, or **ERROR** if the table of RTP Init-Complete routines is full. RETURNS N/A. ERRNO SEE ALSO rtpHookLib, rtpInitCompleteHookDelete()

2

# rtpInitCompleteHookDelete()

NAME	rtpInitCompleteHookDelete() – delete a previously added RTP init-complete hook				
SYNOPSIS	STATUS rtpInitCompleteHookDelete ( RTP_INIT_COMPLETE_HOOK hook /* routine to be deleted from list */ )				
DESCRIPTION	This routine removes a specified hook routine from the list of RTP Init-Complete hook routines.				
RETURNS	OK on success, or ERROR if the hook routine was not found.				
ERRNO	S_hookLib_HOOK_NOT_FOUND				
SEE ALSO	rtpHookLib, rtpInitCompleteHookAdd( )				

# rtpKill()

NAME	<b>rtpKill( )</b> – send a signal to a RTP
SYNOPSIS	<pre>int rtpKill   (    RTP_ID rtpId,    int signo   )</pre>
DESCRIPTION	This routine sends a kill signal <i>signo</i> to the RTP specified by <i>pRtpld</i> . If signo equals -1, it will use objEach to go through the list of RTP's and send a kill signal to each RTP.
RETURNS	<b>OK</b> (0), or <b>ERROR</b> (-1) if the task ID or signal number is invalid.
ERRNO	EINVAL
SEE ALSO	rtpSigLib

	rtpLkAddr()					
NAME	rtpLkAddr() – list symbols in an RTP whose values are near a specified value					
SYNOPSIS	<pre>void rtpLkAddr  (  UINT addr, /* address around which to look */  RTP_ID rtpId /* RTP to look for symbols in */ )</pre>					
DESCRIPTION	This command lists the symbols in the RTP's symbol table that are near a specified value. The symbols that are displayed include:					
	- symbols whose values are immediately less than the specified value					
	- symbols with the specified value					
	- succeeding symbols, until at least 12 symbols have been displayed					
	The results of this command depend partly on what options were used to spawn the RTP. The default options cause only global symbols from the ELF file used to spawn the RTP to be available for this routine.					
RETURNS	N/A					
ERRNO	N/A					
SEE ALSO	<b>usrRtpLib</b> , <b>rtpLib</b> , <b>symLib</b> , <b>symEach( )</b> , the VxWorks programmer guides, the IDE and host tools guides.					

# rtpLkup()

NAME	<pre>rtpLkup() – list symbols from an RTP's symbol table</pre>
SYNOPSIS	<pre>void rtpLkup  (     char * substr, /* substring to match */     RTP_ID rtpId /* Id of rtp to search for symbol in */ )</pre>
DESCRIPTION	This command lists all symbols in the RTP's symbol table whose names contain the

By default, **rtpLkup()** displays 22 symbols at a time. This can be changed by modifying the global variable **symLkupPgSz**. If this variable is set to 0, **rtpLkup()** displays all the symbols without interruption.

RETURNS	N/A
ERRNO	N/A
SEE ALSO	<b>usrRtpLib</b> , <b>usrLib</b> , <b>rtpLib</b> , <b>symLib</b> , <b>symEach()</b> , the VxWorks programmer guides, the IDE and host tools programer guides.

### rtpMemShow()

**NAME** rtpMemShow() – display memory context information for real time process

```
SYNOPSIS STATUS rtpMemShow
(
char * rtpNameOrId, /* RTP name or ID */
int level /* detail level */
)
```

**DESCRIPTION** This routine displays memory context information for a real time process. This routine takes two parameters, *rtpNameOrId* and *level*. The first parameter can either be an RTP ID or an RTP name string. The second parameter is the level of detail to display the information for the RTPs. If the *level* is 0, then it displays a summary memory context information for the specified RTP. If the *level* is 1, then **rtpMemShow()** displays detailed memory context information. If the *level* is 2, then **rtpMemShow()** also displays POSIX mapped file and mapped object information.

#### SUMMARY INFORMATION EXAMPLE

The following example shows the summary output for all RTPs in the system.

#### DETAILED INFORMATION EXAMPLE

The following example shows the detailed output for an RTP.

#### C-interpreter shell:

-> rtpMemShow 1635018888, 1

Memory Information for 0x61746888 RTP (name = "< in/tmMmanFdLib.vxe"):

Virtual Memory Context:

Virtual	Memory (	Context	ID:	0x606ad2	290
Private	Virtual	Memory	Allocated:	0x6a000	bytes
Private	Virtual	Memory	Mapped:	0x66000	bytes

Private Mappings:

VIRTUAL ADDR	BLOCK LENGTH	PHYSICAL ADDR	PROT (S/U)	CACHE	SPECIAL
0x63000000	0x00012000	0x6174a000	R-X / R-X	CB-//-	
0x63012000	0x00010000	0x6175c000	RWX / RWX	CB-//-	
0x63022000	0x00002000	***unmapped***			
0x63024000	0x00010000	0x6176c000	RWX / RWX	CB-//-	
0x63034000	0x00002000	***unmapped***			
0x63036000	0x00032000	0x6177c000	RWX / RWX	CB-//-	
0x63068000	0x00002000	0x617ae000	R-X / R-X	CB-//-	

Shared Data Mappings:

VIRTUAL ADDR	BLOCK LENGTH	PHYSICAL ADDR	PROT (S/U)	CACHE	SPECIAL
0x6306c000	0x00002000	0x617b8000	RWX / RWX	CB-//-	
Shared Object	(POSIX) Mappi:	ngs:			
VIRTUAL ADDR	BLOCK LENGTH	PHYSICAL ADDR	PROT (S/U)	CACHE	SPECIAL
0x6306a000	0x00002000	0x617b0000	R-X / R-X	CB-//-	

value =  $0 = 0 \times 0$ 

For further information on the fields from the Memory Context section, see **vxContextShow()**.

The private mappings of an RTP consist of the memory mapped objects (mapped with **MAP\_PRIVATE**), the RTP's code segments (text/data/bss), the RTP's heap, and the stacks of the tasks running in the RTP.

The Shared Data Mappings consists of Shared Data regions opened with the **sdLib** API; this also includes shared library text segments.

The Shared Object Mappings consist of mappings obtained with **mmap()** using the **MAP\_SHARED** flag.

#### MAPPED OBJECT INFORMATION EXAMPLE

The following example shows POSIX memory mapped files information.

#### C-interpreter shell:

-> rtpMemShow 1635018888, 2

[removed detailed memory context information, same as above]

Memory Mapped Objects (POSIX):

ADDRESS	LENGTH	PROT	FLAGS	OFFSET	OBJECT
0x63036000	0x00010000	RW-	PRIVATE	N/A	***anonymous***
0x63046000	0x00010000	RW-	PRIVATE	N/A	***anonymous***
0x63056000	0x00012000	RW-	PRIVATE	N/A	***anonymous***
0x63068000	0x00002000	R	PRIVATE	0x0000000000000000000	/pxFs/mmapFd1
0x6306a000	0x00002000	R	SHARED	0x0000000000000000000	/pxFs/mmapFd2
value = 0 =	= 0x0				

Note that the PROT value shown in the Memory Mapped Objects section is the parameter passed to **mmap()** (i.e. the bit values of **PROT\_READ**, **PROT\_WRITE**, **PROT\_EXE** passed via the *prot* parameter). The PROT value displayed in the Memory Context section shows the actual protection that resulted after setting the corresponding MMU protection attributes. Depending on the constraints of the processor architecture and system configuration these may or may not be the same.

#### **COMMAND INTERPRETER**

For the command-interpreter shell, use the **rtp meminfo** command.

**RETURNS** OK if success, ERROR otherwise

**ERRNOS** Possible errnos generated by this function include:

#### S\_objLib\_OBJ\_ID\_ERROR

An incorrect RTP ID was provided.

#### S\_objLib\_ACCESS\_DENIED

Unable to get exclusive access to the RTP or the RTP list.

**SEE ALSO** rtpShow, rtpShow(), rtpLib, rtpUtilLib, vmContextShow(), the VxWorks programmer guides.

### rtpPostCreateHookAdd()

**rtpPostCreateHookAdd()** – add a routine to be called just after RTP creation. NAME SYNOPSIS STATUS rtpPostCreateHookAdd RTP\_POST\_CREATE\_HOOK hook, /\* routine to be called on rtp creation \*/ addToHead /\* add routine to head of list? \*/ BOOL ) DESCRIPTION This routine adds a specified routine to a list of routines that will be called just after an RTP and its initial task are created, but before the newly created RTP starts running. Upon creation, all routines specified by rtpPostCreateHookAdd() will be called in the context of the creating RTP, so any objects created by an post-create hook will be owned by the caller's RTP rather than the newly created RTP. To set the ownership of newly created objects to the new RTP, objOwnerSet() should be used. For example: objOwnerSet (createdObjId, rtpId) Post-create hook routines should have the following prototype: STATUS rtpPostCreateHook ( const RTP\_ID rtpId /\* ID of the created RTP \*/ ) The second parameter *addToHead* specifies the order in which the hook is added to the table. If FALSE, the hook is appended to the list of hooks already installed. If addToHead is TRUE, the new hook is added to the head of the list (i.e. it will be the first hook to execute). RTP Post-creation hooks are called from **rtpSpawn()**, and should return either **OK** or **ERROR**. If the return value from a post-create hook is anything other than **OK**, the created RTP and its initial task are deleted, and **rtpSpawn()** returns ERROR. Post-creation hooks can be used to perform additional application-specific resource allocation etc where the created RTP's details should be known. Should such allocations fail, users have the option of reversing RTP creation by returning ERROR. RETURNS **OK**, or **ERROR** if the table of RTP post-create routines is full. N/A. ERRNO rtpHookLib, rtpPostCreateHookDelete() SEE ALSO

rtı	bР	OS	tC	rea	te	Ho	ok	D	el	ete	(	)
								_			<b>١</b>	/

NAME	<b>rtpPostCreateHookDelete()</b> – delete a previously added RTP post-create hook.			
SYNOPSIS	STATUS rtpPostCreateHookDelete ( RTP_POST_CREATE_HOOK hook /* routine to be deleted from list */ )			
DESCRIPTION	This routine removes a specified hook routine from the list of RTP post-create hook routines.			
RETURNS	OK on success, or ERROR if the hook routine was not found.			
ERRNO	S_hookLib_HOOK_NOT_FOUND			
SEE ALSO	rtpHookLib, rtpPostCreateHookAdd( )			

### rtpPreCreateHookAdd()

rtpPreCreateHookAdd() – add a routine to be called before RTP creation. NAME SYNOPSIS STATUS rtpPreCreateHookAdd ( RTP PRE CREATE HOOK hook, /\* routine to be called on rtp creation \*/ BOOL addToHead /\* add routine to head of list? \*/ ) This routine adds a specified routine to a list of routines that will be called just before an DESCRIPTION RTP is created. The hook routine should have the following prototype: STATUS rtpPreCreateHook ( const char \* rtpFileName, /\* Null-terminated path to executable \*/ const char \* argv[], /\* pointer to NULL terminated argv array \*/ const char \* envp[], /\* pointer to NULL terminated of int priority, /\* priority of initial task \*/ /\* pointer to NULL terminated envp array \*/ uStackSize, /\* User space stack size for intial task \*/ int int options, /\* the options passed to RTP \*/ int taskOptions /\* Task options for RTP's initial task \*/ )

The second parameter *addToHead* specifies the order in which the hook is added to the table. If **FALSE**, the hook is appended to the list of hooks already installed. If addToHead is **TRUE**, the new hook is added to the head of the list (i.e. it will be the first hook to execute).

VxWorks Kernel API Reference, 6.6 rtpPreCreateHookDelete()

	RTP Pre-creation hooks are called from <b>rtpSpawn()</b> before RTP creation begins. Hooks should return either <b>OK</b> or <b>ERROR</b> . If the return value from any hook is anything other than <b>OK</b> , RTP creation does not proceed. Pre-creation hooks can be used to implement rudimentary authentication schemes by rejecting RTP spawn requests before any action is taken.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the table of RTP pre-create routines is full.
ERRNO	N/A.
SEE ALSO	rtpHookLib, rtpPreCreateHookDelete()

## rtpPreCreateHookDelete()

NAME	rtpPreCreateHookDelete() – delete a previously added RTP pre-create hook.			
SYNOPSIS	STATUS rtpPreCreateHookDelete ( RTP_PRE_CREATE_HOOK hook /* routine to be deleted from list */ )			
DESCRIPTION	This routine removes a specified hook routine from the list of RTP pre- create hook routines			
RETURNS	OK on success, or ERROR if the hook routine was not found.			
ERRNO	S_hookLib_HOOK_NOT_FOUND			
SEE ALSO	rtpHookLib, rtpPreCreateHookAdd()			

## rtpShlShow()

NAME	<b>rtpShlShow()</b> – Display shared library information for an RTP
SYNOPSIS	STATUS rtpShlShow ( RTP_ID rtpId /* RTP to display SHLs on */

	RTP. The <b>SHL_ID</b> information displayed by this routine can be used by <b>shlShow()</b> to display detail information for a specific SHL.							
	<ul> <li>The <i>options</i> field may be specified to not print the header for the SHLs displayed.</li> <li>SHL_SHOW_NO_HDR (0x0001) Do not display the header for this routine.</li> </ul>							
EXAMPLE	Below is an example display	<sup>7</sup> of a sha	red library.					
	-> rtpShlShow 0x8d1ea9c SHL NAME	ID	TEXT_ADDR	TEXT_SIZE	DATA_ADDR	DATA_SIZE		
	< tty/slDfw/libSo.so	1	0xff435000	0x574	0xff46d000	0x628		
RETURNS	<b>OK</b> , or <b>ERROR</b> if invalid RTF	P id						
ERRNOS	Possible errnos generated by	y this fun	ction include	e:				
<b>S_objLib_OBJ_ID_ERROR</b> An incorrect SHL ID was provided.								
	S_objLib_ACCESS_DENIE Unable to get exclusive	D access to	the SHL list					
SEE ALSO	shlShow, shlShow( ), rtpSh	ow( )						

This routine takes an RTP as parameter and displays SHL information of all SHLs for the

### rtpShow()

DESCRIPTION

 NAME
 rtpShow() – display information for real time proceses

 SYNOPSIS
 BOOL rtpShow

 (char \* rtpNameOrId, /\* RTP name or ID \*/
 int level /\* 0 = summary, 1 = detailed, 2 = all in details \*/
 )

 DESCRIPTION
 This routine displays information for a real time process. This routine takes two parameters,
 *rtpNameOrId* and *level*. The first parameter can either be an RTP ID or an RTP name string.
 The second parameter is the level of detail to display the information for the RTPs.
 Depending on the level and the RTP ID specified, the information displayed differs. If the

Depending on the level and the KTP ID specified, the information displayed differs. If the *level* is 0, then it displays the summary information for either the specified RTP or all RTPs in the system. If the *level* is 1, then **rtpShow()** displays the detailed information, for the specified RTP or the current task's home RTP (if RTP ID is **NULL**, and called from the

command shell). If *level* is 2, **rtpShow()** displays the detailed information for all RTPs in the system, regardless of the RTP ID you specify. Refer to the table for more information.

	RTP Name		Cmd Shell
Level	or ID	Meaning	Equivalent
0	0	Display summary information for all RTPs.	rtp
0	RTP	Display summary information for specified RTP.	rtp rtpId
1	RTP	Display detailed information for specified RTP.	rtp info rtpId
2	ANY	Display detailed information for all RTP.	rtp info

In summary mode, **rtpShow()** only displays the RTP name (including the path) up to a maximum of 20 characters long. If the name is more than 20 characters, it will be truncated to 20 characters for displaying purposes. Preceding the truncated name, a "<" will be displayed to indicate that the name is more than 20 characters long. To get a display of the full RTP name, display the RTP with the *level* set to 1.

#### SUMMARY INFORMATION EXAMPLE

The following example shows the summary output for all RTPs in the system. If a RTP ID (or name) is specified, only the information for that RTP will be displayed.

#### C-interpreter shell:

-> rtpShow

	NAME	ID	STATE	ENTRY	ADDR	OPTIONS	TASK	CNT
<	/apps/myApp.vxe	0x4a9450	STATE_NORMAL	0xa000	0148	0x11		1

value = 1 = 0x1

For the command-interpreter shell, use the command **rtp**.

The display contains the following fields:

Field	Meaning
NAME	The name of the RTP, using the executable filename.
ID	The numeric ID associated with the RTP.
STATUS	State of the RTP. Refer to the table below for more information.
ENTRY ADRS	The entry routine address of the application.
OPTIONS	Options specified for the RTP. Refer to the options below.
TASK CNT	Number of tasks in the RTP.

#### STATE\_CREATE

The RTP is currently in its create phase.

#### STATE\_CREATE+S

The RTP is in **RTP\_STATE\_CREATE** state and its status is in **RTP\_STATUS\_STOP**. This indicates that the RTP currently in the create phase but is suspended during this phase.

#### STATE\_CREATE+D

The RTP is in **RTP\_STATE\_CREATE** state and its status is in **RTP\_STATUS\_ELECTED\_DELETER**. This indicates that the RTP has encountered an error during its create phase and has started its deletion.

#### STATE\_NORMAL

The RTP has completed initializing and its currently running normally.

#### STATE\_NORMAL+S

The RTP is in **RTP\_STATE\_NORMAL** state and its status is in **RTP\_STATUS\_STOP**. This indicates that the RTP currently has all its task in **TASK\_STOP** state.

#### STATE\_NORMAL+D

The RTP is in **RTP\_STATE\_NORMAL** state and its status is in **RTP\_STATUS\_ELECTED\_DELETER**. This indicates that the RTP has initiated its delete phase.

#### STATE\_DELETE

The RTP is in the processing of terminating and cleaning up.

#### STATE\_DELETE+S

The RTP is in **RTP\_STATE\_DELETE** state and its status is in **RTP\_STATUS\_STOP**. This indicates that the RTP is in the delete phase but tasks in the RTP has been stopped.

#### STATE\_DELETE+D

The RTP is in RTP\_STATE\_DELETE state and its status is RTP\_STATUS\_ELECTED\_DELETER. This indicates that the RTP has initiated its delete phase. The deletion process can not be undone.

#### DETAILED INFORMATION EXAMPLE

The following example shows the detailed output for a single RTP (i.e. the level was specified as 1). If the level is specified as 2, the detailed information is displayed for all RTP and the user is prompted to press **return** or **Q** between each RTP.

#### C-interpreter shell:

-> rtpShow 0x4a9450, 1

Ν	IAME	ID	STATE	ENTRY ADDR	OPTIONS	TASK CNT
< /apps/	/myApp.vxe	0x4a9450	STATE_NORMAL	0xa0000148	0x11	1
Full Nam Options: rtpId->p rtpId->p Initial Symbol t	ne: /usi : ( DArgv ptr: DEnv ptr: Task ID: table:	c/apps/myApp 0x11) RT 0xa001ef8c 0xa001ef94 0x4ab020 0x47b6c0	D.VXE FP_GLOBAL_SYMB	OLS RTP_DEBU	G	
SEGMENT  text data	START ADDR 0xa0000080 0xa000d02c	SIZE  49068 1424				

# VxWorks Kernel API Reference, 6.6 rtpSigqueue()

For the command-interpreter shell, use the command **rtp info**.

The summary line contains the same fields as explained above. The additional information is explained in the following table:

	Field	Meaning				
	Full Name	The complete name for the RTP.				
	Options	Detailed breakdown of the options word (see <b>rtpCreate( )</b> ).				
	rtpId->pArgv ptr	The address in user space where the arguments are stored.				
	rtpId->pEnv ptr	Address in user space of the environment string.				
	SEGMENT Information The text and data information for the RTP execu					
	Initial Task ID The Initial task ID of the RTP.					
	Symbol table The ID of the symbol table holding this RTP's symbol					
	Shared Libraries	The summary list of shared libraries used by the RTP.				
RETURNS	TRUE if success, FALSE oth Possible errnos generated l	<b>FRUE</b> if success, <b>FALSE</b> otherwise Possible errnos generated by this function include:				
	S_objLib_OBJ_ID_ERRO An incorrect RTP ID w	<b>R</b> 7as provided.				
<b>S_objLib_ACCESS_DENIED</b> Unable to get exclusive access to the RTP or the RTP list.						
SEE ALSO	<b>rtpShow</b> , <b>rtpMemShow( )</b> , <b>rtpLib</b> , <b>rtpUtilLib</b> , <b>vmContextShow( )</b> , the VxWorks programmer guides.					

### rtpSigqueue()

NAME rtpSigqueue() – send a queued signal to a RTP

rtpId,

signo,

SYNOPSIS int rtpSigqueue ( RTP\_ID int

	const union sigval value )
DESCRIPTION	The function <b>sigqueue()</b> sends the queued signal specified by <i>signo</i> with the signal-parameter value specified by <i>value</i> to the process specified by <i>pRtpId</i> .
RETURNS	<b>OK</b> (0), or <b>ERROR</b> (-1) if the task ID or signal number is invalid, or if there are no queued-signal buffers available.
ERRNO	EINVAL EAGAIN
SEE ALSO	rtpSigLib

# rtpSp()

NAME	<b>rtpSp( )</b> – launch a RTP with default options.		
SYNOPSIS	<pre>int rtpSp   (     char * execAndArgs,    /* path to the executable file + arguments */     int initTaskPrio,    /* priority of RTP's initial task */     int userStackSize,    /* size of the initial task's user stack */     int launchOptions,    /* options to apply to the RTP at launch */     int launchTaskOptions    /* task option for the RTP's initial task */     )</pre>		
DESCRIPTION	This command is a short form of the underlying <b>rtpSpawn()</b> routine, convenient for launching Real-Time Process (RTP) from the shell.		
	The executable file used to launch the application as well as the arguments to be passed to the <b>main()</b> routine of the application are all contained in the routine's first argument: the string <i>execAndArgs</i> . For instance:		
	rtpSp "/folk/me/myVxWorksApp.vxe firstArgument secondArgument"		
	The space character is interpreted as the separator between each element of the string. Any application argument can be a sub-string of its own providing that it is surrounded with escaped double quote characters. For instance:		
	rtpSp "/folk/me/myVxWorksApp.vxe \"first argument\" \"second argument\""		
	Although it is possible to specify properties and options to apply to the RTP's initial task using the parameters <i>initTaskPrio</i> , <i>userStackSize</i> , <i>launchOptions</i> and <i>launchTaskOptions</i> these parameters may be left unspecified (i.e. left null). In this case, default values for these properties and options will be applied. These default values may be overriden by updating the values of the following global variables:		

2

SpStackSize: 65,535 - initial task's user stack size, in bytes. SpOptions: 0x1 ( <b>RTP_GLOBAL_SYMBOLS</b> ) - OR'ed options to apply to the RTP, see <b>rtpSpawn()</b> details. rou to want to spawn an RTP with no option bits set but do not want to modify the val the rtpSpOptions variable, then rtpSp should be called with <i>launchOptions</i> set to -1. The cial value -1 forces the default value in rtpSpOptions to be ignored. This results in th P to be started without any option set (the same as if the rtpSpOptions variable held a
SpOptions: 0x1 ( <b>RTP_GLOBAL_SYMBOLS</b> ) - OR'ed options to apply to the RTP, see <b>rtpSpawn()</b> details. rou to want to spawn an RTP with no option bits set but do not want to modify the val the rtpSpOptions variable, then rtpSp should be called with <i>launchOptions</i> set to -1. The cial value -1 forces the default value in rtpSpOptions to be ignored. This results in th P to be started without any option set (the same as if the rtpSpOptions variable held a
rou to want to spawn an RTP with no option bits set but do not want to modify the val the rtpSpOptions variable, then rtpSp should be called with <i>launchOptions</i> set to -1. The cial value -1 forces the default value in rtpSpOptions to be ignored. This results in th P to be started without any option set (the same as if the rtpSpOptions variable held a
ıl value).
SpTaskOptions: VX_FP_TASK - OR'ed options to apply to the initial task of the RTP, see also taskSpawn().
in the case of <i>launchOptions</i> , if <i>rtpSpTaskOptions</i> is set to -1, the default value in SpTaskOptions is ignored and the initial task of the RTP is spawned with no task opti s (i.e. no floating-point support etc).
ditionally, the delay to wait before terminating the execution of the command after a plication has been launched can be contolled using:
SpDelay: 100 - number of ticks to wait.
te that these global variables also apply to the command interpreter mode of the she
content of the string <i>execAndArgs</i> is modified by this routine when parsed.
RTP_ID on success, ERROR otherwise.
A
RtpLib, rtpLib, rtpSpawn( ), taskSpawn( )

# rtpSpawn()

NAME	rtpSpawn() –	spawns a new Rea	l Time Process	(RTP) in the s	ystem
------	--------------	------------------	----------------	----------------	-------

SYNOPSIS RTP\_ID rtpSpawn

CAVEAT

RETURNS

ERRNO

SEE ALSO

```
(
const char * rtpFileName, /* Null terminated path to executable */
const char * argv[], /* Pointer to NULL terminated argv array */
const char * envp[], /* Pointer to NULL terminated envp array */
int priority, /* Priority of initial task */
int uStackSize, /* User space stack size for initial task */
int options, /* The options passed to the RTP */
int taskOptions /* Task options for the RTP's initial task */
)
```

**DESCRIPTION** This routine creates and initializes a Real Time Process (RTP) in the system, with the specified file as the executable for the RTP.

Each RTP is named. The name is based on the specified executable filename, via the *rtpFileName* argument, loaded in the RTP. This executable file must reside in a filesystem. The filesystem may be external or media-less and bundled (ROMFS) into the VxWorks system.

The first element to the argv[] array, by convention, should be the filename path of the executable. **rtpSpawn()** does not automatically populate argv[0] to be the executable pathname; the user must set it. Not providing argv[0] with the executable pathname may cause unexpected results if dynamic shared libraries are involved. Below is an example:

```
char * argv[] = {"/usr/test.vxe", NULL};
rtpSpawn (argv[0], argv, NULL, 100, 0x10000, 0, 0);
```

An RTP is a container for resources of the RTP application. Resources that may be associated with an RTP are: tasks, heap memory, and objects. Memory allocated for an RTP is unique in the system. Memory allocated to an RTP are task stacks, heap memory to be used by the user level heap manager, and memory allocated for the text and data segments of the application.

RTPs provide symbol name isolation. An executable may be spawned more than once in the system and the execution of the applications will not interfer with each other.

Tasks in an RTP are scheduled as part of the global scheduling scheme in the system. RTPs are not schedulable entities; only tasks within the RTPs are schedulable. Thus, for an RTP to exist, tasks must exist in it.

The *envp* environment array may be used to pass specific RTP environment variable settings to the application. Environment variables, such as **LD\_LIBRARY\_PATH**, may be set for an RTP. To obtain environment information for an RTP, use the **getenv()** routine or the **extern char** \*\***environ** variable in the application. Other reserved environment variables can be used to pass information used by the RTP when it initializes:

#### HEAP\_INITIAL\_SIZE

Set the initial size of the RTP's heap to a value other than the default (0x10000).

#### HEAP\_MAX\_SIZE

Set the maximum size that the RTP's heap may grow to.

#### HEAP\_INCR\_SIZE

Set the growth increment when it should be different from the default (a virtual memory page size).

See the application-side **memLib** documentation for more details. Such variables can be used as follows:

```
char * argv[] = {"/usr/test.vxe", NULL};
char * envp[] = {"HEAP_INITIAL_SIZE=0x20000", "HEAP_MAX_SIZE=0x100000",
NULL);
rtpSpawn (argv[0], argv, envp, 100, 0x10000, 0, 0);
```

The creation and initialization of an RTP also creates the initial task of the RTP. This initial task initializes the VxWorks user level library, libc support or **taskLib** support, of the RTP. Three of **rtpSpawn()**'s parameters are dedicated to setting the initial task's priority, user-side stack and options:

- *priority*:this parameter sets the priority of the RTP's initial task and care should be taken in setting a priority appropriate for an application (i.e. do not leave this parameter set to zero as this would create an initial task of the highest priority in VxWorks, possibly disturbing the functioning the rest of the system. A value between 200 and 220 is usually adequate).
- *uStackSize*:this parameter sets the size of the initial task's user-side stack. If this parameter is left null this size is set to the default value (0x4000 bytes).
- taskOptions: this parameter allows to pass options to the initial task created with the RTP. The taskOptions parameter has exactly the same value and meaning as the options parameter passed to taskSpawn(). Some task options available for kernel tasks are prohibited for RTP tasks, and will be ignored if set. These are the VX\_SUPERVISOR\_MODE and VX\_UNBREAKABLE options. The initial task of every RTP is created with the VX\_DEALLOC\_STACK option.

Options may be passed to the **rtpSpawn()** API to specify the behavior of the RTP:

#### RTP\_GLOBAL\_SYMBOLS (0x01)

The global symbols of the executable file will be registered in the RTP's symbol table. This is required when debugging using the embedded debugging facility.

#### RTP\_ALL\_SYMBOLS (0x03)

Both the global and local symbols of the executable file will be registered in the RTP's symbol table. This can be helpful when debugging using the embedded debugging facility.

#### RTP\_DEBUG (0x10)

The execution of the RTP will be stopped at startup in order to enable debugging the application.

#### RTP\_BUFFER\_VAL\_OFF (0x20)

User buffer passed to system calls will not be validated for this RTP. This will reduce the system call overhead, to the detriment of security. This option should be used only once the application code was properly debugged.

#### RTP\_LOADED\_WAIT (0x40)

**rtpSpawn()** will not return until the RTP has been instantiated, all code loaded, the RTP's state is **RTP\_STATE\_NORMAL**, and execution is about to transfer to user mode.

#### RTP\_CPU\_AFFINITY\_NONE (0x80)

By default the RTP's initial task inherits the CPU affinity of the task that spawned the RTP. This option removes any CPU affinity that would have applied to the initial task (i.e. this task will migrate from one CPU to another). Applies to SMP only.

A set of hooks are provided for users to extend the capabilities of the **rtpSpawn()** routine. Hooks may be added prior to the creation of the RTP, after the creation of the RTP object and VM context, and also after the loading of the RTP executable file. To add hooks at each of these points, use the following routines:

#### rtpPreCreateHookAdd()

Add a hook to be called prior to the creation of the RTP. The hook is executed in the caller's context. This hook facility is useful for validations prior to RTP creation, such as restricting the creation of RTPs to certain RTPs. Note, when this hook fails, RTP delete hook is not called.

#### rtpPostCreateHookAdd()

Add a hook to be called after the creation of the RTP object and VM context. Hooks are executed in the context of the caller. Any objects created by the hooks are owned by the caller unless the owner is reset via the **objOwnerSet()** API. This hook should not attempt to delete the RTP.

The post create hook is useful for extending the initialization of the RTP object such as initializing any user defined structures or objects.

An error returned from the registered hook routines will invoke a termination of the RTP. The RTP state is **RTP\_STATE\_CREATE**. Please refer to the routine **rtpPostCreateHookAdd()** for more information.

#### rtpInitCompleteHookAdd()

Add a hook to be called after loading of the RTP executable file and before the RTP's initial task starts executing in user mode. The RTP state is **RTP\_STATE\_NORMAL**.

#### rtpDeleteHookAdd()

Add a hook to be called when an RTP is terminated. The hooks are called at the beginning of the **rtpDelete()** routine. The RTP state is **RTP\_STATE\_NORMAL** 

For more detailed information on RTP hooks, refer to rtpHookLib.

The default behavior when an RTP task encounters an error, such as an exception, is that the system will terminate the faulty RTP. However, for debugging purposes, the system may be configured to behave in a **lab** mode where an exception would not terminate the RTP. Instead the faulty task and RTP will be suspended for debugging. To turn on the lab mode refer to the **edrLib** documentation.

**WARNING rtpSpawn()** may not be called from an Interrupt Service Routine (ISR).

VxWorks Kernel API Reference, 6.6 rtpSymTblldGet()

#### SMP CONSIDERATIONS

	This API is spinlock and intCpuLock restricted.
RETURNS	<b>RTP_ID</b> of the new RTP, or <b>ERROR</b> otherwise.
ERRNOS	Possible errnos returned from this routine are:
	<b>S_rtpLib_INVALID_FILE</b> The path to the executable file is not valid. The <i>rtpFileName</i> parameter is either null or the executable file cannot be found via the provided path. A valid path is a path that can be successfully accessed via the kernel shell.
	S_rtpLib_INVALID_TASK_OPTION One or more of the options specified for the initial task are not supported for a user-mode task.
	S_rtpLib_INSTANTIATE_FAILED The RTP object was created but failed to load and reach RTP_STATE_NORMAL
SEE ALSO	<b>rtpLib, rtpDelete( ), rtpInfoGet( ), rtpHookLib, memLib</b> , the VxWorks programmer guides.

# rtpSymTblIdGet()

NAME	<b>rtpSymTblIdGet( )</b> – Get the symbol table ID of an RTP	
SYNOPSIS	SYMTAB_ID rtpSymTblIdGet ( RTP_ID rtpId /* RTP ID whose symbol table ID is needed */ )	
DESCRIPTION	This routine gets and returns the symbol table ID for an RTP's symbol table.	
SMP CONSIDERATIO	<b>NS</b> This API is spinlock and intCpuLock restricted.	
RETURNS	The RTP's symbol table ID or NULL if the ID could not be retrieved	
ERRNOS	<b>S_objLib_OBJ_ID_ERROR</b> RTP object is not valid.	
SEE ALSO	rtpUtilLib, rtpLib	

### rtpSymsAdd()

rtpSymsAdd() – add symbols from an executable file to a RTP symbol table NAME SYNOPSIS STATUS rtpSymsAdd ( RTP\_ID rtpId, /\* RTP the symbols should be added to \*/ UINT32 regPolicy, /\* symbol registration policy \*/ char \* filePath /\* path and name of the executable file \*/ ) DESCRIPTION This command is provided as a help in case a RTP needs to be debugged but has been launched with an empty symbol table. It forces the registration of the symbols from an executable file into a RTP symbol table. Note that this command does not verify whether the symbols are already in the symbol table and does not prevent the creation of multiple occurences of these symbols. It is important to understand that symbols are added to the symbol table in the order of their registration and that the most recent entry will hide symbols of same name already registered. The **rtpLkup()** command will show all occurences of the symbols of a given name so it is possible to use their addresses instead of their names if there is a risk of confusion. The only required information is the RTP ID (*rtpId* parameter). The *regPolicy* parameter sets the symbol registration policy. The policy can be one of the following: 0x01 (RTP\_GLOBAL\_SYMBOLS) Add only global symbols to the symbol table. This is the default when the parameter is left null. 0x02 (RTP LOCAL SYMBOLS) Add only local symbols to the symbol table. 0x03 (RTP\_ALL\_SYMBOLS) Add both local and global symbols to the symbol table. The *filePath* parameter overrides the path recorded for the RTP. It may be left null if the symbols should be read from the same file as the one used to start the RTP with. This parameter must be used when the symbols should be read from a file stored in a different location than what was recorded when the RTP has been launched. Note that there is no runtime verification that the file corresponds to the executable used to launch the RTP. SMP CONSIDERATIONS This API is spinlock and intCpuLock restricted. OK if the symbols could be read and recorded, ERROR otherwise. RETURNS

VxWorks Kernel API Reference, 6.6 rtpSymsOverride()

ERRNO N/A

SEE ALSO usrRtpLib, rtpSymsRemove(), shlSymsAdd(), shlSymsRemove(), rtpSymsForce()

## rtpSymsOverride()

NAME	<b>rtpSymsOverride()</b> – override the RTP symbol registration policy	
SYNOPSIS	<pre>STATUS rtpSymsOverride   (    int overridePolicy /* override the symbol registration policy */   )</pre>	
DESCRIPTION	This command is provided as a help for debugging and monitoring RTPs launched by othe applications. It allows to temporarily bypass the symbol registration policy encoded in RTF and forces an alternate symbol registration policy which will apply to all applications launched after the change is made. The change is in effect until cancelled.	
	The <i>overridePolicy</i> parameter sets the alternate symbol registration policy. The policy can be one of the following:	
	- Prevent any symbol to be registered in the RTP's symbol table.	
	1 Add only global symbols to the symbol table.	
	3 Add both local and global symbols to the symbol table.	
	-1 Cancel any override of the symbol registration policy.	
	After the override of the RTP symbol registration policy is canceled, any newly launched application will get applied the symbol registration policy set by the parent application, or the shell.	
RETURNS	<b>OK</b> if the override policy could be set, <b>ERROR</b> otherwise.	
ERRNO	Not Available	
SEE ALSO	usrRtpLib, rtpSymsAdd( ), rtpSymsRemove( ), shlSymsAdd( ), shlSymsRemove( )	

# rtpSymsRemove()

NAME	rtpSymsRemove() – remove symbols from a RTP symbol table		
SYNOPSIS	STATUS rtpSymsRemove ( RTP_ID rtpId, /* RTP the symbols should be removed from */ UINT32 remPolicy /* symbol removal policy */ )		
DESCRIPTION	This command forces the removal of symbols from a RTP symbol table.		
	The <i>remPolicy</i> parameter sets the symbol removal policy. The policy can be one of the following:		
	0x02 ( <b>RTP_LOCAL_SYMBOLS</b> ) Remove only local symbols from the symbol table.		
	0x03 ( <b>RTP_ALL_SYMBOLS</b> ) Removes both local and global symbols from the symbol table.		
	Note: in the current implementation, this command will also remove the symbols related to shared libraries bound to the RTP. In order to remove only the symbols related to a specific shared library use <b>shlSymsRemove()</b> .		
RETURNS	OK if the symbols could be removed, ERROR otherwise.		
ERRNO	N/A		
SEE ALSO	usrRtpLib, rtpSymsAdd( ), shlSymsAdd( ), shlSymsRemove( ), rtpSymsForce( )		

# rtpTaskKill()

NAME	<b>rtpTaskKill( )</b> – send a signal to a task
SYNOPSIS	<pre>int rtpTaskKill   (   TASK_ID tid, /* task to send signal to */   int signo /* signal to send to task */  )</pre>
DESCRIPTION	This routine sends a kill signal <i>signo</i> to the RTP task specified by <i>tid</i> .
RETURNS	<b>OK</b> (0), or <b>ERROR</b> (-1) if the task ID or signal number is invalid.

VxWorks Kernel API Reference, 6.6 rtpTaskSigqueue()

ERRNO EINVAL

SEE ALSO rtpSigLib

# rtpTaskSigqueue()

NAME	<pre>rtpTaskSigqueue() – send a queued signal to a task</pre>	
SYNOPSIS	<pre>int rtpTaskSigqueue   (    TASK_ID tid,    int signo,    const union sigval value   )</pre>	
DESCRIPTION	The function <b>rtpTaskSigqueue()</b> sends the queued signal specified by <i>signo</i> with the signal-parameter value specified by <i>value</i> to the RTP task specified by <i>tid</i> .	
RETURNS	<b>OK</b> (0), or <b>ERROR</b> (-1) if the task ID or signal number is invalid, or if there are no queued-signal buffers available.	
ERRNO	EINVAL EAGAIN	
SEE ALSO	rtpSigLib	

# rtpi()

NAME	<b>rtpi( )</b> – display all tasks within an RTP	
SYNOPSIS	STATUS rtpi ( RTP_ID rtpId /* RTP identifier value, or 0 for task's RTP */ )	
DESCRIPTION	This command displays summary task information on tasks associated with the RT For more detailed information, the routine <b>ti()</b> or <b>taskShow()</b> should be used with specified <b>TASK_ID</b> .	

This routine displays tasks of an RTP in sorted order if the number of tasks is less than 100. If more than 100 tasks are within an RTP, the tasks will be displayed in the order they are created.

- **RETURNS** OK, or ERROR if *rtpId* is invalid
- ERRNO Possible errno values are:
  - S\_objLib\_OBJ\_ID\_ERROR An invalid RTP ID was provided.

S\_objLib\_ACCESS\_DENIED Unable to get exclusive access to the RTP to display tasks.

SEE ALSO usrRtpLib, i(), ti(), taskShow()

### s()

NAME	<b>s( )</b> – single-step a task		
SYNOPSIS	<pre>STATUS s ( int taskNameOrId, /* task to step; 0 = default */ INSTR * addr, /* address to step to; 0 = next instruction */ INSTR * addr1 /* address for npc, 0 = next instruction */ )</pre>		
DESCRIPTION	This routine single-steps a task that is stopped at a breakpoint.		
	To execute, enter:		
	-> s [task[,addr[,addr1]]]		
	If <i>task</i> is omitted or zero, the last task referenced is assumed. If <i>addr</i> is non-zero, then the program counter is changed to <i>addr</i> ; if <i>addr1</i> is non-zero, the next program counter is changed to <i>addr1</i> , and the task is stepped.		
CAVEAT	When a task is continued, <b>s()</b> does not distinguish between a stopped task or a task stopped by the debugger. Therefore, its use should be restricted to only those tasks being debugged.		
NOTE	The next program counter, <i>addr1</i> , is currently supported only by SPARC.		
RETURNS	<b>OK</b> , or <b>ERROR</b> if the debugging package is not installed, the task cannot be found, or the task is not suspended.		
ERRNO	N/A		

**SEE ALSO dbgLib**, **so()**, **c()**, *VxWorks Kernel Programmer's Guide: Kernel Shell*, *VxWorks Command-Line Tools User's Guide 2.2: Host Shell* 

### salCall()

NAME	salCall() – invoke a socket-based server		
SYNOPSIS	<pre>int salCall    (     int sockfd, /* client socket fd */    void * pSendBuf, /* message buffer */    int sendLen, /* size of message buffer */    void * pRecvBuf, /* reply buffer */    int recvLen /* size of reply buffer */    )</pre>		
DESCRIPTION	This routine sends a message to the server associated with the socket descriptor <i>sockfd</i> and waits for a reply. The message consists of the <i>sendLen</i> bytes pointed at by <i>pSendBuf</i> . The reply is placed in the <i>recvLen</i> bytes pointed at by <i>pRecvBuf</i> . If fewer than <i>recvLen</i> bytes are received the unused portion of <i>pRecvBuf</i> is not altered; if more than <i>recvLen</i> bytes are received the unused portion of the reply may be kept or discarded depending on the socket protocol being used.		
	If the socket descriptor is used by multiple clients, mutual exclusion needs to be provided before this routines is called. This is to avoid the case when a reply is intercepted by a higher priority task sharing the same <i>sockfd</i> .		
RETURNS	# of bytes placed in reply buffer, for connection based transport, 0 bytes may returned when the called end closes the connection; <b>ERROR</b> otherwise.		
ERRNO	S_salLib_INVALID_ARGUMENT An invalid argument was passed to this routine.		
SEE ALSO	salClient		

### salCreate()

SYNOPSIS SAL\_SERVER\_ID salCreate

	<pre>(     const char * name, /* service name */     int sockFamily, /* desired socket address family */     int sockType, /* desired socket type */     int sockProtocol, /* desired socket protocol */     const struct salSockopt * options, /* array of socket options */     int numOptions /* number of socket options */ )</pre>			
DESCRIPTION	This routine creates a socket-based server. One or more sockets are created for the server, and the service is registered with SNS using the service name <i>name</i> .			
	name is represented in the following URL format:			
	[SNS:]service_name[@scope]			
	Refer to <b>snsLib</b> for more information on the format.			
	This routine tries to create one or more sockets for the combination defined by <i>sockFamily</i> , <i>sockType</i> , and <i>sockProtocol</i> . If the <i>sockFamily</i> specified is <b>AF_UNSPEC</b> , then a socket creation attempt is made with each family type supported by SAL. If the <i>sockType</i> specified is 0, then a socket creation attempt is made with each socket type. If the <i>sockProtocol</i> specified is 0, then the default protocol for that family is used.			
	The <i>sockFamily, sockType</i> , and <i>sockProtocol</i> parameters can be used to limit the server to a given address family and/or socket type and/or socket protocol. salCreate supports connection-oriented message based socket types only, and creates a passive listening socket.			
	The <i>options</i> parameter points to an array of <i>numOptions</i> socket option values that are applied to each server socket created. If the socket cannot be successfully configured, it is closed and is not incorporated into the server.			
WARNING	Once successfully created, the SAL server must still be configured with one or more processing routines before calling <b>salRun()</b> .			
RETURNS	created server ID, NULL if fails.			
ERRNO	S_salLib_INVALID_ARGUMENT An invalid argument was passed to this routine.			
	S_salLib_SERVER_SOCKET_ERROR Unable to create any sockets with the desired properties			
	S_salLib_SNS_UNAVAILABLE Unable to establish connection to the SNS server task.			
	S_salLib_SNS_DID_NOT_REPLY Did not receive a reply from the SNS server task.			
	S_salLib_SNS_PROTOCOL_ERROR Received an invalid reply from the SNS server task.			

S_salLib_SNS_OUT_OF_MEMORY
The SNS server task has insufficient memory to register the service.
The SNS server task has insufficient memory to register the service.

### S\_salLib\_SERVICE\_ALREADY\_EXISTS

The specified service has already been registered with SNS.

SEE ALSO salServer, salDelete(), salRemove(), salServerRtnSet()

### salDelete()

NAME	salDelete() – delete a named socket-based server	
SYNOPSIS	STATUS salDelete ( SAL_SERVER_ID server /* server structure to use */ )	
DESCRIPTION	This routine deletes the socket-based server specified by <i>server</i> . and frees the server data structure memory. All the sockets associated with <i>server</i> are closed. The associated server is deregistered from SNS.	
RETURNS	OK or ERROR.	
ERRNO	S_salLib_INVALID_ARGUMENT An invalid argument was passed to this routine.	
	S_salLib_SNS_UNAVAILABLE Unable to establish connection to the SNS server task.	
	S_salLib_SNS_DID_NOT_REPLY Did not receive a reply from the SNS server task.	
	S_salLib_SNS_PROTOCOL_ERROR Received an invalid reply from the SNS server task.	
	S_salLib_INVALID_SERVICE_DESCRIPTOR Service descriptor is not registered with SNS, or has a different owner.	
SEE ALSO	salServer, salCreate(), salRemove()	

### salNameFind()

```
salNameFind() - find services with the specified name
NAME
SYNOPSIS
                 int salNameFind
                      const char * pattern,
                                                                   /* services name pattern */
                      char
                                   servName[][SAL_SERV_NAME_MAXSIZE],
                                                                   /* buffer to hold the returned
                 name */
                      int
                                    num,
                                                                    /* number of element in the
                 servNames */
                     void ** ppCookie
                                                                    /* cookie get/return last
                 matching address */
                      )
DESCRIPTION
                 This function returns services with names that match the specified pattern.
                 Applications provide the buffer for storing the returned names. The function returns the
                 number of names found. The function also returns a cookie for follow up searching.
                 pattern is represented in the following URL format:
                 [SNS:]service_name[@scope]
                 If pattern contains wildcard characters, the routine will search for all services that match the
                 pattern.
                 Refer to snsLib for more information on the format and the use of wildcards.
                 The function returns a number of services no greater than num. If more matches are found
                 the function can be called again to retrieve the remaining values. The behavior of the
                 function is determined by the ppCookie field.
                 In order to guarantee all data can be retrieved (possibly through subsequent calls) when the
                 function is called for the first time, the ppCookie field needs to be non-NULL and the value
                 *ppCookie needs to be set to NULL. If the returned value *ppCookie is still NULL, this means
                 all the services matching the pattern have been retrieved. XXX - Yiming to verify If the
                 returned value *ppCookie is not NULL, this means that more matches might be available. In
                 this case, the client application can call salNameFind() again using the returned ppCookie to
                 retrieve further entries.
                 Hence, in order to start a new search, either ppCookie is NULL (in which case the function
                 can not be called again to retrieve more values) or *ppCookie is NULL.
                 >=0: number of services found, -1: error.
RETURNS
ERRNO
                 S salLib INVALID ARGUMENT
                      Invalid argument.
```

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#### S\_salLib\_SNS\_UNAVAILABLE Unable to establish communications with the SNS server task.

SEE ALSO salClient, salSocketFind(), snsLib

### salOpen()

NAME salOpen() – establish communication with a named socket-based server SYNOPSIS int salOpen ( const char \* name /\* service name in URL format \*/ DESCRIPTION This routine establishes a connection to the server application corresponding to the SNS service name *name*. If the specified service exists **salOpen()** tries to connect to each of the server's sockets in turn, until it is successful or all sockets have been tried; it returns the resulting socket descriptor. *name* is represented in the following URL format: [SNS:]service\_name[@scope] If *name* contains wildcard characters, the routine will use the first matching service. Refer to **snsLib** for more information on the format and the use of wildcards. This routine uses the default socket options for the client socket it creates; if special options are required by the client before completing the connection, use **salSocketFind()** to establish communication with the server. User should close the returned socket using **close()**. >=0: the descriptor of the newly connected socket; -1 : cannot establish communication. RETURNS ERRNO S\_salLib\_INVALID\_ARGUMENT An invalid argument was passed to this routine. S\_salLib\_SNS\_UNAVAILABLE Unable to establish connection to the SNS server task. S salLib SNS DID NOT REPLY Did not receive a reply from the SNS server task. S\_salLib\_SNS\_PROTOCOL\_ERROR Received an invalid reply from the SNS server task.

	S_salLib_SERVICE_NOT_FOUND The specified service is not registered with SNS.
	S_salLib_INVALID_SERVICE_DESCRIPTOR The specified service was deregistered from SNS before all socket addresses could be examined.
	S_salLib_CLIENT_SOCKET_ERROR Unable to connect to any of the specified server socket addresses.
SEE ALSO	salClient, close(), salSocketFind(), snsLib

## salRemove()

NAME	salRemove() – Remove service from SNS by name	
SYNOPSIS	STATUS salRemove ( const char * name /* service name */ )	
DESCRIPTION	This function removes a service identified by <i>name</i> from SNS. Unlike <b>salDelete()</b> , which requires the caller and service owner to be in the same memory space, this function can delete any service as long as the service is visible to the caller. Therefore, a service with scope <b>node</b> can be deleted by any task on the same node, and a service with scope <b>private</b> can only be deleted by tasks in the same memory space. Further, services of scope <b>cluster</b> (or larger) can only be deleted by the node that created them.	
<i>name</i> is represented in the following URL format:		
	[SNS:]service_name[@scope]	
Refer to <b>snsLib</b> for more information on the format.		
	name must uniquely identify a service:	
	service_name should not contain any wildcard character	
	scope must refer a specific level (i.e. the "upto_" prefix can not be used)	
NOTE	This routine removes only the service name from SNS. It does not remove the service, nor does it close any of the sockets associated to it. These features are provided by <b>salDelete()</b> .	
RETURNS	<b>OK</b> if the service is removed, <b>ERROR</b> otherwise.	

VxWorks Kernel API Reference, 6.6 salRun()

ERRNO	S_salLib_INVALID_ARGUMENT The service name is invalid
	S_salLib_SERVICE_NOT_FOUND The specified service is not found.

SEE ALSO salServer, salDelete(), salCreate(), snsLib

## salRun()

NAME	salRun() – activate a socket-based server	
SYNOPSIS	<pre>STATUS salRun ( SAL_SERVER_ID server, /* server structure to use */ void * pData /* user private data */ )</pre>	
DESCRIPTION	This routine activates the SAL server specified by <i>server</i> . The server monitors all sockets associated with the server, and calls an appropriate processing routine whenever a socket requires attention.	
	Once invoked, this routine will execute indefinitely and will return only when the server terminates.	
	Server termination occurs automatically if salRun() detects an error.	
	The server can terminate also by the application through the processing routine return value SAL_RUN_TERMINATE. In this case salRun() simply returns OK.	
	In both cases <b>salRun()</b> does not close any socket. <b>salDelete()</b> should be called to perform the cleanup.	
	The parameter $pData$ can be used to pass any user data. This data is passed to the processing routines when they are being called.	
	Processing routines should be configured in the server before this routine is called.	
RETURNS	OK if server is terminated by processing routine, ERROR otherwise.	
ERRNO	S_salLib_INVALID_ARGUMENT An invalid argument was passed to this routine.	
	S_salLib_SERVER_SOCKET_ERROR A server socket has failed unexpectedly.	
	S_salLib_INTERNAL_ERROR The server's internal data structure has become corrupted.	

SEE ALSO salServer, salServerRtnSet()

### salServerRtnSet()

NAME	salServerRtnSet() – configures the processing routine with the SAL server		
SYNOPSIS	<pre>STATUS salServerRtnSet   (    SAL_SERVER_ID svrId, /* server ID */    SAL_RTN_TYPE rtnType, /* type of processing routine to set */    SAL_SERV_RTN routine /* processing routine entry point */   )</pre>		
DESCRIPTION	This routine configures a processing routine with the server $pSrvrId$ . The processing routine is identified by the type $rtnType$ and the <b>SAL_SERV_RTN</b> function pointer <i>routine</i> .		
	It accepts the following <i>rtnType</i> :		
	SAL_RTN_READ read routine		
	SAL_RTN_ACCEPT accept routine		
	If <i>routine</i> is <b>NULL</b> , the processing routine is cleared and the default handler will be used, if available.		
	This function must be called before activating the SAL server, i.e. before the call to <b>salRun()</b> .		
RETURNS	OK or ERROR		
ERRNO	S_salLib_INVALID_ARGUMENT An invalid argument was passed to this routine.		
SEE ALSO	salServer, salRun()		

### salSocketFind()

NAME	<pre>salSocketFind() - find sockets for a named socket-based server</pre>			
SYNOPSIS	STATUS salSocketFind ( const char *	name,	/* service name	in URL format */

VxWorks Kernel API Reference, 6.6 salSocketFind()

	<pre>int sockFamily, /* desired socket address family */ int sockType, /* desired socket type */ int sockProtocol, /* desired socket protocol */ struct addrinfo ** ppSockInfoList /* list of socket entries */ )</pre>			
DESCRIPTION	This routine looks for sockets related to a server application registered with SNS, which matches the specified search criteria. Each socket entry associated with the SNS service name <i>name</i> is examined to see if it is compatible with the restrictions imposed by <i>sockFamily</i> , <i>sockType</i> , and <i>sockProtocol</i> . The search succeeds if at least one matching socket entry is found.			
	<i>name</i> is represented in the following URL format:			
	[SNS:]service_name[@scope]			
	Please refer to <b>snsLib</b> for more information on the format.			
	If <i>name</i> contains wildcard characters, the function will only find the first matching service and retrieve its socket information.			
	To obtain the complete list of service matching the given pattern, use the <b>salNameFind()</b> routine.			
	If <i>sockInfoList</i> is not NULL then a list of the matching socket entries is created, and <i>sockInfoList</i> is set to the start of the list. However if <i>sockInfoList</i> is NULL, or the service specified by <i>name</i> does not exist, then no list of socket entries is created and <i>sockInfoList</i> is left unchanged.			
WARNING	The storage for the socket list created by this routine must be released by calling <b>snsfreeaddrinfo()</b> when the list is no longer required.			
RETURNS	OK or ERROR			
ERRNO	S_salLib_INVALID_ARGUMENT An invalid argument was passed to this routine.			
	S_salLib_SNS_UNAVAILABLE Unable to establish connection to the SNS server task.			
	S_salLib_SNS_DID_NOT_REPLY Did not receive a reply from the SNS server task.			
	S_salLib_SNS_PROTOCOL_ERROR Received an invalid reply from the SNS server task.			
	S_salLib_SERVICE_NOT_FOUND The specified service is not registered with SNS.			
	S_salLib_INVALID_SERVICE_DESCRIPTOR The specified service was deregistered from SNS before all socket entries could be examined.			
#### S\_salLib\_NO\_SOCKET\_FOUND The specified service has no sockets that match the desired criteria.

SEE ALSO salClient, salNameFind(), snsLib

### sbeRegister()

**NAME sbeRegister()** – register with the VxBus subsystem

SYNOPSIS void sbeRegister(void)

**DESCRIPTION** This routine registers the Broadcom driver with VxBus as a child of the PCI bus type.

RETURNS N/A

ERRNO N/A

SEE ALSO sbeVxbEnd

#### scMemValEnable()

NAME	scMemValEnable() – enable or disable pointer/buffer validation in system calls
SYNOPSIS	<pre>void scMemValEnable   (   BOOL enable /* TRUE: enable validation; FALSE: disable validation */ )</pre>
DESCRIPTION	This routine either enables or disables pointer/buffer validations in system calls, system wide. By default when the OS starts, pointer/buffer validation is ON. Pointer validation can also be disabled for a given RTP by passing the option <b>RTP_BUFFER_VAL_OFF</b> as part of the <i>options</i> parameter of <b>rtpSpawn()</b> .
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	scMemVal, scMemValidate( ), rtpSpawn( )

#### scMemValidate()

scMemValidate() – validate an address range passed to a system call routine NAME SYNOPSIS STATUS scMemValidate ( const void \* addr, /\* start address \*/ UINT size, /\* address range size in bytes \* / SC\_PROT\_ATTR access /\* minimal access in supervisor mode \*/ ) DESCRIPTION The routine scMemValidate() should be used by system call validation code to verify that a pointer passed as a parameter to a system call routine points to a memory location that belongs to the calling RTP memory context, and that this memory can be dereferenced by the kernel while it executes the system call. The routine checks if the memory range [addr ... (addr + size -1)] belongs entirely to either: - the calling task's stack - the memory section corresponding to the read-only segment of the RTP: text + rodata segment. the memory section corresponding to the data segment of the RTP (read-write access). - the memory section corresponding to the bss segment of the RTP (read-write access). - a block that was mapped by the calling RTP vith **mmap()**. Note that in addition to memory block obtained by direct call to **mmap()**, this also includes blocks that correspond the RTP heap and shared library private data (data and bss segments) that the RTP is attached to. - a memory section corresponding to a Shared Data region that the RTP has mapped. Note that this also takes care of the read-only code of SLs that the RTP is attached to. - other RTP-private memory, such as stack of other tasks in the same RTP. Note that if *size* is equal to 0, **scMemVal()** simply returns **OK** without validating the zero-length buffer. Once the buffer to validate is matched with one of the memory section listed above, that is the buffer is contained entirely within this memory section, scMemValidate() checks that the access permissions to this memory section authorizes the access defined by access. The type of access requested is defined by *access* parameter and can take the following values: - SC\_PROT\_READ (0x1)- SC\_PROT\_WRITE (0x2) If *access* is set to 0, then it defaults to (SC\_PROT\_READ | SC\_PROT\_WRITE).

	The parameters <i>addr</i> and <i>size</i> do not need to be aligned to a MMU page size.
	Buffer validation accross system calls can be disabled system wide by calling scMemValEnable (FALSE). To re-enable buffer validation system-wide simply call scMemValEnable(TRUE). Buffer validation accross system calls can be disabled for a given RTP, by passing the option RTP_BUFFER_VAL_OFF as part of the <i>options</i> parameter when calling rtpSpawn(). For more details refer to the rtpSpawn() manual entry.
RETURNS	<b>OK</b> if the address range is valid, <b>ERROR</b> otherwise.
ERRNO	Possible errno generated by this routine include:
	EINVAL The access parameter passed is invalid.
	ENOTSUP The routine was called in unsupported context.
	ENOMEM Memory validation failed with due to boundary constraints.
	EACCES Memory validation failed due to access constraints.

SEE ALSO scMemVal, scMemValEnable(), rtpSpawn()

# sched\_get\_priority\_max()

NAME	<pre>sched_get_priority_max() - get the maximum priority (POSIX)</pre>
SYNOPSIS	<pre>int sched_get_priority_max   (    int policy /* scheduling policy */ )</pre>
DESCRIPTION	This routine returns the value of the highest possible task priority for a specified scheduling policy (SCHED_FIFO or SCHED_RR).
NOTE	If the global variable <b>posixPriorityNumbering</b> is <b>FALSE</b> , the VxWorks native priority numbering scheme is used, in which higher priorities are indicated by smaller numbers. This is different than the priority numbering scheme specified by POSIX, in which higher priorities are indicated by larger numbers.
RETURNS	Maximum priority value, or -1 (ERROR) on error.

VxWorks Kernel API Reference, 6.6 sched\_get\_priority\_min()

**ERRNO EINVAL** – invalid scheduling policy.

SEE ALSO schedPxLib

### sched\_get\_priority\_min()

NAME	<pre>sched_get_priority_min() - get the minimum priority (POSIX)</pre>
SYNOPSIS	<pre>int sched_get_priority_min   (    int policy /* scheduling policy */ )</pre>
DESCRIPTION	This routine returns the value of the lowest possible task priority for a specified scheduling policy (SCHED_FIFO or SCHED_RR).
NOTE	If the global variable <b>posixPriorityNumbering</b> is <b>FALSE</b> , the VxWorks native priority numbering scheme is used, in which higher priorities are indicated by smaller numbers. This is different than the priority numbering scheme specified by POSIX, in which higher priorities are indicated by larger numbers.
RETURNS	Minimum priority value, or -1 (ERROR) on error.
ERRNO	EINVAL – invalid scheduling policy.
SEE ALSO	schedPxLib

### sched\_getparam()

NAME	<b>sched_getparam()</b> – get the scheduling parameters for a specified task (POSIX)
SYNOPSIS	int sched_getparam ( pid_t tid, /* task ID */ struct sched_param * param /* scheduling param to store priority */ )
DESCRIPTION	This routine gets the scheduling priority for a specified task, <i>tid</i> . If <i>tid</i> is 0, it gets the priority of the calling task. The task's priority is copied to the <b>sched_param</b> structure pointed to by

param.

NOTE	If the global variable <b>posixPriorityNumbering</b> is <b>FALSE</b> , the VxWorks native priority numbering scheme is used, in which higher priorities are indicated by smaller numbers. This is different than the priority numbering scheme specified by POSIX, in which higher priorities are indicated by larger numbers.
RETURNS	0 (OK) if successful, or -1 (ERROR) on error.
ERRNO	ESRCH – invalid task ID.
SEE ALSO	schedPxLib

# sched\_getscheduler()

NAME	<pre>sched_getscheduler() - get the current scheduling policy (POSIX)</pre>
SYNOPSIS	<pre>int sched_getscheduler   (    pid_t tid /* task ID */   )</pre>
DESCRIPTION	This routine returns the currents scheduling policy (i.e., SCHED_FIFO or SCHED_RR).
RETURNS	Current scheduling policy (SCHED_FIFO or SCHED_RR), or -1 (ERROR) on error.
ERRNO	ESRCH – invalid task ID.
SEE ALSO	schedPxLib

# sched\_rr\_get\_interval()

NAME	<pre>sched_rr_get_interval() - get the cut</pre>	rrent time slice (POSIX)
SYNOPSIS	<pre>int sched_rr_get_interval   (     pid_t tid,     struct timespec * interval   )</pre>	/* task ID */ /* struct to store time slice */
DESCRIPTION	This routine sets <i>interval</i> to the curren currently enabled.	nt time slice period if round-robin scheduling is

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VxWorks Kernel API Reference, 6.6 sched\_setparam()

**RETURNS** 0 (OK) if successful, -1 (ERROR) on error.

ERRNO EINVAL – round-robin scheduling is not currently enabled. ESRCH – invalid task ID.

SEE ALSO schedPxLib

#### sched\_setparam()

NAME sched\_setparam() - set a task's priority (POSIX) SYNOPSIS int sched\_setparam ( pid\_t tid, /\* task ID \*/ const struct sched\_param \* param /\* scheduling parameter \*/ ) DESCRIPTION This routine sets the priority of a specified task, *tid*. If *tid* is 0, it sets the priority of the calling task. Valid priority numbers are 0 through 255. The *param* argument is a structure whose member **sched\_priority** is the integer priority value. For example, the following program fragment sets the calling task's priority to 13 using POSIX interfaces: #include "sched.h" . . . struct sched\_param AppSchedPrio; . . . AppSchedPrio.sched\_priority = 13; if ( sched setparam (0, &AppSchedPrio) != OK ) { ... /\* recovery attempt or abort message \*/ } . . . NOTE If the global variable **posixPriorityNumbering** is **FALSE**, the VxWorks native priority numbering scheme is used, in which higher priorities are indicated by smaller numbers. This is different than the priority numbering scheme specified by POSIX, in which higher priorities are indicated by larger numbers. 0 (OK) if successful, or -1 (ERROR) on error. RETURNS ERRNO EINVAL – scheduling priority is outside valid range. ESRCH - task ID is invalid. schedPxLib SEE ALSO

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### sched\_setscheduler()

NAME	<b>sched_setscheduler()</b> – set scheduling policy and scheduling parameters (POSIX)	
SYNOPSIS	<pre>int sched_setscheduler   (    pid_t tid, /* task ID */    int policy, /* scheduling policy requested */    const struct sched_param * param /* scheduling parameters requested */   )</pre>	
DESCRIPTION	This routine sets the scheduling policy and scheduling parameters for a specified task, <i>tid</i> . If <i>tid</i> is 0, it sets the scheduling policy and scheduling parameters for the calling task.	
	Because VxWorks does not set scheduling policies (e.g., round-robin scheduling) on a task-by-task basis, setting a scheduling policy that conflicts with the current system policy simply fails and errno is set to <b>EINVAL</b> . If the requested scheduling policy is the same as the current system policy, then this routine acts just like <b>sched_setparam()</b> .	
NOTE	If the global variable <b>posixPriorityNumbering</b> is <b>FALSE</b> , the VxWorks native priority numbering scheme is used, in which higher priorities are indicated by smaller numbers. This is different than the priority numbering scheme specified by POSIX, in which higher priorities are indicated by larger numbers.	
RETURNS	The previous scheduling policy (SCHED_FIFO or SCHED_RR), or -1 (ERROR) on error.	
ERRNO	EINVAL – scheduling priority is outside valid range, or it is impossible to set the specified scheduling policy. ESRCH – invalid task ID.	
SEE ALSO	schedPxLib	

### sched\_yield()

NAMEsched\_yield() - relinquish the CPU (POSIX)SYNOPSISint sched\_yield (void)DESCRIPTIONThis routine forces the running task to give up the CPU.RETURNS0 (OK) if successful, or -1 (ERROR) on error.

VxWorks Kernel API Reference, 6.6 scsi2lflnit()

ERRNO Not Available

SEE ALSO schedPxLib

#### scsi2IfInit( )

NAME scsi2IfInit() – initialize the SCSI-2 interface to scsiLib

SYNOPSIS void scsi2IfInit (void)

**DESCRIPTION** This routine initializes the SCSI-2 function interface by adding all the routines in scsi2Lib plus those in scsiDirectLib and scsiCommonLib. It is invoked at startup if the component INCLUDE\_SCSI2 is configured in VxWorks. The calling interface remains the same between SCSI-1 and SCSI-2; this routine simply sets the calling interface function pointers to the SCSI-2 functions.

RETURNS N/A

ERRNO Not Available

SEE ALSO scsi2Lib

#### scsiAutoConfig()

NAME	<pre>scsiAutoConfig() – configure all devices connected to a SCSI controller</pre>
SYNOPSIS	STATUS scsiAutoConfig ( SCSI_CTRL *pScsiCtrl /* ptr to SCSI controller info */ )
DESCRIPTION	This routine cycles through all valid SCSI bus IDs and logical unit numbers (LUNs), attempting a <b>scsiPhysDevCreate()</b> with default parameters on each. All devices which support the INQUIRY command are configured. The <b>scsiShow()</b> routine can be used to find the system table of SCSI physical devices attached to a specified SCSI controller. In addition, <b>scsiPhysDevIdGet()</b> can be used programmatically to get a pointer to the <b>SCSI_PHYS_DEV</b> structure associated with the device at a specified SCSI bus ID and LUN.
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>pScsiCtrl</i> and the global variable <b>pSysScsiCtrl</b> are both <b>NULL</b> .

Not Available ERRNO

SEE ALSO scsiLib

#### scsiBlkDevCreate()

scsiBlkDevCreate() – define a logical partition on a SCSI block device NAME

SYNOPSIS BLK DEV \* scsiBlkDevCreate int int

)

- SCSI\_PHYS\_DEV \* pScsiPhysDev, /\* ptr to SCSI physical device info \*/ numBlocks, /\* number of blocks in block device \*/ blockOffset /\* address of first block in volume \*/
- This routine creates and initializes a **BLK\_DEV** structure, which describes a logical partition DESCRIPTION on a SCSI physical-block device. A logical partition is an array of contiguously addressed blocks; it can be completely described by the number of blocks and the address of the first block in the partition. In normal configurations partitions do not overlap, although such a condition is not an error.
- If *numBlocks* is 0, the rest of device is used. NOTE
- RETURNS A pointer to the created **BLK\_DEV**, or **NULL** if parameters exceed physical device boundaries, if the physical device is not a block device, or if memory is insufficient for the structures.

ERRNO Not Available

scsiLib SEE ALSO

#### scsiBlkDevInit( )

NAME scsiBlkDevInit() - initialize fields in a SCSI logical partition

```
SYNOPSIS
                    void scsiBlkDevInit
                         (
                         SCSI_BLK_DEV * pScsiBlkDev, /* ptr to SCSI block dev. struct */
int blksPerTrack, /* blocks per track */
                                                               /* number of heads */
                         int
                                            nHeads
                         )
```

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VxWorks Kernel API Reference, 6.6 scsiBlkDevShow()

 DESCRIPTION
 This routine specifies the disk-geometry parameters required by certain file systems (for example, dosFs). It is called after a SCSI\_BLK\_DEV structure is created with scsiBlkDevCreate(), but before calling a file system initialization routine. It is generally required only for removable-media devices.

 RETURNS
 N/A

 ERRNO
 Not Available

 SEE ALSO
 scsiLib

#### scsiBlkDevShow()

NAME	<pre>scsiBlkDevShow() - show the BLK_DEV structures on a specified physical device</pre>
SYNOPSIS	<pre>void scsiBlkDevShow   (    SCSI_PHYS_DEV * pScsiPhysDev /* ptr to SCSI physical device info */   )</pre>
DESCRIPTION	This routine displays all of the <b>BLK_DEV</b> structures created on a specified physical device. This routine is called by <b>scsiShow()</b> but may also be invoked directly, usually from the shell.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	scsiLib, scsiShow()

### scsiBusReset()

NAME	<b>scsiBusReset( )</b> – pulse the reset signal on the SCSI bus
SYNOPSIS	STATUS scsiBusReset ( SCSI_CTRL * pScsiCtrl /* ptr to SCSI controller info */ )

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DESCRIPTION	This routine calls a controller-specific routine to reset a specified controller's SCSI bus. If no controller is specified ( <i>pScsiCtrl</i> is 0), the value in the global variable <b>pSysScsiCtrl</b> is used.
RETURNS	OK, or ERROR if there is no controller or controller-specific routine.
ERRNO	Not Available
SEE ALSO	scsiLib

### scsiCacheSnoopDisable()

NAME	<pre>scsiCacheSnoopDisable() - inform SCSI that hardware snooping of caches is disabled</pre>
SYNOPSIS	<pre>void scsiCacheSnoopDisable   (    SCSI_CTRL * pScsiCtrl /* pointer to a SCSI_CTRL structure */   )</pre>
DESCRIPTION	This routine informs the SCSI library that hardware snooping is disabled and that <b>scsi2Lib</b> should execute any neccessary cache coherency code. In order to make <b>scsi2Lib</b> aware that hardware snooping is disabled, this routine should be called after all SCSI-2 initializations, especially after <b>scsi2CtrlInit()</b> .
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	scsi2Lib

# scsiCacheSnoopEnable()

NAME	scsiCacheSnoopEnable() – inform SCSI that hardware snooping of caches is enabled
SYNOPSIS	void scsiCacheSnoopEnable ( SCSI_CTRL * pScsiCtrl /* pointer to a SCSI_CTRL structure */ )
DESCRIPTION	This routine informs the SCSI library that hardware snooping is enabled and that <b>scsi2Lib</b> need not execute any cache coherency code. In order to make <b>scsi2Lib</b> aware that hardware

VxWorks Kernel API Reference, 6.6 scsiCacheSynchronize()

snooping is enabled, this routine should be called after all SCSI-2 initializations, especially after scsi2CtrlInit().

RETURNS N/A

ERRNO Not Available

SEE ALSO scsi2Lib

#### scsiCacheSynchronize()

**scsiCacheSynchronize()** – synchronize the caches for data coherency NAME SYNOPSIS void scsiCacheSynchronize ( SCSI\_THREAD \* pThread, /\* ptr to thread info \*/ SCSI CACHE ACTION action /\* cache action required \*/ ) This routine performs whatever cache action is necessary to ensure cache coherency with DESCRIPTION respect to the various buffers involved in a SCSI command. The process is as follows: 1. The buffers for command, identification, and write data, which are simply written to SCSI, are flushed before the command. 2. The status buffer, which is written and then read, is cleared (flushed and invalidated) before the command. 3. The data buffer for a read command, which is only read, is cleared before the command. The data buffer for a read command is cleared before the command rather than invalidated after it because it may share dirty cache lines with data outside the read buffer. DMA drivers for older versions of the SCSI library have flushed the first and last bytes of the data buffer before the command. However, this approach is not sufficient with the enhanced SCSI library because the amount of data transferred into the buffer may not fill it, which would cause dirty cache lines which contain correct data for the un-filled part of the buffer to be lost when the buffer is invalidated after the command. To optimize the performance of the driver in supporting different caching policies, the routine uses the CACHE\_USER\_FLUSH macro when flushing the cache. In the absence of a CACHE\_USER\_CLEAR macro, the following steps are taken: If there is a non-NULL flush routine in the **cacheUserFuncs** structure, the cache is 1. cleared.

- 2. If there is a non-NULL invalidate routine, the cache is invalidated.
- 3. Otherwise nothing is done; the cache is assumed to be coherent without any software intervention.

Finally, since flushing (clearing) cache line entries for a large data buffer can be time-consuming, if the data buffer is larger than a preset (run-time configurable) size, the entire cache is flushed.

RETURNS N/A

Not Available ERRNO

scsi2Lib SEE ALSO

#### scsiErase()

scsiErase() – issue an ERASE command to a SCSI device NAME

SYNOPSIS STATUS scsiErase (

SCSI\_PHYS\_DEV \*pScsiPhysDev, /\* ptr to SCSI physical device \*/ /\* TRUE for entire tape erase \*/ BOOL longErase

- This routine issues an ERASE command to a specified SCSI device. DESCRIPTION
- RETURNS OK, or ERROR if the command fails.

)

Not Available ERRNO

SEE ALSO scsiSeqLib

#### scsiFormatUnit( )

NAME scsiFormatUnit() - issue a FORMAT\_UNIT command to a SCSI device SYNOPSIS STATUS scsiFormatUnit SCSI\_PHYS\_DEV \* pScsiPhysDev, /\* ptr to SCSI physical device \*/ cmpDefectList, /\* whether defect list is complete \*/ BOOL

VxWorks Kernel API Reference, 6.6 scsildentMsgBuild()

int defListFormat, /\* defect list format \*/
int vendorUnique, /\* vendor unique byte \*/
int interleave, /\* interleave factor \*/
char \* buffer, /\* ptr to input data buffer \*/
int bufLength /\* length of buffer in bytes \*/
)

**DESCRIPTION** This routine issues a **FORMAT\_UNIT** command to a specified SCSI device.

**RETURNS** OK, or **ERROR** if the command fails.

ERRNO Not Available

SEE ALSO scsiLib

### scsiIdentMsgBuild()

NAME	scsiIdentMsgBuild() – build an identification message	
SYNOPSIS	<pre>int scsiIdentMsgBuild   (    UINT8 * msg,    SCSI_PHYS_DEV * pScsiPhysDev,    SCSI_TAG_TYPE tagType,    UINT tagNumber   )</pre>	
DESCRIPTION	This routine builds an identification message in the caller's buffer, based on the specified physical device, tag type, and tag number.	
	If the target device does not support messages, there is no identification message to build.	
	Otherwise, the identification message consists of an IDENTIFY byte plus an optional QUEUE TAG message (two bytes), depending on the type of tag used.	
NOTE	This function is not intended for use by application programs.	
RETURNS	The length of the resulting identification message in bytes or -1 for ERROR.	
ERRNO	Not Available	
SEE ALSO	scsi2Lib	

	scsiIdentMsgParse()
NAME	scsildentMsgParse() – parse an identification message
SYNOPSIS	<pre>SCSI_IDENT_STATUS scsiIdentMsgParse (     SCSI_CTRL * pScsiCtrl,     UINT8 * msg,     int msgLength,     SCSI_PHYS_DEV ** ppScsiPhysDev,     SCSI_TAG * pTagNum )</pre>
DESCRIPTION	This routine scans a (possibly incomplete) identification message, validating it in the process. If there is an IDENTIFY message, it identifies the corresponding physical device.
	If the physical device is currently processing an untagged (ITL) nexus, identification is complete. Otherwise, the identification is complete only if there is a complete QUEUE TAG message.
	If there is no physical device corresponding to the IDENTIFY message, or if the device is processing tagged (ITLQ) nexuses and the tag does not correspond to an active thread (it may have been aborted by a timeout, for example), then the identification sequence fails.
	The caller's buffers for physical device and tag number (the results of the identification process) are always updated. This is required by the thread event handler (see <b>scsiMgrThreadEvent()</b> .)
NOTE	This function is not intended for use by application programs.
RETURNS	The identification status (incomplete, complete, or rejected).
ERRNO	Not Available
SEE ALSO	scsi2Lib

### scsiInquiry()

NAME scsiInquiry() - issue an INQUIRY command to a SCSI device
SYNOPSIS STATUS scsiInquiry
(
SCSI\_PHYS\_DEV \* pScsiPhysDev, /\* ptr to SCSI physical device \*/
char \* buffer, /\* ptr to input data buffer \*/

	int )	bufLength	/* length of buffer in bytes */
DESCRIPTION	This routine issues an	INQUIRY comm	and to a specified SCSI device.
RETURNS	OK, or ERROR if the co	ommand fails.	
ERRNO	Not Available		
SEE ALSO	scsiLib		

# scsiIoctl()

NAME	scsiIoctl() – perform a device-specific I/O control function
SYNOPSIS	<pre>STATUS scsiloctl   (    SCSI_PHYS_DEV * pScsiPhysDev, /* ptr to SCSI block device info */    int function, /* function code */    int arg /* argument to pass called function */   )</pre>
DESCRIPTION	This routine performs a specified <b>ioctl</b> function using a specified SCSI block device.
RETURNS	The status of the request, or <b>ERROR</b> if the request is unsupported.
ERRNO	Not Available
SEE ALSO	scsiLib

# scsiLoadUnit( )

NAME	scsiLoadUnit() – issu	ie a LOAD/UN	LOAD command to a SCSI device	
SYNOPSIS	STATUS scsiLoadUni	t		
	SCSI_SEQ_DEV * BOOL BOOL BOOL )	pScsiSeqDev, load, reten, eot	/* ptr to SCSI physical device /* TRUE=load, FALSE=unload /* TRUE=retention and unload /* TRUE=end of tape and unload	*/ */ */

**DESCRIPTION** This routine issues a LOAD/UNLOAD command to a specified SCSI device.

**RETURNS** OK, or ERROR if the command fails.

ERRNO Not Available

SEE ALSO scsiSeqLib

### scsiMgrBusReset()

NAME	<b>scsiMgrBusReset( )</b> – handle a controller-bus reset event
SYNOPSIS	<pre>void scsiMgrBusReset   (    SCSI_CTRL * pScsiCtrl /* SCSI ctrlr on which bus reset */ )</pre>
DESCRIPTION	This routine resets in turn: each attached physical device, each target, and the controller-finite-state machine. In practice, this routine implements the SCSI hard reset option.
NOTE	This routine does not physically reset the SCSI bus; see <b>scsiBusReset( )</b> . This routine should not be called by application programs.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	scsiMgrLib

### scsiMgrCtrlEvent()

NAME scsiMgrCtrlEvent() – send an event to the SCSI controller state machine

SYNOPSIS void scsiMgrCtrlEvent ( SCSI\_CTRL \* pScsiCtrl, SCSI\_EVENT\_TYPE eventType ) VxWorks Kernel API Reference, 6.6 scsiMgrEventNotify()

DESCRIPTION	This routine is called by the thread driver whenever selection, reselection, or disconnection occurs or when a thread is activated. It manages a simple finite-state machine for the SCSI controller.
NOTE	This function should not be called by application programs.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	scsiMgrLib

# scsiMgrEventNotify()

NAME	scsiMgrEventNotify() – notify the SCSI manager of a SCSI (controller) event	
SYNOPSIS	<pre>STATUS scsiMgrEventNotify ( SCSI_CTRL * pScsiCtrl, /* pointer to SCSI controller structure */ SCSI_EVENT * pEvent, /* pointer to the SCSI event */ int eventSize /* size of the event information */ )</pre>	
DESCRIPTION	This routine posts an event message on the appropriate SCSI manager queue, then notifies the SCSI manager that there is a message to be accepted.	
NOTE	This routine should not be called by application programs.	
	No access serialization is required, because event messages are only posted by the SCSI controller ISR. See the reference entry for <b>scsiBusResetNotify()</b> .	
RETURNS	OK, or ERROR if the SCSI manager's event queue is full.	
ERRNO	Not Available	
SEE ALSO	scsiMgrLib, scsiBusResetNotify()	

### scsiMgrShow()

NAME	scsiMgrShow() – show status information for the SCSI manager
SYNOPSIS	<pre>void scsiMgrShow   (    SCSI_CTRL * pScsiCtrl, /* SCSI controller to use */ BOOL showPhysDevs, /* TRUE =&gt; show phys dev details */ BOOL showThreads, /* TRUE =&gt; show thread details */ BOOL showFreeThreads /* TRUE =&gt; show free thread IDs */ )</pre>
DESCRIPTION	This routine shows the current state of the SCSI manager for the specified controller, including the total number of threads created and the number of threads currently free.
	Optionally, this routine also shows details for all created physical devices on this controller and all threads for which SCSI requests are outstanding. It also shows the IDs of all free threads.
NOTE	The information displayed is volatile; this routine is best used when there is no activity on the SCSI bus. Threads allocated by a client but for which there are no outstanding SCSI requests are not shown.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	scsiMgrLib

### scsiMgrThreadEvent()

 NAME
 scsiMgrThreadEvent() - send an event to the thread state machine

 SYNOPSIS
 void scsiMgrThreadEvent

 (
 SCSI\_THREAD \* pThread,

 SCSI\_THREAD\_EVENT\_TYPE
 eventType

 )
 This routine forwards an event to the thread's physical device. If the event is completion or deferral, it frees up the tag which was allocated when the thread was activated and either completes or defers the thread.

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NOTE	This function should not be called by application programs.
	The thread passed into this function does not have to be an active client thread (it may be an identification thread).
	If the thread has no corresponding physical device, this routine does nothing. (This occasionally occurs if an unexpected disconnection or bus reset happens when an identification thread has not yet identified which physical device it corresponds to.)
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	scsiMgrLib

### scsiModeSelect()

NAME scsiModeSelect() – issue a MODE\_SELECT command to a SCSI device

```
SYNOPSIS
                STATUS scsiModeSelect
                    (
                    SCSI_PHYS_DEV * pScsiPhysDev, /* ptr to SCSI physical device
                */
                                                    /* value of the page format bit (0-1)
                    int
                                     pageFormat,
                * /
                                                     /* value of the save parameters bit (0-1)
                    int
                                     saveParams,
                * /
                                     buffer,
                                                     /* ptr to output data buffer
                    char *
                */
                    int
                                     bufLength
                                                     /* length of buffer in bytes
                */
                    )
                This routine issues a MODE_SELECT command to a specified SCSI device.
DESCRIPTION
RETURNS
                OK, or ERROR if the command fails.
ERRNO
                Not Available
                scsiLib
SEE ALSO
```

scsiModeSense() - issue a MODE\_SENSE command to a SCSI device NAME SYNOPSIS STATUS scsiModeSense ( SCSI\_PHYS\_DEV \* pScsiPhysDev, /\* ptr to SCSI physical device \*/ pageControl, /\* value of the page control field (0-3) int \*/ /\* value of the page code field (0-0x3f) pageCode, int \*/ char \* buffer, /\* ptr to input data buffer \*/ bufLength /\* length of buffer in bytes \*/ int ) DESCRIPTION This routine issues a MODE\_SENSE command to a specified SCSI device. OK, or ERROR if the command fails. RETURNS Not Available ERRNO scsiLib SEE ALSO

### scsiMsgInComplete()

NAME	<pre>scsiMsgInComplete() - handle a complete SCSI message received from the target</pre>	
SYNOPSIS	STATUS scsiMsgInComplete ( SCSI_CTRL *pScsiCtrl, /* ptr to SCSI controller info */ SCSI_THREAD *pThread /* ptr to thread info */ )	
DESCRIPTION	This routine parses the complete message and takes any necessary action, which may include setting up an outgoing message in reply. If the message is not understood, the routine rejects it and returns an <b>ERROR</b> status.	
NOTE	This function is intended for use only by SCSI controller drivers.	
RETURNS	OK, or ERROR if the message is not supported.	
ERRNO	Not Available	

SEE ALSO scsi2Lib

### scsiMsgOutComplete()

NAME	<pre>scsiMsgOutComplete() - perform post-processing after a SCSI message is sent</pre>
SYNOPSIS	STATUS scsiMsgOutComplete ( SCSI_CTRL *pScsiCtrl, /* ptr to SCSI controller info */ SCSI_THREAD *pThread /* ptr to thread info */ )
DESCRIPTION	This routine parses the complete message and takes any necessary action.
NOTE	This function is intended for use only by SCSI controller drivers.
RETURNS	OK, or ERROR if the message is not supported.
ERRNO	Not Available
SEE ALSO	scsi2Lib

# scsiMsgOutReject()

NAME	<pre>scsiMsgOutReject() – perform post-processing when an outgoing message is rejected</pre>	
SYNOPSIS	<pre>void scsiMsgOutReject (     SCSI_CTRL *pScsiCtrl, /* ptr to SCSI controller info */     SCSI_THREAD *pThread /* ptr to thread info */ )</pre>	
NOTE	This function is intended for use only by SCSI controller drivers.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if the message is not supported.	
ERRNO	Not Available	
SEE ALSO	scsi2Lib	

# scsiPhysDevCreate()

NAME	scsiPhysDevCreate() – create a SCSI physical device structure		
SYNOPSIS	SCSI_PHYS_DEV * scsiPhysDevCreate (		
	<pre>SCSI_CTRL * pScsiCtrl, /* ptr to SCSI controller info */ int devBusId, /* device's SCSI bus ID */ int devLUN, /* device's logical unit number */ int reqSenseLength, /* length of REQUEST SENSE data dev returns */</pre>		
	<pre>int devType, /* type of SCSI device */ BOOL removable, /* whether medium is removable */ int numBlocks, /* number of blocks on device */ int blockSize /* size of a block in bytes */ )</pre>		
DESCRIPTION	This routine enables access to a SCSI device and must be the first routine invoked. It must be called once for each physical device on the SCSI bus.		
	If <i>reqSenseLength</i> is <b>NULL</b> (0), one or more <b>REQUEST_SENSE</b> commands are issued to the device to determine the number of bytes of sense data it typically returns. Note that if the device returns variable amounts of sense data depending on its state, you must consult the device manual to determine the maximum amount of sense data that can be returned.		
	If <i>devType</i> is NONE (-1), an INQUIRY command is issued to determine the device type; as an added benefit, it acquires the device's make and model number. The <b>scsiShow()</b> routine displays this information. Common values of <i>devType</i> can be found in <b>scsiLib.h</b> or in the SCSI specification.		
	If <i>numBlocks</i> or <i>blockSize</i> are specified as <b>NULL</b> (0), a <b>READ_CAPACITY</b> command is issued to determine those values. This occurs only for device types that support <b>READ_CAPACITY</b> .		
RETURNS	A pointer to the created <b>SCSI_PHYS_DEV</b> structure, or <b>NULL</b> if the routine is unable to create the physical-device structure.		
ERRNO	Not Available		
SEE ALSO	scsiLib		

# scsiPhysDevDelete()

NAME	<pre>scsiPhysDevDelete() – delete a SCSI physical-device structure</pre>
SYNOPSIS	STATUS scsiPhysDevDelete ( SCSI_PHYS_DEV *pScsiPhysDev /* ptr to SCSI physical device info */ )
DESCRIPTION	This routine deletes a specified SCSI physical-device structure.
RETURNS	<b>OK</b> , or <b>ERROR</b> if <b>pScsiPhysDev</b> is <b>NULL</b> or SCSI_BLK_DEVs have been created on the device.
ERRNO	Not Available
SEE ALSO	scsiLib

# scsiPhysDevIdGet()

NAME	<pre>scsiPhysDevIdGet() - return a pointer to a SCSI_PHYS_DEV structure</pre>
SYNOPSIS	SCSI_PHYS_DEV * scsiPhysDevIdGet
	SCSI_CTRL * pScsiCtrl, /* ptr to SCSI controller info */ int devBusId, /* device's SCSI bus ID */ int devLUN /* device's logical unit number */ )
DESCRIPTION	This routine returns a pointer to the <b>SCSI_PHYS_DEV</b> structure of the SCSI physical device located at a specified bus ID ( <i>devBusId</i> ) and logical unit number ( <i>devLUN</i> ) and attached to a specified SCSI controller ( <i>pScsiCtrl</i> ).
RETURNS	A pointer to the specified <b>SCSI_PHYS_DEV</b> structure, or <b>NULL</b> if the structure does not exist.
ERRNO	Not Available
SEE ALSO	scsiLib

NAME	<pre>scsiPhysDevShow() - show status information for a physical device</pre>	
SYNOPSIS	<pre>void scsiPhysDevShow   (    SCSI_PHYS_DEV * pScsiPhysDev, /* physical device to be displayed */   BOOL showThreads, /* show IDs of associated threads */   BOOL noHeader /* do not print title line */   )</pre>	
DESCRIPTION	This routine shows the state, the current nexus type, the current tag number, the number of tagged commands in progress, and the number of waiting and active threads for a SCSI physical device. Optionally, it shows the IDs of waiting and active threads, if any. This routine may be called at any time, but note that all of the information displayed is volatile.	
RETURNS	N/A	
ERRNO	Not Available	
SEE ALSO	scsi2Lib	

### scsiRdSecs()

NAME	<pre>scsiRdSecs() - read sector(s) from</pre>	n a SCSI block device
SYNOPSIS	STATUS scsiRdSecs ( SCSI_BLK_DEV * pScsiBlkD int sector, int numSecs, char * buffer )	ev, /* ptr to SCSI block device info */ /* sector number to be read */ /* total sectors to be read */ /* ptr to input data buffer */
DESCRIPTION	This routine reads the specified p	hysical sector(s) from a specified physical device
RETURNS	<b>OK</b> , or <b>ERROR</b> if the sector(s) can	not be read.
ERRNO	Not Available	

SEE ALSO scsiLib

	scsiRdTape()		
NAME	<b>scsiRdTape( )</b> – read bytes or blocks fro	om a SCSI tape device	
SYNOPSIS	<pre>int scsiRdTape   (    SCSI_SEQ_DEV *pScsiSeqDev, /    UINT count, /    char *buffer, /    BOOL fixedSize /   )</pre>	* ptr to SCSI sequential device info * total bytes or blocks to be read * ptr to input data buffer * if variable size blocks	*/ */ */
DESCRIPTION	This routine reads the specified number of bytes or blocks from a specified physical device. If the boolean <i>fixedSize</i> is true, then <i>numBytes</i> represents the number of blocks of size <i>blockSize</i> , defined in the <b>pScsiPhysDev</b> structure. If variable block sizes are used ( <i>fixedSiz</i> = <b>FALSE</b> ), then <i>numBytes</i> represents the actual number of bytes to be read.		l device. ize (fixedSize
RETURNS	Number of bytes or blocks actually rea	d, 0 if EOF, or ERROR.	
ERRNO	Not Available		
SEE ALSO	scsiSeqLib		

# scsiReadCapacity()

NAME	<pre>scsiReadCapacity() – issue a READ_CAPACITY command to a SCSI device</pre>
SYNOPSIS	<pre>STATUS scsiReadCapacity (     SCSI_PHYS_DEV * pScsiPhysDev, /* ptr to SCSI physical device */     int *    pLastLBA, /* where to return last */</pre>
DESCRIPTION	This routine issues a <b>READ_CAPACITY</b> command to a specified SCSI device.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the command fails.
ERRNO	Not Available
SEE ALSO	scsiLib

#### scsiRelease()

NAME	scsiRelease() – issue a RELEASE command to a SCSI device
SYNOPSIS	STATUS scsiRelease ( SCSI_PHYS_DEV *pScsiPhysDev /* ptr to SCSI physical device */ )
DESCRIPTION	This routine issues a RELEASE command to a specified SCSI device.
RETURNS	OK, or ERROR if the command fails.
ERRNO	Not Available
SEE ALSO	scsiDirectLib

#### scsiReleaseUnit()

NAMEscsiReleaseUnit() – issue a RELEASE UNIT command to a SCSI deviceSYNOPSISSTATUS scsiReleaseUnit<br/>(<br/>SCSI\_SEQ\_DEV \*pScsiSeqDev /\* ptr to SCSI sequential device \*/<br/>)DESCRIPTIONThis routine issues a RELEASE UNIT command to a specified SCSI device.RETURNSOK, or ERROR if the command fails.ERRNONot AvailableSEE ALSOscsiSeqLib

### scsiReqSense()

NAME scsiReqSense() – issue a REQUEST\_SENSE command to a SCSI device and read results

SYNOPSIS STATUS scsiReqSense

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(
SCSI\_PHYS\_DEV \* pScsiPhysDev, /\* ptr to SCSI physical device \*/
char \* buffer, /\* ptr to input data buffer \*/
int bufLength /\* length of buffer in bytes \*/
)

**DESCRIPTION** This routine issues a **REQUEST\_SENSE** command to a specified SCSI device and reads the results.

ERRNO Not Available

SEE ALSO scsiLib

#### scsiReserve()

NAME	<b>scsiReserve()</b> – issue a RESERVE command to a SCSI device
SYNOPSIS	STATUS scsiReserve ( SCSI_PHYS_DEV *pScsiPhysDev /* ptr to SCSI physical device */ )
DESCRIPTION	This routine issues a RESERVE command to a specified SCSI device.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the command fails.
ERRNO	Not Available
SEE ALSO	scsiDirectLib

### scsiReserveUnit()

NAME	<pre>scsiReserveUnit() - issue a RESERVE UNIT command to a SCSI device</pre>		
SYNOPSIS	STATUS scsiReserveUnit ( SCSI_SEQ_DEV *pScsiSeqDev /* ptr to SCSI sequential device */ )		

DESCRIPTIONThis routine issues a RESERVE UNIT command to a specified SCSI device.RETURNSOK, or ERROR if the command fails.ERRNONot AvailableSEE ALSOscsiSeqLib

#### scsiRewind()

NAMEscsiRewind() - issue a REWIND command to a SCSI deviceSYNOPSISSTATUS scsiRewind<br/>(<br/>SCSI\_SEQ\_DEV \*pScsiSeqDev /\* ptr to SCSI Sequential device \*/<br/>)DESCRIPTIONThis routine issues a REWIND command to a specified SCSI device.RETURNSOK, or ERROR if the command fails.ERRNONot AvailableSEE ALSOscsiSeqLib

### scsiSeqDevCreate()

NAME	<pre>scsiSeqDevCreate() - create</pre>	ate a SCSI sequential device
SYNOPSIS	SEQ_DEV *scsiSeqDevCrea ( SCSI_PHYS_DEV *pSc: )	ate siPhysDev /* ptr to SCSI physical device info */
DESCRIPTION	This routine creates a SCS SCSI physical device. The	I sequential device and saves a pointer to this <b>SEQ_DEV</b> in the following functions are initialized in this structure:
	sd_seqRd sd_seqWrt sd_ioctl	scsiRdTape() scsiWrtTape() scsiIoctl() (in scsiLib)

VxWorks Kernel API Reference, 6.6 scsiSeqloctl( )

	sd_seqWrtFileMarks	scsiWrtFileMarks()
	sd_statusChk	scsiSeqStatusCheck()
	sd_reset	(not used)
	sd_rewind	scsiRewind()
	sd_reserve	scsiReserve()
	sd_release	scsiRelease()
	sd_readBlkLim	scsiSeqReadBlockLimits()
	sd_load	scsiLoadUnit( )
	sd_space	scsiSpace()
	sd_erase	scsiErase()
	Only one <b>SEQ_DEV</b> per <b>SCS</b> is maintained. Therefore, the device.	<b>SI_PHYS_DEV</b> is allowed, unlike BLK_DEVs where an entire list nis routine can be called only once per creation of a sequential
RETURNS	A pointer to the SEQ_DEV structure, or NULL if the command fails.	
ERRNO	Not Available	
SEE ALSO	scsiSeqLib	

# scsiSeqIoctl()

NAME	<b>scsiSeqIoctl()</b> – perform an I/O control function for sequential access devices		
SYNOPSIS	<pre>int scsiSeqIoctl   (     SCSI_SEQ_DEV * pScsiSeqDev, /* ptr to SCSI sequential device */     int function, /* ioctl function code */     int arg /* argument to pass to called function */     )</pre>		
DESCRIPTION	This routine issues <b>scsiSeqLib</b> commands to perform sequential device-specific I/O control operations.		
RETURNS	OK or ERROR.		
ERRNO	S_scsiLib_INVALID_BLOCK_SIZE		
SEE ALSO	scsiSeqLib		

#### scsiSeqReadBlockLimits() scsiSeqReadBlockLimits() - issue a READ\_BLOCK\_LIMITS command to a SCSI device NAME SYNOPSIS STATUS scsiSeqReadBlockLimits SCSI\_SEQ\_DEV \* pScsiSeqDev, /\* ptr to SCSI sequential device \*/ int \*pMaxBlockLength, /\* where to return maximum block length \*/ \*pMinBlockLength /\* where to return minimum block length UINT16 \*/ ) DESCRIPTION This routine issues a READ\_BLOCK\_LIMITS command to a specified SCSI device. RETURNS OK, or ERROR if the command fails. ERRNO Not Available

SEE ALSO scsiSeqLib

### scsiSeqStatusCheck()

NAME	<pre>scsiSeqStatusCheck() – detect a change in media</pre>
SYNOPSIS	STATUS scsiSeqStatusCheck ( SCSI_SEQ_DEV *pScsiSeqDev /* ptr to a sequential dev */ )
DESCRIPTION	This routine issues a <b>TEST_UNIT_READY</b> command to a SCSI device to detect a change in media. It is called by file systems before executing <b>open()</b> or <b>creat()</b> .
RETURNS	OK or ERROR.
ERRNO	Not Available
SEE ALSO	scsiSeqLib

### scsiShow()

NAME	<b>scsiShow()</b> – list the physical devices attached to a SCSI controller		
SYNOPSIS	STATUS scsiShow ( SCSI_CTRL *pScsiCtrl /* ptr to SCSI controller info */ )		
DESCRIPTION	This routine displays the SCSI bus ID, logical unit number (LUN), vendor ID, product ID, firmware revision (rev.), device type, number of blocks, block size in bytes, and a pointer to the associated <b>SCSI_PHYS_DEV</b> structure for each physical SCSI device known to be attached to a specified SCSI controller.		
NOTE	If <i>pScsiCtrl</i> is <b>NULL</b> , the value of the global variable <b>pSysScsiCtrl</b> is used, unless it is also <b>NULL</b> .		
RETURNS	<b>OK</b> , or <b>ERROR</b> if both <i>pScsiCtrl</i> and <b>pSysScsiCtrl</b> are <b>NULL</b> .		
ERRNO	Not Available		
SEE ALSO	scsiLib		

### scsiSpace()

NAME	scsiSpace() – move the tape on a specified physical SCSI device		
SYNOPSIS	STATUS scsiSpace ( SCSI_SEQ_DEV * pScsiSeqDev, int count, int spaceCode	/* ptr to SCSI sequential device info */ /* count for space command */ /* code for the type of space command */	
	)		

**DESCRIPTION** This routine moves the tape on a specified SCSI physical device. There are two types of space code that are mandatory in SCSI; currently these are the only two supported:

Code	Description	Support	
000	Blocks	Yes	
001	File marks	Yes	
010	Sequential file marks	No	
011	End-of-data	No	
100	Set marks	No	

	Code	Description	Support
	101	Sequential set marks	No
RETURNS	<b>OK</b> , or <b>ERROR</b> if an error is returned by the device.		
ERRNO	S_scsiLib_ILLEGAL_REQUEST		
SEE ALSO	scsiSeqLib		

### scsiStartStopUnit()

NAME	<pre>scsiStartStopUnit() - issue a START_STOP_UNIT command to a SCSI device</pre>		
SYNOPSIS	<pre>STATUS scsiStartStopUnit   (    SCSI_PHYS_DEV *pScsiPhysDev, /* ptr to SCSI physical device */    BOOL start /* TRUE == start, FALSE == stop */   )</pre>		
DESCRIPTION	This routine issues a <b>START_STOP_UNIT</b> command to a specified SCSI device.		
RETURNS	<b>OK</b> , or <b>ERROR</b> if the command fails.		
ERRNO	Not Available		
SEE ALSO	scsiDirectLib		

### scsiSyncXferNegotiate()

NAME scsiSyncXferNegotiate() – initiate or continue negotiating transfer parameters

 SYNOPSIS
 void scsiSyncXferNegotiate

 (
 SCSI\_CTRL
 \*pScsiCtrl,
 /\* ptr to SCSI controller info \*/

 SCSI\_TARGET
 \*pScsiTarget,
 /\* ptr to SCSI target info
 \*/

 SCSI\_SYNC\_XFER\_EVENT
 eventType
 /\* tells what has just happened \*/

 )
 )

VxWorks Kernel API Reference, 6.6 scsiTapeModeSelect()

DESCRIPTION	This routine manages negotiation by means of a finite-state machine which is driven by "significant events" such as incoming and outgoing messages. Each SCSI target has its own independent state machine.
NOTE	If the controller does not support synchronous transfer or if the target's maximum REQ/ACK offset is zero, attempts to initiate a round of negotiation are ignored.
	This function is intended for use only by SCSI controller drivers.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	scsi2Lib

# scsiTapeModeSelect()

NAME	<pre>scsiTapeModeSelect() - issue a MODE_SELECT command to a SCSI tape device</pre>		
SYNOPSIS	<pre>STATUS scsiTapeMod     (         SCSI_PHYS_DEV */         int */         int */         char */         int */ */         int */</pre>	<pre>deSelect  *pScsiPhysDev,  pageFormat,  saveParams,  *buffer,  bufLength</pre>	<pre>/* ptr to SCSI physical device /* value of the page format bit (0-1) /* value of the save parameters bit (0-1) /* ptr to output data buffer /* length of buffer in bytes</pre>
DESCRIPTION	This routine issues a <b>MODE_SELECT</b> command to a specified SCSI device.		
RETURNS	<b>OK</b> , or <b>ERROR</b> if the command fails.		
ERRNO	Not Available		
SEE ALSO	scsiSeqLib		

NAME	scsiTapeModeSense	e( ) – issue a MODE	_SENSE command to a SCSI tape device	
SYNOPSIS	STATUS scsiTapeMod ( SCSI_PHYS_DEV int int char int )	deSense *pScsiPhysDev, pageControl, pageCode, *buffer, bufLength	<pre>/* ptr to SCSI physical device /* value of the page control field (0-3) /* value of the page code field (0-0x3f) /* ptr to input data buffer /* length of buffer in bytes</pre>	*/ */ */
DESCRIPTION	This routine issues a	MODE_SENSE cor	nmand to a specified SCSI tape device.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if the command fails.			
ERRNO	Not Available			
SEE ALSO	scsiSeqLib			

### scsiTargetOptionsGet()

NAME	<pre>scsiTargetOptionsGet() - get options for one or all SCSI targets</pre>		
SYNOPSIS	<pre>STATUS scsiTargetOptionsGet   (    SCSI_CTRL *pScsiCtrl, /* ptr to SCSI controller info */    int devBusId, /* target to interrogate */    SCSI_OPTIONS *pOptions /* buffer to return options */   )</pre>		
DESCRIPTION	This routine copies the current options for the specified target into the caller's buffer.		
RETURNS	<b>OK</b> , or <b>ERROR</b> if the bus ID is invalid.		
ERRNO	Not Available		
SEE ALSO	scsi2Lib		

	scsiTargetOptions	Set()		
NAME	<pre>scsiTargetOptionsSet( ) - set options for one or all SCSI targets</pre>			
SYNOPSIS	<pre>STATUS scsiTargetOptionsSet   (     SCSI_CTRL *pScsiCtrl,     int devBusId,     SCSI_OPTIONS *pOptions,     UINT which   )</pre>	/* ptr to SCSJ /* target to a /* buffer cont /* which optic	controller info */ affect, or all */ caining new options */ ons to change */	
DESCRIPTION	<ul> <li>This routine sets the options defined by the bitmask which for the specified target (or all targets if devBusId is SCSI_SET_OPT_ALL_TARGETS).</li> <li>The bitmask which can be any combination of the following, bitwise OR'd together (corresponding fields in the SCSI_OPTIONS structure are shown in parentheses):</li> </ul>			
	SCSI_SET_OPT_TIMEOUT SCSI_SET_OPT_MESSAGES SCSI_SET_OPT_DISCONNECT SCSI_SET_OPT_XFER_PARAMS	selTimeOut messages disconnect maxOffset, minPoriod	select timeout period, microseconds FALSE to disable SCSI messages FALSE to disable discon/recon max sync xfer offset, 0=>async min sync xfer period x 4 psoc	
	SCSI_SET_OPT_TAG_PARAMS	tagType, maxTags	default tag type (SCSI_TAG_*) max cmd tags available	
	SCSI_SEI_OPI_WIDE_PAKAMS	xferWidth	xferWidth = 0 ; 8 bits wide xferWidth = 1 ; 16 bits wide	
NOTE	This routine can be used after the target device has already been used; in this case, however, it is not possible to change the tag parameters. This routine must not be used while there is any SCSI activity on the specified target(s).			
RETURNS	OK, or ERROR if the bus ID or options are invalid.			
ERRNO	Not Available			
SEE ALSO	scsi2Lib			
	scsiTargetOptionsShow()			
-------------	---	--	--	
NAME	<pre>scsiTargetOptionsShow() – display options for specified SCSI target</pre>			
SYNOPSIS	<pre>STATUS scsiTargetOptionsShow   (    SCSI_CTRL *pScsiCtrl, /* ptr to SCSI controller info */    int devBusId /* target to interrogate */   )</pre>			
DESCRIPTION	This routine displays the current target options for the specified target in the following format:			
	<pre>Target Options (id <scsi bus="" id="">): selection TimeOut: <timeout> nano secs messages allowed: TRUE or FALSE disconnect allowed: TRUE or FALSE REQ/ACK offset: <negotiated offset=""> transfer period: <negotiated period=""> transfer width: 8 or 16 bits maximum transfer rate: <peak rate="" transfer=""> MB/sec tag type: <tag type=""> maximum tags: <max tags=""></max></tag></peak></negotiated></negotiated></timeout></scsi></pre>			
RETURNS	OK, or ERROR if the bus ID is invalid.			
ERRNO	Not Available			

SEE ALSO scsi2Lib

# scsiTestUnitRdy()

NAME	<pre>scsiTestUnitRdy() - issue a TEST_UNIT_READY command to a SCSI device</pre>
SYNOPSIS	STATUS scsiTestUnitRdy ( SCSI_PHYS_DEV * pScsiPhysDev /* ptr to SCSI physical device */ )
DESCRIPTION	This routine issues a <b>TEST_UNIT_READY</b> command to a specified SCSI device.
RETURNS	OK, or ERROR if the command fails.
ERRNO	Not Available

VxWorks Kernel API Reference, 6.6 scsiThreadInit()

SEE ALSO scsiLib

## scsiThreadInit( )

NAME	scsiThreadInit() – perform generic SCSI thread initialization
SYNOPSIS	STATUS scsiThreadInit ( SCSI_THREAD * pThread )
DESCRIPTION	This routine initializes the controller-independent parts of a thread structure, which are specific to the SCSI manager.
NOTE	This function should not be called by application programs. It is intended to be used by SCSI controller drivers.
RETURNS	OK, or ERROR if the thread cannot be initialized.
ERRNO	Not Available
SEE ALSO	scsi2Lib

# scsiWideXferNegotiate()

NAME	scsiWideXferNegotiate() - in	nitiate or continue	e negotiating wide parameters
SYNOPSIS	<pre>void scsiWideXferNegotiat   (    SCSI_CTRL    SCSI_TARGET    SCSI_WIDE_XFER_EVENT  )</pre>	e *pScsiCtrl, *pScsiTarget, eventType	/* ptr to SCSI controller info */ /* ptr to SCSI target info */ /* tells what has just happened */
DESCRIPTION	This routine manages negotia	ation means of a f	inite-state machine which is driven by

**DESCRIPTION** This routine manages negotiation means of a finite-state machine which is driven by "significant events" such as incoming and outgoing messages. Each SCSI target has its own independent state machine.

NOTE	If the controller does not support wide transfers or the target's transfer width is zero, attempts to initiate a round of negotiation are ignored; this is because zero is the default narrow transfer.
	This function is intended for use only by SCSI controller drivers.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	scsi2Lib

### scsiWrtFileMarks()

**NAME** scsiWrtFileMarks() – write file marks to a SCSI sequential device

SYNOPSIS STATUS scsiWrtFileMarks

(				
SCSI_SEQ_DEV *	pScsiSeqDev,	/*	ptr to SCSI sequential device info	*/
int	numMarks,	/*	number of file marks to write	*/
BOOL	shortMark	/*	TRUE to write short file mark	*/
)				

- **DESCRIPTION** This routine writes file marks to a specified physical device.
- **RETURNS** OK, or ERROR if the file mark cannot be written.
- ERRNO Not Available
- SEE ALSO scsiSeqLib

### scsiWrtSecs()

NAME scsiWrtSecs() - write sector(s) to a SCSI block device
SYNOPSIS STATUS scsiWrtSecs
(
SCSI\_BLK\_DEV \* pScsiBlkDev, /\* ptr to SCSI block device info \*/
int sector, /\* sector number to be written \*/
int numSecs, /\* total sectors to be written \*/

SEE ALSO	scsiLib			
ERRNO	Not Available			
RETURNS	OK, or ERROR if the s	sector(s) cannot l	be written.	
DESCRIPTION	This routine writes the	ne specified phys	sical sector(s) to a specified physical device	e.
	char * )	buffer	/* ptr to input data buffer */	

# scsiWrtTape()

NAME	scsiWrtTape() – wi	rite data to a SCS	I tape device	
SYNOPSIS	STATUS scsiWrtTag ( SCSI_SEQ_DEV int char BOOL )	pe *pScsiSeqDev, numBytes, *buffer, fixedSize	/* ptr to SCSI sequential device info /* total bytes or blocks to be writter /* ptr to input data buffer /* if variable size blocks	*/ 1 */ */ */
DESCRIPTION	This routine writes data to the current block on a specified physical device. If the boolean <i>fixedSize</i> is true, then <i>numBytes</i> represents the number of blocks of size <i>blockSize</i> , defined in the <b>pScsiPhysDev</b> structure. If variable block sizes are used ( <i>fixedSize</i> = <b>FALSE</b> ), then <i>numBytes</i> represents the actual number of bytes to be written. If <i>numBytes</i> is greater than the <b>maxBytesLimit</b> field defined in the <b>pScsiPhysDev</b> structure, then more than one SCSI transaction is used to transfer the data.			
RETURNS	OK, or ERROR if the	e data cannot be	written or zero bytes are written.	
ERRNO	Not Available			
SEE ALSO	scsiSeqLib			

### sdCreate()

NAME

sdCreate() – create a new shared data region

- **DESCRIPTION** This routine creates a new shared data region and maps it into the calling task's memory context. The following table shows each parameter and whether it is required or not:

Parameter	<b>Required?</b>	Default
name	Yes	N/A
options	No	0
size	Yes	N/A
physAddress	No	System Allocated
attr	No	Read/Write, System Default Cache Setting
pVirtAddress	Yes	N/A

Because each shared data region must have a unique name, if the region specified by *name* already exists in the system the creation will fail. **NULL** will be returned.

Currently there are only two possible values of options:

Option name	Value	Meaning
SD_LINGER	0x1	SD region may remain after the last client unmaps.
SD_PRIVATE	0x2	SD region is only available in the owner RTP.

The value of *size* must be greater than 0. It is rounded up to a page aligned size determined by the architecture.

If *physAddress* is specified and the address is not available, **NULL** will be returned. The *physAddress* specified must be aligned on the architecture dependent page size boundary and must not be mapped to any other memory context.

The MMU attributes specified in *attr* will be used as the default attributes of the shared data region. All client applications will use these by default, and may only change the local access permissions to a subset of these. The application which creates the region will have read and write access in addition to the defaults and will be allowed to set local permissions to any allowed by the architecture.

Basic MMU attribute definitions for shared data regions are provided in the **sdLibCommon.h** header file. These include:

Attribute	Meaning
SD_ATTR_RW	Read/Write for both Supervisor and User Modes
SD_ATTR_RO	Read Only for both Supervisor and User Modes
SD_ATTR_RWX	Read/Write/Execute for both Supervisor and User
SD_ATTR_RX SD_CACHE_COPYBACK SD_CACHE_WRITETHROUGH SD_CACHE_OFF	Modes Read/Execute for both Supervisor and User Modes Copyback cache mode Write through cache mode Cache Off
One of each the <b>SD_ATTR</b> and <b>SD_</b> macros can not be combined.	CACHE macros above must be provided. The SD_CACHE
The cache attributes of a shared d of that region will use the value p	ata region can not be changed after creation. All clients rovided at create time, including the owner.
If more specific MMU attributes a list of available MMU attributes.	re required please see <b>vmLibCommon.h</b> for a complete
The MMU_ATTR mask used inter	nally by the shared data library is the combination of:
MMU_ATTR_PROT_MASK	
MMU_ATTR_VALID_MSK	
MMU_ATTR_SPL_MSK	
Care must be taken to provide sur	table values for all these attributes.
The start address of the shared da <i>pVirtAddress</i> . This must be a valid can not be <b>NULL</b> .	ta region is stored at the location specified by address within the context of the calling application. It
The <b>SD_ID</b> returned is private to t within that application but not wi	he calling application. It can be shared between tasks that reside outside that application.
ID of new shared data region, or I	NULL on error.
Possible errno values set by this r	outine are:
S_sdLib_VIRT_ADDR_PTR_IS_NU pVirtAddress is NULL	LL
<b>S_sdLib_ADDR_NOT_ALIGNED</b> <i>physAddress</i> is not properly a	ligned
S_sdLib_PHYS_ADDR_OUT_OF_R physAddress exceeds physical	ANGE address space
S_sdLib_SIZE_IS_NULL size is NULL	

NOTE

RETURNS

ERRNO

- S\_sdLib\_VIRT\_PAGES\_NOT\_AVAILABLE not enough virtual space left in system
- S\_sdLib\_PHYS\_PAGES\_NOT\_AVAILABLE not enough physical memory left in system

SEE ALSO sdLib, sdOpen(), sdUnmap(), sdProtect(), sdDelete()

### sdCreateHookAdd()

sdCreateHookAdd() – add a hook routine to be called at Shared Data creation NAME SYNOPSIS STATUS sdCreateHookAdd ( SD\_CREATE\_HOOK sdCreateHook, /\* hook routine to call \*/ BOOL addToHead /\* add routine to head of list \*/ ) DESCRIPTION This routine adds a specified routine to a list of routines that will be called just after an SD is created. The hook routine should have the following prototype: STATUS sdCreateHook ( const SD\_ID sdId, /\* ID of the created SD \*/ ) The second parameter *addToHead* specifies the order in which the hook is added to the table. If FALSE, the hook is appended to the list of hooks already installed. If addToHead is TRUE, the new hook is added to the head of the list (i.e. it will be the first hook to execute). Shared Data create hooks are called from **sdCreate()** or **sdOpen()** after the creation is done. Create hooks are not expected to return anything (return values if any are not checked). OK, or ERROR if the table of SD create routines is full. RETURNS N/A. ERRNO SEE ALSO sdLib, sdCreateHookDelete()

## sdCreateHookDelete()

NAME	sdCreateHookDelete() – delete a Shared Data creation hook routine
SYNOPSIS	STATUS sdCreateHookDelete ( SD_CREATE_HOOK sdCreateHook /* hook routine to delete */ )
DESCRIPTION	This routine removes a specified hook routine from the list of Shared Data create hook routines.
RETURNS	OK on success, or ERROR if the hook routine was not found.
ERRNO	S_hookLib_HOOK_NOT_FOUND
SEE ALSO	sdLib, sdCreateHookAdd()

# sdDelete()

NAME	sdDelete() – delete a shared data region				
SYNOPSIS	<pre>STATUS sdDelete   (    SD_ID sdId, /* ID of shared data region to delete */    int options /* options field is not used */   )</pre>				
DESCRIPTION	Deletes a shared data region. This is only possible if there are no applications that have the shared data region mapped. Currently there are no options defined for this function, this parameter should be passed as zero always.				
	Unless the option <b>SD_LINGER</b> was specified at creation of the shared data region it will automatically be deleted when the last client application exits or explicitly calls <b>sdUnmap()</b> .				
RETURNS	OK, or ERROR on failure.				
ERRNO	Possible errno values set by this routine are:				
	S_sdLib_INVALID_SD_ID sdId is not valid				
	S_sdLib_CLIENT_COUNT_NOT_NULL sdId still mapped by an application				

## sdDeleteHookAdd()

NAME	sdDeleteHookAdd() – add a hook routine to be called at Shared Data deletion				
SYNOPSIS	STATUS sdDeleteHookAdd ( SD_DELETE_HOOK sdDeleteHook, /* hook routine to call */ BOOL addToHead /* add routine to head of list */ )				
DESCRIPTION	This routine adds a specified routine to a list of routines that will be called just before a SD is deleted. The hook routine should have the following prototype:				
	void sdDeleteHook ( const SD_ID sdId, /* ID of the deleted SD */ )				
	The second parameter <i>addToHead</i> specifies the order in which the hook is added to the table. If <b>FALSE</b> , the hook is appended to the list of hooks already installed. If addToHead is <b>TRUE</b> , the new hook is added to the head of the list (i.e. it will be the first hook to execute).				
	SD delete hooks are called from <b>sdDelete()</b> before any deletion is done. Delete hooks are not expected to return anything (return values if any are not checked).				
RETURNS	<b>OK</b> , or <b>ERROR</b> if the table of SD create routines is full.				
ERRNO	N/A.				
SEE ALSO	sdLib, sdDeleteHookDelete()				

## sdDeleteHookDelete()

NAME	<pre>sdDeleteHookDelete() – delete a Shared Data deletion hook routine</pre>					
SYNOPSIS	<pre>STATUS sdDeleteHookDelete   (    SD_DELETE_HOOK sdDeleteHook /* hook routine to delete */ )</pre>					

VxWorks Kernel API Reference, 6.6 sdGenericHookAdd()

DESCRIPTION	This routine removes a specified hook routine from the list of Shared Data delete hook routines.
RETURNS	OK on success, or ERROR if the hook routine was not found.
ERRNO	S_hookLib_HOOK_NOT_FOUND
SEE ALSO	sdLib, sdDeleteHookAdd()

### sdGenericHookAdd()

**NAME sdGenericHookAdd()** – add a hook routine to be called before Shared Data routine

SYNOPSIS STATUS sdGenericHookAdd

```
(
SD_GENERIC_HOOK sdGenericHook, /* hook routine to call */
BOOL addToHead /* add routine to head of list */
)
```

**DESCRIPTION** This routine adds a specified routine to a list of routines that will be called just before an SD is created, mapped, unmapped, deleted, or has its protection attributes changed. The hook routine should have the following prototype:

The *options* argument is used to identify what routine invoked the hook and whether the first argument is to be treated as a name or an ID. These are specified by the following enumeration:

```
typedef enum sd_routines
{
    SD_HOOK_TYPE_MSK = 0x0000001,
    SD_HOOK_ID = 0x0000000,
    SD_HOOK_NAME = 0x00000001,
    SD_HOOK_ROUTINE_MSK = 0x00000002,
    SD_HOOK_CREATE = 0x00000004,
    SD_HOOK_OPEN = 0x00000006,
    SD_HOOK_MAP = 0x00000008,
    SD_HOOK_WAPP = 0x00000008,
    SD_HOOK_PROTECT = 0x0000000c
    SD_HOOK_OPTIONS;
```

	Only <b>sdCreate()</b> and <b>sdOpen()</b> invoke the generic hook with the <b>SD_HOOK_NAME</b> option specified.			
	The second parameter <i>addToHead</i> specifies the order in which the hook is added to the table. If <b>FALSE</b> , the hook is appended to the list of hooks already installed. If addToHead is <b>TRUE</b> , the new hook is added to the head of the list (i.e. it will be the first hook to execute).			
Shared Data generic hooks are called from <b>sdCreate()</b> , <b>sdOpen()</b> , <b>sdMap()</b> , <b>s</b> and <b>sdProtect()</b> and should return either <b>OK</b> or <b>ERROR</b> . If the return value from hook is anything other than <b>OK</b> the operation is aborted and the routine from invoked returns <b>ERROR</b> .				
RETURNS	<b>OK</b> , or <b>ERROR</b> if the table of SD create routines is full.			
ERRNO	N/A.			
SEE ALSO	sdLib, sdGenericHookDelete()			

# sdGenericHookDelete()

NAME	sdGenericHookDelete() – delete a Shared Data generic hook routine				
SYNOPSIS	STATUS sdGenericHookDelete ( SD_GENERIC_HOOK sdGenericHook /* hook routine to delete */ )				
DESCRIPTION	This routine removes a specified hook routine from the list of Shared Data generic hook routines.				
RETURNS	OK on success, or ERROR if the hook routine was not found.				
ERRNO	S_hookLib_HOOK_NOT_FOUND				
SEE ALSO	sdLib, sdGenericHookAdd( )				

## sdInfoGet()

sdInfoGet() – get specific information about a Shared Data Region NAME SYNOPSIS STATUS sdInfoGet ( SD ID sdId, /\* SD ID to get info \*/ SD\_DESC \* pSdStruct /\* location to store SD info \*/ ) DESCRIPTION This routine obtains the information for a Shared Data region and stores the information in the specified SD descriptor (sdStruct). The information stored in the descriptor is copied from information in the SD object. The descriptor must have been allocated before calling this function, and the memory for it must come from the calling task's RTP space. To allocate the memory for the descriptor from the calling task's RTP space, either use malloc() within the calling task or declare the structure as an automatic variable in the calling task, placing it on the calling task's stack. If the name of the Shared Data region is longer than VX\_SD\_NAME\_LENGTH characters it will be truncated. The sdStruct structure looks like the following: typedef struct { char name[VX\_SD\_NAME\_LENGTH+1]; // name of SD int options; // options, e.g. SD\_LINGER, SD\_PRIVATE MMU ATTR defaultAttr; // default attributes of SD MMU\_ATTR currentAttr; // current attributes of SD // size of SD in bytes UINT size; VIRT\_ADDR startAddr // start address of SD } SD\_DESC; See the header file **vmLibCommon.h** for definitions of the values returned in *defaultAttr* and currentAttr. RETURNS OK, or ERROR on failure. ERRNO Possible errno values set by this routine are: S\_sdLib\_INVALID\_SD\_ID sdId is not valid SEE ALSO sdLib, sdCreate(), sdOpen(), sdMap(), sdUnmap(), sdProtect(), sdDelete()

### sdMap() sdMap() – map a shared data region into an application or the kernel NAME SYNOPSIS VIRT\_ADDR sdMap ( SD\_ID sdId, /\* ID of shared data region to map \*/ MMU\_ATTR attr, /\* MMU attr used to map region \*/ int options /\* reserved - use zero \*/ ) DESCRIPTION This routine maps the shared data region specified by *sdld* into the current calling task's memory context. The region is then available to all tasks within that application, or all tasks in the kernel if the calling task was a kernel task. The shared data region is mapped using the MMU attributes specified by *attr*. These attributes must be equal to, or a subset of the default attributes of *sdId*. If 0 was passed then the default attributes of *sdId* are used. It is possible to use this routine to set the attributes on a shared data region for the calling task's RTP even if *sdId* is currently mapped in its memory context. Basic MMU attribute definitions for shared data regions are provided in the sdLibCommon.h header file. These include: Attribute Meaning Read/Write for both Supervisor and User Modes SD ATTR RW SD\_ATTR\_RO Read Only for both Supervisor and User Modes SD\_ATTR\_RWX Read/Write/Execute for both Supervisor and User Modes SD ATTR RX Read/Execute for both Supervisor and User Modes If more specific MMU attributes are required please see **vmLibCommon.h** for a complete list of available MMU attributes. The MMU\_ATTR mask used internally by the shared data library is the combination of: NOTE MMU\_ATTR\_PROT\_MASK MMU\_ATTR\_VALID\_MSK MMU\_ATTR\_SPL\_MSK Care must be taken to provide suitable values for all these attributes. There are currently no options specified for this function, zero should be passed in the options parameter.

### **RETURNS** The base virtual address of the shared data region, or **NULL** on failure.

**ERRNO** Possible errno values set by this routine are:

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VxWorks Kernel API Reference, 6.6 sdOpen()

S\_sdLib\_INVALID\_SD\_ID sdId is not valid

**S\_sdLib\_SD\_IS\_PRIVATE** *sdId* is private to another application

SEE ALSO sdLib, sdCreate(), sdOpen(), sdUnmap(), sdProtect(), sdDelete()

# sdOpen()

NAME

**sdOpen()** – open a shared data region for use

**DESCRIPTION** This routine takes a shared data region name and looks for the region in the system. If the region does not exist in the system, and the **OM\_CREATE** flag is specified in *mode*, then a new shared data region is created and mapped to the application. If *mode* does not specify **OM\_CREATE** then no shared data region is created and NULL is returned. If the region does already exist in the system it is mapped into the calling task's memory context.

The following table shows each parameter and whether it is required or not:

Parameter	<b>Required?</b>	Default		
name	Yes	N/A		
options	No	0		
mode	No	0		
size	Yes	N/A		
physAddress	No	System Allocated		
attr	No	Read/Write, System Default Cache Setting		
pVirtAddress	Yes	N/A		

If the region specified by *name* already exists in the system all other arguments, except *pVirtAddress* and *attr* (if specified) will be ignored. In this case the region will be mapped into the calling task's memory context and the start address of the region will still be stored at *pVirtAddress* and the **SD\_ID** of the region will be returned.

Currently there are only two possible values of options:

Option name	Value	Meaning
SD_LINGER	0x1	SD region may remain after the last client unmaps.
SD_PRIVATE	0x2	SD region is only available in the owner RTP.

Currently there are only two possible values of *mode* other than the default (0):

Mode	Meaning			
DEFAULT (0)	Do not create an SD region if a matching name was not found.			
OM_CREATE	Create a shared data region if a matching name was not found.			
OM_EXCL	When set jointly with OM_CREATE, create a new shared data region			
	immediately without attempting to open an existing shared data			
	region. An error condition is returned if a shared data region with <i>name</i>			
	already exists. This attribute has no effect if the OM_CREATE attribute			
	is not specified.			

The value of *size* must be greater than 0. It is rounded up to a page aligned size determined by the architecture.

If *physAddress* is specified and the address is not available, **NULL** will be returned. The *physAddress* specified must be aligned on the architecture dependent page size boundary and must not be mapped to any other memory context.

The MMU attributes specified in *attr* will be used as the default attributes of the shared data region. All client applications will use these by default, and may only change the local access permissions to a subset of these. The application which creates the region will have read and write access in addition to the defaults and will be allowed to set local permissions to any allowed by the architecture.

Basic MMU attribute definitions for shared data regions are provided in the **sdLibCommon.h** header file. These include:

Attribute	Meaning		
SD_ATTR_RW	Read/Write for both Supervisor and User Modes		
SD_ATTR_RO	Read Only for both Supervisor and User Modes		
SD_ATTR_RWX	Read/Write/Execute for both Supervisor and User		
	Modes		
SD_ATTR_RX	Read/Execute for both Supervisor and User Modes		
SD_CACHE_COPYBACK	Copyback cache mode		
SD_CACHE_WRITETHROUGH	Write through cache mode		
SD_CACHE_OFF	Cache Off		

One of each the **SD\_ATTR** and **SD\_CACHE** macros above must be provided. The **SD\_CACHE** macros can not be combined.

The cache attributes of a shared data region can not be changed after creation. All clients of that region will use the value provided at create time, including the owner.

If more specific MMU attributes are required please see **vmLibCommon.h** for a complete list of available MMU attributes.

VxWorks Kernel API Reference, 6.6 sdProtect()

NOTE	The MMU_ATTR mask used internally by the shared data library is the combination of:			
	MMU_ATTR_PROT_MASK			
	MMU_ATTR_VALID_MSK			
	MMU_ATTR_SPL_MSK			
	Care must be taken to provide suitable values for all these attributes.			
	The start address of the shared data region is stored at the location specified by $pVirtAddress$ . This must be a valid address within the context of the calling application. It can not be <b>NULL</b> .			
	The <b>SD_ID</b> returned is private to the calling application. It can be shared between tasks within that application but not with tasks that reside outside that application.			
RETURNS	SD_ID of opened Shared Data region, or NULL on failure.			
ERRNO	Possible errno values set by this routine are:			
	S_sdLib_VIRT_ADDR_PTR_IS_NULL pVirtAddress is NULL			
	S_sdLib_ADDR_NOT_ALIGNED physAddress is not properly aligned			
	S_sdLib_PHYS_ADDR_OUT_OF_RANGE physAddress exceeds physical address space			
	S_sdLib_SIZE_IS_NULL size is NULL			
	S_sdLib_INVALID_OPTIONS <i>options</i> is not a valid combination			
	S_sdLib_VIRT_PAGES_NOT_AVAILABLE not enough virtual space left in system			
	S_sdLib_PHYS_PAGES_NOT_AVAILABLE not enough physical memory left in system			
SEE ALSO	<pre>sdLib, sdCreate( ), sdUnmap( ), sdProtect( ), sdDelete( )</pre>			

# sdProtect()

NAME	sdProtect()	- change the	protection	attributes	of a mapped	d SD
------	-------------	--------------	------------	------------	-------------	------

SYNOPSIS STATUS sdProtect

	( SD_ID so MMU_ATTR at )	AId, /* ID of shared data region */ .tr /* new attributes to set */
DESCRIPTION	This routine allow memory context.	vs the caller to change the protection of a mapped shared data region in its The shared data must be mapped in the context of the calling task.
	These attributes r then the default a	nust be equal to, or a subset of the default attributes of <i>sdId</i> . If 0 was passed attributes of <i>sdId</i> are used.
	The default attrik	putes of <i>sdld</i> may be retrieved by calling the routine <b>sdInfoGet()</b> .
	Basic MMU attril sdLibCommon.h	pute definitions for shared data regions are provided in the header file. These include:
	Attribute	Meaning
	SD_ATTR_RW	Read/Write for both Supervisor and User Modes
	SD_ATTR_RO	Read /Write/Execute for both Supervisor and User Modes
	SD_ATTR_RX	Read/Execute for both Supervisor and User Modes
NOTE	The MMU_ATTR	mask used internally by the shared data library is the combination of:
	MMU_ATTR_PRC	DT_MASK
	MMU_ATTR_VAL	ID_MSK
	MMU_ATTR_CAC	CHE_MSK
	MMU_ATTR_SPL	_MSK
	Care must be tak	en to provide suitable values for all these attributes.
RETURNS	OK, or ERROR or	a failure.
ERRNO	Possible errno va	lues set by this routine are:
	<b>S_sdLib_INVALIE</b> <i>sdId</i> is not va	D_SD_ID Ilid
	S_sdLib_NOT_MA sdId is not m	APPED apped to the current application
SEE ALSO	sdLib.sdCreate(	), sdOpen(), sdMap(), sdUnmap(), sdDelete()

### sdShow()

NAME

sdShow() – display information for shared data regions

DESCRIPTION

This routine displays information for a shared data region. This routine takes two parameters, *sdNameOrId* and *level*. The first parameter can either be an SD ID or an SD name string. The second parameter is the level of detail to display the information for the SDs.

Depending on the level and the SD ID specified, the information displayed differs. If the *level* is 0, then it displays the summary information for either the specified SD or all SDs in the system. If the *level* is 1, then **sdShow()** displays the detailed information, including the client information, for the specified SD or all SDs in the system (if **SD\_ID** is **NULL**). If *level* is 2, **sdShow()** displays the detailed information for all SDs in the system, regardless of the SD ID you specify. Refer to the table for more information.

Level	SD Name or ID	Meaning
0	0	Display summary information for all SDs.
0	SD	Display summary information for specified SD.
1	0	Display detailed information for all SDs.
1	SD	Display detailed information for specified SD.
2	ANY	Display detailed information for all SD.

**sdShow(**) only displays the SD name up to a maximum of 12 characters long. If the name is more than 12 characters, the name will be truncated to 10 characters for displaying purposes. Following the truncated name, a ">" will be display to indicate that the name is more than 12 characters long. To get a display of the full SD name, display the SD with the *level* set to 1.

#### SUMMARY INFORMATION EXAMPLE

The following example shows the summary output for all SDs in the system. If a SD ID (or name) is specified, only the information for that SD will be displayed.

-> sdShow

NAMEIDVIRT ADDRPHYS ADDRSIZECLIENT CNTmySharedDa >0x4c18200xa00000000x017fa0000x10001

value =  $0 = 0 \times 0$ 

The display contains the following fields:

Field	Meaning
NAME	The name of the SD.
ID	The numeric ID associated with the SD in the kernel.
VIRT ADRS	The virtual start address of the SD.
PHYS ADRS	The physical start address of the SD.
SIZE	SD size in bytes.
CLIENT CNT	Number of clients of the SD.

#### DETAILED INFORMATION EXAMPLE

The following example shows the detailed output for a single SD (i.e. the level was specified as 1). If the level is specified as 2, the detailed information is displayed for all SDs in the system and the user is prompted to press **return** or **Q** between each SD.

-> sdShow 0x4c1820, 1

N	IAME	ID	VIRT ADDR	PHYS ADDR	SIZE	CLIENT CNT
mySha	aredDa >	0x4c1820	0xa0000000	0x017fa000	0x1000	1
Full Optic	Name: ons (0x1):	mySl SD_1	naredDataRegion LINGER			
Defau	ilt MMU Ati	tributes	(0x85b):			
	ACCESS		CACHE			
	Sup.: RW-	User: RI	v− default			
Clier	its:					
	NAME	ID	ACCESS		CACHE	
	kernel	0x2!	5a7c0 Sup.: RW	X User: RWX	CB //	

value =  $0 = 0 \times 0$ 

The summary line contains the same fields as explained above. The additional information is explained in the following table:

Field	Meaning
Full Name	The complete name for the SD.
Options	Detailed breakdown of the options word (see
	sdCreate()).
Default MMU Attributes	Default MMU attributes for the SD.
Clients	Complete list of current clients and their access rights.



N/A

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VxWorks Kernel API Reference, 6.6 sdUnmap()

SEE ALSO	sdShow, sdLib, rtpLib, vmBaseLib, the VxWorks programmer guides.
	S_objLib_OBJ_ID_ERROR An incorrect SD ID was provided.
ERRNOS	Possible errnos generated by this function include:

### sdUnmap()

**ERRNO** Possible errno values set by this routine are:

S\_sdLib\_INVALID\_SD\_ID sdId is not valid

**S\_sdLib\_NOT\_MAPPED** *sdId* is not mapped to the current application

SEE ALSO sdLib, sdCreate(), sdOpen(), sdMap(), sdProtect(), sdDelete()

### selNodeAdd()

NAME	<pre>selNodeAdd() – add a wake-up node to a select() wake-up list</pre>
SYNOPSIS	STATUS selNodeAdd

SEL\_WAKEUP\_LIST \*pWakeupList, /\* list of tasks to wake up \*/

 SEL\_WAKEUP\_NODE \*pWakeupNode /\* node to add to list \*/

 DESCRIPTION
 This routine adds a wake-up node to a device's wake-up list. It is typically called from a driver's FIOSELECT function.

 RETURNS
 OK, or ERROR if memory is insufficient.

 ERRNO
 N/A

 SEE ALSO
 selectLib

## selNodeDelete()

NAME	<pre>selNodeDelete() - find and delete a node from a select() wake-up list</pre>
SYNOPSIS	STATUS selNodeDelete ( SEL_WAKEUP_LIST *pWakeupList, /* list of tasks to wake up */ SEL_WAKEUP_NODE *pWakeupNode /* node to delete from list */ )
DESCRIPTION	This routine deletes a specified wake-up node from a specified wake-up list. Typically, it is called by a driver's FIOUNSELECT function.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the node is not found in the wake-up list.
ERRNO	N/A
SEE ALSO	selectLib

## selWakeup()

NAME	<pre>selWakeup() - wake up a task pended in select()</pre>
SYNOPSIS	<pre>void selWakeup   (    SEL_WAKEUP_NODE *pWakeupNode /* node to wake up */   )</pre>

VxWorks Kernel API Reference, 6.6 selWakeupAll()

DESCRIPTION	This routine wakes up a task pended in <b>select()</b> . Once a driver's FIOSELECT function installs a wake-up node in a device's wake-up list (using <b>selNodeAdd()</b> ) and checks to make sure the device is ready, this routine ensures that the <b>select()</b> call does not pend.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	selectLib

# selWakeupAll()

NAME	<pre>selWakeupAll() – wake up all tasks in a select() wake-up list</pre>
SYNOPSIS	<pre>void selWakeupAll   (    SEL_WAKEUP_LIST *pWakeupList, /* list of tasks to wake up */   FAST SELECT_TYPE type /* readers (SELREAD) or writers   (SELWRITE) */   )</pre>
DESCRIPTION	This routine wakes up all tasks pended in <b>select()</b> that are waiting for a device; it is called by a driver when the device becomes ready. The <i>type</i> parameter specifies the task to be awakened, either reader tasks (SELREAD) or writer tasks (SELWRITE).
RETURNS	N/A
ERRNO	N/A
SEE ALSO	selectLib

# selWakeupListInit()

NAME	<pre>selWakeupListInit() - initialize a select() wake-up list</pre>
SYNOPSIS	<pre>void selWakeupListInit   (    SEL_WAKEUP_LIST *pWakeupList /* wake-up list to initialize */ )</pre>

DESCRIPTION	This routine should be called in a device's create routine to initialize the <b>SEL_WAKEUP_LIST</b> structure.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	selectLib

# selWakeupListLen()

NAME	<pre>selWakeupListLen() - get the number of nodes in a select() wake-up list</pre>
SYNOPSIS	<pre>int selWakeupListLen   (    SEL_WAKEUP_LIST *pWakeupList /* list of tasks to wake up */ )</pre>
DESCRIPTION	This routine returns the number of nodes in a specified <b>SEL_WAKEUP_LIST</b> . It can be used by a driver to determine if any tasks are currently pended in <b>select( )</b> on this device, and whether these tasks need to be activated with <b>selWakeupAll( )</b> .
RETURNS	The number of nodes currently in a <b>select( )</b> wake-up list, or <b>ERROR</b> .
ERRNO	N/A
SEE ALSO	selectLib

# selWakeupListTerm()

NAME	<pre>selWakeupListTerm() - terminate a select() wake-up list</pre>
SYNOPSIS	<pre>void selWakeupListTerm   (    SEL_WAKEUP_LIST *pWakeupList /* wake-up list to terminate */ )</pre>
DESCRIPTION	This routine should be called in a device's terminate routine to terminate the <b>SEL_WAKEUP_LIST</b> structure.

VxWorks Kernel API Reference, 6.6 selWakeupType()

RETURNS N/A

ERRNO N/A

SEE ALSO selectLib

# selWakeupType()

NAME	<b>selWakeupType()</b> – get the type of a <b>select()</b> wake-up node
SYNOPSIS	SELECT_TYPE selWakeupType ( SEL_WAKEUP_NODE *pWakeupNode /* node to get type of */ )
DESCRIPTION	This routine returns the type of a specified <b>SEL_WAKEUP_NODE</b> . It is typically used in a device's FIOSELECT function to determine if the device is being selected for read or write operations.
RETURNS	SELREAD (read operation) or SELWRITE (write operation).
ERRNO	N/A
SEE ALSO	selectLib

### select()

NAME	<b>select( )</b> – pend on a set of file descriptors
SYNOPSIS	<pre>int select   (    int width, /* number of bits to examine from 0 */ FAST fd_set *pReadFds, /* read fds */ FAST fd_set *pWriteFds, /* write fds */ fd_set *pExcFds, /* exception fds */ struct timeval *pTimeOut /* max time to wait, NULL = forever */ )</pre>
DESCRIPTION	This routine permits a task to pend until one of a set of file descriptors becomes ready. The

**ESCRIPTION** This routine permits a task to pend until one of a set of file descriptors becomes ready. Three parameters -- *pReadFds*, *pWriteFds*, and *pExceptFds* -- point to file descriptor sets in which each bit corresponds to a particular file descriptor. Bits set in the read file descriptor set

(*pReadFds*) will cause **select()** to pend until data is available on any of the corresponding file descriptors, while bits set in the write file descriptor set (*pWriteFds*) will cause **select()** to pend until any of the corresponding file descriptors become writable.

The following macros are available for setting the appropriate bits in the file descriptor set structure:

```
FD_SET(fd, &fdset)
FD_CLR(fd, &fdset)
FD_ZERO(&fdset)
```

If either *pReadFds*, *pWriteFds*, or *pExceptFds* is **NULL**, they are ignored. The *width* parameter defines how many bits will be examined in the file descriptor sets, and should be set to either the maximum file descriptor value in use plus one, or simply to **FD\_SETSIZE**. When **select()** returns, it zeros out the file descriptor sets, and sets only the bits that correspond to file descriptors that are ready. The **FD\_ISSET** macro may be used to determine which bits are set.

If *pTimeOut* is **NULL**, **select()** will block indefinitely. If *pTimeOut* is not **NULL**, but points to a **timeval** structure with an effective time of zero, the file descriptors in the file descriptor sets will be polled and the results returned immediately. If the effective time value is greater than zero, **select()** will return after the specified time has elapsed, even if none of the file descriptors are ready.

Applications can use **select()** with pipes and serial devices, in addition to sockets. Select now has the capability to support exception reports, but note that most devices do not provide exception notification for select activity. Refer to the manual for each particular driver to learn about its **select()** support, if any.

The value for the maximum number of file descriptors configured in the system (**NUM\_FILES**) should be less than or equal to the value of **FD\_SETSIZE** (2048).

Driver developers should consult the VxWorks programmer guides for details on writing drivers that will use **select()**.

- **RETURNS** The number of file descriptors with activity, 0 if timed out, or **ERROR** if an error occurred when the driver's **select()** routine was invoked via **ioctl()**.
- **ERRNOS** Possible errnos generated by this routine include:
  - S\_selectLib\_NO\_SELECT\_SUPPORT\_IN\_DRIVER A driver associated with one or more fds does not support select().
  - S\_selectLib\_NO\_SELECT\_CONTEXT

The task's select context was not initialized at task creation time.

#### S\_selectLib\_WIDTH\_OUT\_OF\_RANGE The width parameter is greater than the maximum possible fd.

### S\_memLib\_NOT\_ENOUGH\_MEMORY

Heap allocation failure has caused select to fail.

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	EBADF
	An invalid file descriptor was specified in one of the sets, or a valid file descriptor which was specified was closed by another task while the <b>select()</b> call was in progress. (Note, closing a file descriptor in use by another task is NOT recommended.)
SEE ALSO	selectLib, the VxWorks programmer guides.

selectInit( )

NAME	<b>selectInit()</b> – initialize the select facility
SYNOPSIS	<pre>void selectInit   (    int numFiles /* no longer used */ )</pre>
DESCRIPTION	This routine initializes the UNIX BSD 4.3 select facility. It is initialized automatically when the INCLUDE_SELECT component is configured. It installs a task create hook such that a select context is initialized for each task.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	selectLib

## semBCreate()

NAME	<b>semBCreate( )</b> – create and initialize a binary semaphore
SYNOPSIS	<pre>SEM_ID semBCreate   (    int options, /* semaphore options */    SEM_B_STATE initialState /* initial semaphore state */   )</pre>
DESCRIPTION	This routine allocates and initializes a binary semaphore. The semaphore is initialized to the <i>initialState</i> of either <b>SEM_FULL</b> (1) or <b>SEM_EMPTY</b> (0).
	The <i>options</i> parameter specifies the queuing style blocked tasks and response on signals for blocked RTP tasks. Tasks may be queued on a priority basis or a first-in-first-out basis. The

	queuing style options are <b>SEM_Q_PRIORITY</b> (0x1) and <b>SEM_Q_FIFO</b> (0x0), respectively. That parameter also specifies if <b>semGive()</b> should return <b>ERROR</b> when the semaphore fails to send events. This option is turned off by default; it is activated by doing a bitwise-OR of <b>SEM_EVENTSEND_ERR_NOTIFY</b> (0x10) with the queuing style of the semaphore. <b>SEM_INTERRUPTIBLE</b> (0x20) is the option which makes the blocked RTP task on the semaphore ready and return <b>ERROR</b> with errno set to <b>EINTR</b> when a signal is generated to that task. This option has no affect when a kernel task blocks on the same semaphore created with this option. This option is turned off by default.
SMP CONSIDERATIO	INS
	This API is spinlock and intCpuLock restricted. This restriction is not strictly enforced.
RETURNS	The semaphore ID, or NULL if memory cannot be allocated or if error.
ERRNO	S_semLib_INVALID_OPTION Invalid option was specified.
	<b>S_memLib_NOT_ENOUGH_MEMORY</b> Not enough memory available to create the semaphore.
	S_semLib_INVALID_STATE Invalid initial state.
	S_semLib_INVALID_QUEUE_TYPE Invalid type of semaphore queue specified.
	S_spinLockLib_NOT_SPIN_LOCK_CALLABLE This API is spinlock restricted and can not be called taking a spinlock.
SEE ALSO	semBLib, semLib

# semBInitialize()

NAME	<b>semBInitialize( )</b> – initialize a pre-allocated binary semaphore.
SYNOPSIS	SEM_ID semBInitialize (
	char * pSemMem, /* pointer to allocated storage */ int options, /* semaphore options */ SEM_B_STATE initialState /* initial semaphore state */ )
DESCRIPTION	This routine initializes a binary semaphore that has been pre-allocated (i.e. by the <b>VX_BINARY_SEMAPHORE</b> macro). The semaphore is initialized and an ID is returned for further operations on this semaphore.

VxWorks Kernel API Reference, 6.6 semBSmCreate()

The *options* and *initialState* parameters have the same meaning as those for **semBCreate()**. Please see the documentation for **semBCreate()** for more details.

The following example illustrates use of the VX\_BINARY\_SEMAPHORE macro and this function together to instantiate a binary semaphore statically (without using any dynamic memory allocation):

```
#include <vxWorks.h>
#include <semLib.h>
VX_BINARY_SEMAPHORE(mySemB); /* declare the semaphore */
SEM_ID mySemBId; /* semaphore ID for further operations */
STATUS initializeFunction (void)
{
    if ((mySemBId = semBInitialize (mysemB, options, 0)) == NULL)
        return (ERROR); /* initialization failed */
    else
        return (OK);
}
```

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. This restriction is not strictly enforced.

RETURNS	The semaphore ID, or NULL on error.
ERRNO	S_spinLockLib_NOT_SPIN_LOCK_CALLABLE This API is spinlock restricted and can not be called taking a spinlock.

SEE ALSO semBLib

### semBSmCreate()

NAME	semBSmCreate() - create and initialize a shared memory binary semaphore (VxMP
	Option)

SYNOPSIS	SEM_ID semBSmCre	eate			
	( int SEM_B_STATE )	options, initialState	/* /*	<pre>semaphore options */ initial semaphore state *</pre>	/

**DESCRIPTION** This routine allocates and initializes a shared memory binary semaphore. The semaphore is initialized to an *initialState* of either **SEM\_FULL** (available) or **SEM\_EMPTY** (not available). The shared semaphore structure is allocated from the shared semaphore dedicated memory partition.

	The semaphore ID returned by this routine can be used directly by the generic semaphore-handling routines in <b>semLib semGive()</b> , <b>semTake()</b> , and <b>semFlush()</b> and the show routines, such as <b>show()</b> and <b>semShow()</b> .
	The queuing style for blocked tasks is set by <i>options;</i> the only supported queuing style for shared memory semaphores is first-in-first-out, selected by <b>SEM_Q_FIFO</b> .
	Before this routine can be called, the shared memory objects facility must be initialized (see <b>semSmLib</b> ).
	The maximum number of shared memory semaphores (binary plus counting) that can be created is <b>SM_OBJ_MAX_SEM</b> , a configurable parameter.
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory support option, VxMP.
RETURNS	The semaphore ID, or NULL if memory cannot be allocated from the shared semaphore dedicated memory partition.
ERRNO	S_intLib_NOT_ISR_CALLABLE Routine has been called from ISR.
	S_objLib_OBJ_ID_ERROR The shared memory semaphore partition has not been initialized properly.
	S_memLib_NOT_ENOUGH_MEMORY Can't allocate shared memory semaphore object.
	S_semLib_INVALID_QUEUE_TYPE Incorrect semaphore pend queue type specified.
	S_semLib_INVALID_STATE Incorrect initial semaphore state specified.
	S_smObjLib_LOCK_TIMEOUT Can't get the lock on the shared memory semaphore partition in time.
SEE ALSO	semSmLib, semLib, semBLib, smObjLib, semShow, the VxWorks programmer guides.

## semCCreate()

(

NAME	<b>semCCreate()</b> – create and initialize a counting semaphore
SYNOPSIS	SEM_ID semCCreate

```
int options, /* semaphore option modes */
```

VxWorks Kernel API Reference, 6.6 semClnitialize()

```
int initialCount /* initial count */
)
```

**DESCRIPTION** This routine allocates and initializes a counting semaphore. The semaphore is initialized to the specified initial count.

The *options* parameter specifies the queuing style and response on signals for blocked RTP tasks. Tasks may be queued on a priority basis or a first-in-first-out basis. The queuing style options are **SEM\_Q\_PRIORITY** (0x1) and **SEM\_Q\_FIFO** (0x0), respectively. That parameter also specifies if **semGive()** should return **ERROR** when the semaphore fails to send events. This option is turned off by default; it is activated by doing a bitwise-OR of **SEM\_EVENTSEND\_ERR\_NOTIFY** (0x10) with the queuing style of the semaphore. **SEM\_INTERRUPTIBLE**(0x20) is the option which makes the blocked RTP task on the semaphore ready and return **ERROR** with errno set to **EINTR** when a signal is generated to that task. This option has no affect when a kernel task blocks on the same semaphore created with this option. This option is turned off by default.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted.

RETURNS	The semaphore ID, or NULL if memory cannot be allocated or error.
ERRNO	S_semLib_INVALID_INITIAL_COUNT The specified initial count is negative
	S_semLib_INVALID_OPTION Options not applicable to counting semaphores were specified.
	S_memLib_NOT_ENOUGH_MEMORY There is not enough memory to create the semaphore.
SEE ALSO	semCLib, semLib

### semCInitialize()

NAME	<pre>semCInitialize( ) – initialize</pre>	e a pre-allocated counting semaphore.
SYNOPSIS	<pre>SEM_ID semCInitialize   (     char * pSemMem,     int options,     int initialCount</pre>	/* pointer to allocated storage */ /* semaphore options */ /* initial count */

2

**DESCRIPTION** This routine initializes a counting semaphore that has been pre-allocated (i.e. by the **VX\_COUNTING\_SEMAPHORE** macro). The semaphore is initialized and an ID is returned for further operations on this semaphore.

The *options* and *initialCount* parameters have the same meaning as those for **semCCreate()**. Please see the documentation for **semCCreate()** for more details.

The following example illustrates use of the VX\_COUNTING\_SEMAPHORE macro and this function together to instantiate a counting semaphore statically (without using any dynamic memory allocation):

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted.

SEE ALSO	semCLib
ERRNO	N/A
RETURNS	The semaphore ID, or NULL on error.

### semCSmCreate()

NAME	<b>semCSmCreate()</b> – create and initialize a shared memory counting semaphore (VxMP Option)
SYNOPSIS	<pre>SEM_ID semCSmCreate   (     int options, /* semaphore options */     int initialCount /* initial semaphore count */   )</pre>

VxWorks Kernel API Reference, 6.6 semClose()

DESCRIPTION	This routine allocates and initializes a shared memory counting semaphore. The initial count value of the semaphore is specified by <i>initialCount</i> .
	The semaphore ID returned by this routine can be used directly by the generic semaphore-handling routines in <b>semLib semGive()</b> , <b>semTake()</b> and <b>semFlush()</b> and the show routines, such as <b>show()</b> and <b>semShow()</b> .
	The queuing style for blocked tasks is set by <i>options;</i> the only supported queuing style for shared memory semaphores is first-in-first-out, selected by <b>SEM_Q_FIFO</b> .
	Before this routine can be called, the shared memory objects facility must be initialized (see <b>semSmLib</b> ).
	The maximum number of shared memory semaphores (binary plus counting) that can be created is <b>SM_OBJ_MAX_SEM</b> , a configurable paramter.
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory support option, VxMP.
RETURNS	The semaphore ID, or <b>NULL</b> if memory cannot be allocated from the shared semaphore dedicated memory partition.
ERRNO	S_intLib_NOT_ISR_CALLABLE Routine has been called from ISR.
	S_objLib_OBJ_ID_ERROR The shared memory semaphore partition has not been initialized properly.
	S_memLib_NOT_ENOUGH_MEMORY Can't allocate shared memory semaphore object.
	S_semLib_INVALID_QUEUE_TYPE Incorrect semaphore pend queue type specified.
	S_semLib_INVALID_COUNT Incorrect initial count (negative) specified.
	S_smObjLib_LOCK_TIMEOUT Can't get the lock on the shared memory semaphore partition in time.
SEE ALSO	semSmLib, semLib, semCLib, smObjLib, semShow, the VxWorks programmer guides.

# semClose()

NAME semClose() – close a named semaphore

SYNOPSIS STATUS semClose

```
(
SEM_ID semId /* semaphore ID to close */
)
```

#### **DESCRIPTION** This routine closes a named semaphore. It decrements the semaphore's reference counter. In case it becomes zero, the semaphore is deleted if:

- It has been already removed from the name space by a call to **semUnlink()**.
- It was created with the OM\_DESTROY\_ON\_LAST\_CLOSE option.

This routine is not ISR callable.

#### SMP CONSIDERATIONS

ERRNO

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

- **RETURNS** OK, or ERROR if unsuccessful.
  - S\_objLib\_OBJ\_ID\_ERROR Semaphore ID is invalid.
    - S\_objLib\_OBJ\_INVALID\_ARGUMENT Semaphore ID is NULL.
    - **S\_objLib\_OBJ\_OPERATION\_UNSUPPORTED** Semaphore is not named.
    - **S\_objLib\_OBJ\_DESTROY\_ERROR** Error while deleting the semaphore.
    - S\_intLib\_NOT\_ISR\_CALLABLE This routine must not be called from an ISR.

SEE ALSO semOpen, semOpen, semUnlink

### semDelete()

NAME	<b>semDelete( )</b> – delete a semaphore
SYNOPSIS	STATUS semDelete ( SEM_ID semId /* semaphore ID to delete */ )

VxWorks Kernel API Reference, 6.6 semEvStart()

DESCRIPTION	This routine terminates and deallocates any memory associated with a specified semaphore. All tasks pending on the semaphore or pending for the reception of events meant to be sent from the semaphore will unblock and return <b>ERROR</b> .
SMP CONSIDERATIO	NS
	This API is spinlock and intCpuLock restricted.
WARNING	Take care when deleting semaphores, particularly those used for mutual exclusion, to avoid deleting a semaphore out from under a task that already has taken (owns) that semaphore. Applications should adopt the protocol of only deleting semaphores that the deleting task has successfully taken.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the semaphore ID is invalid.
ERRNOS	S_intLib_NOT_ISR_CALLABLE Routine cannot be called from ISR.
	S_objLib_OBJ_ID_ERROR Semaphore ID is invalid.
	S_smObjLib_NO_OBJECT_DESTROY Deleting a shared semaphore is not permitted
	S_objLib_OBJ_OPERATION_UNSUPPORTED Deleting a named semaphore is not permitted.
SEE ALSO	semLib, semBLib, semCLib, semMLib, semSmLib

### semEvStart()

 NAME
 semEvStart() - start the event notification process for a semaphore

 SYNOPSIS
 STATUS semEvStart

 (
 SEM\_ID semId, /\* semaphore on which to register events \*/
 UINT32 events, /\* 32 possible events to register \*/
 UINT8 options /\* event-related semaphore options \*/
 This routine turns on the event notification process for a given semaphore, registering the calling task on that semaphore. When the semaphore becomes available but no task is pending on it, the events specified will be sent to the registered task. A task can always overwrite its own registration.
 The events are user-defined. For more information, see the reference entry for eventLib.

The *option* parameter is used for 3 user options:

- Specify if the events are to be sent only once or every time the semaphore becomes free until **semEvStop()** is called.
- Specify if another task can subsequently register itself while the calling task is still registered. If so specified, the existing task registration will be overwritten without any warning.
- Specify if events are to be sent at the time of the registration in the case the semaphore is free.

Here are the respective values to be used to form the options field:

#### EVENTS\_SEND\_ONCE (0x1)

The semaphore will send the events only once.

#### EVENTS\_ALLOW\_OVERWRITE (0x2)

Subsequent registrations from other tasks may overwrite the current one.

#### EVENTS\_SEND\_IF\_FREE (0x4)

The registration process will send events if the semaphore is free at the time **semEvStart()** is called.

#### EVENTS\_OPTIONS\_NONE (0x0)

Must be passed to the options parameter if none of the other three options are used.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted.

- **WARNING** This routine cannot be called from interrupt level.
- WARNING Task preemption can allow a semDelete() to be performed between the calls to semEvStart() and eventReceive(). This prevents the task from ever receiving the events wanted from the semaphore.
- **RETURNS** OK on success, or ERROR.
- ERRNO S\_objLib\_OBJ\_ID\_ERROR The semaphore ID is invalid.
  - S\_eventLib\_ALREADY\_REGISTERED

A task is already registered on the semaphore.

### S\_intLib\_NOT\_ISR\_CALLABLE

This routine cannot be called from interrupt level.

#### S\_eventLib\_EVENTSEND\_FAILED

The user chose to send events immediately and that operation failed.

#### S\_eventLib\_ZERO\_EVENTS

The user passed in a value of zero to the *events* parameter.

VxWorks Kernel API Reference, 6.6 semEvStop()

SEE ALSO semEvLib, eventLib, semLib, semEvStop()

# semEvStop()

NAME	<b>semEvStop( )</b> – stop the event notification process for a semaphore
SYNOPSIS	STATUS semEvStop ( SEM_ID semId )
DESCRIPTION	This routine turns off the event notification process for a given semaphore. It thus allows another task to register itself for event notification on that particular semaphore. It must be called from the task that is already registered on that particular semaphore.
SMP CONSIDERATIO	INS
	This API is spinlock and intCpuLock restricted.
RETURNS	OK on success, or ERROR.
ERRNO	S_objLib_OBJ_ID_ERROR The semaphore ID is invalid.
	S_intLib_NOT_ISR_CALLABLE The routine cannot be called from interrupt level.
	S_eventLib_TASK_NOT_REGISTERED The routine was not called by the registered task.
SEE ALSO	semEvLib, eventLib, semLib, semEvStart()

# semExchange()

NAME	<b>semExchange( )</b> – atomically give and take a pair of semaphores
SYNOPSIS	<pre>STATUS semExchange (    SEM_ID giveSemId, /* semaphore ID to give */    SEM_ID takeSemId, /* semaphore ID to take */    int timeout /* timeout in ticks */ )</pre>
**DESCRIPTION** This routine atomically performs a give operation on a sempahore and a take operation on another semaphore. The semaphore specified to be given will be released when the caller acquires or pends attempting to acquire the semaphore specified to be taken.

This routine performs the give operation on a semaphore specified by the *giveSemId* argument. Depending on the type of this semaphore, the state of the semaphore and of the pending tasks may be affected. If no tasks are pending on the semaphore and a task has previously registered to receive events from the semaphore, these events are sent in the context of this call. This may result in the unpending of the task waiting for the events. If the semaphore fails to send events and if it was created using the

SEM\_EVENTSEND\_ERR\_NOTIFY option, ERROR is returned even though the give operation was successful. The behavior of **semGive()** is discussed fully in the library description of the specific semaphore type being used.

If the give operation returns **ERROR** for any reason the subsequent take operation will not be performed.

This routine performs the take operation on a semaphore specified by the *takeSemId* argument. Depending on the type of this semaphore, the state of the semaphore and the calling task may be affected. The behavior of **semTake()** is discussed fully in the library description of the specific semaphore type being used.

A timeout in ticks may be specified for the **semTake()** portion of the **semExchange()** operation. If a task times out, **semExchange()** will return **ERROR**. Timeouts of **WAIT\_FOREVER** (-1) and **NO\_WAIT** (0) indicate to wait indefinitely or not to wait at all.

When **semExchange()** returns due to timeout, it sets the errno to **S\_objLib\_OBJ\_TIMEOUT** (defined in **objLib.h**).

Because it completes when the caller pends during the **semTake()** operation the **semGive()** operation will occur regardless of timeout. It is possible for the caller to release the specified give semaphore and not acquire the semaphore specified to be taken.

The **semExchange()** routine is not callable from interrupt service routines.

Currently only binary and mutex semaphore types are supported by semExchange().

An attempt to specify a semaphore of another type for either the give or take operation of **semExchange()** will result in a return value of **ERROR**. Neither the give or take operation will be performed.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted.

- **RETURNS** OK, or ERROR if the semaphore ID is invalid or the task timed out.
- ERRNOS S\_intLib\_NOT\_ISR\_CALLABLE Routine was called from an ISR.
  - S\_objLib\_OBJ\_ID\_ERROR Semaphore ID is invalid.

VxWorks Kernel API Reference, 6.6 semFlush()

S\_objLib\_OBJ\_TIMEOUT Timeout occured while pending on sempahore.

S\_objLib\_OBJ\_UNAVAILABLE Would have blocked but NO\_WAIT was specified.

#### S\_semLib\_INVALID\_OPERATION

Current task not owner of semaphore.

#### S\_eventLib\_EVENTSEND\_FAILED

Semaphore failed to send events to the registered task. This errno value can only exist if the semaphore was created with the SEM\_EVENTSEND\_ERR\_NOTIFY option.

SEE ALSO semExchange, semLib, semBLib, semMLib

### semFlush( )

NAME	<b>semFlush( )</b> – unblock every task pended on a semaphore	
SYNOPSIS	STATUS semFlush ( SEM_ID semId /* semaphore ID to unblock everyone for */ )	
DESCRIPTION	This routine atomically unblocks all tasks pended on a specified semaphore, i.e., all tasks will be unblocked before any is allowed to run. The state of the underlying semaphore is unchanged. All pended tasks will enter the ready queue before having a chance to execute.	
	The flush operation is useful as a means of broadcast in synchronization applications. Its use is illegal for mutual-exclusion semaphores created with <b>semMCreate()</b> or with reader/writer semaphores created with <b>semRWCreate()</b> .	
SMP CONSIDERATIONS		
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.	
RETURNS	OK, or ERROR if the semaphore ID is invalid or the operation is not supported.	
ERRNO	S_objLib_OBJ_ID_ERROR	
SEE ALSO	semLib, semBLib, semCLib, semMLib, semRWLib, semSmLib	

## semGive()

NAME	<b>semGive( )</b> – give a semaphore
SYNOPSIS	STATUS semGive ( SEM_ID semId /* semaphore ID to give */ )
DESCRIPTION	This routine performs the give operation on a specified semaphore. Depending on the type of semaphore, the state of the semaphore and of the pending tasks may be affected. If no tasks are pending on the semaphore and a task has previously registered to receive events from the semaphore, these events are sent in the context of this call. This may result in the unpending of the task waiting for the events. If the semaphore fails to send events and if it was created using the <b>SEM_EVENTSEND_ERR_NOTIFY</b> option, <b>ERROR</b> is returned even though the give operation was successful. The behavior of <b>semGive()</b> is discussed fully in the library description of the specific semaphore type being used.
SMP CONSIDERATIO	NS
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	OK on success or ERROR otherwise
ERRNO	S_intLib_NOT_ISR_CALLABLE Routine was called from an ISR for a semaphore.
	S_objLib_OBJ_ID_ERROR Semaphore ID is invalid.
	S_semLib_INVALID_OPERATION Current task not owner of semaphore.
	S_semLib_COUNT_OVERFLOW Counting semaphore was given when count was already at maximum.
	<b>S_eventLib_EVENTSEND_FAILED</b> Semaphore failed to send events to the registered task. This errno value can only exist if the semaphore was created with the <b>SEM_EVENTSEND_ERR_NOTIFY</b> option.
SEE ALSO	semLib, semBLib, semCLib, semMLib, semRWLib, semSmLib, semEvStart()

### semInfo()

```
semInfo() - get information about tasks blocked on a semaphore
NAME
SYNOPSIS
                int semInfo
                     (
                    int idList[], /* array of task IDs to be filled in */
                    int maxTasks /* max tasks idList can accommodate */
                    )
DESCRIPTION
                This routine returns the number of tasks that are blocked on the specified semaphore, semId.
                If a non-NULL array is passed in idList, then up to maxTasks task IDs are copied into the
                array. In this case, this routine returns the number of task IDs that could be copied into the
                array. The array is unordered.
WARNING
                There is no guarantee that all listed tasks are still valid or that new tasks have not been
                blocked by the time semInfo() returns.
                The actual number of blocked tasks if idList is NULL
RETURNS
                     or the number of task IDs placed in idList.
ERRNO
                S_objLib_OBJ_ID_ERROR
                    Invalid semaphore ID.
                S_intLib_NOT_ISR_CALLABLE
                    This routine is not callable from an ISR.
SEE ALSO
                semInfo, semInfoGet()
```

### semInfoGet()

NAME	<b>semInfoGet( )</b> – get information about a semaphore
SYNOPSIS	STATUS semInfoGet ( SEM_ID semId, /* semaphore to query */ SEM_INFO * pInfo /* where to return semaphore info */ )
DESCRIPTION	This routine gets information about the state of a semaphore. The parameter <i>pInfo</i> pointer to a structure of type <b>SEM_INFO</b> defined in <b>semLibCommon.h</b> as follows:

*pInfo* is a

```
typedef struct
                             /* SEM INFO */
   {
   UINT numTasks;
                             /* OUT: number of blocked tasks */
   SEM_TYPE semType;
                             /* OUT: semaphore type */
   int options;
                             /* OUT: options with which sem was created */
   union
       {
                         /* OUT: semaphore count (counting sems) */
       UINT count;
                            /* OUT: binary semaphore FULL? */
       BOOL full;
       int owner;
                            /* OUT: task ID of mutex semaphore owner */
       } state;
   int taskIdListMax; /* IN: max tasks to fill in taskIdLi
int * taskIdList; /* PTR: array of pending task IDs */
           taskIdListMax; /* IN: max tasks to fill in taskIdList */
    SEM_INFO;
```

The semaphore type is determined by examining *semType*. Based on this information the appropriate field in the *state* union can be examined to determine a) the current count of a counting semaphore *state.count*, b) whether a binary semaphore is full *state.full*, or c) the owner of a mutex semaphore *state.owner*.

If a binary semaphore is not full *state.full* = **FALSE**, or if a counting semaphore's count is 0 *state.count* = 0, or a mutex semaphore is already owned *state.owner* != **NULL**, then there may be tasks blocked on **semTake()**. The *numTasks* field indicates the number of blocked tasks.

A list of the task IDs of tasks blocked on the semaphore can be obtained by setting *taskIdList* to the address of an array to receive the list, and setting *taskIdListMax* to the maximum number of elements in that array. If *taskIdList* is **NULL**, then no task IDs are returned. No more than *taskIdListMax* task IDs are returned, although *numTasks* will always be returned with the actual number of tasks blocked.

For example, if the caller supplies a *taskIdList* with room for 10 task IDs and sets *taskIdListMax* to 10, but there are 20 tasks blocked on the semaphore, then the IDs of the first 10 tasks blocked on the semaphore will be returned in *taskIdList*, but *numTasks* will be returned with the value 20.

The *options* field is the parameter with which the semaphore was created.

**WARNING** The information returned by this routine is not static and may be obsolete by the time it is examined. In particular, the list of task IDs may no longer be valid. However, the information is obtained atomically, thus it will be an accurate snapshot of the state of the semaphore at the time of the call. This information is generally used for debugging purposes only.

If *taskIdList* is non-NULL, i.e. the caller is requesting the list of pended tasks, the execution time of **semInfoGet()** may be non-deterministic.

This routine cannot be used to extract information on shared semaphores.

RETURNS OK or ERROR

ERRNO S\_objLib\_OBJ\_ID\_ERROR Invalid semaphore ID. 2

#### S\_semLib\_INVALID\_OPERATION

Specified semaphore is a shared semaphore, or the semaphore is an unknown type.

SEE ALSO semInfo, semInfo()

### semMCreate()

NAME semMCreate() - create and initialize a mutual-exclusion semaphore SYNOPSIS SEM\_ID semMCreate ( int options /\* mutex semaphore options \*/ This routine allocates and initializes a mutual-exclusion semaphore. The semaphore state DESCRIPTION is initialized to full. Semaphore options include the following: **SEM\_Q\_PRIORITY** (0x1)Queue pended tasks on the basis of their priority.  $SEM_Q_FIFO$  (0x0) Queue pended tasks on a first-in-first-out basis. **SEM\_DELETE\_SAFE** (0x4)Protect a task that owns the semaphore from unexpected deletion. This option enables an implicit taskSafe() for each semTake(), and an implicit taskUnsafe() for each semGive(). **SEM\_INVERSION\_SAFE** (0x8) Protect the system from priority inversion. With this option, the task owning the semaphore will execute at the highest priority of the tasks pended on the semaphore, if it is higher than its current priority. This option must be accompanied by the SEM\_Q\_PRIORITY queuing mode. SEM\_EVENTSEND\_ERR\_NOTIFY (0x10) When the semaphore is given, if a task is registered for events and the actual sending of events fails, a value of ERROR is returned and the errno is set accordingly. This option is off by default. **SEM\_INTERRUPTIBLE** (0x20) Signal sent to an RTP task blocked on a semaphore created with this option, would make the task ready and return with ERROR and errno set to EINTR. This option has no affect for a kernel task blocked on the same semaphore created with this option. This option is off by default.

SMP CONSIDERATIONS	
	This API is spinlock and intCpuLock restricted. This restriction is not strictly enforced.
RETURNS	The semaphore ID, or NULL if the semaphore cannot be created.
ERRNO	S semLib INVALID OPTION
	Invalid option was passed to semMCreate.
	S_memLib_NOT_ENOUGH_MEMORY
	Not enough memory available to create the semaphore.
SEE ALSO	semMLib, semLib, semBLib, taskSafe( ), taskUnsafe( )

## semMGiveForce()

NAME	<b>semMGiveForce( )</b> – give a mutual-exclusion semaphore without restrictions	
SYNOPSIS	STATUS semMGiveForce ( FAST SEM_ID semId /* semaphore ID to give */ )	
DESCRIPTION	This routine gives a mutual-exclusion semaphore, regardless of semaphore ownership. It is intended as a debugging aid only.	
	The routine is particularly useful when a task dies while holding some mutual-exclusion semaphore, because the semaphore can be resurrected. The routine will give the semaphore to the next task in the pend queue or make the semaphore full if no tasks are pending. In effect, execution will continue as if the task owning the semaphore had actually given the semaphore.	
SMP CONSIDERATIONS		
	This API is spinlock and intCpuLock restricted. This restriction is not strictly enforced.	
CAVEATS	This routine should be used only as a debugging aid, when the condition of the semaphore is known.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if the semaphore ID is invalid.	
ERRNO	N/A	
SEE ALSO	semMLib, semGive()	

### semMInitialize()

**NAME semMInitialize()** – initialize a pre-allocated mutex semaphore.

**DESCRIPTION** This routine initializes a mutual exclusion semaphore that has been pre- allocated (i.e. by the VX\_MUTEX\_SEMAPHORE macro). The semaphore is initialized and an ID is returned for further operations on this semaphore.

The *options* parameter has the same meaning as that for **semMCreate()**. Please see the documentation for **semBCreate()** for more details.

The following example illustrates use of the **VX\_MUTEX\_SEMAPHORE** macro and this function together to instantiate a mutex semaphore statically (without using any dynamic memory allocation):

```
#include <vxWorks.h>
#include <semLib.h>
VX_MUTEX_SEMAPHORE(mySemM); /* declare the semaphore */
SEM_ID mySemMId; /* semaphore ID for further operations */
STATUS initializeFunction (void)
{
    if ((mySemMId = semMInitialize (mySemM, options, 0)) == NULL)
        return (ERROR); /* initialization failed */
    else
        return (OK);
}
```

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. This restriction is not strictly enforced.

RETURNS	The semaphore ID, or NULL on error.
ERRNO	S_spinLockLib_NOT_SPIN_LOCK_CALLABLE This API is spinlock restricted and can not be called after taking a spinlock.

```
SEE ALSO semMLib
```

### semOpen()

NAME

semOpen() – open a named semaphore

SYNOPSIS SEM\_ID semOpen

**DESCRIPTION** This routine either opens an existing semaphore or creates a new semaphore if the appropriate flags in the *mode* parameter are set. A semaphore with the name specified by the *name* parameter is searched for, and if found the **SEM\_ID** of the semaphore is returned. A new semaphore may only be created if the search of existing semaphores fails (ie. the name must be unique).

There are two name spaces in which **semOpen()** can perform a search in, the "private to the application" name space and the "public" name space. Which is selected depends on the first character in the *name* parameter. When this character is a forward slash /, the "public" name space is used, otherwise the the "private to the application" name space is used.

Semaphores created by this routine can not be deleted with **semDelete()**. Instead, a **semClose()** must be issued for every **semOpen()**. Then the semaphore is deleted when it is removed from the name space by a call to **semUnlink()**. Alternatively, the semaphore can be previously removed from the name space, and deleted during the last **semClose()**.

The parameters to the semOpen function are as follows:

name

A mandatory text string which represents the name by which the semaphore is known by. NULL or empty strings can not be used.

#### type

When creating a semaphore, it specifies which type of semaphore is to be created. The valid types are:

SEM_TYPE_BINARY	create a binary semaphore
SEM_TYPE_MUTEX	create a mutual exclusion semaphore
SEM_TYPE_COUNTING	create a counting semaphore

initState

When a binary or counting semaphore is created, the initial state of the semaphore is set according to the value of *initState*. For binary semaphores the value of *initState* must be either **SEM\_FULL** or **SEM\_EMPTY**. For counting semaphores the semaphore count is set to the value of *initState*.

2

#### options

Semaphore creation options as decribed in semLib.

mode

The mode parameter consists of the access rights (which are currently ignored) and the opening flags which are bitwise-OR'd together. The flags available are:

#### OM\_CREATE

Create a new semaphore if a matching semaphore name is not found.

#### OM\_EXCL

When set jointly with the **OM\_CREATE** flag, creates a new semaphore immediately without trying to open an existing semaphore. The call fails if the semaphore's name causes a name clash. This flag has no effect if the **OM\_CREATE** flag is not specified.

#### OM\_DELETE\_ON\_LAST\_CLOSE

Only used when a semaphore is created. If set, the semaphore will be deleted during the last **semClose()** call, independently on whether **semUnlink()** was previously called or not.

context

Context value assigned to the created semaphore. This value is not actually used by VxWorks. Instead, the context value can be used by OS extensions to implement object permissions, for example.

Unlike private objects, a public semaphore is not automatically reclaimed when an application terminates. Note that nevertheless, a **semClose()** is issued on every application's outstanding **semOpen()**. Therefore, a public semaphore can effectively be deleted, if during this process it is closed for the last time, and it is already unlinked or it was created with the OM\_DELETE\_ON\_LAST\_CLOSE flag.

This routine is not ISR callable.

name was not found.

#### SMP CONSIDERATIONS

 

 This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

 RETURNS
 The SEM\_ID of the opened semaphore, or NULL if unsuccessful.

 ERRNO
 S\_objLib\_OBJ\_INVALID\_ARGUMENT An invalid option was specified in the *mode* argument or *name* is invalid.

 S\_semLib\_INVALID\_INITIAL\_COUNT The specified initial count for counting semaphore is negative
 S\_objLib\_OBJ\_NOT\_FOUND The OM\_CREATE flag was not set in the *mode* argument and a semaphore matching

	S_objLib_OBJ_NAME_CLASH The OM_CREATE and OM_EXCL flags were set and a name clash was detected when creating the semaphore.
	S_intLib_NOT_ISR_CALLABLE This routine must not be called from an ISR.
SEE ALSO	semOpen, semUnlink(), semClose()

### semOpenInit()

NAME	<b>semOpenInit()</b> – initialize the semaphore open facility
SYNOPSIS	void semOpenInit (void)
DESCRIPTION	This routine links the semaphore creation routine with open facility into the VxWorks system. It is called automatically when the semaphore facility is configured into VxWorks by either defining INCLUDE_OBJ_OPEN and INCLUDE_SEM_BINARY in config.h or selecting INCLUDE_OBJ_OPEN and INCLUDE_SEM_BINARY in the project facility.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	semOpen

### semPxLibInit( )

**NAME semPxLibInit()** – initialize POSIX semaphore support

SYNOPSIS STATUS semPxLibInit (void)

**DESCRIPTION** This routine must be called before using POSIX semaphores. If POSIX semaphores are included, this routine will be called during system initialization.

**RETURNS** OK, or ERROR if there is an error installing the semaphore library.

ERRNO None

SEE ALSO semPxLib

## semPxShow()

NAME	<b>semPxShow()</b> – display semaphore internals
SYNOPSIS	<pre>STATUS semPxShow   (    sem_t * semDesc,    int level   )</pre>
DESCRIPTION	This routine displays POSIX semaphore information. Currently, only a <i>level</i> of 0 is supported. This function prints the semaphore name, how many times sem_open has been called, and the value of the semaphore. If the semaphore value is greater than zero, than the number of available semaphores is printed. If the semaphore value is equal to 0, then the number of blocked tasks are also printed.
RETURNS	<b>OK</b> or <b>ERROR</b> if the descriptor is invalid.
ERRNO	N/A
SEE ALSO	semPxShow

## semPxShowInit()

NAME	<pre>semPxShowInit() - initialize the POSIX semaphore show facility</pre>
SYNOPSIS	STATUS semPxShowInit (void)
DESCRIPTION	This routine links the POSIX semaphore show routine into the VxWorks system. It is called automatically when the this show facility is configured into VxWorks using the INCLUDE_POSIX_SEM_SHOW component.
RETURNS	ОК.
ERRNO	N/A
SEE ALSO	semPxShow

## semRTake()

NAME	<b>semRTake( )</b> – take a semaphore as a reader
SYNOPSIS	STATUS semRTake ( SEM_ID semId, /* semaphore ID to take */ int timeout /* timeout in ticks */ )
DESCRIPTION	Takes the semaphore. If the semaphore is held by another task in "write" mode (or another task has attempted to take the semaphore in "write" mode and pended) the task will become pended until the semaphore becomes available. If the semaphore is already available or held by other tasks in "read" mode (with no tasks pended in "write" mode) the caller will gain ownership.
	After a successful call to this routine the caller is granted concurrent access along with those tasks that have also taken the semaphore in this mode. Mutual exclusion is maintained between these tasks and tasks that have taken the semaphore in "write" mode.
	This routine may be called recursively. However, it should not be called by a task that holds the semaphore in "write" mode. Calling <b>semRTake()</b> in such circumstances will result in a return value of <b>ERROR</b> .
	If deletion safe option is enabled, an implicit <b>taskSafe()</b> operation will occur.
	If priority inversion safe option is enabled, and the calling task blocks, and the priority of the calling task is greater than the semaphore owner, the owner will inherit the caller's priority.
SMP CONSIDERATIO	NS
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
WARNING	This routine must not be used from interrupt level.
RETURNS	OK, or ERROR if the semaphore ID is invalid or the task timed out
ERRNO	S_intLib_NOT_ISR_CALLABLE Routine was called from an ISR.
	S_objLib_OBJ_ID_ERROR Semaphore ID is invalid.
	S_objLib_OBJ_TIMEOUT Timeout occured while pending on sempahore.

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S\_objLib\_OBJ\_UNAVAILABLE Would have blocked but NO\_WAIT was specified.

S\_semLib\_INVALID\_OPERATION Task already holds the semaphore as a writer.

SEE ALSO semRWLib

### semRWCreate( )

semRWCreate() - create and initialize a reader/writer semaphore NAME SYNOPSIS SEM\_ID semRWCreate ( int options, /\* reader/writer semaphore options \*/ int maxReaders /\* maximum concurrent readers \*/ DESCRIPTION This routine allocates and initializes a reader/writer semaphore. Semaphore options include the following: **SEM\_Q\_PRIORITY** (0x1) Queue pended tasks on the basis of their priority.  $SEM_Q_FIFO$  (0x0) Queue pended tasks on a first-in-first-out basis. **SEM\_DELETE\_SAFE** (0x4) Protect a task that owns the semaphore from unexpected deletion. This option enables an implicit taskSafe() for each semTake(), and an implicit taskUnsafe() for each semGive(). SEM\_INVERSION\_SAFE (0x8) Protect the system from priority inversion. With this option, the task or tasks owning the semaphore will execute at the highest priority of the tasks pended on the semaphore, if it is higher than its current priority. This option must be accompanied by the **SEM\_Q\_PRIORITY** queuing mode. The maxReaders argument specifies the maximum number of tasks that may concurrently hold a read/write semaphore in **read** mode. It is an error to specify a value of **0** for maxReaders. If the value of maxReaders exceeds the system maximum value (specified in the

system specific maximum will be used instead of maxReaders.

component configuration option SEM\_RW\_MAX\_CONCURRENT\_READERS) then that

SMP CONSIDERATIO	This API is spinlock and intCpuLock restricted.
RETURNS	The semaphore ID, or <b>NULL</b> if the semaphore cannot be created.
ERRNO	<b>S_semLib_INVALID_OPTION</b> Invalid option was passed to semRWCreate or <i>maxReaders</i> is 0.
	S_memLib_NOT_ENOUGH_MEMORY Not enough memory available to create the semaphore.
SEE ALSO	<pre>semLib, semRWLib, semMLib, semBLib, taskSafe( ), taskUnsafe( )</pre>

## semRWGiveForce()

NAME	<b>semRWGiveForce()</b> – give a reader/writer semaphore without restrictions
SYNOPSIS	STATUS semRWGiveForce ( FAST SEM_ID semId /* semaphore ID to give */ )
DESCRIPTION	This routine gives a reader/writer semaphore, regardless of semaphore ownership. It is intended as a debugging aid only.
	The routine is particularly useful when a task dies while holding some reader/writer semaphore, because the semaphore can be resurrected. The routine will give the semaphore to the next task in the pend queue or make the semaphore full if no tasks are pending. In effect, execution will continue as if the task owning the semaphore had actually given the semaphore.
SMP CONSIDERATIO	NS
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
CAVEATS	This routine should be used only as a debugging aid, when the condition of the semaphore is known.
RETURNS	OK, or ERROR if the semaphore ID is invalid
ERRNO	S_intLib_NOT_ISR_CALLABLE This routine is not callable from an ISR.

VxWorks Kernel API Reference, 6.6 semRWInitialize()

SEE ALSO semRWLib, semGive()

#### semRWInitialize()

```
semRWInitialize() – initialize a pre-allocated read/write semaphore.
NAME
SYNOPSIS
                 SEM ID semRWInitialize
                      (
                      char * pSemMem, /* pointer to allocated storage */
int options, /* RW semaphore options */
int maxReaders /* maximum concurrent readers */
                      )
DESCRIPTION
                 This routine initializes a reader/writer semaphore that has been pre- allocated (i.e. by the
                  VX_READ_WRITE_SEMAPHORE macro). The semaphore is initialized and an ID is returned
                  for further operations on this semaphore.
                 The options parameter has the same meaning as that for semRWCreate(). Please see the
                 documentation for semRWCreate() for more details.
                  The maxReaders parameter specifies the maximum concurrent readers for the semaphores.
                 If this value exceeds that of the system defined maximum, specified in
                 SEM_RW_MAX_CONCURRENT_READERS, then that system specified value will be used
                 instead of maxReaders. It is worth noting that memory allocated in this case will still be that
                  of a semaphore created with maxReaders number of maximum readers. It is an error to
                  specify 0 as the value of maxReaders.
                 The following example illustrates use of the VX_READ_WRITE_SEMAPHORE macro and
                  this function together to instantiate a read/write semaphore statically (without using any
                  dynamic memory allocation):
                       #include <vxWorks.h>
                       #include <semLib.h>
                       #define NUM READERS 0x20
                       /* declare the semaphore */
                       VX READ WRITE SEMAPHORE (mySemRW, NUM READERS);
                       SEM ID mySemRWId;
                                                          /* semaphore ID for further operations */
                       STATUS initializeFunction (void)
                            {
                            if ((mySemRWId =
                                          semRWInitialize (mySemRW, options, NUM READERS)) ==
                 NULL)
                                  return (ERROR); /* initialization failed */
                            else
```

return (OK);

RETURNS	The semaphore ID, or NULL on error
ERRNO	S_semLib_INVALID_OPTION Invalid options were provided.
SEE ALSO	semRWLib

}

### semShow()

semShow() - show information about a semaphore NAME SYNOPSIS STATUS semShow ( SEM\_ID semId, /\* semaphore to display \*/ int level /\* 0 = summary, 1 = details \*/ ) This routine displays the state and optionally the pended tasks of a semaphore. DESCRIPTION A summary of the state of the semaphore is displayed as follows: Semaphore Id : 0x585f2 Semaphore Type : BINARY Task Queuing : PRIORITY Pended Tasks : 1 State State : EMPTY {Count if COUNTING, Owner if MUTEX} Options : 0x1 SEM\_Q\_PRIORITY VxWorks Events \_\_\_\_\_ Registered Task : 0x594f0 (t1) Event(s) to Send : 0x1 Options : 0x7 EVENTS\_SEND\_ONCE EVENTS\_ALLOW\_OVERWRITE EVENTS\_SEND\_IF\_FREE

If *level* is 1, then more detailed information will be displayed. If tasks are blocked on the queue, they are displayed in the order in which they will unblock, as follows:

Pended Tasks				
NAME	TID	PRI	DELAY	
tExcTask	3fd678	0	21	
tLogTask	3f8ac0	0	611	

VxWorks Kernel API Reference, 6.6 semTake()

RETURNS	OK or ERROR.
ERRNO	S_smObjLib_NOT_INITIALIZED The shared memory object library is not initialized.
SEE ALSO	<b>semShow, windsh</b> , <i>VxWorks Programmer's Guide</i> , <i>VxWorks Command-Line Tools User's Guide</i> .

### semTake()

NAME semTake() – take a semaphore SYNOPSIS STATUS semTake ( SEM\_ID semId, /\* semaphore ID to take \*/ int timeout /\* timeout in ticks \*/ ) DESCRIPTION This routine performs the take operation on a specified semaphore. Depending on the type of semaphore, the state of the semaphore and the calling task may be affected. The behavior of **semTake()** is discussed fully in the library description of the specific semaphore type being used. A timeout in ticks may be specified. If a task times out, **semTake()** will return **ERROR**. Timeouts of WAIT\_FOREVER (-1) and NO\_WAIT (0) indicate to wait indefinitely or not to wait at all. When **semTake()** returns due to timeout, it sets the errno to **S\_objLib\_OBJ\_TIMEOUT** (defined in objLib.h). The semTake() routine is not callable from interrupt service routines. SMP CONSIDERATIONS This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions. RETURNS OK, or ERROR if the semaphore ID is invalid or the task timed out. ERRNO S\_intLib\_NOT\_ISR\_CALLABLE Routine was called from an ISR. S\_objLib\_OBJ\_ID\_ERROR Semaphore ID is invalid.

SEE ALSO	semLib, semBLib, semCLib, semMLib, semRWLib, semSmLib
	S_objLib_OBJ_UNAVAILABLE Would have blocked but NO_WAIT was specified.
	S_objLib_OBJ_TIMEOUT Timeout occured while pending on sempahore.

## semUnlink()

NAME	<pre>semUnlink() – unlink a named semaphore</pre>
SYNOPSIS	STATUS semUnlink ( const char * name /* name of semaphore to unlink */ )
DESCRIPTION	This routine removes a semaphore from the name space, and marks it as ready for deletion on the last <b>semClose()</b> . In case there are already no outstanding <b>semOpen()</b> calls, the semaphore is deleted. After a semaphore is unlinked, subsequent calls to <b>semOpen()</b> using <i>name</i> will not be able to find the semaphore, even if it has not been deleted yet. Instead, a new semaphore could be created if <b>semOpen()</b> is called with the <b>OM_CREATE</b> flag.
	This routine is not ISR callable.
SMP CONSIDERATIO	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	OK, or ERROR if unsuccessful.
ERRNO	S_objLib_OBJ_INVALID_ARGUMENT <i>name</i> is NULL or empty.
	S_objLib_OBJ_NOT_FOUND No semaphore with <i>name</i> was found.
	S_objLib_OBJ_OPERATION_UNSUPPORTED Semaphore is not named.
	S_objLib_OBJ_DESTROY_ERROR Error while deleting the semaphore.
	S_intLib_NOT_ISR_CALLABLE This routine must not be called from an ISR.

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VxWorks Kernel API Reference, 6.6 semWTake()

SEE ALSO semOpen, semOpen(), semClose()

## semWTake()

NAME	<b>semWTake( )</b> – take a semaphore in write mode
SYNOPSIS	<pre>STATUS semWTake   (    SEM_ID semId, /* semaphore ID to take */    int timeout /* timeout in ticks */   )</pre>
DESCRIPTION	Takes the semaphore. If the semaphore is not available, i.e., it is held in either "read" or "write" mode by another task, this task will become pended until the semaphore becomes available. If the semaphore is already available this call will take the semaphore and continue running.
	After a successful call to this routine the caller is granted exclusive access to the resource.
	This routine may be called recursively. However, it should not be called by a task that holds the semaphore in "read" mode. Calling <b>semWTake()</b> in such circumstances will result in a return value of <b>ERROR</b> .
	If deletion safe option is enabled, an implicit <b>taskSafe()</b> operation will occur.
	If priority inversion safe option is enabled, and the calling task blocks, and the priority of the calling task is greater than the semaphore owner, the owner will inherit the caller's priority.
SMP CONSIDERATIO	NS
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
WARNING	This routine must not be used from interrupt level.
RETURNS	OK, or ERROR if the semaphore ID is invalid or the task timed out
ERRNO	S_intLib_NOT_ISR_CALLABLE Routine was called from an ISR.
	S_objLib_OBJ_ID_ERROR Semaphore ID is invalid.
	<b>S_objLib_OBJ_TIMEOUT</b> Timeout occured while pending on sempahore.

**S\_semLib\_INVALID\_OPERATION** Task already holds the semaphore as a writer.

SEE ALSO semRWLib

# sem\_close()

NAME	<pre>sem_close() - close a named semaphore (POSIX)</pre>
SYNOPSIS	<pre>int sem_close   (    sem_t * sem /* semaphore descriptor */   )</pre>
DESCRIPTION	This routine is called to indicate that the calling task is finished with the specified named semaphore, <i>sem</i> . It deallocates any system resources allocated by the system for use by this task for this semaphore. Calling <b>sem_close()</b> with an unnamed semaphore will result in an <b>EINVAL</b> error.
	If the semaphore has not been removed with a call to <b>sem_unlink()</b> , then <b>sem_close()</b> has no effect on the state of the semaphore. However, if the semaphore has been unlinked, it is destroyed when the last reference to it is closed.
WARNING	Take care to avoid risking the deletion of a semaphore that another task has already locked. Applications should only close semaphores that the closing task has opened.
SMP CONSIDERATIO	NS
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	0 (OK), or -1 (ERROR) if unsuccessful.
ERRNO	EINVAL – invalid semaphore descriptor – the semaphore is unnamed
SEE ALSO	semPxLib, sem_unlink( ), sem_open( ), sem_init( )

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## sem\_destroy()

NAME	<pre>sem_destroy() - destroy an unnamed semaphore (POSIX)</pre>	
SYNOPSIS	<pre>int sem_destroy   (    sem_t * sem /* semaphore descriptor */ )</pre>	
DESCRIPTION	This routine is used to destroy the unnamed semaphore indicated by <i>sem</i> .	
	The <b>sem_destroy()</b> call can only destroy a semaphore created by <b>sem_init()</b> . Calling <b>sem_destroy()</b> with a named semaphore will cause an <b>EINVAL</b> error. Subsequent use of the <i>sem</i> semaphore after destruction will cause an <b>EINVAL</b> error.	
	If one or more tasks is blocked on the semaphore, the semaphore is not destroyed, and the routine returns with <b>EBUSY</b> error.	
WARNING	Take care when deleting semaphores, particularly those used for mutual exclusion, to avoid deleting a semaphore out from under a task that has already locked that semaphore. Applications should adopt the protocol of only deleting semaphores that the deleting task has successfully locked.	
SMP CONSIDERATIONS		
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.	
RETURNS	0 ( <b>OK</b> ), or -1 ( <b>ERROR</b> ) if unsuccessful.	
ERRNO	EINVAL – invalid semaphore descriptor – the specified semaphore, <i>sem</i> , is named EBUSY – one or more tasks is blocked on the semaphore	
SEE ALSO	semPxLib, sem init()	

## sem\_getvalue()

**NAME sem\_getvalue()** – get the value of a semaphore (POSIX)

SYNOPSIS int sem\_getvalue

(

sem\_t \* sem, /\* semaphore descriptor \*/

int \* sval /\* buffer by which the value is returned \*/
)

**DESCRIPTION** This routine updates the location referenced by the *sval* argument to have the value of the semaphore referenced by *sem* without affecting the state of the semaphore. The updated value represents an actual semaphore value that occurred at some unspecified time during the call, but may not be the actual value of the semaphore when it is returned to the calling task.

If *sem* is locked, the value returned by **sem\_getvalue()** will either be zero or a negative number whose absolute value represents the number of tasks waiting for the semaphore at some unspecified time during the call.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

SEE ALSO	<pre>semPxLib, sem_post( ), sem_trywait( ), sem_trywait( )</pre>
ERRNO	EINVAL – invalid semaphore descriptor – invalid sval pointer
RETURNS	0 (OK), or -1 (ERROR) if unsuccessful.

### sem\_init( )

NAME	<pre>sem_init() - initialize an unnamed semaphore (POSIX)</pre>
Synopsis	<pre>int sem_init   (    sem_t * sem, /* semaphore to be initialized */    int pshared, /* RTP sharing :ignored */    unsigned int value /* semaphore initialization value */   )</pre>
DESCRIPTION	This routine is used to initialize the unnamed semaphore <i>sem</i> . The value of the initialized semaphore is <i>value</i> . Following a successful call to <b>sem_init()</b> the semaphore may be used in subsequent calls to <b>sem_wait()</b> , <b>sem_trywait()</b> , and <b>sem_post()</b> . This semaphore remains usable until the semaphore is destroyed.
	The value of <i>pshared</i> parameter is ignored.

Only *sem* itself maybe used for performing synchronization. The result of referring to copies of *sem* in calls to sem\_wait, **sem\_trywait()**, **sem\_post()** and **sem\_destroy()** is undefined.

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RETURNS	0 ( <b>OK</b> ), or -1 ( <b>ERROR</b> ) if unsuccessful.
ERRNO	EINVAL – <i>value</i> exceeds SEM_VALUE_MAX – <i>sem</i> points to an invalid buffer ENOSPC – unable to initialize semaphore due to resource constraints
SEE ALSO	<pre>semPxLib, sem_wait( ), sem_trywait( ), sem_post( ), sem_destroy( )</pre>

#### sem\_open()

NAME sem\_open() - initialize/open a named semaphore (POSIX) SYNOPSIS sem\_t \* sem\_open ( const char \* name, /\* semaphore name \*/ int oflag, /\* semaphore creation flags \*/

DESCRIPTION

This routine establishes a connection between a named semaphore and a task. Following a call to **sem\_open()** with a semaphore name *name*, the task may reference the semaphore associated with *name* using the address returned by this call. This semaphore may be used in subsequent calls to **sem\_wait()**, **sem\_trywait()**, and **sem\_post()**. The semaphore remains usable until the semaphore is closed by a successful call to **sem\_close()**.

The *oflag* argument controls whether the semaphore is created or merely accessed by the call to **sem\_open()**. The following flag bits may be set in *oflag*:

/\* extra optional parameters \*/

#### O\_CREAT

•••• )

Use this flag to create a semaphore if it does not already exist. If **O\_CREAT** is set and the semaphore already exists, **O\_CREAT** has no effect except as noted below under **O\_EXCL**. Otherwise, **sem\_open()** creats a semaphore. **O\_CREAT** requires a third and fourth argument: *mode*, which is of type mode\_t, and *value*, which is of type unsigned int. *mode* has no effect in this implementation. The semaphore is created with an initial value of *value*. Valid initial values for semaphores must be less than or equal to **SEM\_VALUE\_MAX**.

#### O\_EXCL

If O\_EXCL and O\_CREAT are set, **sem\_open()** will fail if the semaphore name exists. If O\_EXCL is set and O\_CREAT is not set, the named semaphore is not created.

To determine whether a named semaphore already exists in the system, call **sem\_open()** with the flags **O\_CREAT | O\_EXCL**. If the **sem\_open()** call fails, the semaphore exists.

The semaphore must have a name. **NULL** or empty strings result in **EINVAL**. If the semaphore *name* begins with the slash character, then it is treated as a public semaphore.

	RTPs can open their own references to the public semaphore by using its name. If the <i>name</i> does not begin with the slash character, then it is treated as a private semaphore and RTPs cannot get access to it.
	If a task makes multiple calls to <b>sem_open()</b> with the same value for <i>name</i> , then a reference to the same semaphore is returned for each such call, provided that there have been no calls to <b>sem_unlink()</b> for this semaphore.
	References to copies of the semaphore will produce undefined results.
NOTE	The current implementation has the following limitations:
	<ul> <li>A semaphore cannot be closed with calls to _exit() or exec().</li> <li>A semaphore cannot be implemented as a file.</li> <li>Semaphore names will not appear in the file system.</li> </ul>
RETURNS	A pointer to sem_t, or -1 (ERROR) if unsuccessful.
ERRNO	EEXIST - O_CREAT and O_EXCL are set and the semaphore already exists EINVAL – <i>value</i> exceeds SEM_VALUE_MAX – the semaphore name is invalid ENAMETOOLONG – the semaphore name is too long ENOENT – the named semaphore does not exist and O_CREAT is not set ENOSPC – the semaphore could not be initialized due to resource constraints
SEE ALSO	<pre>semPxLib, sem_unlink(), sem_close()</pre>

# sem\_post()

NAME	<pre>sem_post() - unlock (give) a semaphore (POSIX)</pre>
SYNOPSIS	<pre>int sem_post   (     sem_t * sem /* semaphore descriptor */ )</pre>
DESCRIPTION	This routine unlocks the semaphore referenced by <i>sem</i> by performing the semaphore unlock operation on that semaphore.

	If the semaphore value resulting from the operation is positive, then no tasks were blocked waiting for the semaphore to become unlocked; the semaphore value is simply incremented.	
	If the value of the semaphore resulting from this semaphore is zero, then one of the tasks blocked waiting for the semaphore will return successfully from its call to <b>sem_wait()</b> .	
NOTE	The _POSIX_PRIORITY_SCHEDULING functionality is not yet supported.	
	Note that the POSIX terms <i>unlock</i> and <i>post</i> correspond to the term <i>give</i> used in other VxWorks semaphore documentation.	
SMP CONSIDERATIONS		
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.	
RETURNS	0 ( <b>OK</b> ), or -1 ( <b>ERROR</b> ) if unsuccessful.	
ERRNO	EINVAL – invalid semaphore descriptor	
SEE ALSO	<pre>semPxLib, sem_wait( ), sem_trywait( )</pre>	

## sem\_timedwait()

NAME	<b>sem_timedwait( )</b> – lock (take) a semaphore with a timeout (POSIX)
SYNOPSIS	<pre>int sem_timedwait   (     sem_t * sem,     const struct timespec * abs_timeout   )</pre>
DESCRIPTION	This routine locks the semaphore referenced by <i>sem</i> . If the semaphore cannot be locked immediately, the calling process will wait till the absolute time specified by <i>abs_timeout</i> passes. If the semaphore cannot be locked before <i>abs_timeout</i> has passed, an error is returned.
	Upon successful return, the state of the semaphore is always locked (either as a result of this call or by a previous <b>sem_wait()</b> or <b>sem_trywait()</b> ). The semaphore will remain locked until <b>sem_post()</b> is executed and returns successfully.
	Deadlock detection is not implemented.

Note that the POSIX term lock corresponds to the term take used in other VxWorks semaphore documentation. SMP CONSIDERATIONS This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions. RETURNS 0 (OK), or -1 (ERROR) if unsuccessful. ERRNO **ETIMEDOUT** The semaphore could not be locked before the timeout expired. EINVAL The semaphore descriptor is invalid, or the nanosecond field of the timeout value is greater than 1 billion. EINTR A signal interrupted this function. SEE ALSO semPxLib, sem\_wait( ), sem\_trywait( ), sem\_post( )

#### sem\_trywait()

NAME	<pre>sem_trywait() - lock (take) a semaphore, returning error if unavailable (POSIX)</pre>
SYNOPSIS	<pre>int sem_trywait   (    sem_t * sem /* semaphore descriptor */ )</pre>
DESCRIPTION	This routine locks the semaphore referenced by <i>sem</i> only if the semaphore is currently not locked; that is, if the semaphore value is currently positive. Otherwise, it does not lock the semaphore. In either case, this call returns immediately without blocking.
	Upon successful return, the state of the semaphore is always locked (either as a result of this call or by a previous <b>sem_wait()</b> or <b>sem_trywait()</b> ). The semaphore will remain locked until <b>sem_post()</b> is executed and returns successfully.
	Deadlock detection is not implemented.
	Note that the POSIX term <i>lock</i> corresponds to the term <i>take</i> used in other VxWorks semaphore documentation.

VxWorks Kernel API Reference, 6.6 sem\_unlink()

#### SMP CONSIDERATIONS

SEE ALSO	<pre>semPxLib, sem_wait( ), sem_post( )</pre>
ERRNO	EAGAIN – semaphore is already locked EINVAL – invalid semaphore descriptor
RETURNS	0 ( <b>OK</b> ), or -1 ( <b>ERROR</b> ) if unsuccessful.
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

## sem\_unlink()

NAME	<pre>sem_unlink() – remove a named semaphore (POSIX)</pre>
SYNOPSIS	<pre>int sem_unlink   (    const char * name /* semaphore name */   )</pre>
DESCRIPTION	This routine removes the string <i>name</i> from the semaphore name table, and marks the corresponding semaphore for destruction. An unlinked semaphore is destroyed when the last reference to it is removed by <b>sem_close()</b> . After a name is unlinked, calls to <b>sem_open()</b> using the same name cannot connect to the same semaphore, even if other tasks are still using it. Instead, such calls refer to a new semaphore with the same name.
SMP CONSIDERATIO	NS
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	0 ( <b>OK</b> ), or -1 ( <b>ERROR</b> ) if unsuccessful.
ERRNO	ENAMETOOLONG – semaphore name too long ENOENT – a semaphore with the specified <i>name</i> does not exist
SEE ALSO	semPxLib, sem_open(), sem_close()

# sem\_wait()

NAME	<pre>sem_wait() - lock (take) a semaphore, blocking if not available (POSIX)</pre>
SYNOPSIS	<pre>int sem_wait   (    sem_t * sem /* semaphore descriptor */   )</pre>
DESCRIPTION	This routine locks the semaphore referenced by <i>sem</i> by performing the semaphore lock operation on that semaphore. If the semaphore value is currently zero, the calling task will not return from the call to <b>sem_wait()</b> until it either locks the semaphore or the call is interrupted by a signal.
	On return, the state of the semaphore is locked and will remain locked until <b>sem_post()</b> is executed and returns successfully.
	Deadlock detection is not implemented.
	Note that the POSIX term <i>lock</i> corresponds to the term <i>take</i> used in other VxWorks documentation regarding semaphores.
SMP CONSIDERATIO	NS
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	0 (OK), or -1 (ERROR) if unsuccessful.
ERRNO	EINVAL – invalid semaphore descriptor EINTR – got a signal while blocking on the semaphore. Applicable only for RTP task
SEE ALSO	semPxLib, sem_trywait( ), sem_post( )

# set\_new\_handler()

NAME	<pre>set_new_handler() - set new_handler to user-defined function (C++)</pre>
SYNOPSIS	<pre>extern void (*set_new_handler (void(* pNewNewHandler)())) (void)</pre>
DESCRIPTION	This function is used to define the function that will be called when operator new cannot allocate memory.

VxWorks Kernel API Reference, 6.6 set\_terminate()

	The new_handler acts for all threads in the system; you cannot set a different handler for different tasks.
RETURNS	A pointer to the previous value of new_handler.
ERRNO	Not Available
SEE ALSO	cplusLib

## set\_terminate()

NAME	<b>set_terminate()</b> – set terminate to user-defined function (C++)
SYNOPSIS	<pre>extern void (*set_terminate (void(* terminate_handler)())) (void)</pre>
DESCRIPTION	This function is used to define the terminate_handler which will be called when an uncaught exception is raised.
	The terminate_handler acts for all threads in the system; you cannot set a different handler for different tasks.
RETURNS	The previous terminate_handler.
ERRNO	Not Available
SEE ALSO	cplusLib

# shConfig()

NAME	<b>shConfig( )</b> – display or set the shell configuration
SYNOPSIS	<pre>void shConfig   (     const char * config /* configuration string */ )</pre>
DESCRIPTION	This routine displays or sets the shell configuration of the current shell session. If <i>config</i> is <b>NULL</b> , the routine displays the cofiguration variables; otherwise, it sets the configuration. The format of the string <i>config</i> is:

	<variable> = <value> , <variable> = <value> ,</value></variable></value></variable>
	The variable name or value can contain , and = if these characters are escaped or quoted.
RETURNS	N/A.
ERRNO	N/A
SEE ALSO	usrLib, shellConfigGet(), shellConfigSet(), the VxWorks programmer guides.

## shellAbort()

NAME	shellAbort() – abort a shell session
SYNOPSIS	STATUS shellAbort ( SHELL_ID shellId /* shell session Id */ )
DESCRIPTION	This routine aborts the shell session <i>shellId</i> . Before the shell task is restarted, its task trace is printed.
	It <i>shellld</i> is equal to <b>ALL_SHELL_SESSIONS</b> , all the shell session are aborted. If <i>shellld</i> is equal to <b>CURRENT_SHELL_SESSION</b> , the current shell session is aborted.
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>shellId</i> is not a valid shell session Id.
ERRNO	N/A
SEE ALSO	shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

## shellCmdAdd()

NAME	<pre>shellCmdAdd() - add a shell command</pre>
SYNOPSIS	<pre>STATUS shellCmdAdd  (     const char * topic, /* topic name */     const SHELL_CMD * pShellCmd /* pointer on the command structure */   )</pre>

VxWorks Kernel API Reference, 6.6 shellCmdAdd()

 

 DESCRIPTION
 This routine adds the shell command pointed by *pShellCmd*. The command is added under the topic *topic*. If the topic does not already exist, an error is returned.

 COMMAND FORMAT
 A command is defined as a structure: {*cmdFullname, func, opt, shortDesc, fullDesc, synopsis*} *cmdFullname* is the name of the command to add. It may be a composed command name, that means a command name like "foo bar". In that example, "foo" is the top command name, and "bar" is a sub-command of "foo".

 *func* is the function to call for that command name.
 If the *opt* parameter is not equal to NULL, the declaration of that function is: int func (

SHELL\_OPTION options[] /\* options array \*/

options is a pointer on the argument array of the command.

If the *opt* parameter is **NULL**, the declaration of the function is:

*opt* is a string that describes the possible options that the command accepts. Each option is a single character (case sensitive) which defines a single option. If an option expects an extra argument, an extra : character has to be added after the option character. For example, "avf:" means that the command accepts options "-a", "-v" and "-f *extraArg*". The order of the option character matters: it defines the order of the options in the option array passed to *func*.

Each cell of the options array passed to *func* is composed of a boolean value (**TRUE** if the option is set, **FALSE** otherwise) and of a pointer on a string (pointer on an extra argument). Another boolean indicates if it is the last cell of the array.

If the option string *opt* is ":", the argument string of the command is passed to *func* without any process, as the string field of the first cell of the options array.

*shortDesc* is a short description of the command. A sequence of "%s" characters will be replaced by the function name. The string should not be ended by a \n character.

*fullDesc* is the full description of the command. A sequence of "%s" characters will be replaced by the function name. This description should contain the explanation of the command options. The string should not be ended by a \n character.

*synopsis* is the synopsis of the command. A sequence of "%s" characters will be replaced by the function name. The string should not be ended by a \n character.

**RETURNS** OK, or ERROR if an error occured

ERRNO	S_shellInterpCmdLib_UNKNOWN_TOPIC The topic is not registered yet.
	S_shellLib_UNMATCHED_QUOTE A quote character is missing in the command name
	S_shellInterpCmdLib_WRONG_CMD A wrong command name was supplied
	malloc(), calloc() and memPartCreate() errnos.
SEE ALSO	<pre>shellInterpCmdLib, shellCmdTopicAdd(), shellCmdArrayAdd().</pre>

# shellCmdAliasAdd()

NAME	shellCmdAliasAdd() – add an alias string
SYNOPSIS	<pre>STATUS shellCmdAliasAdd  (     const char * aliasName, /* alias name */     const char * string, /* string aliased */     BOOL allocate /* TRUE to allocate memory */  )</pre>
DESCRIPTION	This routine adds the alias string <i>aliasName</i> . During the parsing of the shell imput line by the command interpreter, this <i>aliasName</i> string will be replaced by <i>string</i> .
	If <i>allocate</i> is set to <b>FALSE</b> , both strings are not copied into an internal database; only their pointers are stored for future use. If <i>allocate</i> is <b>TRUE</b> , a buffer is allocated to store the strings <i>aliasName</i> and <i>string</i> . This buffer is freed when the alias is removed.
NOTE	It is not possible to add an alias if a command with the same name exists.
RETURNS	OK, or ERROR if the alias cannot be added.
ERRNO	N/A
SEE ALSO	shellInterpCmdLib, shellCmdAliasArrayAdd(), shellCmdAliasDelete().

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## shellCmdAliasArrayAdd()

NAME	<pre>shellCmdAliasArrayAdd() - add an array of alias strings</pre>
SYNOPSIS	STATUS shellCmdAliasArrayAdd ( const SHELL_CMD_ALIAS aliasArray[], /* array of aliases */ BOOL allocate /* TRUE to allocate memory */ )
DESCRIPTION	This routine adds the alias stored in the array <i>aliasArray</i> . The end of the array is defined by an entry with the alias name set to NULL.
	If <i>allocate</i> is set to <b>FALSE</b> , both strings are not copied into an internal database; only their pointers are stored for future use. If <i>allocate</i> is <b>TRUE</b> , a buffer is allocated to store the strings <i>aliasName</i> and <i>string</i> . This buffer is freed when the alias is removed.
NOTE	It is not possible to add an alias if a command with the same name exists.
RETURNS	OK, or ERROR if one alias cannot be added.
ERRNO	N/A

SEE ALSO shellInterpCmdLib, shellCmdAliasAdd()

## shellCmdAliasDelete()

NAME	<pre>shellCmdAliasDelete( ) - delete an alias</pre>
SYNOPSIS	<pre>STATUS shellCmdAliasDelete   (     const char * alias /* alias to delete */ )</pre>
DESCRIPTION	This routine deletes the alias <i>alias</i> .
RETURNS	OK, or ERROR if the alias cannot be found.
ERRNO	N/A
SEE ALSO	shellInterpCmdLib, shellCmdAliasAdd()

### shellCmdArrayAdd()

shellCmdArrayAdd() - add an array of shell commands NAME SYNOPSIS STATUS shellCmdArrayAdd ( const char \* topic, /\* topic name \*/ const SHELL\_CMD shellCmdArray[] /\* array of commands \*/ ) DESCRIPTION This routine adds the list of shell commands stored in the array pointed by *pShellCmdArray*. These commands are added under the topic topic. If the topic does not already exist, an error is returned. **COMMAND ARRAY FORMAT** An element of the array is a command structure as described by the routine **shellCmdAdd()**. The end of the array is marked by a *cmdFullname* equals to NULL. OK, or ERROR if one command cannot be added. RETURNS S\_shellInterpCmdLib\_UNKNOWN\_TOPIC ERRNO The topic is not registered yet. S\_shellLib\_UNMATCHED\_QUOTE A quote character is missing in the command name

> S\_shellInterpCmdLib\_WRONG\_CMD A wrong command name was supplied

malloc(), calloc() and memPartCreate() errno.

SEE ALSO shellInterpCmdLib, shellCmdTopicAdd(), shellCmdAdd().

### shellCmdExec( )

<pre>shellCmdExec() – execute a shell command</pre>
<pre>STATUS shellCmdExec   (     const char * name, /* command name to execute */     const char * args /* arguments of the command */   )</pre>

VxWorks Kernel API Reference, 6.6 shellCmdMemRegister()

DESCRIPTION	This routine executes the shell command <i>name</i> with the argument string <i>args</i> . The argument string <i>args</i> is parsed to extract all arguments and options.
	The arguments are separated by blank characters. If a blank character has to be set as part of an argument, the argument needs to be quoted (with simple or double quote) or the space has to be escaped (with the "\" character). The argument string is parsed using the option string set for the command.
	If the pointer to the option string of the command is <b>NULL</b> , the arguments are only split into an array of strings. This array and the number of strings is passed to the command function (as argc/argv parameters).
	This routine has to be called from within a shell task. It cannot be called by another task.
RETURNS	OK, or ERROR if the command cannot be executed.
ERRNO	S_shellInterpCmdLib_UNKNOWN_CMD The command name <i>name</i> is unknown.
	S_shellLib_UNKNOWN_OPT A options of the command is not valid.
	S_shellLib_UNMATCHED_QUOTE A quote string is not ended with a quote character.
	S_shellLib_MISSING_ARG An option is missing its extra argument.
	malloc() and memPartCreate() errno.
SEE ALSO	shellInterpCmdLib

# shellCmdMemRegister()

NAME	<pre>shellCmdMemRegister() - register a buffer against the command interpreter</pre>
SYNOPSIS	STATUS shellCmdMemRegister ( void * pMem, /* memory block to register */
	BOOL shellPool /* TRUE for shell memory pool */ )
DESCRIPTION	This routine registers the memory block pointed by $pMem$ into the memory list of the command interpreter. When the shell task is restarted or ended, the registered memory blocks are freed by the shell.
	If the memory block has been allocated from the shell memory pool using <b>shellMemMalloc()</b> or <b>shellMemCalloc()</b> routines, <i>shellPool</i> must be set to <b>TRUE</b> . Otherwise, <i>shellPool</i> must be set to <b>FALSE</b> .
----------	--
RETURNS	OK, or ERROR if an error occured.
ERRNO	<pre>S_shellLib_NO_SHELL_CMD The current interpreter of the current shell session is not the command interpreter. malloc() and memPartCreate() errno.</pre>
SEE ALSO	shellInterpCmdLib, shellCmdMemRegister()

# shellCmdMemUnregister()

NAME	shellCmdMemUnregister() – unregister a buffer
SYNOPSIS	STATUS shellCmdMemUnregister ( void * pMem /* memory block to unregister */ )
DESCRIPTION	This routine unregisters the memory block pointed by $pMem$ from the memory list of the command interpreter.
RETURNS	OK, or ERROR if the memory was not registered
ERRNO	S_shellLib_NO_SHELL_CMD The current interpreter of the current shell session is not the command interpreter.
SEE ALSO	shellInterpCmdLib, shellCmdMemRegister().

## shellCmdPreParseAdd()

NAME	<pre>shellCmdPreParseAdd() – defir</pre>	ne a comn	nand	to be pre	e-pars	ed		
SYNOPSIS	STATUS shellCmdPreParseAdd							
	const char *	name,	/*	command	name	to	pre-parse	*/

```
char * (*preParseRtn) (char * line) /* pre-parse routine address */
)
```

**DESCRIPTION** This routine is used to define that the command *name* needs a special handling. During the command line parsing, if *name* string is found at the beginning of a line, the routine *preParseRtn* is called before any processing by the interpreter. *name* must be a single command name.

The prototype of *preParseRtn* is:

char \* preParseRtn
{
 char \* line /\* complete line to be parsed \*/
};

The *line* string is the line parsed by the interpreter, the initial blank characters being stripped. This string begins with the command *name*.

If *preParseRtn* returns **NULL**, the interpreter interupts the line parsing. Otherwise, it continues the parsing of the returned pre-parsed line.

The returned buffer must be allocated by the routine *preParseRtn* using either a **malloc()**, **calloc()** or **strdup()** call. It will be freed by the interpreter using **free()**.

- **RETURNS** OK, or ERROR if the command *name* is not registered yet
- ERRNO N/A
- SEE ALSO shellInterpCmdLib

### shellCmdSymTabIdGet()

NAME	<pre>shellCmdSymTabIdGet() - get symbol table Id of a shell session</pre>
SYNOPSIS	SYMTAB_ID shellCmdSymTabIdGet ( SHELL_ID shellId )
DESCRIPTION	This routine returns the symbol table Id associated to the current memory context of the shell session <i>shellId</i> . <i>shellId</i> can be equal to <b>CURRENT_SHELL_SESSION</b> .
RETURNS	the symbol table Id, or 0 if the symbol table is not available
ERRNO	N/A

SEE ALSO shellInterpCmdLib

# shellCmdTopicAdd()

NAME	<pre>shellCmdTopicAdd() - add a shell command topic</pre>
SYNOPSIS	<pre>STATUS shellCmdTopicAdd   (     const char * topic, /* topic name */     const char * desc /* topic description */   )</pre>
DESCRIPTION	This routine adds the topic <i>topic</i> with the description <i>desc</i> to the list of topics.
NOTE	The topic name and the description are not copied internaly. Only the string pointers are stored in the internal structure.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the topic cannot be added.
ERRNO	N/A
SEE ALSO	shellInterpCmdLib

# shellCompatibleCheck()

NAME	<pre>shellCompatibleCheck() - check the compatibility mode of the shell</pre>
SYNOPSIS	BOOL shellCompatibleCheck (void)
DESCRIPTION	This routine checks if the shell is configured to be VxWorks 5.5 compatible or not.
RETURNS	TRUE if the shell is VxWorks 5.5 compatible, FALSE otherwise.
ERRNO	N/A
SEE ALSO	shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

VxWorks Kernel API Reference, 6.6 shellConfigDefaultGet()

# shellConfigDefaultGet( )

NAME	<pre>shellConfigDefaultGet( ) - get default shell configuration</pre>
SYNOPSIS	char * shellConfigDefaultGet (void)
DESCRIPTION	This routine returns the default shell configuration variables. The format of the string returned is:
	<variable>=<value>,<variable>=<value>,</value></variable></value></variable>
	The returned string has to be freed by the caller with a <b>free()</b> .
RETURNS	a pointer on the default shell configuration string, or NULL if an error occured.
ERRNO	N/A
SEE ALSO	shellConfigLib, shellConfigGet(), VxWorks Kernel Programmer's Guide: Kernel Shell

# shellConfigDefaultSet( )

NAME	<pre>shellConfigDefaultSet() - set default shell configuration</pre>
SYNOPSIS	<pre>STATUS shellConfigDefaultSet   (    const char * config /* default configuration string */ )</pre>
DESCRIPTION	This routine sets the default values of the shell configuration variables. The format of the string <i>config</i> is:
	<variable> = <value> , <variable> = <value> ,</value></variable></value></variable>
	The variable name or value can contain "," and "=" characters if they are escaped or quoted.
	These configuration variables are seen by all shell sessions and can be superseed by a configuration variable defined localy to a shell session.
RETURNS	OK, ERROR if the configuration cannot be set or there was a problem
ERRNO	S_shellLib_CONFIG_ERROR The format of the configuration string is wrong.
	malloc() and memPartCreate() errno.

# shellConfigDefaultValueGet( )

NAME	<pre>shellConfigDefaultValueGet() – get a default configuration variable value</pre>
SYNOPSIS	<pre>const char * shellConfigDefaultValueGet   (    const char * name /* variable name */ )</pre>
DESCRIPTION	This routine gets the value of the default configuration variable <i>name</i> . If the variable does not exist, <b>NULL</b> is returned.
RETURNS	a pointer on the string value, or NULL if the variable does not exist or if there was a problem.
ERRNO	N/A
SEE ALSO	shellConfigLib, shellConfigValueGet(), VxWorks Kernel Programmer's Guide: Kernel Shell

# shellConfigDefaultValueSet()

NAME	<pre>shellConfigDefaultValueSet( ) - set a default configuration variable value</pre>
SYNOPSIS	STATUS shellConfigDefaultValueSet ( const char * name, /* variable name */ const char * value /* variable value or NULL */ )
DESCRIPTION	This routine sets the value of the default configuration variable named <i>name</i> . If the variable does not exist, it is added to the default configuration variable list. The strings <i>name</i> and <i>values</i> are copied into memory. <i>value</i> can be a <b>NULL</b> pointer.
RETURNS	OK, ERROR if the configuration cannot be set or there was a problem
ERRNO	N/A
SEE ALSO	shellConfigLib, shellConfigValueSet(), VxWorks Kernel Programmer's Guide: Kernel Shell

# shellConfigDefaultValueUnset()

NAME	<pre>shellConfigDefaultValueUnset() – unset a default configuration variable value</pre>
SYNOPSIS	<pre>void shellConfigDefaultValueUnset   (     const char * name /* variable name */   )</pre>
DESCRIPTION	This routine unsets the default configuration variable <i>name</i> .
RETURNS	N/A
ERRNO	N/A
SEE ALSO	shellConfigLib, shellConfigValueUnset(), VxWorks Kernel Programmer's Guide: Kernel Shell

# shellConfigGet()

NAME	<pre>shellConfigGet() - get the shell configuration</pre>
SYNOPSIS	<pre>char * shellConfigGet   (    SHELL_ID shellId /* shell session Id */   )</pre>
DESCRIPTION	This routine returns the shell configuration variables of the shell session <i>shellId</i> . The format of the string returned is:
	<variable>=<value>,<variable>=<value>,</value></variable></value></variable>
	<i>shellId</i> can be <b>CURRENT_SHELL_SESSION</b> .
	The returned string has to be freed by the caller with a <b>free()</b> .
RETURNS	a pointer on the shell configuration string, or NULL if an error occured.
ERRNO	S_shellLib_NOT_SHELL_TASK The shell session is invalid.
	<pre>shellDataAdd(), shellDataGet(), malloc() and memPartCreate() errno.</pre>

# SEE ALSO shellConfigLib, shellConfigDefaultGet(), shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

# shellConfigSet()

NAME	<pre>shellConfigSet() - set shell configuration</pre>
SYNOPSIS	<pre>STATUS shellConfigSet   (    SHELL_ID shellId, /* shell session identifier */    const char * config /* default configuration string */   )</pre>
DESCRIPTION	This routine sets the values of the shell configuration variables of the shell session <i>shellId</i> . The format of the string <i>config</i> is:
	<variable> = <value> , <variable> = <value> ,</value></variable></value></variable>
	The variable name or value can contain "," and "=" characters if they are escaped or quoted.
	The configuration variables will only be seen by the shell session <i>shellId</i> . They superseed any default configuration values with same name.
	shellId can be CURRENT_SHELL_SESSION.
RETURNS	<b>OK</b> , <b>ERROR</b> if the configuration cannot be set or if <i>shellId</i> is not a valid shell session identifier.
ERRNO	S_shellLib_CONFIG_ERROR The format of the configuration string is wrong.
	S_shellLib_NOT_SHELL_TASK The shell session is invalid.
	<pre>shellDataAdd(), shellDataGet(), malloc() and memPartCreate() errno.</pre>
SEE ALSO	shellConfigLib, shellConfigDefaultSet(), shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

# shellConfigValueGet()

NAME shellConfigValueGet() – get a shell configuration variable value

SYNOPSIS const char \* shellConfigValueGet

VxWorks Kernel API Reference, 6.6 shellConfigValueSet()

(

```
      SHELL_ID
      shellId, /* shell session Id */

      const char * name
      /* variable name */

      )
      This routine gets, from the shell session shellId, the value of the configuration variable name.

      If the variable does not exist or if shellId is not a valid shell session, NULL is returned.

      shellId can be CURRENT_SHELL_SESSION.

      a pointer on the string value, or NULL if the variable does not exist or an error occurred.

      ERRNO
      N/A

      SEE ALSO
      shellConfigLib, shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell
```

## shellConfigValueSet()

NAME	<pre>shellConfigValueSet( ) - set a shell configuration variable value</pre>
SYNOPSIS	<pre>STATUS shellConfigValueSet   (    SHELL_ID shellId, /* shell session Id */    const char * name, /* variable name */    const char * value /* variable value */   )</pre>
DESCRIPTION	This routine sets, for the shell session <i>shellId</i> , the value of the configuration variable <i>name</i> to <i>value</i> . If the variable does not exist, it is added to the internal list of the shell session.
	The strings name and values are copied into memory. value can be a NULL pointer.
	<i>shellId</i> can be <b>CURRENT_SHELL_SESSION</b> .
RETURNS	OK, ERROR if the configuration cannot be set
ERRNO	S_shellLib_NOT_SHELL_TASK The shell session is invalid.
	<pre>shellDataAdd(), shellDataGet(), malloc() and memPartCreate() errno.</pre>
SEE ALSO	shellConfigLib, shellConfigDefaultValueSet(), shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

shellConfigValueUnset()
-------------------------

NAME	<pre>shellConfigValueUnset() - unset a shell configuration variable value</pre>
SYNOPSIS	STATUS shellConfigValueUnset ( SHELL_ID shellId, /* shell session Id */ const char * name /* variable name */ )
DESCRIPTION	This routine unsets, for the shell session <i>shellId</i> , the configuration variable <i>name</i> . <i>shellId</i> can be <b>CURRENT_SHELL_SESSION</b> .
RETURNS	OK, or ERROR if the variable cannot be unset.
ERRNO	S_shellLib_NOT_SHELL_TASK The shell session is invalid. shellDataAdd( ), shellDataGet( ), malloc( ) and memPartCreate( ) errno.
SEE ALSO	shellConfigLib, shellConfigDefaultValueUnset(), shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

# shellDataAdd()

NAME	shellDataAdd() – add user data to a specified shell
SYNOPSIS	<pre>STATUS shellDataAdd   (    SHELL_ID shellId, /* the shell identifier */    const char * key, /* the key for the value to add */    void * pData, /* the value to add */    SHELL_DATA_FUNCPTR finalizeRtn /* finalize routine of the data, or NULL */    )</pre>
DESCRIPTION	This routine adds to the user data list of the shell <i>shellId</i> a new data <i>pData</i> with the key named <i>key</i> . If the key already exists, the previous data is overwritten.
	The routine <i>finalizeRtn</i> is called when the shell session is terminated to free any internal value associated with the user data <i>pData</i> . <i>finalizeRtn</i> may be <b>NULL</b> .
	The prototype of <i>finalizeRtn</i> is <b>SHELL_DATA_FUNCPTR</b> :
	void finalizeRtn

VxWorks Kernel API Reference, 6.6 shellDataFirst()

(		
SHELL_ID	shellId,	/* shell session Id */
const char *	key,	/* data key */
void *	pData	/* data value */
)		

with *shellId* the identifier of the shell session terminated and *pData* the user data associated with the key *key*.

RETURNS	OK, or ERROR if an error occured
ERRNO	S_shellLib_NOT_SHELL_TASK shellId is not a valid shell session.
	malloc() and memPartCreate() errno.
SEE ALSO	shellDataLib, shellFromNameDataAdd(), shellDataRemove(), shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

## shellDataFirst()

NAME	<b>shellDataFirst()</b> – get the first user data that matchs a key
SYNOPSIS	SHELL_ID shellDataFirst ( const char * key, /* key data to search */ void ** ppData /* where to store data value */ )
DESCRIPTION	This routine returns the first user data named <i>key</i> in all shell sessions. The user data value is stored into the location pointed by <i>ppData</i> .
RETURNS	the shell session Id, or 0 if none shell session contains the key key.
ERRNO	N/A
SEE ALSO	shellDataLib, shellDataNext(), shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

#### shellDataFromNameAdd() shellDataFromNameAdd() - add user data to a specified shell NAME SYNOPSIS STATUS shellDataFromNameAdd ( const char \*taskName,/\* the shell task name \*/const char \*key,/\* the key for the valuevoid \*pData,/\* the value to add \*/ /\* the key for the value to add \*/ SHELL\_DATA\_FUNCPTR finalizeRtn /\* finalize routine of the data, or NULL \*/ ) This routine adds to the user data list of shell session associated to the shell task named DESCRIPTION taskName a new data pData with the key named key. If the key already exists, the previous data is overwritten. The routine *finalizeRtn* is called when the shell session is terminated to free any internal value associated with the user data *pData*. *finalizeRtn* may be **NULL**. The prototype of *finalizeRtn* is **SHELL\_DATA\_FUNCPTR**: void finalizeRtn ( SHELL\_IDshellId,/\* shell session Id \*/const char \*key,/\* data key \*/void \*pData/\* data value \*/ with *shellId* the identifier of the shell session terminated and *pData* the user data associated with the key key. OK, or ERROR if an error occured. RETURNS ERRNO S\_shellLib\_NOT\_SHELL\_TASK *taskName* is not a shell task. malloc() and memPartCreate() errno. shellDataLib, shellDataAdd(), shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell SEE ALSO

### shellDataFromNameGet( )

NAME shellDataFromNameGet() – get user data from a specified shell

SYNOPSIS STATUS shellDataFromNameGet

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VxWorks Kernel API Reference, 6.6 shellDataGet()

	( const char * taskName, /* the shell task name */ const char * key, /* the key for the value to get */ void ** ppData /* get the data there */ )
DESCRIPTION	This routine returns the data named <i>key</i> stored within the shell session whose associated shell task is name <i>taskName</i> . The user data value is stored into the location pointed by <i>ppData</i> .
RETURNS	OK, or ERROR if an error occured.
ERRNO	S_shellLib_NOT_SHELL_TASK <i>taskName</i> is not a shell task.
	S_shellLib_NO_USER_DATA Key <i>key</i> does not exist within the shell session context.
SEE ALSO	shellDataLib, shellDataGet(), shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

# shellDataGet( )

NAME	<pre>shellDataGet( ) - get user data from a specified shell</pre>	
SYNOPSIS	<pre>STATUS shellDataGet   (    SHELL_ID shellId, /* the shell identifier */    const char * key, /* the key for the value to get */    void ** ppData /* where to store the data value */   )</pre>	
DESCRIPTION	This routine returns the data named <i>key</i> stored within the shell session <i>shellId</i> . The user data value is stored into the location pointed by <i>ppData</i> .	
RETURNS	<b>OK</b> , or <b>ERROR</b> if an error occured.	
ERRNO	S_shellLib_NOT_SHELL_TASK <i>taskName</i> is not a shell task.	
	S_shellLib_NO_USER_DATA the key name <i>key</i> does not exist within the shell session context <i>shellId</i> .	
SEE ALSO	shellDataLib, shellFromNameDataGet( ), shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell	

# shellDataNext()

NAME	<b>shellDataNext( )</b> – get the next user data that matchs a key
SYNOPSIS	<pre>SHELL_ID shellDataNext (     const char * key,    /* key data to search */     SHELL_ID shellId, /* previous shell session id */     void **    ppData    /* where to store data value */ )</pre>
DESCRIPTION	This routine returns the next user data named <i>key</i> in all shell session. <i>shellId</i> is the previous shell session checked. The user data value is stored into the location pointed by <i>ppData</i> .
RETURNS	the shell session Id, or 0 if none shell session contains the key key.
ERRNO	N/A
SEE ALSO	shellDataLib, shellDataFirst(), shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

# shellDataRemove()

NAME	shellDataRemove() – remove user data from a specified shell
SYNOPSIS	<pre>void shellDataRemove   (    SHELL_ID shellId, /* the shell identifier */    const char * key, /* the key for the value to remove */    BOOL finalize /* TRUE to call the finalize routine */   )</pre>
DESCRIPTION	This routine removes the data defined by the key <i>key</i> from the shell session <i>shellId</i> . If <i>finalize</i> is <b>TRUE</b> , the finalize routine associated with the data is called before the data is removed from the list.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	shellDataLib, shellDataAdd(), shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

# shellErrnoGet( )

NAME	<b>shellErrnoGet( )</b> – get the shell session errno
SYNOPSIS	<pre>int shellErrnoGet   (    SHELL_ID shellId /* shell session Id */ )</pre>
DESCRIPTION	This routine returns the errno value for the shell session <i>shellId</i> . The shell task errno is set to this value before calling a VxWorks function. <i>shellId</i> can be equal to <b>CURRENT_SHELL_SESSION</b> .
RETURNS	the errno value, or -1 if <i>shellId</i> is not a valid shell session
ERRNO	N/A
SEE ALSO	shellLib, shellErrnoSet(), VxWorks Kernel Programmer's Guide: Kernel Shell

# shellErrnoSet( )

NAME	<pre>shellErrnoSet() - set the shell session errno</pre>
SYNOPSIS	<pre>void shellErrnoSet   (   SHELL_ID shellId, /* shell session Id */   int errNo /* errno number */  )</pre>
DESCRIPTION	This routine sets the errno value for the shell session <i>shellId</i> to <i>errNo</i> . The shell task errno is set to this value before calling a VxWorks function; when the VxWorks function returns, the current errno value of the shell task is save using <b>shellErrnoSet()</b> .
RETURNS	N/A
ERRNO	N/A
SEE ALSO	shellLib, shellErrnoGet(), VxWorks Kernel Programmer's Guide: Kernel Shell

## shellFirst()

NAME	<pre>shellFirst() - get the first shell session</pre>
SYNOPSIS	SHELL_ID shellFirst (void)
DESCRIPTION	This routine returns the Id of the first running shell session.
RETURNS	the first shell Id, or 0 if no shell session is running.
ERRNO	N/A
SEE ALSO	<pre>shellLib, shellNext( ), shellFromTaskGet( ), shellFromNameGet( ), VxWorks Kernel Programmer's Guide: Kernel Shell</pre>

## shellFromNameGet()

NAME	<pre>shellFromNameGet() – get a shell session Id from a task name</pre>
SYNOPSIS	SHELL_ID shellFromNameGet ( const char * taskName /* the shell task name */ )
DESCRIPTION	This routine returns the shell session Id of the shell task whose name is <i>taskName</i> .
RETURNS	the shell session Id, or 0 if the task is not a shell task.
ERRNO	N/A
SEE ALSO	shellLib, shellFirst(), shellFromTaskGet(), VxWorks Kernel Programmer's Guide: Kernel Shell

# shellFromTaskGet( )

NAME shellFromTaskGet() – get a shell session Id from its task Id

SYNOPSIS SHELL\_ID shellFromTaskGet

VxWorks Kernel API Reference, 6.6 shellGenericInit()

```
      (int taskId /* the shell task ID or 0 */)

      DESCRIPTION

      This routine returns the shell session identifier whose task identifier is taskId. If taskId is 0, the current task Id is used.

      RETURNS
      the shell session identifier, or 0 if taskId is not a shell task.

      ERRNO
      N/A

      SEE ALSO
      shellLib, shellFirst(), shellFromNameGet(), VxWorks Kernel Programmer's Guide: Kernel Shell
```

## shellGenericInit( )

```
shellGenericInit() – start a shell session
NAME
SYNOPSIS
               STATUS shellGenericInit
                  (
                  const char * config,
                                          /* configuration string or NULL */
                  int stackSize, /* shell stack (0 = default value) */
                  const char * shellName,
                                           /* shell task name or NULL for def. base name
               */
                  char **
                                            /* pointer on the shell task name or NULL */
                             pShellName,
                              interactive, /* interactive mode if TRUE */
                  BOOL
                              loginAccess, /* login access */
                  BOOL
                              fdin,
                  int
                                            /* input file descriptor */
                              fdout,
                                            /* output file descriptor */
                  int
                  int
                               fderr
                                            /* error file descriptor */
                  )
```

DESCRIPTION

This routine starts a shell session. This is a generic routine to start a shell session.

*config* is a string that holds the values of the configuration variables this new shell session. *stackSize* defines the size of the stack allocated for the shell task. A value of 0 is used to define the default value. *shellName* is a pointer on the desired shell task name. If this parameter is **NULL**, a generic name is used for each shell task. *pShellName*, if not **NULL**, will return a pointer on the shell task name. *interpName* is the desired interpreter name to use with this shell session. *interactive* is set to **TRUE** if this shell session is interactive; **FALSE** otherwise. *loginAccess* is set to **TRUE** if one wants the user to identify itself with a login and password before accessing the shell (this feature is only usefull if a login function has been previously installed). *fdin*, *fdout* and *fderr* are the file descriptors to use respectively for the standard input, output and error of the shell task.

	If <i>loginAccess</i> is set to <b>TRUE</b> , and if the <b>INCLUDE_SECURITY</b> component is installed, the user will be asked for a login and a password before the shell start.
NOTE 1	If the shell is configured to use only one shell task for all the connections, (compatibility mode), this routine only changes the I/O of the only shell session, and restarts it. On termination of the shell task (see <b>shellTerminate()</b> ), the previous I/O of the shell task is restored, and the shell task is restarted.
NOTE 2	If the shell is configured to use only one shell task for all the connections, <i>pShellName</i> may return a task name different from the one specified by <i>shellName</i> .
RETURNS	OK, or ERROR if the shell session cannot be created.
ERRNO	S_shellLib_NO_INTERP The interpreter specified is not registered.
	S_shellLib_SHELL_TASK_EXISTS A shell session with the same name or the same standard input already exists.
	S_shellLib_SHELL_TASK_MAX The maximun number of shell session has been reached.
	S_shellLib_INTERNAL_ERROR An internal error prevents the shell session to be created.
	taskSpawn() errnos, malloc() and memPartCreate() errno.
SEE ALSO	shellLib, shellLock(), shellTerminate(), VxWorks Kernel Programmer's Guide: Kernel Shell

# shellHistory()

NAME	<b>shellHistory()</b> – display or set the size of the shell history (vxWorks 5.5 compatibility)
SYNOPSIS	<pre>STATUS shellHistory (     int size /* 0 = display, &gt;0 = set history to new size */ )</pre>
DESCRIPTION	This routine displays shell history, or resets the default number of commands displayed by shell history to <i>size</i> . By default, history size is 20 commands. Shell history is actually maintained by <b>ledLib</b> . If <i>size</i> is 0, the routine displays the line history.
IMPORTANT NOTE	This routine is backward compatible with previous version of the kernel shell. It only changes history size of the current shell session. If the current task is not a shell task, the routine will use the shell task attached to the console (if one exists).

VxWorks Kernel API Reference, 6.6 shellldVerify()

**RETURNS** OK, or ERROR if there is not a shell task attached to the console.

ERRNO N/A

**SEE ALSO** shellLib, ledLib, ledControl(), h(), VxWorks Kernel Programmer's Guide: Kernel Shell, Wind River Workbench Command-Line User's Guide 2.2: Host Shell

## shellIdVerify()

NAME	<b>shellIdVerify()</b> – verify the validity of a shell session Id
SYNOPSIS	STATUS shellIdVerify ( SHELL_ID shellId /* shell session Id to verify */ )
DESCRIPTION	This routine checks the validity of the shell session identifier <i>shellId</i> .
RETURNS	OK if the shell identifier is valid, ERROR otherwise.
ERRNO	N/A
SEE ALSO	shellLib, shellFirst(), shellFromNameGet(), VxWorks Kernel Programmer's Guide: Kernel Shell

## shellInterpByNameFind()

NAME	<pre>shellInterpByNameFind() - Find an interpreter based on its name</pre>
SYNOPSIS	<pre>SHELL_INTERP * shellInterpByNameFind   (     const char * interpName /* shell interpreter name to find */   )</pre>
DESCRIPTION	This routine checks if an interpreter with name <i>interpName</i> does exist.
RETURNS	a pointer on the interpreter structure if an interpreter is found, or <b>NULL</b> if the interpreter with specified name is not registered.
ERRNO	N/A

#### **SEE ALSO shellInterpDefaultNameGet()**, **shellInterpRegister()**, **shellLib**, **shellInterpLib**, *VxWorks Kernel Programmer's Guide:* **Kernel Shell**

# shellInterpCtxGet()

NAME	<pre>shellInterpCtxGet() - get the interpreter context</pre>
SYNOPSIS	SHELL_INTERP_CTX * shellInterpCtxGet ( SHELL_ID shellId /* shell session identifier */ )
DESCRIPTION	This routine returns a pointer to the context of the current interpreter used by the shell session <i>shellld</i>
RETURNS	a pointer to the interpreter context, or NULL if none interpreter is active.
ERRNO	N/A
SEE ALSO	shellInterpLib, shellInterpNameGet( ), shellInterpRegister( ), shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

# shellInterpDefaultNameGet( )

NAME	shellInterpDefaultNameGet() – get the name of the default interpreter
SYNOPSIS	const char * shellInterpDefaultNameGet (void)
DESCRIPTION	This routine returns the name of the default interpreter.
RETURNS	a pointer on the interpreter name, or NULL if none interpreter is registered.
ERRNO	N/A
SEE ALSO	shellInterpLib, shellInterpNameGet(), shellInterpRegister(), shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

### shellInterpEvaluate()

```
shellInterpEvaluate() – interpret a string by an interpreter
NAME
SYNOPSIS
                  STATUS shellInterpEvaluate
                      (
                      char * arg, /* argument to evaluate *,
const char * interpreterName, /* or NULL for default */
                                                                /* argument to evaluate */
                      SHELL_EVAL_VALUE * pValue
                                                                /* interpreter return value */
                      )
DESCRIPTION
                  This routine interprets the string arg by the interpreter named interpreterName. The result
                  value is returned in pValue. If interpreterName is NULL, the function uses the default
                  interpreter, the one which was registered first. Note that the string arg may be modified by
                 the interpreter, you need to save it before calling this routine if you want to use it later.
RETURNS
                  OK, or ERROR if an error occured
                 S shellLib_NO_INTERP
ERRNO
                      The interpreter specified is not registered or it does not have an evaluation routine.
                  S_shellLib_NOT_SHELL_TASK
                      The current task is not a shell task.
                  S_shellLib_INTERNAL_ERROR
                      An internal error occured and prevents the evaluation.
                  S_shellLib_INTERP_INIT_ERROR
                      The context of the interpreter cannot be initialized.
                  Interpreter's evaluation routine errnos.
SEE ALSO
                  shellInterpLib, VxWorks Kernel Programmer's Guide: Kernel Shell
```

### shellInterpNameGet()

NAME	<pre>shellInterpNameGet() - get the name of the current interpreter</pre>
SYNOPSIS	<pre>const char * shellInterpNameGet   (    SHELL_ID shellId /* shell session ID */ )</pre>
DESCRIPTION	This routine returns the name of the current interpreter of the shell session <i>shellId</i> .

ERRNO N/A

**SEE ALSO** shellInterpLib, shellInterpDefaultNameGet(), shellInterpRegister(), shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

# shellInterpRegister()

NAME	shellInterpRegister() – register a new interpreter
SYNOPSIS	STATUS shellInterpRegister ( FUNCPTR interpInitRtn /* interpreter init routine */ )
DESCRIPTION	This routine is used to register a new interpreter against the kernel shell.
	<i>interpInitRtn</i> is the initialization routine of the interpreter to register. Its definition is:
	STATUS init ( SHELL_INTERP * pInterp /* interpreter structure */ )
	with <i>pInterp</i> a pointer on an interpreter structure.
	The <i>interpInitRtn</i> routine is called at registration time. This routine must complete the fields of the <b>SHELL_INTERP</b> structure with the addresses of the interpreter routines:
	- the interpreter context initialization function,
	- the parser function,
	- the evaluation function,
	- the completion function,
	- the restart function (called whenever the shell is restarted, by CTRL-C key combination for example),
	- the interpreter context finalizer function (to release any resources).
	Any of these routine addresses (except the init function) can be NULL.
	The <i>interpInitRtn</i> routine also has to set the interpreter name and default prompt. The interpreter name must be unique among the registered interpreters.
	The definitions of the interpreter routines are:
	STATUS ctxInit

```
SHELL_INTERP_CTX * pInterpCtx /* interpreter context */
)
```

with *pInterpCtx* a pointer to the interpreter context, which is unique per shell session. This routine is called either when a new shell session is started, when a statement is evaluated or when the shell is switched to a new interpreter for which the interpreter context does not exist yet. This routine returns **ERROR** if an error occured, **OK** otherwise.

```
STATUS parser
(
SHELL_INTERP_CTX * pInterpCtx, /* interpreter context */
const char * inputLine, /* input line to parse */
BOOL isInteractive /* TRUE for interactive session */
)
```

with *plnterpCtx* a pointer to the interpreter context, *inputLine* a pointer to the input string to interpret. *islnteractive* is **TRUE** if the parsing is interactive. This routine returns **ERROR** if an error occured, **OK** otherwise.

```
STATUS evaluate
  (
   SHELL_INTERP_CTX * pInterpCtx, /* interpreter context */
   const char * inputLine, /* input line to parse */
   SHELL_EVAL_VALUE * pValue /* where to store return value */
  )
```

This routine is used to evaluate a string *inputLine* by the interpreter. The resulting value is stored into the shell value pointed by *pValue*. *pInterpCtx* is a pointer to the interpreter context. This routine returns **ERROR** if an error occured, **OK** otherwise.

```
STATUS completion
(
SHELL_INTERP_CTX * pInterpCtx, /* interpreter context */
LED_ID ledId, /* LED identifier */
char * line, /* line to complete */
UINT lineSize, /* size of line buffer */
UINT * pCursorPos, /* cursor position in the line */
char completionChar /* completion character */
)
```

This routine is called when a completion character is typed on the input. The completion characters are defined by the Line EDiting mode used. *ledId* identifies the Line EDiting session. *line* is the input line currently printed. Its maximal size is *lineSize*. The terminal EOS is not counted with this size. *pCursorPos* is a pointer to the position of the cursor in *line*. *completionChar* is the character typed for completion. *pInterpCtx* is a pointer to the interpreter context. This routine returns **ERROR** if the completion does not succeed, **OK** otherwise.

```
void ctxRestart
  (
  SHELL_INTERP_CTX * pInterpCtx /* interpreter context */
 )
```

with *pInterpCtx* a pointer to the interpreter context. This routine is called whenever the shell session is restarted. It is up to the interpreter to release any allocated resources.

	STATUS ctxFinalize
	( SHELL_INTERP_CTX * pInterpCtx /* interpreter context */ )
	with <i>pInterpCtx</i> a pointer to the interpreter context. This routine is used to free any internal resource used by the interpreter for a shell session. It is called when a shell session is terminated. This routine returns <b>ERROR</b> if an error occured, <b>OK</b> otherwise.
	The <i>ctxInit</i> routine does not need to set up the <b>currentPrompt</b> field of the <i>pInterpCtx</i> structure, but only the <b>prompt</b> field with the prompt string of the interpreter.
	The <i>pInterpParam</i> field of <i>pInterpCtx</i> can be used to store the address of an internal structure of the interpreter.
	The first interpreter registered becomes the default interpreter. The default interpreter is used by a shell session if no name is defined for its initialisation.
RETURNS	OK, or ERROR if an error occured.
ERRNO	S_shellLib_INTERP_EXISTS An interpreter with the same name is already registered
	S_shellLib_INTERNAL_ERROR An internal error occurs. The interpreter is not registered.
	memPartAlloc(), malloc() errnos
SEE ALSO	shellInterpLib, shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

# shellLock()

always returns TRUE.

NAME	<b>shellLock()</b> – lock access to the shell (vxWorks 5.5 compatibility)
SYNOPSIS	BOOL shellLock ( BOOL request /* TRUE = lock, FALSE = unlock */ )
	VxWorks 5.5 behavior
	This routine locks or unlocks access to the shell. When locked, cooperating tasks, such as <b>telnetdTask()</b> and <b>rlogindTask()</b> , will not be able to control the shell.
	If the shell is configured to use a unique task (compatibility mode), this routine reacts as the VxWorks 5.5 version. But for normal mode of the shell (multiple shell sessions), the routine

VxWorks Kernel API Reference, 6.6 shellNext()

RETURNS	If the shell session is unique, <b>TRUE</b> if <i>request</i> is "lock" and the routine successfully locks the shell, otherwise <b>FALSE</b> ; <b>TRUE</b> if request is "unlock" and the routine successfully unlocks the shell, otherwise <b>FALSE</b> . In multiple shell session mode, <b>TRUE</b> is always returned.
ERRNO	N/A
SEE ALSO	shellLib, shellCompatibleSet(), VxWorks Kernel Programmer's Guide: Kernel Shell

# shellNext()

NAME	<pre>shellNext() - get the next shell session</pre>
SYNOPSIS	SHELL_ID shellNext ( SHELL_ID shellId /* shell ID whose successor is to be found */ )
DESCRIPTION	This routine returns the Id of the next shell session, compared to <i>shellId</i> .
RETURNS	the next shell Id or 0 if there is no more running shell.
ERRNO	N/A
SEE ALSO	shellLib, shellFirst(), shellFromTaskGet(), shellFromNameGet(), VxWorks Kernel Programmer's Guide: Kernel Shell

# shellPromptFmtDftSet( )

NAME	<pre>shellPromptFmtDftSet( ) - set the default prompt format string</pre>
SYNOPSIS	STATUS shellPromptFmtDftSet ( const char * interp, /* interpreter name or NULL for default */ const char * promptFmt /* prompt format string or NULL for initial */ )
DESCRIPTION	This routine sets the format of the default prompt of the interpreter named <i>interp</i> to <i>promptFmt</i> . If <i>interp</i> is NULL, the default interpreter is used (the first registered). If <i>promptFmt</i> is NULL, the initial interpreter prompt, the one defined by the interpreter init function, is restored.

NOTE	Setting the default prompt of an interpreter after its first use will not modify the current prompt for a shell session. The change will be visible for new shell sessions or if it is done prior the first use of the interpreter. In order to change the current prompt, the routine <b>shellPromptFmtSet()</b> should be used.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the memory for the prompt cannot be allocated or if the interpreter is not registered.
ERRNO	<pre>S_shellLib_NO_INTERP     none interpreter named interp exists. malloc() errnos, memPartCreate() errnos.</pre>
SEE ALSO	shellPromptLib, shellPromptFmtSet(), VxWorks Kernel Programmer's Guide: Kernel Shell

# shellPromptFmtSet()

NAME	<pre>shellPromptFmtSet() - set the current prompt format string</pre>
SYNOPSIS	<pre>STATUS shellPromptFmtSet   (    SHELL_ID shellId, /* CURRENT_SHELL_SESSION for current shell */    const char * interp, /* interpreter name or NULL for current */    const char * promptFmt /* prompt format string or NULL for default */   )</pre>
DESCRIPTION	This routine sets the format of the current prompt of the interpreter named <i>interp</i> , associated with the shell session <i>shellld</i> , to <i>promptFmt</i> . If <i>shellId</i> is <b>CURRENT_SHELL_SESSION</b> , the current shell session is used. If <i>interp</i> is <b>NULL</b> , the current interpreter of <i>shellId</i> is used. If <i>promptFmt</i> is <b>NULL</b> , the default interpreter prompt is restored.
NOTE	Setting the current prompt for a shell session does not affect the other running sessions or new shell sessions. To change the default prompt of an interpreter, the routine <b>shellPromptFmtDftSet()</b> should be used instead.
RETURNS	OK, or ERROR if the interpreter does not exist or a memory error occured.
ERRNO	S_shellLib_NOT_SHELL_TASK shellId is an invalid shell session.
	S_shellLib_NO_INTERP none interpreter named <i>interp</i> exists.
	malloc() errnos, memPartCreate() errnos.

VxWorks Kernel API Reference, 6.6 shellPromptFmtStrAdd()

SEE ALSO shellPromptLib, shellPromptFmtDftSet(), VxWorks Kernel Programmer's Guide: Kernel Shell

### shellPromptFmtStrAdd()

shellPromptFmtStrAdd() – add a new prompt format string NAME SYNOPSIS STATUS shellPromptFmtStrAdd ( fmt, /\* format character \*/ char VOIDFUNCPTR fmtRtn, /\* format routine \*/ BOOL force /\* TRUE to superseed a previous definition \*/ ) DESCRIPTION This routine adds a new format string which manage the format character *fmt*. The display function associated with the format is *fmtRtn*. If *force* is **TRUE**, the format routine *fmtRtn* will superseed a previous definition of the format routine. When the shell prompt is printed, each format routine associated with the format strings used in the prompt string are called with the prototype: void fmtRtn ( SHELL\_ID shellId /\* shell session Id \*/ ) with *shellId* the identifier of the shell session managed. The format routine can use regular output functions (printf(), printErr()...) as the IO of the current task is correctly redirected to the dedicated shell terminal. OK, or ERROR if the format cannot be added. RETURNS ERRNO shellPromptLib, VxWorks Kernel Programmer's Guide: Kernel Shell SEE ALSO

### shellPromptSet()

NAME shellPromptSet() – change the shell prompt (vxWorks 5.5 compatibility)

SYNOPSIS STATUS shellPromptSet

	( const char * newPrompt /* string to become new shell prompt */ )
DESCRIPTION	This routine changes the format of the current shell prompt of the C interpreter to <i>newPrompt</i> . If <i>newPrompt</i> is <b>NULL</b> , the default C interpreter prompt is restored instead.
IMPORTANT NOTE	This routine is backward compatible with previous version of the kernel shell. This routine changes the prompt of the C interpreter only. If the shell is configured as compatible with VxWorks 5.5 (SHELL_COMPATIBLE parameter is TRUE), the default C interpreter prompt is also modified, not only the current prompt. If the current task is not a shell task, the routine will use the shell task attached to the console (if one exists).
RETURNS	OK, or ERROR if there is not a shell task attached to the console.
ERRNO	N/A
SEE ALSO	shellLib, shellPromptFmtSet(), shellPromptFmtDftSet(), VxWorks Kernel Programmer's Guide: Kernel Shell, Wind River Workbench Command-Line User's Guide 2.2: Host Shell

# shellResourceReleaseHookAdd()

NAME	<pre>shellResourceReleaseHookAdd() - add a resource-releasing hook to the shell</pre>
SYNOPSIS	STATUS shellResourceReleaseHookAdd ( SHELL_RES_RELEASE_HOOK hook /* hook routine to add */ )
DESCRIPTION	The function registers a routine which will be run when the shell is terminated or aborted. The routine is meant to release mutexes, semaphores or related data structures used within the shell components. When an abnormal shell exit occurs, this is used for deadlock prevention.
	Each hook routine must be of the following form :
	<pre>void releaseHook   (   SHELL_ID shellId, /* shell session Id */   BOOL force /* whether to force-release */   )</pre>
	The hooks are meant to release global resources, otherwise the release would be

The hooks are meant to release global resources, otherwise the release would be implemented in context-related termination routines. However, it is possible that session-related operations be needed. For that purpose, the hook takes a session ID as a parameter.



### shellRestart()

NAME	shellRestart() – restart a shell session
SYNOPSIS	STATUS shellRestart ( SHELL_ID shellId /* the ID of the shell session to restart */ )
DESCRIPTION	This routine restarts the shell session <i>shellId</i> . <i>shellId</i> can be a valid session identifier or <b>CURRENT_SHELL_SESSION</b> .
RETURNS	OK, or ERROR it not possible to restart the shell session.
ERRNO	N/A
SEE ALSO	shellLib, shellTaskGet(), taskRestart(), VxWorks Kernel Programmer's Guide: Kernel Shell

### shellScriptAbort()

**NAME** shellScriptAbort() – signal the shell to stop processing a script (vxWorks 5.5 compatibility)

SYNOPSIS STATUS shellScriptAbort (void)

**DESCRIPTION** This routine signals the current shell session to abort processing a script file. It can be called from within a script if an error is detected.

IMPORTANT NOTE	This routine is kept for backward compatibility reason with previous version of the kernel shell. This routine aborts scripting of the current shell session. If the current task is not a shell task, the function will use the shell task attached to the console (if one exists).
RETURNS	OK, or ERROR if there is not a shell task attached to the console.
ERRNO	N/A
SEE ALSO	shellLib, shellScriptNoAbort(), VxWorks Kernel Programmer's Guide: Kernel Shell

# shellTaskGet()

NAME	<pre>shellTaskGet() - get the task Id of a shell session</pre>
SYNOPSIS	<pre>int shellTaskGet   (    SHELL_ID shellId /* shell session Id */ )</pre>
DESCRIPTION	This routine returns the task Id for the shell session <i>shellId</i> .
RETURNS	the shell task Id, or <b>ERROR</b> if the shell session is not valid.
ERRNO	N/A
SEE ALSO	shellLib, shellIdVerify(), VxWorks Kernel Programmer's Guide: Kernel Shell

# shellTaskIdDefault()

NAME	<b>shellTaskIdDefault( )</b> – set the default task for a given shell session		
SYNOPSIS	<pre>int shellTaskIdDefault   (    SHELL_ID shellId, /* shell session Id */    int taskId /* default task Id or 0 */ )</pre>		
DESCRIPTION	This routine set the default task associated with the shell session <i>shellId</i> to <i>taskId</i> .		

	If <i>shellId</i> is equal to <b>CURRENT_SHELL_SESSION</b> , the current shell session is used. If this function is called outside of a shell session, and <i>shellId</i> is <b>CURRENT_SHELL_SESSION</b> , the shell session attached to the console is used (if one exists).
	If <i>shellId</i> is equal to <b>ALL_SHELL_SESSIONS</b> , the default task <i>taskId</i> is set for all shell sessions.
	If <i>taskId</i> is 0, the default task of the shell session <i>shellId</i> is returned.
RETURNS	the last default task Id (may be 0 if not previously set)
ERRNO	N/A
SEE ALSO	shellLib, taskIdDefault(), VxWorks Kernel Programmer's Guide: Kernel Shell

# shellTerminate( )

NAME	shellTerminate() – terminate a shell task
SYNOPSIS	<pre>void shellTerminate   (    SHELL_ID shellId /* the shell ID to terminate */ )</pre>
DESCRIPTION	This routine kills a shell session, based on its Id <i>shellId</i> . Suicide is prohibited (a shell session cannot kill itself by this call; use <b>exit()</b> instead).
NOTE	If the shell is configured to use only one shell session for all the connections, (compatible mode), this routine only restores the I/O of the shell task (see <b>shellGenericInit()</b> ), and then restarts it.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	shellLib, VxWorks Kernel Programmer's Guide: Kernel Shell

### shlShow()

NAME

**shlShow()** – display information for shared libraries

**DESCRIPTION** This routine displays information for shared libraries (SHL). This routine takes two parameters, *shlId* and *level*. The first parameter is the shared library ID, which may be obtained by using **rtpShlShow()**. The second parameter is the level of detail to display the SHL information.

Depending on the level and the SHL ID specified, the information displayed differs. If the *level* is 0, then it displays the summary information for either the specified SHL or all SHLs in the system. If the *level* is 1, then **shlShow()** displays the detailed information for the specified SHL. If *level* is 2, **shlShow()** displays the detailed information for all SHLs in the system, regardless of the SHL ID you specify. Refer to the table for more information.

Level	SHL ID	Meaning
0	0	Display summary information for all SHLs.
0	SHL	Display summary information for specified SHL.
1	0	Invalid Display of the SHL, must specify SHL.
1	SHL	Display detailed information for specified SHL.
2	ANY	Display detailed information for all SHLs.

**shlShow()** only displays the SHL name up to a maximum of 20 characters long. If the name is more than 20 characters, the name will be truncated to 18 characters for displaying purposes. Prepended to the truncated name, a "<" will be display to indicate that the name is more than 20 characters long. To get a display of the full SHL name, display the SHL with the *level* set to 1.

For the command interpreter shell, use the **shl** and the **shl info** commands to display SHL information.

Use this routine to display information on SHLs in the system. For information on SHLs associated with an RTP, use the **rtpShlShow()** routine to display this information.

**EXAMPLE** Below is an example display of a shared library.

-> shlShow 0x11c0724

	SHL NAME	ID	TEXT ADDR	TEXT SIZE	DATA SIZE	REF CN
<	tty/slDfw/libSo.so	1	0xff435000	0x574	0x628	2

**RETURNS TRUE**, or **FALSE** if the SL is invalid.

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VxWorks Kernel API Reference, 6.6 shlSymsAdd()

ERRNO	Possible errnos generated by this function include:		
	S_objLib_OBJ_ID_ERROR An incorrect SHL ID was provided.		
	S_objLib_ACCESS_DENIED Unable to get exclusive access to the SHL list.		
SEE ALSO	shlShow, rtpLib, rtpShow, shlLib, the VxWorks programmer guides.		

# shlSymsAdd()

NAME	shlSymsAdd() – add symbols from a shared object file to a RTP symbol table		
SYNOPSIS	<pre>STATUS shlSymsAdd   (     void * shlId, /* ID of the shared library */     RTP_ID rtpId, /* RTP the symbols should be added to */     UINT32 regPolicy, /* symbol registration policy */     char * filePath /* path and name of the shared object file */    )</pre>		
DESCRIPTION	This command is provided as a help in case a RTP needs to be debugged but has been launched with an empty symbol table. It forces the registration of the symbols from a shared object file into a RTP symbol table.		
	Note that this command does not verify whether the symbols are already in the symbol table and does not prevent the creation of multiple occurences of these symbols.		
	It is important to understand that symbols are added to the symbol table in the order of their registration and that the most recent entry will hide symbols of same name already registered. The <b>rtpLkup()</b> command will show all occurences of the symbols of a given name so it is possible to use their addresses instead of their names if there is a risk of confusion.		
	The only required information are the shared library ID ( <i>shlld</i> parameter) and the RTP ID ( <i>rtpId</i> parameter).		
	The <i>regPolicy</i> parameter sets the symbol registration policy. The policy can be one of the following:		
	0x01 ( <b>RTP_GLOBAL_SYMBOLS</b> ) Add only global symbols to the symbol table. This is the default when the parameter is left null.		
	0x02 ( <b>RTP_LOCAL_SYMBOLS</b> ) Add only local symbols to the symbol table.		

	0x03 ( <b>RTP_ALL_SYMBOLS</b> ) Add both local and global symbols to the symbol table.
	The <i>filePath</i> parameter overrides the path recorded for the shared library. It may be left null if the symbols should be read from the same file as the one used to create the shared library with. This parameter must be used when the symbols should be read from a file stored in a different location than what was recorded when the shared library has been created.
RETURNS	OK if the symbols could be read and recorded, ERROR otherwise.
ERRNO	N/A
SEE ALSO	usrRtpLib, shlSymsRemove( ), rtpSymsAdd( ), rtpSymsRemove( )

# shlSymsRemove()

NAME	shlSymsRemove() – remove shared library symbols from a RTP symbol table		
SYNOPSIS	<pre>STATUS shlSymsRemove   (    void * shlId, /* ID of the shared library */    RTP_ID rtpId, /* RTP the symbols should be removed from */    UINT32 remPolicy /* symbol removal policy */   )</pre>		
DESCRIPTION	This command forces the removal of symbols related to a shared library from a RTP symbol table.		
	The <i>remPolicy</i> parameter sets the symbol removal policy. The policy can be one of the following:		
	0x02 ( <b>RTP_LOCAL_SYMBOLS</b> ) Remove only the shared library's local symbols from the symbol table.		
	0x03 ( <b>RTP_ALL_SYMBOLS</b> ) Removes both the shared library's local and global symbols from the symbol table.		
RETURNS	OK if the symbols could be removed, ERROR otherwise.		
ERRNO	N/A		
SEE ALSO	usrRtpLib, shlSymsAdd( ), rtpSymsAdd( ), rtpSymsRemove( )		

# show()

NAME	<b>show()</b> – print information on a specified object		
SYNOPSIS	<pre>void show   (    int objId, /* object ID */    int level /* information level */   )</pre>		
DESCRIPTION	This command prints information on the specified object. System objects include tasks, local and shared semaphores, local and shared message queues, local and shared memory partitions, watchdogs, and symbol tables. An information level is interpreted by the objects show routine on a class by class basis. Refer to the object's library manual page for more information.		
RETURNS	N/A		
ERRNO	N/A		
SEE ALSO	usrLib, i(), ti(), lkup(), the VxWorks programmer guides.		

# sigInit()

NAME	<pre>sigInit() - initialize the signal facilities</pre>
SYNOPSIS	int sigInit ( BOOL posixMode )
DESCRIPTION	This routine initializes the signal facilities. It is usually called from the system start-up routine <b>usrInit()</b> in usrConfig, before interrupts are enabled.
	If the boolean parameter <i>posixMode</i> is <b>TRUE</b> then the signals sent to a faulting task will be POSIX conformant, if it is <b>FALSE</b> the signals will be backwards compatible with previous versions of VxWorks.
RETURNS	OK, or ERROR if the delete hooks cannot be installed.
ERRNO	<b>S_taskLib_TASK_HOOK_TABLE_FULL</b> Task hook table is full and signal delete hook can not be added.

SEE ALSO sigLib

## sigaction()

**sigaction()** – examine and/or specify the action associated with a signal (POSIX) NAME SYNOPSIS int sigaction ( signo, /\* signal of handler of interest \*/ int const struct sigaction \*pAct, /\* location of new handler \*/ struct sigaction \*pOact /\* location to store old handler \*/ DESCRIPTION This routine allows the calling process to examine and/or specify the action to be associated with a specific signal. OK (0), or ERROR (-1) if the signal number is invalid. RETURNS EINVAL ERRNO sigLib SEE ALSO

### sigaddset()

NAME	<b>sigaddset( )</b> – add a signal to a signal set (POSIX)
SYNOPSIS	<pre>int sigaddset   (     sigset_t *pSet, /* signal set to add signal to */     int signo /* signal to add */ )</pre>
DESCRIPTION	This routine adds the signal specified by <i>signo</i> to the signal set specified by $pSet$ .
RETURNS	<b>OK</b> (0), or <b>ERROR</b> (-1) if the signal number is invalid.
ERRNO	EINVAL
SEE ALSO	sigLib

# sigblock()

NAME	<b>sigblock( )</b> – add to a set of blocked signals
SYNOPSIS	<pre>int sigblock   (    int mask /* mask of additional signals to be blocked */ )</pre>
DESCRIPTION	This routine adds the signals in <i>mask</i> to the task's set of blocked signals. A one (1) in the bit mask indicates that the specified signal is blocked from delivery. Use the macro SIGMASK to construct the mask for a specified signal number. This routine has been deprecated, instead use <b>sigprocmask()</b> .
RETURNS	The previous value of the signal mask.
ERRNO	N/A
SEE ALSO	sigLib, sigprocmask()

# sigdelset()

NAME	<b>sigdelset( )</b> – delete a signal from a signal set (POSIX)
SYNOPSIS	<pre>int sigdelset   (     sigset_t *pSet, /* signal set to delete signal from */     int signo /* signal to delete */   )</pre>
DESCRIPTION	This routine deletes the signal specified by <i>signo</i> from the signal set specified by <i>pSet</i> .
RETURNS	<b>OK</b> (0), or <b>ERROR</b> (-1) if the signal number is invalid.
ERRNO	EINVAL
SEE ALSO	sigLib
## sigemptyset()

NAME	<b>sigemptyset()</b> – initialize a signal set with no signals included (POSIX)
SYNOPSIS	<pre>int sigemptyset    (     sigset_t *pSet /* signal set to initialize */ )</pre>
DESCRIPTION	This routine initializes the signal set specified by $pSet$ , such that all signals are excluded.
RETURNS	<b>OK</b> (0), or <b>ERROR</b> (-1) if the signal set cannot be initialized.
ERRNO	N/A
SEE ALSO	sigLib

## sigfillset()

NAME	<b>sigfillset( )</b> – initialize a signal set with all signals included (POSIX)
SYNOPSIS	<pre>int sigfillset   (    sigset_t *pSet /* signal set to initialize */ )</pre>
DESCRIPTION	This routine initializes the signal set specified by $pSet$ , such that all signals are included.
RETURNS	<b>OK</b> (0), or <b>ERROR</b> (-1) if the signal set cannot be initialized.
ERRNO	N/A
SEE ALSO	sigLib

## sigismember()

**NAME** sigismember() – test to see if a signal is in a signal set (POSIX)

SYNOPSIS int sigismember

VxWorks Kernel API Reference, 6.6 signal()

```
(
const sigset_t *pSet, /* signal set to test */
int signo /* signal to test for */
)
```

**DESCRIPTION** This routine tests whether the signal specified by *signo* is a member of the set specified by *pSet*.

**RETURNS** 1 if the specified signal is a member of the specified set, **OK** (0) if it is not, or **ERROR** (-1) if the test fails.

ERRNO EINVAL

SEE ALSO sigLib

### signal()

NAME	<b>signal( )</b> – specify the handler associated with a signal
SYNOPSIS	<pre>void (*signal     (     int signo,     void (*pHandler) ()     )) ()</pre>
DESCRIPTION	This routine chooses one of three ways in which receipt of the signal number <i>signo</i> is to be subsequently handled. If the value of <i>pHandler</i> is <b>SIG_DFL</b> , default handling for that signal will occur. If the value of <i>pHandler</i> is <b>SIG_IGN</b> , the signal will be ignored. Otherwise, <i>pHandler</i> must point to a function to be called when that signal occurs.
	A signal handler associated with <i>signo</i> as a result of a call to this routine will be reset to <b>SIG_DFL</b> upon entry into the signal handler. Subsequent instances of <i>signo</i> will thus be handled with the default action. The <b>sigaction()</b> routine must be used if this behavior is not desired.
RETURNS	The value of the previous signal handler, or <b>SIG_ERR</b> .
ERRNO	EINVAL
SEE ALSO	sigLib, sigaction()

#### sigpending() **sigpending()** – retrieve the set of pending signals blocked from delivery (POSIX) NAME SYNOPSIS int sigpending ( sigset\_t \*pSet /\* location to store pending signal set \*/ ) DESCRIPTION This routine stores the set of signals that are blocked from delivery and that are pending for the calling process in the space pointed to by *pSet*. **OK** (0), or **ERROR** (-1) if the signal TCB cannot be allocated. RETURNS ERRNO **ENOMEM** sigLib SEE ALSO

sigprocmask()

NAME	<b>sigprocmask()</b> – examine and/or change the signal mask (POSIX)
SYNOPSIS	<pre>int sigprocmask   (    int   how,  /* how signal mask will be changed */    const sigset_t *pSet,  /* location of new signal mask */    sigset_t *pOset  /* location to store old signal mask */   )</pre>
DESCRIPTION	This routine allows the calling process to examine and/or change its signal mask. If the value of <i>pSet</i> is not <b>NULL</b> , it points to a set of signals to be used to change the currently blocked set. The value of <i>how</i> indicates the manner in which the set is changed and consists of one of the following, defined in <b>signal.h</b> :
	<ul> <li>SIG_BLOCK the resulting set is the union of the current set and the signal set pointed to by <i>pSet</i>.</li> <li>SIG_UNBLOCK the resulting set is the intersection of the current set and the complement of the signal set pointed to by <i>pSet</i>.</li> <li>SIG_SETMASK the resulting set is the signal set pointed to by <i>pSet</i>.</li> </ul>

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VxWorks Kernel API Reference, 6.6 sigqueue()

**RETURNS** OK (0), or ERROR (-1) if *how* is invalid.

ERRNO EINVAL

SEE ALSO sigLib, sigsetmask(), sigblock()

### sigqueue()

NAME	<b>sigqueue( )</b> – send a queued signal to a task
SYNOPSIS	<pre>int sigqueue   (    int tid,    int signo,    const union sigval value   )</pre>
DESCRIPTION	The function <b>sigqueue()</b> sends the signal specified by <i>signo</i> with the signal-parameter value specified by <i>value</i> to the process specified by <i>tid</i> .
RETURNS	<b>OK</b> (0), or <b>ERROR</b> (-1) if the task ID or signal number is invalid, or if there are no queued-signal buffers available.
ERRNO	EINVAL EAGAIN
SEE ALSO	sigLib, taskSigqueue()

## sigqueueInit()

NAME	<pre>sigqueueInit( ) - initialize the queued signal facilities</pre>
SYNOPSIS	<pre>int sigqueueInit    (    int nQueues )</pre>
DESCRIPTION	This routine initializes the queued signal facilities. It must be called before any call to <b>sigqueue()</b> . It is usually called from the system start-up routine <b>usrInit()</b> in usrConfig, after <b>sysInit()</b> is called.

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It allocates *nQueues* buffers to be used by **sigqueue()**. A buffer is used by each call to **sigqueue()** and freed when the signal is delivered (thus if a signal is block, the buffer is unavailable until the signal is unblocked.)

**RETURNS** OK, or ERROR if memory could not be allocated.

ERRNO	N/A
ERRNO	N/A

SEE ALSO sigLib

#### sigsetmask()

NAME	sigsetmask() – set the signal mask
SYNOPSIS	<pre>int sigsetmask   (    int mask /* new signal mask */   )</pre>
DESCRIPTION	This routine sets the calling task's signal mask to a specified value. A one (1) in the bit mask indicates that the specified signal is blocked from delivery. Use the macro SIGMASK to construct the mask for a specified signal number. This routine has been deprecated, instead use <b>sigprocmask()</b> .
RETURNS	The previous value of the signal mask.
ERRNO	N/A
SEE ALSO	sigLib, sigprocmask()

### sigsuspend()

NAME	<b>sigsuspend()</b> – suspend the task until delivery of a signal (POSIX)	
SYNOPSIS	<pre>int sigsuspend   (     const sigset_t *pSet /* signal mask while suspended */ )</pre>	

VxWorks Kernel API Reference, 6.6 sigtimedwait()

DESCRIPTION	This routine suspends the task until delivery of a signal. While suspended, $pSet$ is used as the set of masked signals.
NOTE	Since the <b>sigsuspend()</b> function suspends thread execution indefinitely, there is no successful completion return value.
RETURNS	-1, always.
ERRNO	EINTR
SEE ALSO	sigLib

## sigtimedwait()

NAME	<b>sigtimedwait( )</b> – wait for a signal
SYNOPSIS	<pre>int sigtimedwait   (     const sigset_t *pSet, /* the signal mask while suspended */     siginfo_t *pInfo, /* return value */     const struct timespec *pTimeout   )</pre>
DESCRIPTION	The function <b>sigtimedwait()</b> selects the pending signal from the set specified by <i>pSet</i> . If multiple signals in <i>pSet</i> are pending, it will remove and return the lowest numbered one. If no signal in <i>pSet</i> is pending at the time of the call, the task will be suspend until one of the signals in <i>pSet</i> become pending, it is interrupted by an unblocked caught signal, or until the time interval specified by <i>pTimeout</i> has expired. If <i>pTimeout</i> is <b>NULL</b> , then the timeout interval is forever.
	If the <i>plnfo</i> argument is non- <b>NULL</b> , the selected signal number is stored in the <b>si_signo</b> member, and the cause of the signal is stored in the <b>si_code</b> member. If the signal is a queued signal, the value is stored in the <b>si_value</b> member of <i>plnfo</i> ; otherwise the content of <b>si_value</b> is undefined.
	The following values are defined in <b>signal.h</b> for <b>si_code</b> :
	SI_USER the signal was sent by the kill() function.
	SI_QUEUE the signal was sent by the <b>sigqueue( )</b> function.
	<b>SI_TIMER</b> the signal was generated by the expiration of a timer set by <b>timer_settime()</b> .

	SI_ASYNCIO the signal was generated by the completion of an asynchronous I/O request.
	SI_MESGQ the signal was generated by the arrival of a message on an empty message queue.
	The function <b>sigtimedwait()</b> provides a synchronous mechanism for tasks to wait for asynchromously generated signals. A task should use <b>sigprocmask()</b> to block any signals it wants to handle synchronously and leave their signal handlers in the default state. The task can then make repeated calls to <b>sigtimedwait()</b> to remove any signals that are sent to it.
RETURNS	Upon successful completion (that is, one of the signals specified by <i>pSet</i> is pending or is generated) <b>sigtimedwait()</b> will return the selected signal number. Otherwise, a value of -1 is returned and <b>errno</b> is set to indicate the error.
ERRNO	EINTR The wait was interrupted by an unblocked, caught signal.
	<b>EAGAIN</b> No signal specified by <i>pSet</i> was delivered within the specified timeout period.
	<b>EINVAL</b> The <i>pTimeout</i> argument specified a <b>tv_nsec</b> value less than zero or greater than or equal to 1000 million.
SEE ALSO	sigLib, sigwait( )

## sigvec()

NAME	sigvec() – install a signal handler
SYNOPSIS	<pre>int sigvec   (    int sig, /* signal to attach handler to */    const struct sigvec *pVec, /* new handler information */    struct sigvec *pOvec /* previous handler information */ )</pre>
DESCRIPTION	This routine binds a signal handler routine referenced by $pVec$ to a specified signal <i>sig</i> . It can also be used to determine which handler, if any, has been bound to a particular signal: <b>sigvec()</b> copies current signal handler information for <i>sig</i> to <i>pOvec</i> and does not install a signal handler if <i>pVec</i> is set to <b>NULL</b> (0).
	Both <i>pVec</i> and <i>pOvec</i> are pointers to a structure of type <b>struct sigvec</b> . The information passed includes not only the signal handler routine, but also the signal mask and additional option bits. The structure <b>sigvec</b> and the available options are defined in <b>signal.h</b> .

VxWorks Kernel API Reference, 6.6 sigwait()

**RETURNS** OK (0), or ERROR (-1) if the signal number is invalid or the signal TCB cannot be allocated.

ERRNO EINVAL

ENOMEM

SEE ALSO sigLib

### sigwait()

NAME	sigwait() – wait for a signal to be delivered (POSIX)
SYNOPSIS	<pre>int sigwait   (    const sigset_t *pSet,    int *pSig   )</pre>
DESCRIPTION	This routine waits until one of the signals specified in <i>pSet</i> is delivered to the calling thread. It then stores the number of the signal received in the the location pointed to by <i>pSig</i> .
	The signals in <i>pSet</i> must not be ignored on entrance to <b>sigwait()</b> . If the delivered signal has a signal handler function attached, that function is not called.
RETURNS	OK, or ERROR on failure.
ERRNO	N/A
SEE ALSO	sigLib, sigtimedwait()

### sigwaitinfo()

NAME	<b>sigwaitinfo( )</b> – wait for real-time signals
SYNOPSIS	<pre>int sigwaitinfo   (     const sigset_t *pSet, /* the signal mask while suspended */     siginfo_t *pInfo /* return value */   )</pre>
DESCRIPTION	The function <b>sigwaitinfo()</b> is equivalent to calling <b>sigtimedwait()</b> with <i>pTimeout</i> equal to <b>NULL</b> . See that reference entry for more information.

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RETURNS	Upon successful completion (that is, one of the signals specified by <i>pSet</i> is pending or is generated) <b>sigwaitinfo()</b> returns the selected signal number. Otherwise, a value of -1 is returned and <b>errno</b> is set to indicate the error.
ERRNO	EINTR The wait was interrupted by an unblocked, caught signal.
SEE ALSO	sigLib

### sil31xxBIST()

NAME	sil31xxBIST() – Controller Built-In Self Test
SYNOPSIS	STATUS sil31xxBIST (
	int ctrlNum
	)

- DESCRIPTION /NOMANUAL
- **RETURNS** Not Available
- ERRNO Not Available
- SEE ALSO vxbSI31xxStorage

### sil31xxBISTShow()

NAME sil31xxBISTShow() – Show the results of the power-on BIST

SYNOPSIS VOID sil31xxBISTShow

)

- DESCRIPTION none
- RETURNS Nothing
- ERRNO Not Available

VxWorks Kernel API Reference, 6.6 sil31xxDiskPresent()

SEE ALSO vxbSI31xxStorage

#### sil31xxDiskPresent()

NAME	sil31xxDiskPresent() – Return OK if disk exists.
SYNOPSIS	STATUS sil31xxDiskPresent ( int ctrlNum, int devNum )
DESCRIPTION	none
RETURNS	OK or ERROR
ERRNO	Not Available
SEE ALSO	vxbSI31xxStorage

### sil31xxDrvVxbInit()

NAME	<pre>sil31xxDrvVxbInit() - Initialize the driver.</pre>
SYNOPSIS	<pre>void sil31xxDrvVxbInit   (    BUS_DEVICE_ID pDev, /* vxbus DeviceID */    int ctrlNum /* assigned instance (controller number) */   )</pre>
DESCRIPTION	Initialize the driver structure for a single instance of the controller. This routine would get called once for each 31xx device.
RETURNS	N/A
ERRNO	
SEE ALSO	vxbSI31xxStorage

#### sil31xxIsr()

**NAME** sil31xxIsr() – Interrupt service routine.

SYNOPSIS VOID sil31xxIsr

int arg

**DESCRIPTION** none

RETURNS Not Available

ERRNO Not Available

SEE ALSO vxbSI31xxStorage

#### sil31xxRegisterPortCallback()

NAME sil31xxRegisterPortCallback() – register the port call back for a PHYRdyChg

SYNOPSIS STATUS sil31xxRegisterPortCallback ( int ctrlNum, int portNum, VOIDFUNCPTR myCallbackPtr, VOID \*myParam )

DESCRIPTION

event

RETURNS OK or ERROR

ERRNO Not Available

SEE ALSO vxbSI31xxStorage

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VxWorks Kernel API Reference, 6.6 sil31xxSectorRW()

sil31xxSectorRW()	
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sil31xxSectorRW() - read a single sector NAME SYNOPSIS STATUS sil31xxSectorRW ( ctrl, int int port, sector\_t sector, uint32\_t numSecs, char \*data, BOOL isRead ) DESCRIPTION This routine is called to read a single sector and dump the output on the console. ctrl controller number port port number sector starting sector for I/O operation numSecs number of sectors to read *data* pointer to data buffer RETURNS OK, or ERROR if the parameters are not valid. Not Available ERRNO SEE ALSO vxbSI31xxStorage

#### sil31xxXbdCreate()

NAME	<b>sil31xxXbdCreate()</b> – Create an XBD for the specified port.
SYNOPSIS	<pre>device_t sil31xxXbdCreate   (    int ctrlNum,    int devNum,    char *devName /* NULL for default, override with value */   )</pre>
DESCRIPTION	none
RETURNS	OK or ERROR
ERRNO	Not Available
SEE ALSO	vxbSI31xxStorage

#### sil31xxXbdDelete()

NAME sil31xxXbdDelete() – Delete an XBD for a specified port

SYNOPSIS STATUS sil31xxXbdDelete

)

int ctrlNum, int devNum

- **DESCRIPTION** none
- RETURNS OK or ERROR
- ERRNO Not Available
- SEE ALSO vxbSI31xxStorage

### sincos()

NAME	<b>sincos( )</b> – compute both a sine and cosine
SYNOPSIS	<pre>void sincos   (     double x, /* angle in radians */     double *sinResult, /* sine result buffer */     double *cosResult /* cosine result buffer */ )</pre>
DESCRIPTION	This routine computes both the sine and cosine of <i>x</i> in double precision. The sine is copied to <i>sinResult</i> and the cosine is copied to <i>cosResult</i> .
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	mathALib

## sincosf()

NAME	sincosf() – compute both a sine and cosine
SYNOPSIS	<pre>void sincosf   (    float x,</pre>
DESCRIPTION	This routine computes both the sine and cosine of $x$ in single precision. The sine is copied to $sinResult$ and the cosine is copied to $cosResult$ . The angle $x$ is expressed in radians.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	mathALib

## sinf()

NAME	<pre>sinf() - compute a sine (ANSI)</pre>
SYNOPSIS	float sinf ( float x /* angle in radians */ )
DESCRIPTION	This routine returns the sine of $x$ in single precision. The angle $x$ is expressed in radians.
RETURNS	The single-precision sine of <i>x</i> .
ERRNO	Not Available
SEE ALSO	mathALib

## sinhf()

NAME	<pre>sinhf() - compute a hyperbolic sine (ANSI)</pre>			
SYNOPSIS	<pre>float sinhf   (    float x /* number whose hyperbolic sine is required */ )</pre>			
DESCRIPTION	This routine returns the hyperbolic sine of $x$ in single precision.			
RETURNS	The single-precision hyperbolic sine of <i>x</i> .			
ERRNO	Not Available			
SEE ALSO	mathALib			

## sleep()

NAME	<b>sleep()</b> – delay for a specified amount of time		
SYNOPSIS	unsigned int sleep ( unsigned int secs )		
DESCRIPTION	This routine causes the calling task to be blocked for <i>secs</i> seconds.		
	The time the task is blocked for may be longer than requested due to the rounding up of the request to the timer's resolution or to other scheduling activities (e.g., a higher priority task intervenes).		
RETURNS	Zero if the requested time has elapsed, or the number of seconds remaining if it was interrupted.		
ERRNO	EINVAL EINTR		
SEE ALSO	timerLib, nanosleep( ), taskDelay( )		

#### smMemAddToPool()

NAME	<b>smMemAddToPool()</b> – add memory to shared memory system partition (VxMP Option)		
SYNOPSIS	<pre>STATUS smMemAddToPool   (     char * pPool, /* pointer to memory pool */     unsigned poolSize /* block size in bytes */   )</pre>		
DESCRIPTION	This routine adds memory to the shared memory system partition after the initial allocation of memory. The memory added need not be contiguous with memory previously assigned, but it must be in the same address space.		
	<i>pPool</i> is the global address of shared memory added to the partition. The memory area pointed to by <i>pPool</i> must be in the same address space as the shared memory anchor and shared memory pool.		
	<i>poolSize</i> is the size in bytes of shared memory added to the partition.		
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.		
RETURNS	OK, or ERROR if access to the shared memory system partition fails.		
ERRNO	S_smObjLib_LOCK_TIMEOUT		
SEE ALSO	smMemLib		

#### smMemCalloc( )

2

AVAILABILITY	This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.
RETURNS	A pointer to the block, or NULL if the memory cannot be allocated.
ERRNO	S_memLib_NOT_ENOUGH_MEMORY S_smObjLib_LOCK_TIMEOUT
SEE ALSO	smMemLib

#### smMemFindMax()

NAME	<b>smMemFindMax()</b> – find largest free block in shared memory system partition (VxMP Option)
SYNOPSIS	int smMemFindMax (void)
DESCRIPTION	This routine searches for the largest block in the shared memory system partition free list and returns its size.
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.
RETURNS	The size (in bytes) of the largest available block, or <b>ERROR</b> if the attempt to access the partition fails.
ERRNO	S_smObjLib_LOCK_TIMEOUT
SEE ALSO	smMemLib

### smMemFree()

smMemFree() - free a shared memory system partition block of memory (VxMP Option) NAME STATUS smMemFree SYNOPSIS ( void \* ptr /\* pointer to block of memory to be freed \*/ )

VxWorks Kernel API Reference, 6.6 smMemMalloc()

DESCRIPTION	This routine takes a block of memory previously allocated with <b>smMemMalloc()</b> or <b>smMemCalloc()</b> and returns it to the free shared memory system pool.
	It is an error to free a block of memory that was not previously allocated.
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the block is invalid.
ERRNO	S_memLib_BLOCK_ERROR S_smObjLib_LOCK_TIMEOUT
SEE ALSO	<pre>smMemLib, smMemMalloc( ), smMemCalloc( )</pre>

## smMemMalloc( )

NAME	<b>smMemMalloc()</b> – allocate block of memory from shared memory system partition (VxMP Option)	
SYNOPSIS	<pre>void * smMemMalloc   (     unsigned nBytes /* number of bytes to allocate */ )</pre>	
DESCRIPTION	This routine allocates a block of memory from the shared memory system partition whose size is equal to or greater than <i>nBytes</i> . The return value is the local address of the allocated shared memory block.	
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.	
RETURNS	A pointer to the block, or NULL if the memory cannot be allocated.	
ERRNO	S_memLib_NOT_ENOUGH_MEMORY S_smObjLib_LOCK_TIMEOUT	
SEE ALSO	smMemLib	

	<pre>smMemOptionsSet() smMemOptionsSet() - set debug options for shared memory system partition (VxMP Option)</pre>				
NAME					
SYNOPSIS	<pre>STATUS smMemOptionsSet   (    unsigned options /* options for system partition */ )</pre>				
DESCRIPTION	This routine sets the debug options for the shared system memory partition. Two kinds of errors are detected: attempts to allocate more memory than is available, and bad blocks found when memory is freed or reallocated. In both cases, the following options can be selected for actions to be taken when an error is detected: (1) return the error status, (2) log an error message and return the error status, or (3) log an error message and suspend the calling task. These options are discussed in detail in the library manual entry for <b>smMemLib</b> .				
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.				
RETURNS	OK or ERROR.				

- ERRNO S\_smObjLib\_LOCK\_TIMEOUT
- SEE ALSO smMemLib

### smMemRealloc()

NAME	<b>smMemRealloc()</b> – reallocate block of memory from shared memory system partition (VxMP Option)		
SYNOPSIS	<pre>void * smMemRealloc   (    void * pBlock, /* block to be reallocated */    unsigned newSize /* new block size */ )</pre>		
DESCRIPTION	This routine changes the size of a specified block and returns a pointer to the new block o shared memory. The contents that fit inside the new size (or old size, if smaller) remain unchanged. The return value is the local address of the reallocated shared memory block.		

VxWorks Kernel API Reference, 6.6 smMemShow()



#### smMemShow()

NAME	${\bf smMemShow}($ ) – show the shared memory system partition blocks and statistics (VxMP Option)				
SYNOPSIS	<pre>void smMemShow   (    int type /* 0 = statistics, 1 = statistics &amp; list */   )</pre>				
DESCRIPTION	This routine displays the total amount of free space in the shared memory system partition, including the number of blocks, the average block size, and the maximum block size. It also shows the number of blocks currently allocated, and the average allocated block size.				
	If <i>type</i> is 1, it displays a list of all the blocks in the free list of the shared memory system partition.				
WARNING	This routine locks access to the shared memory system partition while displaying the information. This can compromise the access time to the partition from other CPUs in the system. Generally, this routine is used for debugging purposes only.				
EXAMPLE	-> smMemShow 1				
	FREE LIST: num addr size				
	1 0x4ffef0 264 2 0x4fef18 1700				
	SUMMARY:				
	current free 1964 2 982 1700 alloc 2356 1 2356 -				

Not Available					
Not Available					
N/A					
This routine is distribut option, VxMP.	ted as a compone	ent of the	unbundled sl	hared memory object	s support
cumulative alloc value = 0 = 0x0	2620	2	1310	-	
	cumulative alloc value = 0 = 0x0 This routine is distribu option, VxMP. N/A	alloc 2620 value = 0 = 0x0 This routine is distributed as a compone option, VxMP.	alloc 2620 2 value = 0 = 0x0 This routine is distributed as a component of the option, VxMP.	alloc 2620 2 1310 value = 0 = 0x0 This routine is distributed as a component of the unbundled s option, VxMP.	alloc 2620 2 1310 - value = 0 = 0x0 This routine is distributed as a component of the unbundled shared memory objects option, VxMP.

**SEE ALSO** smMemShow, windsh, the VxWorks programmer guides, the, VxWorks Command-Line Tools User's Guide

#### smNameAdd()

**smNameAdd()** – add a name to the shared memory name database (VxMP Option) NAME SYNOPSIS STATUS smNameAdd ( char \* name, /\* name string to enter in database \*/ void \* value, /\* value associated with name \*/ /\* type associated with name \*/ int type ) This routine adds a name of specified object type and value to the shared memory objects DESCRIPTION name database. The name parameter is an arbitrary null-terminated string with a maximum of 20 characters, including EOS. By convention, type values of less than 0x1000 are reserved by VxWorks; all other values are user definable. The following types are predefined in smNameLib.h : Name Value Type T\_SM\_SEM\_B = 0shared binary semaphore T\_SM\_SEM\_C = 1 shared counting semaphore т

T_SM_MSG_Q	= 2	shared message queue
T_SM_PART_ID	= 3	shared memory Partition
T_SM_BLOCK	= 4	shared memory allocated block

A name can be entered only once in the database, but there can be more than one name associated with an object ID.

**AVAILABILITY** This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.

VxWorks Kernel API Reference, 6.6 smNameFind()

**RETURNS** OK, or ERROR if there is insufficient memory for *name* to be allocated, if *name* is already in the database, or if the database is already full.

ERRNO S\_smNameLib\_NOT\_INITIALIZED S\_smNameLib\_NAME\_TOO\_LONG S\_smNameLib\_NAME\_ALREADY\_EXIST S\_smNameLib\_DATABASE\_FULL S\_smObjLib\_LOCK\_TIMEOUT

SEE ALSO smNameLib, smNameShow

### smNameFind()

NAME	<b>smNameFind()</b> – look up a shared memory object by name (VxMP Option)					
SYNOPSIS	<pre>STATUS smNameFind   (     char * name, /* name to search for */     void ** pValue, /* pointer where to return value */     int * pType, /* pointer where to return object type */     int waitType /* NO_WAIT or WAIT_FOREVER */   )</pre>					
DESCRIPTION	This routine searches the shared memory objects name database for an object matching a specified <i>name</i> . If the object is found, its value and type are copied to the addresses pointed to by <i>pValue</i> and <i>pType</i> . The value of <i>waitType</i> can be one of the following:					
	<b>NO_WAIT (0)</b> The call returns immediately, even if <i>name</i> is not in the database.					
	<b>WAIT_FOREVER (-1)</b> The call returns only when <i>name</i> is available in the database. If <i>name</i> is not already in, the database is scanned periodically as the routine waits for <i>name</i> to be entered.					
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.					
RETURNS	<b>OK</b> , or <b>ERROR</b> if the object is not found, if <i>name</i> is too long, or the wait type is invalid.					
ERRNO	S_smNameLib_NOT_INITIALIZED S_smNameLib_NAME_TOO_LONG S_smNameLib_NAME_NOT_FOUND S_smNameLib_INVALID_WAIT_TYPE S_smObjLib_LOCK_TIMEOUT					

SEE ALSO smNameLib, smNameShow

## smNameFindByValue()

NAME	<b>smNameFindByValue( )</b> – look up a shared memory object by value (VxMP Option)				
SYNOPSIS	<pre>STATUS smNameFindByValue  (     void * value, /* value to search for */     char * name, /* pointer where to return name */     int * pType, /* pointer where to return object type */     int waitType /* NO_WAIT or WAIT_FOREVER */ )</pre>				
DESCRIPTION	This routine searches the shared memory name database for an object matching a specified value. If the object is found, its name and type are copied to the addresses pointed to by <i>name</i> and <i>pType</i> . The value of <i>waitType</i> can be one of the following:				
	<b>NO_WAIT (0)</b> The call returns immediately, even if the object value is not in the database.				
	WAIT_FOREVER (-1) The call returns only when the object value is available in the database.				
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.				
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>value</i> is not found or if the wait type is invalid.				
ERRNO	S_smNameLib_NOT_INITIALIZED S_smNameLib_VALUE_NOT_FOUND S_smNameLib_INVALID_WAIT_TYPE S_smObjLib_LOCK_TIMEOUT				
SEE ALSO	smNameLib, smNameShow				

#### smNameRemove()

**smNameRemove()** – remove an object from the shared memory objects name database NAME (VxMP Option) SYNOPSIS STATUS smNameRemove ( char \* name /\* name of object to remove \*/ This routine removes an object called *name* from the shared memory objects name database. DESCRIPTION AVAILABILITY This routine is distributed as a component of the unbundled shared memory objects support option, VxMP. RETURNS OK, or ERROR if the object name is not in the database or if *name* is too long. ERRNO S\_smNameLib\_NOT\_INITIALIZED S\_smNameLib\_NAME\_TOO\_LONG S\_smNameLib\_NAME\_NOT\_FOUND S\_smObjLib\_LOCK\_TIMEOUT smNameLib, smNameShow SEE ALSO

#### smNameShow()

NAME smNameShow() – show the contents of the shared memory objects name database (VxMP Option)

```
SYNOPSIS STATUS smNameShow
(
int level /* information level */
)
```

**DESCRIPTION** This routine displays the names, values, and types of objects stored in the shared memory objects name database. Predefined types are shown, using their ASCII representations; all other types are printed in hexadecimal.

The *level* parameter defines the level of database information displayed. If *level* is 0, only statistics on the database contents are displayed. If *level* is greater than 0, then both statistics and database contents are displayed.

**WARNING** This routine locks access to the shared memory objects name database while displaying its contents. This can compromise the access time to the name database from other CPUs in the system. Generally, this routine is used for debugging purposes only.

EXAMPLE -> smNameShow

Names in Database Max: 30 Current: 6 Free: 24 -> smNameShow 1 Names in Database Max: 30 Current: 6 Free: 24 Value Name Type \_\_\_\_\_ --- -----inputImage 0x802340 SM\_MEM\_BLOCK 0x806340 SM\_MEM\_BLOCK 0x802001 SM\_MEM\_PART ouputImage imagePool imagePool imageInSem 0x8e0001 SM\_SEM\_B imageOutSem 0x8e0101 SM\_SEM\_C actionQ 0x8e0201 SM\_MSG\_Q userObject 0x8e0400 0x1b0

- AVAILABILITY This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.
- **RETURNS** OK, or **ERROR** if the name facility is not initialized.
- ERRNO S\_smNameLib\_NOT\_INITIALIZED S\_smObjLib\_LOCK\_TIMEOUT
- SEE ALSO smNameShow, smNameLib

#### smObjAttach()

NAME	<b>smObjAttach()</b> – attach the calling CPU to the shared memory objects facility (VxMP Option)
SYNOPSIS	STATUS smObjAttach ( SM_OBJ_DESC * pSmObjDesc /* pointer to shared memory descriptor */ )
DESCRIPTION	This routine "attaches" the calling CPU to the shared memory objects facility. The shared memory area is identified by the shared memory descriptor with an address specified by

*pSmObjDesc.* The descriptor must already have been initialized by calling **smObjInit()**.

	smObjGlobalToLocal( )
	This routine is called automatically when the component INCLUDE_SM_OBJ is included.
	This routine will complete the attach process only if and when the shared memory has been initialized by the master CPU. If the shared memory is not recognized as active within the timeout period (10 minutes), this routine returns <b>ERROR</b> .
	The <b>smObjAttach()</b> routine connects the shared memory objects handler to the shared memory interrupt. Note that this interrupt may be shared between the shared memory network driver and the shared memory objects facility when both are used at the same time.
WARNING	Once a CPU has attached itself to the shared memory objects facility, it cannot be detached. Since the shared memory network driver and the shared memory objects facility use the same low-level attaching mechanism, a CPU cannot be detached from a shared memory network driver if the CPU also uses shared memory objects.
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the shared memory objects facility is not active or the number of CPUs exceeds the maximum.
ERRNO	S_smLib_INVALID_CPU_NUMBER
SEE ALSO	smObjLib, smObjSetup(), smObjInit()

# smObjGlobalToLocal()

VxWorks Kernel API Reference, 6.6

NAME	<b>smObjGlobalToLocal()</b> – convert a global address to a local address (VxMP Option)
SYNOPSIS	<pre>void * smObjGlobalToLocal    (    void * globalAdrs /* global address to convert */ )</pre>
DESCRIPTION	This routine converts a global shared memory address <i>globalAdrs</i> to its corresponding local value. This routine does not verify that <i>globalAdrs</i> is really a valid global shared memory address.
	All addresses stored in shared memory are global. Any access made to shared memory by the local CPU must be done using local addresses. This routine and <b>smObjLocalToGlobal()</b> are used to convert between these address types.
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.

RETURNS The local shared memory address pointed to by *globalAdrs*.

ERRNO Not Available

SEE ALSO smObjLib, smObjLocalToGlobal()

#### smObjInit()

**smObiInit()** – initialize a shared memory objects descriptor (VxMP Option) NAME

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SYNOPSIS void smObjInit

(			
SM_OBJ_DESC *	pSmObjDesc,	/*	ptr to shared memory descriptor */
SM_ANCHOR *	anchorLocalAdrs,	/*	shared memory anchor local adrs */
int	ticksPerBeat,	/*	cpu ticks per heartbeat */
int	smObjMaxTries,	/*	<pre>max no. of tries to obtain spinLock */</pre>
int	intType,	/*	interrupt method */
int	intArg1,	/*	interrupt argument #1 */
int	intArg2,	/*	interrupt argument #2 */
int	intArg3	/*	interrupt argument #3 */
)			

DESCRIPTION This routine initializes a shared memory descriptor. The descriptor must already be allocated in the CPU's local memory. Once the descriptor has been initialized by this routine, the CPU may attach itself to the shared memory area by calling smObjAttach().

This routine is called automatically when the component INCLUDE\_SM\_OBJ is included.

Only the shared memory descriptor itself is modified by this routine. No structures in shared memory are affected.

Parameters:

pSmObjDesc

The address of the shared memory descriptor to be initialized; this structure must be allocated before **smObjInit()** is called.

anchorLocalAdrs

The memory address by which the local CPU may access the shared memory anchor. This address may vary among CPUs in the system because of address offsets (particularly if the anchor is located in one CPU's dual-ported memory).

ticksPerBeat

Specifies the frequency of the shared memory anchor's heartbeat. The frequency is expressed in terms of how many CPU ticks on the local CPU correspond to one heartbeat period.

	<i>smObjMaxTries</i> Specifies the maximum number of tries to obtain access to an internal mutually exclusive data structure.
	<i>intType, intArg1, intArg2, intArg3</i> Allow a CPU to announce the method by which it is to be notified of shared memory events. See the manual entry for <b>if_sm</b> for a discussion about interrupt types and their associated parameters.
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	<pre>smObjLib, smObjSetup(), smObjAttach()</pre>

## smObjLibInit( )

NAME	<pre>smObjLibInit() – install the shared memory objects facility (VxMP Option)</pre>
SYNOPSIS	STATUS smObjLibInit (void)
DESCRIPTION	This routine installs the shared memory objects facility. It is called automatically when the component <b>INCLUDE_SM_OBJ</b> is included.
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.
RETURNS	OK, or ERROR if the shared memory objects facility has already been installed.
ERRNO	Not Available
SEE ALSO	smObjLib

## smObjLocalToGlobal()

NAME	<pre>smObjLocalToGlobal() - convert a local address to a global address (VxMP Option)</pre>
SYNOPSIS	<pre>void * smObjLocalToGlobal   (    void * localAdrs /* local address to convert */ )</pre>
DESCRIPTION	This routine converts a local shared memory address <i>localAdrs</i> to its corresponding global value. This routine does not verify that <i>localAdrs</i> is really a valid local shared memory address.
	All addresses stored in shared memory are global. Any access made to shared memory by the local CPU must be done using local addresses. This routine and <b>smObjGlobalToLocal()</b> are used to convert between these address types.
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.
RETURNS	The global shared memory address pointed to by <i>localAdrs</i> .
ERRNO	Not Available
SEE ALSO	smObjLib, smObjGlobalToLocal( )

## smObjSetup()

NAME	<b>smObjSetup()</b> – initialize the shared memory objects facility (VxMP Option)				
SYNOPSIS	STATUS smObjSetup ( SM_OBJ_PARAMS * smObjParams /* setup parameters */ )				
DESCRIPTION	This routine initializes the shared memory objects facility by filling the shared memory header. It must be called only once by the shared memory master CPU. It is called automatically only by the master CPU, when the component <b>INCLUDE_SM_OBJ</b> is included.				
	Any CPU on the system backplane can use the shared memory objects facility; however, the facility must first be initialized on the master CPU. Then before other CPUs are attached to the shared memory area by <b>smObjAttach()</b> , each must initialize its own shared memory				

objects descriptor using **smObjInit()**. This mechanism is similar to the one used by the shared memory network driver.

The *smObjParams* parameter is a pointer to a structure containing the values used to describe the shared memory objects setup. This structure is defined as follows in **smObjLib.h**:

typedef struct s	sm_obj_params	/*	setup parameters */	
{				
BOOL	allocatedPool;	/*	TRUE if shared memory pool is malloced	*/
SM_ANCHOR *	pAnchor;	/*	shared memory anchor	*/
char *	smObjFreeAdrs;	/*	start address of shared memory pool	*/
int	smObjMemSize;	/*	memory size reserved for shared memory	*/
int	maxCpus;	/*	max number of CPUs in the system	*/
int	maxTasks;	/*	max number of tasks using smObj	*/
int	maxSems;	/*	max number of shared semaphores	*/
int	<pre>maxMsgQueues;</pre>	/*	max number of shared message queues	*/
int	<pre>maxMemParts;</pre>	/*	max number of shared memory partitions	*/
int	maxNames;	/*	max number of names of shared objects	*/
} SM OBJ PAR	RAMS;			

## **AVAILABILITY** This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.

## **RETURNS** OK, or ERROR if the shared memory pool cannot hold all the requested objects or the number of CPUs exceeds the maximum.

ERRNO S\_smObjLib\_TOO\_MANY\_CPU S\_smObjLib\_SHARED\_MEM\_TOO\_SMALL

SEE ALSO smObjLib, smObjInit(), smObjAttach()

#### smObjShow()

NAME	<b>smObjShow()</b> – display the current status of shared memory objects (VxMP Option)
SYNOPSIS	STATUS smObjShow (void)
DESCRIPTION	This routine displays useful information about the current status of shared memory objects facilities.
WARNING	The information returned by this routine is not static and may be obsolete by the time it is examined. This information is generally used for debugging purposes only.
EXAMPLE	-> smObjShow Shared Mem Anchor Local Addr: 0x600. Shared Mem Hdr Local Addr: 0xb1514.

Attached CPU : Max Tries to Take Lo	5 ck: 1		
Shared Object Type	Current	Maximum	Available
 Tasks	1	20	19
Binary Semaphores	8	30	20
Counting Semaphores	2	30	20
Messages Queues	3	10	7
Memory Partitions	1	4	3
Names in Database	16	100	84

- **AVAILABILITY** This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.
- **RETURNS** OK, or **ERROR** if no shared memory objects are initialized.
- ERRNO S\_smObjLib\_NOT\_INITIALIZED
- SEE ALSO smObjShow, smObjLib

## smObjTimeoutLogEnable( )

NAME	<b>smObjTimeoutLogEnable()</b> – control logging of failed attempts to take a spin-lock (VxMP Option)
SYNOPSIS	void smObjTimeoutLogEnable ( BOOL timeoutLogEnable /* TRUE to enable, FALSE to disable */ )
DESCRIPTION	This routine enables or disables the printing of a message when an attempt to take a shared memory spin-lock fails.
	By default, message logging is enabled.
AVAILABILITY	This routine is distributed as a component of the unbundled shared memory objects support option, VxMP.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	smObjLib

VxWorks Kernel API Reference, 6.6 smeRegister()

## smeRegister()

SEE ALSO	vxbSmscLan9118End
ERRNO	N/A
RETURNS	N/A
DESCRIPTION	This routine registers the SMSC driver with VxBus as a child of the PLB bus type.
SYNOPSIS	void smeRegister(void)
NAME	<b>smeRegister()</b> – register with the VxBus subsystem

## smpLockDemo()

NAME	<pre>smpLockDemo() - smpLockDemo entry point (shell command)</pre>		
SYNOPSIS	<pre>STATUS smpLockDemo   (     unsigned int secs,</pre>		
DESCRIPTION	<pre>Invoke the smpLockDemo by calling this routine from the kernel shell: -&gt; smpLockDemo <number of="" secs="">, <number of="" tasks="">, &lt;[TRUE, FALSE]; (affinity) The secs argument is optional. It represents the number of seconds the demo spends updating the local and global counts for each synchronization mechanism mentioned in the module description. When invoked with no arguments the default is two seconds.</number></number></pre>		
RETURNS	OK if worker tasks were spawn without failure otherwise returns ERROR		
ERRNO	N/A		
SEE ALSO	smpLockDemo		

## snprintf()

NAME	<b>snprintf()</b> – write a formatted string to a buffer, not exceeding buffer size (ANSI)
SYNOPSIS	<pre>int snprintf  (     char * buffer, /* buffer to write to */     size_t count, /* max number of characters to store in buffer */     const char * fmt, /* format string */     /* optional arguments to format */ )</pre>
DESCRIPTION	This routine copies a formatted string to a specified buffer, up to a given number of characters. The formatted string will be null terminated. This routine guarantees never to write beyond the provided buffer regardless of the format specifier or the arguments to be formatted. The <i>count</i> argument specifies the maximum number of characters to store in the buffer, including the null terminator.
	Its function and syntax are otherwise identical to <b>printf()</b> .
RETURNS	The number of characters copied to <i>buffer</i> , not including the <b>NULL</b> terminator. Even when the supplied <i>buffer</i> is too small to hold the complete formatted string, the return value represents the number of characters that would have been written to <i>buffer</i> if <i>count</i> was sufficiently large.
ERRNO	Not Available
SEE ALSO	fioBaseLib, sprintf(), printf(), "International Organization for Standardization, ISO/IEC 9899:1999, ", "Programming languages - C: Input/output (stdio.h)"

## snsShow()

NAME	<b>snsShow()</b> – show information about services in the SNS directory
SYNOPSIS	void snsShow ( const char * servName /* service name prefix */ )
DESCRIPTION	This routine displays information about the services registered with SNS. <i>servName</i> is represented in URL format:
	[SNS:]service_name[@scope]

where the parts in brackets, [], are optional.

**SNS:** represent the URL service, i.e. the Socket Name Service. It is the only value accepted and can be omitted. *@scope* represents the visibility of the service name within the system. It can take several values, depending from the context and the application needs.

If the the scope is not specified, "@node" is assumed.

The URL representation is case insensitive.

All services whose name begins with the string specified by **service\_name** are listed. **service\_name** may contain wildcard characters \* or ? for name pattern matching, where \* denotes matching as many characters as possible, including zero number of character, ? denotes matching any single character.

If *servName* is **NULL**, or points to a null string, then all services in the SNS directory are listed.

The information displayed for each service listed includes:

- the service name
- the service scope
- socket address family associated with the service
- socket type associated with the service
- socket protocol number associated with the service: 0 represents the dummy value, since the socket address family and type are known.
- socket address associated with the service

Examples,

NAME	SCOPE	FAMILY	.TYPE	PROTO	ADDR
webAdmin	node	LOCAL	SEQPKT	0	/comp/socket/0x4
eventBlog	priv	LOCAL	SEQPKT	0	/comp/socket/0x8
clusterTimeServer	clust	TIPC	SEQPKT	0 *	1.1.5,1086717964

RETURNS N/A.

ERRNO Not Available

SEE ALSO snsShow, snsLib

## **so( )**

NAME	<b>so( )</b> – single-step, but step over a subroutine	
SYNOPSIS	<pre>STATUS so   (     int taskNameOrId /* task to step; 0 = default */   )</pre>	
DESCRIPTION	This routine single-steps a task that is stopped at a breakpoint. However, if the next instruction is a branch call to a subroutine, <b>so()</b> executes the subroutine and stops after.	
	To execute, enter:	
	-> so [task]	
	If <i>task</i> is omitted or zero, the last task referenced is assumed.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if the debugging package is not installed, the task cannot be found, or the task is not suspended.	
ERRNO	N/A	
SEE ALSO	<b>dbgLib</b> , <b>s()</b> , <b>cret()</b> , the VxWorks programmer guides, the , <i>VxWorks Command-Line Tools User's Guide</i> .	

## sp()

NAME	<b>sp( )</b> – spawn a task with default parameters
SYNOPSIS	<pre>int sp   (   FUNCPTR func, /* function to call */   int arg1, /* first of nine args to pass to spawned task */   int arg2,   int arg3,   int arg4,   int arg5,   int arg6,   int arg7,   int arg8,   int arg9   )</pre>

VxWorks Kernel API Reference, 6.6 spinLockIsrGive()

DESCRIPTION	This command spawns a specified function as a task with the following defaults. These default priorities may be overriden by updating the specified shell variable:
	priority (spTaskPriority): 100
	stack size (spTaskStackSize): 20,000 bytes
	task options (spTaskOptions): COPROCS_ALL (execute with all coprocessors support)
	task name: a name of the form <b>tN</b> where N is an integer which increments as new tasks are spawned, e.g., <b>t1</b> , <b>t2</b> , <b>t3</b> , etc.
	task ID: highest not currently used
	The task ID is displayed after the task is spawned.
	This command is a short form of the underlying <b>taskSpawn()</b> routine, convenient for spawning tasks in which the default parameters are satisfactory. If the default parameters are unacceptable, <b>taskSpawn()</b> should be called directly.
RETURNS	a task ID, or <b>ERROR</b> if the task cannot be spawned.
ERRNO	EINVAL the address <i>func</i> is NULL
	taskSpawn() errnos.
SEE ALSO	usrLib, taskLib, taskSpawn(), the VxWorks programmer guides.

## spinLockIsrGive()

NAME	<pre>spinLockIsrGive() – release an ISR-callable spinlock</pre>
SYNOPSIS	void spinLockIsrGive
	spinlockisr_t *plock /* pointer to isk-callable spinlock */ )
DESCRIPTION	This routine releases the ISR-callable spinlock pointed to by <i>pLock</i> . Furthermore, it re-enables interrupts that had been disabled on the local CPU when the lock was acquired using <b>spinLockIsrTake()</b> . Calling this routine under the following circumstances is
	considered to be an error condition and has undefined behaviour:
	- The calling task or ISR is not the one that acquired the spinlock.
----------	---
	- The pLock argument does not point to a properly initialized ISR-callable spinlock.
	This function forces a read/write memory barrier before releasing the lock.
	If <b>INCLUDE_SPINLOCK_DEBUG</b> is defined the following scenarios will cause a ED&R kernel fatal error which may reboot the target depending on ED&R policy in place:
	- If a CPU different than the owner of a spinlock attempts to release it
	- If a CPU attempts to release a spinlock that was never acquired
RETURNS	N/A
ERRNO	N/A
SEE ALSO	<pre>spinLockLib, spinLockIsrGive( ), spinLockIsrInit( )</pre>

# spinLockIsrHeld()

NAME	<b>spinLockIsrHeld()</b> – is an ISR-callable spinlock held by the current CPU?
SYNOPSIS	BOOL spinLockIsrHeld ( spinlockIsr_t *pLock /* pointer to ISR-callable spinlock */ )
DESCRIPTION	This routine returns <b>TRUE</b> if the ISR-callable spinlock pointed to by <i>pLock</i> is currently held by the calling CPU, or <b>FALSE</b> if it is not.
	Calling this routine with a <i>pLock</i> that points to anything that is not a properly initialized ISR-callable spinlock has undefined behaviour.
RETURNS	BOOL
ERRNO	N/A
SEE ALSO	spinLockLib, spinLockIsrGive( ), spinLockIsrTake( )

	spinLockIsrInit()	
NAME	<pre>spinLockIsrInit() - initialize an ISR-callable spinlock</pre>	
SYNOPSIS	VxWorks Architecture Supplements	
	<pre>void spinLockIsrInit   (    spinlockIsr_t *pLock, /* pointer to ISR-callable spinlock */    int flags /* spinlock attributes */   )</pre>	
DESCRIPTION	This routine initializes the ISR-callable spinlock pointed to by <i>pLock</i> , using the <i>flags</i> specified. Currently, no flags are defined; this argument is a placeholder for future enhancements. A spinlock must be initialized before it is used for the first time. A spinl is built on the ability of a processor to perform an atomic read-modify-write access to memory. Some CPUs may have cache attributes and memory alignment restrictions on use of these instructions. It is the responsibility of the caller to ensure the memory locat where the spinlock is located respects these restrictions, if any.	
	This routine must not be called from interrupt level.	
	If <b>INCLUDE_SPINLOCK_DEBUG</b> is defined, the following scenarios will cause a ED&R kernel fatal error which may reboot the target depending on the ED&R policy in place:	
	- If this routine is called from interrupt level	
RETURNS	N/A	
ERRNO	N/A	
SEE ALSO	spinLockLib, spinLockIsrTake( ), spinLockIsrGive( ), spinLockTaskInit( )	

## spinLockIsrTake()

```
NAME spinLockIsrTake() – take an ISR-callable spinlock
```

SYNOPSIS

```
void spinLockIsrTake
  (
#ifdef SPIN_LOCK_TRACE
  spinlockIsr_t *pLock, /* pointer to ISR-callable spinlock */
  char *file,
  int line
```

```
#else
   spinlockIsr_t *pLock /* pointer to ISR-callable spinlock */
#endif
    )
RETURNS: N/A
ERRNO: N/A
SEE ALSO: spinLockIsrGive()
<section>
<heading>DESCRIPTION
This routine acquires the ISR-callable spinlock pointed to by pLock.
If the lock is available at the time of the call, this routine
marks the spinlock as being in use and returns immediately.
                                                            If the
spinlock is unavailable, the routine busy-waits for the lock to become
available. Because of this busy-wait characteristic, recursive acquisition
of a spinlock causes a live lock situation where the acquiring task or
ISR busy-waits forever for a spinlock it already holds.
Acquisition of an ISR-callable spinlock causes interrupts to be masked
on the local CPU until the lock is released using spinLockIsrGive().
It is therefore recommended that this type of spinlock be held for
a minimal amount of time as it increases interrupt latency on
the local CPU.
Calling this routine with a pLock that points to anything that is not a
properly initialized ISR-callable spinlock has undefined behaviour.
This routine provides a memory barrier mechanism to prevent memory access
reordering that may be performed by the hardware.
If INCLUDE_SPINLOCK_DEBUG is defined the following scenario will inject a
kernel fatal error message in ED&R and may reboot the target depending on
the policy in place:
- If this routine is called within the context of a CPU already holding the
  same spinlock. A lock can not be taken recursively.
- If this routine is called while any another ISR-callable is already held.
  ISR-callable spinlocks can not be nested.
```

#### 

If INCLUDE\_SPINLOCK\_DEBUG is defined the following scenarios will cause a

VxWorks Kernel API Reference, 6.6 spinLockIsrTake()

ED&R kernel fatal error which may reboot the target depending on ED&R policy in place:

<returns> <heading>RETURNS Not Available

<errno> <heading>ERRNO Not Available

<seealso> <heading>SEE ALSO spinLockLib

<routinedoc> <rtnhead>spinLockTaskGive( ) <rtnname> <heading>NAME <rtnshort> spinLockTaskGive() - release a task-only spinlock

```
<synopsis>
<heading>SYNOPSIS
<code>
```

void spinLockTaskGive
 (
 spinlockTask\_t \*pLock /\* pointer to task-only spinlock \*/
 )

DESCRIPTION

This routine releases the task-only spinlock pointed to by *pLock*. Furthermore, it re-enables task pre-emption that had been disabled on the local CPU when the lock was acquired using **spinLockTaskTake()**. Calling this routine under the following circumstances is considered to be an error condition and has undefined behaviour:

- The calling task is not the one that acquired the spinlock.
- The caller is an ISR.
- The pLock argument does not point to a properly initialized task-only spinlock.

This function forces a read/write memory barrier before releasing the lock.

If **INCLUDE\_SPINLOCK\_DEBUG** is defined the following scenarios will cause a ED&R kernel fatal error which may reboot the target depending on the ED&R policy in place:

ner of		

2

	- If this routine is called from interrupt level	
	- If this routine is called within the context of CPU other than owr the spinlock	
RETURNS	N/A	
ERRNO	N/A	

SEE ALSO spinLockLib, spinLockTaskTake(), spinLockTaskInit()

## spinLockTaskInit()

NAME	<pre>spinLockTaskInit() - initialize a task-only spinlock</pre>		
SYNOPSIS	<pre>void spinLockTaskInit   (    spinlockTask_t *pLock, /* pointer to task-only spinlock */    int flags</pre>		
DESCRIPTION	This routine initializes the task-only spinlock pointed to by <i>pLock</i> , using the <i>flags</i> specified Currently, no flags are defined; this argument is a placeholder for future enhancements. A spinlock must be initialized before it is used for the first time. A spinlock is build on the ability of a processor to perform an atomic read-modify-write access to memory. Some CPUs may have cache attributes and memory alignment restrictions on the use of these instructions. It is the responsibility of the caller to ensure the memory location where the spinlock is located respects these restrictions, if any.		
	This routine must not be called from interrupt level.		
	If <b>INCLUDE_SPINLOCK_DEBUG</b> is defined the following scenarios will cause a ED&R kernel fatal error which may reboot the target depending on the ED&R policy in place:		
	- If this routine is called from interrupt level		
RETURNS	N/A		
ERRNO	N/A		
SEE ALSO	<pre>spinLockLib, spinLockTaskTake( ), spinLockTaskGive( ), spinLockIsrInit( ), VxWorks Architecture Supplements</pre>		

### spinLockTaskTake()

```
NAME
                 spinLockTaskTake() – take a task-only spinlock
SYNOPSIS
                 void spinLockTaskTake
                 #ifdef SPIN LOCK TRACE
                    spinlockTask_t *pLock, /* pointer to task-only spinlock */
                                    *file,
                    char
                    int
                                   line
                 #else
                     spinlockTask_t *pLock /* pointer to task-only spinlock */
                 #endif
                     )
DESCRIPTION
                 This routine acquires the task-only spinlock pointed to by pLock. If the lock is available at
                 the time of the call, this routine marks the spinlock as being in use and returns immediately.
                 If the spinlock is unavailable, the routine busy-waits for the lock to become available.
                 Because of this busy-wait characteristic, recursive acquisition of a spinlock causes a live lock
                 situation where the acquiring task busy-waits forever for a spinlock it already holds.
                 Acquisition of an task-only spinlock causes task pre-emption to be disabled on the local
                 CPU until the lock is released using spinLockTaskGive(). It is therefore recommended that
                 this type of spinlock be held for a minimal amount of time as it prevents scheduling on the
                 local CPU.
                 Calling this routine under the following circumstances is considered to be an error
                 condition and has undefined behaviour:
                      Calling this routine with a pLock that points to anything that is not a properly
                      initialized task-only spinlock.
                      The caller is an ISR.
                 This routine provides a memory barrier mechanism to prevent memory access reordering
                 that may be performed by the hardware.
                 If INCLUDE_SPINLOCK_DEBUG is defined the following scenarios will cause a ED&R
                 kernel fatal error which may reboot the target depending on the ED&R policy in place:

    If this routine is called from interrupt level

                 - If this routine is called within the context of a CPU already holding the
                  target spinlock
                 - If any other spinlock (any type) is held when this routine is called
                 N/A
RETURNS
ERRNO
                 N/A
```

SEE ALSO spinLockLib, spinLockTaskGive()

# sprintf()

NAME	<b>sprintf()</b> – write a formatted string to a buffer (ANSI)		
SYNOPSIS	<pre>int sprintf   (     char * buffer, /* buffer to write to */     const char * fmt, /* format string */     /* optional arguments to format */   )</pre>		
DESCRIPTION	This routine copies a formatted string to a specified buffer, which is null-terminated. Its function and syntax are otherwise identical to <b>printf()</b> .		
RETURNS	The number of characters copied to <i>buffer</i> , not including the NULL terminator.		
ERRNO	Not Available		
SEE ALSO	<b>fioBaseLib, printf( )</b> , American National Standard for Information Systems -, Programming Language - C, ANSI X3.159-1989: Input/Output ( <b>stdio.h</b> )		

# spy()

NAME	<b>spy()</b> – begin periodic task activity reports	
SYNOPSIS	<pre>void spy   (     int freq,    /* reporting freq in sec, 0 = default of 5 */     int ticksPerSec    /* interrupt clock freq, 0 = default of 100 */   )</pre>	
DESCRIPTION	This routine collects task activity data and periodically runs <b>spyReport()</b> . Data is gathe <i>ticksPerSec</i> times per second, and a report is made every <i>freq</i> seconds. If <i>freq</i> is zero, it defaults to 5 seconds. If <i>ticksPerSec</i> is omitted or zero, it defaults to 100.	
	This routine spawns <b>spyTask()</b> to do the actual reporting.	
	It is not necessary to call <b>spyClkStart()</b> before running <b>spy()</b> .	
RETURNS	N/A	

VxWorks Kernel API Reference, 6.6 spyClkStart()

#### ERRNO N/A

**SEE ALSO** usrLib, spyLib, spyClkStart(), spyTask(), the VxWorks programmer guides.

## spyClkStart()

NAME	<pre>spyClkStart() - start collecting task activity data</pre>	
SYNOPSIS	STATUS spyClkStart ( int intsPerSec /* timer interrupt freq, 0 = default of 100 */ )	
DESCRIPTION	This routine begins data collection by enabling the auxiliary clock interrupts at a frequency of <i>intsPerSec</i> interrupts per second. If <i>intsPerSec</i> is omitted or zero, the frequency will be 100. Data from previous collections is cleared.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if the CPU has no auxiliary clock, or if task create and delete hooks cannot be installed.	
ERRNO	N/A	
SEE ALSO	usrLib, spyLib, sysAuxClkConnect(), the VxWorks programmer guides.	

## spyClkStop()

NAME	<b>spyClkStop()</b> – stop collecting task activity data	
SYNOPSIS	void spyClkStop (void)	
DESCRIPTION	This routine disables the auxiliary clock interrupts. Data collected remains valid until the next <b>spyClkStart()</b> call.	
RETURNS	N/A	
ERRNO	N/A	
SEE ALSO	usrLib, spyLib, spyClkStart(), the VxWorks programmer guides.	

## spyHelp()

NAME	<b>spyHelp( )</b> – display task monitoring help menu		
SYNOPSIS	void spyHelp (void)		
DESCRIPTION	This routine displays a summary of <b>spyLib</b> utilities:		
	spyHelp spyClkStart [ticksPerSec]	Print this list Start task activity monitor running at ticksPerSec ticks per second	
	spyClkStop spyReport	Stop collecting data Prints display of task activity statistics	
	spyStop spy [freq[,ticksPerSec]]	Stop collecting data and reports Start spyClkStart and do a report every freq seconds	
	ticksPerSec defaults to 100.	freq defaults to 5 seconds.	
RETURNS	N/A		
ERRNO	N/A		
SEE ALSO	<b>usrLib</b> , <b>spyLib</b> , the VxWorks programmer guides.		

## spyLibInit()

NAMEspyLibInit() – initialize task cpu utilization tool packageSYNOPSISvoid spyLibInit (void)DESCRIPTIONThis routine initializes the task cpu utilization tool package. If the configuration macro<br/>INCLUDE\_SPY is defined, it is called by the root task, usrRoot(), in usrConfig.c.RETURNSN/AERRNONot AvailableSEE ALSOspyLib, usrLib

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NAME	<b>spyReport( )</b> – display task activity data
SYNOPSIS	void spyReport (void)
DESCRIPTION	This routine reports on data gathered at interrupt level for the amount of CPU time utilized by each task, the amount of time spent at interrupt level, the amount of time spent in the kernel, and the amount of idle time. Time is displayed in ticks and as a percentage, and the data is shown since both the last call to <b>spyClkStart()</b> and the last <b>spyReport()</b> . If no interrupts have occurred since the last <b>spyReport()</b> , nothing is displayed.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	usrLib, spyLib, spyClkStart(), the VxWorks programmer guides.

## spyStop()

NAME	<pre>spyStop() - stop spying and reporting</pre>
SYNOPSIS	void spyStop (void)
DESCRIPTION	This routine calls <b>spyClkStop()</b> . Any periodic reporting by <b>spyTask()</b> is terminated.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	usrLib, spyLib, spyClkStop(), spyTask(), the VxWorks programmer guides.

## spyTask()

NAME **spyTask()** – run periodic task activity reports

SYNOPSIS void spyTask

	( int freq /* reporting frequency, in seconds */ )
DESCRIPTION	This routine is spawned as a task by <b>spy()</b> to provide periodic task activity reports. It prints a report, delays for the specified number of seconds, and repeats.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	usrLib, spyLib, spy(), the VxWorks programmer guides.

# sqrtf()

NAME	<b>sqrtf()</b> – compute a non-negative square root (ANSI)
SYNOPSIS	float sqrtf ( float x /* value to compute the square root of */ )
DESCRIPTION	This routine returns the non-negative square root of $x$ in single precision.
RETURNS	The single-precision square root of <i>x</i> .
ERRNO	Not Available
SEE ALSO	mathALib

# sr()

NAME	$\mathbf{sr}(\mathbf{)}$ – return the contents of the status register (SH)
SYNOPSIS	<pre>int sr (     int taskId /* task ID, 0 means default task */ )</pre>

VxWorks Kernel API Reference, 6.6 sr( )

DESCRIPTION	This command extracts the contents of the status register from the TCB of a specified task. If <i>taskId</i> is omitted or zero, the last task referenced is assumed.
RETURNS	The contents of the status register.
ERRNO	Not Available
SEE ALSO	dbgArchLib, the VxWorks programmer guides.

## sr()

NAME	sr() – return the contents of control register $sr$ (also $gbr$ , $vbr$ ) (SH)
SYNOPSIS	<pre>int sr (     int taskId /* task ID, 0 means default task */ )</pre>
DESCRIPTION	This command extracts the contents of register sr from the TCB of a specified task. If <i>taskId</i> is omitted or zero, the last task referenced is assumed. Similar routines are provided for all control registers ( <b>gbr</b> , <b>vbr</b> ): <b>gbr()</b> , <b>vbr()</b> .
RETURNS	The contents of register sr (or the requested control register).
ERRNO	Not Available
SEE ALSO	dbgArchLib, the VxWorks programmer guides.

## sscanf()

NAME	sscanf() – read and convert characters from an ASCII string (ANSI)
SYNOPSIS	<pre>int sscanf   (     const char * str, /* string to scan */     const char * fmt, /* format string */     /* optional arguments to format string */     )</pre>

If there are insufficient arguments for the format, the behavior is undefined. If the format is exhausted while arguments remain, the excess arguments are evaluated but are otherwise ignored.

The format is a multibyte character sequence, beginning and ending in its initial shift state. The format is composed of zero or more directives: one or more white-space characters; an ordinary multibyte character (neither % nor a white-space character); or a conversion specification. Each conversion specification is introduced by the % character. After the %, the following appear in sequence:

- An optional assignment-suppressing character \*.
- An optional non-zero decimal integer that specifies the maximum field width.
- An optional h, l (ell) or ll (ell-ell) indicating the size of the receiving object. The conversion specifiers d, i, and n should be preceded by h if the corresponding argument is a pointer to `short int' rather than a pointer to int, or by l if it is a pointer to long int, or by ll if it is a pointer to long long int. Similarly, the conversion specifiers o, u, and x shall be preceded by h if the corresponding argument is a pointer to unsigned short int rather than a pointer to `unsigned int, or by l if it is a pointer to unsigned int, or by ll if it is a pointer to unsigned long int, or by ll if it is a pointer to `unsigned long int, or by ll if it is a pointer to unsigned long int, or by ll if it is a pointer to `unsigned long int, or by ll if it is a pointer to `unsigned long int, or by ll if it is a pointer to `unsigned long int, or by ll if it is a pointer to `unsigned long int, or by ll if it is a pointer to `unsigned long int, or by ll if it is a pointer to `unsigned long int, or by ll if it is a pointer to `unsigned long long int'. Finally, the conversion specifiers e, f, and g shall be preceded by l if the corresponding argument is a pointer to double rather than a pointer to float. If a h, l or ll appears with any other conversion specifier, the behavior is undefined.
- WARNING: ANSI C also specifies an optional L in some of the same contexts as I above, corresponding to a **long double** \* argument. However, the current release of the VxWorks libraries does not support **long double** data; using the optional L gives unpredictable results.
- A character that specifies the type of conversion to be applied. The valid conversion specifiers are described below.

The **sscanf()** routine executes each directive of the format in turn. If a directive fails, as detailed below, **sscanf()** returns. Failures are described as input failures (due to the unavailability of input characters), or matching failures (due to inappropriate input).

A directive composed of white-space character(s) is executed by reading input up to the first non-white-space character (which remains unread), or until no more characters can be read.

A directive that is an ordinary multibyte character is executed by reading the next characters of the stream. If one of the characters differs from one comprising the directive, the directive fails, and the differing and subsequent characters remain unread.

A directive that is a conversion specification defines a set of matching input sequences, as described below for each specifier. A conversion specification is executed in the following steps:

Input white-space characters (as specified by the **isspace()** function) are skipped, unless the specification includes a [, **c**, or **n** specifier.

An input item is read from the stream, unless the specification includes an **n** specifier. An input item is defined as the longest matching sequence of input characters, unless that exceeds a specified field width, in which case it is the initial subsequence of that length in the sequence. The first character, if any, after the input item remains unread. If the length of the input item is zero, the execution of the directive fails: this condition is a matching failure, unless an error prevented input from the stream, in which case it is an input failure.

Except in the case of a % specifier, the input item is converted to a type appropriate to the conversion specifier. If the input item is not a matching sequence, the execution of the directive fails: this condition is a matching failure. Unless assignment suppression was indicated by a \*, the result of the conversion is placed in the object pointed to by the first argument following the *fmt* argument that has not already received a conversion result. If this object does not have an appropriate type, or if the result of the conversion cannot be represented in the space provided, the behavior is undefined.

The following conversion specifiers are valid:

d

Matches an optionally signed decimal integer whose format is the same as expected for the subject sequence of the **strtol()** function with the value 10 for the *base* argument. The corresponding argument should be a pointer to **int**.

i

Matches an optionally signed integer, whose format is the same as expected for the subject sequence of the **strtol()** function with the value 0 for the *base* argument. The corresponding argument should be a pointer to **int**.

0

Matches an optionally signed octal integer, whose format is the same as expected for the subject sequence of the **strtoul()** function with the value 8 for the *base* argument. The corresponding argument should be a pointer to **unsigned int**.

u

Matches an optionally signed decimal integer, whose format is the same as expected for the subject sequence of the **strtoul()** function with the value 10 for the *base* argument. The corresponding argument should be a pointer to **unsigned int**.

x

Matches an optionally signed hexadecimal integer, whose format is the same as expected for the subject sequence of the **strtoul()** function with the value 16 for the *base* argument. The corresponding argument should be a pointer to **unsigned int**.

#### e, f, g

Match an optionally signed floating-point number, whose format is the same as expected for the subject string of the **strtod()** function. The corresponding argument should be a pointer to **float**.

 $\mathbf{s}$ 

Matches a sequence of non-white-space characters. The corresponding argument should be a pointer to the initial character of an array large enough to accept the sequence and a terminating null character, which will be added automatically.

[

Matches a non-empty sequence of characters from a set of expected characters (the **scanset**). The corresponding argument should be a pointer to the initial character of an array large enough to accept the sequence and a terminating null character, which is added automatically. The conversion specifier includes all subsequent character in the format string, up to and including the matching right bracket (]). The characters between the brackets (the **scanlist**) comprise the scanset, unless the character after the left bracket is a circumflex (^) in which case the scanset contains all characters that do not appear in the scanlist between the circumflex and the right bracket. If the conversion specifier begins with "[]" or "[^]", the right bracket character is in the scanlist and the next right bracket character is the matching right bracket that ends the specification; otherwise the first right bracket character is the one that ends the specification.

С

Matches a sequence of characters of the number specified by the field width (1 if no field width is present in the directive). The corresponding argument should be a pointer to the initial character of an array large enough to accept the sequence. No null character is added.

p

Matches an implementation-defined set of sequences, which should be the same as the set of sequences that may be produced by the %p conversion of the **fprintf()** function. The corresponding argument should be a pointer to a pointer to **void**. VxWorks defines its pointer input field to be consistent with pointers written by the **fprintf()** function ("0x" hexadecimal notation). If the input item is a value converted earlier during the same program execution, the pointer that results should compare equal to that value; otherwise the behavior of the %p conversion is undefined.

n

No input is consumed. The corresponding argument should be a pointer to **int** into which the number of characters read from the input stream so far by this call to **sscanf()** is written. Execution of a %n directive does not increment the assignment count returned when **sscanf()** completes execution.

%

Matches a single %; no conversion or assignment occurs. The complete conversion specification is %%.

	If a conversion specification is invalid, the behavior is undefined.
	The conversion specifiers <b>E</b> , <b>G</b> , and <b>X</b> are also valid and behave the same as <b>e</b> , <b>g</b> , and <b>x</b> , respectively.
	If end-of-file is encountered during input, conversion is terminated. If end-of-file occurs before any characters matching the current directive have been read (other than leading white space, where permitted), execution of the current directive terminates with an input failure; otherwise, unless execution of the current directive is terminated with a matching failure, execution of the following directive (if any) is terminated with an input failure.
	If conversion terminates on a conflicting input character, the offending input character is left unread in the input stream. Trailing white space (including new-line characters) is left unread unless matched by a directive. The success of literal matches and suppressed assignments is not directly determinable other than via the %n directive.
RETURNS	The number of input items assigned, which can be fewer than provided for, or even zero, in the event of an early matching failure; or <b>EOF</b> if an input failure occurs before any conversion.
ERRNO	Not Available
SEE ALSO	<b>fioLib</b> , <b>fscanf()</b> , <b>scanf()</b> , American National Standard for Information Systems -, Programming Language - C, ANSI X3.159-1989: Input/Output ( <b>stdio.h</b> )

## ssiDbInit( )

NAME	ssiDbInit() – Initialize SSI database.
SYNOPSIS	STATUS ssiDbInit ( ssiCompRegInfo_t * pComps )
DESCRIPTION	This routine initializes the SSI database. If <b>pComps</b> is not <b>NULL</b> , it registers all the components in <b>pComps</b> to its database. If the information provided by <b>pComps</b> is complete, the SSI dependency tree is also genrated.
	<b>pComps</b> points to component registration table, which is used to generate the dependency tree. The component dependency table is considered complete if each of the components in the table have either no dependency or if they have one or many, then all the dependency components are also in the table.
RETURNS	ERROR if event cannot be sent for one of the following reason: - specified descriptor is not valid

	<ul> <li>descriptor not specified and could not establish connection</li> <li>cnsWrite() fails</li> </ul>
ERRNO	N/A
SEE ALSO	ssiDb
	ssiShow()
NAME	ssiShow() – Display SSI information
SYNOPSIS	STATUS ssiShow ( char * pArgs )
DESCRIPTION	This routine displays information of a specified or all component(s) registered with SSI. The routine can be called with NULL argument or with an empty string, in which case all SSI groups and components are displayed.
ERRNO	N/A
RETURNS	ОК
ERRNO	Not Available

## ssmCompInfoGet( )

ssiDb

SEE ALSO

<pre>ssmCompInfoGet( ) – Get component information.</pre>
STATUS ssmCompInfoGet ( char * pName, ssmCompInfo_t ** ppInfo )

VxWorks Kernel API Reference, 6.6 ssmCompRegister()

DESCRIPTION	This routine retrieves the information of a component specified by <b>pName</b> .	
ARGUMENTS	<b>pName</b> - specifies the component.	
	<b>ppInfo</b> - points to a pointer to the component information structure.	
RETURNS	ERROR if the component cannot be found, OK otherwise.	
ERRNO	N/A	
SEE ALSO	ssiDb	

# ssmCompRegister()

NAME	<b>ssmCompRegister()</b> – Register a component with SSI Manager.		
SYNOPSIS	<pre>STATUS ssmCompRegister   (     char *</pre>		
DESCRIPTION	This routine registers a component with the SSI manager.		
	To participate in the SSI process, components have to be registered with the SSI Manager.		
ARGUMENTS	<b>pCompName</b> is a <b>NULL</b> terminated ASCII string that uniquely identifies the component in the system. The string has to be at least 2 byte long (including the terminator) and can be up to <b>SSM_NAME_LEN</b> long.		
	<b>pDependencyList</b> is a list of zero or more names of components that the component requires to wait on before being started. A comma (,) separates each name from another.		
	<b>launchFunc</b> is a pointer to a function to launch or initialize component. Depending upon whether the component is SSI-aware or not, this routine either starts the component to interact with the CSM or it actually starts		

the component.

The following is the optional usage of <b>launchFunc</b> . o For an SSI-aware component, <b>launchFunc</b> can be <b>NULL</b> if the component does not expect to be automatically launched by SSM. If <b>launchFunc</b> is provided, SSM launches the component, which will then wait until it receives <b>CSM_EVENT_INIT</b> before starting actual execution.		
o For a non-compliant component, <b>launchFunc</b> does the actual component initialization. There are two ways this can be implemented: - If the SSM wrapper function is to be used (default for non-compliant comps), the routine is passed to the wrapper function to be called later when the CSM allows the component to initialize.		
- If the component opts not to use the wrapper routine, SSM calls the routine directly when the CSM decides the component can be started.		
The following table summarizes this usage scenario: /ts SSI Aware   USE SSI Wrapper   Init Routine   Init Routine Usage		
FALSE	FALSE 	Required   SSM calls routine to   initialize the component.
FALSE	TRUE 	Reuired   Wrapper task calls routine   to initialize the component.
TRUE	FALSE   TR 	UE   NULL   Component to be launched by   some other means.
TRUE	FALSE   TR 	UE   Defined   SSM calls routine to launch   the component.
/te		, , ,

The routine should return its state after the initialization and should not block even if the "ready" state is not returned. The implication of this is that dependent components may not be able to proceed. VxWorks Kernel API Reference, 6.6 ssmCompRegister()

If the component is running in its own thread, it should send the CSM\_EVENT\_READY eventually. Otherwise, it is recommended that the component uses the default SSM wrapper.

For RTPs, the launch routine is the name of the RTP executable.

For kernel components, if the registration data is obtained from a storage medium, the launch routine field can be 0, in which case the process that is parsing the data has to lookup the component's name in a symbol table to get the initialization function address. If the registration data is input dynamically, the caller of the registration routine may have to pass a function pointer. This is especially true for downloadable kernel modules.

The initialization routine has the following function prototype:

```
STATUS compLaunch (csmStatus_t * pStatus);
```

**pArgs** points to an ASCCII character string. For kernel components, string format varies from component to component. For RTP components, the string format should include the RTP executable path and name and other RTP attributes including initial arguments.

**pOpts** represents 0 or more comma-separated strings indicating options. The strings could be one of the following:

"nowarpper" implies that the component does not opt to use the default SSM wrapper routine. Unless the "ssmaware" option is specified, the component's launch/initialization routine is directly called when the CSM allows the component to initialize.

"compliant" implies that component is SSI aware, meaning that the component is compliant with the rules of the SSI system. For SSI aware components, the SSM will simply call the component's registered initialization/launch routine when it is time to launch the component. Subsequently, the initialization process proceeds in the component's own thread or in the thread of the SSM.

"multinst" implies that the component can be started multiple times. This option affects the component's state table. By default, the CSM considers it an error if it receives a CSM\_EV\_CREATED event.

**RETURNS ERROR** if the component cannot be registered.

ERRNO N/A

SEE ALSO ssiDb

## startupScriptFieldSplit()

NAME	<pre>startupScriptFieldSplit( ) – Split the startup script field of the bootline</pre>
SYNOPSIS	char * startupScriptFieldSplit ( char * field )
DESCRIPTION	This routine splits the startup script field of the bootline at the first occurence of a <b>#</b> character and null-terminates it at that location. The text before the <b>#</b> is the name of a traditional startup script file containing shell commands. Everything following the first <b>#</b> is part of a list of RTP's to startup.
RETURNS	Pointer to a string containing the name of a shell startup script or NULL.
ERRNO	N/A.
SEE ALSO	usrRtpStartup, the VxWorks programmer guides.

## stat()

NAME	<b>stat()</b> – get file status information using a pathname (POSIX)	
SYNOPSIS	STATUS stat ( const char * name, struct stat *pStat )	/* name of file to check */ /* pointer to stat structure */

VxWorks Kernel API Reference, 6.6 statfs()

DESCRIPTION	This routine obtains various characteristics of a file (or directory). This routine is equivalent to <b>fstat()</b> , except that the <i>name</i> of the file is specified, rather than an open file descriptor.
	The <i>pStat</i> parameter is a pointer to a <b>stat</b> structure (defined in <b>stat.h</b> ). This structure must have already been allocated before this routine is called.
NOTE	When used with <b>netDrv</b> devices (FTP or RSH), <b>stat()</b> returns the size of the file and always sets the mode to regular; <b>stat()</b> does not distinguish between files, directories, links, etc.
	Upon return, the fields in the <b>stat</b> structure are updated to reflect the characteristics of the file.
RETURNS	OK or ERROR, from the underlying io commands <b>open()</b> , <b>ioctl()</b> , or <b>close()</b> .
ERRNO	See open(), ioctl(), and close().
SEE ALSO	dirLib, fstat(), ls()

# statfs()

NAME	<b>statfs( )</b> – get file status information us	ing a pathname (POSIX)
SYNOPSIS	STATUS statfs ( char *name, struct statfs *pStat )	/* name of file to check */ /* pointer to statfs structure */
DESCRIPTION	This routine obtains various characteristics of a file system. This routine is equivalent to <b>fstatfs()</b> , except that the <i>name</i> of the file is specified, rather than an open file descriptor.	
	The <i>pStat</i> parameter is a pointer to a <b>sta</b> have already been allocated before this	<b>.tfs</b> structure (defined in <b>stat.h</b> ). This structure must s routine is called.
	Upon return, the fields in the <b>statfs</b> stru file.	acture are updated to reflect the characteristics of the
RETURNS	OK or ERROR, from the underlying IO	commands <b>open()</b> , <b>ioctl()</b> , <b>close()</b> .
ERRNO	EBADF Bad file descriptor number.	
	S_ioLib_UNKNOWN_REQUEST (ENOS) Device driver does not support the	( <b>S</b> ) e ioctl command.

ELOOP Circular symbolic link, too many links.
EMFILE Maximum number of files already open.
<b>S_iosLib_DEVICE_NOT_FOUND</b> (ENODEV) No valid device name found in path.
Other Other errors reported by device driver.

SEE ALSO dirLib, fstatfs(), ls()

## strFree()

NAME	<b>strFree( )</b> – free shell strings
SYNOPSIS	<pre>void strFree   (    char * string /* shell string pointer to free, or 0, or -1 */ )</pre>
DESCRIPTION	This command free strings allocated within the shell. If <i>string</i> is <b>NULL</b> , all allocated strings are displayed. If <i>string</i> is -1, all allocated strings are freed.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	usrLib, the VxWorks programmer guides.

## swab()

NAME	<b>swab()</b> – swap bytes		
SYNOPSIS	void swab ( char *source, char *destination, int nbytes )	<pre>/* pointer to source buffer /* pointer to destination buffer /* number of bytes to exchange</pre>	*/ */ */

VxWorks Kernel API Reference, 6.6 symAdd()

DESCRIPTION	This routine gets the specified number of bytes from <i>source</i> , exchanges the adjacent even and odd bytes, and puts them in <i>destination</i> . The buffers <i>source</i> and <i>destination</i> should not overlap.
	NOTE: On some CPUs, <b>swab()</b> will cause an exception if the buffers are unaligned. In such cases, use <b>uswab()</b> for unaligned swaps. On ARM family CPUs, <b>swab()</b> may reorder the bytes incorrectly without causing an exception if the buffers are unaligned. Again, use <b>uswab()</b> for unaligned swaps.
	The value of <i>nBytes</i> must not be odd. Failure to adhere to this may yield incorrect results.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	bLib, uswab( )

# symAdd()

NAME	symAdd() – create and add a symbol to a symbol table, including a group number	
SYNOPSIS	<pre>STATUS symAdd  (    SYMTAB_ID symTblId, /* symbol table to add symbol to */    char *name, /* pointer to symbol name string */    char *value, /* symbol address */    SYM_TYPE type, /* symbol type */    UINT16 group /* symbol group */   )</pre>	
DESCRIPTION	This routine allocates a symbol with the specified <i>name, value, type,</i> and <i>group</i> and adds it to the symbol table specified by the <i>symTblId</i> parameter.	
	The <i>group</i> parameter specifies the group number assigned to a module when it is loaded; see the manual entry for <b>moduleLib</b> .	
RETURNS	<b>OK</b> , or <b>ERROR</b> if the symbol table is invalid there is insufficient memory for the symbol to be allocated, or any other fatal error occurs.	
ERRNO	Possible errnos set by this routine include:	
	+ S_symLib_INVALID_SYMTAB_ID	
	+ S_symLib_INVALID_SYMBOL_NAME	
	+ S_symLib_NAME_CLASH	

For a complete description of the errnos, see the reference documentation for **symLib**.

SEE ALSO symLib, moduleLib

## symByValueAndTypeFind()

NAME	symByValueAndTypeFind() – look up a symbol by value and type
SYNOPSIS	<pre>STATUS symByValueAndTypeFind   (    SYMTAB_ID symTblId, /* ID of symbol table to look in */    UINT value, /* value of symbol to find */    char ** pName, /* where to return symbol name string */    int * pValue, /* where to put symbol value */    SYM_TYPE * pType, /* where to put symbol type */    SYM_TYPE sType, /* symbol type to look for */    SYM_TYPE mask /* bits in <stype> to pay attention to */   )</stype></pre>
DESCRIPTION	This routine searches a symbol table for a symbol matching both the specified value and the specified type ( <i>value</i> and <i>sType</i> ). If there is no matching entry, it returns the table entry with the next lowest value (among entries with matching type). A pointer to the symbol name string (with terminating EOS) is returned in <i>pName</i> . The actual value and the type are copied to <i>pValue</i> and <i>pType</i> . The <i>mask</i> parameter can be used to match sub-classes of type. <i>pName</i> is a pointer to memory allocated by symByValueAndTypeFind; the memory must be freed by the caller after the use of <i>pName</i> . To search the global VxWorks symbol table, specify <b>sysSymTbl</b> as the <i>symTblId</i> parameter.
RETURNS	<b>OK</b> or <b>ERROR</b> if <i>symTblld</i> is invalid, <i>pName</i> is <b>NULL</b> , or <i>value</i> is less than the lowest value in the table.
ERRNO	Possible errnos set by this routine include:
	+ S_symLib_INVALID_SYMTAB_ID
	+ S_symLib_INVALID_SYM_ID_PTR
	+ S_symLib_SYMBOL_NOT_FOUND
	For a complete description of the errnos, see the reference documentation for <b>symLib</b> .
SEE ALSO	symLib, symFindSymbol()

	symByValueFind()
NAME	<pre>symByValueFind() – look up a symbol by value</pre>
SYNOPSIS	<pre>STATUS symByValueFind   (    SYMTAB_ID symTblId, /* ID of symbol table to look in */    UINT value, /* value of symbol to find */    char ** pName, /* where return symbol name string */    int * pValue, /* where to put symbol value */    SYM_TYPE * pType /* where to put symbol type */   )</pre>

This routine searches a symbol table for a symbol whose value matches the specified value. DESCRIPTION If there is no matching entry, it chooses the table entry with the next lowest value. A pointer to the symbol name string (with terminating EOS) is returned in *pName*. The actual value and the type are copied to the memory pointed to by *pValue* and *pType*.

> *pName* is a pointer to memory allocated by symByValueFind, not to an internal copy of the symbol's name; the memory must be freed by the caller after the use of *pName*.

> To search the global VxWorks symbol table, specify **sysSymTbl** as the *symTblId* parameter.

- RETURNS **OK** or **ERROR** if *symTblld* is invalid, *pName* is **NULL**, or *value* is less than the lowest value in the table.
- ERRNO Possible errnos set by this routine include:
  - S\_symLib\_INVALID\_SYMTAB\_ID +
  - + S\_symLib\_INVALID\_SYM\_ID\_PTR
  - S\_symLib\_SYMBOL\_NOT\_FOUND +

For a complete description of the errnos, see the reference documentation for symLib.

SEE ALSO symLib, symByValueAndTypeFind( )

#### symEach()

NAME **symEach()** – call a routine to examine each entry in a symbol table SYNOPSIS SYMBOL \* symEach (

```
SYMTAB_ID symTblId, /* pointer to symbol table */
FUNCPTR routine,
                    /* func to call for each tbl entry */
```

int routineArg /\* arbitrary user-supplied arg \*/
)

**DESCRIPTION** This routine calls a user-supplied routine to examine each entry in the symbol table; it calls the specified routine once for each entry. The routine should be declared as follows:

```
BOOL routine

(

char * name, /* symbol/entry name */

int val, /* symbol/entry value */

SYM_TYPE type, /* symbol/entry type */

int arg, /* arbitrary user-supplied arg */

UINT16 group /* symbol/entry group number */

)
```

The user-supplied routine should return TRUE if **symEach()** is to continue calling it for each entry, or **FALSE** if it is done and **symEach()** can exit.

**RETURNS** A pointer to the last symbol reached, or **NULL** if all symbols are reached or there is an error.

**ERRNO** Possible errnos set by this routine include:

+ S\_symLib\_INVALID\_SYMTAB\_ID

For a complete description of the errnos, see the reference documentation for **symLib**.

SEE ALSO symLib

### symFindByName()

NAME	<b>symFindByName( )</b> – look up a symbol by name
SYNOPSIS	<pre>STATUS symFindByName  (    SYMTAB_ID symTblId, /* ID of symbol table to look in */    char * name, /* symbol name to look for */    char ** pValue, /* where to return symbol value */    SYM_TYPE * pType /* where to return symbol type */ )</pre>
DESCRIPTION	This routine searches a symbol table for a symbol matching the specified name. If a symbol is found, its value and type are copied to the memory pointed to by <i>pValue</i> and <i>pType</i> .
	If multiple symbols have the same name, the routine returns the matching symbol most recently added to the symbol table.

To search the global VxWorks (kernel) symbol table, specify **sysSymTbl** as the *symTblId*.

VxWorks Kernel API Reference, 6.6 symFindByNameAndType()

**RETURNS** OK, or ERROR if the symbol table ID is invalid or the symbol cannot be found.

ERRNO

+ S symLib INVALID SYMTAB ID

+ S\_symLib\_INVALID\_SYMTAB\_ID

Possible errnos set by this routine include:

- + S\_symLib\_INVALID\_SYM\_ID\_PTR
- + S\_symLib\_SYMBOL\_NOT\_FOUND

For a complete description of the errnos, see the reference documentation for symLib.

SEE ALSO symLib

### symFindByNameAndType()

**symFindByNameAndType()** – look up a symbol by name and type NAME SYNOPSIS STATUS symFindByNameAndType ( SYMTAB\_ID symTblId, /\* ID of symbol table to look in \*/ char \* name, /\* symbol name to look for \*/ char \*\* pValue, /\* where to put symbol value \*/ SYM\_TYPE \* pType,/\* where to put symbol type \*/SYM\_TYPE sType,/\* symbol type to look for \*/ /\* bits in <sType> to pay attention to \*/ SYM\_TYPE mask ) DESCRIPTION This routine searches a symbol table for a symbol with matching name and type (name and *sType*). If the symbol is found, its value and type are copied to the memory pointed to by the pointers *pValue* and *pType*. The *mask* parameter can be used to match sub-classes of type. To search the global VxWorks (kernel) symbol table, specify **sysSymTbl** as the *symTblld* parameter. RETURNS OK, or ERROR if the symbol table ID is invalid or the symbol is not found. ERRNO Possible errnos set by this routine include: S\_symLib\_INVALID\_SYMTAB\_ID + + S\_symLib\_INVALID\_SYM\_ID\_PTR + S\_symLib\_SYMBOL\_NOT\_FOUND For a complete description of the errnos, see the reference documentation for **symLib**.

SEE ALSO symLib

# symFindByValue()

NAME	<b>symFindByValue( )</b> – look up a symbol by value
SYNOPSIS	<pre>STATUS symFindByValue  (  SYMTAB_ID symTblId, /* ID of symbol table to look in */ UINT value, /* value of symbol to find */ char * name, /* where to put symbol name string */ int * pValue, /* where to put symbol value */ SYM_TYPE * pType /* where to put symbol type */ )</pre>
DESCRIPTION	This routine is obsolete. It is replaced by <b>symByValueFind()</b> and will be removed in the next version of VxWorks.
	This routine searches a symbol table for a symbol matching a specified value. If there is no matching entry, it chooses the table entry with the next lowest value. The symbol name (with terminating EOS), the actual value, and the type are copied to <i>name</i> , <i>pValue</i> , and <i>pType</i> .
	For the <i>name</i> buffer, allocate MAX_SYS_SYM_LEN + 1 bytes. The value MAX_SYS_SYM_LEN is defined in <b>sysSymTbl.h</b> . If the name of the symbol is longer than MAX_SYS_SYM_LEN bytes, it will be truncated to fit into the buffer. Whether or not the name was truncated, the string returned in the buffer will be null-terminated.
	To search the global VxWorks symbol table, specify <b>sysSymTbl</b> as the <i>symTblId</i> parameter.
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>symTblId</i> is invalid or <i>value</i> is less than the lowest value in the table.
ERRNO	Possible errnos set by this routine include:
	+ S_symLib_INVALID_SYMTAB_ID
	+ S_symLib_INVALID_SYM_ID_PTR
	+ S_symLib_SYMBOL_NOT_FOUND
	For a complete description of the errnos, see the reference documentation for <b>symLib</b> .
SEE ALSO	symLib

	symFindByValueAndType()
NAME	<pre>symFindByValueAndType() – look up a symbol by value and type</pre>
SYNOPSIS	<pre>STATUS symFindByValueAndType   (    SYMTAB_ID symTblId, /* ID of symbol table to look in */    UINT value, /* value of symbol to find */    char * name, /* where to put symbol name string */    int * pValue, /* where to put symbol value */    SYM_TYPE * pType, /* where to put symbol type */    SYM_TYPE sType, /* symbol type to look for */    SYM_TYPE mask /* bits in <stype> to pay attention to */   )</stype></pre>
DESCRIPTION	This routine is obsolete. It is replaced by the routine <b>symByValueAndTypeFind()</b> and will be removed in the next version of VxWorks.
	This routine searches a symbol table for a symbol matching both the specified value and type ( <i>value</i> and <i>sType</i> ). If there is no matching entry, it returns the symbol table entry with the next lowest value. The symbol name (with terminating EOS), the actual value, and the type are copied to the memory pointed to by <i>name</i> , <i>pValue</i> , and <i>pType</i> . The <i>mask</i> parameter can be used to match sub-classes of type.
	For the <i>name</i> buffer, allocate MAX_SYS_SYM_LEN + 1 bytes. The value MAX_SYS_SYM_LEN is defined in <b>sysSymTbl.h</b> . If the name of the symbol is longer than MAX_SYS_SYM_LEN bytes, it will be truncated to fit into the buffer. Whether or not the name was truncated, the string returned in the buffer will be null-terminated.
	To search the global VxWorks symbol table, specify <b>sysSymTbl</b> as the <i>symTblId</i> parameter.
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>symTblId</i> is invalid or <i>value</i> is less than the lowest value in the table.
ERRNO	<ul> <li>Possible errnos set by this routine include:</li> <li>+ S_symLib_INVALID_SYMTAB_ID</li> <li>+ S_symLib_INVALID_SYM_ID_PTR</li> <li>+ S_symLib_SYMBOL_NOT_FOUND</li> <li>For a complete description of the errnos, see the reference documentation for symLib.</li> </ul>
SEE ALSO	symLib

NAME	<pre>symLibInit() – initialize the symbol table library</pre>
SYNOPSIS	STATUS symLibInit (void)
DESCRIPTION	This routine initializes the symbol table library. If the configuration macro INCLUDE_SYM_TBL is defined, <b>symLibInit()</b> is called by the root task, <b>usrRoot()</b> , in <b>usrConfig.c</b> .
RETURNS	<b>OK</b> , or <b>ERROR</b> if the library could not be initialized.
ERRNO	Not Available
SEE ALSO	symLib

# symRemove()

NAME	<b>symRemove( )</b> – remove a symbol from a symbol table
SYNOPSIS	STATUS symRemove ( SYMTAB_ID symTblId, /* symbol tbl to remove symbol from */ char *name, /* name of symbol to remove */ SYM_TYPE type /* type of symbol to remove */ )
DESCRIPTION	This routine removes a symbol with matching name and type from a specified symbol table. The symbol is deallocated if found.
	Note that VxWorks symbols in a standalone VxWorks image (where the symbol table is linked in) cannot be removed.
RETURNS	OK, or ERROR if the symbol is not found or could not be deallocated.
ERRNO	Possible errnos set by this routine include:
	+ S_symLib_INVALID_SYMTAB_ID
	+ S_symLib_INVALID_SYM_ID_PTR
	+ S_symLib_SYMBOL_NOT_FOUND
	For a complete description of the errnos, see the reference documentation for <b>symLib</b> .

SEE ALSO symLib

## symShow()

NAME	<b>symShow()</b> – show the symbols of specified symbol table with matching substring
SYNOPSIS	<pre>STATUS symShow   (    SYMTAB_ID pSymTbl, /* ID of symbol table involved */    char * substr /* substring to match */   )</pre>
DESCRIPTION	This routine lists all symbols in the specified symbol table whose names contain the string <i>substr</i> . If <i>substr</i> is an empty string (""), all symbols in the table will be listed. If <i>substr</i> is <b>NULL</b> then the symbol table structure will be summarized
EXAMPLES	The system symbol table ID is stored in the global variable <i>sysSymTbl</i> .
	Look for a symbol containing the "vxWorks" substring (C shell):
	-> symShow (sysSymTbl, "vxWorks")
	Print out general information from the system symbol table (C shell):
	-> symShow (sysSymTbl, 0)
	Print all symbols contained in the system symbol table (C shell):
	-> symShow (sysSymTbl,"")
RETURNS	OK, or ERROR if the symbol table ID is invalid
ERRNO	Possible errnos set by this routine include:
	+ S_symLib_INVALID_SYMTAB_ID
	For a complete description of the errnos, see the reference documentation for <b>symShow</b> .
SEE ALSO	symShow, symLib, symEach()

# symShowInit()

NAME	<pre>symShowInit() - initialize symbol table show routine</pre>
SYNOPSIS	void symShowInit (void)
DESCRIPTION	This routine links the symbol table show facility into the VxWorks system. It is called automatically when the symbol table show facility is configured into VxWorks by including the INCLUDE_SYM_TBL_SHOW component.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	symShow

# symTblCreate()

NAME	symTblCreate() – create a symbol table
SYNOPSIS	<pre>SYMTAB_ID symTblCreate ( int hashSizeLog2, /* size of hash table as a power of 2 */ BOOL sameNameOk, /* allow 2 symbols of same name &amp; type */ PART_ID symPartId /* memory part ID for symbol allocation */ )</pre>
DESCRIPTION	This routine creates and initializes a symbol table with a hash table of a specified size. The size of the hash table is specified as a power of two. For example, if <i>hashSizeLog2</i> is 6, a 64-entry hash table is created.
	If the <i>sameNameOk</i> parameter is <b>FALSE</b> , attempting to add a symbol with the same name and type as an already-existing symbol in the symbol table will result in an error. This behavior cannot be changed once the symbol table has been created.
	Memory for storing symbols as they are added to the symbol table will be allocated from the memory partition <i>symPartId</i> . Note: the ID of the system memory partition is stored in the global variable <b>memSysPartId</b> , which is declared in <b>memLib.h</b> .
RETURNS	Symbol table ID, or <b>NULL</b> if sufficient memory is not available or another fatal error occurred.
ERRNO	Not Available

SEE ALSO symLib

## symTblDelete()

NAME	symTblDelete( ) – delete a symbol table
SYNOPSIS	STATUS symTblDelete ( SYMTAB_ID symTblId /* ID of symbol table to delete */ )
DESCRIPTION	This routine deletes a specified symbol table. It deallocates all associated memory, including the hash table, and marks the table as invalid.
	An attempt to delete a table that still contains symbols will return <b>ERROR</b> . Successful deletion includes the deletion of the internal hash table and the deallocation of memory associated with the table. The table is marked invalid to prohibit any future references.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the symbol table ID is invalid or if there was a problem.
ERRNO	Possible errnos set by this routine include:
	+ S_symLib_INVALID_SYMTAB_ID
	+ S_symLib_TABLE_NOT_EMPTY
	For a complete description of the errnos, see the reference documentation for <b>symLib</b> .
SEE ALSO	symLib

## sysAuxClkConnect()

NAME	<pre>sysAuxClkConnect() - connect a routine to the auxiliary clock interrupt</pre>
SYNOPSIS	STATUS sysAuxClkConnect ( FUNCPTR routine, /* routine called at each aux clock interrupt */ int arg /* argument to auxiliary clock interrupt routine */ )
DESCRIPTION	This routine specifies the interrupt service routine to be called at each auxiliary clock interrupt. It does not enable auxiliary clock interrupts.

NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	OK, or ERROR if the routine cannot be connected to the interrupt.
ERRNO	Not Available
SEE ALSO	<b>sysLib</b> , <b>intConnect( )</b> , <b>sysAuxClkEnable( )</b> , and BSP-specific reference pages for this routine.

## sysAuxClkDisable()

NAME	<pre>sysAuxClkDisable() - turn off auxiliary clock interrupts</pre>
SYNOPSIS	void sysAuxClkDisable (void)
DESCRIPTION	This routine disables auxiliary clock interrupts.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	sysLib, sysAuxClkEnable(), and BSP-specific reference pages for this routine.

## sysAuxClkEnable()

- NAME sysAuxClkEnable() turn on auxiliary clock interrupts
- SYNOPSIS void sysAuxClkEnable (void)
- **DESCRIPTION** This routine enables auxiliary clock interrupts.

VxWorks Kernel API Reference, 6.6 sysAuxClkRateGet()

NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	<b>sysLib</b> , <b>sysAuxClkConnect( )</b> , <b>sysAuxClkDisable( )</b> , <b>sysAuxClkRateSet( )</b> , and BSP-specific reference pages for this routine.

## sysAuxClkRateGet()

NAME	<b>sysAuxClkRateGet()</b> – get the auxiliary clock rate
SYNOPSIS	int sysAuxClkRateGet (void)
DESCRIPTION	This routine returns the interrupt rate of the auxiliary clock.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	The number of ticks per second of the auxiliary clock.
ERRNO	Not Available
SEE ALSO	<pre>sysLib, sysAuxClkEnable(), sysAuxClkRateSet(), and BSP-specific reference pages for this routine.</pre>

## sysAuxClkRateSet()

NAME	<pre>sysAuxClkRateSet() - set the auxiliary clock rate</pre>
SYNOPSIS	<pre>STATUS sysAuxClkRateSet   (    int ticksPerSecond /* number of clock interrupts per second */ )</pre>
DESCRIPTION	This routine sets the interrupt rate of the auxiliary clock. It does not enable auxiliary clock interrupts.
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NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	OK, or ERROR if the tick rate is invalid or the timer cannot be set.
ERRNO	Not Available
SEE ALSO	<pre>sysLib, sysAuxClkEnable(), sysAuxClkRateGet(), and BSP-specific reference pages for this routine.</pre>

## sysBspRev()

NAME	<b>sysBspRev()</b> – return the BSP version and revision number
SYNOPSIS	char * sysBspRev (void)
DESCRIPTION	This routine returns a pointer to a BSP version and revision number, for example, 1.0/1. <b>BSP_REV</b> is concatenated to <b>BSP_VERSION</b> and returned.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	A pointer to the BSP version/revision string.
ERRNO	Not Available
SEE ALSO	sysLib, and BSP-specific reference pages for this routine.

## sysBusIntAck()

NAME sysBusIntAck() – acknowledge a bus interrupt

SYNOPSIS int sysBusIntAck

VxWorks Kernel API Reference, 6.6 sysBusIntGen()

	( int intLevel /* interrupt level to acknowledge */ )
DESCRIPTION	This routine acknowledges a specified VMEbus interrupt level.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	NULL.
ERRNO	Not Available
SEE ALSO	sysLib, sysBusIntGen(), and BSP-specific reference pages for this routine.

# sysBusIntGen()

NAME	sysBusIntGen() – generate a bus interrupt	
SYNOPSIS	<pre>STATUS sysBusIntGen   (    int intLevel, /* bus interrupt level to generate */    int vector /* interrupt vector to generate (0-255) */ )</pre>	
DESCRIPTION	This routine generates a bus interrupt for a specified level with a specified vector.	
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.	
RETURNS	OK, or ERROR if <i>intLevel</i> is out of range or the board cannot generate a bus interrupt.	
ERRNO	Not Available	
SEE ALSO	sysLib, sysBusIntAck(), and BSP-specific reference pages for this routine.	

# sysBusTas()

NAME	<b>sysBusTas( )</b> – test and set a location across the bus
SYNOPSIS	BOOL sysBusTas ( char * adrs /* address to be tested and set */ )
DESCRIPTION	This routine performs a test-and-set instruction across the backplane.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
NOTE	This routine is equivalent to <b>vxTas( )</b> .
RETURNS	TRUE if the value had not been set but is now, or FALSE if the value was set already.
ERRNO	Not Available
SEE ALSO	sysLib, vxTas(), and BSP-specific reference pages for this routine.

# sysBusToLocalAdrs()

NAME	sysBusToLocalAdrs( ) – convert a bus address to a local address	
SYNOPSIS	<pre>STATUS sysBusToLocalAdrs   (     int adrsSpace, /* bus address space in which busAdrs resides     char * busAdrs, /* bus address to convert     char ** pLocalAdrs /* where to return local address )</pre>	* / * / * /
DESCRIPTION	This routine gets the local address that accesses a specified bus memory address.	
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for info specific to your BSP's version of this routine, see the reference pages for your BSP.	ormation
RETURNS	OK, or ERROR if the address space is unknown or the mapping is not possible.	

VxWorks Kernel API Reference, 6.6 sysClkConnect()

ERRNO Not Available

**SEE ALSO** sysLib, sysLocalToBusAdrs(), and BSP-specific reference pages for this routine.

## sysClkConnect()

NAME	<b>sysClkConnect()</b> – connect a routine to the system clock interrupt
SYNOPSIS	<pre>STATUS sysClkConnect   (   FUNCPTR routine, /* routine called at each system clock interrupt */   int arg /* argument with which to call routine */   )</pre>
DESCRIPTION	This routine specifies the interrupt service routine to be called at each clock interrupt. Normally, it is called from <b>usrRoot()</b> in <b>usrConfig.c</b> to connect <b>usrClock()</b> to the system clock interrupt.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURN	OK, or ERROR if the routine cannot be connected to the interrupt.
RETURNS	Not Available
ERRNO	Not Available
SEE ALSO	<b>sysLib</b> , <b>intConnect()</b> , <b>usrClock()</b> , <b>sysClkEnable()</b> , and BSP-specific reference pages for this routine.

## sysClkDisable()

NAME sysClkDisable()	- turn off system clock interrupts
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- SYNOPSIS void sysClkDisable (void)
- **DESCRIPTION** This routine disables system clock interrupts.

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NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	sysLib, sysClkEnable(), and BSP-specific reference pages for this routine.

## sysClkEnable()

NAME	<pre>sysClkEnable() - turn on system clock interrupts</pre>
SYNOPSIS	void sysClkEnable (void)
DESCRIPTION	This routine enables system clock interrupts.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	<b>sysLib</b> , <b>sysClkConnect()</b> , <b>sysClkDisable()</b> , <b>sysClkRateSet()</b> , and BSP-specific reference pages for this routine.

# sysClkRateGet()

- NAME sysClkRateGet() get the system clock rate
- SYNOPSIS int sysClkRateGet (void)
- **DESCRIPTION** This routine returns the system clock rate.

VxWorks Kernel API Reference, 6.6 sysClkRateSet()

NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	The number of ticks per second of the system clock.
ERRNO	Not Available
SEE ALSO	<pre>sysLib, sysClkEnable(), sysClkRateSet(), and BSP-specific reference pages for this routine.</pre>

# sysClkRateSet()

NAME	<b>sysClkRateSet( )</b> – set the system clock rate
SYNOPSIS	STATUS sysClkRateSet ( int ticksPerSecond /* number of clock interrupts per second */ )
DESCRIPTION	This routine sets the interrupt rate of the system clock. It is called by <b>usrRoot( )</b> in <b>usrConfig.c</b> .
	There may be interactions between this routine and the POSIX <b>clockLib</b> routines. Refer to the <b>clockLib</b> reference entry.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	OK, or ERROR if the tick rate is invalid or the timer cannot be set.
ERRNO	Not Available
SEE ALSO	<pre>sysLib, sysClkEnable(), sysClkRateGet(), clockLib, and BSP-specific reference pages for this routine.</pre>

# sysHwInit()

NAME	<pre>sysHwInit() - initialize the system hardware</pre>
SYNOPSIS	void sysHwInit (void)
DESCRIPTION	This routine initializes various features of the board. It is called from <b>usrInit()</b> in <b>usrConfig.c</b> .
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
NOTE	This routine should not be called directly by the user application.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	sysLib, and BSP-specific reference pages for this routine.

# sysIntDisable()

NAME	<b>sysIntDisable( )</b> – disable a bus interrupt level
SYNOPSIS	STATUS sysIntDisable ( int intLevel /* interrupt level to disable */ )
DESCRIPTION	This routine disables a specified bus interrupt level.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>intLevel</i> is out of range.
ERRNO	Not Available
SEE ALSO	sysLib, sysIntEnable(), and BSP-specific reference pages for this routine.

# sysIntEnable()

NAME	sysIntEnable( ) – enable a bus interrupt level
SYNOPSIS	STATUS sysIntEnable ( int intLevel /* interrupt level to enable (1-7) */ )
DESCRIPTION	This routine enables a specified bus interrupt level.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>intLevel</i> is out of range.
ERRNO	Not Available
SEE ALSO	sysLib, sysIntDisable(), and BSP-specific reference pages for this routine.

# sysLocalToBusAdrs()

NAME	sysLocalToBusAdrs( ) – convert a local address to a bus address
SYNOPSIS	STATUS sysLocalToBusAdrs ( int adrsSpace, /* bus address space in which busAdrs resides */ char * localAdrs, /* local address to convert */ char ** pBusAdrs /* where to return bus address */ )
DESCRIPTION	This routine gets the bus address that accesses a specified local memory address.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	OK, or ERROR if the address space is unknown or not mapped.
ERRNO	Not Available

## sysMailboxConnect()

NAME	<pre>sysMailboxConnect() - connect a routine to the mailbox interrupt</pre>
SYNOPSIS	<pre>STATUS sysMailboxConnect   (    FUNCPTR routine, /* routine called at each mailbox interrupt */    int arg /* argument with which to call routine */   )</pre>
DESCRIPTION	This routine specifies the interrupt service routine to be called at each mailbox interrupt.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	OK, or ERROR if the routine cannot be connected to the interrupt.
ERRNO	Not Available
SEE ALSO	<b>sysLib</b> , <b>intConnect( )</b> , <b>sysMailboxEnable( )</b> , and BSP-specific reference pages for this routine.

## sysMailboxEnable()

NAME	<pre>sysMailboxEnable( ) - enable the mailbox interrupt</pre>
SYNOPSIS	STATUS sysMailboxEnable ( char * mailboxAdrs /* address of mailbox (ignored) */ )
DESCRIPTION	This routine enables the mailbox interrupt.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.

VxWorks Kernel API Reference, 6.6 sysMemTop()

**RETURNS OK**, always.

ERRNO Not Available

**SEE ALSO** sysLib, sysMailboxConnect(), and BSP-specific reference pages for this routine.

## sysMemTop()

NAME	<b>sysMemTop()</b> – get the address of the top of logical memory
SYNOPSIS	char * sysMemTop (void)
DESCRIPTION	This routine returns the address of the top of memory.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	The address of the top of memory.
ERRNO	Not Available
SEE ALSO	sysLib, and BSP-specific reference pages for this routine.

## sysModel()

NAME	sysModel() – return the model name of the CPU board
SYNOPSIS	char * sysModel (void)
DESCRIPTION	This routine returns the model name of the CPU board.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	A pointer to a string containing the board name.
ERRNO	Not Available

## sysNanoDelay()

NAME	<pre>sysNanoDelay() – delay for specified number of nanoseconds</pre>
SYNOPSIS	void sysNanoDelay ( UINT32 nanoseconds /* nanoseconds to delay */ )
DESCRIPTION	This is an optional API for BSPs to provide. Some, but not all, drivers do require the BSP to implement this function.
	When implemented, this function implements a spin loop type delay for at least the specified number of nanoseconds. This is not a task delay, control of the processor is not given up to another task. The actual delay must be equal to or greater than the requested number of nanoseconds.
	The purpose of this function is to provide a reasonably accurate time delay of very short duration. It should not be used for any delays that are much greater than two system clock ticks in length. For delays of a full clock tick, or more, the use of <b>taskDelay()</b> is recommended.
	This routine should be implemented as interrupt safe.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	N/A.
ERRNO	Not Available
SEE ALSO	sysLib, and BSP-specific reference pages for this routine.

## sysNetMacNVRamAddrGet()

NAME sysNetMacNVRamAddrGet() – get network MAC address from NVRAM

SYNOPSIS STATUS sysNetMacNVRamAddrGet

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VxWorks Kernel API Reference, 6.6 sysNvRamGet()

( char \* ifName, int ifUnit, UINT8 \* ifMacAddr, ifMacAddrLen int ) This routine gets the current MAC address from the DESCRIPTION Non Volatile RAM, and store it in the ifMacAddr buffer provided by the caller. It is not required for the BSP to provide NVRAM to store the MAC address. Also, some interfaces do not allow the MAC address to be set by software. In either of these cases, this routine simply returns ERROR. Given a MAC address m0:m1:m2:m3:m4:m5, the byte order of ifMacAddr is: m0@ifMacAddr m1@ifMacAddr + 1 m2@ifMacAddr+2 m3@ifMacAddr + 3 m4@ifMacAddr+4 m5@ifMacAddr+5 OK, if MAC address available, ERROR otherwise RETURNS ERRNO Not Available SEE ALSO vxbNonVolLib

### sysNvRamGet()

DESCRIPTION	This routine copies the contents of non-volatile memory into a specified string. The string will be terminated with an EOS.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	OK, or ERROR if access is outside the non-volatile RAM address range.
ERRNO	Not Available
SEE ALSO	sysLib, sysNvRamSet(), and BSP-specific reference pages for this routine.

# sysNvRamSet()

NAME	sysNvRamSet() – write to non-volatile RAM
SYNOPSIS	<pre>STATUS sysNvRamSet   (     char * string, /* string to be copied into non-volatile RAM */     int strLen, /* maximum number of bytes to copy */     int offset /* byte offset into non-volatile RAM */   )</pre>
DESCRIPTION	This routine copies a specified string into non-volatile RAM.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	OK, or ERROR if access is outside the non-volatile RAM address range.
ERRNO	Not Available
SEE ALSO	sysLib, sysNvRamGet(), and BSP-specific reference pages for this routine.

# sysPhysMemTop()

NAME	<pre>sysPhysMemTop() - get the address of the top of memory</pre>
SYNOPSIS	char * sysPhysMemTop (void)
DESCRIPTION	This routine returns the address of the first missing byte of memory, which indicates the top of memory.
	Normally, the amount of physical memory is specified with the macro LOCAL_MEM_SIZE. BSPs that support run-time memory sizing do so only if the macro LOCAL_MEM_AUTOSIZE is defined. If not defined, then LOCAL_MEM_SIZE is assumed to be, and must be, the true size of physical memory.
NOTE	Do no adjust LOCAL_MEM_SIZE to reserve memory for application use. See <b>sysMemTop()</b> for more information on reserving memory.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	The address of the top of physical memory.
ERRNO	Not Available
SEE ALSO	sysLib, sysMemTop(), and BSP-specific reference pages for this routine.

# sysProcNumGet()

NAME	sysProcNumGet() – get the processor number
SYNOPSIS	int sysProcNumGet (void)
DESCRIPTION	This routine returns the processor number for the CPU board, which is set with <b>sysProcNumSet( )</b> .
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	The processor number for the CPU board.

ERRNO Not Available

**SEE ALSO** sysLib, sysProcNumSet(), and BSP-specific reference pages for this routine.

## sysProcNumSet()

NAME	<pre>sysProcNumSet() - set the processor number</pre>
SYNOPSIS	<pre>void sysProcNumSet   (    int procNum /* processor number */   )</pre>
DESCRIPTION	This routine sets the processor number for the CPU board. Processor numbers should be unique on a single backplane.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	sysLib, sysProcNumGet(), and BSP-specific reference pages for this routine.

## sysScsiBusReset()

NAME	<b>sysScsiBusReset()</b> – assert the RST line on the SCSI bus (Western Digital WD33C93 only)
SYNOPSIS	<pre>void sysScsiBusReset   (    FAST WD_33C93_SCSI_CTRL * pSbic /* ptr to SBIC info */ )</pre>
DESCRIPTION	This routine asserts the RST line on the SCSI bus, which causes all connected devices to return to a quiescent state.

VxWorks Kernel API Reference, 6.6 sysScsiConfig()

 NOTE
 This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.

 RETURNS
 N/A

 ERRNO
 Not Available

 SEE ALSO
 sysLib, and BSP-specific reference pages for this routine.

## sysScsiConfig()

NAME sysScsiConfig() – system SCSI configuration

SYNOPSIS STATUS sysScsiConfig (void)

**DESCRIPTION** This is an example SCSI configuration routine.

Most of the code found here is an example of how to declare a SCSI peripheral configuration. You must edit this routine to reflect the actual configuration of your SCSI bus. This example can also be found in **src/config/usrScsi.c**.

If you are just getting started, you can test your hardware configuration by defining **SCSI\_AUTO\_CONFIG**, which will probe the bus and display all devices found. No device should have the same SCSI bus ID as your VxWorks SCSI port (default = 7), or the same as any other device. Check for proper bus termination.

There are two configuration examples here. They demonstrate configuration of a SCSI hard disk (any type) and an OMTI 3500 floppy disk.

Hard Disk The hard disk is divided into two 32-Mbyte partitions and a third partition with the remainder of the disk.

It is recommended that the first partition (**BLK\_DEV**) on a block device be a dosFs device, if the intention is eventually to boot VxWorks from the device. This will simplify the task considerably.

### Floppy Disk

The floppy, since it is a removable medium device, is allowed to have only a single partition.

In contrast to the hard disk configuration, the floppy setup in this example is more intricate. Note that the **scsiPhysDevCreate()** call is issued twice. The first time is merely to get a "handle" to pass to **scsiModeSelect()**, since the default media type is sometimes inappropriate (in the case of generic SCSI-to-floppy cards). After the hardware is correctly configured, the handle is discarded via **scsiPhysDevDelete()**, after which the peripheral is

	correctly configured by a second call to <b>scsiPhysDevCreate()</b> . (Before the <b>scsiModeSelect()</b> call, the configuration information was incorrect.) Note that after the <b>scsiBlkDevCreate()</b> call, the correct values for <i>sectorsPerTrack</i> and <i>nHeads</i> must be set via <b>scsiBlkDevInit()</b> . This is necessary for IBM PC compatibility.
NOTE	The variable <b>pSbdFloppy</b> is global to allow the above calls to be made from the VxWorks shell, for example:
	-> dosFsMkfs "/fd0", pSbdFloppy
	If a disk is new, use <b>diskFormat( )</b> to format it.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	OK or ERROR.
ERRNO	Not Available
SEE ALSO	sysLib, and BSP-specific reference pages for this routine.

# sysScsiInit()

NAME	sysScsiInit() – initialize an on-board SCSI port
SYNOPSIS	STATUS sysScsiInit (void)
DESCRIPTION	This routine creates and initializes a SCSI control structure, enabling use of the on-board SCSI port. It also connects the proper interrupt service routine to the desired vector, and enables the interrupt at the desired level.
	If SCSI DMA is supported by the board and <b>INCLUDE_SCSI_DMA</b> is defined, the DMA is also initialized.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the control structure cannot be connected, the controller cannot be initialized, or the DMA's interrupt cannot be connected.
ERRNO	Not Available

## sysSerialChanGet()

NAME	sysSerialChanGet() – get the SIO_CHAN device associated with a serial channel
SYNOPSIS	<pre>SIO_CHAN * sysSerialChanGet   (    int channel /* serial channel */ )</pre>
DESCRIPTION	This routine gets the SIO_CHAN device associated with a specified serial channel.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	A pointer to the <b>SIO_CHAN</b> structure for the channel, or <b>ERROR</b> if the channel is invalid.
ERRNO	Not Available
SEE ALSO	sysLib, and BSP-specific reference pages for this routine.

# sysSerialHwInit()

NAME	<b>sysSerialHwInit()</b> – initialize the BSP serial devices to a quiesent state
SYNOPSIS	void sysSerialHwInit (void)
DESCRIPTION	This routine initializes the BSP serial device descriptors and puts the devices in a quiesent state. It is called from <b>sysHwInit()</b> with interrupts locked.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	N/A
ERRNO	Not Available

## sysSerialHwInit2()

NAME	sysSerialHwInit2() – connect BSP serial device interrupts
SYNOPSIS	void sysSerialHwInit2 (void)
DESCRIPTION	This routine connects the BSP serial device interrupts. It is called from <b>sysHwInit2()</b> . Serial device interrupts could not be connected in <b>sysSerialHwInit()</b> because the kernel memory allocator was not initialized at that point, and <b>intConnect()</b> calls <b>malloc()</b> .
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	sysLib, and BSP-specific reference pages for this routine.

## sysSerialReset()

NAME	<b>sysSerialReset( )</b> – reset all SIO devices to a quiet state
SYNOPSIS	void sysSerialReset (void)
DESCRIPTION	This routine is called from <b>sysToMonitor()</b> to reset all SIO device and prevent them from generating interrupts or performing DMA cycles.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	N/A
ERRNO	Not Available

## sysToMonitor()

NAME	sysToMonitor() – transfer control to the ROM monitor
SYNOPSIS	STATUS sysToMonitor ( int startType /* parameter passed to ROM to tell it how to boot */ )
DESCRIPTION	This routine transfers control to the ROM monitor. Normally, it is called only by <b>reboot()</b> which services ^Xand by bus errors at interrupt level. However, in some circumstances, the user may wish to introduce a <i>startType</i> to enable special boot ROM facilities.
NOTE	This is a generic page for a BSP-specific routine; this description contains general information only. To determine if this routine is supported by your BSP, or for information specific to your BSP's version of this routine, see the reference pages for your BSP.
RETURNS	Does not return.
ERRNO	Not Available
SEE ALSO	sysLib, and BSP-specific reference pages for this routine.

## syscallDispatch()

NAME	<b>syscallDispatch( )</b> – dispatch a system call request to its system call handler
SYNOPSIS	STATUS syscallDispatch ( SYSCALL_ENTRY_STATE * pState )
DESCRIPTION	This routine is the system call dispatcher. It decodes a system call number into the group and routine numbers, and locates the proper handler function to call for the given system call request. The system call group must have been previously configured at build time, or registered by calling <b>syscallGroupRegister()</b> .

	Users can hook into the system call dispatch process by attaching entry and exit hook funcitons. These hooks are installed via functions <b>syscallEntryHookAdd()</b> and <b>syscallExitHookAdd()</b> respectively.
	Entry hooks are called after decoding the system call number and verifying that this is a valid system call, but before the handler is called. Entry hooks are passed a pointer to the system call entry state structure as a parameter. Since the state structure contains vital machine information about the trap conditions, it is recommended that its contents not be changed. No explicit barriers prevent such a change from being made. However it is the authors duty to warn you that an ill-considered hook modifying machine state information can crash the system easily. So please do so only if you must, and with the utmost caution. If an entry hook does modify state information, the system call handler function sees the changed values. The state information is architecture-specific. Entry hook functions must return a STATUS value. Any return value other than <b>OK</b> causes the dispatcher to return <b>ERROR</b> back to user code instead of calling the system call handler function. Thus, entry hooks can prevent an otherwise valid system call from being executed.
	Exit hooks are called after the system call handler function returns. They are passed the handler's return value as a parameter. Exit hooks are not expected to return anything.
RETURNS	The return value from the system call handler called, or <b>ERROR</b> .
ERRNO	ENOSYS invalid system call request, or no handler function set.
SEE ALSO	syscallLib, the VxWorks programmer guides.

# syscallEntryHookAdd()

NAME	<pre>syscallEntryHookAdd() - add a routine to be called on each system call entry</pre>			
SYNOPSIS	STATUS syscallEntryHookAdd			
	SYSCALL_ENTRY_HOOK hook, /* routine to call upon system call entry			
	*/ BOOL addToHead /* add routine to head of list? */ )			
DESCRIPTION	This routine adds a specified routine to a list of routines that will be called when a system call is made. The hook routine should have the following prototype:			
	STATUS syscallEntryHook ( SYSCALL_ENTRY_STATE * pState /* system call entry state */			

The second parameter *addToHead* specifies the order in which the hook is added to the table. If **FALSE**, the hook is appended to the list of hooks already installed. If addToHead is **TRUE**, the new hook is added to the head of the list (in other words, it will be the first hook to execute).

System call entry hooks are called from the system call dispatcher after the system call is decoded, but before the handler function is called. Hook functions should return either **OK** or **ERROR**. If the return value from any hook is anything other than **OK**, the system call is aborted and **ERROR** is returned back to the user. Entry hooks can be used to implement rudimentary authentication schemes by rejecting otherwise valid system calls.

RETURNS	OK, or ERROR if the table of hook routines table is full.
---------	---

ERRNO N/A.

SEE ALSO syscallHookLib, syscallEntryHookDelete()

### syscallEntryHookDelete()

NAME	<pre>syscallEntryHookDelete() – delete a previously added entry hook</pre>		
SYNOPSIS	STATUS syscallEntryHookDelete ( SYSCALL_ENTRY_HOOK hook /* routine to be deleted from list */ )		
DESCRIPTION	This routine removes a specified hook routine from the list of system call entry hook routines.		
RETURNS	OK on success, or ERROR if the hook routine was not found.		
ERRNO	S_hookLib_HOOK_NOT_FOUND		
SEE ALSO	syscallHookLib, syscallPreCreateHookAdd()		

### syscallExitHookAdd()

NAME syscallExitHookAdd() – add a routine to be called on each system call exit

SYNOPSIS STATUS syscallExitHookAdd

```
(

SYSCALL_EXIT_HOOK hook, /* routine to call upon system call exit */

BOOL addToHead /* add routine to head of list? */

)
```

**DESCRIPTION** This routine adds a specified routine to a list of routines that will be called when a system call is about to return back to the user. The hook routine should have the following prototype:

```
void syscallExitHook
  (
    int returnValue /* system call return value */
  )
```

The second parameter *addToHead* specifies the order in which the hook is added to the table. If **FALSE**, the hook is appended to the list of hooks already installed. If addToHead is **TRUE**, the new hook is added to the head of the list (in other words, it will be the first hook to execute).

System call exit hooks are called from the system call dispatcher before the system call exits back to the user. Exit hooks are not expected to return anything (return values are not checked).

- **RETURNS** OK, or ERROR if the table of hook routines table is full.
- ERRNO N/A.
- SEE ALSO syscallHookLib, syscallExitHookDelete()

### syscallExitHookDelete()

NAME	<pre>syscallExitHookDelete() – delete a previously added exit hook</pre>		
SYNOPSIS	STATUS syscallExitHookDelete ( SYSCALL_EXIT_HOOK hook /* routine to be deleted from list */ )		
DESCRIPTION	This routine removes a specified hook routine from the list of system call exit hook routines.		
RETURNS	OK on success, or ERROR if the hook routine was not found.		
ERRNO	S_hookLib_HOOK_NOT_FOUND		
SEE ALSO	syscallHookLib, syscallExitHookAdd( )		

	syscallGroupRegister()		
NAME	<b>syscallGroupRegister()</b> – register a system call group with the SCI		
SYNOPSIS	<pre>STATUS syscallGroupRegister   (     int groupNum,     char * groupName,     int numRoutines,     SYSCALL_RTN_TBL_ENTRY * pRoutineTbl,     BOOL force /* forcibly overwrite an existing */     /* entry helps when debugging */   )</pre>		
DESCRIPTION	<ul> <li>This routine registers a system call group with the System Call Infrastructure. Registration is a must, without which user-level code cannot make any system calls to the group in question.</li> <li>Users can hook into the registration process by adding a registration hook function using syscallRegisterHookAdd(). Any attached hooks functions will be called prior to actually registering the group in question. The hook functions are passed the same parameters as are passed to this function with the exception of the force parameter. Hook functions are expected to return a STATUS value. Any return value other than OK causes this function to return ERROR instead of performing the actual registration. Thus, the registration hook can prevent an otherwise valid registration operation from proceeding. This function performs parameter validation prior to calling the registration hooks if any.</li> </ul>		
RETURNS	OK on success, ERROR otherwise.		
ERRNO	S_syscallLib_UNKNOWN_GROUP S_syscallLib_GROUP_EXISTS S_syscallLib_TOO_MANY_ROUTINES S_syscallLib_NO_ROUTINES_TBL		
SEE ALSO	<b>syscallLib</b> , <b>syscallRegisterHookAdd()</b> , <b>syscallRegisterHookDelete()</b> , the VxWorks programmer guides.		

NAME	<b>syscallHookShow()</b> – display all installed system call infrastructure hooks		
SYNOPSIS	void syscallHookShow (void)		
DESCRIPTION	This routine displays the contents of all three system call infrastructure hook tables - the entry, exit and registration hook tables.		
EXAMPLE	The following example shows hypothetical system call hook table contents:		
	-> syscallHookShow		
	System Call Entry Hook Table:		
	entryHook1 entryHook2 entryHook3		
	System Call Exit Hook Table:		
	exitHook		
	System Call Registration Hook Table:		
	registrationHookA registrationHookB value = 1 = 0x1 ->		
RETURNS	N/A		
ERRNOS	N/A		
SEE ALSO	syscallShow, rtpHookShow(), hookShow()		

# syscallMonitor()

NAME	<b>syscallMonitor()</b> – monitor system call activity		
SYNOPSIS	void syscallMonitor		
	int level,		
	RTP_ID rtpId		
	)		

VxWorks Kernel API Reference, 6.6 syscallRegisterHookAdd()

DESCRIPTION	This routine enables/disables system call monitoring activity. It behaves a little like the BSD ktrace or Solaris truss utilities. Users can monitor either a single or all RTPs in the system. When monitoring is turned on, all system calls made by the target RTP are displayed on the console with their argument and return values. This helps monitor all system calls made by applications.		
	syscallMonitor works via the system call hook facility (i.e. <b>syscallHookLib</b> ) It preserves other hooks previously installed. Enabling this facility more than once has no effect.		
	The first parameter enables monitoring (level = 1), or disables it (level = 0). The second parameter is the <b>RTP_ID</b> of the RTP to monitor system calls for. If it is 0, all RTPs are monitored.		
	In SMP version of VxWorks, all actions regarding system calls are guaranteed to be logged only once syscallMonitor returns.		
RETURNS	N/A.		
ERRNO	None.		
SEE ALSO	syscallShow, the VxWorks programmer guides.		

### syscallRegisterHookAdd()

NAME syscallRegisterHookAdd() - add hook for system call group registration requests
SYNOPSIS STATUS syscallRegisterHookAdd
(
SYSCALL\_REGISTER\_HOOK hook, /\* routine to call upon group
registration \*/
BOOL addToHead /\* add routine to head of list \*/

**DESCRIPTION** This routine adds a specified routine to a list of routines that will be called just before an system call group is registered. The hook routine should have the following prototype:

STATUS syscallRegistrationHook

(
 for the second seco

The second parameter *addToHead* specifies the order in which the hook is added to the table. If **FALSE**, the hook is appended to the list of hooks already installed. If addToHead is **TRUE**,

)

	the new hook is added to the head of the list (in other words, it will be the first hook to execute).		
	System call registration hooks are called from <b>syscallGroupRegister()</b> before the registration is actually done. Each hook function should return either <b>OK</b> or <b>ERROR</b> . If the return value from any hook is anything other than <b>OK</b> , registration is aborted and <b>ERROR</b> is returned from <b>syscallGroupRegister()</b> . Registration hooks can be used to implement rudimentary authentication schemes by rejecting otherwise valid group registration requests.		
RETURNS	OK, or ERROR if the table of hook routines table is full.		
ERRNO	N/A.		
SEE ALSO	syscallHookLib, syscallRegisterHookDelete()		

## syscallRegisterHookDelete()

NAME	syscallRegisterHookDelete() – delete a previously added registration hook.		
SYNOPSIS	STATUS syscallRegisterHookDelete ( SYSCALL_REGISTER_HOOK hook /* routine to be deleted from list */ )		
DESCRIPTION	This routine removes a specified hook routine from the list of system call registration hook routines.		
RETURNS	OK on success, or ERROR if the hook routine was not found.		
ERRNO	S_hookLib_HOOK_NOT_FOUND		
SEE ALSO	syscallHookLib, syscallRegisterHookAdd( )		

## syscallShow()

NAME syscallShow() – show registered System Call Groups, or a specific group

SYNOPSIS void syscallShow

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( int grp, int level )

#### DESCRIPTION

This routine shows the registered System Call Groups, as well as details of routines exported by a specific group. When the level parameter is 0, the *grp* parameter is ignored, and all registered groups are shown. The level 0 output (equivalent to **syscallShow** (0,0)) on a target shell looks like this:

-> syscallShow			
Group Name	GroupNo	NumRtns	Rtn Tbl Addr
STANDARDGroup	8	49	0x00274bc0
VXWORKSGroup	9	51	0x00274ed0
value = $53 = 0x35 = '!$	5'		
->			

When the level parameter is 1, details of the System Call Group *grp* are shown. The level 1 output for one the above group looks like this:

-> syscallShow 8,1 System Call Group name: STANDARDGroup Group Number : 8

Routines provided :			
Rtn#	Name	Address	# Arguments
0		0v001f88fc	1
1	creat	0x001e631c	2
2	open	0x001e62b8	3
3	close	0x001e643c	1
4	read	0x001e6444	3
5	write	0x001e63d8	3
6	ioctl	0x001e64a8	3
7	dup	0x001e66d0	1
8	dup2	0x001e66d8	2
9	pipe	0x001e66f0	1
10	remove	0x001e637c	1
11	select	0x001e64c4	5
12	socket	0x001aa5e4	3
13	bind	0x001aa6e4	3
14	listen	0x001aa7c4	2
15	accept	0x001aa85c	3
16	connect	0x001aa9a0	3
17	sendto	0x001aac38	6
18	send	0x001aad58	4
19	sendmsg	0x001ab718	3
20	recvfrom	0x001aae34	6
21	recv	0x001aaf7c	4
22	recvmsg	0x001ab5c8	3
23	setsockopt	0x001ab064	5
24	getsockopt	0x001ab164	5
25	getsockname	0x001ab278	3
26	getpeername	0x001ab378	3

### 2 Routines sysctl()

	27	shutdown	0x001ab478	2
	28	mmap	0x001ed47c	8
	29	munmap	0x001ed4a4	2
	30	mprotect	0x001ed4bc	3
	31	kill	0x001f8c54	2
	32	pause	0x001f8bdc	0
	33	sigpending	0x001f8bcc	1
	34	sigprocmask	0x001f8adc	3
	35	_sigqueue	0x001f8cb0	3
	36	sigsuspend	0x001f8bd4	1
	37	sigtimedwait	0x001f8d34	3
	38	_sigaction	0x001f8a2c	4
	39	_sigreturn	0x001f8d10	0
	40	chdir	0x001e65f4	1
	41	_getcwd	0x001e6650	2
	42	symlink	0x001e6584	2
	43	getpid	0x001f8e64	0
	44	getppid	0x001f8e88	0
	45	waitpid	0x001f8dd8	3
	46	sysctl	0x00164a34	6
	47	_schedPxInfoGet	0x001dab00	2
	48	sigaltstack	0x001f8af8	2
	value	= 50 = 0x32 = '2'		
	->			
RETURNS	N/A.			
ERRNO	None.			
SEE ALSO	syscallShow, the VxWorks programmer guide.			

## sysctl()

NAME	<b>sysctl()</b> – get or set the the values of objects in the sysctl tree
SYNOPSIS	<pre>int sysctl     (     int * pName, /* Name vector of object in MIB style */     u_int nameLen, /* Number of elements in the name vector */     void * pOld, /* Buffer to place the current value of object */     size_t * pOldLen, /* Buffer for the size of current value */     void * pNew, /* Buffer containing a value to set, if needed */     size_t newLen /* Size of the buffer containing new value */     )</pre>

DESCRIPTION This routine retrieves system state information and allows the setting of system information, provided that they have appropriate privileges. The information that sysctl returns will be either an integer, string or table. The state description, hold by the pName

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parameter, is in a MIB or Management Information Base style: a vector of integers. The number of elements in the name vector is specified via the *nameLen* parameter.

The information is copied into the buffer specified by *pOld*. The size of the buffer is given by the location specified by *pOldLen* before the call, and that location gives the amount of data copied after a successful call and after a call that returns with the error code **ENOMEM**. If the amount of data available is greater than the size of the buffer supplied, the call supplies as much data as fits in the buffer provided and returns with the error code **ENOMEM**. If the old value is not desired, *pOld* and *pOldLen* should be set to **NULL**.

The size of the available data can be determined by calling **sysctl()** with a **NULL** parameter for *pOld*. The size of the available data will be returned in the location pointed to by *pOldLen*. For some operations, the amount of space may change often. For these operations, the system attempts to round up so that the returned size is large enough for a call to return the data shortly thereafter.

To set a new value, *pNew* is set to point to a buffer of length *newLen*.

If a new value doesn't need to be set, *pNew* should be set to **NULL**, and *newLen* should be set to 0.

The name vector's elements correspond to a hierarchy of integer values which description can be found in **sys/sysctl.h**. The top level names start with the **CTL**\_ prefix, for instance **CTL\_KERN**. The second level names start with a prefix referring to the top level name they are related to, for instance **KERN\_OSTYPE**, etc.

For instance to get the name of the CPU family on a target board:

**RETURNS** 0 on success, or -1 if an error occurred.

**ERRNO** Beside the errnos possibly set by the OID's non-default handler the following errnos may be set:

ME elements, or the OID ew buffer is too small.	

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### EPERM

An attempt is made to set a read-only value.

### EINVAL

The name vector has less than two or more than **CTL\_MAXNAME** elements, or the OID is not a node and has no handler, or the *newLen* size of the *pNew* buffer is too small.

#### ENOMEM

the *pOldLen* size of the *pOld* buffer is too small for the requested information to be stored in this buffer.

### ENOENT

The OID does not exist.

### EISDIR

The OID is a node without a handler so no information can be set or retreived.

### ENOTDIR

One of the OID numbers in the name vector, except for the last element, does not correspond to a node OID.

### SEE ALSO kern\_sysctl, sysctlbyname(), sysctlnametomib()

### sysctl\_add\_oid()

**NAME** sysctl\_add\_oid() – add a parameter into the sysctl tree during run-time

SYNOPSIS struct sysctl\_oid \* sysctl\_add\_oid ( struct sysctl\_ctx\_list \* clist, struct sysctl\_oid\_list \* parent, int number, const char \* name, int kind, void \* arg1, int arg2, int (\*handler)(SYSCTL\_HANDLER\_ARGS), const char \* fmt, const char \* descr )

**DESCRIPTION** This routine allows the dynamic addition of a parameter that needs to be accessed via sysctl. To use this API, the following arguments are needed.

clist

This is always set to NULL since user-defined contexts are currently not supported.

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parent

The node under which this object is to be registered. When connecting to the static node available for user extensions, **usr\_ext**, use the **SYSCTL\_NODE\_CHILDREN(usr\_ext)** macro to get the pointer to this static OID. When connecting to a dynamic node use the **SYSCTL\_CHILDREN()** macro with the name of the pointer of type **struct sysctl\_oid** representing the parent OID.

number

the OID number that will be assigned to this object. It is highly recommended to use the **OID\_AUTO** macro to avoid conflict with already registered OIDs.

name

A user-specified name for this object.

kind

The kind of object this OID represents as well as the access permissions it holds. For instance CTLTYPE\_NODE | CTLFLAG\_RD, CTLTYPE\_STRING | CTLFLAG\_RW, CTLTYPE\_INT | CTLFLAG\_RW, etc. See sys/sysctl.h for the full list of the CTLTYPE\_... and CTLFLAG\_... macros.

arg1

A pointer to any data that the OID should reference, or NULL. See the SYSCTL\_ADD\_... macros below for specific details.

arg2

The size of *arg1* or 0 if *arg1* is **NULL**.

#### handler

A pointer to the function that will handle read and write requests to this OID. A set of standard handlers are provided that support operations on integers (sysctl\_handle\_int()), strings (sysctl\_handle\_string()) and opaque objects (sysctl\_handle\_opaque()). New handlers can be created (see SYSCTL\_ADD\_PROC() below).

fmt

A pointer to a string that specifies the format of this object. The string must hold "N" for nodes, "A" for strings, "I" for integers, "IU" for unsigned integers, "L" for longs, "LU" for unsigned longs and "S, <type>" for structures (see **SYSCTL\_ADD\_STRUCT()** below for details).

descr

An optional description string.

Utility macros are provided to simplify the creation of OIDs. These macros are:

#### SYSCTL\_ADD\_OID()

Creates a raw OID. This is equivalent to calling **sysctl\_add\_oid()**. The parameters are therefore the same as for the **sysctl\_add\_oid()** routine:

 int arg2, int (\*handler)(SYSCTL\_HANDLER\_ARGS), const char \*
fmt, const char \* descr);

### SYSCTL\_ADD\_NODE()

Creates an OID of type **CTLTYPE\_NODE**. Other OIDs can be added to nodes, as children:

The *access* parameter represents a combination of access permission flags (CTLFLAG\_ macros).

#### SYSCTL\_ADD\_STRING()

Creates an OID of type **CTLTYPE\_STRING** that handles a null-terminated character string:

The *arg* parameter is the address of the string variable and *len* is the maximum length of the string referred to by this OID should it be changed later.

### SYSCTL\_ADD\_INT()

Creates an OID of type CTLTYPE\_INT that handles an integer type variable:

The *access* parameter represents a combination of access permission flags (CTLFLAG\_ macros). The *arg* parameter is the address of the integer variable and *len* is its length.

### SYSCTL\_ADD\_UINT()

Creates an OID of type CTLTYPE\_INT and format "IU" that handles an unsigned integer type variable. The parameters are the same as for SYSCTL\_ADD\_INT() except for the *arg* parameter which is of type "unsigned int \*".

### SYSCTL\_ADD\_LONG()

Creates an OID of type **CTLTYPE\_INT** and format "L" that handles a long integer type variable. The parameters are the same as for **SYSCTL\_ADD\_INT()** except for the *arg* parameter which is of type "long \*".

### SYSCTL\_ADD\_ULONG()

Creates an OID of type **CTLTYPE\_INT** and format "LU" that handles an unsigned long integer type variable. The parameters are the same as for **SYSCTL\_ADD\_INT()** except for the *arg* parameter which is of type "unsigned long \*".

### SYSCTL\_ADD\_OPAQUE()

Creates an OID of type **CTLTYPE\_OPAQUE** that handles an unspecified block of data:

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void \* arg, int len, const char \* format, const char \*
descr);

The *access* parameter represents a combination of access permission flags (CTLFLAG\_ macros). The *arg* parameter is the address of the opaque data and *len* is its length.

### SYSCTL\_ADD\_STRUCT()

Creates an OID that handles a structure type variable. The type of the OID will be CTLTYPE\_OPAQUE and its format will indicates the type name "S, <type>":

The *access* parameter represents a combination of access permission flags (CTLFLAG\_ macros). The *arg* parameter is the address of the stucture variable data and *TYPE* is its type name (without the struct keyword) used in the format string.

#### SYSCTL\_ADD\_PROC()

Creates an OID that specifies a handler routine:

The *arg1* and *arg2* parameters are passed to the *handler* routine and may be set to zero.

#### **EXAMPLES** Creation of a node OID:

static struct sysctl\_oid \* myNode = NULL;

#### Alternatively with **SYSCTL\_ADD\_NODE()**:

static struct sysctl\_oid \* myNode = NULL;

Creation of a string OID under this node:

"This is my own string OID");

Note the usage of the *aString* variable. Its address will be stored in the OID. Since this OID has the CTLFLAG\_RW permission this address must be writable.

Alternatively with SYSCTL\_ADD\_STRING():

	#define MAX_STRING_LENGTH 40	
	<pre>static struct sysctl_oid * myString = NULL; static char aString[MAX_STRING_LENGTH];</pre>	
	<pre>strcpy (aString, "Initial string"); myString = SYSCTL_ADD_STRING (NULL, SYSCTL_CHILDREN (myNode), OID_AUTO,</pre>	1
	This node and this string will be shown as follows by the <b>Sysctl</b> command:	
	-> Sysctl "usr_ext" usr_ext.myNode.myString: Initial string	
RETURNS	Pointer to a <b>sysctl_oid</b> structure on success, or <b>NULL</b> if an error occurred	
ERRNO	EEXIST An OID with the same name vector already exists.	
	EINVAL The <i>parent</i> parameter is NULL.	
	ENOMEM Not enough memoty available to create the OID.	
SEE ALSO	kern_sysctl, sysctl_remove_oid( ), sysctl( )	

### sysctl\_remove\_oid()

NAME sysctl\_remove\_oid() - remove dynamically created sysctl trees SYNOPSIS int sysctl\_remove\_oid ( struct sysctl\_oid \* oidp, int del, int recurse )

DESCRIPTION This routine can be used to remove an object from the sysctl tree. It will only remove objects that were registered dynamically.

	oidp
	Pointer to the object that needs to be removed
	<i>del</i> If 0, it just de-registers this object; otherwise it frees up all resources associated with the object, such as memory for the name, and so on.
	<i>recurse</i> If the value is 0, this routine will return <b>ENOTEMPTY</b> for objects that are nodes and have children.
RETURNS	0 If the entry was successfully removed.
	<b>EINVAL</b> If <i>oidp</i> is an invalid pointer or if this object was not registered dynamically.
	<b>ENOTEMPTY</b> If <i>recurse</i> is 0 and the object has children.
ERRNO	N/A
SEE ALSO	kern_sysctl

### sysctlbyname()

NAME	<b>sysctlbyname()</b> – get or set the values of objects in the sysctl tree by name
SYNOPSIS	<pre>int sysctlbyname   (     char * pName, /* Name of object i.e "net.inet.tcp.delacktime" */     void * pOld, /* Pointer to buffer to get the current values */     size_t * pOldLen, /* Buffer to get the size of current values */     void * pNew, /* Buffer containing the new values to be set */     size_t newLen /* Size of buffer containing new values */    )</pre>
DESCRIPTION	This function accepts an ASCII representation of the name and internally looks up the integer name vector. Apart from that, it behaves the same as the standard <b>sysctl()</b> function.
	The information is copied into the buffer specified by <i>pOld</i> . The size of the buffer is given

The information is copied into the buffer specified by *pOld*. The size of the buffer is given by the location specified by *pOldLen* before the call, and that location gives the amount of data copied after a successful call and after a call that returns with the error code **ENOMEM**. If the amount of data available is greater than the size of the buffer supplied, the call supplies as much data as fits in the buffer provided and returns with the error code **ENOMEM**. If the old value is not desired, *pOld* and *pOldLen* should be set to **NULL**.
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To set a new value, *pNew* is set to point to a buffer of length *newLen* from which the requested value is to be taken. If a new value is not to be set, *pNew* should be set to **NULL** and *newLen* set to 0.

For more information, please refer sysctl()

**RETURNS** 0 on success, or -1 if an error occurred

**ERRNO** Beside the errnos possibly set by the OID's non-default handler the following errnos may be set:

### EPERM

An attempt is made to set a read-only value.

### EINVAL

The name vector has less than two or more than **CTL\_MAXNAME** elements, or the OID is not a node and has no handler, or the *newLen* size of the *pNew* buffer is too small.

#### ENOMEM

the *pOldLen* size of the *pOld* buffer is too small for the requested information to be stored in this buffer.

#### ENOENT

The OID does not exist.

### EISDIR

The OID is a node without a handler so no information can be set or retreived.

### ENOTDIR

One of the OID numbers in the name vector, except for the last element, does not correspond to a node OID.

### SEE ALSO kern\_sysctl

### sysctlnametomib()

**NAME** sysctlnametomib() – return the numeric representation of sysctl object

SYNOPSIS int sysctlnametomib ( const char \* name, int \* mibp, size\_t \* sizep

**DESCRIPTION** This function accepts an ASCII representation of an object in *name*, looks up the integer name vector, and returns the numeric representation in the mib array pointed to by *mibp*.

# tanf()

NAME	tanf() – compute a tangent (ANSI)
SYNOPSIS	<pre>float tanf   (    float x /* angle in radians */   )</pre>
DESCRIPTION	This routine returns the tangent of $x$ in single precision. The angle $x$ is expressed in radians.
RETURNS	The single-precision tangent of <i>x</i> .
ERRNO	Not Available
SEE ALSO	mathALib

# tanhf()

NAME	<pre>tanhf() - comp</pre>	ute a hyperbolic	tangent (ANSI)
NAME	<pre>tanhf() - comp</pre>	ute a hyperbolic	tangent (ANSI)

SYNOPSIS float tanhf

```
(
float x /* number whose hyperbolic tangent is required */
)
```

- **DESCRIPTION** This routine returns the hyperbolic tangent of *x* in single precision.
- **RETURNS** The single-precision hyperbolic tangent of *x*.
- ERRNO Not Available
- SEE ALSO mathALib

## tarArchive()

NAME	<b>tarArchive()</b> – archive named file/dir onto tape in tar format
SYNOPSIS	<pre>STATUS tarArchive  (     char * pTape, /* tape device name */     int bfactor, /* requested blocking factor */     BOOL verbose, /* if TRUE print progress info */     char * pName /* file/dir name to archive */ )</pre>
DESCRIPTION	This function creates a UNIX compatible tar formatted archives which contain entire file hierarchies from disk file systems. Files and directories are archived with mode and time information as returned by <b>stat()</b> .
	The <i>tape</i> argument can be any tape drive device name or a name of any file that will be created if necessary, and will contain the archive. If <i>tape</i> is set to "-", standard output will be used. If <i>tape</i> is <b>NULL</b> (unspecified from Shell), the default archive file name stored in global variable <i>TAPE</i> will be used.
	Each <b>write()</b> of the archive file will be exactly <i>bfactor</i> *512 bytes long, hence on tapes in variable mode, this will be the physical block size on the tape. With Fixed Mode tapes this is only a performance matter. If <i>bfactor</i> is 0, or unspecified from Shell, it will be set to the default value of 20.
	The <i>verbose</i> argument is a boolean, if set to 1, will cause informative messages to be printed to standard error whenever an action is taken, otherwise, only errors are reported.
	The <i>name</i> argument is the path of the hierarchy to be archived. if <b>NULL</b> (or unspecified from the Shell), the current directory path "." will be used. This is the path as seen from the target, not from the Tornado host.
	All informative and error message are printed to standard error.

NOTE	Refrain from specifying absolute path names in <i>path</i> , such archives tend to be either difficult to extract or can cause unexpected damage to existing files if such exist under the same absolute name.
	There is no way of specifying a number of hierarchies to dump.
RETURNS	Not Available
ERRNO	Not Available
SEE ALSO	tarLib

## tarExtract()

NAME	<b>tarExtract()</b> – extract all files from a tar formatted tape
SYNOPSIS	<pre>STATUS tarExtract   (     char * pTape, /* tape device name */     int bfactor, /* requested blocking factor */     BOOL verbose /* if TRUE print progress info */   )</pre>
DESCRIPTION	This is a UNIX-tar compatible utility that extracts entire file hierarchies from tar-formatted archive. The files are extracted with their original names and modes. In some cases a file cannot be created on disk, for example if the name is too long for regular DOS file name conventions, in such cases entire files are skipped, and this program will continue with the next file. Directories are created in order to be able to create all files on tape.
	The <i>tape</i> argument may be any tape device name or file name that contains a tar formatted archive. If <i>tape</i> is equal "-", standard input is used. If <i>tape</i> is <b>NULL</b> (or unspecified from Shell) the default archive file name stored in global variable <i>TAPE</i> is used.
	The <i>bfactor</i> dictates the blocking factor the tape was written with. If 0, or unspecified from the shell, a default of 20 is used.
	The <i>verbose</i> argument is a boolean, if set to 1, will cause informative messages to be printed to standard error whenever an action is taken, otherwise, only errors are reported.
	All informative and error message are printed to standard error.
	There is no way to selectively extract tar archives with this utility. It extracts entire archives.
RETURNS	Not Available
ERRNO	Not Available

SEE ALSO tarLib

## tarToc()

NAME	<b>tarToc( )</b> – display all contents of a tar formatted tape
SYNOPSIS	<pre>STATUS tarToc   (     char * tape, /* tape device name */     int bfactor /* requested blocking factor */   )</pre>
DESCRIPTION	This is a UNIX-tar compatible utility that displays entire file hierarchies from tar-formatted media, e.g. tape.
	The <i>tape</i> argument may be any tape device name or file name that contains a tar formatted archive. If <i>tape</i> is equal "-", standard input is used. If <i>tape</i> is <b>NULL</b> (or unspecified from Shell) the default archive file name stored in global variable <i>TAPE</i> is used.
	The <i>bfactor</i> dictates the blocking factor the tape was written with. If 0, or unspecified from Shell, default of 20 is used.
	Archive contents are displayed on standard output, while all informative and eror message are printed to standard error.
RETURNS	Not Available
ERRNO	Not Available
SEE ALSO	tarLib

# taskActivate()

NAME	taskActivate() – activate a task that has been initialized
SYNOPSIS	STATUS taskActivate ( int tid /* task ID of task to activate */ )
DESCRIPTION	This routine activates tasks created by <b>taskInit()</b> . Without activation, a task is ineligible for CPU allocation by the scheduler.

The *tid* (task ID) argument is simply the address of the **WIND\_TCB** for the task (the **taskInit()** *pTcb* argument), cast to an integer:

tid = (int) pTcb;

The **taskSpawn()** routine is built from **taskActivate()** and **taskInit()**. Tasks created by **taskSpawn()** do not require explicit task activation.

### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

RETURNS	<b>OK</b> , or <b>ERROR</b> if the task cannot be activated.
ERRNO	<b>S_objLib_OBJ_ID_ERROR</b> The <i>tid</i> parameter is an invalid task ID.

```
SEE ALSO taskLib, taskInit()
```

### taskClose()

NAME	taskClose() – close a task
SYNOPSIS	STATUS taskClose ( int tid /* task to close */ )
DESCRIPTION	This routine closes a task. It decrements the task's reference counter.
	This routine does not delete a task. <b>taskDelete( )</b> should be called to terminate and delete a task.
	This routine is not ISR callable.
SMP CONSIDERATIO	NS
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	OK, or ERROR if tid is invalid.
ERRNO	<b>S_objLib_OBJ_ID_ERROR</b> Task ID is invalid.

### S\_intLib\_NOT\_ISR\_CALLABLE This routine must not be called from an ISR.

SEE ALSO taskOpen, taskOpen()

### taskCpuAffinityGet()

NAME taskCpuAffinityGet() – get the CPU affinity of a task

**DESCRIPTION** This routine provides the caller with the CPU affinity of task *tid*. This affinity is represented using a CPU set that is copied in the user supplied *pAffinity*. Passing a null task ID causes the affinity of the caller to be provided. If tid has no affinity the resulting pAffinity CPU set contains no CPU index. If tid has an affinity, the resulting pAffinity CPU set is identical to the CPU set that was passed on the last invocation of **taskCpuAffinitySet()** for that task. This behaviour also applies when calling this routine in the uniprocessor version of VxWorks.

This routine must not be called from interrupt level.

The following code example shows how a task can determine if it has an affinity:

```
{
cpuset_t affinity;
if (taskCpuAffinityGet (0, &affinity) == OK)
    {
        if (CPUSET_ISZERO(affinity))
            {
            /* Task has no affinity */
            }
        else
            {
            /* Task has an affinity */
            }
        }
}
```

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

VxWorks Kernel API Reference, 6.6 taskCpuAffinitySet()

RETURNS	<b>OK</b> or <b>ERROR</b> if the task ID is invalid.
ERRNO	S_objLib_OBJ_ID_ERROR S_intLib_NOT_ISR_CALLABLE
SEE ALSO	<pre>taskLib, taskCpuAffinitySet( ), cpuset</pre>

### taskCpuAffinitySet()

taskCpuAffinitySet() – set the CPU affinity of a task NAME SYNOPSIS STATUS taskCpuAffinitySet ( /\* task ID \*/ int tid, cpuset\_t newAffinity /\* new affinity set \*/ ) DESCRIPTION This routine sets the CPU affinity of task *tid* to the CPU specified in *newAffinity*. From that point on the scheduler ensures the task is only executed on the specified CPU. Passing a tid equal to zero causes an affinity to be set for the calling task. Should the tid argument refer to a task presently running on a CPU other than the one listed in the newAffinity argument, this routine causes the task to cease execution and be rescheduled, based on its priority, on the CPU it has an affinity for. Therefore calling this routine can cause a scheduling event to take place. Calling this routine with an empty CPU set as the newAffinity argument effectively removes any affinity for task tid. If the CPU set identifies a CPU index that is not one of the CPUs configured in the system or if the set contains more than one CPU an error is returned. Once a task has an affinity set, all other tasks it creates have the same affinity except for the case where the child task is the init task of an RTP created using the rtpSpawn() API. Calling this routine in the uniprocessor version of VxWorks is permitted but the newAffinity argument must specify that CPU 0 is the one the task has an affinity for. This action has no effect whatsoever on the scheduling of the task. The only visible effect on uniprocessor VxWorks is that a subsequent call to taskCpuAffinityGet() would indicate the task has an affinity to CPU 0. This routine must not be called from interrupt level.

The following sample code illustrates the sequence to set the affinity of a newly created task to CPU 1:

```
STATUS affinitySetExample (void)
  {
   cpuset_t affinity;
   int tid;
```

```
/* Create the task but only activate it after setting its affinity */
tid = taskCreate ("myCpulTask", 100, 0, 5000, printf,
                  (int) "myCpulTask executed on CPU 1 !", 0, 0, 0,
                  0, 0, 0, 0, 0, 0);
if (tid == NULL)
   return (ERROR);
/* Clear the affinity CPU set and set index for CPU 1 */
CPUSET_ZERO (affinity);
CPUSET_SET (affinity, 1);
if (taskCpuAffinitySet (tid, affinity) == ERROR)
    {
    /* Ooops, looks like we're running on a uniprocessor */
    taskDelete (tid);
    return (ERROR);
    }
/* Now let the task run on CPU 1 */
taskActivate (tid);
return (OK);
}
```

The following example shows how a task can remove its affinity to a CPU:

```
{
cpuset_t affinity;
CPUSET_ZERO (affinity);
taskCpuAffinitySet (0, affinity);
}
```

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

RETURNS	<b>OK</b> , or <b>ERROR</b> if the task ID or affinity is invalid.
ERRNO	S_taskLib_ILLEGAL_OPERATION S_objLib_OBJ_ID_ERROR S_intLib_NOT_ISR_CALLABLE
SEE ALSO	<pre>taskLib, taskCpuAffinityGet( ), vxCpuConfiguredGet( ), cpuset</pre>

VxWorks Kernel API Reference, 6.6 taskCpuLock()

## taskCpuLock()

NAME taskCpuLock() – disable local CPU task rescheduling

SYNOPSIS STATUS taskCpuLock (void)

DESCRIPTION This routine disables scheduling on the CPU the calling task is running on. This effectively prevents any other task from running on the local CPU and prevents the calling task from migrating to another CPU until the lock is released by calling taskCpuUnlock(). This could prove useful when used in conjunction with the vxCpuIndexGet() to ensure a CPU index stays valid for a short period of time while a per-CPU object needs to be read or modified. This routine can be recursively called but there is no effect on scheduling other than an equal number of calls to **taskCpuUnlock()** needs to be done to re-enable scheduling on the local CPU. Execution on other CPUs in the SMP system is not affected by this routine. Because of this behaviour this routine is not a suitable task mutual exclusion mechanism unless all tasks participating in the mutual exclusion scenario have a single CPU affinity for the very same CPU. A task that calls a blocking API such as **semTake()** after calling **taskCpuLock()** constitutes an error condition that results in an error to be returned and reported through ED&R regardless of the fact the task may not block at all. When a task is in a task locked state, its priority cannot be modified by a task or ISR running on another CPU nor can it be suspended or stopped by a task or ISR running on another CPU. Calling this routine on the uniprocessor version of VxWorks is equivalent to calling taskLock().

This routine is not callable from interrupt level. This is not enforced and it is the user's responsibility to adhere to this restriction.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

RETURNS	OK, or ERROR
ERRNO	S_taskLib_ILLEGAL_OPERATION S_objLib_OBJ_ID_ERROR
SEE ALSO	taskLib, taskCpuUnlock(), vxCpuIndexGet()

## taskCpuUnlock()

**NAME** taskCpuUnlock() – enable local CPU task rescheduling

SYNOPSIS STATUS taskCpuUnlock (void)

**DESCRIPTION** This routine removes the lock established using **taskCpuLock()**. It re-enables context task switching on the CPU the calling task is running on. Calling this routine on the uniprocessor version of VxWorks is equivalent to calling **taskUnlock()**.

This routine is not callable from interrupt level. This is not enforced and it is the user's responsibility to adhere to this restriction.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

RETURNS OK, or ERROR

ERRNO S\_taskLib\_ILLEGAL\_OPERATION S\_objLib\_OBJ\_ID\_ERROR

SEE ALSO taskLib, taskCpuLock()

### taskCreate()

**NAME** taskCreate() – allocate and initialize a task without activation

SYNOPSIS int taskCreate

(				
char *	name,	/*	name of new task (stored at pStackBase)	*/
int	priority,	/*	priority of new task */	
int	options,	/*	task option word */	
int	stackSize,	/*	size (bytes) of stack needed */	
FUNCPTR	entryPt,	/*	entry point of new task */	
int	arg1,	/*	1st of 10 req'd args to pass to entryPt	*/
int	arg2,			
int	arg3,			
int	arg4,			
int	arg5,			
int	arg6,			
int	arg7,			
int	arg8,			

int	arg9,
int	arg10
)	

**DESCRIPTION** This routine creates, but does not activate, a new task with a specified priority and options and returns a system-assigned ID. Activate the newly created task by invoking **taskActivate()**.

To create and activate a new task, use the taskSpawn() routine instead of taskCreate().

See **taskSpawn()** for an explanation of all the parameters.

### SMP CONSIDERATIONS

	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	Task ID, or NULL if out of memory or unable to create task.
ERRNO	<b>S_memLib_NOT_ENOUGH_MEMORY</b> There is not enough memory to spawn the task.
	<b>S_taskLib_ILLEGAL_PRIORITY</b> A priority outside the range 0 to 255 was specified.
	S_taskLib_ILLEGAL_OPTIONS The following illegal options were set: VX_DEALLOC_STACK, VX_DEALLOC_EXC_STACK, or VX_DEALLOC_TCB for taskCreate().
	S_taskLib_ILLEGAL_STACK_INFO An invalid stack size has been specified
SEE ALSO	taskLib, taskSpawn(), taskCreat(), taskActivate(), The VxWorks Programmer's Guide

### taskCreateHookAdd()

NAME	<pre>taskCreateHookAdd() – add a routine to be called at every task create</pre>
SYNOPSIS	STATUS taskCreateHookAdd ( FUNCPTR createHook /* routine to be called when a task is created */ )
DESCRIPTION	This routine adds a specified routine to a list of routines that will be called whenever a task is created. Upon creation, all routines specified by <b>taskCreateHookAdd()</b> will be called in

the context of the creating task, so any objects created by a task create hook will be owned

by the caller's RTP rather than the newly created task's RTP. To set the ownership of newly created objects to the new task's RTP, **objOwnerSet()** should be used, e.g.

objOwnerSet (createdObjId, pNewTcb->rtpId)

The routine should be declared as follows:

**RETURNS** OK, or ERROR if the table of task create routines is full.

ERRNO Not Available

SEE ALSO taskHookLib, taskCreateHookDelete()

### taskCreateHookDelete()

NAME	taskCreateHookDelete() – delete a previously added task create routine
SYNOPSIS	STATUS taskCreateHookDelete ( FUNCPTR createHook /* routine to be deleted from list */ )
DESCRIPTION	This routine removes a specified routine from the list of routines to be called at each task create.
RETURNS	OK, or ERROR if the routine is not in the table of task create routines.
ERRNO	Not Available
SEE ALSO	taskHookLib, taskCreateHookAdd()

### taskCreateHookShow()

NAME taskCreateHookShow() – show the list of task create routines

SYNOPSIS void taskCreateHookShow (void)

VxWorks Kernel API Reference, 6.6 taskDelay()

 DESCRIPTION
 This routine shows all the task create routines installed in the task create hook table, in the order in which they were installed.

 RETURNS
 N/A

ERRNO Not Available

SEE ALSO taskHookShow, taskCreateHookAdd()

## taskDelay()

NAME	taskDelay() – delay a task from executing
SYNOPSIS	STATUS taskDelay ( int ticks /* number of ticks to delay task */ )
DESCRIPTION	This routine causes the calling task to relinquish the CPU for the duration specified (in ticks). This is commonly referred to as manual rescheduling, but it is also useful when waiting for some external condition that does not have an interrupt associated with it.
	If the calling task receives a signal that is not being blocked or ignored, <b>taskDelay()</b> returns <b>ERROR</b> and sets <b>errno</b> to <b>EINTR</b> after the signal handler is run.
SMP CONSIDERATIO	NS
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	<b>OK</b> , or <b>ERROR</b> if called from interrupt level or if the calling task receives a signal that is not blocked or ignored.
ERRNO	S_intLib_NOT_ISR_CALLABLE
SEE ALSO	taskLib

## taskDelete()

NAME	taskDelete() – delete a task
SYNOPSIS	STATUS taskDelete ( int tid /* task ID of task to delete */ )
DESCRIPTION	This routine causes a specified task to cease to exist and deallocates the stack and WIND_TCB memory resources. Upon deletion, all routines specified by taskDeleteHookAdd() will be called in the context of the deleting task. This routine is the companion routine to taskSpawn().
WARNING	Deleting individual user tasks, as opposed to deleting the entire RTP, may result in unpredictable RTP behavior. The deletion of individual user tasks should only be performed for debugging purposes.
SMP CONSIDERATIO	NS
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the task cannot be deleted.
ERRNO	S_intLib_NOT_ISR_CALLABLE
	S_objLib_OBJ_DELETED
	S_objLib_OBJ_UNAVAILABLE
	S_objLib_OBJ_ID_ERROR
SEE ALSO	taskLib, excLib, taskDeleteHookAdd(), taskSpawn(), the VxWorks programmer's guides

## taskDeleteForce()

NAME	<b>taskDeleteForce()</b> – delete a task without restriction
SYNOPSIS	STATUS taskDeleteForce ( int tid /* task ID of task to delete */ )

VxWorks Kernel API Reference, 6.6 taskDeleteHookAdd()

DESCRIPTION	This routine deletes a task even if the task is protected from deletion. It is similar to <b>taskDelete()</b> . Upon deletion, all routines specified by <b>taskDeleteHookAdd()</b> will be called in the context of the deleting task.
CAVEATS	This routine is intended as a debugging aid, and is generally inappropriate for applications. Disregarding a task's deletion protection could leave the the system in an unstable state or lead to system deadlock.
	The system does not protect against simultaneous <b>taskDeleteForce( )</b> calls. Such a situation could leave the system in an unstable state.
	Deleting individual user tasks, as opposed to deleting the entire RTP, may result in unpredictable RTP behavior.
SMP CONSIDERATIO	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the task cannot be deleted.
ERRNO	S_intLib_NOT_ISR_CALLABLE
	S_objLib_OBJ_DELETED
	S_objLib_OBJ_UNAVAILABLE
	S_objLib_OBJ_ID_ERROR
SEE ALSO	taskLib, taskDeleteHookAdd(), taskDelete()

## taskDeleteHookAdd()

NAME	taskDeleteHookAdd() – add a routine to be called at every task delete
SYNOPSIS	STATUS taskDeleteHookAdd ( FUNCPTR deleteHook /* routine to be called when a task is deleted */ )
DESCRIPTION	This routine adds a specified routine to a list of routines that will be called whenever a task is deleted. Upon deletion, all routines specified by <b>taskDeleteHookAdd()</b> will be called in the context of the deleting task.
	The routine should be declared as follows:
	void deleteHook

```
(
WIND_TCB *pTcb /* pointer to deleted task's WIND_TCB */
)
```

- **RETURNS** OK, or **ERROR** if the table of task delete routines is full.
- ERRNO Not Available
- SEE ALSO taskHookLib, taskDeleteHookDelete()

### taskDeleteHookDelete()

NAME	taskDeleteHookDelete() – delete a previously added task delete routine
SYNOPSIS	STATUS taskDeleteHookDelete ( FUNCPTR deleteHook /* routine to be deleted from list */ )
DESCRIPTION	This routine removes a specified routine from the list of routines to be called at each task delete.
RETURNS	OK, or ERROR if the routine is not in the table of task delete routines.
ERRNO	Not Available
SEE ALSO	taskHookLib, taskDeleteHookAdd()

## taskDeleteHookShow()

- NAME taskDeleteHookShow() show the list of task delete routines
- SYNOPSIS void taskDeleteHookShow (void)
- **DESCRIPTION** This routine shows all the delete routines installed in the task delete hook table, in the order in which they were installed. Note that the delete routines will be run in reverse of the order in which they were installed.

RETURNS N/A

VxWorks Kernel API Reference, 6.6 taskExit()

ERRNO Not Available

SEE ALSO taskHookShow, taskDeleteHookAdd()

## taskExit()

NAME	taskExit() – exit a task
SYNOPSIS	<pre>void taskExit   (    int exitCode</pre>
DESCRIPTION	This routine is called by a task to cease to exist as a task. The <i>exitCode</i> parameter will be stored in the <b>WIND_TCB</b> for possible use by the delete hooks, or post-mortem debugging.
	This function is currently aliased to <b>exit()</b> , and is provided as a convenience to achieve uniform meaning across both kernel and user-mode code.
SMP CONSIDERATIO	NS
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	<b>taskLib</b> , <b>taskDelete()</b> , "American National Standard for Information Systems - "Programming Language - C, ANSI X3.159-1989: Input/Output (stdlib.h),", the VxWorks programmer's guides

## taskHookShowInit()

NAME taskHookShowInit() – initialize the task hook show f	facility
---	----------

SYNOPSIS void taskHookShowInit (void)

 DESCRIPTION
 This routine links the task hook show facility into the VxWorks system. It is called automatically when the task hook show facility is configured into VxWorks using the INCLUDE\_TASK\_HOOK\_SHOW component.

 RETURNS
 N/A

ERRNO Not Available

SEE ALSO taskHookShow

### taskIdDefault()

NAME	taskIdDefault() – set the default task ID	
SYNOPSIS	<pre>int taskIdDefault (     int tid /* user supplied task ID; if 0, return default */ )</pre>	
<b>DESCRIPTION</b> This routine maintains a global default task ID. This ID is used by libraries th allow a task ID argument to take on a default value if the user did not explicitly		
	If <i>tid</i> is not zero (i.e., the user did specify a task ID), the default ID is set to that value, and that value is returned. If <i>tid</i> is zero (i.e., the user did not specify a task ID), the default ID is not changed and its value is returned. Thus the value returned is always the last task ID the user specified.	
RETURNS	The most recent non-zero task ID.	
ERRNOS	N/A	
SEE ALSO	<b>taskInfo</b> , <b>dbgLib</b> , <b>windsh</b> , the VxWorks programmer guides, the, <i>VxWorks Command-Line</i> <i>Tools User's Guide</i> .	

### taskIdListGet()

NAME taskIdListGet() – get a list of active task IDs

SYNOPSIS int taskIdListGet

VxWorks Kernel API Reference, 6.6 taskIdSelf()

( int idList[], /\* array of task IDs to be filled in \*/ int maxTasks /\* max tasks <idList> can accommodate \*/ DESCRIPTION This routine provides the calling task with a list of all active tasks. An unsorted list of task IDs for no more than *maxTasks* tasks is put into *idList*. The provided list is a snapshot of the system. There is no guarantee that this snapshot will still represent the state of the system by the time execution returns to the caller. This is especially true on VxWorks SMP because of the concurrent execution environment. This routine is provided if the INCLUDE\_TASK\_LIST component is present in the configuration. The number of tasks put into the ID list. RETURNS ERRNO N/A SEE ALSO taskInfo, taskNameToId

### taskIdSelf()

NAMEtaskIdSelf() – get the task ID of a running taskSYNOPSISint taskIdSelf (void)DESCRIPTIONThis routine gets the task ID of the calling task. The task ID will be invalid if called at<br/>interrupt level.RETURNSThe task ID of the calling task.ERRNON/ASEE ALSOtaskLib

# taskIdVerify()

NAME	<b>taskIdVerify()</b> – verify the existence of a task	
SYNOPSIS	STATUS taskIdVerify ( int tid /* task ID */ )	
DESCRIPTION	This routine verifies the existence of a specified task by validating the specified ID as a task ID. Note that an exception occurs if the task ID parameter points to an address not located in physical memory.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if the task ID is invalid.	
ERRNO	S_objLib_OBJ_ID_ERROR	
SEE ALSO	taskLib	

# taskInfoGet( )

NAME	<b>taskInfoGet( )</b> – get information about a task		
SYNOPSIS	<pre>STATUS taskInfoGet   (    int tid, /* ID of task for which to get info */   TASK_DESC * pTaskDesc /* task descriptor to be filled in */   )</pre>		
DESCRIPTION	This routine fills in a specified task descriptor (TASK_DESC) for a specified task. The information in the task descriptor is, for the most part, a copy of information kept in the task control block (WIND_TCB). The TASK_DESC structure is useful for common information and avoids dealing directly with the unwieldy WIND_TCB.		
NOTE	Examination of WIND_TCBs should be restricted to debugging aids.		
RETURNS	OK, or ERROR if the task ID is invalid or access to task info denied.		
ERRNO	N/A		
SEE ALSO	taskShow		

### taskInit()

STATUS taskInit

NAME

taskInit() - initialize a task with a stack at a specified address

SYNOPSIS

(		
FAST WIND_TCB	*pTcb,	/* address of new task's TCB */
char *	name,	/* name of new task (stored at pStackBase) */
int	priority,	/* priority of new task */
int	options,	/* task option word */
char *	pStackBase,	/* base of new task's execution stack */
int	stackSize,	/* size (bytes) of stack needed */
FUNCPTR	entryPt,	/* entry point of new task */
int	arg1,	/* 1st of 10 req'd args to pass to entryPt */
int	arg2,	
int	arg3,	
int	arg4,	
int	arg5,	
int	arg6,	
int	arg7,	
int	arg8,	
int	arg9,	
int	arg10	
)		

DESCRIPTION

TION This routine initializes user-specified regions of memory for a task stack and control block instead of allocating them from memory as **taskSpawn()** does. This routine will utilize the specified pointers to the **WIND\_TCB** and stack as the components of the task. This allows, for example, the initialization of a static **WIND\_TCB** variable. It also allows for special stack positioning as a debugging aid.

As in **taskSpawn()**, a task may be given a name. While **taskSpawn()** automatically names unnamed tasks, **taskInit()** permits the existence of tasks without names. The task ID required by other task routines is simply the address *pTcb*, cast to an integer.

Unlike **taskSpawn()**, **taskInit()** allows one to control the activation of the **VX\_DEALLOC\_STACK** bit in *options*; **taskSpawn()** always sets this bit. Setting this bit causes the stack to be automatically deallocated when a **taskDelete()** is performed, or when a return command is issued from the entry function.

It is not recommended to use the VX\_DEALLOC\_STACK option for taskInit() if RTP or KERNEL\_HARDENING is configured into the system. A system configured with RTP or KERNEL\_HARDENING expects the task stacks to have guard zones, governed by the configuration parameters TASK\_KERNEL\_EXEC\_STACK\_UNDERFLOW\_SIZE and TASK\_KERNEL\_EXEC\_STACK\_OVERFLOW\_SIZE. Setting the VX\_DEALLOC\_STACK option but not set the appropriate guard zones in these configuration might cause corruption in the system. If corruption occurs, use the Error Detection and Reporting mechanism to help you detect the corruption.

If it is necessary to use the VX\_DEALLOC\_STACK option in the above configurations, the user is responsible for setting up the guard zones corresponding to the same value of the configuration parameters or the user can specify the VX\_NO\_PROTECT option for the task.

The *pStackBase* parameter specifies the base of the execution stack. The stack may grow up or down from *pStackBase* depending on the target architecture. The caller is responsible for setting up any guard zones around the specified stack area. The following code fragment illustrates how to specify the stack base location:

For architectures where the stack grows down:

```
pStackMem = (char *) memalign (_STACK_ALIGN_SIZE, stackSize);
if (pStackMem != NULL)
    status = taskInit ( ... , pStackMem + stackSize, stackSize, ... );
```

For architectures where the stack grows up:

```
pStackMem = (char *) memalign (_STACK_ALIGN_SIZE, stackSize);
if (pStackMem != NULL)
    status = taskInit ( ... , pStackMem, stackSize, ... );
```

Please note that **memalign()** is used in the above code fragment for illustrative purposes only since it's a well-known API. The stack memory can be obtained by any mechanism that ensures allocation of aligned memory region.

The *stackSize* parameter specifies the size in bytes of the execution stack area. This API does not check against illegal stack size, since it is assumed that the user has allocated the stack memory with a valid stack size, before calling this API.

It is assumed that the caller passes valid pointers for *pTcb*, *pStackBase* and *entryPt* while calling this API. No validity checks for these parameters are done here.

Other arguments are the same as in **taskSpawn()**. Unlike **taskSpawn()**, **taskInit()** does not activate the task. This must be done by calling **taskActivate()** after calling **taskInit()**.

Normally, tasks should be started using **taskSpawn()** rather than **taskInit()**, except when additional control is required for task memory allocation or a separate task activation is desired.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

- NOTE For future releases of VxWorks, each task has an exception stack. To continue providing the taskInit() API, this routine now carves memory for the exception stack of a task . To use a specific region of memory for the exception stack, use the routine taskInitExcStk() instead.
- **RETURNS** OK, or **ERROR** if the task cannot be initialized.

VxWorks Kernel API Reference, 6.6 taskInitExcStk()

ERRNO	S_intLib_NOT_ISR_CALLABLE Routine is not callable from an ISR.
	<b>S_taskLib_ILLEGAL_PRIORITY</b> Priority specified is not within 0-255.
	S_objLib_OBJ_ID_ERROR
SEE ALSO	taskLib, taskActivate( ), taskSpawn( ), taskInitExcStk( )

### taskInitExcStk()

NAME taskInitExcStk() – initialize a task with stacks at specified addresses

SYNOPSIS

STATUS	taskInitExcStk	

	(			
	FAST WIND_TCB	*pTcb,	/*	address of new task's TCB */
	char *	name,	/*	name of new task (stored at pStackBase)
*/				
	int	priority,	/*	priority of new task */
	int	options,	/*	task option word */
	char *	pStackBase,	/*	base of new task's execution stack */
	int	stackSize,	/*	size (bytes) of stack needed */
	char *	pExcStackBase,	/*	base of new task's exception stack */
	int	excStackSize,	/*	size (bytes) of exception stack needed
*/				
	FUNCPTR	entryPt,	/*	entry point of new task */
	int	arg1,	/*	first of ten task args to pass to func
*/				
	int	arg2,		
	int	arg3,		
	int	arg4,		
	int	arg5,		
	int	arg6,		
	int	arg7,		
	int	arg8,		
	int	arg9,		
	int	arg10		
	)			

**DESCRIPTION** This routine initializes user-specified regions of memory for a task stack, exception stack, and control block instead of allocating them from memory as **taskSpawn()** does. This routine will utilize the specified pointers to the **WIND\_TCB** and stacks as the components of the task. This allows, for example, the initialization of a static **WIND\_TCB** variable. It also allows for special stack positioning as a debugging aid.

As in **taskSpawn()**, a task may be given a name. While **taskSpawn()** automatically names unnamed tasks, **taskInitExcStk()** permits the existence of tasks without names. The task ID required by other task routines is simply the address *pTcb*, cast to an integer.

Unlike taskSpawn(), taskInitExcStk() allows one to control the activation of both the VX\_DEALLOC\_STACK and VX\_DEALLOC\_EXC\_STACK bits in *options*; Setting the VX\_DEALLOC\_STACK bit causes the stack (aka execution stack) to be automatically deallocated when a taskDelete() is performed, or when a return command is issued from the entry function. Setting the VX\_DEALLOC\_EXC\_STACK bit causes the exception stack to be automatically deallocated when a taskDelete() is performed, or when a return command is issued from the entry function.

It is not recommended to use the VX\_DEALLOC\_STACK option for taskInit() if RTP or KERNEL\_HARDENING is configured into the system. A system configured with RTP or KERNEL\_HARDENING expects the task stacks to have guard zones, governed by the configuration parameters TASK\_KERNEL\_EXEC\_STACK\_UNDERFLOW\_SIZE and TASK\_KERNEL\_EXEC\_STACK\_OVERFLOW\_SIZE. Setting the VX\_DEALLOC\_STACK option but not set the appropriate guard zones in these configuration might cause corruption in the system. If corruption occurs, use the Error Detection and Reporting mechanism to help you detect the corruption.

If it is necessary to use the VX\_DEALLOC\_STACK option in the above configurations, the user is responsible for setting up the guard zones corresponding to the same value of the configuration parameters or the user can specify the VX\_NO\_PROTECT option for the task.

The *pStackBase* parameter specifies the base of the execution stack, and *pExcStackBase* specified the base of the exception stack. The stacks may grow up or down from *pStackBase/pExcStackBase* depending on the target architecture. The caller is responsible for setting up any guard zones around the specified stack areas. The following code fragment illustrates how to specify the stack base location:

For architectures where the stack grows down:

For architectures where the stack grows up:

Please note that <b>memalign()</b> is used in the above code fragment for illustrative pu only since it's a well-known API. Typically, the stack memory would be obtained b other mechanism.				
Other arguments are the same as in <b>taskSpawn()</b> . Unlike <b>taskSpawn(), taskI</b> r does not activate the task. This must be done by calling <b>taskActivate()</b> after ca <b>taskInitExcStk()</b> .				
	Normally, tasks should be started using <b>taskSpawn()</b> rather than <b>taskInit()</b> or <b>taskInitExcStk()</b> , except when additional control is required for task memory allocation or a separate task activation is desired.			
SMP CONSIDERATIONS				
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.			
RETURNS	<b>OK</b> , or <b>ERROR</b> if the task cannot be initialized.			
ERRNO	S_intLib_NOT_ISR_CALLABLE Routine is not callable from an ISR.			
	S_taskLib_ILLEGAL_PRIORITY Priority specified is not within 0-255.			
	S_objLib_OBJ_ID_ERROR			
SEE ALSO	taskLib, taskActivate( ), taskSpawn( ), taskInit( )			

# taskIsPended()

NAME	taskIsPended() – check if a task is pended	
SYNOPSIS	BOOL taskIsPended ( int tid /* task ID */ )	
DESCRIPTION	This routine tests the status field of a task to determine if it is pended. No indication is given regarding the timeout, if any, associated with the pending operation.	
RETURNS	TRUE if the task is pended, otherwise FALSE.	
ERRNOS	N/A	

SEE ALSO taskInfo

# taskIsReady()

NAME	taskIsReady() – check if a task is ready to run
SYNOPSIS	BOOL taskIsReady ( int tid /* task ID */ )
DESCRIPTION	This routine tests the status field of a task to determine if it is ready to run.
RETURNS	TRUE if the task is ready, otherwise FALSE.
ERRNOS	N/A
SEE ALSO	taskInfo

# taskIsStopped()

NAME	<b>taskIsStopped()</b> – check if a task is stopped by the debugger
SYNOPSIS	BOOL taskIsStopped ( int tid /* task ID */ )
DESCRIPTION	This routine tests the status field of a task to determine if it is stopped by the debugger.
RETURNS	TRUE if the task is stopped by the debugger, otherwise FALSE.
ERRNOS	N/A /NOMANUAL
SEE ALSO	taskInfo

VxWorks Kernel API Reference, 6.6 taskIsSuspended()

## taskIsSuspended()

 NAME
 taskIsSuspended() - check if a task is suspended

 SYNOPSIS
 BOOL taskIsSuspended

 (int tid /\* task ID \*/)
 This routine tests the status field of a task to determine if it is suspended.

 DESCRIPTION
 This routine tests the status field of a task to determine if it is suspended.

 RETURNS
 TRUE if the task is suspended, otherwise FALSE.

 ERRNOS
 N/A

 SEE ALSO
 taskInfo

## taskKill()

NAME	taskKill() – send a signal to a task
SYNOPSIS	<pre>int taskKill   (    int tid,    int signo   )</pre>
DESCRIPTION	This routine sends a signal <i>signo</i> to the task specified by <i>tid</i> . This function is currently aliased to <b>kill()</b> , and is provided as a convenience to achieve uniform meaning across both kernel and user-mode code.
RETURNS	<b>OK</b> (0), or <b>ERROR</b> (-1) if the task ID or signal number is invalid.
ERRNO	EINVAL
SEE ALSO	sigLib, kill()

# taskLock()

NAME	taskLock() – disable task rescheduling
SYNOPSIS	STATUS taskLock ( void )
DESCRIPTION	This routine disables task context switching. The task that calls this routine will be the only task that is allowed to execute, unless the task explicitly gives up the CPU by making itself no longer ready. Typically this call is paired with <b>taskUnlock()</b> ; together they surround a critical section of code. These preemption locks are implemented with a counting variable that allows nested preemption locks. Preemption will not be unlocked until <b>taskUnlock()</b> has been called as many times as <b>taskLock()</b> .
	This routine does not lock out interrupts; use <b>intLock()</b> to lock out interrupts.
	A <b>taskLock()</b> is preferable to <b>intLock()</b> as a means of mutual exclusion, because interrupt lock-outs add interrupt latency to the system.
	A <b>semTake()</b> is preferable to <b>taskLock()</b> as a means of mutual exclusion, because preemption lock-outs add preemptive latency to the system.
WARNINGS	The <b>taskLock()</b> routine is not callable from interrupt service routines.
	Invoking a VxWorks system routine with preemption locked may result in preemption being unlocked for an unspecified period of time. If the called routine blocks or suspends the calling task, the scheduler will always select the highest priority ready task to execute (or become idle) regardless of whether the task has locked preemption via <b>taskLock()</b> .
	The preemption lock state is an attribute of a task, i.e. it's part of the task context. Thus, if a task disables preemption and subsequently invokes a VxWorks system routine that causes the calling task to block or cause a higher priority task to be ready, the preemption lock state will be restored when the task is later rescheduled for execution.
SMP CONSIDERATIO	NS
	This API is not available in VxWorks SMP.
RETURNS	N/A
RETURNS	OK or ERROR.
ERRNO	S_objLib_OBJ_ID_ERROR
	S_intLib_NOT_ISR_CALLABLE
SEE ALSO	taskLib, taskUnlock( ), taskCpuLock( ), intLock( ), taskSafe( ), semTake( )

## taskName()

NAME	<b>taskName( )</b> – get the name associated with a task ID
SYNOPSIS	char * taskName ( int tid /* ID of task whose name is to be found */ )
DESCRIPTION	This routine returns a pointer to the name of a task of a specified ID, if the task has a name. If the task has no name, it returns an empty string.
RETURNS	A pointer to the task name, or NULL if the task ID is invalid.
ERRNOS	N/A
SEE ALSO	taskInfo

# taskNameToId()

NAME	<b>taskNameToId()</b> – look up the task ID associated with a task name
SYNOPSIS	<pre>int taskNameToId   (    char * name /* task name to look up */ )</pre>
DESCRIPTION	This routine returns the ID of the task matching a specified name. Referencing a task in this way is inefficient, since it involves a search of the task list. This routine is provided if the INCLUDE_TASK_LIST component is present in the configuration.
RETURNS	The task ID, or <b>ERROR</b> if the task is not found.
ERRNO	S_taskLib_NAME_NOT_FOUND
SEE ALSO	taskInfo, taskName

### taskOpen()

NAME taskOpen() – open a task

SYNOPSIS	int	taskOpen				
		<pre>(     const char     int     int     char *     int     void     FUNCPTR     int     int</pre>	* *	<pre>name, priority, options, mode, pStackBase, stackSize, context, entryPt, arg1, arg2, arg3, arg4, arg5, arg6, arg7, arg8, arg9, arg10</pre>	/ / / / / / / / / / / / /	<pre>task name - default name will be chosen */ task priority */ VX_ task option bits */ object management mode bits */ base of new task's execution stack */ execution stack size */ context value */ application entry point */ lst of 10 req'd args to pass to entryPt */</pre>

**DESCRIPTION** The **taskOpen()** API is the most general purpose task creation routine. It can also be used to obtain a task ID to an already existing task, typically a public task with an RTP. It searches the task name space for a matching task. If a matching task is found, it returns the task ID of the matched task. If a matching task is not found but the **OM\_CREATE** flag is specified in the *mode* parameter, then it creates a task. This routine is not ISR callable.

There are two name spaces available in which **taskOpen()** can perform the search. The name space searched is dependent upon the first character in the *name* parameter. When this character is a forward slash /, the **public** name space is searched; otherwise the **private** name space is searched. Similarly, if a task is created, the *name*'s first character specifies the name space that contains the task.

Unlike other objects in VxWorks, private task names are not unique. Thus a search on a private name space finds the first matching task. However, this task may not be the only task with the specified name. Public task names on the other hand, are unique.

A description of the taskOpen() arguments follows:

name

This is a mandatory argument. Unlinke taskSpawn(), NULL or empty strings are not allowed when using this routine. The task's name appears in various kernel shell facilities such as i(). The name may be of arbitrary length and content. Public task names are unique, private task names are not.

### priority

The VxWorks kernel schedules tasks on the basis of priority. Tasks may have priorities ranging from 0 (highest) to 255 (lowest). The priority of a task in VxWorks is dynamic, and the priority of an existing task can be changed using **taskPrioritySet()**. Also, a task can inherit a priority as a result of the acquisition of a priority-inversion-safe mutex semaphore.

### options

Bits in the options argument may be set to run with the following modes:

VX_UNBREAKABLE	do not allow breakpoint debugging
VX_FP_TASK	execute with floating-point coprocessor support
VX_ALTIVEC_TASK	execute with Altivec support (PowerPC only)
VX_SPE_TASK	execute with SPE support (PowerPC only)
VX_DSP_TASK	execute with DSP support (SuperH only)
VX_PRIVATE_ENV	the task has a private environment area
VX_NO_STACK_FILL	do not fill the stack with 0xee (for debugging)
VX_TASK_NOACTIVATE	do not activate the task upon creation
VX_NO_STACK_PROTECT	do not provide overflow/underflow stack protection, stack
	remains executable.

mode

This parameter specifies the various object management attribute bits as follows:

#### OM\_CREATE

Create a new task if a matching task name is not found.

#### OM\_EXCL

When set jointly with **OM\_CREATE**, create a new task immediately without attempting to open an existing task. The call fails if the task is public and its name causes a name clash. This flag has no effect if the **OM\_CREATE** attribute is not specified.

### OM\_DELETE\_ON\_LAST\_CLOSE

This bit is ignored on tasks because it would allow a task to be deleted from another RTP.

### pStackBase

Base of the execution stack. When a **NULL** pointer is specified, the kernel allocates a page-aligned stack area.

The stack may grow up or down from *pStackBase* depending on the target architecture. The caller is responsible for setting up any guard zones around the specified stack area. The following code fragment illustrates how to specify the stack base location:

For architectures where the stack grows down:

```
pStackMem = (char *) malloc (stackSize);
if (pStackMem != NULL)
    taskId = taskOpen ( ... , pStackMem + stackSize, stackSize, ... );
```

#### For architectures where the stack grows up:

```
pStackMem = (char *) malloc (stackSize);
if (pStackMem != NULL)
    taskId = taskOpen ( ... , pStackMem, stackSize, ... );
```

Please note that **malloc()** is used in the above code fragment for illustrative purposes only since it's a well-known API. Typically, the stack memory would be obtained by some other mechanism.

It is assumed that if the caller passes a non-NULL pointer as *pStackBase*, it is valid. No validity check for this parameter is done here.

#### stackSize

The size in bytes of the execution stack area. If **NULL** pointer is specified as *pStackBase* and a negative value is specified for this parameter, the API returns **ERROR** considering it an illegal stack size. However, the API does not check against illegal stack size, if a non-**NULL** pointer is specified as *pStackBase*, since it is assumed that the user has allocated the stack memory with a valid stack size, before calling this API.

Every byte of the stack is filled with 0xee (unless the VX\_NO\_STACK\_FILL option is specifed or the global kernel configuration parameter VX\_GLOBAL\_NO\_STACK\_FILL is set to TRUE) for the **checkStack()** kernel shell facility.

#### context

Context value assigned to the created task. This value is not actually used by VxWorks. Instead, the context value can be used by OS extensions to implement object permissions, for example.

#### entryPt

The entry point is the address of the **main** routine of the task. The routine is called once the C environment has been set up. The specified routine is called with the ten arguments *arg1* to *arg10*. Should the specified **main** routine return, a call to **exit()** is automatically made.

It is assumed that the caller passes a valid function pointer as *entryPt*. No validity check for this parameter is done here.

To delete a task created via the **taskOpen()** API, **taskDelete()** must be called. A call to **taskClose()** will not perform the task deletion.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

**RETURNS** The task ID, or **NULL** if unsuccessful.

#### ERRNO S\_memLib\_NOT\_ENOUGH\_MEMORY

There is not enough memory to spawn the task.

VxWorks Kernel API Reference, 6.6 taskOpenInit()

	<b>S_taskLib_ILLEGAL_PRIORITY</b> A priority outside the range 0 to 255 was specified.
	S_taskLib_ILLEGAL_OPERATION The operation attempted to specify a location for the stack (not supported in TAR).
	S_taskLib_ILLEGAL_STACK_INFO An invalid stack size has been specified.
	S_objLib_OBJ_INVALID_ARGUMENT An invalid option was specified in the <i>mode</i> argument or <i>name</i> is invalid.
	S_objLib_OBJ_NOT_FOUND The OM_CREATE flag was not set in the <i>mode</i> argument and a task matching <i>name</i> was not found.
	S_intLib_NOT_ISR_CALLABLE This routine must not be called from an ISR.
SEE ALSO	<b>taskOpen, taskSpawn( ), taskCreate( ), taskActivate( ), taskClose( ),</b> the VxWorks programmer guides.

# taskOpenInit()

NAME	taskOpenInit() – initialize the task open facility
SYNOPSIS	void taskOpenInit (void)
DESCRIPTION	This routine links the task creation routine with the open facility into the VxWorks system. It is called automatically when the task facility is configured into VxWorks by either defining INCLUDE_OBJ_OPEN and INCLUDE_TASK_CREATE_DELETE in config.h or selecting INCLUDE_OBJ_OPEN and INCLUDE_TASK_CREATE_DELETE in the project facility.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	taskOpen

## taskOptionsGet()

NAME	taskOptionsGet() – examine task options						
SYNOPSIS	STATUS taskOptionsGet ( int tid, /* task ID */ int * pOptions /* task's options */ )						
DESCRIPTION	This routine gets the current execution options of the specified task. The option bits returned by this routine indicate the following modes:						
	VX_FP_TASK execute with floating-point coprocessor support.						
	VX_PRIVATE_ENV include private environment support (see <b>envLib</b> ).						
	VX_NO_STACK_FILL do not fill the stack for use by <b>checkstack()</b> .						
	VX_UNBREAKABLE do not allow breakpoint debugging.						
	For definitions, see taskLib.h.						
RETURNS	<b>OK</b> , or <b>ERROR</b> if the task ID is invalid.						
ERRNOS	N/A						
SEE ALSO	taskInfo, taskOptionsSet( )						

## taskOptionsSet()

NAME	taskOptionsSet() – change task options						
SYNOPSIS	<pre>STATUS taskOptionsSet    (     int tid,    /* task ID */     int mask,    /* bit mask of option bits to unset */     int newOptions   /* bit mask of option bits to set */    )</pre>						

**DESCRIPTION** This routine changes the execution options of a task. The only option that can be changed after a task has been created is:

VxWorks Kernel API Reference, 6.6 taskPriNormalGet()

### VX\_UNBREAKABLE do not allow breakpoint debugging.

For definitions, see taskLib.h.

### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

RETURNS	OK, or	ERROR	if the	task	ID	is	invalio	ł.
---------	--------	-------	--------	------	----	----	---------	----

ERRNOS N/A

SEE ALSO taskInfo, taskOptionsGet()

### taskPriNormalGet()

NAME	<b>taskPriNormalGet()</b> – get the normal priority of the task
SYNOPSIS	STATUS taskPriNormalGet ( int tid, /* task ID */ int* pPriNormal /* where to return priority */ )
DESCRIPTION	This routine gets the normal priority of the specified task, which is the priority assigned at task creation time or subsequently assigned using <b>taskPrioritySet()</b> . A task executes at its normal assigned priority unless priority inheritance has occurred.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the task ID is invalid.
ERRNO	N/A
SEE ALSO	taskInfo, taskSpawn( ), taskCreate( ), taskPrioritySet( )
## taskPriorityGet()

NAME	<b>taskPriorityGet( )</b> – examine the priority of a task
SYNOPSIS	<pre>STATUS taskPriorityGet   (    int tid,   /* task ID */    int * pPriority /* return priority here */   )</pre>
DESCRIPTION	This routine determines the current priority of a specified task. The current priority is copied to the integer pointed to by $pPriority$ .
RETURNS	<b>OK</b> , or <b>ERROR</b> if the task ID is invalid.
ERRNO	S_objLib_OBJ_ID_ERROR
SEE ALSO	taskstack remains executablestack remains executable <b>skLib</b> , <b>taskPrioritySet( )</b>

# taskPrioritySet()

NAME	<b>taskPrioritySet( )</b> – change the priority of a task
SYNOPSIS	<pre>STATUS taskPrioritySet   (    int tid,</pre>
DESCRIPTION	This routine changes a task's priority to a specified priority. Priorities range from 0, the highest priority, to 255, the lowest priority.
	A request to lower the priority of a task that has acquired a priority inversion safe mutex semaphore will not take immediate effect. To prevent a priority inversion situation, the requested lower priority will take effect, in general, only after the task relinquishes all priority inversion safe mutex semaphores.
	A request to raise the priority of a task will take immediate effect.
SMP CONSIDERATIO	NS
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

VxWorks Kernel API Reference, 6.6 taskRaise()

RETURNS	<b>OK</b> , or <b>ERROR</b> if the task ID is invalid.
ERRNO	S_taskLib_ILLEGAL_PRIORITY S_objLib_OBJ_ID_ERROR
SEE ALSO	taskLib, taskPriorityGet( )

## taskRaise()

NAME	<b>taskRaise( )</b> – send a signal to the caller's task
SYNOPSIS	int taskRaise ( int signo )
DESCRIPTION	This routine sends the signal <i>signo</i> to the task invoking the call. This function is currently aliased to <b>raise()</b> , and is provided as a convenience to achieve uniform meaning across both kernel and user-mode code.
RETURNS	<b>OK</b> (0), or <b>ERROR</b> (-1) if the signal number or task ID is invalid.
ERRNO	EINVAL
SEE ALSO	sigLib, raise()

## taskRegsGet()

NAME	<b>taskRegsGet( )</b> – get a task's registers from the TCB
SYNOPSIS	STATUS taskRegsGet ( int tid, /* task ID */ REG_SET * pRegs /* put register contents here */ )
DESCRIPTION	This routine gathers task information kept in the TCB. It copies the contents of the task's registers to the register structure $pRegs$ .

NOTE	This routine only works well if the task is known to be in a stable, non-executing state. Self-examination, for instance, is not advisable, as results are unpredictable.
SMP CONSIDERATIO	DNS
	Because of the concurrent execution environment of VxWorks SMP the specified task <b>must</b> explicitely be put in a non-executing state before calling this routine.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the task ID is invalid.
ERRNOS	N/A
SEE ALSO	taskInfo, taskSuspend(), taskRegsSet()

## taskRegsSet()

NAME	taskRegsSet() – set a task's registers
SYNOPSIS	STATUS taskRegsSet ( int tid, /* task ID */ REG_SET * pRegs /* get register contents from here */ )
DESCRIPTION	This routine loads a specified register set $pRegs$ into a specified task's TCB.
NOTE	This routine only works well if the task is known not to be in the ready state. Suspending the task before changing the register set is recommended.
SMP CONSIDERATIO	NS
	Because of the concurrent execution environment of VxWorks SMP the specified task <b>must</b> explicitely be put in a non-executing state before calling this routine.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the task ID is invalid.
ERRNOS	N/A
SEE ALSO	taskInfo, taskSuspend( ), taskRegsGet( )

	taskR	legsShow(	)			
NAME	taskRegs	Show() – display t	he contents c	of a task's registers	3	
Synopsis	void tas) ( int t )	KRegsShow	*/			
DESCRIPTION	This routi	ne displays the reg	ister content	s of a specified tas	sk on standar	d output.
EXAMPLE	The follov -> taskRe	ving example displ	lays the regis eToId ("tSh	ter of the shell tas	k (PowerPC	family):
	r0 r3 r6 r9 r12 r15 r18 r21 r24 r27 r30 lr cr scSrTblPP value = 2	= 0x0000000 = 0x0000000 = 0x0000000 = 0x0000000 = 0x0000000 = 0x0000000 = 0x02f93c22 = 0x02f93e24 = 0x00c30b88 = 0x007877a4 = 0x0000000 = 0x2000480 cr = 0x01	sp r4 r7 r10 r13 r16 r19 r22 r25 r28 r31 ctr xer srTblPtr	<pre>= 0x00c30ae0 = 0x0000000 = 0x00000000 = 0x00000000 = 0x00000000 = 0x00000000 = 0x00000000 = 0x00000003 = 0x00000003 = 0x00348494 = 0x0200b030 = 0x0000000 = 0x00000000 = 0x00000000</pre>	r2 r5 r8 r11 r14 r17 r20 r23 r26 r29 msr pc pgTblPtr	<pre>= 0x00000000 = 0x00000000 = 0x00000000 = 0x00000000 = 0x00000000 = 0x00000000 = 0x00000000 = 0x00000000 = 0x00073f80 = 0xfffffff = 0x0200b030 = 0x0025d66c = 0x00745000</pre>
RETURNS	N/A					
ERRNO	N/A					
SEE ALSO	taskShow	r				

## taskRestart()

```
NAME taskRestart() - restart a task
SYNOPSIS STATUS taskRestart
(
int tid /* task ID of task to restart */
)
```

DESCRIPTION	This routine "restarts" a task. The task is first terminated, and then reinitialized with the same ID, priority, options, original entry point, stack size, and parameters it had when it was terminated. Self-restarting of a calling task is performed by a newly spawned "tRestart" task. The shell utilizes this routine to restart itself when aborted.		
NOTE	If the task has modified any of its start-up parameters, the restarted task will start with the changed values.		
SMP CONSIDERATIO	INS		
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.		
WARNING	Restarting user mode tasks is not supported from kernel space, and may have unpredictable behavior. User mode tasks may be restarted from within an RTP (except for the initial task).		
RETURNS	<b>OK</b> , or <b>ERROR</b> if the task ID is invalid or the task could not be restarted.		
ERRNO	S_intLib_NOT_ISR_CALLABLE		
	S_objLib_OBJ_DELETED		
	S_objLib_OBJ_UNAVAILABLE		
	S_objLib_OBJ_ID_ERROR		
	S_smObjLib_NOT_INITIALIZED		
	S_memLib_NOT_ENOUGH_MEMORY		
	S_memLib_BLOCK_ERROR		
	S_taskLib_ILLEGAL_PRIORITY		
SEE ALSO	taskLib		

## taskResume()

NAME	taskResume() – resume a task
SYNOPSIS	STATUS taskResume ( int tid /* task ID of task to resume */ )
DESCRIPTION	This routine resumes a specified task. Suspension is cleared, and the task operates in the remaining state.

#### SMP CONSIDERATIONS

	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	OK, or ERROR if the task cannot be resumed.
ERRNO	S_objLib_OBJ_ID_ERROR
SEE ALSO	taskLib

## taskRotate()

NAME	taskRotate() – rotate ready queue for a given task priority
SYNOPSIS	STATUS taskRotate ( int priority /* VX_TASK_PRIORITY_MIN to VX_TASK_PRIORITY_MAX */ )
DESCRIPTION	This routine rotates the ready queue of tasks that are ready to run for the priority specified by the <i>priority</i> parameter. In the special case that <i>priority</i> is set to <b>TASK_PRIORITY_SELF</b> , the ready queue for the caller's normal (spawned) priority is rotated. If no tasks are ready, or only one task is ready at the specified priority, no action is taken and this routine returns <b>OK</b> and leaves errno unchanged.
SMP CONSIDERATIO	NS
	This routine is not supported for SMP. This routine will return <b>ERROR</b> if called from a SMP system.
NOTE	The ITRON API <b>rot_rdq()</b> can be implemented using the following macro.
	#define rot_rdq(p) taskRotate (p == $0$ ? TASK_PRIORITY_SELF : p)
	or using the following function definition.
	STATUS rot_rdq
	( UINT priority
	return taskRotate ( priority == 0 ? TASK_PRIORITY_SELF: priority); }

RETURNS	OK, or ERROR		
ERRNO	S_intLib_NOT_ISR_CALLABLE Routine is not callable from an ISR.		
	S_taskLib_ILLEGAL_PRIORITY Priority specified is not within VX_TASK_PRIORITY_MIN to VX_TASK_PRIORITY_MAX.		
SEE ALSO	taskRotate		

## taskSRInit()

NAME	taskSRInit() – initialize the default task status register (MIPS)
SYNOPSIS	ULONG taskSRInit ( ULONG newSRValue /* new default task status register */ )
DESCRIPTION	This routine sets the default status register for system-wide tasks. All tasks will be spawned with the status register set to this value; thus, it must be called before <b>kernelInit()</b> .
RETURNS	The previous value of the default status register.
ERRNO	Not Available
SEE ALSO	taskArchLib

## taskSRSet()

NAME taskSRSet() – set the task status register (MC680x0, MIPS, x86)

SYNOPSIS STATUS taskSRSet ( int tid, /\* task ID \*/ UINT16 sr /\* new SR \*/ ) VxWorks Kernel API Reference, 6.6 taskSafe()

DESCRIPTION	This routine sets the status register of a task that is not running (i.e., the TCB must not be that of the calling task). Debugging facilities use this routine to set the trace bit in the status
	register of a task that is being single-stepped.
	x86: The second parameter represents EFLAGS register and the size is 32 bit.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the task ID is invalid.
ERRNO	Not Available
SEE ALSO	taskArchLib

## taskSafe()

NAME	taskSafe( ) – make the calling task safe from deletion
SYNOPSIS	STATUS taskSafe (void)
DESCRIPTION	This routine protects the calling task from deletion. Tasks that attempt to delete a protected task will block until the task is made unsafe, using <b>taskUnsafe()</b> . When a task becomes unsafe, the deleter will be unblocked and allowed to delete the task.
	The <b>taskSafe()</b> primitive utilizes a count to keep track of nested calls for task protection. When nesting occurs, the task becomes unsafe only after the outermost <b>taskUnsafe()</b> is executed.
RETURNS	OK.
ERRNO	N/A
SEE ALSO	taskLib, taskUnsafe(), the VxWorks programmer guides.

## taskShow()

NAME	<b>taskShow()</b> – display task information from TCBs
SYNOPSIS	STATUS taskShow
	int tid, /* task ID */

```
int level /* 0 = summary, 1 = details, 2 = all tasks */
)
```

**DESCRIPTION** This routine displays the contents of a task control block (TCB) for a specified task. If *level* is 1, it also displays task options and registers. If *level* is 2, it displays all tasks in sorted order of the number of tasks are less than 500. If more than 500 tasks are in the system, *level* equal to 2 will display all tasks in the system unsorted and in the order they are created.

The TCB display contains the following fields:

Field	Meaning
NAME	Task name (truncated, if ending with a > character)
ENTRY	Symbol name or address where task began execution
TID	Task ID
PRI	Priority
STATUS	Task status, as formatted by <b>taskStatusString()</b>
PC	Program counter
SP	Stack pointer
ERRNO	Most recent error code for this task
DELAY	If task is delayed, number of clock ticks remaining in delay (0 otherwise)
CPU #	For SMP systems, CPU index the task is running on ("-" otherwise)

Stack and register information for the specified task are also displayed.

#### SMP CONSIDERATIONS

Specifying a *level* of 2 will display a "CPU #" column instead of "DELAY". Specifying a *level* of 3 will result in the same output format as uniprocessor VxWorks, i.e. the "DELAY" column will be displayed.

EXAMPLE

The following example shows the TCB contents for the network task (PowerPC family):

UP version:

-> taskShow tNetTask, 1

NAME DELAY	ENTRY	TID	PRI	STATUS	PC	SP	ERRNO
 tNetTask 0	netTask	7b3c50	50 P	END	25d66c	b056d0	0
task stack: 8592	base 0xb057a	a0 end	0xb0309	0 size	e 10000 hi	gh 1408	margin
exc. stack: exc. stack:	base 0xb067a size 4080	10 end high	0xb057k 1 272	o0 star marg	rt 0xb067b0 gin 3808	)	
proc id: 0x options: 0x VX_SUPERVIS	36ffa8 ((nul] 9007 OR_MODE VX_U	.)) INBREAKAI	3LE	VX_DEAI	LLOC_STACK	VX_DEA	LLOC_TCE

2

## VxWorks Kernel API Reference, 6.6 taskShow()

VX\_DEALLOC\_EXC\_STACK

VxWorks Events \_\_\_\_\_ Events Pended on : Not Pended Received Events : 0x0 Options : N/A r0  $= 0 \times 00000000$  sp  $= 0 \times 00 b 056 d 0 r 2$  $= 0 \times 00000000$  r4  $= 0 \times 00000000$  r5 r3 r6 = 0x0000000 r7 = 0x0000000 r8 r9 = 0x0000000 r10 = 0x0000000 r11  $= 0 \times 00000000$ = 0x00000000 r14 = 0x0000000 r13  $= 0 \times 00000000$ r12 = 0x00000000 r17 = 0x00000000 r20  $\begin{array}{rcrcrcrc} = & 0x00000000 & r17 & = & 0x00000000 \\ = & 0x00000000 & r20 & = & 0x00000000 \\ = & 0x00000000 & r23 & = & 0x00000000 \\ = & 0x00000000 & r26 & = & 0x00371adc \\ = & 0x00348494 & r29 & = & 0xfffffff \\ = & 0x0000b030 & msr & = & 0x0000b030 \\ = & 0x00000000 & pc & = & 0x0025d66c \\ \end{array}$ = 0x0000000 r16 r15  $= 0 \times 00000000$ = 0x0000000 r19 r18  $= 0 \times 00000000$  r22 r21 r24  $= 0 \times 00000000$  r25  $= 0 \times 0000 b 030 r 28$ r27 r30 = 0x00371a40 r31 = 0x00000000 ctr lr  $\begin{array}{rcl} \mbox{lr} &= 0 \\ \mbox{cv} &= 0 \\ \mbox{v0000000} & \mbox{cr} &= 0 \\ \mbox{v00000000} & \mbox{pc} \\ \mbox{cr} &= 0 \\ \mbox{v00000000} & \mbox{pc} \\ \mbox{pc} \\ \mbox{pc} \\ \mbox{cv} \\ \mbox$ scSrTblPtr = 0x007841f4 srTblPtr = 0x007841f4 coprocTaskShow: TaskId 0x7b3c50 has no coprocessors selected value =  $0 = 0 \times 0$ SMP version (shows TCB contents for the log task): -> taskShow tLogTask, 1 NAME ENTRY TID PRI STATUS PC SP ERRNO DELAY \_\_\_\_\_ \_\_ \_\_\_\_\_ \_\_\_\_ tLogTask logTask 2ce5a0 0 PEND 1ecad8 2ce480 0 0 task affinity: 0x0 task cpuIndex: -1 (Task Not Running) task stack: base 0x2ce5a0 end 0x2cd210 size 5008 high 384 margin 4624 exc. stack: base 0x2cf7d0 end 0x2ce840 start 0x2cf840 exc. stack: size 3984 high 0 margin 3984 proc id: 0x245910 ((null)) options: 0x9003 VX\_SUPERVISOR\_MODE VX\_UNBREAKABLE VX\_DEALLOC\_TCB VX\_DEALLOC\_EXC\_STACK VxWorks Events \_\_\_\_\_ Events Pended on : Not Pended Received Events : 0x0 Options : N/A r0 = 0x0000000 sp = 0x002ce480 r2 = 0x00000000

r3 =	$= 0 \times 00000000$	r4		$= 0 \times 000000000$	r5	=	0x00000000
r6 =	$= 0 \times 000000000$	r7		$= 0 \times 000000000$	r8	=	0x00000000
r9 =	$= 0 \times 000000000$	r10		$= 0 \times 000000000$	r11	=	0x00000000
r12 =	$= 0 \times 000000000$	r13		$= 0 \times 000000000$	r14	=	0x00000000
r15 =	$= 0 \times 000000000$	r16		$= 0 \times 000000000$	r17	=	0x00000000
r18 =	$= 0 \times 000000000$	r19		$= 0 \times 00000000$	r20	=	0x00000000
r21 =	$= 0 \times 000000000$	r22		$= 0 \times 00000000$	r23	=	0x00000000
r24 =	$= 0 \times 00000000$	r25		= 0x002ce4b8	r26	=	0x00000000
r27 =	= 0xfffffff	r28		$= 0 \times 00000001$	r29	=	0x002cc830
r30 :	= 0x002cc7f0	r31		= 0xfffffff	msr	=	0x0000b030
lr :	$= 0 \times 00000000$	ctr		$= 0 \times 00000000$	pc	=	0x001ecad8
cr :	= 0x20000000	xer		$= 0 \times 00000000$	pgTblPt	r =	0x00279000
scSrTblPtr :	= 0x00278154	srTblPt	r	$= 0 \times 00278154$			
coprocTaskSl	how: TaskId	0x2ce5a0 1	has	no coprocess	ors select	ed	
value = 0 =	0x0						
SMP level 3 -> taskShow	taskShow exa 0, 3	mple:					
NAME	ENTRY	ΨTD	PRT	STATUS	PC.	SP	ERRNO
DELAY							
tExcTask	192c6c	26ce00	0	PEND	1ef424	26ef80	0
0							
tJobTask 0	193bc4	2cb440	0	PEND	1ef424	2cb380	0
tLogTask	logTask	2ce5a0	0	PEND	lecad8	2ce480	0
U tNhiotog	195070	241-20	0	סואיזס	1of121	241410	0
0	193070	Zuiezu	0	FEND	161424	Zuiuiu	0
0 + 9hol 10	abol 1maak	204130	1	DEADV	1 = 0 1 6 1	202300	0
0	SHELLIASK	264130	T	KEAD1	110104	262360	0
miiBusMoni> 50	140fc8	2c3010	254	DELAY	1f5844	2c2f8	0 0
tIdleTask0	idleTaskEntr	272c30	287	READY	1eeda8	272bb0	0
u tIdleTask1 0	idleTaskEntr	276250	287	READY	1eedb4	2761d0	0
value = 0 =	0x0						

If the specified task uses coprocessors, such as a floating point coprocessor, this routine will also display the registers for the corresponding coprocessor.

If this routine is called with the current task as the argument, register information will not be displayed. Use **taskRegsShow()** to display the register information for the calling task instead.

RETURNS N/A

ERRNO N/A

**SEE ALSO taskShow**, **taskStatusString()**, **windsh**, the VxWorks programmer guides, the, *VxWorks Command-Line Tools User's Guide*.

## taskShowInit()

NAME	taskShowInit() – initialize the task show routine facility
SYNOPSIS	void taskShowInit (void)
DESCRIPTION	This routine links the task show routines into the VxWorks system. It is called automatically when the task show facility is configured into VxWorks using either of the following methods:
	- If you use the configuration header files, define INCLUDE_SHOW_ROUTINES in config.h.
	- If you use the project facility, select INCLUDE_TASK_SHOW.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	taskShow

## taskSigqueue()

NAME	<b>taskSigqueue( )</b> – send a queued signal to a task			
SYNOPSIS	<pre>int taskSigqueue   (    int    int    const union sigval   )</pre>	tid, signo, value		
DESCRIPTION	The function <b>sigqueue()</b> sends the signal specified by <i>signo</i> with the signal-parameter value specified by <i>value</i> to the process specified by <i>tid</i> .			
	This function is currently a achieve uniform meaning a	liased to <b>sigqueue()</b> , and is provided as a convenience to across both kernel and user-mode code.		

RETURNS	<b>OK</b> (0), or <b>ERROR</b> (-1) if the task ID or signal number is invalid, or if there are no queued-signal buffers available.
ERRNO	EINVAL
SEE ALSO	sigLib.sigqueue()

### taskSpareFieldGet()

NAME	<b>taskSpareFieldGet( )</b> – get the spare field of a TCB
SYNOPSIS	<pre>int taskSpareFieldGet   (    int tid, /* task ID */    SPARE_NUM numAllotted /* spare field to get */   )</pre>
DESCRIPTION	This routine gets the value of a spare field. The spare field was gotten by calling <b>taskSpareNumAllot()</b> to get the available spare field to use.
RETURNS	value of the spare field, or <b>ERROR</b> if task or numAllotted is invalid
ERRNO	N/A
SEE ALSO	taskUtilLib, taskSpareNumAllot( ), taskSpareFieldSet( )

## taskSpareFieldSet()

NAME	<b>taskSpareFieldSet( )</b> – set the s	pare field of a TCB
SYNOPSIS	STATUS taskSpareFieldSet ( int tid, SPARE_NUM numAllotted, int value )	/* task ID */ /* spare field to set */ /* value to set */
DESCRIPTION	This routine sets the value of a <b>taskSpareNumAllot( )</b> to get th	spare field. The spare field is gotten by calling ne available spare field to use.

VxWorks Kernel API Reference, 6.6 taskSpareNumAllot()

An example:

```
int spareNum;

taskSpareNumAllot (t1, &spareNum);

if (spareNum != ERROR)

taskSpareFieldSet (t1, spareNum, 0x12345678);

RETURNS OK, or ERROR if task is invalid or numAllotted is invalid

ERRNOS N/A

SEE ALSO taskUtilLib, taskSpareFieldGet(), taskSpareNumAllot()
```

### taskSpareNumAllot()

NAME	taskSpareNumAllot() – Allocate the first available spare field in the TCB
SYNOPSIS	<pre>void taskSpareNumAllot   (    int tid, /* task ID */    SPARE_NUM * numAllotted /* where to return SPARE_NUM */   )</pre>
DESCRIPTION	This routine allots the first available spare field in the TCB. Once a spare field is allotted, this same number is used to reference the same spare field on all WIND_TCBs in the system. In other words, two different tasks, t1 and t2, will get different <b>SPARE_NUM</b> values when called.
	Once a field has been allotted, the field is reserved for the life of the system. A field may not be un-allotted or unreserved.
RETURNS	SPARE_NUM, or ERROR if fields are not available
ERRNOS	N/A
SEE ALSO	taskUtilLib, taskSpareFieldGet( ), taskSpareFieldSet( )

### taskSpawn()

NAME taskSpawn() – spawn a task

#### SYNOPSIS int taskSpawn ( char \* name, /\* name of new task (stored at pStackBase) \*/ int priority, /\* priority of new task \*/ int options, /\* task option word \*/ stackSize, /\* size (bytes) of stack needed plus name \*/ int FUNCPTR entryPt, /\* entry point of new task \*/ int arg1, /\* 1st of 10 reg'd args to pass to entryPt \*/ int arg2, int arg3, int arg4, int arg5, int arg6, int arg7, int arg8, int arg9, int arg10 )

**DESCRIPTION** This routine creates and activates a new task with a specified priority and options and returns a system-assigned ID. See **taskInit()** and **taskActivate()** for the building blocks of this routine.

A task may be assigned a name as a debugging aid. This name will appear in displays generated by various system information facilities such as **i()**. The name may be of arbitrary length and content, but the current VxWorks convention is to limit task names to ten characters and prefix them with a "t". If *name* is specified as **NULL**, an ASCII name will be assigned to the task of the form "t*n*" where *n* is an integer which increments as new tasks are spawned.

VxWorks schedules tasks on the basis of *priority*. Tasks may have priorities ranging from 0, the highest priority, to 255, the lowest priority. The priority of a task in VxWorks in dynamic and one may change an existing task's priority with **taskPrioritySet()**.

The only resource allocated to a spawned task is a stack of a specified size *stackSize*, which is allocated from the system memory partition. Stack size should be an even integer. A task control block (TCB) is carved from the stack, as well as any memory required by the task name. The remaining memory is the task's stack and every byte is filled with the value 0xEE (unless the VX\_NO\_STACK\_FILL option is specifed or the global kernel configuration parameter VX\_GLOBAL\_NO\_STACK\_FILL is set to TRUE) for the **checkStack()** facility. See the manual entry for **checkStack()** for stack-size checking aids.

The entry address *entryPt* is the address of the "main" routine of the task. The routine will be called once the C environment has been set up. The specified routine will be called with the ten given arguments. Should the specified main routine return, a call to **exit()** will automatically be made.

Note that ten (and only ten) arguments must be passed for the spawned function.

Bits in the options argument may be set to run with the following modes:

#### VX\_FP\_TASK

execute with floating-point coprocessor support. A task which performs floating point operations or calls any functions which either return or take a floating point value as arguments must be created with this option. Some routines perform floating point operations internally. The VxWorks documentation for these clearly state the need to use the VX\_FP\_TASK option.

#### VX\_ALTIVEC\_TASK

execute with Altivec support (PowerPC only)

#### VX\_SPE\_TASK

execute with SPE support (PowerPC only)

#### VX\_DSP\_TASK

execute with DSP support (SuperH only)

#### VX\_PRIVATE\_ENV

include private environment support (see envLib).

VX\_NO\_STACK\_FILL

do not fill the stack for use by checkStack().

VX\_UNBREAKABLE do not allow breakpoint debugging.

#### VX\_NO\_STACK\_PROTECT

do not provide stack protection: no overflow or underflow detection, stack remains executable.

The option bits VX\_DEALLOC\_STACK and VX\_DEALLOC\_EXC\_STACK are not options available for the **taskSpawn()** API. **taskSpawn()** internally sets these option bits by default depending on the configuration of the system. Specifying these options to **taskSpawn()** results in an **ERROR** and the ERRNO, **S\_taskLib\_ILLEGAL\_OPTIONS**, will be returned. See the definitions in **taskLib.h**.

It is assumed that the caller passes a valid function pointer as *entryPt* while calling this API. No validity check for this parameter is done here.

#### SMP CONSIDERATIONS

	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	The task ID, or <b>ERROR</b> if memory is insufficient or the task cannot be created.
ERRNO	S_intLib_NOT_ISR_CALLABLE Routine is not callable from an ISR.

	S_smObjLib_NOT_INITIALIZED S_memLib_NOT_ENOUGH_MEMORY Out of memory for allocation of stack or TCB.
	S_memLib_BLOCK_ERROR
	<b>S_taskLib_ILLEGAL_PRIORITY</b> Priority specified is not within 0-255.
	S_taskLib_ILLEGAL_OPTIONS The following illegal options were set: VX_DEALLOC_STACK, VX_DEALLOC_EXC_STACK, or VX_DEALLOC_TCB for taskSpawn().
	S_taskLib_ILLEGAL_STACK_INFO An invalid stack size has been specified
SEE ALSO	taskLib, taskInit(), taskActivate(), sp(), the VxWorks programmer guides.

## taskStackAllot()

S\_objLib\_OBJ\_ID\_ERROR

NAME	taskStackAllot() – allot memory from a task's exception stack
SYNOPSIS	<pre>void * taskStackAllot   (    int tid, /* task whose stack will be allotted from */    unsigned nBytes /* number of bytes to allot */  )</pre>
DESCRIPTION	This routine allots the specified amount of memory from the start of the exception stack of the task specified by <i>tid</i> . This is a non-blocking operation meant to be used by task create hooks that need to allocate small amounts of memory on a per-task basis. Since the memory is carved from the exception stack calling this routine essentially causes the amount of stack space available for execution to be reduced. Hence the requirement to only allocate small amounts of memory. The exception stack size of a kernel task cannot be modified. The exception stack size of tasks that run in real-time processes is controlled by the <b>USER_TASK_EXC_STACK_SIZE</b> configuration parameter of the <b>INCLUDE_KERNEL</b> component.
	It is an error condition for a task to call this routine to attempt to allocate memory from its own exception stack. Attempting to do this results in this routine returning <b>NULL</b> . Attempting to allocate memory from the exception stack of a task that has started execution has undefined results. This routine is meant to only ever be called before a task initially starts executing.

VxWorks Kernel API Reference, 6.6 taskStatusString( )

	The memory allocated with this routine cannot be added back to the task's exception stack. It will be reclaimed as part of the reclamation of the exception stack when the task is deleted.
	Note that an exception stack underrun will overwrite the allotments made from this routine because all portions are carved from the start of the exception stack.
	This routine returns NULL if the requested size exceeds available stack memory.
RETURNS	pointer to block, or NULL if unsuccessful.
ERRNO	N/A
SEE ALSO	taskLib, taskCreateHookAdd()

## taskStatusString()

NAME	taskStatusString() – get a task's status as a string
SYNOPSIS	STATUS taskStatusString ( int tid, /* task to get string for */ char * pString /* where to return string */ )
DESCRIPTION	This routine deciphers the WIND task status word in the TCB for a specified task, and copies

the appropriate string to *pString*.

The formatted string is one of the following:

String	Meaning
READY	Task is not waiting for any resource other than the CPU.
PEND	Task is blocked due to the unavailability of some resource.
DELAY	Task is asleep for some duration.
SUSPEND	Task is unavailable for execution (but not delayed, or pended).
STOP	Task is stopped by the debugger.
DELAY+S	Task is both delayed and suspended.
PEND+S	Task is both pended and suspended.
PEND+T	Task is pended with a timeout.
STOP+P	Task is both pended and stopped by the debugger.
STOP+S	Task is both stopped by the debugger and suspended.
STOP+T	Task is both delayed and stopped by the debugger.
PEND+S+T	Task is pended with a timeout, and also suspended.
STOP+P+S	Task is pended, suspended, and also stopped by the debugger.
STOP+P+T	Task is pended with a timeout, and also stopped by the debugger.
STOP+S+T	Task is suspended with a timeout, and also stopped by the debugger.

String	Meaning
ST+P+S+T	Task is pended with a timeout, suspended, and stopped by the debugger.
+I	Task has inherited priority (+I may be appended to any string above).
DEAD	Task no longer exists.
-> taskStatusString (taskNameToId ("tShell0"), xx=malloc (10)) new symbol "xx" added to symbol table. value = 0 = 0x0 -> printf ("shell status = <%s>\n", xx) shell status = <ready> value = 2 = 0x2</ready>	
OK, or ERROR if	the task ID is invalid.
N/A	
taskShow	
	<pre>String ST+P+S+T +I DEAD -&gt; taskStat new symbol value = 0 = -&gt; printf shell statu value = 2 = OK, or ERROR if N/A taskShow</pre>

## taskSuspend()

NAME	taskSuspend() – suspend a task
SYNOPSIS	STATUS taskSuspend ( int tid /* task ID of task to suspend */ )
DESCRIPTION	This routine suspends a specified task. A task ID of zero results in the suspension of the calling task. Suspension is additive, thus tasks can be delayed and suspended, or pended and suspended. Suspended, delayed tasks whose delays expire remain suspended. Likewise, suspended, pended tasks that unblock remain suspended only.
	Care should be taken with asynchronous use of this facility. The specified task is suspended regardless of its current state. The task could, for instance, have mutual exclusion to some system resource, such as the network or system memory partition. If suspended during such a time, the facilities engaged are unavailable, and the situation often ends in deadlock.
	This facility should be rejected as a synchronization mechanism in favor of the more general semaphore facility.
SMP CONSIDERATIO	

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

VxWorks Kernel API Reference, 6.6 taskSwitchHookAdd()

**RETURNS** OK, or ERROR if the task cannot be suspended.

ERRNO S\_objLib\_OBJ\_ID\_ERROR

SEE ALSO taskLib

### taskSwitchHookAdd()

NAME taskSwitchHookAdd() – add a routine to be called at every task switch SYNOPSIS STATUS taskSwitchHookAdd ( FUNCPTR switchHook /\* routine to be called at every task switch \*/ ) DESCRIPTION This routine adds a specified routine to a list of routines that will be called at every task switch. The routine should be declared as follows: void switchHook ( WIND\_TCB \*pOldTcb, /\* pointer to old task's WIND\_TCB \*/ /\* pointer to new task's WIND\_TCB \*/ WIND\_TCB \*pNewTcb ) User-installed switch hooks are called within the kernel context. Therefore, switch hooks do NOTE not have access to all VxWorks facilities. The following routines can be called from within a task switch hook: Library Routines All routines bLib fppArchLib fppSave(), fppRestore() intLib intContext(), intCount(), intVecSet(), intVecGet() All routines lstLib mathALib All routines, if **fppSave()**/**fppRestore()** are used rngLib All routines except rngCreate() taskLib taskIdVerify(), taskIdDefault(), taskIsReady() taskIsSuspended(), taskTcb() vxLib vxTas() **OK**, or **ERROR** if the table of task switch routines is full. RETURNS Not Available ERRNO taskHookLib, taskSwitchHookDelete() SEE ALSO

taskSwitchHookDelete()	
taskSwitchHookDelete() – delete a previously added task switch routine	
STATUS taskSwitchHookDelete ( FUNCPTR switchHook /* routine to be deleted from list */ )	

- DESCRIPTION This routine removes the specified routine from the list of routines to be called at each task switch.
- RETURNS OK, or ERROR if the routine is not in the table of task switch routines.
- Not Available ERRNO

NAME

SYNOPSIS

SEE ALSO taskHookLib, taskSwitchHookAdd()

### taskSwitchHookShow()

taskSwitchHookShow() - show the list of task switch routines NAME

- SYNOPSIS void taskSwitchHookShow (void)
- DESCRIPTION This routine shows all the switch routines installed in the task switch hook table, in the order in which they were installed.
- N/A RETURNS
- Not Available ERRNO
- SEE ALSO taskHookShow, taskSwitchHookAdd()

### taskTcb()

NAME taskTcb() – get the task control block for a task ID

WIND\_TCB \* taskTcb SYNOPSIS

VxWorks Kernel API Reference, 6.6 taskUnlink()

```
(
int tid /* task ID */
)
```

**DESCRIPTION** This routine returns a pointer to the task control block (WIND\_TCB) for a specified task. Although all task state information is contained in the TCB, users must not modify it directly. To change registers, for instance, use taskRegsSet() and taskRegsGet().

**RETURNS** A pointer to a **WIND\_TCB**, or **NULL** if the task ID is invalid.

ERRNO S\_objLib\_OBJ\_ID\_ERROR

SEE ALSO taskLib

### taskUnlink()

NAME	taskUnlink() – unlink a task	
SYNOPSIS	STATUS taskUnlink ( const char * name /* name of task to unlink */ )	
DESCRIPTION	This routine removes a task from its name space. The use of this routine on private tasks, which support duplicate names, is not recomended. After a task is unlinked, subsequent calls to <b>taskOpen()</b> using <i>name</i> will not be able to find the task, even if it has not been deleted yet. Instead, a new task could be created if <b>taskOpen()</b> is called with the <b>OM_CREATE</b> flag.	
	This routine is not ISR callable.	
SMP CONSIDERATIO	NS	
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.	
RETURNS	OK, or ERROR if unsuccessful.	
ERRNO	S_objLib_OBJ_INVALID_ARGUMENT name is NULL or empty.	
	S_objLib_OBJ_NOT_FOUND No task with <i>name</i> was found.	

### S\_intLib\_NOT\_ISR\_CALLABLE This routine must not be called from an ISR.

SEE ALSO taskOpen, taskOpen(), taskClose()

### taskUnlock()

NAME taskUnlock() – enable task rescheduling

SYNOPSIS STATUS taskUnlock

void

)

**DESCRIPTION** This routine decrements the preemption lock count. Typically this call is paired with **taskLock()** and concludes a critical section of code. Preemption will not be unlocked until **taskUnlock()** has been called as many times as **taskLock()**. When the lock count is decremented to zero, any tasks that were eligible to preempt the current task will execute.

The taskUnlock() routine is not callable from interrupt service routines.

#### SMP CONSIDERATIONS

This API is not available in VxWorks SMP.

RETURNS	OK or ERROR.
	010 01 210010

ERRNO S\_intLib\_NOT\_ISR\_CALLABLE

SEE ALSO taskLib, taskLock(), taskCpuUnlock()

### taskUnsafe()

NAME taskUnsafe() – make the calling task unsafe from deletion

SYNOPSIS STATUS taskUnsafe (void)

**DESCRIPTION** This routine removes the calling task's protection from deletion. Tasks that attempt to delete a protected task will block until the task is unsafe. When a task becomes unsafe, the deleter will be unblocked and allowed to delete the task.

The **taskUnsafe()** primitive utilizes a count to keep track of nested calls for task protection. When nesting occurs, the task becomes unsafe only after the outermost **taskUnsafe()** is executed.

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.

RETURNS	ОК.
ERRNO	N/A
SEE ALSO	taskLib, taskSafe(), the VxWorks programmer's guides

### taskVarAdd()

taskVarAdd() – add a task variable to a task NAME SYNOPSIS STATUS taskVarAdd ( int tid, /\* ID of task to have new variable \*/ int \*pVar /\* pointer to variable to be switched for task \*/ DESCRIPTION This routine adds a specified variable pVar (4-byte memory location) to a specified task's context. After calling this routine, the variable will be private to the task. The task can access and modify the variable, but the modifications will not appear to other tasks, and other tasks' modifications to that variable will not affect the value seen by the task. This is accomplished by saving and restoring the variable's initial value each time a task switch occurs to or from the calling task. This facility can be used when a routine is to be spawned repeatedly as several independent tasks. Although each task will have its own stack, and thus separate stack variables, they will all share the same static and global variables. To make a variable *not* shareable, the routine can call **taskVarAdd()** to make a separate copy of the variable for each task, but all at the same physical address. Note that task variables increase the task switch time to and from the tasks that own them.

Therefore, it is desirables increase the task switch time to and from the tasks that own them. Therefore, it is desirable to limit the number of task variables that a task uses. One efficient way to use task variables is to have a single task variable that is a pointer to a dynamically allocated structure containing the task's private data. **EXAMPLE** Assume that three identical tasks were spawned with a routine called **operator()**. All three use the structure **OP\_GLOBAL** for all variables that are specific to a particular incarnation of the task. The following code fragment shows how this is set up:

```
OP_GLOBAL *opGlobal; /* ptr to operator task's global variables */
void operator
    (
                    /* number of this operator task */
   int opNum
   )
    {
   if (taskVarAdd (0, (int *)&opGlobal) != OK)
       printErr ("operator%d: can't taskVarAdd opGlobal\\n", opNum);
       taskSuspend (0);
        }
   if ((opGlobal = (OP_GLOBAL *) malloc (sizeof (OP_GLOBAL))) == NULL)
        {
       printErr ("operator%d: can't malloc opGlobal\\n", opNum);
       taskSuspend (0);
       }
    . . .
   }
```

SMP CONSIDERATIONS

This routine is not available in VxWorks SMP. Use <u>thread</u> variables instead

RETURNS OK, or ERROR i	memory is insufficient for the task	variable descriptor or semaphore.
------------------------	-------------------------------------	-----------------------------------

**ERRNOS** no errnos for this routine

SEE ALSO taskVarLib, taskVarDelete(), taskVarGet(), taskVarSet()

### taskVarDelete()

private value of that variable is lost.

NAME	taskVarDelete() – remove a task variable from a task
SYNOPSIS	STATUS taskVarDelete ( int tid, /* ID of task whose variable is to be removed */ int *pVar /* pointer to task variable to be removed */ )
DESCRIPTION	This routine removes a specified task variable, $pVar$ , from the specified task's context. The

VxWorks Kernel API Reference, 6.6 taskVarGet()

#### SMP CONSIDERATIONS

SEE ALSO	taskVarLib, taskVarAdd( ), taskVarGet( ), taskVarSet( )
	<b>S_taskLib_TASK_VAR_NOT_FOUND</b> - address specified in <i>pVar</i> is not a task variable for <i>tid</i>
ERRNOS	Possible errno values set by this routine are:
RETURNS	OK, or ERROR if the task variable does not exist for the specified task.
	This routine is not available in VxWorks SMP. Use <u>thread</u> variables instead.

## taskVarGet()

NAME	taskVarGet() – get the value of a task variable		
SYNOPSIS	<pre>int taskVarGet   (    int tid, /* ID of task whose task variable is to be retrieved */    int *pVar /* pointer to task variable */   )</pre>		
DESCRIPTION	This routine returns the private value of a task variable for a specified task. The specified task is usually not the calling task, which can get its private value by directly accessing the variable. This routine is provided primarily for debugging purposes.		
SMP CONSIDERATIONS			
	This routine is not available in VxWorks SMP. Use <i>thread</i> variables instead.		
RETURNS	The private value of the task variable, or <b>ERROR</b> if the task is not found or it does not own the task variable.		
ERRNOS	Possible errno values set by this routine are:		
	<b>S_taskLib_TASK_VAR_NOT_FOUND</b> - address specified in <i>pVar</i> is not a task variable for <i>tid</i>		
SEE ALSO	taskVarLib, taskVarAdd( ), taskVarDelete( ), taskVarSet( )		

NAME

taskVarInfo() – get a list of task variables of a task

SYNOPSIS	int	taskVarIr (	nfo		
		int TASK_VAR int )	tid, varList[], maxVars	/* /* /*	ID of task whose task variable is to be set */ array to hold task variable addresses */ maximum variables varList can accommodate */

**DESCRIPTION** This routine provides the calling task with a list of all of the task variables of a specified task. The unsorted array of task variables is copied to *varList*.

#### SMP CONSIDERATIONS

This routine is not available in VxWorks SMP.

- **RETURNS** The number of task variables in the list or **ERROR** if the specified task ID is not valid.
- **ERRNOS** no errnos for this routine

SEE ALSO taskVarLib

### taskVarInit()

NAME	taskVarInit() – initialize the task variables facility		
SYNOPSIS	STATUS taskVarInit (void)		
DESCRIPTION	This routine initializes the task variables facility. It installs task switch and delete hooks used for implementing task variables. If <b>taskVarInit()</b> is not called explicitly, <b>taskVarAdd()</b> will call it automatically when the first task variable is added. After the first invocation of this routine, subsequent invocations have no effect.		
WARNING	Order dependencies in task delete hooks often involve task variables. If a facility uses task variables and has a task delete hook that expects to use those task variables, the facility's delete hook must run before the task variables' delete hook. Otherwise, the task variables will be deleted by the time the facility's delete hook runs.		
	VxWorks is careful to run the delete hooks in reverse of the order in which they were installed. Any facility that has a delete hook that will use task variables can guarantee proper ordering by calling <b>taskVarInit()</b> before adding its own delete hook.		

Note that this is not an issue in normal use of task variables. The issue only arises when adding another task delete hook that uses task variables.

Caution should also be taken when adding task variables from within create hooks. If the task variable package has not been installed via **taskVarInit()**, the create hook attempts to create a create hook, and that may cause system failure. To avoid this situation, **taskVarInit()** should be called during system initialization from the root task, **usrRoot()**, in **usrConfig.c**.

#### SMP CONSIDERATIONS

This routine is not available in VxWorks SMP.

**RETURNS** OK, or ERROR if the task switch/delete hooks could not be installed.

**ERRNOS** no errnos for this routine

SEE ALSO taskVarLib

### taskVarSet()

NAME	<b>taskVarSet( )</b> – set the value of a task variable		
SYNOPSIS	<pre>STATUS taskVarSet   (    int tid, /* ID of task whose task variable is to be set */    int *pVar, /* pointer to task variable to be set for this task */    int value /* new value of task variable */   )</pre>		
DESCRIPTION	This routine sets the private value of the task variable for a specified task. The specified task is usually not the calling task, which can set its private value by directly modifying the variable. This routine is provided primarily for debugging purposes.		
SMP CONSIDERATIO	NS		
	This routine is not available in VxWorks SMP. Use <u>thread</u> variables instead.		
RETURNS	OK, or ERROR if the task is not found or it does not own the task variable.		
ERRNOS	Possible errno values set by this routine are:		
	<b>S_taskLib_TASK_VAR_NOT_FOUND</b> - address specified in <i>pVar</i> is not a task variable for <i>tid</i>		
SEE ALSO	taskVarLib, taskVarAdd(), taskVarDelete(), taskVarGet()		

### td()

NAME	td() – delete a task		
SYNOPSIS	void td ( int taskNameOrId /* task name or task ID */ )		
DESCRIPTION	This command deletes a specified task. It simply calls <b>taskDelete()</b> .		
RETURNS	N/A		
ERRNO	N/A		
SEE ALSO	usrLib, taskDelete(), the VxWorks programmer guides.		

### tffsDevCreate()

tffsDevCreate() - create a TrueFFS block device suitable for use with dosFs NAME SYNOPSIS BLK\_DEV \* tffsDevCreate ( int tffsDriveNo, /\* TFFS drive number (0 - DRIVES-1) \*/ int removableMediaFlag /\* 0 - nonremovable flash media \*/ ) DESCRIPTION This routine creates a TFFS block device on top of a flash device. It takes as arguments a drive number, determined from the order in which the socket components were registered, and a flag integer that indicates whether the medium is removable or not. A zero indicates a non removable medium. A one indicates a removable medium. If you intend to mount dosFs on this block device, you probably do not want to call tffsDevCreate(), but should call usrTffsConfig() instead. Internally, usrTffsConfig() calls tffsDevCreate() for you. It then does everything necessary (such as calling the **dosFsDevInit()** routine) to mount dosFs on the just created block device. RETURNS BLK\_DEV pointer, or NULL if it failed. ERRNO Not Available SEE ALSO tffsDrv

### tffsDevFormat()

tffsDevFormat() - format a flash device for use with TrueFFS NAME SYNOPSIS STATUS tffsDevFormat ( int tffsDriveNo, /\* TrueFFS drive number (0 - DRIVES-1) \*/ int arg /\* pointer to tffsDevFormatParams structure \*/ ) This routine formats a flash device for use with TrueFFS. It takes two parameters, a drive DESCRIPTION number and a pointer to a device format structure. This structure describes how the volume should be formatted. The structure is defined in **dosformt.h**. The drive number is assigned in the order that the socket component for the device was registered. The format process marks each erase unit with an Erase Unit Header (EUH) and creates the physical and virtual Block Allocation Maps (BAM) for the device. The erase units reserved for the "boot-image" are skipped and the first EUH is placed at number (boot-image length - 1). To write to the boot-image region, call tffsBootImagePut(). WARNING If any of the erase units in the boot-image region contains an erase unit header from a previous format call (this can happen if you reformat a flash device specifying a larger boot region) TrueFFS fails to mount the device. To fix this problem, use tffsRawio() to erase the problem erase units (thus removing the outdated EUH). The macro TFFS\_STD\_FORMAT\_PARAMS defines the default values used for formatting a flask disk device. If the second argument to this routine is zero, tffsDevFormat() uses these default values. RETURNS OK, or ERROR if it failed. ERRNO Not Available tffsDrv SEE ALSO

### tffsDevOptionsSet()

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DESCRIPTION This routine is intended to set various TrueFFS volume options. At present it only disables FAT monitoring. If VxWorks long file names are to be used with TrueFFS, FAT monitoring must be turned off.

**RETURNS OK**, or **ERROR** if it failed.

ERRNO Not Available

SEE ALSO tffsDrv

### tffsDrv()

NAME tffsDrv() – initialize the TrueFFS system

SYNOPSIS STATUS tffsDrv (void)

**DESCRIPTION** This routine sets up the structures, the global variables, and the mutual exclusion semaphore needed to manage TrueFFS. This call also registers socket component drivers for all the flash devices attached to your target.

Because tffsDrv() is the call that initializes the TrueFFS system, this function must be called (exactly once) before calling any other TrueFFS utilities, such as tffsDevFormat() or tffsDevCreate(). Typically, the call to tffsDrv() is handled for you automatically. If you defined INCLUDE\_TFFS in your BSP's config.h, the call to tffsDrv() is made from usrRoot(). If your BSP's config.h defines INCLUDE\_PCMCIA, the call to tffsDrv() is made from pccardTffsEnabler().

- **RETURNS** OK, or **ERROR** if it fails.
- ERRNO Not Available

SEE ALSO tffsDrv

### tffsDrvOptionsSet()

NAME tffsDrvOptionsSet() – set TrueFFS volume options

SYNOPSIS STATUS tffsDrvOptionsSet

```
(
int tffsDriveNo /* TFFS drive number (0 - DRIVES-1) */
)
```

- This routine is intended to set various TrueFFS volume options. At present it only disables DESCRIPTION FAT monitoring. If VxWorks long file names are to be used with TrueFFS, FAT monitoring must be turned off. If Datalite's Reliance file file system is to be used with TrueFFS, FAT monitoring must be turned off.
- RETURNS OK, or ERROR if it failed.

Not Available ERRNO

tffsDrv SEE ALSO

### tffsRawio()

NAME	tffsRawio() – low level I/O	access to flash comj	ponents	
SYNOPSIS	<pre>STATUS tffsRawio   (     int tffsDriveNo, /*     int functionNo, /*     int arg0, /*     int arg1, /*     int arg2 /*     )</pre>	TrueFFS drive nu TrueFFS function argument 0 */ argument 1 */ argument 2 */	mber (0 - DRI . code */	/ES-1) */
DESCRIPTION	Use the utilities provided by thisroutine with the utmost care. If you use these routines carelessly, you risk data loss as well as permanent physical damage to the flash device.			
	This routine is a gateway to a series of utilities (listed below). Functions such as <b>mkbootTffs()</b> and <b>tffsBootImagePut()</b> use these <b>tffsRawio()</b> utilities to write boot sector information. The functions for physical read, write, and erase are made available with the intention that they be used on erase units allocated to the boot-image region by <b>tffsDevFormat()</b> . Using these functions elsewhere could be dangerous.			
	The <i>arg0</i> , <i>arg1</i> , and <i>arg2</i> parameters to <b>tffsRawio()</b> are interpreted differently depending on the function number you specify for <i>functionNo</i> . The drive number is determined by the order in which the socket components were registered.			
	Function Name	arg0	arg1	arg2
	TFFS_GET_PHYSICAL_INFO	user buffer	N/A	N/A

address

Function Name	arg0	arg1	arg2
TFFS_PHYSICAL_READ	address to read	byte count	user buffer
TFFS_PHYSICAL_WRITE	address to write	byte count	address user buffer address
TFFS_PHYSICAL_ERASE	first unit	number of units	N/A
TFFS_ABS_READ	sector number	number of sectors	user buffer
			address
TFFS_ABS_WRITE	sector number	number of sectors	user buffer
			address
TFFS_ABS_DELETE	sector number	number of sectors	N/A
TFFS_DEFRAGMENT_VOLUME	number of sectors	user buffer	N/A
		address	

#### TFFS\_GET\_PHYSICAL\_INFO

writes the flash type, erasable block size, and media size to the user buffer specified in *arg0*.

#### TFFS\_PHYSICAL\_READ

reads *arg1* bytes from *arg0* and writes them to the buffer specified by *arg2*.

#### TFFS\_PHYSICAL\_WRITE

copies *arg1* bytes from the *arg2* buffer and writes them to the flash memory location specified by *arg0*. This aborts if the volume is already mounted to prevent the versions of translation data in memory and in flash from going out of synchronization.

#### TFFS\_PHYSICAL\_ERASE

erases *arg1* erase units, starting at the erase unit specified in *arg0*. This aborts if the volume is already mounted to prevent the versions of translation data in memory and in flash from going out of synchronization.

#### TFFS\_ABS\_READ

reads *arg1* sectors, starting at sector *arg0*, and writes them to the user buffer specified in *arg2*.

#### TFFS\_ABS\_WRITE

takes data from the *arg2* user buffer and writes *arg1* sectors of it to the flash location starting at sector *arg0*.

#### TFFS\_ABS\_DELETE

deletes arg1 sectors of data starting at sector arg0.

#### TFFS\_DEFRAGMENT\_VOLUME

calls the defragmentation routine with the minimum number of sectors to be reclaimed, *arg0*, and writes the actual number reclaimed in the user buffer by *arg1*. Calling this function through some low priority task will make writes more deterministic. No validation is done of the user specified address fields, so the functions assume they are writable. If the address is invalid, you could see bus errors or segmentation faults.

VxWorks Kernel API Reference, 6.6 ti()

**RETURNS** OK, or ERROR if it failed.

ERRNO Not Available

SEE ALSO tffsDrv

### ti()

ti() – print complete information from a task's TCB NAME SYNOPSIS void ti ( int taskNameOrId /\* task name or task ID; 0 = use default \*/ ) This command prints the task control block (TCB) contents, including registers, for a DESCRIPTION specified task. If taskNameOrId is omitted or zero, the last task referenced is assumed. The ti() routine uses taskShow(); see the documentation for taskShow() for a description of the output format. The following shows the TCB contents for the shell task: EXAMPLE -> ti ENTRY TID PRI STATUS PC ERRNO DELAY NAME SP tShell0 shellTask 60351ba8 1 READY 6015fe68 603508d0 ad0007 0 task stack: base 0x60351ba8 end 0x6033eba8 size 77824 high 14144 margin 63680 exc. stack: base 0x60354f18 end 0x60351fd8 start 0x60354fd8 exc. stack: size 12096 high 4600 margin 7496 proc id: 0x60187008 ((null)) options: 0x1001007 VX\_SUPERVISOR\_MODE VX\_UNBREAKABLE VX\_DEALLOC\_STACK VX\_DEALLOC\_EXC\_STACK VX\_FP\_TASK VxWorks Events \_\_\_\_\_ Events Pended on : Not Pended Received Events : 0x0 Options : N/A value =  $0 = 0 \times 0$ RETURNS N/A

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ERRNO N/A

**SEE ALSO usrLib**, **taskShow()**, the VxWorks programmer guides.

### tick64Get( )

NAME tick64Get() – get the value of the kernel's tick counter as a 64 bit value

SYNOPSIS UINT64 tick64Get (void)

**DESCRIPTION** This routine returns the current value of the 64 bit absolute tick counter. This value is set to zero at startup, incremented by **tickAnnounce()**, and can be changed using **tickSet()** or **tick64Set()**.

#### SMP CONSIDERATIONS

	In SMP configuration this API is spinLock restricted meaning that calling this API while holding a spinLock is not allowed. Not comforming to this restriction will potentially lead to a live-lock scenerio.
RETURNS	The most recent <b>tickSet()/tick64Set()</b> value, plus all <b>tickAnnounce()</b> calls since.
ERRNO	N/A
SEE ALSO	tickLib, tickGet( ), tick64Set( ), tickSet( ), tickAnnounce( )

### tick64Set( )

NAME	tick64Set() – set the value of the kernel's tick counter in 64 bits
SYNOPSIS	void tick64Set ( UINT64 ticks /* new time in ticks */ )
DESCRIPTION	This routine sets the internal tick counter to a specified value in ticks. The new count will be reflected by <b>tick64Get()</b> and <b>tickGet()</b> (only the lower 32 bits), but will not change any delay fields or timeouts selected for any tasks. For example, if a task is delayed for ten ticks, and this routine is called to advance time, the delayed task will still be delayed until ten <b>tickAnnounce()</b> calls have been made.

VxWorks Kernel API Reference, 6.6 tickAnnounce()

#### SMP CONSIDERATIONS

	In SMP configuration this API is spinLock restricted meaning that calling this API while holding a spinLock is not allowed. Not comforming to this restriction will potentially lead to a live-lock scenerio.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	tickLib, tick64Get( ), tickGet( ), tickSet( ), tickAnnounce( )

## tickAnnounce()

NAME	tickAnnounce() – announce a clock tick to the kernel
SYNOPSIS	void tickAnnounce (void)
DESCRIPTION	This routine informs the kernel of the passing of time. It should be called from an interrupt service routine that is connected to the system clock. The most common frequencies are 60Hz or 100Hz. Frequencies in excess of 600Hz are an inefficient use of processor power because the system will spend most of its time advancing the clock. By default, this routine is called by <b>usrClock()</b> in <b>usrConfig.c</b> .
RETURNS	N/A
ERRNO	N/A
SEE ALSO	tickLib, kernelLib, taskLib, semLib, wdLib, the VxWorks programmer's guides

## tickAnnounceHookAdd()

NAME	tickAnnounceHookAdd() – add a hook routine to be called at each tick interrupt

SYNOPSIS STATUS tickAnnounceHookAdd (

FUNCPTR pFunc
**DESCRIPTION** This routine adds a hook to perform operations, such as round robin policy implementation, at each tick interrupt. The hooked function must follow the same ISR restrictions and must be callable at interrupt context. The user provided hook routine should be declared as follows:

```
void mySchedulerTickHook
  (
   int tid /* interrupt task ID */
  )
```

The user specified hook routines must not access the task structure fields directly. Access routines, such as **taskPriorityGet()**, should be used to access data structure fields.

ERRNO N/A

SEE ALSO tickLib

### tickGet()

NAME tickGet() – get the value of the kernel's tick counter

SYNOPSIS ULONG tickGet (void)

**DESCRIPTION** This routine returns the current value of the tick counter. This value is set to zero at startup, incremented by **tickAnnounce()**, and can be changed using **tickSet()**.

#### SMP CONSIDERATIONS

In SMP configuration this API is spinLock restricted meaning that calling this API while holding a spinLock is not allowed. Not comforming to this restriction will potentially lead to a live-lock scenerio.

**RETURNS** The most recent **tickSet()** value, plus all **tickAnnounce()** calls since.

ERRNO N/A

SEE ALSO tickLib, tickSet(), tickAnnounce()

## tickSet()

NAME	tickSet() – set the value of the kernel's tick counter	
SYNOPSIS	void tickSet ( ULONG ticks /* new time in ticks */ )	
DESCRIPTION	This routine sets the internal tick counter to a specified value in ticks. The new count will be reflected by <b>tickGet()</b> , but will not change any delay fields or timeouts selected for any tasks. For example, if a task is delayed for ten ticks, and this routine is called to advance time, the delayed task will still be delayed until ten <b>tickAnnounce()</b> calls have been made.	
SMP CONSIDERATIONS		
	In SMP configuration this API is spinLock restricted meaning that calling this API while holding a spinLock is not allowed. Not comforming to this restriction will potentially lead to a live-lock scenerio.	
RETURNS	N/A	
ERRNO	N/A	
SEE ALSO	tickLib, tickGet( ), tickAnnounce( )	

# timerOpenInit()

NAME	<b>timerOpenInit( )</b> – initialize the timer open facility
SYNOPSIS	void timerOpenInit (void)
DESCRIPTION	This routine links the timer creation routine with the open facility into the VxWorks system. It is called automatically when the timer facility is configured into VxWorks by either defining INCLUDE_OBJ_OPEN INCLUDE_POSIX_TIMERS in config.h or selecting INCLUDE_OBJ_OPEN and INCLUDE_POSIX_TIMERS in the project facility.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	timerOpen

## timerShowInit()

NAME	<b>timerShowInit( )</b> – initialize the timer show routine facility
SYNOPSIS	void timerShowInit (void)
DESCRIPTION	This routine links the timer show routines into the VxWorks system. It is called automatically when the timer show facility is configured into VxWorks using either of the following methods:
	- If you use the configuration header files, define INCLUDE_POSIX_TIMER_SHOW in config.h.
	- If you use the project facility, select INCLUDE_POSIX_TIMER_SHOW.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	timerShow

## timer\_cancel()

NAME	timer_cancel() – cancel a timer
SYNOPSIS	<pre>int timer_cancel   (    timer_t timerid /* timer ID */ )</pre>
DESCRIPTION	This routine is a shorthand method of invoking <b>timer_settime()</b> , which stops a timer.
NOTE	This is a non-POSIX API.
RETURNS	0 ( <b>OK</b> ), or -1 ( <b>ERROR</b> ) if <i>timerid</i> is invalid.
ERRNO	EINVAL
SEE ALSO	timerLib

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	timer_close()
NAME	timer_close() – close a named timer
SYNOPSIS	STATUS timer_close ( timer_t timerId /* timer ID to close */ )
DESCRIPTION	This routine closes a named timer and decrements its reference counter. In case where the counter becomes zero, the timer is deleted if:
	- It has been already removed from the name space by a call to <b>timer_unlink()</b> .
	- It was created with the <b>OM_DESTROY_ON_LAST_CALL</b> option.
NOTE	This is a non-POSIX API. This routine is not ISR callable.
SMP CONSIDERATI	ONS
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.
RETURNS	OK, or ERROR if unsuccessful.
ERRNO	S_objLib_OBJ_ID_ERROR The timer ID is invalid.
	S_objLib_OBJ_OPERATION_UNSUPPORTED The timer is not named.
	S_objLib_OBJ_DESTROY_ERROR An error was detected while deleting the timer.
	S_intLib_NOT_ISR_CALLABLE This routine must not be called from an ISR.
SEE ALSO	<pre>timerOpen, timer_open(), timer_unlink()</pre>

# timer\_connect()

NAME timer\_connect() – connect a user routine to the timer signal

SYNOPSIS int timer\_connect

(
timer\_t timerid, /\* timer ID \*/
VOIDFUNCPTR routine, /\* user routine \*/
int arg /\* user argument \*/
)

**DESCRIPTION** This routine sets the specified *routine* to be invoked with *arg* when fielding a signal indicated by the timer's *evp* signal number, or if *evp* is **NULL**, when fielding the default signal (SIGALRM).

The signal handling routine should be declared as:

SEE ALSO timerLib

### timer\_create( )

**DESCRIPTION** This routine returns a value in *pTimer* that identifies the timer in subsequent timer requests. The *evp* argument, if non-NULL, points to a **sigevent** structure, which is allocated by the application and defines the signal number and application-specific data to be sent to the task when the timer expires. If *evp* is NULL, a default signal (SIGALRM) is queued to the task, and the signal data is set to the timer ID. Initially, the timer is disarmed.

# **NOTE** If the task that created the timer goes away before the timer expires, the timer expiration process will display a warning message. However, if **INCLUDE\_OBJ\_OWNERSHIP** is

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VxWorks Kernel API Reference, 6.6 timer\_delete()

	configured, the timer will be deleted at the time the task is deleted and the message will not be displayed.
RETURNS	0 ( <b>OK</b> ), or -1 ( <b>ERROR</b> ) if too many timers already are allocated or the signal number is invalid.
ERRNO	EINVAL
	ENOSYS
	EAGAIN
	S_memLib_NOT_ENOUGH_MEMORY
SEE ALSO	timerLib, timer_delete()

## timer\_delete()

NAME	<b>timer_delete( )</b> – remove a previously created timer (POSIX)
SYNOPSIS	STATUS timer_delete ( timer_t timerid /* timer ID */ )
DESCRIPTION	This routine removes a timer.
RETURNS	0 ( <b>OK</b> ), or -1 ( <b>ERROR</b> ) if <i>timerid</i> is invalid.
ERRNO	EINVAL
SEE ALSO	timerLib, timer_create( )

# timer\_getoverrun()

NAME	<b>timer_getoverrun()</b> – return the timer expiration overrun (POSIX)
SYNOPSIS	<pre>int timer_getoverrun   (    timer_t timerid /* timer ID */ )</pre>

DESCRIPTION	This routine returns the timer expiration overrun count for <i>timerid</i> , when called from a timer expiration signal catcher. The overrun count is the number of extra timer expirations that have occurred, up to the implementation-defined maximum <b>DELAYTIMER_MAX</b> . If the count is greater than the maximum, it returns the maximum.
RETURNS	The number of overruns, or <b>DELAYTIMER_MAX</b> if the count equals or is greater than <b>DELAYTIMER_MAX</b> , or -1 ( <b>ERROR</b> ) if <i>timerid</i> is invalid.
ERRNO	EINVAL ENOSYS
SEE ALSO	timerLib

## timer\_gettime()

NAME	<b>timer_gettime()</b> – get the remaining time before expiration and the reload value (POSIX)
SYNOPSIS	<pre>int timer_gettime   (    timer_t timerid, /* timer ID */    struct itimerspec * value /* where to return remaining time */   )</pre>
DESCRIPTION	This routine gets the remaining time and reload value of a specified timer. Both values are copied to the <i>value</i> structure.
RETURNS	0 ( <b>OK</b> ), or -1 ( <b>ERROR</b> ) if <i>timerid</i> is invalid.
ERRNO	EINVAL
SEE ALSO	timerLib

## timer\_modify()

NAME timer\_modify() - modify a timer
SYNOPSIS STATUS timer\_modify
(
timer\_t timerId, /\* timer ID \*/

	struct sigevent * pSigev /* sigevent describing the notification */ )
DESCRIPTION	This routine updates the timer <i>timerId</i> with the new notification mechanism as indicated by $pSigev$ . This routine should be called in the context of RTP task only, that is, as a system call. This routine should be called with the timer disarmed.
NOTE	This is a non-POSIX API.
RETURNS	<b>ERROR</b> if the <i>timerId</i> is invalid or armed, or if there is an error in the notification mechanisim. Otherwise <b>OK</b> .
ERRNO	EINVAL
SEE ALSO	timerLib

## timer\_open()

NAME	timer_open() – open a timer					
SYNOPSIS	<pre>timer_t timer_open   (     const char * name, /* name of timer */     int mode, /* OM_CREATE, */     clockid_t clockId, /* clock ID */     struct sigevent * evp, /* user event handler */     void * context /* context value */   )</pre>					
DESCRIPTION	This routine opens a timer, which means that it will search the name space and will return the timer_id of an existent timer with same name as <i>name</i> , and if none is found, then creates a new one with that name depending on the flags set in the mode parameter. Note that there are two name spaces available to the calling routine in which <b>timer_open()</b> can perform the search, and which are selected depending on the first character in the <i>name</i> parameter. When this character is a forward slash <i>I</i> , the <b>public</b> name space is searched; otherwise the <b>private</b> name space is searched. Similarly, if a timer is created, the first character in <i>name</i> specifies the name space that contains the timer.					
	The argument <i>name</i> is mandatory. NULL or empty strings are not allowed.					
	Timers created by this routine can not be deleted with <b>timer_delete()</b> . Instead, a <b>timer close()</b> must be issued for every <b>timer open()</b> . Then the timer is deleted when it is					

**timer\_close()** must be issued for every **timer\_open()**. Then the timer is deleted when it is removed from the name space by a call to **timer\_unlink()**. Alternatively, the timer can be previously removed from the name space, and deleted during the last **timer\_close()**.

A description of the *mode* and *context* arguments follows. See the reference entry for **timer\_create()** for a description of the remaining arguments.

mode

This parameter specifies the timer permissions (not implemented) along with various object management attribute bits as follows:

#### OM\_CREATE

Create a new timer if a matching timer name is not found.

#### OM\_EXCL

When set jointly with **OM\_CREATE**, create a new timer immediately without attempting to open an existing timer. An error condition is returned if a timer with *name* already exists. This attribute has no effect if the **OM\_CREATE** attribute is not specified.

#### OM\_DELETE\_ON\_LAST\_CLOSE

Only used when a timer is created. If set, the timer will be deleted during the last **timer\_close()** call, independently on whether **timer\_unlink()** was previously called or not.

#### context

Context value assigned to the created timer. This value is not actually used by VxWorks. Instead, the context value can be used by OS extensions to implement object permissions, for example.

The *clockId* and *evp* are used only when creating a new timer. The clock used by the timer *clockId* is the one defined in *time.h*. The *evp* argument, if non-NULL, points to a **sigevent** structure, which is allocated by the application and defines the signal number and application-specific data to be sent to the task when the timer expires. If *evp* is NULL, a default signal (SIGALRM) is queued to the task, and the signal data is set to the timer ID. Initially, the timer is disarmed.

**NOTE** This is a non-POSIX API. This routine is not ISR callable.

#### SMP CONSIDERATIONS

ERRNO

This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the
implementation so it is the responsibility of the caller to ensure they are complied with.
Future implementations may enforce these restrictions.

**RETURNS** The timer ID on success; otherwise **ERROR**.

EINVAL EAGAIN

#### S\_intLib\_NOT\_ISR\_CALLABLE

This routine must not be called from an ISR.

```
SEE ALSO timerOpen, timer_close(), timer_unlink()
```

	timer_settime()			
NAME	<b>timer_settime()</b> – set the time until the next expiration and arm timer (POSIX)			
SYNOPSIS	<pre>int timer_settime   (   timer_t timerid, /* timer ID */   int flags, /* absolute or relative */   const struct itimerspec * value, /* time to be set */   struct itimerspec * ovalue /* prev time set (NULL=no result) */   )</pre>			
DESCRIPTION	This routine sets the next expiration of the timer, using the <b>.it_value</b> of <i>value</i> , thus arming the timer. If the timer is already armed, this call resets the time until the next expiration. If <b>.it_value</b> is zero, the timer is disarmed.			
	If <i>flags</i> is not equal to <b>TIMER_ABSTIME</b> , the interval is relative to the current time, the interval being the <b>.it_value</b> of the <i>value</i> parameter. If <i>flags</i> is equal to <b>TIMER_ABSTIME</b> , the expiration is set to the difference between the absolute time of <b>.it_value</b> and the current value of the clock associated with <i>timerid</i> . If the time has already passed, then the timer expiration notification is made immediately with the signal delivered to the task that created the timer.			
	The reload value of the timer is set to the value specified by the <b>.it_interval</b> field of <i>value</i> . When a timer is armed with a nonzero <b>.it_interval</b> a periodic timer is set up.			
	Time values that are between two consecutive non-negative integer multiples of the resolution of the specified timer are rounded up to the larger multiple of the resolution.			
	If <i>ovalue</i> is non-NULL, the routine stores a value representing the previous amount of time before the timer would have expired. Or if the timer is disarmed, the routine stores zero, together with the previous timer reload value. The <i>ovalue</i> parameter is the same value as that returned by <b>timer_gettime()</b> and is subject to the timer resolution.			
WARNING	If <b>clock_settime()</b> is called to reset the absolute clock time after a timer has been set with <b>timer_settime()</b> , and if <i>flags</i> is equal to <b>TIMER_ABSTIME</b> , then the timer will behave unpredictably. If you must reset the absolute clock time after setting a timer, do not use <i>flags</i> equal to <b>TIMER_ABSTIME</b> .			
RETURNS	0 ( <b>OK</b> ), or -1 ( <b>ERROR</b> ) if <i>timerid</i> is invalid, the number of nanoseconds specified by <i>value</i> is less than 0 or greater than or equal to 1,000,000,000, or the time specified by <i>value</i> exceeds the maximum allowed by the timer.			
ERRNO	EINVAL			
SEE ALSO	timerLib			

## timer\_show()

NAME	<b>timer_show()</b> – show information on a specified timer			
SYNOPSIS	<pre>int timer_show   (    (    timer_t timerid, /* timer ID */    int verbose /* Verbose mode: 0, 1 */   )</pre>			
DESCRIPTION	This routine shows information about the timer specified in <i>timerid</i> . If <i>timerid</i> is 0 then a list of all timers will be printed. Verbose mode will show additional information about <i>timerid</i> including the owner's task name, the timers type, and the state of the timer.			
EXAMPLE	->timer_show (0,0)			
	timerid taskId evp routine arg Remaining Period			
	0x604ce330 1617182668 0.000000 0.000000 0x603b56b8 0x6170bba0 0x603b570c 0x604ce330 1617182668 0.000000 0.000000 0x617115e0 0x6170bba0 0x61711634 0x604ce330 1617182668 0.000000 0.000000 0x617115e0 0x6170bba0 0x61711634 0x604ce330 1617182668 0.000000 0.000000			
->timer_show (0x6170bf20,0) timerid taskId evp routine arg Remaining Period 				
->timer_show (0x6170bf20,1)				
	timerid taskId evp routine arg Remaining Period			
	0x604ce330 1617182668 0.000000 0.000000 Owners Task Name: tTimer_gettimeTest3 Type of Timer: CLOCK_REALTIME State: Active			
WARNING	This is a non-POSIX API.			
RETURNS	0 (OK), or -1 (ERROR) if <i>timerid</i> is invalid, or the context is invalid.			
ERRNO	N/A			
SEE ALSO	timerShow			

timer unlink()

	timer_ummk()			
NAME	timer_unlink() – unlink a named timer			
SYNOPSIS	STATUS timer_unlink ( const char * name /* name of the timer to unlink */ )			
DESCRIPTION	This routine removes a timer from the name space, and marks it as ready for deletion on the last <b>timer_close()</b> . In case there are already no outstanding <b>timer_open()</b> calls, the timer is deleted. After a timer is unlinked, subsequent calls to <b>timer_open()</b> using <i>name</i> will not be able to find the timer, even if it has not been deleted yet. Instead, a new timer could be created if <b>timer_open()</b> is called with the <b>OM_CREATE</b> flag.			
NOTE	This is a non-POSIX API. This routine is not ISR callable.			
SMP CONSIDERATIONS				
	This API is spinlock and intCpuLock restricted. These restrictions are not enforced by the implementation so it is the responsibility of the caller to ensure they are complied with. Future implementations may enforce these restrictions.			
RETURNS	OK, or ERROR if unsuccessful			
ERRNO	S_objLib_OBJ_INVALID_ARGUMENT <i>name</i> is NULL or empty.			
	S_objLib_OBJ_NOT_FOUND No timer with <i>name</i> was found.			
	S_objLib_OBJ_DESTROY_ERROR Error while deleting the timer.			
	S_intLib_NOT_ISR_CALLABLE This routine must not be called from an ISR.			
SEE ALSO	timerOpen, timer_open(), timer_close()			

## timex()

NAME timex() – time a single execution of a function or functions

SYNOPSIS void timex

	<pre>( FUNCPTR func, /* function to time (optional) */ int arg1, /* first of up to 8 args to call function with (opt) */ int arg2, int arg3, int arg4, int arg5, int arg6, int arg7, int arg8 )</pre>
DESCRIPTION	This routine times a single execution of a specified function with up to eight of the function's arguments. If no function is specified, it times the execution of the current list of functions to be timed, which is created using <b>timexFunc()</b> , <b>timexPre()</b> , and <b>timexPost()</b> . If <b>timex()</b> is executed with a function argument, the entire current list is replaced with the single specified function.
	When execution is complete, <b>timex()</b> displays the execution time. If the execution was so fast relative to the clock rate that the time is meaningless (error > 50%), a warning message is printed instead. In such cases, use <b>timexN()</b> .
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	<pre>timexLib, timexFunc(), timexPre(), timexPost(), timexN()</pre>

# timexClear()

SEE ALSO	timexLib
ERRNO	Not Available
RETURNS	N/A
DESCRIPTION	This routine clears the current list of functions to be timed.
SYNOPSIS	void timexClear (void)
NAME	<b>timexClear()</b> – clear the list of function calls to be timed

## timexFunc()

NAME

timexFunc() – specify functions to be timed

SYNOPSIS	void timexFu	inc							
	( int FUNCPTR int int int int int int int )	i, func, arg1, arg2, arg3, arg4, arg5, arg6, arg7, arg8	/* functio /* functio /* first o	n number n to be f up to	in list added (N 8 args t	z (03) NULL if to call	to be functi	deleted on with	*/ .) */ . */
DESCRIPTION	This routine a to <b>timex()</b> or specifies the f deleted by sp	dds or d timexN unction' ecifying	leletes funct ( ). Up to fo s position in its sequence	ions in the ur function the seque number i	e list of fu ns can be ence of ex i and <b>NU</b>	inctions include cecution LL for th	to be tin ed in the (0, 1, 2, le functi	ned as a e list. Th or 3). A on argu	group by calls e argument <i>i</i> function is ment <i>func</i> .
RETURNS	N/A								
ERRNO	Not Available								
SEE ALSO	timexLib, tim	nex(), tir	nexN()						

## timexHelp()

NAME	timexHelp() – display synopsis of execution timer facilities			
SYNOPSIS	void timexHelp (void)			
DESCRIPTION	This routine displ timexHelp timex timexN timexClear timexFunc timexPre	<pre>ays the following sum [func, [args]] [func, [args]] i, func, [args] i, func, [args]</pre>	<pre>mary of the available execution timer functions:     Print this list.     Time a single execution.     Time repeated executions.     Clear all functions.     Add timed function number i (0,1,2,3).     Add pre-timing function number i.</pre>	
	timexPost timexShow	i,func,[args]	Add post-timing function number i. Show all functions to be called.	

	Notes:
	<ol> <li>timexN() will repeat calls enough times to get timing accuracy to approximately 2%.</li> <li>A single function can be specified with timex() and timexN(); or, multiple functions can be pre-set with timexFunc().</li> <li>Up to 4 functions can be pre-set with timexFunc(), timexPre() and timexPost() i.e. i in the range 0 = 3</li> </ol>
	<ul> <li>4) timexPre() and timexPost() allow locking/unlocking, or raising/lowering priority before/after timing.</li> </ul>
RETURNS	N/A
ERRNO	Not Available

SEE ALSO timexLib

### timexInit()

**NAME** timexInit() – include the execution timer library

- SYNOPSIS void timexInit (void)
- DESCRIPTION This null routine is provided so that timexLib can be linked into the system. If the configuration macro INCLUDE\_TIMEX is defined, it is called by the root task, usrRoot(), in usrConfig.c.

RETURNS N/A

ERRNO Not Available

SEE ALSO timexLib

### timexN()

NAME timexN() – time repeated executions of a function or group of functions

SYNOPSIS void timexN
(
FUNCPTR func, /\* function to time (optional) \*/
int arg1, /\* first of up to 8 args to call function with \*/

VxWorks Kernel API Reference, 6.6 timexPost()

int arg2, int arg3, int arg4, int arg5, int arg6, int arg7, int arg8)

DESCRIPTION	This routine times the execution of the current list of functions to be timed in the same				
	manner as timex(); however, the list of functions is called a variable number of times until				
	sufficient resolution is achieved to establish the time with an error less than 2%. (Since each				
	iteration of the list may be measured to a resolution of +/- 1 clock tick, repetitive timings				
	decrease this error to $1/N$ ticks, where N is the number of repetitions.)				
RETURNS	N/A				

ERRNO	Not Available
SEE ALSO	<pre>timexLib, timexFunc( ), timex( )</pre>

### timexPost()

NAME	timexPost() – specify functions to be called after timing
SYNOPSIS	<pre>void timexPost   (    int i, /* function number in list (03) */ FUNCPTR func, /* function to be added (NULL if to be deleted) */    int arg1, /* first of up to 8 args to call function with */    int arg3,    int arg4,    int arg5,    int arg6,    int arg7,    int arg8   )</pre>
DESCRIPTION	This routine adds or deletes functions in the list of functions to be called immediately following the timed functions. A maximum of four functions may be included. Up to eight arguments may be passed to each function.
RETURNS	N/A
ERRNO	Not Available

SEE ALSO timexLib

### timexPre()

NAME	<b>timexPre( )</b> – specify functions to be called prior to timing
SYNOPSIS	<pre>void timexPre   (     int i, /* function number in list (03) */   FUNCPTR func, /* function to be added (NULL if to be deleted) */     int arg1, /* first of up to 8 args to call function with */     int arg3,     int arg4,     int arg5,     int arg6,     int arg7,     int arg8   )</pre>
DESCRIPTION	This routine adds or deletes functions in the list of functions to be called immediately prior to the timed functions. A maximum of four functions may be included. Up to eight arguments may be passed to each function.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	timexLib

### timexShow()

NAME timexShow() – display the list of function calls to be timed

SYNOPSIS void timexShow (void)

**DESCRIPTION** This routine displays the current list of function calls to be timed. These lists are created by calls to **timexPre()**, **timexFunc()**, and **timexPost()**.

RETURNS N/A

VxWorks Kernel API Reference, 6.6 tlsTaskInit()

ERRNO Not Available

SEE ALSO timexLib, timexPre(), timexFunc(), timexPost()

### tlsTaskInit()

NAME tlsTaskInit() – Thread Local Storage init routine

SYNOPSIS STATUS tlsTaskInit (void)

**DESCRIPTION** This routine initializes the current task thread local storage for all available modules. Once this is called, any access of a \_\_thread variable of any of those modules will be deterministic. If this routine is not called, access to the \_\_thread variable of a module may take longer the first time because of time needed to allocate memory to manage the module's \_\_thread variables.

**RETURNS** OK, or ERROR

ERRNO N/A

SEE ALSO tlsLib

### tr( )

NAME	tr() – resume a task
SYNOPSIS	<pre>void tr   (   int taskNameOrId /* task name or task ID */ )</pre>
DESCRIPTION	This command resumes the execution of a suspended task. It simply calls <b>taskResume()</b> .
RETURNS	N/A
ERRNO	N/A
SEE ALSO	usrLib, ts(), taskResume(), the VxWorks programmer guides.

### traceTmrResolutionGet()

NAME	<b>traceTmrResolutionGet()</b> – get resolution of timestamp source, in nanoseconds
SYNOPSIS	STATUS traceTmrResolutionGet ( struct timespec * pTimestamp /* destination for resolution */ )
DESCRIPTION	Write the resolution of the timestamp source into the supplied timespec. If the timestamp resolution cannot be obtained (maybe the timestamp driver is not available) then the timespec structure will be filled with zeroes.
RETURNS	OK, or ERROR.
ERRNO	Not Available
SEE ALSO	wvTmrLib

## transCommit()

NAME	transCommit() – externally-callable function to do a commit
SYNOPSIS	<pre>int transCommit   (    TRANS_XBD * pDev )</pre>
DESCRIPTION	The $pDev$ argument is the device handle (as provided to the warning and panic hooks).
	This function is only for special cases (such as "automatic commits" on a <b>WARN_OUT_OF_UNITS</b> warning, which are generally advised against for data-control reasons).
RETURNS	error number, or 0 ( <b>OK</b> ) on success.
ERRNO	N/A
SEE ALSO	xbdTrans, usrTransCommit( ), usrTransCommitFd( ).

VxWorks Kernel API Reference, 6.6 transDevCreate()

### transDevCreate()

NAME	transDevCreate() – create a transactional XBD.
SYNOPSIS	TRANS_XBD *transDevCreate ( device_t subDev /* lower level device */ )
DESCRIPTION	THIS ROUTINE IS EXTERNALLY-VISIBLE FOR BACKWARDS COMPATIBILITY ONLY.
	This is essentially an internal version that gives you the pointer to the device. The pointer becomes invalid when the device is ejected, so one should use xbdTransDevCreate instead (to get a device_t that can be checked).
RETURNS	TRANS_XBD *, or NULL on failure.
ERRNO	Not Available
SEE ALSO	xbdTrans

# trgAdd()

NAME	trgAdd() – add	a new trigger to the trigger list
SYNOPSIS	<pre>TRIGGER_ID try (     event_t     int     int     UINT32     OBJ_ID     int     int     int     int     TRIGGER *     int     FUNCPTR     BOOL     int     )</pre>	pAdd event, status, contextType, contextId, objId, conditional, condType, condEx1, condOp, condEx2, disable, chain, actionType, actionFunc, actionArg

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**DESCRIPTION** This routine creates a new trigger and adds it to the proper trigger list. It takes the following parameters:

#### event

as defined in eventP.h for System Viewer, if given.

#### status

the initial status of the trigger (enabled or disabled).

#### contextType

the type of context where the event occurs.

#### contextId

the ID (if any) of the context where the event occurs.

#### objectId

if given and applicable.

#### conditional

the indicator that there is a condition on the trigger.

#### condType

the indicator that the condition is either a variable or a function.

#### condEx1

the first element in the comparison.

#### condOp

the type of operator (==, !=, , <=, , >=, |, &).

#### condEx2

the second element in the comparison (a constant).

#### disable

the indicator of whether the trigger must be disabled once it is hit.

#### chain

a pointer to another trigger associated to this one (if any).

#### actionType

the type of action associated with the trigger (none, func, lib).

#### actionFunc

the action associated with the trigger (the function).

#### actionDef

the indicator of whether the action can be deferred (deferred is the default).

#### actionArg

the argument passed to the function, if any.

Attempting to call trgAdd whilst triggering is enabled is not allowed and will return NULL.

VxWorks Kernel API Reference, 6.6 trgChainSet()

 RETURNS
 TRIGGER\_ID, or NULL if either the trigger ID can not be allocated, or if called whilst triggering is enabled.

 ERRNO
 S\_intLib\_NOT\_ISR\_CALLABLE

 S\_objLib\_OBJ\_ID\_ERROR

 S\_memLib\_NOT\_ENOUGH\_MEMORY

 S\_memLib\_BLOCK\_ERROR

SEE ALSO trgLib, trgDelete()

### trgChainSet()

S\_taskLib\_ILLEGAL\_PRIORITY

NAME	<b>trgChainSet( )</b> – chains two triggers
SYNOPSIS	STATUS trgChainSet ( TRIGGER_ID fromId, TRIGGER_ID toId )
DESCRIPTION	This routine chains two triggers together. When the first trigger fires, it calls <b>trgEnable()</b> for the second trigger. The second trigger must be created disabled in order to maintain the correct sequence.
RETURNS	OK or ERROR.
ERRNO	
SEE ALSO	trgLib, trgEnable()

### trgDelete()

NAME	<b>trgDelete( )</b> – delete a trigger from the trigger list
SYNOPSIS	STATUS trgDelete ( TRIGGER_ID trgId
	)

DESCRIPTIONThis routine deletes a trigger by removing it from the trigger list. It also checks that no other<br/>triggers are still active. If there are no active triggers and triggering is still on, it turns<br/>triggering off.RETURNSOK, or ERROR if the trigger is not found.ERRNOS\_objLib\_OBJ\_ID\_ERROR<br/>trgLib, trgAdd()

### trgDisable()

NAME	trgDisable() – turn a trigger off
SYNOPSIS	STATUS trgDisable ( TRIGGER_ID trgId )
DESCRIPTION	This routine disables a trigger. It also checks to see if there are triggers still active. If this is the last active trigger it sets triggering off.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the trigger ID is not found.
ERRNO	
SEE ALSO	trgLib, trgEnable()

### trgEnable()

NAME	trgEnable() – enable a trigger
SYNOPSIS	STATUS trgEnable ( TRIGGER_ID trgId )
DESCRIPTION	This routine enables a trigger that has been created with <b>trgAdd()</b> . A counter is incremented to keep track of the total number of enabled triggers so that <b>trgDisable()</b> knows when to

VxWorks Kernel API Reference, 6.6 trgEvent()

	set triggering off. If the maximum number of enabled triggers is reached, an error is returned.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the trigger ID is not found or if the maximum number of triggers has already been enabled.
ERRNO	
SEE ALSO	trgLib, trgDisable()

## trgEvent()

NAME	trgEvent() – trigger a user-defined event
SYNOPSIS	<pre>void trgEvent   (    event_t evtId /* event */   )</pre>
DESCRIPTION	This routine triggers a user event. A trigger must exist and triggering must have been started with <b>trgOn()</b> or from the triggering GUI to use this routine. The <i>evtId</i> should be in the range 40000-65535.
RETURNS	N/A
ERRNO	
SEE ALSO	trgLib, dbgLib, e()

# trgLibInit()

NAME	<b>trgLibInit()</b> – initialize the triggering library
SYNOPSIS	STATUS trgLibInit (void)
DESCRIPTION	This routine initializes the trigger class. Triggers are VxWorks objects and therefore require a class to be initialized.
RETURNS	OK or ERROR.

ERRNO

SEE ALSO trgLib

# trgOff()

NAME	<b>trgOff()</b> – set triggering off
SYNOPSIS	void trgOff (void)
DESCRIPTION	This routine turns triggering off. From this time on, when an event point is hit, no search on triggers is performed.
RETURNS	N/A
ERRNO	
SEE ALSO	trgLib, trgOn()

# trgOn()

NAME	<b>trgOn()</b> – set triggering on
SYNOPSIS	STATUS trgOn (void)
DESCRIPTION	This routine activates triggering. From this time on, any time an event point is hit, a check for the presence of possible triggers is performed. Start triggering only when needed since some overhead is introduced.
NOTE	If <b>trgOn()</b> is called when there are no triggers in the trigger list, it immediately sets triggering off again. If <b>trgOn()</b> is called with at least one trigger in the list, triggering begins. Triggers should not be added to the list while triggering is on since this can create instability.
RETURNS	OK or ERROR.
ERRNO	
SEE ALSO	trgLib, trgOff()

# trgReset()

NAME	<b>trgReset()</b> – Reset a trigger in the trigger list
SYNOPSIS	STATUS trgReset ( TRIGGER_ID trgId )
DESCRIPTION	This routine resets a trigger. It sets the triggers hit count to zero and sets its state back to the initial state saved when the trigger was downloaded
RETURNS	<b>OK</b> , or <b>ERROR</b> if the trigger is not found.
ERRNO	S_objLib_OBJ_ID_ERROR
SEE ALSO	trgLib

# trgShow()

NAME	<b>trgShow()</b> – show trigger information
SYNOPSIS	STATUS trgShow ( TRIGGER_ID trgId, /* trigger id to show, or NULL for all triggers */ int level /* detail level: 1 gives more detail */ )
DESCRIPTION	This routine displays trigger information. If <i>trgId</i> is passed, only the summary for that trigger is displayed. If no parameter is passed, the list of existing triggers is displayed with a summary of their state. For example:
	Oxffedfc disabled 101 3 Ox14e7a4 Y Oxffe088
	0xffe088 enabled 55 1 0x10db58 Y 0x0
	If <i>level</i> is 1, then more detailed information is displayed.
EXAMPLE	-> trgShow trgId, 1
RETURNS	OK.
ERRNO	Not Available

SEE ALSO trgShow, trgLib

# trgShowInit()

NAME	<b>trgShowInit()</b> – initialize the trigger show facility
SYNOPSIS	void trgShowInit (void)
DESCRIPTION	This routine links the trigger show facility into the VxWorks system. These routines are included automatically when INCLUDE_TRIGGER_SHOW is defined.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	trgShow

### trgWorkQReset()

**NAME** trgWorkQReset() – Resets the trigger work queue task and queue

SYNOPSIS STATUS trgWorkQReset (void)

**DESCRIPTION** When a trigger fires, if the assocated action requires a function to be called in "safe" mode, a pointer to the required function will be placed on a queue known as the "triggering work queue". A system task "tActDef" is spawned to action these requests at task level. Should the user have need to reset this work queue (e.g. if a called task causes an exception which causes the trgActDef task to be SUSPENDED, or if the queue gets out of sync and becomes unresponsive), **trgWorkQReset()** may be called.

Its effect is to delete the trigger work queue task and its associated resources and then recreate them. Any entries pending on the triggering work queue will be lost. Calling this function with triggering on will result in triggering being turned off before the queue reset takes place. It is the responsibility of the user to turn triggering back on.

This function may not be called from interrupt.

**RETURNS** OK, or ERROR if the triggering task and its associated resources cannot be deleted and recreated.

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ERRNO S\_taskLib\_NAME\_NOT\_FOUND S\_taskLib\_ILLEGAL\_PRIORITY S\_intLib\_NOT\_ISR\_CALLABLE S\_objLib\_OBJ\_ID\_ERROR S\_memLib\_NOT\_ENOUGH\_MEMORY S\_memLib\_BLOCK\_ERROR

```
SEE ALSO trgLib
```

## trunc()

NAME	<b>trunc()</b> – truncate to integer
SYNOPSIS	<pre>double trunc   (     double x /* value to truncate */ )</pre>
DESCRIPTION	This routine discards the fractional part of a double-precision value $x$ .
RETURNS	The integer portion of $x$ , represented in double-precision.
ERRNO	Not Available
SEE ALSO	mathALib

### truncf()

NAME	<b>truncf()</b> – truncate to integer
SYNOPSIS	<pre>float truncf   (    float x /* value to truncate */   )</pre>
DESCRIPTION	This routine discards the fractional part of a single-precision value $x$ .
RETURNS	The integer portion of $x$ , represented in single precision.
ERRNO	Not Available

SEE ALSO mathALib

## ts()

NAME	ts() – suspend a task
SYNOPSIS	<pre>void ts   (    int taskNameOrId /* task name or task ID */ )</pre>
DESCRIPTION	This command suspends the execution of a specified task. It simply calls <b>taskSuspend()</b> .
RETURNS	N/A
ERRNO	N/A
SEE ALSO	usrLib, tr(), taskSuspend(), the VxWorks programmer guides.

# tsecRegister()

SEE ALSO	tsecVxbEnd
ERRNO	N/A
RETURNS	N/A
DESCRIPTION	This routine registers the TSEC driver with VxBus as a child of the PLB bus type.
SYNOPSIS	void tsecRegister(void)
NAME	<b>tsecRegister()</b> – register with the VxBus subsystem

# tt()

NAME	tt() – display a stack trace of a task
SYNOPSIS	STATUS tt ( int taskNameOrId /* task name or task ID */ )
DESCRIPTION	This routine displays a list of the nested routine calls that the specified task is in. Each routine call and its parameters are shown.
	If <i>taskNameOrId</i> is not specified or zero, the last task referenced is assumed. The <b>tt()</b> routine can only trace the stack of a task other than itself and unbreakable tasks. For instance, when <b>tt()</b> is called from the shell, it cannot trace the shell's stack.
EXAMPLE	<pre>-&gt; tt "tAioIoTask1"</pre>
	This indicates that <b>tAioIoTask1()</b> is currently in <b>semTake()</b> (with two parameters) and was called by <b>aioIoTask()</b> (with five parameters).
CAVEAT	In order to do the trace, some assumptions are made. In general, the trace will work for all C language routines and for assembly language routines. Depending of the architecture and at which point the task is suspended, the trace facility may produce inaccurate results or fail completely. Moreover, if the routine is written in a language other than C, the routine's entry point is non-standard, or the task's stack is corrupted, the trace facility may produce inaccurate results too. Also, all parameters are assumed to be 32-bit quantities, so structures passed as parameters will be displayed as <i>long</i> integers.
RETURNS	OK, or ERROR if the task does not exist.
ERRNO	N/A
SEE ALSO	<b>dbgLib</b> , the VxWorks programmer guides, <i>VxWorks Command-Line Tools User's Guide</i> .

	ttyDevCreate()
NAME	<b>ttyDevCreate( )</b> – create a VxWorks device for a serial channel
SYNOPSIS	<pre>STATUS ttyDevCreate   (     char * name, /* name to use for this device */     SIO_CHAN * pSioChan, /* pointer to core driver structure */     int rdBufSize, /* read buffer size, in bytes */     int wrtBufSize /* write buffer size, in bytes */   )</pre>
DESCRIPTION	This routine creates a device on a specified serial channel. Each channel to be used should have exactly one device associated with it by calling this routine.
	For instance, to create the device "/tyCo/0", with buffer sizes of 512 bytes, the proper call would be:
	<pre>ttyDevCreate ("/tyCo/0", pSioChan, 512, 512);</pre>
	Where <i>pSioChan</i> is the address of the underlying <b>SIO_CHAN</b> serial channel descriptor (defined in <b>sioLib.h</b> ). This routine is called automatically when <b>INCLUDE_TTY_DEV</b> is configured in VxWorks. It initializes two channels using the default names /tyCo/0 and /tyCo/1.
RETURNS	OK, or ERROR if the driver is not installed, or the device already exists.
ERRNO	S_ioLib_NO_DRIVER (ENXIO) The <b>ttyDrv</b> driver is not installed in system.
	S_iosLib_DUPLICATE_DEVICE_NAME (EINVAL) Device name already in use.
SEE ALSO	ttyDrv

## ttyDrv()

NAME	<b>ttyDrv()</b> – initialize the <i>tty</i> driver
SYNOPSIS	STATUS ttyDrv (void)

**DESCRIPTION** This routine initializes the *tty* driver, which is the OS interface to core serial channel(s). It is called automatically when **INCLUDE\_TTY\_DEV** is configured in VxWorks.

	After this routine is called, <b>ttyDevCreate()</b> is typically called to bind serial channels to VxWorks devices.
RETURNS	OK, or ERROR if the driver cannot be installed.
ERRNO	N/A
SEE ALSO	ttyDrv
	tw()
NAME	<b>tw()</b> – print info about the object the given task is pending on
SYNOPSIS	<pre>void tw   (    int taskNameOrId /* task name or task ID */ )</pre>
DESCRIPTION	This routine shows task's pending information of the task <i>taskNameOrId</i> .
	This routine doesn't support POSIX semaphores and message queues. This command doesn't support pending signals.
	List of object types that are recognized:
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	usrLib, w(), the VxWorks programmer guides.

# tyAbortFuncSet()

NAME	<b>tyAbortFuncSet( )</b> – set the abort function			
SYNOPSIS	void tyAbortFuncSet ( FUNCPTR func /* routine to call when abort char received */ )			

DESCRIPTION	This routine sets the function that will be called when the abort character is received on a <i>tty</i> . There is only one global abort function, used for any <i>tty</i> on which <b>OPT_ABORT</b> is enabled. When the abort character is received from a <i>tty</i> with <b>OPT_ABORT</b> set, the function specified in <i>func</i> will be called, with no parameters, from interrupt level.
	Setting an abort function of NULL will disable the abort function.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	tyLib, tyAbortSet()

# tyAbortGet()

NAME	<b>tyAbortGet( )</b> – get the abort character
SYNOPSIS	char tyAbortGet (void)
DESCRIPTION	This routine returns the abort character.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	tyLib, tyAbortFuncSet( ), tyAbortSet( )

# tyAbortSet()

NAME	<b>tyAbortSet( )</b> – change the abort character
SYNOPSIS	<pre>void tyAbortSet   (    char ch /* char to be abort */ )</pre>
DESCRIPTION	This routine sets the abort character to <i>ch</i> . The default abort character is CTRL-C.

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Typing the abort character to any device whose **OPT\_ABORT** option is set will cause the shell task to be killed and restarted. Note that the character set by this routine applies to all devices whose handlers use the standard *tty* package **tyLib**.

SEE ALSO	tyLib, tyAbortFuncSet( ), tyAbortGet( )
ERRNO	N/A
RETURNS	N/A

### tyBackspaceSet()

NAME	<b>tyBackspaceSet()</b> – change the backspace character
SYNOPSIS	<pre>void tyBackspaceSet   (    char ch /* char to be backspace */ )</pre>
DESCRIPTION	This routine sets the backspace character to <i>ch</i> . The default backspace character is CTRL-H.
	Typing the backspace character to any device operating in line protocol mode ( <b>OPT_LINE</b> set) will cause the previous character typed to be deleted, up to the beginning of the current line. Note that the character set by this routine applies to all devices whose handlers use the standard <i>tty</i> package <b>tyLib</b> .
RETURNS	N/A
ERRNO	N/A
SEE ALSO	tyLib

## tyDeleteLineSet()

NAME	<b>tyDeleteLineSet( )</b> – change the line-delete character
SYNOPSIS	void tyDeleteLineSet
	, char ch /* char to be line-delete */

)

DESCRIPTION	This routine sets the line-delete character to <i>ch</i> . The default line-delete character is CTRL-U. Typing the delete character to any device operating in line protocol mode ( <b>OPT_LINE</b> set) will cause all characters in the current line to be deleted. Note that the character set by this routine applies to all devices whose handlers use the standard <i>tty</i> package <b>tyLib</b> .
RETURNS	N/A
ERRNO	N/A
SEE ALSO	tyLib

# tyDevInit()

NAME	<b>DevInit()</b> – initialize the <i>tty</i> device descriptor	
SYNOPSIS	TATUS tyDevInit ( FAST TY_DEV_ID pTyDev, /* ptr to tty dev descriptor to init */ int rdBufSize, /* size of read buffer in bytes */ int wrtBufSize, /* size of write buffer in bytes */ FUNCPTR txStartup /* device transmit start-up routine */ )	
DESCRIPTION	This routine initializes a <i>tty</i> device descriptor according to the specified parameters. The initialization includes allocating read and write buffers of the specified sizes from the memory pool, and initializing their respective buffer descriptors. The semaphores are initialized and the write semaphore is given to enable writers. Also, the transmitter start-routine pointer is set to the specified routine. All other fields in the descriptor are zeroed	
	his routine should be called only by serial drivers.	
RETURNS	<b>PK</b> , or <b>ERROR</b> if there is not enough memory to allocate data structures.	
ERRNO	J/A	
SEE ALSO	yLib	

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# tyDevRemove()

NAME	<b>tyDevRemove()</b> – remove the <i>tty</i> device descriptor
SYNOPSIS	STATUS tyDevRemove ( TY_DEV_ID pTyDev /* ptr to tty dev descriptor to remove */ )
DESCRIPTION	This routine removes an existing <i>tty</i> device descriptor. It releases the read and write buffers and the descriptor data structure.
RETURNS	OK, or ERROR if expected data structures are not found
ERRNO	N/A.
SEE ALSO	tyLib

# tyDevTerminate()

NAME	<b>tyDevTerminate()</b> – terminate the <i>tty</i> device descriptor
SYNOPSIS	STATUS tyDevTerminate ( TY_DEV_ID pTyDev /* ptr to tty dev descriptor to terminate */ )
DESCRIPTION	This routine terminates a <i>tty</i> device descriptor. The termination includes freeing memory for the read and write buffers, and terminating the various semaphores. This routine should be called only by serial drivers.
RETURNS	ОК
ERRNO	Not Available
SEE ALSO	tyLib
# tyEOFGet()

SEE ALSO	tyLib, tyEOFSet()
ERRNO	N/A
RETURNS	N/A
DESCRIPTION	This routine returns the current end-of-file character.
SYNOPSIS	char tyEOFGet (void)
NAME	<b>tyEOFGet()</b> – get the current end-of-file character

# tyEOFSet()

NAME	<b>tyEOFSet()</b> – change the end-of-file character
SYNOPSIS	void tyEOFSet ( char ch /* char to be EOF */ )
DESCRIPTION	This routine sets the EOF character to <i>ch</i> . The default EOF character is CTRL-D.
	Typing the <b>EOF</b> character to any device operating in line protocol mode ( <b>OPT_LINE</b> set) will cause no character to be entered in the current line, but will cause the current line to be terminated (thus without a newline character). The line is made available to reading tasks. Thus, if the <b>EOF</b> character is the first character input on a line, a line length of zero characters is returned to the reader. This is the standard end-of-file indication on a read call. Note that the <b>EOF</b> character set by this routine will apply to all devices whose handlers use the standard <i>tty</i> package <b>tyLib</b> .
RETURNS	N/A
ERRNO	N/A
SEE ALSO	tyLib

# tyIRd()

NAME	tyIRd() – interrupt-level input
SYNOPSIS	STATUS tyIRd ( FAST TY_DEV_ID pTyDev, /* ptr to tty device descriptor */ FAST char inchar /* character read */ )
DESCRIPTION	This routine handles interrupt-level character input for <i>tty</i> devices. A device driver calls this routine when it has received a character. This routine adds the character to the ring buffer for the specified device, and gives a semaphore if a task is waiting for it.
	This routine also handles all the special characters, as specified in the option word for the device, such as X-on, X-off, NEWLINE, or backspace.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the ring buffer is full.
ERRNO	N/A
SEE ALSO	tyLib

# tyITx()

NAME	<b>tyITx()</b> – interrupt-level output
SYNOPSIS	STATUS tyITx ( FAST TY_DEV_ID pTyDev, /* pointer to tty device descriptor */ char *pChar /* where to put character to be output */ )
DESCRIPTION	This routine gets a single character to be output to a device. It looks at the ring buffer for $pTyDev$ and gives the caller the next available character, if there is one. The character to be output is copied to <i>pChar</i> .
RETURNS	OK if there are more characters to send, or ERROR if there are no more characters.
ERRNO	N/A
SEE ALSO	tyLib

#### tyIoctl()

NAME	<b>tyIoctl()</b> – handle device control requests
SYNOPSIS	<pre>STATUS tyloctl   (   FAST TY_DEV_ID pTyDev, /* ptr to device to control */   int request, /* request code */   int arg /* some argument */   )</pre>
DESCRIPTION	This routine handles <b>ioctl()</b> requests for <i>tty</i> devices. The I/O control functions for <i>tty</i> devices are described in the reference entry for <b>tyLib</b> .
BUGS	In line protocol mode ( <b>OPT_LINE</b> option set), the FIONREAD function actually returns the number of characters available plus the number of lines in the buffer. Thus, if five lines consisting of just NEWLINEs were in the input buffer, the FIONREAD function would return the value ten (five characters + five lines).
RETURNS	OK or ERROR.
ERRNO	N/A
SEE ALSO	tyLib

#### tyLibInit()

NAME tyLibInit() – initialize the *tty* library SYNOPSIS STATUS tyLibInit ( int xoffPercent, /\* default Buffer percentage for sending XOFF \*/ int xonPercent, /\* default Buffer percentage for sending XON \*/ int wrtThreshold /\* default Buffer count for enabling other senders \*/ ) DESCRIPTION This routine initialized the library, and set the threshold values that will be used for xoff/xon control and for enabling new writer tasks. The xoff/xon threshold values are specified as percentages of buffer fill when the related action will occur. Normally xoff is sent out the transmit port when the receive buffer is 85% full. The xon character will be sent to the transmit port when the buffer has drained down to 50% the normal default value for the xonThreshold.

	The wrtThreshold is specified in the number of free bytes that must exist in the transmit
	buffer before a new writer task will be awakened. If the transmit buffer is very full, enabling
	a new transmit task can potentially awaken all waiting tasks and cause them all to priority
	arbitrate for the device only to find it blocked again. This threshold prevents the waking up
	of all pended tasks unless there is some minimum amount of space left in the transmit
	buffer. Pending tasks will be awakened when the buffer has drained to a point below the
	threshold number of characters.
RETURNS	<b>OK</b> or <b>ERROR</b> if arguments are invalid. The xon percentage must be smaller than the xoff percentage. Both must be in the range of 1 to 99.
ERRNO	Not Available

see also tyLib

tyMonitorTrapSet()

NAME	<b>tyMonitorTrapSet()</b> – change the trap-to-monitor character
SYNOPSIS	<pre>void tyMonitorTrapSet    (    char ch /* char to be monitor trap */   )</pre>
DESCRIPTION	This routine sets the trap-to-monitor character to <i>ch</i> . The default trap-to-monitor character is CTRL-X.
	Typing the trap-to-monitor character to any device whose <b>OPT_MON_TRAP</b> option is set will cause the resident ROM monitor to be entered, if one is present. Once the ROM monitor is entered, the normal multitasking system is halted.
	Note that the trap-to-monitor character set by this routine will apply to all devices whose handlers use the standard <i>tty</i> package <b>tyLib</b> . Also note that not all systems have a monitor trap available.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	tyLib

# tyRead()

NAME	<b>tyRead()</b> – do a task-level read for a <i>tty</i> device
SYNOPSIS	<pre>int tyRead   (   FAST TY_DEV_ID pTyDev, /* device to read */   char *buffer, /* buffer to read into */   int maxbytes /* maximum length of read */   )</pre>
DESCRIPTION	This routine handles the task-level portion of the <i>tty</i> handler's read function. It reads into the buffer up to <i>maxbytes</i> available bytes.
	This routine should only be called from serial device drivers.
RETURNS	The number of bytes actually read into the buffer.
ERRNO	N/A
SEE ALSO	tyLib

# tyWrite()

NAME	<b>tyWrite( )</b> – do a task-level write for a <i>tty</i> device
SYNOPSIS	<pre>int tyWrite   (   FAST TY_DEV_ID pTyDev, /* ptr to device structure */   char *buffer, /* buffer of data to write */   FAST int nbytes /* number of bytes in buffer */   )</pre>
DESCRIPTION	This routine handles the task-level portion of the <i>tty</i> handler's write function.
RETURNS	The number of bytes actually written to the device.
ERRNO	N/A
SEE ALSO	tyLib

# tyXoffHookSet()

NAME	<b>tyXoffHookSet( )</b> – install a hardware flow control function
SYNOPSIS	<pre>STATUS tyXoffHookSet   (    TY_DEV_ID pTyDev, /* pointer to device structure */   FUNCPTR func, /* Hardware flow control routine */    int arg /* First argument to func routine */   )</pre>
DESCRIPTION	This routine installs a hook routine to implement incoming flow control, to replace the default software XOFF/XON flow control method.
	The installed function will be called with two arguments. The first is the arg provided when the hook is set, and the second is a boolean value to indicate that incoming characters should be stopped ( <b>TRUE</b> means to disable input, <b>FALSE</b> means to allow input).
	VOID func (int arg, BOOLEAN xoffValue);
	Installing a NULL function pointer will restore the default software XOFF/XON method. This will nullify any previous XoffHook installation.
	With any change to the flow control routine, the old flow control routine is invoked to enable flow, before the new routine is actually installed. This insures that the incoming flow is not locked up when the method is changed.
RETURNS	OK if successful, or ERROR if the TY_DEV_ID is invalid.
ERRNO	Not Available
SEE ALSO	tyLib

#### unixDiskDevCreate()

NAME	<pre>unixDiskDevCreate( ) - create a UNIX disk device</pre>
SYNOPSIS	<pre>BLK_DEV * unixDiskDevCreate   (     char * unixFile, /* name of the UNIX file     int bytesPerBlk, /* number of bytes per block     int blksPerTrack, /* number of blocks per track     int nBlocks /* number of blocks on this device     )</pre>

\*/ \*/ \*/

DESCRIPTION	This routine creates a UNIX disk device.
	The <i>unixFile</i> parameter specifies the name of the UNIX file to use for the disk device.
	The <i>bytesPerBlk</i> parameter specifies the size of each logical block on the disk. If <i>bytesPerBlk</i> is zero, 512 is the default.
	The <i>blksPerTrack</i> parameter specifies the number of blocks on each logical track of the disk. If <i>blksPerTrack</i> is zero, the count of blocks per track is set to <i>nBlocks</i> (i.e., the disk is defined as having only one track).
	The <i>nBlocks</i> parameter specifies the size of the disk, in blocks. If <i>nBlocks</i> is zero, a default size is used. The default is calculated as the size of the UNIX disk divided by the number of bytes per block.
	This routine is only applicable to VxSim for Solaris.
IMPORTANT NOTE	This routine is obsolete, but is kept for backward compatibility with previous version.
RETURNS	A pointer to block device (BLK_DEV) structure, or NULL, if unable to open the UNIX disk.
ERRNO	Not Available

SEE ALSO unixDrv

### unixDiskInit()

NAME	unixDiskInit() – initialize a dosFs disk on top of UNIX		
SYNOPSIS	<pre>void unixDiskInit   (     char * unixFile, /* UNIX file name */     char * volName, /* dosFs name */     int diskSize /* number of bytes */   )</pre>		
DESCRIPTION	This routine provides some convenience for a user wanting to create a UNIX disk-based dosFs file system under VxWorks. The user only specifes the UNIX file to use, the dosFs volume name, and the size of the volume in bytes, if the UNIX file needs to be created.		
	This routine is only applicable to VxSim for Solaris.		
IMPORTANT NOTE	This routine is obsolete, but is kept for backward compatibility with previous version.		
RETURNS	N/A		

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ERRNO Not Available

SEE ALSO unixDrv

# unixDrv()

NAME	unixDrv( ) – install UNIX disk driver		
SYNOPSIS	STATUS unixDrv (void)		
DESCRIPTION	This routine is to cause the UNIX disk driver to be linked in when building VxWorks wh <b>INCLUDE_DOS_DISK</b> component is included in VxWorks image. Otherwise, it is not necessary to call this routine before using the UNIX disk driver.		
	This routine is only applicable to VxSim for Solaris.		
IMPORTANT NOTE	This routine is obsolete, but is kept for backward compatibility.		
RETURNS	OK (always).		
ERRNO	Not Available		
SEE ALSO	unixDrv		

# unld()

NAME	<b>unld()</b> – unload an object module by specifying a file name or module ID (shell command)
SYNOPSIS	<pre>STATUS unld   (     void * nameOrId, /* name or ID of the object module file */     int options /* Options to control behavior */   )</pre>
DESCRIPTION	This routine unloads the specified object module from the system. The module can be specified by name or by module ID. Unloading does the following:
	(1) It frees the space allocated for text, data, and BSS segments, unless the module was loaded using <b>loadModuleAt()</b> with user-specified addresses, in which case the user is responsible for freeing the space.

	(2) It removes all symbols associated with the object module from the system symbol table.
	(3) It removes the module descriptor from the module list.
	Before any modules are unloaded, all breakpoints in the system are deleted. If you need to keep breakpoints, set the options parameter to UNLD_KEEP_BREAKPOINTS. To use this option successfully, no breakpoints can be set in the code that is being unloaded.
	This routine is a <b>shell command</b> . That is, it is designed to be used only in the shell, and not in code running on the target. In future releases, calling <b>unld()</b> directly from code may not be supported.
	Note that using this command with an argument that is neither a module name or an ID can cause unpredictable behavior.
RETURNS	OK or ERROR.
ERRNO	Not Available
SEE ALSO	usrLib, loadLib, ld(), reld(), the VxWorks programmer guides.

# unldByGroup()

NAME	<b>unldByGroup()</b> – unload an object module by specifying a group number		
SYNOPSIS	<pre>STATUS unldByGroup  (  UINT16 group, /* group number to unload */  int options /* options */ )</pre>		
DESCRIPTION	This routine unloads an object module that has a group number matching the <i>group</i> parameter.		
	The <i>options</i> parameter may be set to any of the options that are available to the <b>unldByModuleId()</b> API. See its reference for more information.		
	See the manual entries for <b>unldLib</b> for more information on module unloading.		
RETURNS	<b>OK</b> , or <b>ERROR</b> if there is a problem.		
ERRNO	Not Available		
SEE ALSO	unldLib, symLib, unldByModuleId( )		

#### unldByModuleId()

unldByModuleId() – unload an object module by specifying a module ID NAME SYNOPSIS STATUS unldByModuleId ( MODULE\_ID moduleId, /\* module ID to unload \*/ options /\* options \*/ int ) This routine unloads an object module that has a module ID matching the moduleId DESCRIPTION parameter. Unloading does the following: (1) Frees the space allocated for the code module segments (text, data, and BSS), unless **loadModuleAt()** was used to specify the locations where the segments were to be loaded, in which case the user is responsible for freeing the space. (2) It removes all symbols associated with the code module from the symbol table. (3) It removes the code module descriptor, its list of segment descriptors, and its list of section descriptors from the kernel's code module list. The unloader accepts the following options which may be combined by a binary **OR** (UNLD\_CPLUS\_XTOR\_AUTO and UNLD\_CPLUS\_XTOR\_MANUAL are mutually exclusive): UNLD KEEP BREAKPOINTS Before any modules are unloaded, all breakpoints in the system are deleted. If you need to keep breakpoints, set the options parameter to UNLD\_KEEP\_BREAKPOINTS. To use this option safely, there should be no breakpoints set in the code that is being unloaded. UNLD\_FORCE By default, the unloader does not remove the text sections when they are used by some hooks in the system (see the manual of **unldLib** for the list of hooks). Using UNLD\_FORCE will force the unloader to remove the sections anyway, at the risk of unpredictable results. UNLD CPLUS XTOR AUTO This option specifies that the unloader should call the code module's C++ destructor routines. UNLD\_CPLUS\_XTOR\_MANUAL This option prevents the unloader from calling the code module's C++ destructor routines. If using this option, the user should be sure that the destructor routines do not perform the release of any resources back to the system, such as memory or semaphores. Or the caller may first cause any static destructors to be run by using the function **cplusDtors()**. OK, or ERROR if there is a problem. RETURNS

SEE ALSO	unldLib				
	<ul><li>+ S_moduleLib_INVALID_MODULE_ID</li><li>For a complete description of the errnos, see the reference documentation for moduleLib.</li></ul>				
ERRNO	Possible errnos set by this routine include:				

# unldByNameAndPath()

NAME	unldByNameAndPath() – unload an object module by specifying a name and path			
SYNOPSIS	<pre>STATUS unldByNameAndPath   (     char * name, /* name of the object module to unload */     char * path, /* path to the object module to unload */     int options /* options */   )</pre>			
DESCRIPTION	This routine unloads an object module specified by the <i>name</i> and <i>path</i> parameters. The <i>name</i> and <i>path</i> correspond to the parameters that were passed to the load routine <i>when the module was loaded</i> .			
	The <i>options</i> parameter may be set to any of the options that are available to the <b>unldByModuleId()</b> API. See its reference for more information.			
	See the manual entries for <b>unldLib</b> for more information on module unloading.			
EXAMPLES	If the module was loaded using the following name and path:			
	<pre>fd = open ("path/to/the/module/to/load/moduleName", O_RDONLY); moduleLoad (fd, LOAD_GLOBAL_SYMBOLS);</pre>			
	then the call to <b>unldByNameAndPath()</b> would be done as:			
	unldByNameAndPath ("moduleName", "path/to/the/module/to/load", 0);			
	The path field should be left empty if the module was loaded without any path specified:			
	fd = open ("moduleName", O_RDONLY); moduleLoad (fd, LOAD_GLOBAL_SYMBOLS); unldByNameAndPath ("moduleName", "", 0);			
RETURNS	<b>OK</b> , or <b>ERROR</b> if there is a problem.			
ERRNO	Not Available			
SEE ALSO	unldLib, unldByModuleId()			

#### unlink()

NAME	unlink() – unlink a file
SYNOPSIS	<pre>int unlink   (    const char *name /* name of the file to remove */   )</pre>
DESCRIPTION	This routine removes a link to a file. It shall remove the link named by <i>name</i> and decrease the link count of the file referenced by the link.
RETURNS	OK if successful; ERROR otherwise.
ERRNO	
SEE ALSO	fsPxLib, link()

#### unstatShow()

NAME unstatShow() – display all AF\_LOCAL sockets

SYNOPSIS void unstatShow (void)

**DESCRIPTION** This routine displays a list of all **AF\_LOCAL** family sockets in a format similar to the UNIX **netstat -f unix** command.

Sample output:

AF\_LOCAL/COMP protocol sockets

_ so#	address (self/peer)	State	bytes of data	high watermark	packets dropped	pending connections (cur/max)
7	0001/NONE	LISTENING	0	N/A	N/A	1/ 20
9	0001/0002	EXCHANGING	30012	65535	10	N/A
8	0002/0001	EXCHANGING	20014	20014	0	N/A
10	0003/NONE	LISTENING	0	N/A	N/A	5/ 5
12	0003/0004	DONE_RECV	0	65535	132	N/A
11	0004/0003	DONE_SEND	16	16	0	N/A
4	0405/NONE	LISTENING	0	N/A	N/A	0/ 5
5	0000/NONE	CLOSED	0	0	0	N/A
6	0000/NONE	CLOSED	0	0	0	N/A

	Socket identifier, relative to owner's RTP.
	address Socket's and peer's addresses (/comp/socket/0xWXYZ). <b>peer</b> is only valid in the following state: <b>EXCHANGING</b> , <b>DONE_RECV</b> , <b>DONE_SEND</b> . <b>NONE</b> is printed otherwise.
	state
	CLOSEDstarting and ending state: no data can flow
	LISTENING socket is a listening one, cannot be used to transfer data
	EXCHANGING data can be sent in either direction
	DONE_SENDsocket can only receive, shutdown(write) has been called on it
	DONE_RECVsocket can only send, shutdown(write) has been done on peer
	bytes of data Amount of data that is pending on the socket, waiting to be received.
	high watermark Largest amount of pending data at one given time.
	dropped packets Number of packets dropped due to lack of space in the receiver's buffer space.
	<pre>pending connections Current and maximum number of unaccepted connections on a listening socket, i.e.: (number of connect() calls) - (number of accept() calls). N/A is printed if not in the LISTENING state.</pre>
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	unShow

# usrClock()

NAME usrClock() – user-defined system clock interrupt routine

SYNOPSIS void usrClock (void)

so#

VxWorks Kernel API Reference, 6.6 usrFdiskPartCreate()

 DESCRIPTION
 This routine is called at interrupt level on each clock interrupt. It is installed by usrRoot() with a sysClkConnect() call. It calls all the other packages that need to know about clock ticks, including the kernel itself.

 If the application needs anything to happen at the system clock interrupt level, it can be added to this routine.

 RETURNS
 N/A

 ERRNO
 Not Available

 SEE ALSO
 usrConfig

#### usrFdiskPartCreate()

NAME	usrFdiskPartCreate() – create an FDISK-like partition table on a disk				
SYNOPSIS	<pre>STATUS usrFdiskPartCreate   (     CBIO_DEV_ID cDev, /* device representing the entire disk */     int nPart, /* how many partitions needed, default=1, max=4 */     int size1, /* space percentage for second partition */     int size2, /* space percentage for third partition */     int size3 /* space percentage for fourth partition */     )</pre>				
DESCRIPTION	This function may be used to create a basic PC partition table. Such a partition table is not intended to be compatible with other operating systems; it is intended for disks connected to a VxWorks target, but without the access to a PC which may be used to create the partition table. This function is capable of creating only one partition table - the MBR, and will not create				
	any Bootable or Extended partitions. Therefore, only 4 partitions are supported.				
	<i>cDev</i> is a CBIO device handle for an entire disk, e.g. a handle returned by <b>dcacheDevCreate()</b> , or if dpartCbio is used, it can be either the Master partition manager handle, or the one of the 0th partition if the disk does not contain a partition table at all.				
	The <i>nPart</i> argument contains the number of partitions to create. If <i>nPart</i> is 0 or 1, a single partition covering the entire disk is created. If <i>nPart</i> is between 2 and 4, the arguments <i>size1</i> , <i>size2</i> and <i>size3</i> contain (as integers) the <i>percentage</i> of disk space to be assigned to the 2nd, 3rd, and 4th partitions respectively. The first partition (partition 0) will be assigned the remaining space. Thus, the sum of the three sizes should be less than 100.				
	Partition sizes will be rounded down to be multiple of whole tracks so that partition Cylinder/Head/Track fields will be initialized as well as the LBA fields. Although the CHS				

fields are written they are not used in VxWorks, and can not be guaranteed to work correctly on other systems. **OK** or **ERROR** writing a partition table to disk RETURNS Not Available

usrFdiskPartLib SEE ALSO

ERRNO

#### usrFdiskPartRead()

NAME	usrFdiskPartRead() – read an FDISK-style partition table			
SYNOPSIS	<pre>STATUS usrFdiskPartRead   (    CBIO_DEV_ID</pre>			
DESCRIPTION	This function will read and decode a PC formatted partition table on a disk, and fill the appropriate partition table array with the resulting geometry, which should be used by the dpartCbio partition manager to access a partitioned disk with a shared disk cache.			
EXAMPLE	The following example shows how a hard disk which is expected to have up to two partitions might be configured, assuming the physical level initialization resulted in the <i>blkIoDevId</i> handle:			
	<pre>devCbio = dcacheDevCreate( blkIoDevId, 0, 0x20000, "Hard Disk"); mainDevId = dpartDevCreate( devCbio, 2, usrFdiskPartRead ) dosFsDevCreate( "/disk0a", dpartPartGet (mainDevId, 0), 0,0,0); dosFsDevCreate( "/disk0b", dpartPartGet (mainDevId, 1), 0,0,0);</pre>			
RETURNS	OK or ERROR if partition table is corrupt			
ERRNO	Not Available			
SEE ALSO	usrFdiskPartLib			

#### usrFdiskPartShow()

NAME	usrFdiskPartShow() – parse and display partition data	
SYNOPSIS	<pre>STATUS usrFdiskPartShow   (     CBI0_DEV_ID cbio,</pre>	
DESCRIPTION	This routine is intended to be user callable.	
	A device dependent partition table show routine, this routine outputs formatted data for all partition table fields for every partition table found on a given disk, starting with the MBR sectors partition table. This code can be removed to reduce code size by undefining: <b>INCLUDE_PART_SHOW</b> and rebuilding this library and linking to the new library.	
	This routine takes three arguments. First, a CBIO pointer (assigned for the entire physical disk) usually obtained from <b>dcacheDevCreate()</b> . It also takes two block_t type arguments and one signed int. The user shall pass zero in these paramaters.	
	For example:	
	sp usrFdiskPartShow (pCbio,0,0,0)	
	Developers may use sizearch to view code size.	
RETURNS	OK or ERROR	
ERRNO	Not Available	

SEE ALSO usrFdiskPartLib

# usrFormatTrans()

NAME usrFormatTrans() – Perform a low-level trans XBD format opera	tion
--	------

SYNOPSIS STATUS usrFormatTrans

( char \*dev, int overhead, int type )

DESCRIPTION	This routine formats a trans XBD with the specified parameters.
	The <i>dev</i> parameter is the path name of the device (which will have any existing file system ejected and a rawFS put on it during formatting). The <i>overhead</i> parameter specifies the amount of media to use for uncommitted workspace, in parts-per-thousand. The <i>type</i> parameter specifies the type of format to use. The value of this parameter will be either FORMAT_REGULAR (0) or FORMAT_TFFS (1). FORMAT_REGULAR initializes a system with 2 master records at the beginning and end of the disk. FORMAT_TFFS initializes a system with the first sector unused, and a master record at the end of the disk.
	This routine then waits for the device to re-instantiate as a TRFS device, at which point it is safe to format it for dosFs.
EXAMPLE	usrFormatTrans ("/trans", 100, 0); This formats the device referred to by "/trans" to use 10% of the disk as workspace, and places master records on the first and last sectors of the disk.
RETURNS	<b>OK</b> , or <b>ERROR</b> if an error occurs formatting the device.
ERRNO	Not Available
SEE ALSO	usrTransLib

# usrIdeConfig()

NAME	usrIdeConfig() – mount a DOS file system from an IDE hard disk
SYNOPSIS	STATUS usrIdeConfig ( int drive, /* drive number of hard disk (0 or 1) */ char * fileName /* mount point */ )
DESCRIPTION	This routine mounts a DOS file system from an IDE hard disk.
	The <i>drive</i> parameter is the drive number of the hard disk; 0 is <b>C</b> : and 1 is <b>D</b> :.
	The <i>fileName</i> parameter is the mount point, e.g., /ide0/.
NOTE	Because VxWorks does not support partitioning, hard disks formatted and initialized on VxWorks are not compatible with DOS machines. This routine does not refuse to mount a hard disk that was initialized on VxWorks. The hard disk is assumed to have only one partition with a partition record in sector 0.
RETURNS	OK or ERROR.

VxWorks Kernel API Reference, 6.6 usrInit()

ERRNO Not Available

**SEE ALSO** usrIde, the VxWorks programmer guides, the architecture supplement.

#### usrInit()

NAME	usrInit() – user-defined system initialization routine
SYNOPSIS	<pre>void usrInit    (    int startType    )</pre>
DESCRIPTION	This is the first C code executed after the system boots. This routine is called by the assembly language start-up routine <b>sysInit()</b> which is in the <b>sysALib</b> module of the target-specific directory. It is called with interrupts locked out. The kernel is not multitasking at this point.
	This routine starts by clearing BSS; thus all variables are initialized to 0, as per the C specification. It then initializes the hardware by calling <b>sysHwInit()</b> , sets up the interrupt/exception vectors, and starts kernel multitasking with <b>usrRoot()</b> as the root task.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	usrConfig, kernelLib

#### usrRoot()

NAME	<b>usrRoot( )</b> – the root task	
SYNOPSIS	void usrRoot ( char * pMemPoolStart, unsigned memPoolSize )	/* start of system memory partition */ /* initial size of mem pool */

**DESCRIPTION** This is the first task to run under the multitasking kernel. It performs all final initialization and then starts other tasks.

It initializes the I/O system, installs drivers, creates devices, and sets up the network, etc., as necessary for a particular configuration. It may also create and load the system symbol table, if one is to be included. It may then load and spawn additional tasks as needed. In the default configuration, it simply initializes the VxWorks shell.

RETURNS N/A

ERRNO Not Available

SEE ALSO usrConfig

#### usrScsiConfig()

NAME	usrScsiConfig() – configure SCSI peripherals
SYNOPSIS	STATUS usrScsiConfig (void)
DESCRIPTION	This code configures the SCSI disks and other peripherals on a SCSI controller chain.
	The macro <b>SCSI_AUTO_CONFIG</b> will include code to scan all possible device/lun id's and to configure a scsiPhysDev structure for each device found. Of course this doesn't include final configuration for disk partitions, floppy configuration parameters, or tape system setup. All of these actions must be performed by user code, either through <b>sysScsiConfig()</b> , the startup script, or by the application program.
	The user may customize this code on a per BSP basis using the <b>SYS_SCSI_CONFIG</b> macro. If defined, then this routine will call the routine <b>sysScsiConfig()</b> . That routine is to be provided by the BSP, either in <b>sysLib.c</b> or <b>sysScsi.c</b> . If <b>SYS_SCSI_CONFIG</b> is not defined, then <b>sysScsiConfig()</b> will not be called as part of this routine.
	An example <b>sysScsiConfig()</b> routine can be found in <b>target/src/config/usrScsi.c</b> . The example code contains sample configurations for a hard disk, a floppy disk and a tape unit.
RETURNS	OK or ERROR.
ERRNO	Not Available
SEE ALSO	usrScsi, VxWorks Programmer's Guide: I/O System, Local File Systems

### usrTransCommit()

NAME	usrTransCommit() – Set a transaction point on a trans XBD
SYNOPSIS	STATUS usrTransCommit ( char *volume )
DESCRIPTION	This routine sets a transaction point using the volume name of the device. The <i>volume</i> parameter is the name of the device on which TRFS is instantiated.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the transaction point is not set.
ERRNO	Not Available
SEE ALSO	usrTransLib

#### usrTransCommitFd()

NAME	usrTransCommitFd() – set a transaction point using a file descriptor
SYNOPSIS	STATUS usrTransCommitFd ( int fd )
DESCRIPTION	This routine sets a transaction point on the device which contains the filesystem containing the file to which the parameter refers.
	The <i>fd</i> parameter is a file descriptor, which must refer to a file whose backing media uses a TRFS XBD.
RETURNS	OK, or ERROR if the transaction point is not set.
ERRNO	Not Available
SEE ALSO	usrTransLib

#### uswab() uswab() – swap bytes with buffers that are not necessarily aligned NAME SYNOPSIS void uswab ( char \*source, /\* pointer to source buffer \*/ char \*destination, /\* pointer to destination buffer \*/ int nbytes /\* number of bytes to exchange \*/ ) This routine gets the specified number of bytes from source, exchanges the adjacent even and DESCRIPTION odd bytes, and puts them in *destination*. NOTE: Due to speed considerations, this routine should only be used when absolutely necessary. Use swab() for aligned swaps. The value of *nBytes* must not be odd. Failure to adhere to this may yield incorrect results. N/A RETURNS N/A ERRNO bLib, swab() SEE ALSO

#### utf16ToCP()

NAME utf16ToCP() – Convert a UTF-16 encoded Unicode character to a codepoint. SYNOPSIS int utf16ToCP ( const unsigned short \* utf16, const int length, /\* length is in 16-bit words \*/ littleEndian, const int unsigned long \* codePoint DESCRIPTION This routine converts a character encoded as UTF-16 to an unsigned long which represents the value of a Unicode characters codepoint. RETURNS If positive, the return value is the number of UTF-16 words converted. If negative, the value

RETURNS If positive, the return value is the number of UTF-16 words converted. If negative, the value UC\_NOSRC indicates that insufficient words were given to represent a codepoint. The value UC\_FORMAT indicates that the UTF-16 vector was of an invalid format. VxWorks Kernel API Reference, 6.6 utf16ToUtf8String()

ERRNO Not Available

SEE ALSO utfLib

#### utf16ToUtf8String()

NAME	utf16ToUtf8String() – Convert a UTF-16 string to a UTF-8 String
SYNOPSIS	<pre>int utf16ToUtf8String   (    const unsigned short * utf16,    int littleEndian,    unsigned char * utf8,    const int len8   )</pre>
DESCRIPTION	This routine converts a Zero terminated, UTF-16 encoded string of the indicated endianess to a NULL terminated UTF-8 encoded string.
RETURNS	If positive, returns the number of bytes used by the resulting UTF-8 encoded string. If non-positive, <b>UC_FORMAT</b> indicates that the UTF-16 string is of an invalid format; <b>UC_BUFFER</b> indicates that the buffer provided for the UTF-8 string is too small.
ERRNO	Not Available
SEE ALSO	utfLib

## utf16ToUtf8StringBOM()

NAME	utf16ToUtf8StringBOM() – Convert UTF-16 to UTF-8 based on a Byte Order Mark
SYNOPSIS	<pre>int utf16ToUtf8StringBOM   (    const unsigned short * utf16,    unsigned char * utf8,    const int len8   )</pre>
DESCRIPTION	This routine handles UTF-16 in its standard form. If the first word is a Byte Order Mark Code Point 0xFEFF, then it is examined for endianness, and the rest of the string is interpreted accordingly. If there is no Byte Order Mark, then the string is interpreted as

 big-endian representation. Note that the Byte Order Mark is a legitimate, though deprecated, character.

 RETURNS
 If positive, returns the number of bytes used by the resulting UTF-8 encoded string. If non-positive, UC\_FORMAT indicates that the UTF-16 string is of an invalid format; UC\_BUFFER indicates that the buffer provided for the UTF-8 string is too small.

 ERRNO
 Not Available

 SEE ALSO
 utfLib

#### utf8ToCP()

NAME utf8ToCP() – Convert a UTF-8 encoded Unicode character to the Unicode codepoint. SYNOPSIS int utf8ToCP ( const unsigned char \* utf8, const int length, unsigned long \* codePoint ) This routine converts UTF-8 to an unsigned long which represents the value of the Unicode DESCRIPTION codepoint. If positive, the return value is the number of characters converted to this codepoint. If RETURNS non-positive, the return value of UC\_NOSRC indicates that there are insufficient characters for a valid conversion, and a return value of UC\_FORMAT indicates that the format of the input string is not valid UTF-8. ERRNO Not Available utfLib SEE ALSO

#### utf8ToUtf16String()

NAME utf8ToUtf16String() – convert a UTF-8 string to a UTF-16 string

SYNOPSIS int utf8ToUtf16String

VxWorks Kernel API Reference, 6.6 utf8ToUtf16StringBOM()

	( const unsigned char * utf8, unsigned short * utf16, const int len16, int littleEndian )
DESCRIPTION	This routine converts a NULL terminated UTF-8 encoded string to a ZERO terminated UTF-16 string of the indicated endianess. It does not prepend a Byte Order Marker to the beginning of the string - this must be done before conversion if it is required.
RETURNS	If positive the number of 16-bit words actually converted, including the terminating Zero. If non-positive, UC_FORMAT indicates that the UTF-8 string is not of a legal format, and UC_BUFFER indicates that the provided buffer for containing the UTF-16 string is too small to perform the conversion
ERRNO	Not Available
SEE ALSO	utfLib

#### utf8ToUtf16StringBOM()

utf8ToUtf16StringBOM() - Convert UTF-8 to UTF16 with a Byte Order Mark NAME SYNOPSIS int utf8ToUtf16StringBOM ( const unsigned char \* utf8, unsigned short \* utf16, const int len16, int littleEndian ) This routine first writes out the Unicode Byte Order Mark Character, and then converts the DESCRIPTION UTF-8 encoded string to UTF-16 based on the given endianness If positive the number of 16-bit words actually converted, including the terminating Zero RETURNS and the Byte Order Marker . If non-positive, UC\_FORMAT indicates that the UTF-8 string is not of a legal format, and UC\_BUFFER indicates that the provided buffer for containing the UTF-16 string is too small to perform the conversion. ERRNO Not Available SEE ALSO utfLib

#### utfLibInit()

NAME utfLibInit() – initialize the UTF library

SYNOPSIS void utfLibInit (void)

**DESCRIPTION** none

RETURNS Not Available

ERRNO Not Available

SEE ALSO utfLib

#### utflen16()

utflen16() – Return the number of 16-bit words used by a UTF-16 encoding. NAME SYNOPSIS int utflen16 ( const unsigned short \* utf16 ) DESCRIPTION This routine returns the number of 16-bit words, including the terminating Zero, used by a Zero terminated UTF-16 encoded string. Then number of 16-bit words utilized by a UTF-16 encoding. RETURNS Not Available ERRNO utfLib SEE ALSO

#### utflen8()

NAME utflen8() – return the encoding length of a NULL terminated UTF-8 string

SYNOPSIS int utflen8

2

VxWorks Kernel API Reference, 6.6 utime()

(
const unsigned char \* utf8
)

**DESCRIPTION** This routine returns the length occupied by the encoding, as opposed to the number of Unicode characters actually encoded, including the terminating **NULL** 

**RETURNS** The total number of chars up to and including the terminating NULL.

ERRNO Not Available

SEE ALSO utfLib

#### utime()

NAME	utime() – update time on a file
SYNOPSIS	<pre>int utime   (     const char * file,     const struct utimbuf * newTimes   )</pre>
DESCRIPTION	Update the timestamp on a file. For filesystems that support this command, the timestamp of the file is updated to the current time.
RETURNS	OK or ERROR.
ERRNO	N/A
SEE ALSO	dirLib, stat(), fstat(), ls()

# valloc()

DESCRIPTION	This routine allocates a buffer of <i>size</i> bytes from the system memory partition (kernel heap). Additionally, it insures that the allocated buffer begins on a page boundary. Page sizes are architecture-dependent.
RETURNS	A pointer to the newly allocated block, or NULL if the buffer could not be allocated or the memory management unit (MMU) support library has not been initialized.
ERRNO	S_memLib_PAGE_SIZE_UNAVAILABLE Could not obtain the size of a virtual page. Possible error is that virtual memory support is not included (INCLUDE_MMU_BASIC).
	<b>S_memLib_NOT_ENOUGH_MEMORY</b> There is no free block large enough to satisfy the allocation request.
SEE ALSO	memLib

# version()

NAME	<b>version( )</b> – print VxWorks version information
SYNOPSIS	void version (void)
DESCRIPTION	This command prints the VxWorks version number, the date this copy of VxWorks was made, and other pertinent information.
EXAMPLE	-> version VxWorks (for SunOS 5.8 [sun4u]) version 6.0. Kernel: WIND version 2.7. Made on May 13 2004, 13:23:14. Boot line: passDev(0,0)river:/wind/river/target/proj/solaris_diab/default/vxWorks u=user tn=vxTarget value = 0 = 0x0
RETURNS	N/A
ERRNO	N/A
SEE ALSO	usrLib, the VxWorks programmer guides.

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NAME	<b>vfdprintf()</b> – write a string formatted with a variable argument list to a file descriptor	
SYNOPSIS	<pre>int vfdprintf   (    int fd, /* file descriptor to print to */    const char * fmt, /* format string for print */    va_list vaList /* optional arguments to format */   )</pre>	
DESCRIPTION	This routine prints a string formatted with a variable argument list to a specified file descriptor. It is identical to <b>fdprintf()</b> , except that it takes the variable arguments to be formatted as a list <i>vaList</i> of type <b>va_list</b> rather than as in-line arguments.	
SMP CONSIDERATIO	<b>NS</b> This API is spinlock and intCpuLock restricted.	
RETURNS	The number of characters output, or <b>ERROR</b> if there is an error during output.	
ERRNO	Not Available	
SEE ALSO	fioLib, fdprintf()	

### virtualDiskClose()

NAME	<b>virtualDiskClose( )</b> – close a virtual disk block device.
SYNOPSIS	STATUS virtualDiskClose ( BLK_DEV * blkDev /* virtual disk block device */ )
DESCRIPTION	This routine closes a virtual disk block device by closing the host file associated with the virtual disk. The <i>blkDev</i> parameter specifies the virtual disk to close.
RETURNS	OK on success, else ERROR.
ERRNO	Not Available
SEE ALSO	virtualDiskLib

#### virtualDiskCreate()

NAME	<b>virtualDiskCreate( )</b> – create a virtual disk device.	
SYNOPSIS	<pre>BLK_DEV * virtualDiskCreate   (     char * hostFile, /* name of the host file */     int bytesPerBlk, /* number of bytes per block */     int blksPerTrack, /* number of blocks per track */     int nBlocks /* number of blocks on this device */ )</pre>	
DESCRIPTION	This routine creates a virtual disk device. The host file is created if it does not exist. If it already exists, only the <i>hostFile</i> parameter is taken in account, the others parameters are extracted from the host file.	
	The <i>hostFile</i> parameter specifies the name of the host file used for the virtual disk. The host file pathname is a standard host pathname without the host name. For Windows VxSim, the path separator to use is \ or / (i.e. c:/myDir/myFile or c:\myDir\myFile).	
	The <i>bytesPerBlk</i> parameter specifies the size of each logical block on the disk. If <i>bytesPerBlk</i> is zero, 512 is the default.	
	The <i>blksPerTrack</i> parameter specifies the number of blocks on each logical track of the disk. If <i>blksPerTrack</i> is zero, the count of blocks per track is set to <i>nBlocks</i> (i.e., the disk is defined as having only one track).	
	The <i>nBlocks</i> parameter specifies the size of the disk, in blocks. If <i>nBlocks</i> is zero, 512 is the default.	
RETURNS	A pointer to block device (BLK_DEV) structure, or NULL if the virtual disk creation failed.	
ERRNO	Not Available	
SEE ALSO	virtualDiskLib	

#### virtualDiskInit()

NAME virtualDiskInit() – install the virtual disk driver

SYNOPSIS STATUS virtualDiskInit (void)

**DESCRIPTION** This routine is used to initialize the virtual disk driver. This routine is automatically called when the **INCLUDE\_VIRTUAL\_DISK** component is included.

VxWorks Kernel API Reference, 6.6 vmArch32LibInit()

**RETURNS** OK, always.

ERRNO Not Available

SEE ALSO virtualDiskLib

#### vmArch32LibInit( )

**NAME vmArch32LibInit()** – initialize the arch specific unbundled VM library (VxVMI Option)

SYNOPSIS void vmArch32LibInit (void)

**DESCRIPTION** This routine links the arch specific unbundled VM library into the VxWorks system. It is called automatically when INCLUDE\_MMU\_FULL and INCLUDE\_MMU\_P6\_32BIT are both defined in the BSP.

RETURNS N/A

ERRNO Not Available

SEE ALSO vmArch32Lib

#### vmArch32Map()

NAME	<b>vmArch32Map()</b> – map 32bit physical space into 32bit virtual space (VxVMI Option)
SYNOPSIS	<pre>STATUS vmArch32Map   (     VM_CONTEXT_ID context, /* context - NULL == currentContext */     void * virtAddr, /* virtual address */     void * physAddr, /* physical address */     UINT32 stateMask, /* state mask */     UINT32 state, /* state */     UINT32 len /* len of virtual and physical spaces */     )</pre>
DESCRIPTION	vmArch32Map maps 32bit physical pages into a contiguous block of 32bit virtual memory. <i>virtAddr</i> and <i>physAddr</i> must be on page boundaries, and <i>len</i> must be evenly divisible by the page size. After the mapping the specified state is set to all pages in the newly mapped virtual memory.

The vmArch32Map() routine can fail if the specified virtual address space conflicts with the
translation tables of the global virtual memory space. The global virtual address space is
initialized at boot time. If a conflict results, errno is set to
<b>S_vmLib_ADDR_IN_GLOBAL_SPACE</b> . To avoid this conflict, use <b>vmGlobalInfoGet()</b> to ascertain which portions of the virtual address space are reserved for the global virtual address space. If <i>context</i> is specified as <b>NULL</b> , the current virtual memory context is used.
This routine should not be called from interrupt level.
This routine is distributed as a component of the unbundled virtual memory support

- **RETURNS OK**, or **ERROR** if *virtAddr* or *physAddr* are not on page boundaries, *len* is not a multiple of the page size, the validation fails, or the mapping fails.
- ERRNO S\_vmLib\_NOT\_PAGE\_ALIGNED S\_vmLib\_ADDR\_IN\_GLOBAL\_SPACE

option, VxVMI.

SEE ALSO vmArch32Lib

AVAILABILITY

#### vmArch32Translate()

NAME	<b>vmArch32Translate()</b> – translate a 32bit virtual address to a 32bit physical address (VxVMI Option)	
SYNOPSIS	<pre>STATUS vmArch32Translate   (     VM_CONTEXT_ID context, /* context - NULL == currentContext */     void * virtAddr, /* virtual address */     void ** physAddr /* place to put result */   )</pre>	
DESCRIPTION	vmArch32Translate retrieves mapping information for a 32bit virtual address from the page translation tables. If the specified virtual address has never been mapped, the returned status is <b>ERROR</b> . If <i>context</i> is specified as <b>NULL</b> , the current context is used. This routine is callable from interrupt level.	
AVAILABILITY	This routine is distributed as a component of the unbundled virtual memory support option, VxVMI.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if the validation or translation fails.	
ERRNO	Not Available	

VxWorks Kernel API Reference, 6.6 vmArch36LibInit()

SEE ALSO vmArch32Lib

#### vmArch36LibInit()

**NAME vmArch36LibInit()** – initialize the arch specific unbundled VM library (VxVMI Option)

SYNOPSIS void vmArch36LibInit (void)

- **DESCRIPTION** This routine links the arch specific unbundled VM library into the VxWorks system. It is called automatically when INCLUDE\_MMU\_FULL and INCLUDE\_MMU\_P6\_36BIT are both defined in the BSP.
- RETURNS N/A

ERRNO Not Available

SEE ALSO vmArch36Lib

#### vmArch36Map()

NAME	<b>vmArch36Map()</b> – map 36bit physical space into 32bit virtual space (VxVMI Option)		
SYNOPSIS	<pre>STATUS vmArch36Map   (    VM_CONTEXT_ID context, /* context - NULL == currentContext */    void * virtAddr, /* 32bit virtual address */    LL_INT physAddr, /* 36bit physical address */    UINT32 stateMask, /* state mask */    UINT32 state, /* state */    UINT32 len /* len of virtual and physical spaces */    )</pre>		
DESCRIPTION	vmArch36Map maps 36bit physical pages into a contiguous block of 32bit virtual memory. <i>virtAddr</i> and <i>physAddr</i> must be on page boundaries, and <i>len</i> must be evenly divisible by the page size. After the mapping the specified state is set to all pages in the newly mapped virtual memory.		
	The <b>vmArch36Map()</b> routine can fail if the specified virtual address space conflicts with the translation tables of the global virtual memory space. The global virtual address space is		

translation tables of the global virtual memory space. The global virtual address space is initialized at boot time. If a conflict results, **errno** is set to **S\_vmLib\_ADDR\_IN\_GLOBAL\_SPACE**. To avoid this conflict, use **vmGlobalInfoGet()** to

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	ascertain which portions of the virtual address space are reserved for the global virtual address space. If <i>context</i> is specified as <b>NULL</b> , the current virtual memory context is used.
	This routine should not be called from interrupt level.
AVAILABILITY	This routine is distributed as a component of the unbundled virtual memory support option, VxVMI.
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>virtAddr</i> or <i>physAddr</i> are not on page boundaries, <i>len</i> is not a multiple of the page size, the validation fails, or the mapping fails.
ERRNO	S_vmLib_NOT_PAGE_ALIGNED S_vmLib_ADDR_IN_GLOBAL_SPACE
SEE ALSO	vmArch36Lib

# vmArch36Translate()

NAME	<b>vmArch36Translate( )</b> – translate a 32bit virtual address to a 36bit physical address (VxVMI Option)		
SYNOPSIS	<pre>STATUS vmArch36Translate   (     VM_CONTEXT_ID context, /* context - NULL == currentContext */     void * virtAddr, /* 32bit virtual address */     LL_INT * physAddr /* place to put 36bit result */   )</pre>		
DESCRIPTION	vmArch36Translate retrieves mapping information for a 32bit virtual address from the page translation tables. If the specified virtual address has never been mapped, the returned status is <b>ERROR</b> . If <i>context</i> is specified as <b>NULL</b> , the current context is used. This routine is callable from interrupt level.		
AVAILABILITY	This routine is distributed as a component of the unbundled virtual memory support option, VxVMI.		
RETURNS	<b>OK</b> , or <b>ERROR</b> if the validation or translation fails.		
ERRNO	Not Available		
SEE ALSO	vmArch36Lib		

#### vmAttrShow()

NAME	<b>vmAttrShow()</b> – display the text representation of a MMU attribute value		
SYNOPSIS	STATUS vmAttrShow ( UINT pgAttr, /* MMU attributes */ UINT pgAttrMask /* MMU attributes mask */ )		
DESCRIPTION	This routine will display the text value of the attributes passed in <i>pgAttr</i> . This information includes the supervisor and user RWX values, cache mode (CB/WT/OFF), coherency, and I/O settings.		
	Note that this routine cannot report non-standard architecture-dependent states.		
RETURNS	ERROR if <i>pgAttrMask</i> is not a valid mask or NULL, else OK.		
ERRNO	Not Available		
SEE ALSO	vmShow		

## vmBaseArch32LibInit( )

NAME	<b>vmBaseArch32LibInit( )</b> – initialize the arch specific bundled VM library
SYNOPSIS	void vmBaseArch32LibInit (void)
DESCRIPTION	This routine links the arch specific bundled VM library into the VxWorks system. It is called automatically when INCLUDE_MMU_BASIC and INCLUDE_MMU_P6_32BIT are both defined in the BSP.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	vmBaseArch32Lib

# vmBaseArch32Map() vmBaseArch32Map() - map 32bit physical to the 32bit virtual memory STATUS vmBaseArch32Map (

void *	virtAddr,	/*	32bit virtual address */
void *	physAddr,	/*	32bit physical address */
UINT32	stateMask,	/*	state mask */
UINT32	state,	/*	state */
UINT32	len	/*	length */
)			

**DESCRIPTION** vmBaseArch32Map maps 32bit physical pages into a contiguous block of 32bit virtual memory. *virtAddr* and *physAddr* must be on page boundaries, and *len* must be evenly divisible by the page size. After the mapping the specified state is set to all pages in the newly mapped virtual memory.

This routine should not be called from interrupt level.

- **RETURNS** OK, or ERROR if *virtAddr* or *physAddr* are not on page boundaries, *len* is not a multiple of the page size, the validation fails, or the mapping fails.
- ERRNO S\_vmLib\_NOT\_PAGE\_ALIGNED
- SEE ALSO vmBaseArch32Lib

NAME

SYNOPSIS

#### vmBaseArch32Translate()

NAME	vmBaseArch32Translate() – translate a 32bit virtual address to a 32bit physical address		
SYNOPSIS	STATUS vmBaseArch32Translate ( void * virtAddr, /* virtual address */ void ** physAddr /* place to put result */ )		
DESCRIPTION	vmBaseArch32Translate retrieves mapping information for a 32bit virtual address from the page translation tables. If the specified virtual address has never been mapped, the returned status is <b>ERROR</b> . This routine is callable from interrupt level.		
RETURNS	<b>OK</b> , or <b>ERROR</b> if validation or translation fails.		

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VxWorks Kernel API Reference, 6.6 vmBaseArch36LibInit()

ERRNO Not Available

SEE ALSO vmBaseArch32Lib

#### vmBaseArch36LibInit( )

NAME vmBaseArch36LibInit() – initialize the arch specific bundled VM library

SYNOPSIS void vmBaseArch36LibInit (void)

**DESCRIPTION** This routine links the arch specific bundled VM library into the VxWorks system. It is called automatically when INCLUDE\_MMU\_BASIC and INCLUDE\_MMU\_P6\_36BIT are both defined in the BSP.

RETURNS N/A

ERRNO Not Available

SEE ALSO vmBaseArch36Lib

#### vmBaseArch36Map()

NAME vmBaseArch36Map() – map 36bit physical to the 32bit virtual memory

**DESCRIPTION** vmBaseArch36Map maps 36bit physical pages into a contiguous block of 32bit virtual memory. *virtAddr* and *physAddr* must be on page boundaries, and *len* must be evenly divisible by the page size. After the mapping the specified state is set to all pages in the newly mapped virtual memory.

This routine should not be called from interrupt level.
**RETURNS** OK, or ERROR if *virtAddr* or *physAddr* are not on page boundaries, *len* is not a multiple of the page size, the validation fails, or the mapping fails.

ERRNO S\_vmLib\_NOT\_PAGE\_ALIGNED

SEE ALSO vmBaseArch36Lib

### vmBaseArch36Translate( )

NAME	vmBaseArch36Translate() – translate a 32bit virtual address to a 36bit physical address
SYNOPSIS	STATUS vmBaseArch36Translate ( void * virtAddr, /* 32bit virtual address */ LL_INT * physAddr /* place to put 36bit result */ )
DESCRIPTION	vmBaseArch36Translate retrieves mapping information for a 32bit virtual address from the page translation tables. If the specified virtual address has never been mapped, the returned status is <b>ERROR</b> . This routine is callable from interrupt level.
RETURNS	OK, or ERROR if validation or translation fails.
ERRNO	Not Available
SEE ALSO	vmBaseArch36Lib

### vmBaseGlobalMapInit()

NAME	<b>vmBaseGlobalMapInit()</b> – initialize global mapping (obsolete)		
SYNOPSIS	VM_CONTEXT_ID v	MBaseGlobalMapInit	
	PHYS_MEM_DE int	SC *pMemDescArray, numDescArrayElements,	<pre>/* pointer to array of mem descs */ /* no. of elements in pMemDescArray</pre>
	*/ BOOL int )	enable, cacheDefault	/* enable virtual memory */ /* default data cache mode */

VxWorks Kernel API Reference, 6.6 vmBasePageSizeGet()

DESCRIPTION	This function will be replaced by <b>vmGlobalMapInit( )</b>	
RETURNS	A pointer to a newly created virtual memory context, or <b>NULL</b> if memory cannot be mapped.	
ERRNO	Not Available	
SEE ALSO	vmGlobalMap, vmBaseLibInit( ), vmGlobalMapInit( )	

### vmBasePageSizeGet( )

NAME vmBasePageSizeGet() – return the MMU page size (obsolete)

**SYNOPSIS** int vmBasePageSizeGet (void)

**DESCRIPTION** This routine is to be replaced by **vmPageSizeGet()**.

**RETURNS** The MMU page size of the current architecture.

ERRNO Not Available

SEE ALSO vmBaseLib, vmPageSizeGet( )

#### vmBaseStateSet( )

NAME	<b>vmBaseStateSet()</b> – change the state of a block of virtual memory (obsolete)
SYNOPSIS	<pre>STATUS vmBaseStateSet   (</pre>
DESCRIPTION	This function will be replaced by <b>vmStateSet( )</b> .
RETURNS	<b>OK</b> , or <b>ERROR</b> if the validation fails, <i>virtAdrs</i> is not on a page boundary, <i>len</i> is not a multiple of the page size, or the architecture-dependent state set fails for the specified virtual address.

ERRNO	<b>S_vmLib_NOT_PAGE_ALIGNED</b> <i>virtualAddr</i> must be aligned on a page boundary.
	<b>S_vmLib_BAD_STATE_PARAM</b> <i>state</i> is not a valid combination of MMU states.
	<b>S_vmLib_BAD_MASK_PARAM</b> <i>stateMask</i> is not a valid combination of MMU state masks.
SEE ALSO	vmBaseLib, vmStateSet( )

### vmContextShow()

NAME	vmContextSho	<b>w( )</b> – display th	e translation table	e for a context		
SYNOPSIS	STATUS vmCont ( VM_CONTEX )	extShow T_ID context	/* VM context -	NULL == cur	rentContex	t */
DESCRIPTION	This routine dis NULL, the curre memory with c the read/write, Only virtual me	splays the transla ent context is dis onsecutive physic /execute status f emory that has it	ation table for a sp played. Output i ical addresses and or both USR and s s valid state bit se	pecified contex s formatted to l the same stat SUP modes as et is displayed	ct. If <i>context</i> show block te. State info well as the l.	t is specified as is of virtual prmation shows cacheablity.
	This routine she	ould be used for	debugging purpo	oses only.		
EXAMPLE	The following example shows the output of <b>vmContextShow( )</b> using the shell's C-interpreter:			hell's		
	-> vmContextS VIRTUAL ADDR	how BLOCK LENGTH	PHYSICAL ADDR	PROT (S/U)	CACHE	SPECIAL
	0x60000000 0x60010000 0x6015c000 0x6056a000 0x6056e000 0x60570000 0x60571000 For the comman	0x00010000 0x0014c000 0x0040e000 0x00004000 0x00002000 0x00001000 0x00001000 nd-interpreter sh	0x60000000 0x60010000 0x6015c000 0x6056a000 0x6056e000 0x60570000 0x60571000 mell, use the vm co	RWX         /           R-X         /           RWX         /	CB-//- CB-//- CB-//- CB-//- CB-/CO/- WT-//- OFF//- md.	    NB
	The second second		(TAT.: 1 . / . V )			

The protection attributes (Read/Write/eXecute) are listed separately for supervisor and user mode (S/U).

VxWorks Kernel API Reference, 6.6 vmGlobalMapInit()

	Cache attributes are listed with the following notation:	
	Attribute	Meaning
	СВ	Copyback
	WT	Write-through
	OFF	Cache disabled
	CO	Coherency enabled
	G	Guarded
	Special attributes	are listed with the following notation:
	Attribute	Meaning
	NB	No-block. See note blow.
	S0-S6	Special attributes 0 to 6. See the Architecure Supplement for usage.
	The no-block attri more information	bute has meaning only on systems where page optimization is used. For see <b>vmPageOptimize()</b> .
AVAILABILITY	This routine is dis	tributed as a component of the bundled virtual memory support option.
RETURNS	<b>OK</b> , or <b>ERROR</b> if t	he virtual memory context is invalid.
ERRNO	Not Available	
SEE ALSO	vmShow	

# vmGlobalMapInit()

NAME	vmGlobalMapInit( )	<ul> <li>initialize global mapping</li> </ul>		
SYNOPSIS	VM_CONTEXT_ID vmGl	obalMapInit		
	, PHYS_MEM_DESC */	* pMemDescArray,	/* pointer to array of mem descs	
	' int	numDescArrayElements,	/* no. of elements in	
	BOOL int	enable, cacheDefault	/* enable virtual memory */ /* default data cache mode */	
DESCRIPTION	This routine creates a contiguous memory configuration, an insi <b>sysLib.c</b> ; the variable	nd installs a virtual memory segment defined in <i>pMemDe</i> tance of <b>PHYS_MEM_DESC</b> ( e is passed to <b>vmGlobalMag</b>	y context with mappings defined for eacl <i>escArray.</i> In the standard VxWorks (called <b>sysPhysMemDesc</b> ) is defined in <b>pInit( )</b> by the system configuration	h

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mechanism.

	This routine is called only once during system initialization. It should never be called by application code.
	If <i>enable</i> is <b>TRUE</b> , the MMU is enabled upon return.
RETURNS	A pointer to a newly created virtual memory context, or <b>NULL</b> if memory cannot be mapped.
ERRNO	Not Available
SEE ALSO	vmGlobalMap, vmBaseLibInit( )

# vmMap()

NAME	<b>vmMap()</b> – map physical space into virtual space		
SYNOPSIS	<pre>STATUS vmMap  (     VM_CONTEXT_ID context,  /* context - NULL == currentContext */     VIRT_ADDR virtualAddr, /* virtual address */     PHYS_ADDR physicalAddr, /* physical address */     msize_t len /* len of virtual and physical spaces */     )</pre>		
DESCRIPTION	This routine maps physical pages into a contiguous block of virtual memory. <i>virtualAddr</i> and <i>physicalAddr</i> must be on page boundaries, and <i>len</i> must be evenly divisible by the page size. After the call to <b>vmMap()</b> , the state of all pages in the the newly mapped virtual memory is valid, accessible in SUP mode, and cacheable. Note: If mapping a particular page within the given range fails , then the pages that have already been mapped is not restored back.		
	If <i>context</i> is specified as <b>NULL</b> , the current virtual memory context is used.		
	The physicalAddr has to be of type ( <b>PHYS_ADDR</b> ) since on some architectures the physical address could represent more than 32 bits.		
	This routine should not be called from interrupt level.		
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>virtualAddr</i> or <i>physicalAddr</i> are not on page boundaries, <i>len</i> is not a multiple of the page size, the validation fails, or the mapping fails.		
ERRNO	S_vmLib_NOT_PAGE_ALIGNED <i>virtualAddr</i> must be aligned on a page boundary.		
SEE ALSO	vmBaseLib		

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	vmPageLock()		
NAME	<b>vmPageLock()</b> – lock the pages.		
SYNOPSIS	<pre>STATUS vmPageLock   (     VM_CONTEXT_ID context, /* context - NULL == currentContext */     VIRT_ADDR virtualAddr, /* virtual address */     msize_t len, /* len of virtual address */     UINT option /* unused. (for future if needed) */   )</pre>		
DESCRIPTION	This routine will lock the pages by using a static TLB entry if possible.		
	If <i>context</i> is specified as <b>NULL</b> , the current virtual memory context is used.		
	This routine should not be called from interrupt level.		
IMPORTANT	The support for this routine is not available on many CPUs and architectures. You should reference your architecture supplement to see if this is available.		
	Locking of the vxWorks image text section is configurable by using INCLUDE_LOCK_TEXT_SECTION.		
	This routine currently only will lock a valid page present in the kernel context. It will return an <b>ERROR</b> if this is not the case. Also if a page is locked, it can no longer can have its state changed; therefore if a call from vmStateSet is made on a locked page it will return an <b>ERROR</b> . Finally, there is an additional errno for vmStateSet if a page is locked: <b>S_mmuLib_TLB_LOCKED_PAGE</b> .		
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>virtualAddr</i> is not on page boundaries, <i>len</i> is not a multiple of the page size or if the locking of the pages cannot be done.		
ERRNO	S_vmLib_FUNCTION_UNSUPPORTED page locking function not supported.		
	<b>S_vmLib_NOT_PAGE_ALIGNED</b> <i>virtualAddr</i> must be aligned on a page boundary.		
	S_mmuLib_TLB_LOCKED_PAGE Already a locked page		
	S_mmuLib_NOT_CONTIGUOUS_ADDR Requires contiguous phys addr		
	S_mmuLib_NOT_CONTIGUOUS_STATE Requires contiguous MMU state		
	S_mmuLib_INVALID_DESCRIPTOR Bad address		

#### S\_mmuLib\_NOT\_GLOBAL\_PAGE Page must be shared by all contexts

S\_mmuLib\_LOCK\_NO\_MORE\_TLB\_RESOURCES No more TLB entries available

SEE ALSO vmBaseLib

#### vmPageMap()

NAME vmPageMap() – map physical space into virtual space SYNOPSIS STATUS vmPageMap ( \*/ \* / PHYS\_ADDR physicalAddr, /\* physical address \* / msize\_t len, /\* len of virtual and physical spaces \*/ UINT stateMask, /\* combination of MMU state masks. \*/ UINT state /\* combination of MMU states. \*/ ) DESCRIPTION This routine maps physical pages into a contiguous block of virtual memory. *virtualAddr* and *physicalAddr* must be on page boundaries, and *len* must be evenly divisible by the page size. After the call to **vmMap()**, the state of all pages in the the newly mapped virtual memory is set to the default value (valid, sup rwx & cache default) if stateMask is passed as NULL, or else it is set to whatever is passed via statMask/state Note: If mapping a particular page within the given range fails , then the pages that have already been mapped is not restored back. If *context* is specified as NULL, the current virtual memory context is used. The physicalAddr has to be of type (PHYS\_ADDR) since on some architectures the physical address could represent more than 32 bits. This routine should not be called from interrupt level. This routine cannot be called via the macro VM\_PAGE\_MAP. RETURNS **OK**, or **ERROR** if *virtualAddr* or *physicalAddr* are not on page boundaries, *len* is not a multiple of the page size, the validation fails, or the mapping fails, or if invalid state or stateMask are passed. ERRNO S\_vmLib\_NOT\_PAGE\_ALIGNED *virtualAddr* must be aligned on a page boundary. SEE ALSO vmBaseLib

vmPage	Optin	mize()
--------	-------	--------

NAME	<b>vmPageOptimize()</b> – Optimize the address range if possible.	
SYNOPSIS	STATUS vmPageOptimize ( VM_CONTEXT_ID context, /* context - NULL == currentContext */ VIRT_ADDR virtualAddr, /* virtual address */ msize_t len, /* len of address range in bytes */ UINT option /* unused. for future if needed */ )	
DESCRIPTION	This routine will try to optimize the passed address range by modifing pages to use MMU page sizes larger than the default one <b>VM_PAGE_SIZE</b> , if possible.	
	If <i>context</i> is specified as NULL, the current virtual memory context is used.	
	This routine should not be called from interrupt level.	
	The support for this routine is not available on many CPUs and architectures.	
IMPORTANT	You should reference your architecture supplement to see if this is available.	
WARNING	One side affect after using this routine, when supported, is that vmStateSet will possibl block with a semaphore except in interrupt where it returns <b>ERROR</b> . To prevent a call t vmStateSet from returning an <b>ERROR</b> when in an ISR you must preempt with a call, for same address range, to vmStateSet setting the special state <b>MMU_ATTR_NO_BLOCK</b> . Th should be treated as a special purpose attribute.	
	There is an additional errno for vmStateSet if optimization is enabled: S_mmuLib_ISR_CALL_BLOCKED - StateSet needed to block because of optimization so returned error. Should have used MMU_ATTR_NO_BLOCK state set on address.	
	Initial optimization of the whole of kernel context can be done by just configuring in INCLUDE_PAGE_SIZE_OPTIMIZATION.	
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>virtualAddr</i> is not on page boundary, <i>len</i> is not a multiple of the page size or if the optimization is not possible.	
ERRNO	S_vmLib_FUNCTION_UNSUPPORTED page optimization function not supported.	
	S_vmLib_NOT_PAGE_ALIGNED <i>virtualAddr</i> must be aligned on a page boundary.	
SEE ALSO	vmBaseLib	

# vmPageSizeGet()

NAME	<b>vmPageSizeGet( )</b> – return the page size
SYNOPSIS	int vmPageSizeGet (void)
DESCRIPTION	This routine returns the architecture-dependent MMU page size. This routine is callable from interrupt level.
RETURNS	The page size of the current architecture.
ERRNO	Not Available
SEE ALSO	vmBaseLib

# vmPageUnlock()

NAME	<b>vmPageUnlock()</b> – unlock the pages.		
SYNOPSIS	<pre>STATUS vmPageUnlock   (</pre>		
DESCRIPTION	This routine will lock the pages that were locked by a previous vmPageLock.		
	If <i>context</i> is specified as NULL, the current virtual memory context is used.		
	This routine should not be called from interrupt level.		
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>virtualAddr</i> is not on page boundaries, if the pages were not previously locked.		
ERRNO	S_vmLib_FUNCTION_UNSUPPORTED page unlocking function not supported.		
	S_vmLib_NOT_PAGE_ALIGNED <i>virtualAddr</i> must be aligned on a page boundary.		
	S_mmuLib_TLB_PAGE_NOT_LOCKED Can only unlock a locked page.		

VxWorks Kernel API Reference, 6.6 vmPhysTranslate()

SEE ALSO vmBaseLib

# vmPhysTranslate()

NAME	<b>vmPhysTranslate()</b> – translate a physical address to a virtual address		
SYNOPSIS	<pre>STATUS vmPhysTranslate   (     VM_CONTEXT_ID context,</pre>		
DESCRIPTION	This routine retrieves mapping information for a physical address from the page translation tables. If <i>context</i> is specified as <b>NULL</b> , the current context is used. The physicalAddr has to be of type ( <b>PHYS_ADDR</b> ) since on some architectures the physical address could represent more than 32 bits. This routine is callable from interrupt level.		
RETURNS	OK, or ERROR if the validation or translation failed.		
ERRNO	Not Available		
SEE ALSO	vmBaseLib		

### vmStateGet()

NAME	<b>vmStateGet( )</b> – get the state of a page of virtual memory
SYNOPSIS	<pre>STATUS vmStateGet   (     VM_CONTEXT_ID context, /* VM context; use NULL for current context */     VIRT_ADDR pageAddr, /* virtual page addr */     UINT * pState /* where to return state */   )</pre>
DESCRIPTION	This routine gets the MMU attributes of a page mapped in a virtual memory context. For a description of the supported page attributes see the <b>vmStateSet()</b> API guide.

If *context* is **NULL**, the current virtual memory context is used.

This routine is callable from interrupt level.

For example, to see if a page is writable in supervisor mode, the following code may be used:

```
if (vmStateGet (context, pageAddr, &attr) == OK)
{
    if (((attr & MMU_ATTR_VALID_MSK) == MMU_ATTR_VALID) &&
        ((attr & MMU_ATTR_PROT_SUP_WRITE) == MMU_ATTR_PROT_SUP_WRITE))
        ...
```

- **RETURNS** OK, or ERROR if *pageAddr* is not on a page boundary, the validity check fails, or the architecture-dependent state get fails for the specified virtual address.
- ERRNO S\_vmLib\_NOT\_PAGE\_ALIGNED pageAddr is not aligned on a page boundary.
- SEE ALSO vmBaseLib, vmStateSet()

#### vmStateSet()

NAME	<b>vmStateSet( )</b> – change	e the state of	a block of virtual memory
SYNOPSIS	STATUS vmStateSet ( VM_CONTEXT_ID cc VIRT_ADDR vi msize_t le UINT st UINT st )	ontext, irtAdrs, en, tateMask, tate	/* context - NULL == currentContext */ /* virtual address to modify state of */ /* len of virtual space to modify state of */ /* state mask */ /* state */
DESCRIPTION	This routine changes the MMU attributes of a block of virtual memory. Each page of virtual memory has at least three types of state information: validity, protection, and cacheability. Some architectures define additional state information. The following MMU attributes are supported and may be OR'ed together in the <i>state</i> parameter:		
Protection attributes:			
	Attribute		Description
	MMU_ATTR_PROT_SU MMU_ATTR_PROT_SU MMU_ATTR_PROT_SU MMU_ATTR_PROT_US MMU_ATTR_PROT_US MMU_ATTR_PROT_US	P_READ P_WRITE P_EXE R_READ R_WRITE R_EXE	read access in supervisor mode write access in supervisor mode executable access in supervisor mode read access in user mode write access in user mode executable access in user mode

Validity attribute. Memory accesses to a page set invalid will result in an exception.

Attribute	Description
MMU_ATTR_VALID	page is valid
MMU_ATTR_VALID_NOT	page is not valid

Cache attributes:

Attribute	Description
MMU_ATTR_CACHE_OFF	cache turned off
MMU_ATTR_CACHE_COPYBACK	cache in copy-back mode
MMU_ATTR_CACHE_WRITETHRU	cache set in writethrough mode
MMU_ATTR_CACHE_DEFAULT	default cache value, USER_D_CACHE_MODE
MMU_ATTR_CACHE_GUARDED	page access set to guarded
MMU_ATTR_CACHE_COHERENCY	page access set to cache coherent

The *stateMask* parameter is used to specify which MMU attribute groups are being modified. This should be an inclusive OR of one or more of the following masks:

Mask	Description
MMU_ATTR_PROT_MSK	set protection attributes
MMU_ATTR_VALID_MSK	set valid attribute
MMU_ATTR_CACHE_MSK	set cache attributes
MMU_ATTR_SPL_MSK	set architecture specific attributes

The following restrictions must be respected when setting page attributes:

- only one of MMU\_ATTR\_CACHE\_OFF, MMU\_ATTR\_CACHE\_COPYBACK, MMU\_ATTR\_CACHE\_WRITETHRU or MMU\_ATTR\_CACHE\_DEFAULT can be set at any time.
- not all combinations of the protection attributes are supported by various architectures. For more information see the respective Architecture Supplement documentation.

Refer to the archecture specific mmuLib man pages for specific details.

If *context* is NULL, the current context is used.

This routine is callable from interrupt level.

# **RETURNS** OK, or ERROR if the validation fails, *virtAdrs* is not on a page boundary, *len* is not a multiple of the page size, or the architecture-dependent state set fails for the specified virtual address.

#### ERRNO S\_vmLib\_NOT\_PAGE\_ALIGNED

*virtAdrs* is not aligned on a page boundary.

#### S\_vmLib\_BAD\_STATE\_PARAM

state is not a valid combination of MMU attributes.

#### S\_vmLib\_BAD\_MASK\_PARAM

stateMask is not a valid combination of MMU attribute masks.

SEE ALSO vmBaseLib, vmStateGet( )

### vmTextProtect()

NAME	<b>vmTextProtect()</b> – write-protect kernel text segment	
SYNOPSIS	STATUS vmTextProtect ( BOOL setState )	
DESCRIPTION	This routine enables write-protection of text segments in the VxWorks kernel. This function should not be called by application code; instead, this routine is called automatically at boot time when the <b>INCLUDE_PROTECT_TEXT</b> component is included.	
	If the start of the text segment is not page aligned, text protection starts from the next page boundary. This routine expects that the data segment follows the text segment after a proper alignment padding. The VxWorks build system ensures this condition except for ROM-resident images (i.e. images for which the text segment is directly executed out of ROM or flash memory). For ROM-resident images, if detection of an attempt to write in the ROM is desired, protection should be enabled via the corresponding sysPhysMemDesc[] entry in the BSP.	
	This routine is not setting protection attributes for the MIPS architecture. Text protection for MIPS, when <b>INCLUDE_PROTECT_TEXT</b> is included, is enabled when the initial mappings are created.	
RETURNS	OK, or ERROR if the text segment cannot be write-protected.	
ERRNO	S_vmLib_TEXT_PROTECTION_UNAVAILABLE write-protecting the kernel text segment is not supported.	
SEE ALSO	vmBaseLib	

### vmTranslate()

NAME	<b>vmTranslate( )</b> – translate a virtual address to a physical address		
SYNOPSIS	STATUS vmTranslate		
	<pre>VM_CONTEXT_ID context,</pre>		

	VIRT_ADDR virtualAddr, /* virtual address */ PHYS_ADDR * physicalAddr /* place to put result */ )
DESCRIPTION	This routine retrieves mapping information for a virtual address from the page translation tables. If <i>context</i> is specified as <b>NULL</b> , the current context is used.
	The physicalAddr has to be of type (PHYS_ADDR $^*$ ) since on some architectures the physical address could represent more than 32 bits.
	This routine is callable from interrupt level.
RETURNS	OK, or ERROR if the validation or translation failed.
ERRNO	Not Available
SEE ALSO	vmBaseLib

# voprintf()

NAME	<b>voprintf()</b> – write a formatted string to an output function
SYNOPSIS	<pre>int voprintf  (  FUNCPTR prtFunc, /* pointer to output function */  int prtArg, /* argument for output function */  const char * fmt, /* format string to write */  va_list vaList /* optional arguments to format */ )</pre>
DESCRIPTION	This routine prints a formatted string via the function specified by <i>prtFunc</i> . The function will receive as parameters a pointer to a buffer, an integer indicating the length of the buffer, and the argument <i>prtArg</i> . If <b>NULL</b> is specified as the output function, the output will be sent to stdout.
	This routine is identical to <b>oprintf()</b> , except that it takes the variable arguments to be formatted as a list <i>vaList</i> of type <b>va_list</b> rather than as in-line arguments.
SMP CONSIDERATIONS	
	This API is spinlock and intCpuLock restricted.
RETURNS	The number of characters output, not including the NULL terminator.
ERRNO	Not Available

SEE ALSO fioLib, oprintf(), printf()

# vprintf()

NAME	<b>vprintf()</b> – write a string formatted with a variable argument list to standard output (ANSI)	
SYNOPSIS	<pre>int vprintf (     const char * fmt, /* format string to write */     va_list vaList /* arguments to format */ )</pre>	
DESCRIPTION	This routine prints a string formatted with a variable argument list to standard output. It is identical to <b>printf()</b> , except that it takes the variable arguments to be formatted as a list <i>vaList</i> of type <b>va_list</b> rather than as in-line arguments.	
SMP CONSIDERATIONS		
	This API is spinlock and intCpuLock restricted.	
RETURNS	The number of characters output, or <b>ERROR</b> if there is an error during output.	
ERRNO	Not Available	
SEE ALSO	<b>fioLib</b> , <b>printf()</b> , American National Standard for Information Systems -, Programming Language - C, ANSI X3.159-1989: Input/Output ( <b>stdio.h</b> )	

### vrfsDevCreate()

NAME	vrfsDevCreate() – Instantiate the VRFS
SYNOPSIS	STATUS vrfsDevCreate ( )
DESCRIPTION	This routine creates an instance of the VRFS if one does not exist. That instance will always be installed as device "/".
RETURNS	OK on success, ERROR if failure
ERRNO	EEXIST if this FS is already instantiated.

SEE ALSO vrfsLib

### vrfsInit()

NAME	<b>vrfsInit( )</b> – Initialize the Virtual Root File System Library
SYNOPSIS	STATUS vrfsInit(void)
DESCRIPTION	This routine initializes the Virtual Root File System. It should be called only once, and initializes the vrfs Core IO driver as well as data structures for the library.
RETURNS	Not Available
ERRNO	Not Available
SEE ALSO	vrfsLib

# vsnprintf()

NAME	<b>vsnprintf()</b> – write a string formatted with a variable argument list to a buffer, not exceeding buffer size (ANSI)
SYNOPSIS	<pre>int vsnprintf  (     char * buffer, /* buffer to write to */     size_t count, /* max number of characters to store in buffer */     const char * fmt, /* format string */     va_list vaList /* optional arguments to format */     )</pre>
DESCRIPTION	This routine copies a string formatted with a variable argument list to a specified buffer, up to a given number of characters. The formatted string will be null terminated. This routine guarantees never to write beyond the provided buffer regardless of the format specifier or the arguments to be formatted. The <i>count</i> argument specifies the maximum number of characters to store in the buffer, including the null terminator.
	This routine is identical to <b>snprintf()</b> , except that it takes the variable arguments to be formatted as a list <i>vaList</i> of type <b>va_list</b> rather than as in-line arguments.
RETURNS	The number of characters copied to <i>buffer</i> , not including the NULL terminator.

Even when the supplied *buffer* is too small to hold the complete formatted string, the return value represents the number of characters that would have been written to *buffer* if *count* was sufficiently large.

ERRNO Not Available

**SEE ALSO** fioLib, sprintf(), printf(), "International Organization for Standardization, ISO/IEC 9899:1999, Programming languages - C: Input/output (stdio.h)"

#### vsprintf()

NAME **vsprintf()** – write a string formatted with a variable argument list to a buffer (ANSI) SYNOPSIS int vsprintf ( char \* buffer, /\* buffer to write to \*/ const char \* fmt, /\* format string \*/ va\_list vaList /\* optional arguments to format \*/ ) DESCRIPTION This routine copies a string formatted with a variable argument list to a specified buffer. This routine is identical to **sprintf()**, except that it takes the variable arguments to be formatted as a list *vaList* of type **va\_list** rather than as in-line arguments. The number of characters copied to *buffer*, not including the NULL terminator. RETURNS Not Available ERRNO SEE ALSO fioLib, sprintf(), American National Standard for Information Systems -, Programming Language - C, ANSI X3.159-1989: Input/Output (stdio.h)

#### vxAtomicAdd()

NAME	<b>vxAtomicAdd()</b> – atomically add a value to a memory location
SYNOPSIS	atomicVal_t vxAtomicAdd (
	atomic_t * target, /* memory location to add to */ atomicVal_t value /* value to add */ )

VxWorks Kernel API Reference, 6.6 vxAtomicAnd()

DESCRIPTION	This routine atomically adds <i>*target</i> and <i>value</i> , placing the result in <i>*target</i> . The operation is done using signed integer arithmetic. Various CPU architectures may impose restrictions with regards to the alignment and cache attributes of the atomic_t type.
	This routine can be used from both task and interrupt level.
RETURNS	Contents of <i>*target</i> before the atomic operation
ERRNO	N/A
SEE ALSO	vxAtomicLib

### vxAtomicAnd()

NAME	<b>vxAtomicAnd()</b> – atomically perform a bitwise AND on a memory location
SYNOPSIS	<pre>atomicVal_t vxAtomicAnd   (    atomic_t * target, /* memory location to AND */    atomicVal_t value /* AND with this value */   )</pre>
DESCRIPTION	This routine atomically performs a bitwise AND operation of <i>*target</i> and <i>value</i> , placing the result in <i>*target</i> . Various CPU architectures may impose restrictions with regards to the alignment and cache attributes of the atomic_t type. This routine can be used from both task and interrupt level.
RETURNS	Contents of <i>*target</i> before the atomic operation
ERRNO	N/A
SEE ALSO	vxAtomicLib

### vxAtomicClear()

NAME vxAtomicClear() – atomically clear a memory location

SYNOPSIS atomicVal\_t vxAtomicClear

	( atomic_t * target /* memory location to clear */ )
DESCRIPTION	This routine atomically clears <i>*target</i> and returns the old value that was in <i>*target</i> . Note that all CPU architectures supported by VxWorks can atomically clear a variable of size atomic_t without the need to use this routine. This routine is intended for software that needs to atomically fetch and clear the value of a memory location. Various CPU architectures may impose restrictions with regards to the alignment and cache attributes of the atomic_t type.
	This routine can be used from both task and interrupt level.
RETURNS	Contents of <i>*target</i> before the atomic operation
ERRNO	N/A
SEE ALSO	vxAtomicLib

# vxAtomicDec()

NAME	<b>vxAtomicDec()</b> – atomically decrement a memory location
SYNOPSIS	atomicVal_t vxAtomicDec ( atomic_t * target /* memory location to decrement */ )
DESCRIPTION	This routine atomically decrements the value in <i>*target</i> . The operation is done using unsigned integer arithmetic. Various CPU architectures may impose restrictions with regards to the alignment and cache attributes of the atomic_t type. This routine can be used from both task and interrupt level.
RETURNS	Contents of <i>*target</i> before the atomic operation
ERRNO	N/A
SEE ALSO	vxAtomicLib

## vxAtomicGet()

NAME	<b>vxAtomicGet( )</b> – atomically get a memory location
SYNOPSIS	atomicVal_t vxAtomicGet ( atomic_t * target /* memory location to get */ )
DESCRIPTION	This routine atomically reads <i>*target</i> and returns the value. This routine is intended for software that needs to atomically fetch and replace the value of a memory location. This routine can be used from both task and interrupt level.
RETURNS	Contents of *target.
ERRNO	N/A
SEE ALSO	vxAtomicLib

# vxAtomicInc()

NAME	<b>vxAtomicInc()</b> – atomically increment a memory location
SYNOPSIS	<pre>atomicVal_t vxAtomicInc    (    atomic_t * target /* memory location to increment */ )</pre>
DESCRIPTION	This routine atomically increments the value in <i>*target</i> . The operation is done using unsigned integer arithmetic. Various CPU architectures may impose restrictions with regards to the alignment and cache attributes of the atomic_t type. This routine can be used from both task and interrupt level.
RETURNS	Contents of <i>*target</i> before the atomic operation
ERRNO	N/A
SEE ALSO	vxAtomicLib

## vxAtomicNand()

NAME	<b>vxAtomicNand()</b> – atomically perform a bitwise NAND on a memory location
SYNOPSIS	atomicVal_t vxAtomicNand ( atomic_t * target, /* memory location to NAND */ atomicVal_t value /* NAND with this value */ )
DESCRIPTION	This routine atomically performs a bitwise NAND operation of <i>*target</i> and <i>value</i> , placing the result in <i>*target</i> . Various CPU architectures may impose restrictions with regards to the alignment and cache attributes of the atomic_t type.
	This routine can be used from both task and interrupt level.
RETURNS	Contents of *target before the atomic operation
ERRNO	N/A
SEE ALSO	vxAtomicLib

# vxAtomicOr()

NAME	<b>vxAtomicOr()</b> – atomically perform a bitwise OR on memory location
SYNOPSIS	<pre>atomicVal_t vxAtomicOr   (    atomic_t * target, /* memory location to OR */    atomicVal_t value /* OR with this value */   )</pre>
DESCRIPTION	This routine atomically performs a bitwise OR operation of <i>*target</i> and <i>value</i> , placing the result in <i>*target</i> . Various CPU architectures may impose restrictions with regards to the alignment and cache attributes of the atomic_t type. This routine can be used from both task and interrupt level.
RETURNS	Contents of <i>*target</i> before the atomic operation
ERRNO	N/A
SEE ALSO	vxAtomicLib

# vxAtomicSet()

NAME	vxAtomicSet() – atomically set a memory location
SYNOPSIS	<pre>atomicVal_t vxAtomicSet   (    atomic_t * target, /* memory location to set */    atomicVal_t value /* set with this value */   )</pre>
DESCRIPTION	This routine atomically sets <i>*target</i> to <i>value</i> and returns the old value that was in <i>*target</i> . Note that all CPU architectures supported by VxWorks can atomically write to a variable of size atomic_t without the need to use this routine. This routine is intended for software that needs to atomically fetch and replace the value of a memory location. Various CPU architectures may impose restrictions with regards to the alignment and cache attributes of the atomic_t type.
	This routine can be used from both task and interrupt level.
RETURNS	Contents of <i>*target</i> before the atomic operation
ERRNO	N/A
SEE ALSO	vxAtomicLib

### vxAtomicSub()

NAME	<b>vxAtomicSub()</b> – atomically subtract a value from a memory location
SYNOPSIS	<pre>atomicVal_t vxAtomicSub   (    atomic_t * target, /* memory location to subtract from */    atomicVal_t value /* value to sub */   )</pre>
DESCRIPTION	This routine atomically subtracts <i>value</i> from <i>*target</i> , placing the result in <i>*target</i> . The operation is done using signed integer arithmetic. Various CPU architectures may impose restrictions with regards to the alignment and cache attributes of the atomic_t type.
	This routine can be used from both task and interrupt level.
RETURNS	Contents of <i>*target</i> before the atomic operation

ERRNO N/A

SEE ALSO vxAtomicLib

### vxAtomicXor()

NAME	<b>vxAtomicXor()</b> – atomically perform a bitwise XOR on a memory location
SYNOPSIS	atomicVal_t vxAtomicXor ( atomic_t * target, /* memory location to XOR */ atomicVal_t value /* XOR with this value */ )
DESCRIPTION	This routine atomically performs a bitwise XOR operation of <i>*target</i> and <i>value</i> , placing the result in <i>*target</i> . Various CPU architectures may impose restrictions with regards to the alignment and cache attributes of the atomic_t type. This routine can be used from both task and interrupt level.
RETURNS	Contents of <i>*target</i> before the atomic operation
ERRNO	N/A
SEE ALSO	vxAtomicLib

## vxCas()

NAME	vxCas() – atomically compare-and-swap the contents of a memory location
SYNOPSIS	<pre>BOOL vxCas   (    atomic_t * target, /* memory location to compare-and-swap */    atomicVal_t oldValue, /* compare to this value */    atomicVal_t newValue /* swap with this value */   )</pre>
DESCRIPTION	This routine performs an atomic compare-and-swap; testing that <i>*target</i> contains <i>oldValue</i> , and if it does, setting the value of <i>*target</i> to <i>newValue</i> . Various CPU architectures may impose restrictions with regards to the alignment and cache attributes of the atomic_t type.

This routine can be used from both task and interrupt level.

VxWorks Kernel API Reference, 6.6 vxCpuConfiguredGet()

**RETURNS** TRUE if the swap is actually executed, FALSE otherwise.

ERRNO N/A

SEE ALSO vxAtomicLib

#### vxCpuConfiguredGet( )

NAME vxCpuConfiguredGet() – get the number of configured CPUs in the system

SYNOPSIS unsigned int vxCpuConfiguredGet (void)

**DESCRIPTION** This routine returns the number of CPUs that have been configured in the SMP system, whether they have been enabled or not. This number is set at compile time and stays constant for as long as the system is up and running. This routine can therefore be called at any time, even during the booting sequence of the system. Its purpose is to assist initialization code of a kernel application in determining how many per-CPU objects would need to be allocated in an SMP system.

This routine exists because VxWorks SMP has the flexibility to allow the number of CPUs configured in a VxWorks SMP system to be different than the number of available CPUs on the hardware platform. For example, it would be possible to dedicate two cores of a quad-core platform to run VxWorks SMP while the other two cores are used for another purpose.

Calling this routine in the uniprocessor version of VxWorks returns 1, always. This routine can be called from both task an interrupt level.

- **RETURNS** The number of CPUs configured in the system.
- ERRNO N/A
- SEE ALSO vxCpuLib, vxCpuEnabledGet()

#### vxCpuEnabledGet()

NAME	vxCpuEnabledGet( )	– get a set of running	g CPUs
------	--------------------	------------------------	--------

**SYNOPSIS** cpuset\_t vxCpuEnabledGet (void)

DESCRIPTION This routine returns the set of CPUs that are running in the VxWorks SMP system. This set is updated at run-time as CPUs are enabled by the bootstrap CPU but the number of CPUs in the set can never be larger than the number of CPUs configured in the system. That is, the number of CPUs in the set cannot exceed the value returned by vxCpuConfiguredGet().

The default behaviour of VxWorks SMP is to take all configured CPUs out of reset at boot time. However this behaviour can be modified to only enable additional CPUs at a later point in time. This routine can therefore be used to obtain a true representation of the enabled CPUs as opposed to the number of configured CPUs.

Calling this routine in the uniprocessor version of VxWorks always returns a set that shows CPU0 as being the only enabled CPU. The coding example below shows a test case that could be used to test the expected behaviour of this routine in a uniprocessor environment.

```
STATUS test (void)
{
cpuset_t uniprocessorCpuSet;
/* Get the set of enabled CPUs */
uniprocessorCpuSet = vxCpuEnabledGet();
/* CPU 0 is supposed to be enabled. Check it! */
if (CPUSET_ISSET(uniprocessorCpuSet, 0))
   {
   /*
    * First part of the test passed. Now check that no other CPUs
    * are in the set.
    */
   CPUSET CLR(uniprocessorCpuSet, 0);
   if (CPUSET ISZERO(uniprocessorCpuSet))
       /* No other CPUs in the set. Test passed. */
       return (OK);
       }
   }
/*
 * Test failed. Either CPU 0 was not in the set or other CPUs
 * were in the set.
 */
return (ERROR);
}
This routine can be called from both task or interrupt level.
A set of CPUs that have been enabled.
N/A
vxCpuLib, vxCpuConfiguredGet(), cpuset
```

RETURNS

ERRNO

SEE ALSO

	vxCpuIndexGet()	
NAME	<b>vxCpuIndexGet( )</b> – get the index of the calling CPU	
SYNOPSIS	unsigned int vxCpuIndexGet (void)	
DESCRIPTION	This routine returns the index of the CPU on which the calling task or ISR is running. The index is a number between 0 and N-1, where N is the number of CPUs configured in the SMP system. N is also the figure returned by <b>vxCpuConfiguredGet()</b> . Calling this routine in the uniprocessor version of VxWorks returns 0, always. The value returned by this routine can easily be used as an index into an array of per-CPU objects that would have previously been allocated with the help of the <b>vxCpuConfiguredGet()</b> routine.	
	Since tasks can migrate from one CPU to another in an SMP system, no guarantees are provided that the index is valid by the time program execution returns to the caller of this routine. For example, if a scheduling event takes place immediately after this call returns, it is possible for the caller to be running on a different CPU than it was at the time of the call. It is the responsibility of the caller to prevent task migration to another CPU while the index is being used. This can be done using <b>taskCpuLock()</b> . If this routine is called from interrupt context the caller is guaranteed the index will be valid until the ISR returns. This is because ISRs do not migrate from one CPU to the other while they are running.	
	The purpose of this routine is different than that of <b>sysProcNumGet()</b> , which is used to uniquely identify a node in an asymmetric multiprocessing environment. Should a node running VxWorks SMP exist in such an environment, <b>sysProcNumGet()</b> on that node would always return the same value regardless of the CPU on which the calling thread is running.	
RETURNS	The index of the CPU on which the calling thread executes	
ERRNO	N/A	
SEE ALSO	vxCpuLib, vxCpuConfiguredGet()	

### vxCr0Get()

NAME	<b>vxCr0Get()</b> – get a content of the Control Register 0 (x86)
SYNOPSIS	int vxCr0Get (void)
DESCRIPTION	This routine gets a content of the Control Register 0.

ERRNO Not Available

SEE ALSO vxLib

# vxCr0Set()

NAME	<b>vxCr0Set( )</b> – set a value to the Control Register 0 (x86)
SYNOPSIS	<pre>void vxCr0Set   (    int value /* CR0 value */ )</pre>
DESCRIPTION	This routine sets a value to the Control Register 0.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	vxLib

### vxCr2Get()

NAME	<b>vxCr2Get()</b> – get a content of the Control Register 2 (x86)
SYNOPSIS	int vxCr2Get (void)
DESCRIPTION	This routine gets a content of the Control Register 2.
RETURNS	a value of the Control Register 2
ERRNO	Not Available
SEE ALSO	vxLib

2

#### vxCr2Set()

NAME	<b>vxCr2Set( )</b> – set a value to the Control Register 2 (x86)
SYNOPSIS	<pre>void vxCr2Set   (    int value /* CR2 value */ )</pre>
DESCRIPTION	This routine sets a value to the Control Register 2.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	vxLib

#### vxCr3Get()

NAME	<b>vxCr3Get()</b> – get a content of the Control Register 3 (x86)
SYNOPSIS	int vxCr3Get (void)
DESCRIPTION	This routine gets a content of the Control Register 3.
RETURNS	a value of the Control Register 3
ERRNO	Not Available
SEE ALSO	vxLib

### vxCr3Set()

 NAME
 vxCr3Set() - set a value to the Control Register 3 (x86)

```
SYNOPSIS void vxCr3Set
(
int value /* CR3 value */
)
```

**DESCRIPTION** This routine sets a value to the Control Register 3.

RETURNS N/A

ERRNO Not Available

SEE ALSO vxLib

#### vxCr4Get()

 NAME
 vxCr4Get() - get a content of the Control Register 4 (x86)

 SYNOPSIS
 int vxCr4Get (void)

**DESCRIPTION** This routine gets a content of the Control Register 4.

- **RETURNS** a value of the Control Register 4
- ERRNO Not Available

SEE ALSO vxLib

### vxCr4Set()

NAME	<b>vxCr4Set( )</b> – set a value to the Control Register 4 (x86)
SYNOPSIS	<pre>void vxCr4Set   (    int value /* CR4 value */ )</pre>
DESCRIPTION	This routine sets a value to the Control Register 4.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	vxLib

### vxDrGet()

**NAME vxDrGet()** – get a content of the Debug Register 0 to 7 (x86)

SYNOPSIS

/
/
/
/
/
/
/
/

**DESCRIPTION** This routine gets a content of the Debug Register 0 to 7.

RETURNS	N/A
ERRNO	Not Available
SEE ALSO	vxLib

## vxDrSet()

NAME	<b>vxDrSet( )</b> – set a value to the Debug Register 0 to 7 (x86)
SYNOPSIS	<pre>void vxDrSet     (     int dr0,  /* DR0 */     int dr1,  /* DR1 */     int dr2,  /* DR2 */     int dr3,  /* DR3 */     int dr4,  /* DR4 */     int dr5,  /* DR5 */     int dr6,  /* DR6 */     int dr7    /* DR7 */ )</pre>
DESCRIPTION	This routine sets a value to the Debug Register 0 to 7.
RETURNS	N/A
ERRNO	Not Available

SEE ALSO vxLib

### vxEflagsGet()

**NAME** vxEflagsGet() – get a content of the EFLAGS register (x86)

SYNOPSIS int vxEflagsGet (void)

**DESCRIPTION** This routine gets a content of the EFLAGS register

- **RETURNS** a value of the EFLAGS register
- ERRNO Not Available
- SEE ALSO vxLib

### vxEflagsSet()

NAME	<b>vxEflagsSet( )</b> – set a value to the EFLAGS register (x86)
SYNOPSIS	<pre>void vxEflagsSet   (    int value /* EFLAGS value */ )</pre>
DESCRIPTION	This routine sets a value to the EFLAGS register
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	vxLib

### vxGdtrGet()

NAME	<b>vxGdtrGet()</b> – get a content of the Global Descriptor Table Register (x86)
SYNOPSIS	<pre>void vxGdtrGet   (    long long int * pGdtr /* memory to store GDTR */ )</pre>
DESCRIPTION	This routine gets a content of the Global Descriptor Table Register
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	vxLib

### vxIdtrGet()

NAME	<b>vxIdtrGet( )</b> – get a content of the Interrupt Descriptor Table Register (x86)
SYNOPSIS	<pre>void vxIdtrGet   (    long long int * pIdtr /* memory to store IDTR */ )</pre>
DESCRIPTION	This routine gets a content of the Interrupt Descriptor Table Register
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	vxLib

# vxLdtrGet()

NAME vxLdtrGet() – get a content of the Local Descriptor Table Register (x86)

SYNOPSIS void vxLdtrGet

```
(
long long int * pLdtr /* memory to store LDTR */
)
```

**DESCRIPTION** This routine gets a content of the Local Descriptor Table Register

RETURNS N/A

ERRNO Not Available

SEE ALSO vxLib

#### vxMemArchProbe()

**NAME vxMemArchProbe()** – architecture specific part of vxMemProbe

SYNOPSIS STATUS vxMemArchProbe

(		
FAST char * adrs,	/* address to be probed	*/
int mode,	/* VX_READ or VX_WRITE	*/
int length,	/* 1, 2, 4, or 8	*/
FAST char * pVal	<pre>/* where to return value,</pre>	*/
	/* or ptr to value to be written	*/
)		

- **DESCRIPTION** This is the routine implementing the architecture specific part of the vxMemProbe routine. It traps the relevant exceptions while accessing the specified address. If an exception occurs, then the result will be **ERROR**. If no exception occurs then the result will be **OK**.
- **RETURNS** OK or ERROR if an exception occurred during access.
- ERRNO Not Available

SEE ALSO vxLib

#### vxMemProbe()

NAME vxMemProbe() – probe an address for a bus error

SYNOPSIS STATUS vxMemProbe

```
(
                      FAST char * adrs,/* address to be probedintmode,/* VX_READ or VX_WRITEintlength,/* 1, 2, 4, or 8
                                                                                     */
                                                                                     */
                                                                                     */
                      FAST char * pVal /* where to return value,
                                                                                     */
                                              /* or ptr to value to be written */
                      )
DESCRIPTION
                 This routine probes a specified address to see if it is readable or writable, as specified by
                  mode. The address is read or written as 1, 2, or 4 bytes, as specified by length (values other
                  than 1, 2, or 4 yield unpredictable results). If the probe is a VX_READ (0), the value read is
                  copied to the location pointed to by pVal. If the probe is a VX_WRITE (1), the value written
                  is taken from the location pointed to by pVal. In either case, pVal should point to a value of
                  1, 2, or 4 bytes, as specified by length.
                  Note that only bus errors are trapped during the probe, and that the access must otherwise
                  be valid (i.e., it must not generate an address error).
EXAMPLE
                      testMem (adrs)
                           char *adrs;
                           {
                           char testW = 1;
                           char testR;
                           if (vxMemProbe (adrs, VX_WRITE, 1, &testW) == OK)
                               printf ("value %d written to adrs %x\en", testW, adrs);
                           if (vxMemProbe (adrs, VX READ, 1, &testR) == OK)
                               printf ("value %d read from adrs %x\en", testR, adrs);
                           }
                 The BSP can modify the behaviour of vxMemProbe() by supplying an alternate routine and
MODIFICATION
                  placing the address in the global variable func vxMemProbeHook. The BSP routine will
                  be called instead of the architecture specific routine vxMemArchProbe().
                  OK, or ERROR if the probe caused a bus error or was misaligned.
RETURNS
                 Not Available
ERRNO
                 vxLib, vxMemArchProbe()
SEE ALSO
```

#### vxMemProbe()

NAME vxMemProbe() – probe an address for a bus error

SYNOPSIS STATUS vxMemProbe

```
(
                      FAST char * pAdrs,/* address to be probed */intmode,/* VX_READ or VX_WRITE */intlength,/* 1, 2, or 4 */
                      char *
                                  pVal
                                              /* Data source if VX_WRITE; destination if VX_READ
                  * /
                      )
DESCRIPTION
                  This routine probes a specified address to see if it is readable or writable, as specified by
                  mode. The address will be read or written according to the requested length. The provided
                  pointer must be naturally aligned to the requested length.
                  If the requested mode is VX_READ, the value read will be copied to the location pointed to
                  by pVal. If the requested mode is VX_WRITE, the value written will be taken from the
                  location pointed to by pVal. In either case, pVal should point to a value of length length.
                  Note that only data bus errors (machine check exception, data access exception) are trapped
                  during the probe, and that the access must be otherwise valid (i.e., not generate an address
                  error).
EXAMPLE
                  testMem (adrs)
                     char *adrs;
                     {
                     char testW = 1;
                     char testR;
                     if (vxMemProbe (adrs, VX_WRITE, 1, &testW) == OK)
                          printf ("value %d written to adrs %x\en", testW, adrs);
                     if (vxMemProbe (adrs, VX_READ, 1, &testR) == OK)
                          printf ("value %d read from adrs %x\en", testR, adrs);
                     }
                  The BSP can modify the behaviour of this routine by supplying an alternate routine and
MODIFICATION
                  placing the address of the routine in the global variable _func_vxMemProbeHook. The BSP
                  routine will be called instead of the architecture specific routine vxMemArchProbe().
RETURNS
                   OK if the probe is successful, or
                   ERROR if the probe caused a bus error.
                  Not Available
ERRNO
                 vxMemProbeLib, vxMemArchProbe()
SEE ALSO
```

### vxMemProbeInit()

NAMEvxMemProbeInit() - add vxMemProbeTrap exception handler to exc handler chainSYNOPSISSTATUS vxMemProbeInit (void)DESCRIPTIONAdd the vxMemProbe exception handler hook to the exception handler chain called by<br/>excExcHandleRETURNSOK if initialization OK else ERRORERRNONot AvailableSEE ALSOvxMemProbeLib

#### vxPowerDown()

NAME	<b>vxPowerDown()</b> – place the processor in reduced-power mode (PowerPC, SH)
SYNOPSIS	UINT32 vxPowerDown (void)
DESCRIPTION	This routine activates the reduced-power mode if power management is enabled. It is called by the scheduler when the kernel enters the idle loop. The power management mode is selected by <b>vxPowerModeSet()</b> .
RETURNS	OK, or ERROR if power management is not supported or if external interrupts are disabled.
ERRNO	Not Available
SEE ALSO	vxLib, vxPowerModeSet( ), vxPowerModeGet( )

### vxPowerModeGet()

NAME	<b>vxPowerModeGet( )</b> – get the power management mode (PowerPC, SH, x86)
SYNOPSIS	UINT32 vxPowerModeGet (void)
DESCRIPTION	This routine returns the power management mode set by $\mathbf{vxPowerModeSet(}$ ).
**RETURNS** The power management mode, or **ERROR** if no mode has been selected or if power management is not supported.

 **ERRNO** Not Available

SEE ALSO vxLib, vxPowerModeSet(), vxPowerDown()

## vxPowerModeSet()

NAME	<b>vxPowerModeSet( )</b> – set the power management mode (PowerPC, SH, x86)
SYNOPSIS	STATUS vxPowerModeSet ( UINT32 mode /* power management mode to select */ )
DESCRIPTION	This routine selects the power management mode to be activated when <b>vxPowerDown()</b> is called. <b>vxPowerModeSet()</b> is normally called in the BSP initialization routine <b>sysHwInit()</b> .
USAGE PPC	Power management modes include the following:
	VX_POWER_MODE_DISABLE (0x1) Power management is disabled; this prevents the MSR(POW) bit from being set (all PPC).
	VX_POWER_MODE_FULL (0x2) All CPU units are active while the kernel is idle (PPC603, PPCEC603 and PPC860 only).
	VX_POWER_MODE_DOZE (0x4) Only the decrementer, data cache, and bus snooping are active while the kernel is idle (PPC603, PPCEC603 and PPC860).
	VX_POWER_MODE_NAP (0x8) Only the decrementer is active while the kernel is idle (PPC603, PPCEC603 and PPC604).
	VX_POWER_MODE_SLEEP (0x10) All CPU units are inactive while the kernel is idle (PPC603, PPCEC603 and PPC860 - not recommended for the PPC603 and PPCEC603 architecture).
	VX_POWER_MODE_DEEP_SLEEP (0x20) All CPU units are inactive while the kernel is idle (PPC860 only - not recommended).
	VX_POWER_MODE_DPM (0x40) Dynamic Power Management Mode (PPC603 and PPCEC603 only).

	VX_POWER_MODE_DOWN (0x80) Only a hard reset causes an exit from power-down low power mode (PPC860 only - not recommended).
USAGE SH	Power management modes include the following:
	<b>VX_POWER_MODE_DISABLE</b> (0x0) Power management is disabled.
	<b>VX_POWER_MODE_SLEEP</b> (0x1) The core CPU is halted, on-chip peripherals operating, external memory refreshing.
	VX_POWER_MODE_DEEP_SLEEP (0x2) The core CPU is halted, on-chip peripherals operating, external memory self-refreshing (SH-4 only).
	VX_POWER_MODE_USER (0xff) Set up to three 8-bit standby registers with user-specified values:
	vxPowerModeSet (VX_POWER_MODE_USER   sbr1<<8   sbr2<<16   sbr3<<24);
	The sbr1 value is written to the STBCR or SBYCR1, sbr2 is written to the STBCR2 or SBYCR2, and sbr3 is written to the STBCR3 register (when available), depending on the SH processor type.
USAGE X86	<b>vxPowerModeSet( )</b> is called in the BSP initialization routine <b>sysHwInit( )</b> . Power management modes include the following:
	<b>VX_POWER_MODE_DISABLE</b> (0x1) Power management is disable: this prevents halting the CPU.
	<b>VX_POWER_MODE_AUTOHALT</b> (0x4) Power management is enable: this allows halting the CPU.
RETURNS	<b>OK</b> , or <b>ERROR</b> if <i>mode</i> is incorrect or not supported by the processor.
ERRNO	Not Available
SEE ALSO	vxLib, vxPowerModeGet( ), vxPowerDown( )

# vxSSDisable()

NAME	<b>vxSSDisable()</b> – disable the superscalar dispatch (MC68060)
------	---

SYNOPSIS void vxSSDisable (void)

DESCRIPTION	This function resets the ESS bit of the Processor Configuration Register (PCR) to disable the superscalar dispatch.
RETURNS	N/A
ERRNO	Not Available

SEE ALSO vxLib

## vxSSEnable()

SYNOPSIS void vxSSEnable (void)

**DESCRIPTION** This function sets the ESS bit of the Processor Configuration Register (PCR) to enable the superscalar dispatch.

**vxSSEnable()** – enable the superscalar dispatch (MC68060)

RETURNS N/A

NAME

ERRNO Not Available

SEE ALSO vxLib

## vxTas()

NAME	vxTas() – C-ca	allable atomic test-and-set primitive
SYNOPSIS	BOOL vxTas ( void * a )	ddress /* address to test and set */
DESCRIPTION	This routine painstruction is e	rovides a C-callable interface to a test-and-set instruction. The test-and-set executed on the specified address. The architecture test-and-set instruction is:
	68K x86 SH ARM	tas lock bts tas.b swpb

VxWorks Kernel API Reference, 6.6 vxTssGet()

	This routine is equivalent to <b>sysBusTas( )</b> in <b>sysLib</b> .
MIPS	Because VxWorks does not support the MIPS MMU, only kseg0 and kseg1 addresses are accepted; other addresses return FALSE.
NOTE X86	BTS "Bit Test and Set" instruction is executed with LOCK instruction prefix to lock the Bus during the execution. The bit position 0 is toggled.
NOTE SH	The SH version of <b>vxTas()</b> simply executes the <b>tas.b</b> instruction, and the test-and-set (atomic read-modify-write) operation may require an external bus locking mechanism on some hardware. In this case, wrap the <b>vxTas()</b> with a bus locking and unlocking code in the <b>sysBusTas()</b> .
RETURNS	TRUE if the value had not been set (but is now), or FALSE if the value was set already.
ERRNO	Not Available
SEE ALSO	vxLib, sysBusTas()

# vxTssGet()

	wlib
ERRNO	Not Available
RETURNS	a value of the TASK register
DESCRIPTION	This routine gets a content of the TASK register
SYNOPSIS	int vxTssGet (void)
NAME	<b>vxTssGet()</b> – get a content of the TASK register (x86)

# vxTssSet()

**NAME vxTssSet()** – set a value to the TASK register (x86)

SYNOPSIS void vxTssSet

```
(
int value /* TASK register value */
)
```

**DESCRIPTION** This routine sets a value to the TASK register

RETURNS N/A

ERRNO Not Available

SEE ALSO vxLib

### vxbFileNvRamGet()

NAME	<b>vxbFileNvRamGet( )</b> – get the contents of non-volatile RAM
SYNOPSIS	<pre>STATUS vxbFileNvRamGet   (     char * fileName, /* name of NVRam file */     char *string, /* where to copy non-volatile RAM */     int strLen, /* maximum number of bytes to copy */     int offset /* byte offset into non-volatile RAM */   )</pre>
DESCRIPTION	This routine copies the contents of non-volatile memory into a specified string. The string is terminated with an EOS.
RETURNS	<b>OK</b> , or <b>ERROR</b> if parameters are invalid, the file cannot be opened, or cannot read from the file
ERRNO	Not Available
SEE ALSO	vxbFileNvRam, vxbFileNvRamSet( )

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VxWorks Kernel API Reference, 6.6 vxbFileNvRamRegister()

# vxbFileNvRamRegister()

NAME	<b>vxbFileNvRamRegister()</b> – register vxbFileNvRam driver
SYNOPSIS	<pre>void vxbFileNvRamRegister(void)</pre>
DESCRIPTION	This routine registers the vxbFileNvRam driver and device recognition data with the vxBus subsystem.
RETURNS	none
ERRNO	Not Available
SEE ALSO	vxbFileNvRam

# vxbFileNvRamSet()

NAME	<b>vxbFileNvRamSet( )</b> – write to non-volatile RAM
SYNOPSIS	<pre>STATUS vxbFileNvRamSet   (     char *fileName, /* name of NVRam file */     char *string, /* string to be copied into non-volatile RAM */     int strLen, /* maximum number of bytes to copy */     int offset /* byte offset into non-volatile RAM */     )</pre>
DESCRIPTION	This routine copies a specified string into non-volatile RAM.
RETURNS	<b>OK</b> , or <b>ERROR</b> if parameters are invalid, cannot open the nvram file, or cannot write to the nvram file.
ERRNO	Not Available
SEE ALSO	vxbFileNvRam, vxbFileNvRamGet( )

## vxbFileNvRampDrvCtrlShow()

NAME	<pre>vxbFileNvRampDrvCtrlShow() - show pDrvCtrl for template controller</pre>	
SYNOPSIS	<pre>int vxbFileNvRampDrvCtrlShow   (    VXB_DEVICE_ID pInst,    int verboseLevel   )</pre>	
DESCRIPTION	This routine prints information about the instance to to system console. This is not integrated with <b>vxBusShow</b> . RETURNS: 0, always	
RETURNS	Not Available	
ERRNO	Not Available	
SEE ALSO	vxbFileNvRam	

## vxbIntelIchStorageRegister()

- NAME vxbIntelIchStorageRegister() register driver with vxbus
- SYNOPSIS void vxbIntelIchStorageRegister (void)
- DESCRIPTION none
- RETURNS N/A
- ERRNO
- SEE ALSO vxbIntelIchStorage

	vxbNonVolGet()	
NAME	vxbNonVolGet() – get the contents of non-volatile RAM	
SYNOPSIS	<pre>STATUS vxbNonVolGet   (     char * drvName, /* requestor's name */     int drvUnit, /* requestor's unit number */     char * buff, /* buffer to copy non-volatile RAM into */     int offset, /* offset from start of allocation unit */     int strLen /* maximum number of bytes to copy */    )</pre>	
DESCRIPTION	<ul> <li>This routine reads information from a non-volatile memory device and stores it in the caller-provided buffer.</li> <li>The caller identifies itself by name and unit number. This routine is typically called by a device driver. In this case, the specified name is the name of the device driver, and the unit number is the unit number of the device/driver instance. However, other modules which use NVRam may make use of this routine as well, such as BOOTLINE. In this case, unit number should be set to zero or a number specific to the module.</li> <li>The amount of data copied is the size specified. If the specified size is greater than the size of the NVram segment allocated to the device, the behavior is undefined.</li> </ul>	
RETURNS	Not Available	
ERRNO	Not Available	
SEE ALSO	vxbNonVolLib, vxbNonVolSet( )	

# vxbNonVolLibInit()

- NAME vxbNonVolLibInit() Non Volatile RAM library initialization
- SYNOPSIS void vxbNonVolLibInit(void)

**DESCRIPTION** none

- **RETURNS** Not Available
- ERRNO Not Available

SEE ALSO vxbNonVolLib

## vxbNonVolSet()

NAME	<b>vxbNonVolSet( )</b> – write to non-volatile memory
SYNOPSIS	<pre>STATUS vxbNonVolSet   (     char * drvName, /* requestor's name */     int drvUnit, /* requestor's unit number */     char * buff, /* buffer to copy from, into non-volatile RAM */     int offset, /* offset from start of allocation unit */     int strLen /* maximum number of bytes to copy */    )</pre>
DESCRIPTION	This routine reads information from the caller-provided buffer and stores it in a non-volatile memory device.
	The caller identifies itself by name and unit number. This routine is typically called by a device driver. In this case, the specified name is the name of the device driver, and the unit number is the unit number of the device/driver instance. However, other modules which use NVRam may make use of this routine as well, such as BOOTLINE. In this case, unit number should be set to zero or a number specific to the module.
	The amount of data copied is the size specified. If the specified size is greater than the size of the NVram segment allocated to the device, the behavior is undefined.
RETURNS	Not Available
ERRNO	Not Available
SEE ALSO	vxbNonVolLib, vxbNonVolGet( )

# vxbSI31xxStorageRegister()

- NAME vxbSI31xxStorageRegister() register driver with vxbus
- SYNOPSIS void vxbSI31xxStorageRegister (void)
- **DESCRIPTION** none

VxWorks Kernel API Reference, 6.6 vxsimHostCpuVarsInit()

RETURNS N/A

ERRNO

SEE ALSO vxbSI31xxStorage

# vxsimHostCpuVarsInit()

NAME	vxsimHostCpuVarsInit() – intialize per cpu variable pointers
SYNOPSIS	void vxsimHostCpuVarsInit (void)
DESCRIPTION	This routine initializes VxWorks pointer to addresses in host binary in orde address is the same on every cpu but the value can be different.
RETURNS	N/A
ERRNO	Not Available

SEE ALSO vxsimHostArchLib

# vxsimHostDllLoad()

NAME	<b>vxsimHostDllLoad()</b> – load the given Dll to VxSim.	
SYNOPSIS	STATUS vxsimHostDllLoad ( char * dllName /* name of the DLL to load */ )	
DESCRIPTION	This routine loads the given Dll to VxSim. The Dll is loaded from current directory, using an absolute path name, from the System Dll search path, or from VxSim Dll path <i>WIND_BASE</i> /host/ <i>HOST_TYPE</i> /lib/vxsim.	
SMP CONSIDERATIONS		
RETURNS	OK. or ERROR if load failed.	

ERRNO N/A

SEE ALSO vxsimHostArchLib

## vxsimHostMmuCurrentSet()

NAME	vxsimHostMmuCurrentSet() – set current translation table mapping
SYNOPSIS	STATUS vxsimHostMmuCurrentSet ( MMU_TRANS_TBL * pTransTbl /* translation table to set */ )
DESCRIPTION	Set mapping corresponding to specified translation table. This routine only affects the current CPU. Vxsim host binary makes sure the translation table can not be updated while beeing read.
RETURNS	OK always
ERRNO	Not Available
SEE ALSO	vxsimHostArchLib

# vxsimHostMmuProtect()

NAME	vxsimHostMmuProtect() – set/clear protection on mmu pages
SYNOPSIS	<pre>STATUS vxsimHostMmuProtect   (    MMU_TRANS_TBL * pTransTbl, /* translation table */    VIRT_ADDR   addr,</pre>
DESCRIPTION	This routine sets or clear protection flags on mmu pages. It does not updtae corresponding PTE.
RETURNS	OK or ERROR

VxWorks Kernel API Reference, 6.6 vxsimHostProcAddrGet()

ERRNO Not Available

SEE ALSO vxsimHostArchLib

#### vxsimHostProcAddrGet()

```
    NAME
    vxsimHostProcAddrGet() - return the address of a host API

    SYNOPSIS
    FUNCPTR vxsimHostProcAddrGet

    (
    char * routineName /* host API name */

    )
    The effect of the stable of
```

**DESCRIPTION** This routine returns the address of a host API which name is given in parameter. No error message are displayed if the host API is not found.

#### SMP CONSIDERATIONS

routineName parameter must be in kernel memory.

**RETURNS** The address of the routine, or **NULL** if not found.

ERRNO N/A

SEE ALSO vxsimHostArchLib

#### vxsimHostProcCall()

```
NAME vxsimHostProcCall() - call a host routine

SYNOPSIS UINT32 vxsimHostProcCall

(

FUNCPTR rtnAddr, /* routine to be called */

UINT32 arg0, /* routine arguments */

UINT32 arg1,

UINT32 arg2,

UINT32 arg3,

UINT32 arg5,

UINT32 arg6,

UINT32 arg7,
```

	UINT32 arg8 )
DESCRIPTION	This routine calls a host routine whose address was previously retrieved through <b>vxsimHostProcAddrGet()</b> . On a SMP system it is the only way to use safely the value returned by <b>vxsimHostProcAddrGet()</b> as this value is only guaranteed to be correct on CPUi that performed <b>vxsimHostProcAddrGet()</b> .
RETURNS	routine return value as an UINT32
ERRNO	Not Available
SEE ALSO	vxsimHostArchLib

### vxsimHostSioBaudRateSet()

NAME	vxsimHostSioBaudRateSet() – set SIO device transfert rate

- SYNOPSIS STATUS vxsimHostSioBaudRateSet ( SIO\_ID sioId, int baudRate )
- **DESCRIPTION** This routine sets SIO device transfert rate

RETURNS OK or ERROR

ERRNO N/A

SEE ALSO vxsimHostArchLib

## vxsimHostSioClose()

NAME	vxsimHostSioClose() – close SIO device
SYNOPSIS	STATUS vxsimHostSioClose
	SIO_ID sioId /* sio descriptor */

VxWorks Kernel API Reference, 6.6 vxsimHostSioIntVecGet()

- **DESCRIPTION** This routine closes specified SIO device.
- **RETURNS** number of bytes read

ERRNO N/A

SEE ALSO vxsimHostArchLib

#### vxsimHostSioIntVecGet()

 NAME
 vxsimHostSioIntVecGet() – get SIO device interrupt vector

 SYNOPSIS
 int vxsimHostSioIntVecGet

 (SIO\_ID sioId /\* sio descriptor \*/)
 interrupt vector associated with a SIO device.

 RETURNS
 interrupt vector

 ERRNO
 N/A

 SEE ALSO
 vxsimHostArchLib

#### vxsimHostSioModeSet()

NAME	<pre>vxsimHostSioModeSet( ) - set SIO device mode (poll/interrupt)</pre>
SYNOPSIS	<pre>STATUS vxsimHostSioModeSet   (    SIO_ID sioId, /* sio descriptor */    int mode /* 0 = poll, 1 = interrupt */   )</pre>
DESCRIPTION	This routine sets SIO device interrupt mode.
RETURNS	OK or ERROR
ERRNO	N/A

SEE ALSO vxsimHostArchLib

# vxsimHostSioOpen()

NAME	<pre>vxsimHostSioOpen() – open SIO device</pre>
SYNOPSIS	STATUS vxsimHostSioOpen ( SIO_ID sioId /* sio descriptor */ )
DESCRIPTION	This routine opens specified SIO device.
RETURNS	OK or ERROR
ERRNO	N/A
SEE ALSO	vxsimHostArchLib

## vxsimHostSioRead()

NAME	vxsimHostSioRead() – read SIO device into buffer
SYNOPSIS	<pre>int vxsimHostSioRead   (    SIO_ID sioId,    char * buf,    int len   )</pre>
DESCRIPTION	This routine reads specified SIO device.
RETURNS	number of bytes read
ERRNO	N/A
SEE ALSO	vxsimHostArchLib

VxWorks Kernel API Reference, 6.6 vxsimHostSioWrite()

# vxsimHostSioWrite()

NAME	<b>vxsimHostSioWrite( )</b> – write buffer to SIO device
SYNOPSIS	<pre>int vxsimHostSioWrite   (    SIO_ID sioId,    char * buf,    int len   )</pre>
DESCRIPTION	This routine writes specified buffer to specified SIO device.
RETURNS	number of bytes written
ERRNO	N/A
SEE ALSO	vxsimHostArchLib

# w()

NAME	$\mathbf{w}(0 - \mathbf{print} \ \mathbf{a} \ \mathbf{summary} \ \mathbf{of} \ \mathbf{each} \ \mathbf{task's} \ \mathbf{pending} \ \mathbf{information}, \ \mathbf{task} \ \mathbf{by} \ \mathbf{task}$
SYNOPSIS	<pre>void w   (    int taskNameOrId /* task name or task ID */ )</pre>
DESCRIPTION	This routine shows a summary of each task's pending information, if <i>taskNameOrId</i> is equal to 0. Otherwise, it shows a summary for the specified task.
	This routine doesn't support POSIX semaphores and message queues. This command doesn't support pending signals.
	List of object types that are recognized:
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	usrLib, tw(), the VxWorks programmer guides.

NAME	<b>wdCancel()</b> – cancel a currently counting watchdog
SYNOPSIS	STATUS wdCancel ( WDOG_ID wdId /* ID of watchdog to cancel */ )
DESCRIPTION	This routine cancels a currently running watchdog timer by zeroing its delay count. Watchdog timers may be canceled from interrupt level.
SMP CONSIDERATIONS This API is spinlock and intCpuLock restricted.	
RETURNS	OK, or ERROR if the watchdog timer cannot be canceled.
ERRNO	Not Available
SEE ALSO	wdLib, wdStart( )

# wdCreate()

SEE ALSO	wdLib, wdDelete()
ERRNO	Not Available
RETURNS	The ID for the watchdog created, or NULL if memory is insufficient.
SMP CONSIDERATIO	NS This API is spinlock and intCpuLock restricted.
DESCRIPTION	This fourne creates a watchdog timer by anotating a wbod structure in memory.
DESCRIPTION	This routine creates a watchdog timer by allocating a WDOC structure in memory
SYNOPSIS	WDOG_ID wdCreate (void)
NAME	wdCreate() – create a watchdog timer

# wdDelete()

NAME	wdDelete() – delete a watchdog timer
SYNOPSIS	STATUS wdDelete ( WDOG_ID wdId /* ID of watchdog to delete */ )
DESCRIPTION	This routine de-allocates a watchdog timer. The watchdog will be removed from the timer queue if it has been started. This routine complements <b>wdCreate()</b> .
SMP CONSIDERATIO	<b>NS</b> This API is spinlock and intCpuLock restricted.
RETURNS	OK, or ERROR if the watchdog timer cannot be de-allocated.
ERRNO	Not Available
SEE ALSO	wdLib, wdCreate()

# wdInitialize()

NAME	<b>wdInitialize()</b> – initialize a pre-allocated watchdog.
SYNOPSIS	WDOG_ID wdInitialize ( char * pWdMem )
DESCRIPTION	This routine initializes a watchdog that has been pre-allocated (i.e. by the VX_WDOG macro).
	The following example illustrates use of the <b>VX_WDOG</b> macro and this function together to instantiate a watchdog statically (without using any dynamic memory allocation):
	<pre>#include <vxworks.h> #include <wdlib.h></wdlib.h></vxworks.h></pre>
	VX_WDOG(myWdog); /* declare the watchdog */ WDOG_ID myWdogId; /* watchdog ID for further operations */
	STATUS initializeFunction (void) { if ((myWdogId = wdInitialize (myWdog)) == NULL)

```
return (ERROR); /* initialization failed */
else
    return (OK);
}
```

#### SMP CONSIDERATIONS

This API is spinlock and intCpuLock restricted.

RETURNS	The watchdog ID, or <b>NULL</b> on error.
ERRNO	N/A
SEE ALSO	wdLib

# wdShow()

NAME	wdShow() – show information about a watchdog
SYNOPSIS	STATUS wdShow ( WDOG_ID wdId /* watchdog to display */ )
DESCRIPTION	This routine displays the state of a watchdog.
EXAMPLE	A summary of the state of a watchdog is displayed as follows:
	-> wdShow myWdId Watchdog Id : 0x3dd46c State : OUT_OF_Q Ticks Remaining : 0 Routine : 0 Parameter : 0
SMP CONSIDERATIO	INS
	This API is spinlock and intCpuLock restricted.
RETURNS	OK or ERROR.
ERRNO	Not Available
SEE ALSO	<b>wdShow</b> , <b>windsh</b> , the VxWorks programmer guides, the, <i>VxWorks Command-Line Tools User's Guide</i> .

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### wdShowInit()

NAME	wdShowInit() – initialize the watchdog show facility
SYNOPSIS	void wdShowInit (void)
DESCRIPTION	This routine links the watchdog show facility into the VxWorks system. It is called automatically when the watchdog show facility is configured into VxWorks using either of the following methods:
	- If you use the configuration header files, define INCLUDE_SHOW_ROUTINES in <b>config.h</b> .
	- If you use the Tornado project facility, select INCLUDE_WATCHDOGS_SHOW.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	wdShow

### wdStart()

NAME	wdStart() – start a watchdog timer
SYNOPSIS	<pre>STATUS wdStart   (   WDOG_ID wdId, /* watchdog ID */   int delay, /* delay count, in ticks */   FUNCPTR pRoutine, /* routine to call on time-out */   int parameter /* parameter with which to call routine */   )</pre>

**DESCRIPTION** This routine adds a watchdog timer to the system tick queue. The specified watchdog routine will be called from interrupt level after the specified number of ticks has elapsed. Watchdog timers may be started from interrupt level.

To replace either the timeout *delay* or the routine to be executed, call **wdStart()** again with the same *wdId*; only the most recent **wdStart()** on a given watchdog ID has any effect. (If your application requires multiple watchdog routines, use **wdCreate()** to generate separate a watchdog ID for each.) To cancel a watchdog timer before the specified tick count is reached, call **wdCancel()**.

	Watchdog timers execute only once, but some applications require periodically executing timers. To achieve this effect, the timer routine itself must call <b>wdStart()</b> to restart the timer on each invocation.	
WARNING	The watchdog routine runs in the context of the system-clock ISR; thus, it is subject to all ISR restrictions.	
NOTE	watchdog routine invocation can be deferred. As such isrIdCurrent is either a valid ISR_ID or is NULL in the case of deferral	
SMP CONSIDERATIONS		
	This API is spinlock and intCpuLock restricted.	
RETURNS	OK, or ERROR if the watchdog timer cannot be started.	
ERRNO	Not Available	
SEE ALSO	wdLib, wdCancel()	

# wdbMdlSymSyncLibInit()

- NAME wdbMdlSymSyncLibInit() initialize modules and symbols synchronization library
- SYNOPSIS void wdbMdlSymSyncLibInit (void)
- **DESCRIPTION** This routine initializes the stuff needed by the modules and symbols synchronization.
- RETURNS N/A
- ERRNO N/A
- SEE ALSO wdbMdlSymSyncLib

# wdbSystemSuspend()

NAME wdbSystemSuspend() – suspend the system

SYNOPSIS STATUS wdbSystemSuspend (void)

VxWorks Kernel API Reference, 6.6 wdbUserEvtLibInit()

DESCRIPTION	This routine transfers control from the run time system to the WDB agent running in external mode. In order to give back the control to the system it must be resumed by the external WDB agent.
EXAMPLE	The code below, called in a vxWorks application, suspends the system :
	<pre>if (wdbSystemSuspend () != OK)     printf ("External mode is not supported by the WDB agent.\n");</pre>
	From a host tool, we can detect that the system is suspended.
	First, attach to the target server :
	wtxtcl> wtxToolAttach EP960CX EP960CX_ps@sevre
	Then, you can get the agent mode :
	wtxtcl> wtxAgentModeGet AGENT_MODE_EXTERN
	To get the status of the system context, execute :
	wtxtcl> wtxContextStatusGet CONTEXT_SYSTEM 0 CONTEXT_SUSPENDED
	In order to resume the system, simply execute :
	wtxtcl> wtxContextResume CONTEXT_SYSTEM 0 0
	You will see that the system is now running :
	wtxtcl> wtxContextStatusGet CONTEXT_SYSTEM 0 CONTEXT_RUNNING
RETURNS	<b>OK</b> upon successful completion, <b>ERROR</b> if external mode is not supported by the WDB agent.
ERRNO	Not Available
SEE ALSO	wdbLib

# wdbUserEvtLibInit()

NAME wdbUserEvtLibInit() – include the WDB user event library

SYNOPSIS void wdbUserEvtLibInit (void)

This null routine is provided so that wdbUserEvtLib can be linked into the system. If the

component INCLUDE\_WDB\_USER\_EVENT is included at configuration time,

wdbUserEvtLibInit() is automatically called.

ERRNO Not Available

DESCRIPTION

SEE ALSO wdbUserEvtLib

# wdbUserEvtPost()

NAME	wdbUserEvtPost() – post a user event string to host tools
SYNOPSIS	STATUS wdbUserEvtPost ( char * event /* event string to send */ )
DESCRIPTION	This routine posts the string <i>event</i> to host tools that have registered for it. Host tools will receive a USER WTX event string. The maximum size of the event is <b>WDB_MAX_USER_EVT_SIZE</b> (defined in <b>\$WIND_BASE/target/h/wdb/wdbLib.h</b> ).
EXAMPLE	The code below sends a WDB user event to host tools :
	<pre>char * message = "Alarm: reactor overheating !!!";</pre>
	if (wdbUserEvtPost (message) != OK) printf ("Can't send alarm message to host tools");
	This event will be received by host tools that have registered for it. For example a WTX TCL based tool would do :
	wtxtcl> wtxToolAttach EP960CX EP960CX_ps@sevre wtxtcl> wtxRegisterForEvent "USER.*" 0 wtxtcl> wtxEventGet USER Alarm: reactor overheating !!!
	Host tools can register for more specific user events :
	wtxtcl> wtxToolAttach EP960CX EP960CX_ps@sevre wtxtcl> wtxRegisterForEvent "USER Alarm.*" 0 wtxtcl> wtxEventGet USER Alarm: reactor overheating !!!

VxWorks Kernel API Reference, 6.6 wim()

 In this piece of code, only the USER events beginning with "Alarm" will be received.

 RETURNS
 OK upon successful completion or ERROR if unable to send the event to the host or if the size of the event is greater than WDB\_MAX\_USER\_EVT\_SIZE.

 ERRNO
 Not Available

 SEE ALSO
 wdbUserEvtLib

### wim()

NAME	<b>wim()</b> – return the contents of the window invalid mask register (SimSolaris)
SYNOPSIS	<pre>int wim   (    int taskId /* task ID, 0 means default task */ )</pre>
DESCRIPTION	This command extracts the contents of the window invalid mask register from the TCB of a specified task. If <i>taskId</i> is omitted or 0, the default task is assumed.
RETURNS	The contents of the window invalid mask register.
ERRNO	Not Available
SEE ALSO	dbgArchLib, VxWorks Programmer's Guide: Debugging

## windPwrDownRtnSet()

NAME	<pre>windPwrDownRtnSet() - register a BSP power-down function</pre>
SYNOPSIS	void windPwrDownRtnSet ( WIND_PWR_DOWN_RTN dRtn /* power down function pointer from BSP */ )
WARNING	This routine is deprecated. Calling this function is a nop.
	This routine registers a BSP power-down function with WIND CPU power management. The function registered will be called when the WIND kernel decides that the CPU can be

powered off. Note that the power-down function will not be invoked while the CPU power mode is set to windPwrModeOff.

The power-down function is passed two parameters: a WIND\_PWR\_MODE power mode and a ULONG nTicks.

The power mode parameter is the current kernel power mode that is in effect (as set by **windPwrModeSet()**) when the WIND kernel goes idle. nTicks is the maximum number of ticks that the WIND kernel is willing to sleep for before it must wake up and perform some work, such as scheduling a task. nTicks may be passed as **WAIT\_FOREVER** (0L), which indicates that the kernel has requirements as to when it is woken up next.

The routine registered is invoked with interrupts locked and is not allowed to make any WIND kernel calls either directly or indirectly. If these must be made, the only option is for the power-down function to perform a windPwrModeSet (windPwrModeOff) and arrange for an interrupt (a software or hardware interrupt) to occur to make the WIND kernel calls on its behalf. Such an interrupt will occur after the WIND kernel unlocks interrupts which it does after invoking the registered power-down routine, aborting its subsequent call to **vxArchPowerDown()** to run the interrupt and process any kernel work the interrupt makes.

#### SMP CONSIDERATIONS

In an SMP environment it is possible for some CPUs to be idle and others to be executing tasks or ISRs. The registered power-down routine is called when a CPU goes idle regardless of the state of other CPUs in the SMP system. The routine is executed by the CPU that is going idle. The routine must ensure it does not perform any power-down actions that would disrupt execution on non-idle CPUs.

RETURNS	N/A

ERRNO N/A

SEE ALSO windPwrLib, kernelIsCpuIdle(), kernelIsSystemIdle()

#### windPwrModeGet( )

NAME	<pre>windPwrModeGet( ) - Get the current power mode</pre>
SYNOPSIS	WIND_PWR_MODE windPwrModeGet(void)
DESCRIPTION	This routine is called whenever the BSP needs the current power mode
RETURNS	WIND_PWR_MODE

VxWorks Kernel API Reference, 6.6 windPwrModeSet()

errno N/A

SEE ALSO windPwrLib, windPwrModeSet()

# windPwrModeSet()

NAME	<pre>windPwrModeSet() - Set the BSP power mode</pre>
SYNOPSIS	<pre>void windPwrModeSet   (    WIND_PWR_MODE mode /* new power mode */   )</pre>
DESCRIPTION	This routine is called during initialization and whenever the power mode is set by the BSP.
RETURNS	N/A
ERRNO	N/A
SEE ALSO	windPwrLib, windPwrModeGet()

# windPwrUpRtnSet()

NAME	<pre>windPwrUpRtnSet() - register a BSP power-up function</pre>
SYNOPSIS	void windPwrUpRtnSet ( WIND_PWR_UP_RTN uRtn )
WARNING	This routine is deprecated. Calling this function is a nop. This routine registers a BSP power-up function with WIND CPU power management. The function registered will be called whenever an interrupt exception occurs to wake up the CPU while it was powered off (or it was in the process of powering off).
	The power-up function is passed two parameters: a <b>WIND_PWR_MODE</b> power mode and a pointer to a ULONG nTicks.

SEE ALSO	windPwrLib, kernelIsCpuIdle(), kernelIsSystemIdle()
ERRNO	N/A
RETURNS	N/A
SMP CONSIDERATIO	In an SMP environment it is possible for some CPUs to be idle and others to be executing tasks or ISRs. The registered power-up routine is called when a CPU is awakened regardless of the state of other CPUs in the SMP system. The routine is executed by the CPU that is awakened. The routine must ensure it does not perform any power-up actions that would disrupt execution on other CPUs.
	The routine registered is invoked with interrupts locked and is not allowed to make any WIND kernel calls either directly or indirectly. If these must be made, the only option is for the power-up function to arrange for an interrupt (a software or hardware interrupt) to occur to make the WIND kernel calls on its behalf. Such an interrupt will occur after the WIND kernel unlocks interrupts which it does after invoking the registered power-up routine. If the arranged interrupt is at a higher priority than the interrupt that is waking up the CPU, it will execute immediately, otherwise it will run afterwards.
	The power mode parameter is the current WIND kernel power mode that is in effect (as set by <b>windPwrModeSet( )</b> ) when the kernel wakes up. The pointer to nTicks is to be set by the power-up function to inform the kernel how long it has slept, in ticks.

# write()

NAME	write() – write bytes to a file
SYNOPSIS	<pre>int write   (    int fd, /* file descriptor on which to write */    char *buffer, /* buffer containing bytes to be written */    size_t nbytes /* number of bytes to write */   )</pre>
DESCRIPTION	This routine writes <i>nbytes</i> bytes from <i>buffer</i> to a specified file descriptor <i>fd</i> . It calls the device driver to do the work.
RETURNS	The number of bytes written (if not equal to <i>nbytes</i> , an error has occurred), or <b>ERROR</b> if the file descriptor does not exist, the driver does not have a write routine, or the driver returns <b>ERROR</b> . If the driver does not have a write routine, errno is set to <b>ENOTSUP</b> .

VxWorks Kernel API Reference, 6.6 wvAllObjsSet( )

ERRNO	EBADF Bad file descriptor number.
	ENOTSUP Device driver does not support the write command.
	ENXIO Device and its driver are removed. <b>close( )</b> should be called to release this file descriptor.
	Other Other errors reported by device driver.
SEE ALSO	ioLib

# wvAllObjsSet()

NAME	<pre>wvAllObjsSet() – set instrumented state for all objects and classes</pre>
SYNOPSIS	<pre>void wvAllObjsSet    (    int mode /* INSTRUMENT_ON or INSTRUMENT_OFF */ )</pre>
DESCRIPTION	This routine enables or disables instrumentation for all object classes and instances in the system.
	If <i>mode</i> is <b>INSTRUMENT_ON</b> , instrumentation is turned on; if it is <b>INSTRUMENT_OFF</b> , instrumentation is turned off. Any other value has no effect
	This routine has effect only if <b>INCLUDE_WINDVIEW</b> is defined in <b>configAll.h</b> and event logging has been enabled for system objects.
RETURNS	N/A
ERRNO	
SEE ALSO	wvLib, wvSigInst( ), wvEventInst( ), wvObjInstAllClear( ), wvObjInst( )

#### wvCurrentLogGet( )

 NAME
 wvCurrentLogGet() - return a pointer to the currently active System Viewer log

 SYNOPSIS
 WV\_LOG \* wvCurrentLogGet (void)

 DESCRIPTION
 This routine returns a pointer to the currently active System Viewer log.

 RETURNS
 Pointer to the log, or NULL

 ERRNO
 SEE ALSO

 wvLib, wvCurrentLogSet()

#### wvCurrentLogListGet()

NAME wvCurrentLogListGet() – return a pointer to the System Viewer log list

- SYNOPSIS WV\_LOG\_LIST \* wvCurrentLogListGet (void)
- **DESCRIPTION** This routine returns a pointer to the System Viewer log list. It is not expected that there would be more than one list at any one time.
- **RETURNS** Pointer to the log list, or **NULL**

ERRNO

SEE ALSO wvLib, wvCurrentLogListSet()

#### wvCurrentLogListSet( )

NAME wvCurrentLogListSet() - set the current log list
SYNOPSIS void wvCurrentLogListSet
(
WV\_LOG\_LIST \* pWvLogList
)

VxWorks Kernel API Reference, 6.6 wvCurrentLogSet()

DESCRIPTION	This routine selects a System Viewer log list for subsequent operations. It could be used after reboot, to choose a log in a persistent memory area.
RETURNS	N/A
ERRNO	
SEE ALSO	wvLib, wvCurrentLogListGet()

# wvCurrentLogSet()

NAME wvCurrentLogSet() – select a System Viewer log as currently active

SYNOPSIS	void wvCurrentLogSet (
	WV_LOG * pWvLog )
DESCRIPTION	This routine selects a System Viewer log as the active one.

RETURNS N/A

ERRNO

SEE ALSO wvLib, wvCurrentLogGet()

## wvEdrInst()

wvEdrInst() – instrument ED&R Events
STATUS wvEdrInst ( WV_INSTRUMENTATION_MODE mode /* instrumentation mode */ )
This routine instruments ED&R Event activity.
If <i>mode</i> is <b>INSTRUMENT_ON</b> , instrumentation for ED&R events is turned on; if <i>mode</i> is <b>INSTRUMENT_OFF</b> , instrumentation for ED&R Events is turned off. Any other value causes the current instrumentation state to be returned.

RETURNS OK or ERROR.

ERRNO

SEE ALSO wvLib

#### wvEvent()

NAME wvEvent() – log a user-defined event

- SYNOPSIS STATUS wvEvent
  - ( event\_t usrEventId, /\* event \*/ char \* buffer, /\* buffer \*/ size\_t bufSize /\* buffer size \*/ )
- **DESCRIPTION** This routine logs a user event. Event logging must have been started with **wvEvtLogStart()** or from the System Viewer GUI to use this routine. The *usrEventId* should be in the range 0-25535. A buffer of data can be associated with the event; *buffer* is a pointer to the start of the data block, and *bufSize* is its length in bytes.

#### SMP CONSIDERATIONS

This API is spinlock restricted.

**RETURNS** OK, or ERROR if the event can not be logged.

ERRNO

SEE ALSO wvLib, dbgLib, e( )

#### wvEventInst( )

NAME	<pre>wvEventInst() - instrument VxWorks Events</pre>
SYNOPSIS	<pre>int wvEventInst   (     WV_INSTRUMENTATION_MODE mode /* INSTRUMENT_ON, INSTRUMENT_OFF */ )</pre>
DESCRIPTION	This routine instruments VxWorks Event activity.

VxWorks Kernel API Reference, 6.6 wvEvtClassClear()

If *mode* is **INSTRUMENT\_ON**, instrumentation for VxWorks events is turned on; if it is any other value (including **INSTRUMENT\_OFF**), instrumentation for VxWorks Events is turned off.

This routine has effect only if **INCLUDE\_WINDVIEW** is defined in **configAll.h** and event logging has been enabled for system objects.

Parameters:

	mode The required instrumentation mode. The value INSTRUMENT_ON enables instrumentation for events, INSTRUMENT_OFF disables it, and any other value causes the current state to be returned.
RETURNS	The mode (INSTRUMENT_ON or INSTRUMENT_OFF) currently in force.
ERRNO	
SEE ALSO	wvLib

### wvEvtClassClear()

NAME	wvEvtClassClear() – clear the specified class of events from those being logged
SYNOPSIS	void wvEvtClassClear ( UINT32 classDescription /* description of evt classes to clear */ )
DESCRIPTION	This routine clears the class or classes described by <i>classDescription</i> from the set of classes currently being logged.
RETURNS	N/A
ERRNO	
SEE ALSO	wvLib

NAME	wvEvtClassClearAll() – clear all classes of events from those logged
SYNOPSIS	void wvEvtClassClearAll (void)
DESCRIPTION	This routine clears all classes of events so that no classes are logged if event logging is started.
RETURNS	N/A
ERRNO	
SEE ALSO	wvLib

# wvEvtClassGet()

NAME	<pre>wvEvtClassGet() - get the current set of classes being logged</pre>
SYNOPSIS	UINT32 wvEvtClassGet (void)
DESCRIPTION	This routine returns the set of classes currently being logged.
RETURNS	The class description.
ERRNO	
SEE ALSO	wvLib

# wvEvtClassSet()

NAME	<b>wvEvtClassSet()</b> – set the class of events to log
SYNOPSIS	void wvEvtClassSet (
	UINT32 classDescription /* description of evt classes to set */

VxWorks Kernel API Reference, 6.6 wvEvtLogStart()

 DESCRIPTION
 This routine sets the class of events which are logged when event logging is started.

 classDescription can take the following values:
 WV\_CLASS\_1 /\* Events causing context switches \*/

 WV\_CLASS\_2 /\* Events causing task-state transitions \*/
 WV\_CLASS\_3 /\* Events from object and system libraries \*/

 See wvLib for more information about these classes, particularly Class 3.

 RETURNS
 N/A

 ERRNO

 SEE ALSO
 wvLib, wvObjInst(), wvSigInst(), wvEventInst(), wvSalInst().

## wvEvtLogStart( )

NAME	<b>wvEvtLogStart()</b> – start logging events to the buffer
SYNOPSIS	STATUS wvEvtLogStart (void)
DESCRIPTION	This routine starts event logging. It also resets the timestamp mechanism so that it can be called more than once without stopping event logging.
RETURNS	OK, or ERROR if no buffer in use
ERRNO	
SEE ALSO	wvLib

# wvEvtLogStop()

NAME	<pre>wvEvtLogStop() - stop logging events to the buffer</pre>
SYNOPSIS	void wvEvtLogStop (void)
DESCRIPTION	This routine turns off all event logging, including event-logging of objects and signals specifically requested by the user. In addition, it disables the timestamp facility.
RETURNS	N/A

ERRNO

SEE ALSO wvLib

#### wvFileUploadPathCreate()

NAME wvFileUploadPathCreate() – create a file for depositing event data

SYNOPSIS	UPLOAD_ID wvFileUpl	oadPathCreate	
	(		
	char *fname,	/* name of file to create */	
	int openFlags	/* O_CREAT, O_TRUNC */	
	)		

- **DESCRIPTION** This routine opens and initializes a file to receive uploaded events. The *openFlags* argument is passed on as the flags argument to the actual open call so that the caller can specify things like **O\_TRUNC** and **O\_CREAT**. The file is always opened as **O\_WRONLY**, regardless of the value of *openFlags*.
- **RETURNS** The UPLOAD\_ID, or NULL if the file can not be opened or memory for the ID is not available.
- ERRNO Not Available
- SEE ALSO wvFileUploadPathLib, wvFileUploadPathClose()

#### wvFileUploadPathLibInit()

NAME	<pre>wvFileUploadPathLibInit() - initialize the wvFileUploadPathLib library</pre>
SYNOPSIS	STATUS wvFileUploadPathLibInit (void)
DESCRIPTION	This routine initializes the library by pulling in the routines in this file for use with WindView. It is called during system configuration from <b>usrWindview.c</b> .
RETURNS	OK.
ERRNO	Not Available
SEE ALSO	wvFileUploadPathLib

### wvFileUploadPathWrite()

wvFileUploadPathWrite() - write to the event-destination file NAME SYNOPSIS int wvFileUploadPathWrite ( UPLOAD\_ID pathId, /\* generic upload-path descriptor \*/ char \* pStart, /\* address of data to write \*/ size\_t size /\* number of bytes of data at pStart \*/ ) This routine writes *size* bytes of data beginning at *pStart* to the file indicated by *pathId*. DESCRIPTION The number of bytes written, or ERROR. RETURNS Not Available ERRNO wvFileUploadPathLib SEE ALSO

#### wvLibInit()

NAME	<pre>wvLibInit( ) - initialize wvLib - first step</pre>
SYNOPSIS	void wvLibInit (void)
DESCRIPTION	This routine starts initializing <b>wvLib</b> . Its actions should be performed before object creation, so it is called from <b>usrKernelInit()</b> in <b>usrKernel.c</b> .
RETURNS	N/A
ERRNO	
SEE ALSO	wvLib

#### wvLibInit2()

#### NAME wvLibInit2() – initialize wvLib - final step

SYNOPSIS void wvLibInit2 (void)
This routine is called after wvLibInit() to complete the initialization of wvLib. It should be

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RETURNS N/A

ERRNO

DESCRIPTION

SEE ALSO wvLib

## wvLogCountGet()

called before starting any event logging.

NAME wvLogCountGet() – return the number of logs in the curent log list

SYNOPSIS int wvLogCountGet (void)

- **DESCRIPTION** This routine returns the number of System Viewer logs in the current log list.
- **RETURNS** number of logs

ERRNO

SEE ALSO wvLib, wvLogFirstGet(), wvLogNextGet()

# wvLogCreate()

NAME	wvLogCreate() – Create a System Viewer log
SYNOPSIS	<pre>WV_LOG * wvLogCreate    (    BUFFER_ID evtBuffer /* event buffer to use */ )</pre>
DESCRIPTION	This routine creates a System Viewer log, and then inserts it into the System Viewer log list. If the routine encounters an error, then the insertion is not done, a node is not created, and the caller should delete the event buffer passed in.
RETURNS	Pointer to new log, or <b>NULL</b> if no buffer or log list supplied, or an error occured while allocating memory.

VxWorks Kernel API Reference, 6.6 wvLogDelete()

### ERRNO S\_smObjLib\_NOT\_INITIALIZED S\_memLib\_NOT\_ENOUGH\_MEMORY

SEE ALSO wvLib, wvLogDelete()

# wvLogDelete()

NAME	wvLogDelete() – Delete a System Viewer log
SYNOPSIS	STATUS wvLogDelete ( WV_LOG_LIST * pLogList, /* log list in which log appears */ WV_LOG * pWvLog /* System Viewer log to delete */ )
DESCRIPTION	This routine deletes a System Viewer log, and removes it from the log list. If the chosen log is the current one, and logging is enabled, the operation fails
RETURNS	OK, or ERROR if NULL log or log list supplied, or an error occured freeing memory
ERRNO	S_smObjLib_NOT_INITIALIZED S_memLib_NOT_ENOUGH_MEMORY
SEE ALSO	wvLib, wvLogCreate()

# wvLogFirstGet()

NAME	wvLogFirstGet() – return a pointer to the first log in the System Viewer log list
SYNOPSIS	WV_LOG * wvLogFirstGet (void)
DESCRIPTION	This routine returns a pointer to the first System Viewer log in the list of logs. If there is no log list, the function returns <b>NULL</b>
RETURNS	Pointer to the log, or NULL
ERRNO	
SEE ALSO	wvLib, wvLogNextGet( ), logCountGet( )

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NAME wvLogListCreate() – create a list to hold System Viewer logs

SYNOPSIS WV\_LOG\_LIST \* wvLogListCreate (void)

**DESCRIPTION** This routine creates a list to hold System Viewer logs. The list is created in the memory partition returned by **wvPartitionGet()**. If the partition has not been set, or the required memory could not be allocated, the routine returns **NULL**. If a log list has already been created, the function returns **NULL**, so the old list should be deleted before a new one is created. Otherwise, it returns a pointer to the newly-created list.

**RETURNS** pointer to the created WV\_LOG\_LIST, or NULL

- ERRNO S\_smObjLib\_NOT\_INITIALIZED S\_memLib\_NOT\_ENOUGH\_MEMORY
- SEE ALSO wvLib, wvLogListDelete()

### wvLogListDelete()

NAME	wvLogListDelete( ) – delete a System Viewer log list
SYNOPSIS	STATUS wvLogListDelete ( WV_LOG_LIST * pLogList )
DESCRIPTION	This function deletes a System Viewer log list, and all its contents. If logging is enabled, it returns <b>ERROR</b> .
RETURNS	Ok, or ERROR if no log list supplied, or an error occured while freeing memory
ERRNO	S_smObjLib_NOT_INITIALIZED S_memLib_NOT_ENOUGH_MEMORY
SEE ALSO	wvLib, wvLogListCreate()

## wvLogNextGet()

 

 NAME
 wvLogNextGet() - return a pointer to the next log in the System Viewer log list

 SYNOPSIS
 WV\_LOG \* wvLogNextGet ( WV\_LOG \* pWvLog )

 DESCRIPTION
 This routine returns a pointer to the next System Viewer log in the list of logs.

 RETURNS
 Pointer to the log, or NULL

 ERRNO
 see ALSO

 wvLib, wvLogFirstGet(), wvLogCountGet()

### wvObjInst()

NAME	<pre>wvObjInst() – instrument objects</pre>	
SYNOPSIS	<pre>int wvObjInst   (     enum windObjClassType ol     void * ol     WV_INSTRUMENTATION_MODE ma     )</pre>	ojType, /* object type */ ojId, /* obj ID or NULL for all objs */ ode /* instrumentation mode */
DESCRIPTION	This routine instruments a specifie objects have been enabled for ever WIND objects as listed in the <b>type</b> instrumented:	<pre>d object or set of objects and has effect when system t logging. objType could be set to any type of recognized s/vxWind.h enum list, provided they have been = 0 invalid class type class * Wind native semaphore POSIX semaphore * Wind native message queue POSIX message queue * realtime process * task * watchdog * file descriptor page pool</pre>
	10 windPgMgrClass 11 windGrpClass 12 windVmContextClass	page manager group virtual memory context

	14 windMemPartClass * memory partition 15 windI2oClass I20
	16 windDmsClassdevice management system17 windOmsClassobject management system (HA/FT)18 windSetClassset
	19 windIsrClass*ISR objects20 windTimerClasstimer services21 windSdClassShared data region
	In the list above, the instrumented objects are marked with an asterisk. <i>objId</i> specifies the identifier of the particular object to be instrumented. If <i>objId</i> is <b>NULL</b> , then all objects of <i>objType</i> have instrumentation turned on or off depending on the value of <i>mode</i> .
	If <i>mode</i> is <b>INSTRUMENT_ON</b> , instrumentation is turned on; if it is <b>INSTRUMENT_OFF</b> then instrumentation is turned off. Any other value has no effect, but the current mode ( <b>INSTRUMENTATION_ON</b> or <b>INSTRUMENTATION_OFF</b> ) is returned.
	Use <b>wvSigInst()</b> if you want to enable instrumentation for all signal activity, <b>wvEventInst()</b> for vxWorks Event activity, wvSalInst for all SAL call activities.
	This routine has effect only if the component <b>INCLUDE_WINDVIEW</b> is included in the project.
RETURNS	INSTRUMENT_ON, INSTRUMENT_OFF, or ERROR.
ERRNO	S_objLib_OBJ_ID_ERROR S_objLib_OBJ_ILLEGAL_CLASS_TYPE

trigger

wvLib, wvSigInst(), wvEventInst(), wvObjInstAllClear(), wvAllObjsSet() SEE ALSO

### wvObjInstModeSet( )

13 windTrgClass

NAME wvObjInstModeSet() - set object instrumentation on/off SYNOPSIS STATUS wvObjInstModeSet ( int mode /\* object instrumentation on/off \*/ ) DESCRIPTION This routine causes objects to be created either instrumented or not depending on the value of mode, which can be INSTRUMENT\_ON or INSTRUMENT\_OFF. All objects created after wvObjInstModeSet() is called with INSTRUMENT\_ON and before it is called with **INSTRUMENT\_OFF** are created as instrumented objects.

VxWorks Kernel API Reference, 6.6 wvPartitionGet()

	Use <b>wvObjInst()</b> if you want to enable instrumentation for a specific object or set of objects. Use <b>wvSigInst()</b> if you want to enable instrumentation for all signal activity, and <b>wvEventInst()</b> to enable instrumentation for VxWorks Event activity.
	This routine has effect only if INCLUDE_WINDVIEW is defined in configAll.h.
RETURNS	The previous value of <i>mode</i> or <b>ERROR</b> .
ERRNO	
SEE ALSO	wvLib, wvObjInst( ), wvSigInst( ), wvEventInst( ), wvEdrInst( )

### wvPartitionGet()

NAME	wvPartitionGet() – determine partition in use for System Viewer logging
SYNOPSIS	PART_ID wvPartitionGet (void)
DESCRIPTION	This routine returns the mamory partition id being used for System Viewer logs.
RETURNS	partition id
ERRNO	
SEE ALSO	wvLib, wvPartitionSet()

# wvPartitionSet()

NAME	wvPartitionSet() – specify a partition for use by System Viewer logging
SYNOPSIS	<pre>void wvPartitionSet     (     PART_ID memPart )</pre>
DESCRIPTION	This routine allows the user to specify a memory partition to be used for System Viewer logging. Subsequent calls to create System Viewer log lists will result in the logs being created in this partition.
	If using a post-mortem log, this routine should be called before a new log is created. Then the logs can be read, and a new partition and log list created, if required. Note that in

post-mortem mode, if the target has rebooted, then a log list in the preserved memory should not be deleted, because the memory partition was not created during this run of the target.

RETURNS n/a

ERRNO

SEE ALSO wvLib, wvPartitionGet()

## wvRBuffMgrPrioritySet()

NAME	<b>wvRBuffMgrPrioritySet()</b> – set the priority of the System Viewer rBuff manager
SYNOPSIS	STATUS wvRBuffMgrPrioritySet ( int priority /* new priority */ )
DESCRIPTION	This routine changes the priority of the <b>tWvRBuffMgr</b> task to the value of <i>priority</i> . Priorities range from 0, the highest priority, to 255, the lowest priority. If the task is not yet running, this priority is used when it is spawned.
RETURNS	<b>OK</b> , or <b>ERROR</b> if the priority can not be set.
ERRNO	Not Available
SEE ALSO	rBuffLib, taskPrioritySet(), the VxWorks programmer guides.

### wvSalInst()

NAME	wvSalInst() – instrument SAL
SYNOPSIS	<pre>int wvSalInst   (</pre>
DESCRIPTION	This routine instruments all SAL activity.

	If <i>mode</i> is <b>INSTRUMENT_ON</b> , instrumentation for SAL call is turned on; if it is <b>INSTRUMENT_OFF</b> , instrumentation for SAL call is turned off.
	This routine has effect only if <b>INCLUDE_WINDVIEW</b> is defined in <b>configAll.h</b> and event logging has been enabled for system objects.
	Parameters:
	<i>mode</i> The required instrumentation mode. The value <b>INSTRUMENT_ON</b> enables instrumentation for SAL, <b>INSTRUMENT_OFF</b> disables it, and any other value causes the current state to be returned.
RETURNS	The mode (INSTRUMENT_ON or INSTRUMENT_OFF) currently in force.
ERRNO	Not Available
SEE ALSO	wvLib

# wvSigInst()

NAME	<pre>wvSigInst() - instrument signals</pre>
SYNOPSIS	<pre>int wvSigInst   (</pre>
DESCRIPTION	This routine instruments all signal activity.
	If <i>mode</i> is <b>INSTRUMENT_ON</b> , instrumentation for signals is turned on; if it is <b>INSTRUMENT_OFF</b> , instrumentation for signals is turned off.
	This routine has effect only if <b>INCLUDE_WINDVIEW</b> is defined in <b>configAll.h</b> and event logging has been enabled for system objects.
	Parameters:
	<i>mode</i> The required instrumentation mode. The value <b>INSTRUMENT_ON</b> enables instrumentation for signals, <b>INSTRUMENT_OFF</b> disables it, and any other value causes the current state to be returned.
RETURNS	The mode (INSTRUMENT_ON or INSTRUMENT_OFF) currently in force.

ERRNO

SEE ALSO wvLib

# wvSockUploadPathClose( )

NAME	<pre>wvSockUploadPathClose() - close the socket upload path</pre>
SYNOPSIS	void wvSockUploadPathClose ( UPLOAD_ID upId /* generic upload-path descriptor */ )
DESCRIPTION	This routine closes the socket connection to the event receiver on the host.
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	wvSockUploadPathLib, sockUploadPathCreate()

# wvSockUploadPathCreate()

NAME	wvSockUploadPathCreate() – establish an upload path to the host using a socket
SYNOPSIS	<pre>UPLOAD_ID wvSockUploadPathCreate   (     char *ipAddress, /* server's hostname or IP address innotation */     short port /* port number to bind to */   )</pre>
DESCRIPTION	This routine initializes the TCP/IP connection to the host process that receives uploaded events. It can be retried if the connection attempt fails.
RETURNS	The <b>UPLOAD_ID</b> , or <b>NULL</b> if the connection cannot be completed or memory for the ID is not available.
ERRNO	Not Available
SEE ALSO	wvSockUploadPathLib, sockUploadPathClose()

# wvSockUploadPathLibInit()

NAME	<pre>wvSockUploadPathLibInit() - initialize wvSockUploadPathLib library</pre>
SYNOPSIS	STATUS wvSockUploadPathLibInit (void)
DESCRIPTION	This routine initializes <b>wvSockUploadPathLib</b> by pulling in the routines in this file for use with Wind River System Viewer. It is called during system configuration from <b>usrWindview.c</b> .
RETURN	OK.
RETURNS	Not Available
ERRNO	Not Available
SEE ALSO	wvSockUploadPathLib

# wvSockUploadPathWrite()

NAME	<pre>wvSockUploadPathWrite() - write to the socket upload path</pre>
SYNOPSIS	<pre>int wvSockUploadPathWrite   (    UPLOAD_ID upId, /* generic upload-path descriptor */    char * pStart, /* address of data to write */    size_t size /* number of bytes of data at pStart */   )</pre>
DESCRIPTION	This routine writes <i>size</i> bytes of data beginning at $pStart$ to the upload path between the target and the event receiver on the host.
RETURNS	The number of bytes written, or ERROR.
ERRNO	Not Available
SEE ALSO	wvSockUploadPathLib, wvSockUploadPathCreate()

### wvTmrRegister()

wvTmrRegister() - register a timestamp timer NAME SYNOPSIS void wvTmrRegister ( UINTFUNCPTR wvTmrRtn, /\* timestamp routine \*/ UINTFUNCPTR wvTmrLockRtn, /\* locked timestamp routine \*/ FUNCPTR wvTmrEnable, /\* enable timer routine \*/ FUNCPTR wvTmrDisable, /\* disable timer routine \*/ FUNCPTR wvTmrConnect, /\* connect to timer routine \*/ UINTFUNCPTR wvTmrPeriod, /\* period of timer routine \*/ /\* frequency of timer routine \*/ UINTFUNCPTR wvTmrFreq ) DESCRIPTION This routine registers a timestamp routine for each of the following: wvTmrRtn a timestamp routine, which returns a timestamp when called (must be called with interrupts locked). wvTmrLockRtn a timestamp routine, which returns a timestamp when called (locks interrupts). wvTmrEnable an enable-timer routine, which enables the timestamp timer. wvTmrDisable a disable-timer routine, which disables the timestamp timer. wvTmrConnect a connect-to-timer routine, which connects a handler to be run when the timer rolls over; this routine should return ERROR if the system clock tick is to be used. wvTmrPeriod a period-of-timer routine, which returns the period of the timer. *wvTmrFreq* a frequency-of-timer routine, which returns the frequency of the timer. If any of these routines is set to NULL, the behavior of instrumented code is undefined. RETURNS N/A ERRNO SEE ALSO wvTmrLib

# wvTsfsUploadPathClose()

NAME	<pre>wvTsfsUploadPathClose() – close the TSFS-socket upload path</pre>
SYNOPSIS	void wvTsfsUploadPathClose ( UPLOAD_ID upId /* generic upload-path descriptor */ )
DESCRIPTION	This routine closes the TSFS-socket connection to the event receiver on the host
RETURNS	N/A
ERRNO	Not Available
SEE ALSO	wvTsfsUploadPathLib, wvTsfsUploadPathCreate()

# wvTsfsUploadPathCreate()

NAME	wvTsfsUploadPathCreate() – open an upload path to the host using a TSFS socket
SYNOPSIS	UPLOAD_ID wvTsfsUploadPathCreate ( char *ipAddress, /* server's IP address innotation */ short port /* port number to bind to */ )
DESCRIPTION	This routine opens a TSFS socket to the host to be used for uploading event data. After successfully establishing this connection, an UPLOAD_ID is returned which points to the TSFS_UPLOAD_DESC that is passed to <b>open()</b> , <b>close()</b> , <b>read()</b> , etc. for future operations.
RETURNS	The <b>UPLOAD_ID</b> , or <b>NULL</b> if the connection cannot be completed or not enough memory is available.
ERRNO	Not Available
SEE ALSO	wvTsfsUploadPathLib, wvTsfsUploadPathClose( )

# wwTsfsUploadPathLibInit()NAMEwwTsfsUploadPathLibInit() - initialize wwTsfsUploadPathLib librarySYNOPSISSTATUS wwTsfsUploadPathLibInit (void)DESCRIPTIONThis routine initializes wwTsfsUploadPathLib by pulling in the routines in this file for use<br/>with the Wind River System Viewer. It is called during system configuration from<br/>usrWindview.c.RETURNSOK.ERRNONot AvailableSEE ALSOwvTsfsUploadPathLib

### wvTsfsUploadPathWrite()

NAME	<pre>wvTsfsUploadPathWrite() - write to the TSFS upload path</pre>
SYNOPSIS	<pre>int wvTsfsUploadPathWrite   (     UPLOAD_ID upId, /* generic upload-path descriptor */     char * pStart, /* address of data to write */     size_t size /* number of bytes of data at pStart */   )</pre>
DESCRIPTION	This routine writes <i>size</i> bytes of data beginning at <i>pStart</i> to the upload path connecting the target with the host receiver.
RETURNS	The number of bytes written, or <b>ERROR</b> .
ERRNO	Not Available
SEE ALSO	wvTsfsUploadPathLib, wvTsfsUploadPathCreate()

	wvUploadStart()
NAME	<pre>wvUploadStart() - start upload of events to the host</pre>
SYNOPSIS	<pre>WV_UPLOADTASK_ID wvUploadStart     (     WV_LOG * pWvLog,</pre>
DESCRIPTION	This routine starts uploading events from the System Viewer log to the host. Events can be uploaded either continuously or in one pass until the log is emptied. If <i>uploadContinuously</i> is set to <b>TRUE</b> , the task uploading events pends until more data arrives in the buffer. If <b>FALSE</b> , the buffer is flushed without waiting, but this routine returns immediately with an ID that can be used to kill the upload task. Upload is done by spawning the task <b>tWVUpload</b> . The log to upload is identified by <i>pWvLog</i> , and the upload path to use is identified by <i>pathId</i> .
	This routine blocks if no event data is in the buffer, so it should be called before event logging is started to ensure the buffer does not overflow.
RETURNS	A valid <b>WV_UPLOADTASK_ID</b> if started for continuous upload, a non- <b>NULL</b> value if started for one-pass upload, and <b>NULL</b> if the task can not be spawned or memory for the descriptor can not be allocated.
ERRNO	S_memLib_NOT_ENOUGH_MEMORY
SEE ALSO	wvLib

# wvUploadStop()

NAME	<pre>wvUploadStop() - stop upload of events to host</pre>
SYNOPSIS	STATUS wvUploadStop ( WV_UPLOADTASK_ID upTaskId )
DESCRIPTION	This routine stops continuous upload of events to the host. It does this by making a request to the upload task to terminate after it has emptied the buffer. For this reason it is important to make sure data is no longer being logged to the buffer before calling this routine.



### wvUploadTaskConfig()

NAME	<pre>wvUploadTaskConfig() – set priority and stacksize of tWVUpload task</pre>
SYNOPSIS	<pre>void wvUploadTaskConfig   (    int stackSize, /* the new stack size for tWVUpload */    int priority /* the new priority for tWVUpload */ )</pre>
DESCRIPTION	This routine sets the stack size and priority of future instances of the event-data upload task, created by calling <b>wvUploadStart()</b> . The default stack size for this task is 5000 bytes, and the default priority is 150.
RETURNS	N/A
ERRNO	
SEE ALSO	wvLib

### xattrib()

VxWorks Kernel API Reference, 6.6 xbdBlkDevCreate()

DESCRIPTION	This function is essentially the same as <b>attrib()</b> , but it accepts wildcards in <i>fileName</i> , and traverses subdirectories in order to modify attributes of entire file hierarchies.
	The <i>attr</i> argument string may contain must start with either "+" or "-", meaning the attribute flags which will follow should be either set or cleared. After "+" or "-" any of these four letter will signify their respective attribute flags - "A", "S", "H" and "R".
EXAMPLE	-> xattrib( "/sd0/sysfiles", "+RS") /* write protect "sysfiles" */ -> xattrib( "/sd0/logfiles", "-R") /* unprotect logfiles before deletion */ -> xdelete( "/sd0/logfiles")
CAVEAT	This function may call itself in accordance with the depth of the source directory, and allocates 2 kB of heap memory per stack frame, meaning that to accommodate the maximum depth of subdirectories which is 20, at least 40 kB of heap memory should be available.
RETURNS	OK, or ERROR if the file can not be opened.
ERRNO	Not Available
SEE ALSO	usrFsLib, dosFsLib, the VxWorks programmer guides.

# xbdBlkDevCreate()

NAME	xbdBlkDevCreate() – create an XBD block device wrapper
SYNOPSIS	<pre>device_t xbdBlkDevCreate   (    BLK_DEV * bd, /* pointer to block device */    const char * name /* pointer to device name */   )</pre>
DESCRIPTION	This routine creates an XBD block device wrapper.
RETURNS	a device identifier upon success, or NULLDEV otherwise
ERRNO	
SEE ALSO	xbdBlkDev

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xbdBlkDevCreateSync()	)
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NAME	<b>xbdBlkDevCreateSync()</b> – synchronously create an XBD block device wrapper
SYNOPSIS	<pre>device_t xbdBlkDevCreateSync   (    BLK_DEV * bd, /* pointer to block device */    const char * name /* pointer to device name */   )</pre>
DESCRIPTION	This routine creates an XBD block device wrapper. It returns after the entire XBD stack has been created/initialized.
RETURNS	a device identifier upon success, or NULLDEV otherwise
ERRNO	
SEE ALSO	xbdBlkDev

## xbdBlkDevDelete()

NAME	xbdBlkDevDelete() – deletes an XBD block device wrapper
SYNOPSIS	<pre>STATUS xbdBlkDevDelete   (     device_t d, /* device_t returned from xbdBlkDevCreate */     BLK_DEV ** ppbd /* pointer to block device pointer */   )</pre>
DESCRIPTION	This routine deletes or destroys an XBD block device wrapper.
	The <i>d</i> parameter specifies the XBD block wrapper to delete. This should be the same value that was returned from <b>xbdBlkDevCreate()</b>
	The <i>ppbd</i> parameter is an out parameter that can be used to return the block device pointer used in <b>xbdBlkDevCreate()</b> . If specified as <b>NULL</b> no attempt to return the block device pointer is attempted.
RETURNS	a device identifier upon success, or NULLDEV otherwise
ERRNO	
SEE ALSO	xbdBlkDev

VxWorks Kernel API Reference, 6.6 xbdBlkDevLibInit()

### xbdBlkDevLibInit()

NAME	<b>xbdBlkDevLibInit()</b> – initialize the XBD block device wrapper
SYNOPSIS	STATUS xbdBlkDevLibInit ( int xbdServiceTskPri )
DESCRIPTION	This routine initializes the XBD block device wrapper.
RETURNS	ОК
ERRNO	N/A
SEE ALSO	xbdBlkDev

# xbdCbioDevCreate()

NAME	xbdCbioDevCreate() – create an XBD CBIO device wrapper
SYNOPSIS	<pre>device_t xbdCbioDevCreate   (     CBIO_DEV_ID cbio, /* handle to CBIO device */     const char * name, /* pointer to device name */     unsigned int opts /* options for device */   )</pre>
DESCRIPTION	This routine creates an XBD CBIO device wrapper. It returns after the entire XBD stack has been created/initialized.
	cbio handle to previously created CBIO device
	<i>name</i> base name of the XBD/CBIO wrapper
	The <i>opts</i> argument is a bit-wise or'ed combination of options controlling the operation of this routine as follows:
	XBD_CBIO_NOWAIT Function will not wait until path(s) are instantiated and will return immediately
	<b>XBD_CBIO_NOPART</b> Wrapper is not capabable of supporting partitions. Device will be viewed as one partition spanning the entire media.

	XBD_CBIO_DEFAULT The default behaviour. Partitions are supported and xbdCbioDevCreate will not return until all paths are instantiated.
RETURNS	a device identifier upon success, or NULLDEV otherwise
ERRNO	
SEE ALSO	xbdCbioDev

# xbdCbioDevDelete()

NAME	<b>xbdCbioDevDelete( )</b> – deletes an XBD CBIO device wrapper
SYNOPSIS	<pre>STATUS xbdCbioDelete   (     device_t d, /* device_t returned from xbdCbioDevCreate */     CBIO_DEV_ID* pCbio /* pointer to CBIO handle */   )</pre>
DESCRIPTION	This routine deletes or destroys an XBD CBIO device wrapper.
	The <i>d</i> parameter specifies the XBD CBIO wrapper to delete. This should be the same value that was returned from <b>xbdCbioDevCreate()</b>
	The <i>ppbd</i> parameter is an out parameter that can be used to return the CBIO handle used in <b>xbdCbioDevCreate()</b> . If specified as <b>NULL</b> no attempt to return the handle is attempted.
RETURNS	OK upon success, or ERROR otherwise
ERRNO	
SEE ALSO	xbdCbioDev

# xbdCbioLibInit()

NAME	<b>xbdCbioLibInit()</b> – initialize the XBD block device wrapper
SYNOPSIS	STATUS xbdCbioDevLibInit (void)
DESCRIPTION	This routine initializes the XBD block device wrapper.

VxWorks Kernel API Reference, 6.6 xbdCreatePartition()

RETURNS OK

ERRNO N/A

SEE ALSO xbdCbioDev

### xbdCreatePartition()

xbdCreatePartition() - partition an XBD device NAME SYNOPSIS STATUS xbdCreatePartition ( char \*pathName, /\* name of device to partition \*/ int nPart, /\* number of partitions \*/ /\* space percentage for second partition \*/ int sizel, int size2, /\* space percentage for third partition \*/ /\* space percentage for fourth partition \*/ int size3 ) DESCRIPTION This function is capable of creating only one partition table - the MBR, and will not create any Bootable or Extended partitions. Therefore, only 4 primary partitions are supported. *pathName* is the name the device to be partitioned. The *nPart* argument contains the number of partitions to create. If *nPart* is 0 or 1, a single partition covering the entire disk is created. If *nPart* is between 2 and 4, the arguments *size1*, size2 and size3 contain (as integers) the percentage of disk space to be assigned to the 2nd, 3rd, and 4th partitions respectively. The first partition (partition 0) will be assigned the remaining space. Thus, the sum of the three sizes should be less than 100. Partition sizes will be rounded down to be multiple of whole tracks so that partition Cylinder/Head/Track fields will be initialized as well as the LBA fields. Although the CHS fields are written they are not used in VxWorks, and can not be guaranteed to work correctly on other systems. RETURNS OK upon success, ERROR otherwise Not Available ERRNO partLib SEE ALSO

# xbdRamDiskDevCreate()

NAME	xbdRamDiskDevCreate() – create an XBD ram disk
SYNOPSIS	<pre>device_t xbdRamDiskDevCreate   (     unsigned blockSize, /* block size in bytes */     unsigned totalSize, /* disk size in bytes */     BOOL flag, /* should the disk support partitions? */     const char * name /* name of ram disk */   )</pre>
DESCRIPTION	This routine creates an XBD ram disk. The ram disk links into the file system monitor and eventing framework.
RETURNS	The ID of the XBD created(device_t) or NULLDEV if the routine fails
ERRNO	Not Available
SEE ALSO	xbdRamDisk

# xbdRamDiskDevDelete()

NAME	xbdRamDiskDevDelete() – XBD Ram Disk Deletion routine
SYNOPSIS	STATUS xbdRamDiskDevDelete ( device_t d /* device_t returned from xbdRamDiskDevCreate */ )
DESCRIPTION	This routine deletes or destroy an instantion of an XBD ram disk. The ram disk to be deleted is identified by the supplied device_t. This value must have been previously returned from the corresponding xbdRamDiskDevCreate function. All resource associated with the ram disk are freed. Any file systems sitting on top of the ram disk are ejected.
RETURNS	<b>OK</b> on success or <b>ERROR</b> if the supplied device_t doesn't map to an existing and valid XBD.
ERRNO	Not Available
SEE ALSO	xbdRamDisk

VxWorks Kernel API Reference, 6.6 xbdTransDevCreate()

### xbdTransDevCreate()

**NAME xbdTransDevCreate()** – create a transactional XBD.

SYNOPSIS device\_t xbdTransDevCreate ( device\_t subDev /\* lower level device \*/ )

**RETURNS** device\_t, or NULLDEV on failure.

ERRNO Not Available

SEE ALSO xbdTrans, dosFsDevCreate().

### xbdTransInit()

NAME xbdTransInit() – initialize the transactional XBD subsystem. SYNOPSIS STATUS xbdTransInit ( void ) We just plug ourselves in to the file system monitor so that we get called to probe partitions DESCRIPTION as they are instantiated. OK if all went well, ERROR otherwise. RETURNS ERRNO N/A xbdTrans SEE ALSO

### xcopy()

**NAME xcopy()** – copy a hierarchy of files with wildcards

SYNOPSIS STATUS xcopy

	( const char * source, /* source directory or wildcard name */ const char * dest /* destination directory */ )	
DESCRIPTION	<i>source</i> is a string containing a name of a directory, or a wildcard or both which will cause this function to make a recursive copy of all files residing in that directory and matching the wildcard pattern into the <i>dest</i> directory, preserving the file names and subdirectories.	
CAVEAT	This function may call itself in accordance with the depth of the source directory, and allocates 3 kB of heap memory per stack frame, meaning that to accommodate the maximum depth of subdirectories which is 20, at least 60 kB of heap memory should be available.	
RETURNS	OK, or ERROR if any operation has failed.	
ERRNO	Not Available	
SEE ALSO	usrFsLib, tarLib, cp(), the VxWorks programmer guides.	

# xdelete()

NAME	<b>xdelete()</b> – delete a hierarchy of files with wildcards		
SYNOPSIS	STATUS xdelete ( const char * source /* source directory or wildcard name */ )		
DESCRIPTION	<i>source</i> is a string containing a name of a directory, or a wildcard or both which will cause this function to recursively remove all files and subdirectories residing in that directory and matching the wildcard pattern. When a directory is encountered, all its contents are removed, and then the directory itself is deleted.		
	Note that the wildcard matching is limited to a single directory level.		
	dir is valid *.c is valid dir/*.c is valid *a/*.c is not valid		
RETURNS	OK or ERROR if any operation has failed.		
ERRNO	Not Available		
SEE ALSO	usrFsLib, cp(), copy(), xcopy(), tarLib, the VxWorks programmer guides.		

# y()

NAME	y() – return the contents of the y register (SimSolaris)	
SYNOPSIS	<pre>int y   (    int taskId /* task ID, 0 means default task */ )</pre>	
DESCRIPTION	This command extracts the contents of the y register from the TCB of a specified task. If <i>taskId</i> is omitted or 0, the default task is assumed.	
RETURNS	The contents of the y register.	
ERRNO	Not Available	
SEE ALSO	dbgArchLib, VxWorks Programmer's Guide: Debugging	

# ykRegister()

SEE ALSO	mvYukonVxbEnd
ERRNO	N/A
RETURNS	N/A
DESCRIPTION	This routine registers the Template driver with VxBus as a child of the PCI bus type.
SYNOPSIS	void ykRegister(void)
NAME	<b>ykRegister()</b> – register with the VxBus subsystem

# ynRegister()

NAME	<b>ynRegister()</b> – register with the VxBus subsystem

SYNOPSIS void ynRegister(void)

**DESCRIPTION** This routine registers the Yukon II driver with VxBus as a child of the PCI bus type.

RETURNS N/A

ERRNO N/A

SEE ALSO mvYukonIIVxbEnd